



Performance Evaluation of CL-PHB Admission and Pre-emption Algorithms

draft-zhang-pcn-performance-evaluation-01.pdf

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Outline

- **Simulation results for the example admission control algorithm described in**
<http://www.ietf.org/internet-drafts/draft-briscoe-tsvwg-cl-phb-0>
- **Flow Preemption simulation results for the example algorithm described in**
<http://www.ietf.org/internet-drafts/draft-briscoe-tsvwg-cl-phb-0>

Note: Admission and Preemption were run independently

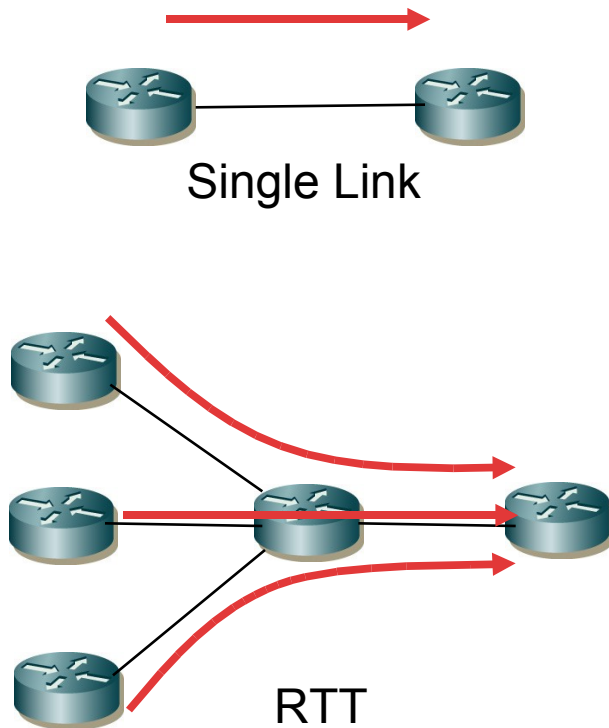
- Reasonable if admission and preemption thresholds are sufficiently apart (10-20%)
- Further work needed to simulate interaction between Admission and Preemption if threshold are closer together

Simulation Environment (Traffic Model)

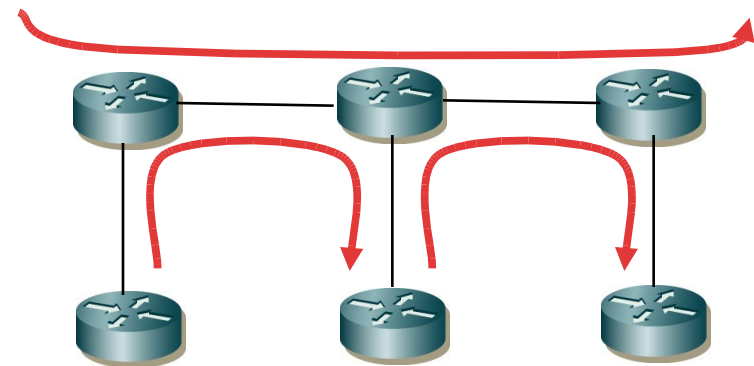
- **CBR**
 - average rate: 64Kbps
- **On-Off Voice (VBR)**
 - Voice w/silence suppression
 - average rate: 21.76Kbps
- **Synthetic "Video" (SVD)**
 - High Peak-to-Mean Ratio (4:1) on-off VBR Traffic
 - Average rate: 12Mbps
- **Real Video Traces (VTR) – new in this version of the draft**
 - Frame size traces of MPEG-4 and H.263 encode video
 - Average rate: 769Kbps
 - <http://www.tkn.tu-berlin.de/research/trace/trace.html>
- **Randomization of Base Traffic Models – new in this version of the draft**
 - Randomly moving the packet by a small amount of time around its transmission time to simulate small queuing delays

Simulation Environment (Topologies)

- **Single Bottleneck**



- **Multiple Bottlenecks (new)**



2-BN Parking Lot (PLT)
(3 and 5 bottleneck
PLT also simulated)

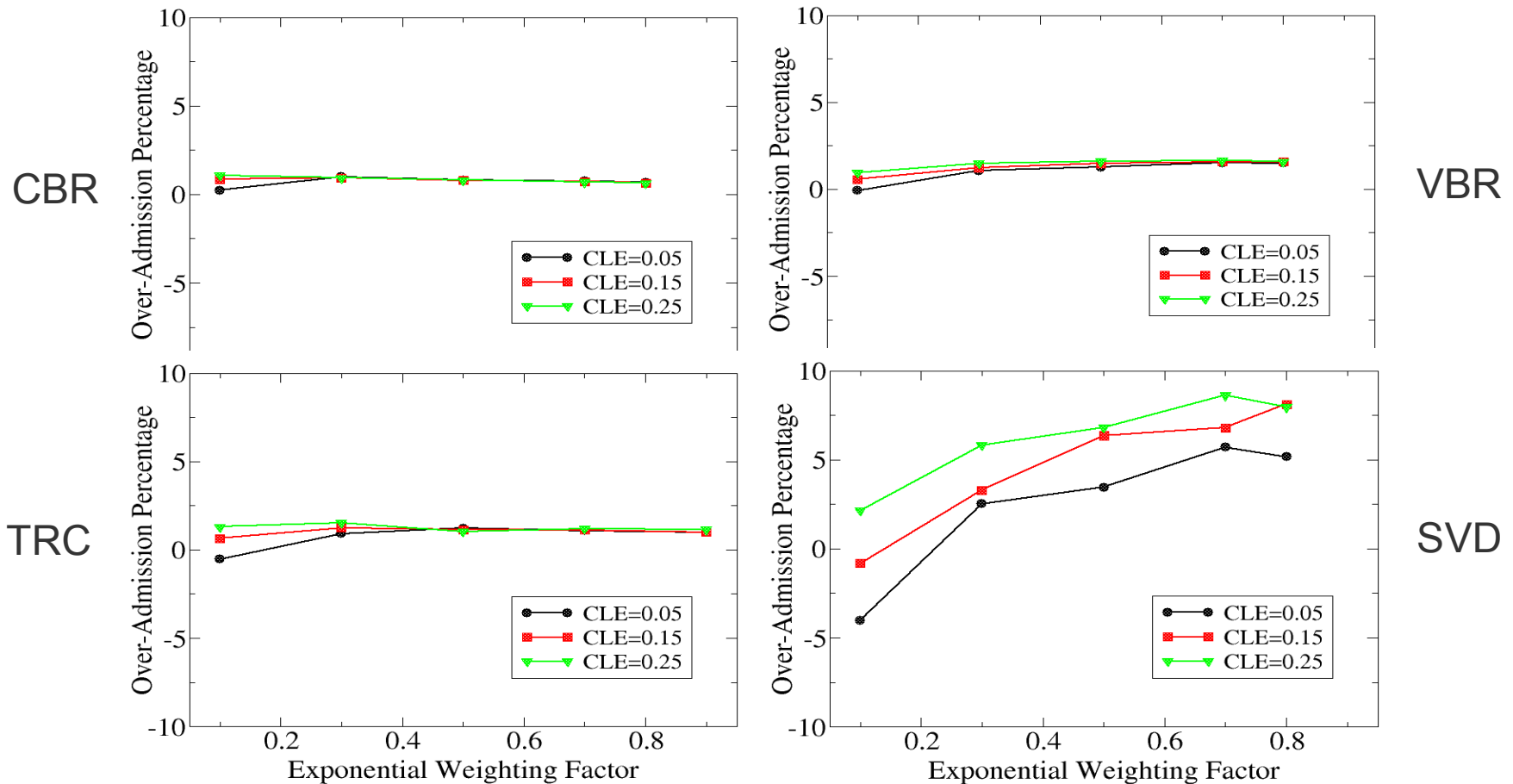
Admission Control Results

- Sensitivity to call arrival assumption (Poisson vs. burstier Batch): **works well for both on links 10M and up**
- Sensitivity to marking thresholds
 - (Ramp vs step): **no substantial difference in the simulated setup**
 - Limited variation of upper/lower marking threshold: **relatively insensitive, but more study needed**
- Sensitivity to RTT: **no effect in the one bottleneck/stable traffic case**
- Sensitivity to EWMA weight and CLE: **Relatively insensitive to parameter change**
- Effect of Ingress-Egress Aggregation: **no effect**
- Effect of Multiple Bottlenecks: **no effect**

Admission Control Results

Sensitivity to EWMA weight and CLE

- Relatively insensitive to parameter change
 - More stressful for SVD, but within -3% to +10% over-admit



Admission Control Results

Effect of Ingress-Egress Aggregation

- **Difference in Ingress-Egress aggregation appears to have no effect on the performance in the simulation we performed**
 - **requires further investigation**

	CBR	VBR	VTR	SVD
SingleLink	1.905	1.948	1.539	8.62
RTT (100 Ingress)	1.956	2.199	1.868	11.26

over-admission-perc. on the bottleneck

Admission Control Results

Effect of Multiple Bottlenecks

- **No visible multiple bottleneck effect**

	CBR	VBR	VTR	SVD
SingleLink	1.905	1.948	1.539	8.62
PLT (5BN)	1.149	1.501	1.117	4.737

over-admission-perc. on the bottleneck

Note: In the case of PLT, the worst among the 5 BN is used

Admission Control Results (Conclusion)

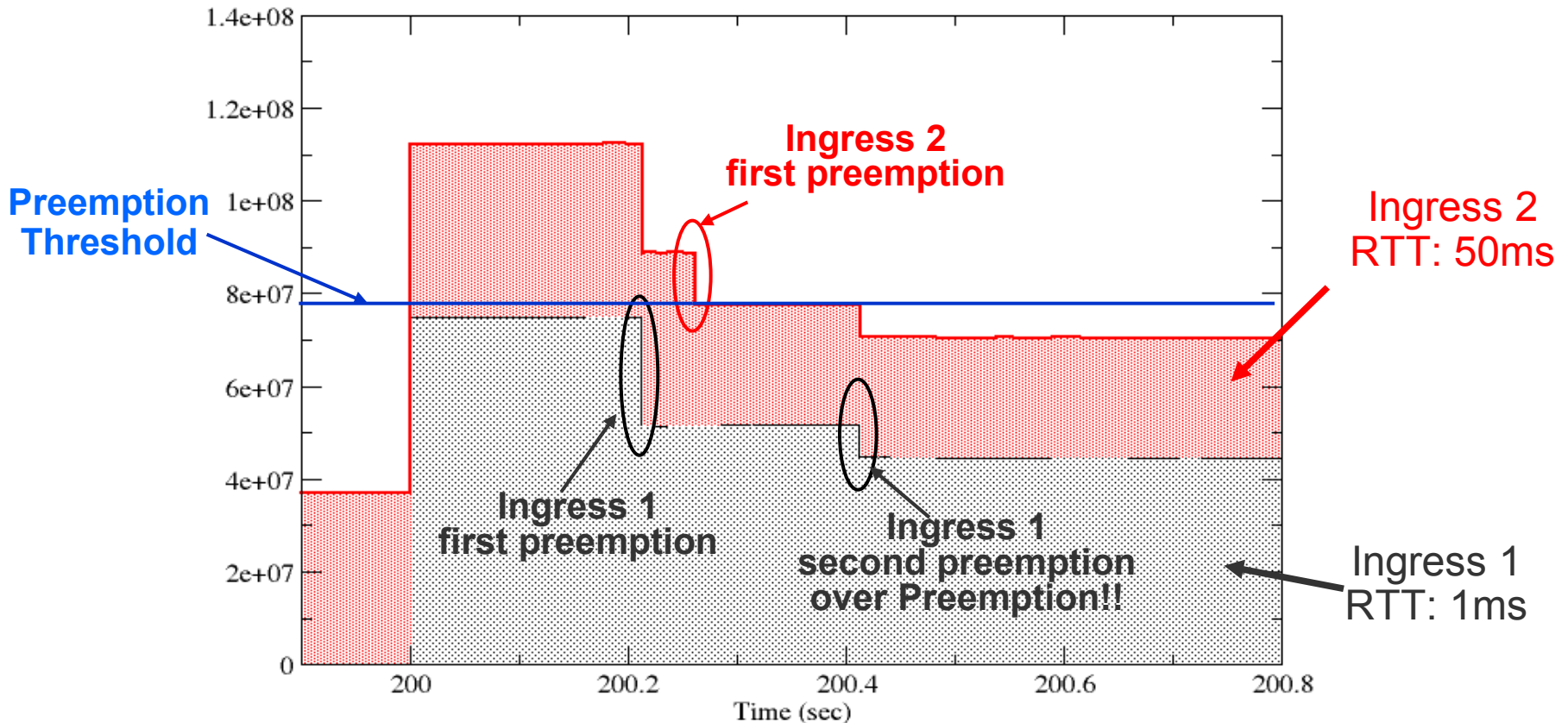
- **Performance remains reasonable even for really low ingress-egress aggregation levels**
- **The algorithm is relatively insensitive to variation of key parameter settings**
- **Synthetic video traffic SVD was the most challenging for all topologies, and the performance of real video traces (VTR) was substantially better**
- **No performance degradation is observed in a multi-bottleneck topology**
- **To summarize, so far all seem performing as advertised**

Preemption Results

- **Sanity check on SingleLink topology: worked as expected**
- **Effect of RTT Difference: Visible effect with relative RTT difference, though not significant**
- **Effect of Ingress-Egress Aggregation: visible effect with low aggregation, though not as feared**
- **Effect of Multiple Bottlenecks: worked as expected**

Preemption Results: Effect of RTT Difference

- Absolute value of RTT has no effect
- Relative difference in RTT causes limited over-preemption
 - 6% – 10% in experiments we run



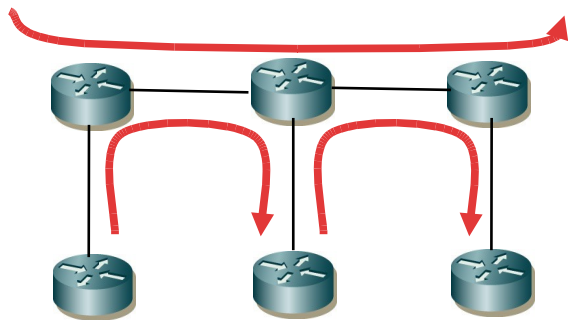
Preemption Results: Effect of Ingress-Egress Aggregation

- **In theory, at low aggregation serious over-preemption may occur**
 - **When ingress-egress aggregate has only one flow, all flows can get marked at the bottleneck and all get preempted**
- **As a rule, over-preemption does occur but much smaller than the worst case behavior**
 - **e.g. for randomized CBR**
 - **0.53% over-preemption with sufficient aggregation**
 - **13.19% over-preemption 1 Flow/Ingress**
- **In some simulations, it does not occur at all**
 - **Due to marking synchronization or specific traffic patterns**
 - **Can not be relied on in general**
- **Over-preemption for low aggregation can not be written off as a “corner case”**

Preemption Control Results

Effect of Multiple Bottlenecks

- **Ingress-Egress Aggregates that travel more bottlenecks will see “beat-down” effect (over-preemption), as theoretically expected**
 - Long-haul flow gets excessively marked at subsequent bottlenecks
 - Upstream bottlenecks become underutilized
 - The absolute amount of “beat-down” depends on traffic matrix
- **Actual simulation results are very close to the theoretically expected ones**



	CBR	VBR	VTR	SVD
Long-haul Flow	13.56	16.30	14.77	23.31
First Bottleneck	9.61	11.59	11.06	18.13
Last Bottleneck	0.92	2.13	2.89	10.85

over-preemption-perc.

Key directions for further evaluation

- **further investigation on call arrival assumptions (burstier than Poisson)**
- **more sophisticated and/or realistic topology and traffic matrix**
- **Mix of traffic types on bottleneck**
- **Interaction between admission and preemption**

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