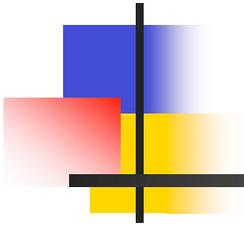


Federated-fs protocol overview

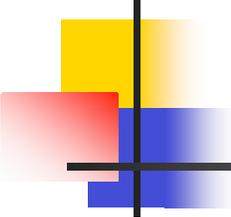


R. Tewari, M. Naik

IBM Almaden

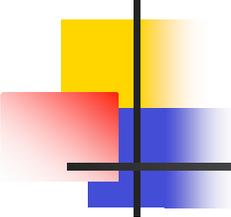
D. Ellard, C. Everhart

Network Appliance



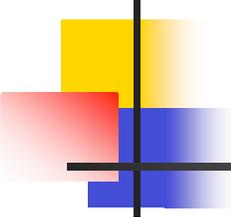
Outline

- Background and motivation
- Terms and definitions
- The protocol
- Next steps



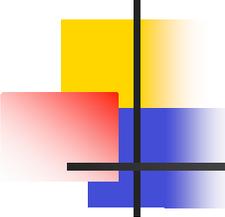
Background

- Idea: build a cross-platform federated file system with a shared common namespace.
 - Independent file servers are federation members
 - Federation members are peers.
 - Could have multiple administrative boundaries
- Use Cases:
 - Single enterprise with multi vendor file servers and file server collections
 - Multiple enterprises sharing independently managed namespace(s)



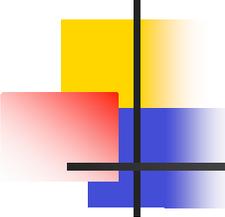
Protocol Goals

- Specify a set of interfaces by which file servers and collections of file servers with different administrators can form a file server federation
 - that provides a namespace composed of the filesystems physically hosted on and exported by the file servers of the federation.
- It should be possible, using a system that implements the interfaces, to share a *common* namespace across all the file servers in the federation.
- It should also be possible for different file servers in the federation to project *different* namespaces and enable clients to traverse them.



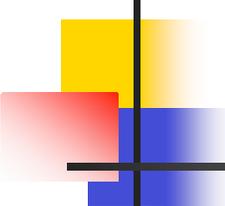
Terms and Definitions

- Fileset: the basic organizational unit for data
 - Container of data abstraction
 - To the user, it's a filesystem or a directory tree
 - Behind the scenes, can be implemented as a replicated, mobile storage container
- FSN: "Fileset Name"
 - Symbolic name of a fileset
 - UUID + Name of the repository responsible for the fileset
 - FSN is globally unique



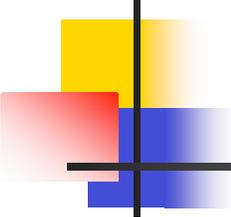
Terms...

- FSL: “Fileset Location”
 - Location of an implementation of a fileset
 - A fileset may be implemented by several FSLs
 - FSLs can come and go, change, etc.
- NSDB: “namespace database”
 - Repository of fileset and fileset relationship information
 - Keeps track of the mapping from FSNs to FSLs
 - Every FSN has a single authoritative NSDB
- Junction: binds a path within a fileset to a target FSN
 - Requires fileserver support
 - a junction can be viewed as a reference from a directory in one fileset to the root of the target fileset



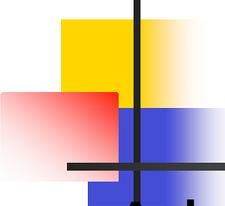
How it all fits together

- Filesets are the building blocks.
- Junctions define the relations between the filesets
- Namespace built by “mounting” filesets
- When an NFSv4 client accesses a junction (which looks to the client like a directory):
 - The fileservers find the FSN of the junction’s target fileset
 - The fileservers find (from the FSN) the NSDB location responsible for that FSN
 - The fileservers query the NSDB for the current set of FSLs
 - The fileservers use the NFSv4 referral mechanism
 - informs the client that the directory has “moved”
 - Returns the FSL information in the fs_locations attribute
 - The client is redirected to the root of the fileset that is the target of the junction



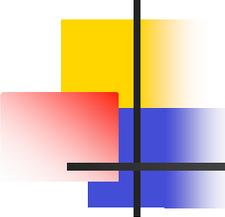
Protocol basic functional requirements

- An interface to add/delete/modify FSN → FSL mappings in the NSDB
- An interface to query the NSDB for FSN → FSL mappings
- An interface to create a junction at a fileserver
 - that a junction to some other fileset exists within a directory of a fileset location hosted by that server



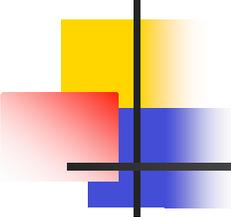
Three players

- Administrator
 - Initiates creation/updates/deletes of filesets, junctions, fileset locations
- NSDB
 - Maintains the fileset state
 - Responds to inserts/updates/queries
 - An NSDB is an LDAP server + persistent database
- NFSv4 server
 - In response to client requests, queries the NSDB for FSLs
 - In response to admin requests, updates/queries server state (junctions)



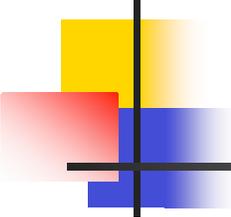
Three protocols

- Admin---NSDB
 - Update the NSDB state
 - Implemented as LDAP using fed-fs schema
- Fileserver---NSDB
 - Query the NSDB
 - Implemented as LDAP queries using fed-fs schema
- Admin---Fileserver
 - Update/query the NFSv4 server state (junctions)
 - Implemented via SUNRPC
 - Why not LDAP?
 - We don't want every NFS server to be an LDAP server as well (it can be an LDAP client). This seems like a simpler approach.



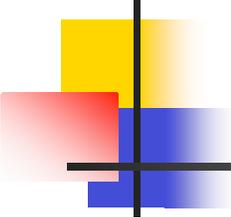
NSDB State

- Each NSDB uses the LDAP fed-fs schema
- Define objects
 - Fileset
 - Fileset location
 - Junctions
- Need fileserver to parse fed-fs schema
 - Query based on FSN not pathname
 - Convert to fs_locations
 - Not identical to `rnfs` experimental schema.



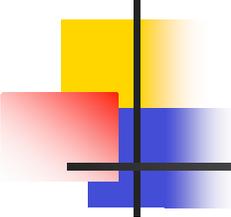
Discussion questions

- Choice of protocol for fileserver—admin interaction
- LDAP schema rdfs variant
- Common namespace
 - finding root fileserver not defined by protocol
- FSL consistency not defined
- Caching of namespace information across NSDBs
 - Default local NSDB not defined
- Security Issues
- Access/authorization ID mapping issues



Open source plans

- IBM Research Glamour base code on Linux to be made available using common public license



More Information at

- list federated-fs@sdsc.edu
- IETF Draft
 - draft-tewari-federated-fs-protocol-01.txt
- Emails:
 - Daniel dot Ellard at netapp dot com
 - tewarir at us dot ibm dot com
- BOF at FAST?