

LISP-Multicast

`draft-farinacci-lisp-multicast-00.txt`

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IETF Dublin - July 2008

Agenda

- Overview of LISP for Unicast Routing
- LISP-Multicast Design Goals
- LISP-Multicast Definitions
- Describe LISP-Multicast
- LISP-Multicast Interworking
- AF Case Studies
- Summary

LISP Internet Drafts

draft-farinacci-lisp-08.txt
draft-fuller-lisp-alt-02.txt
draft-lewis-lisp-interworking-01.txt
draft-farinacci-lisp-multicast-00.txt
draft-meyer-lisp-eid-block-01.txt

draft-mathy-lisp-dht-00.txt
draft-iannone-openlisp-implementation-01.txt
draft-brim-lisp-analysis-00.txt

draft-meyer-lisp-cons-04.txt
draft-lear-lisp-nerd-04.txt
draft-curran-lisp-emacs-00.txt

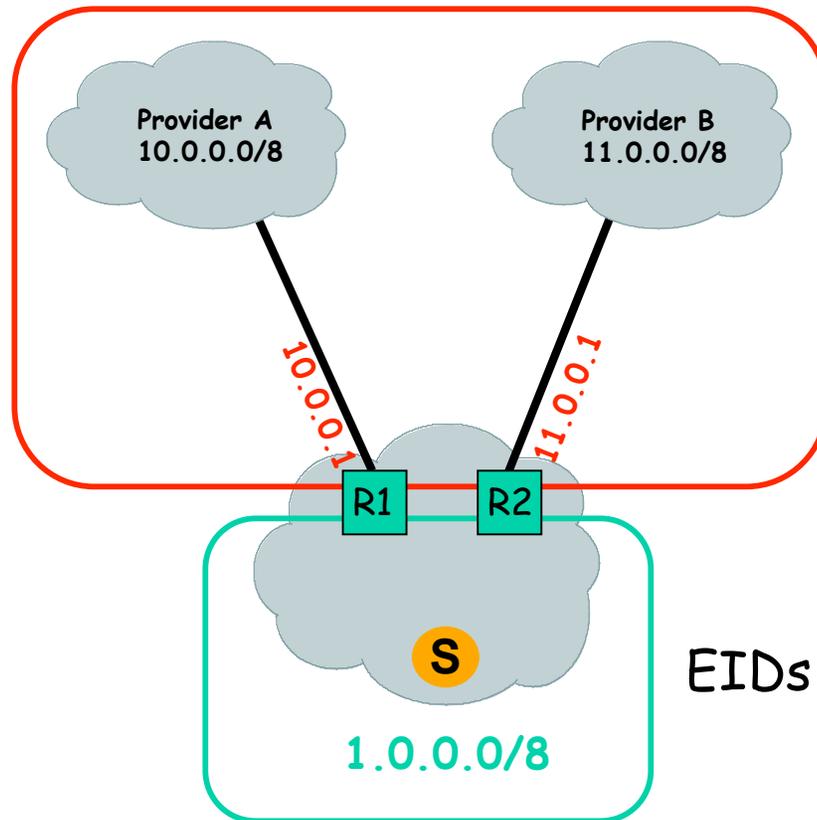
LISP Architecture

- Locator/ID Separation Protocol
 - Network-based solution
 - No changes to hosts whatsoever
 - No new addressing changes to site devices
 - Very few configuration file changes
 - Imperative to be incrementally deployable
 - Address family agnostic

LISP Architecture

- Locator/ID Separation Protocol
 - Defines architecture for 2 address spaces:
 - Endpoint IDs (EIDs)
 - Routing Locators (RLOCs)
 - Defines how Map-n-Encap routers operate and where they reside
 - Defines Variants for "EID routing"
 - Can use other designs for mapping EIDs to RLOCs
 - Defines encapsulation format for IPv4 and IPv6
 - Defines how RLOC reachability is performed
 - Defines how database mapping entries can be updated
 - Defines the Map-Request and Map-Reply format

Multi-Level Addressing



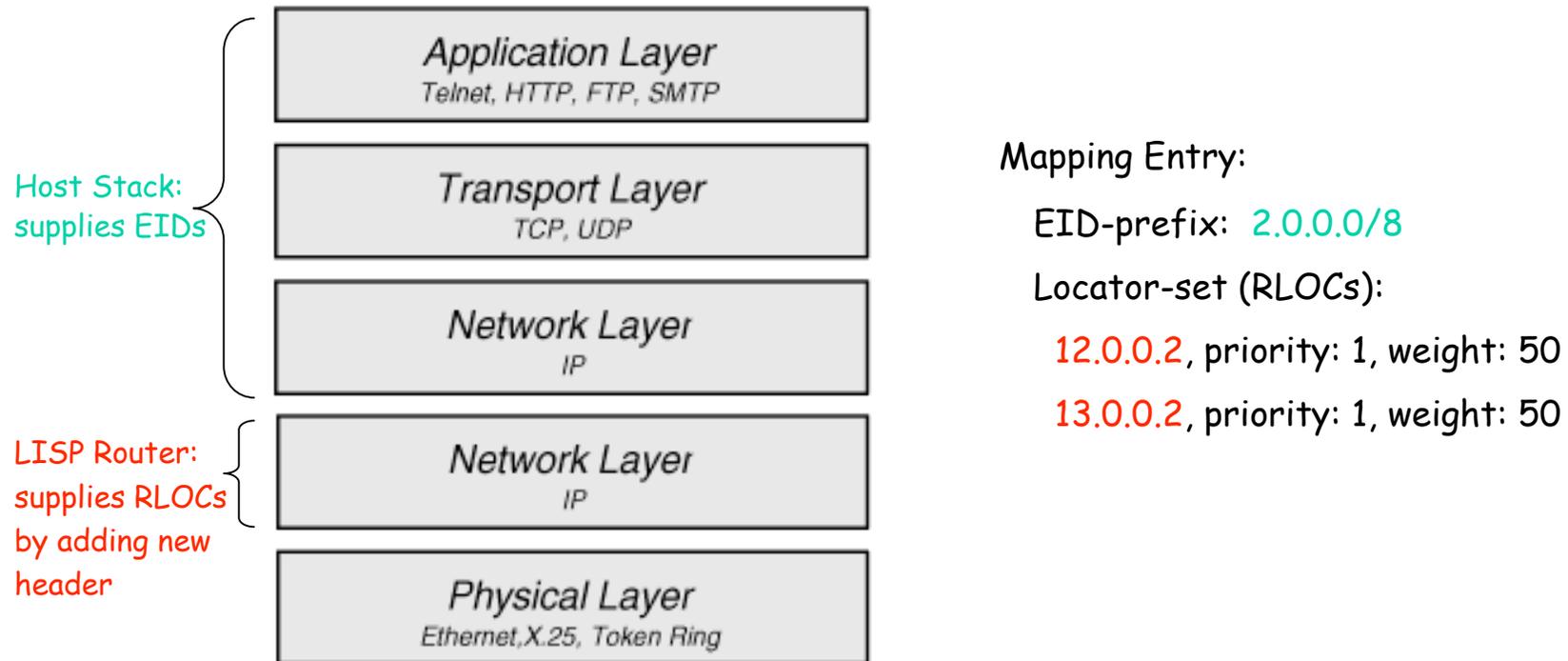
RLOCs used in the core

Mapping Database Entry:

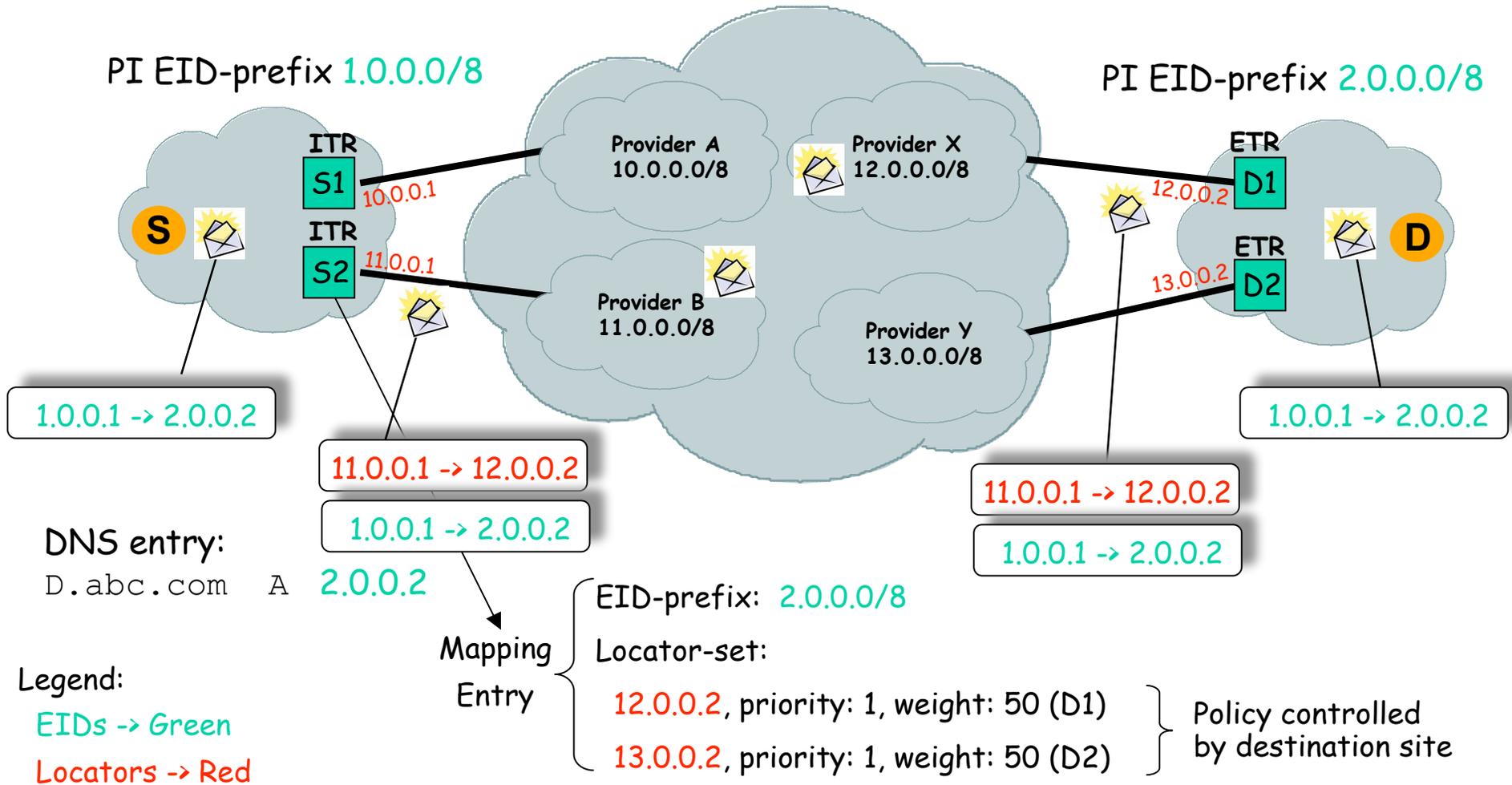
1.0.0.0/8 -> (10.0.0.1, 11.0.0.1)

EIDs are inside of sites

Map-n-Encap



Packet Forwarding



LISP Multicast Design Goals

- Keep EID state out of core network
- No head-end replication at source site
- Packets only go to receiver sites
- No changes to hosts, site routers, core routers
- Use existing protocols
- Support PIM SSM, don't preclude ASM & Bidir
- Have separate unicast and multicast policies

LISP Multicast Defs

- (S-EID, G)
 - S-EID is source host
 - G is group address receivers join to
 - State resides in source and receiver sites
- (S-RLOC, G)
 - S-RLOC is ITR on multicast tree
 - G is group address receivers join to
 - State resides in core
- Group addresses have neither ID or Location semantics
 - G is topologically opaque - can be used everywhere

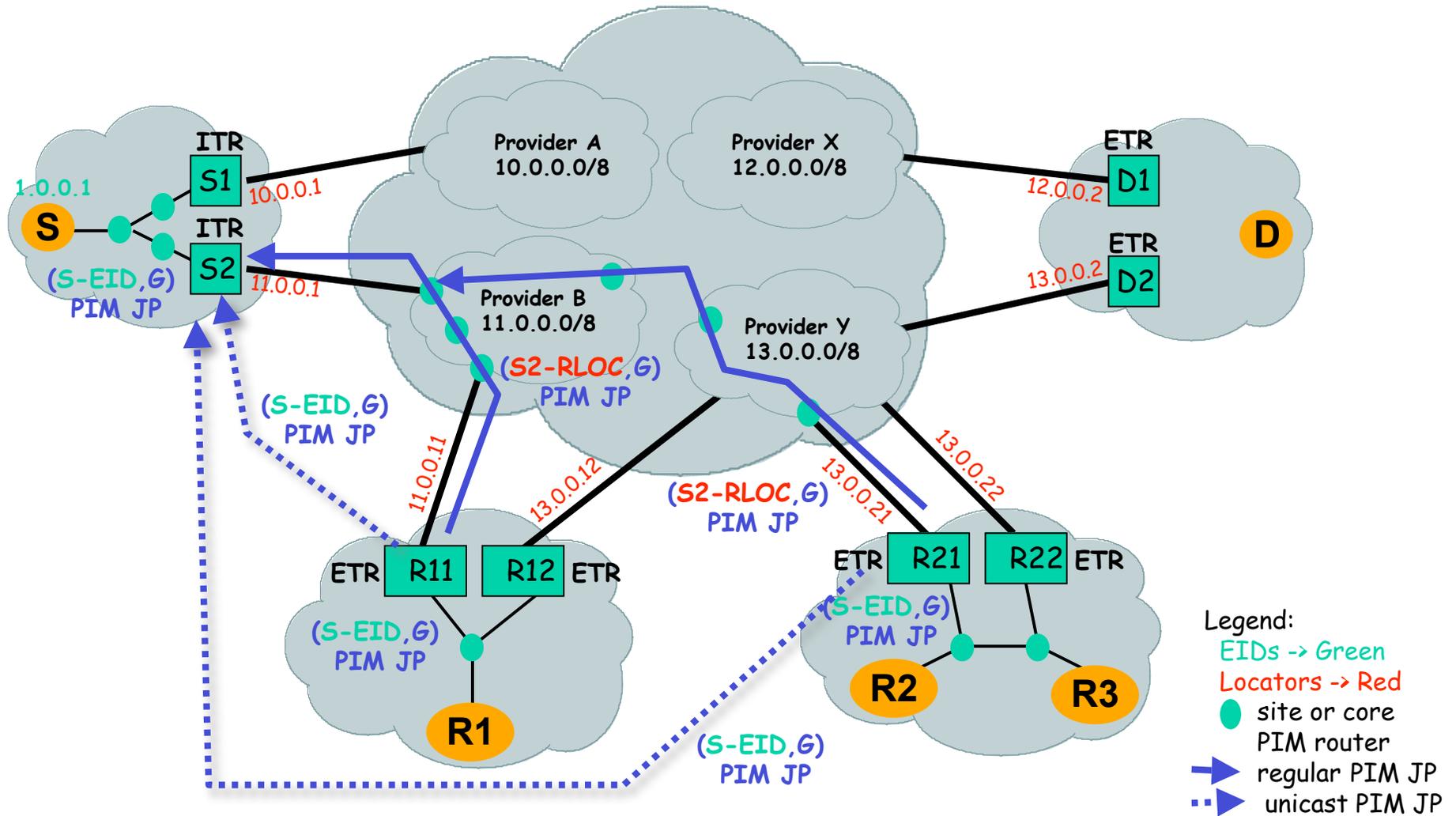
LISP Multicast Defs

- Non-LISP Site
 - Does not support LISP at all
- uLISP Site
 - Support LISP-Unicast only
- LISP-Multicast Site
 - Supports unicast and multicast LISP

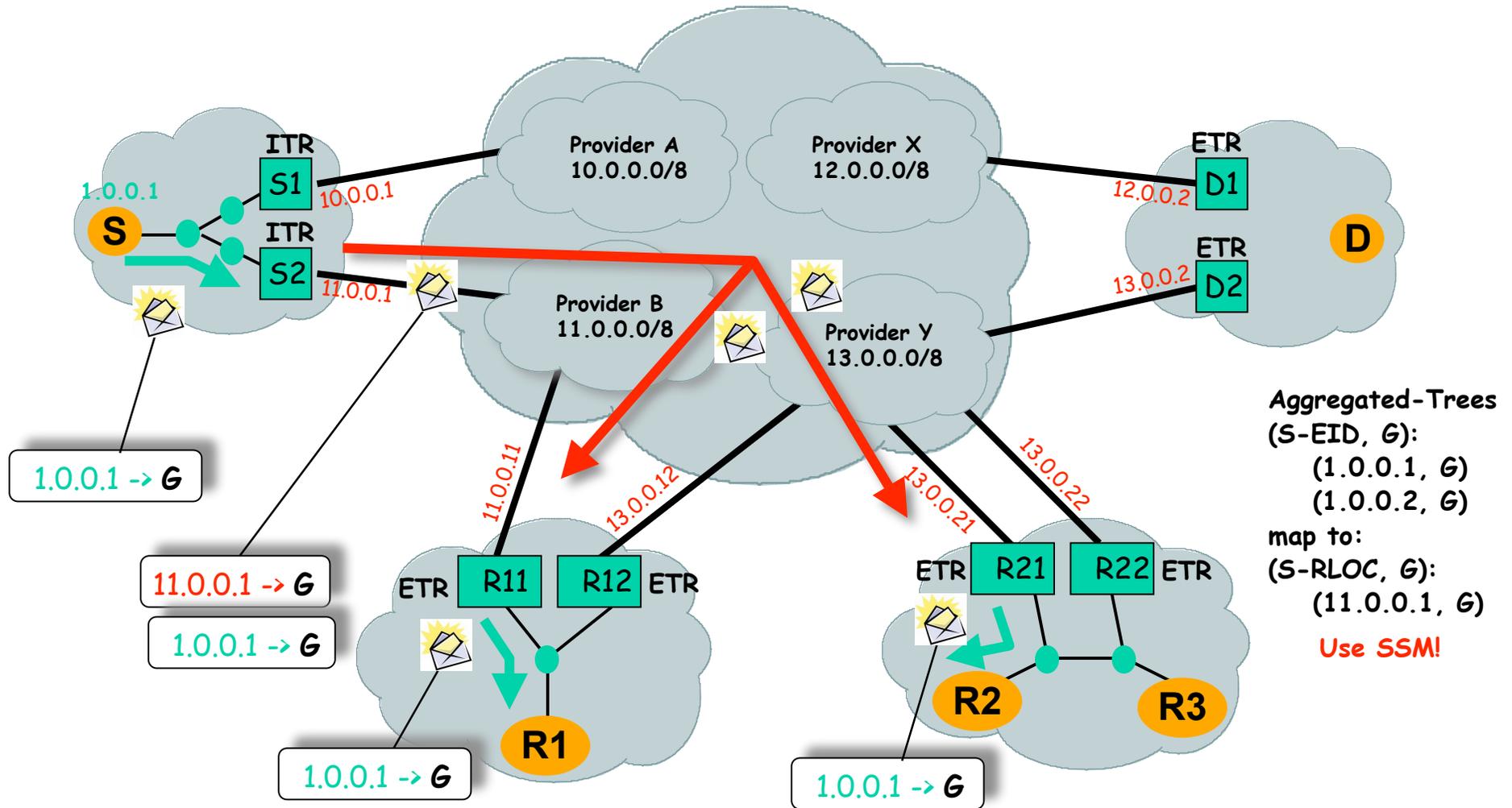
LISP Multicast Defs

- Multicast ETRs at receiver sites
 - Receives PIM JP (S-EID, G) messages from site routers
 - Obtains S-RLOC from mapping for S-EID
 - Sends unicast PIM JP (S-EID, G) to ITR S-RLOC address
 - Sends regular PIM JP (S-RLOC, G) through core
 - Decapsulates multicast (S-RLOC, G) into (S-EID, G)
- Multicast ITRs at source sites
 - Receive unicast PIM JP (S-EID, G) messages and forward into source site
 - Acts as root of (S-RLOC, G) tree for site
 - Encapsulates (S-EID, G) packets into (S-RLOC, G)
- mPTRs
 - Multicast Proxy Tunnel Routers (PTRs)

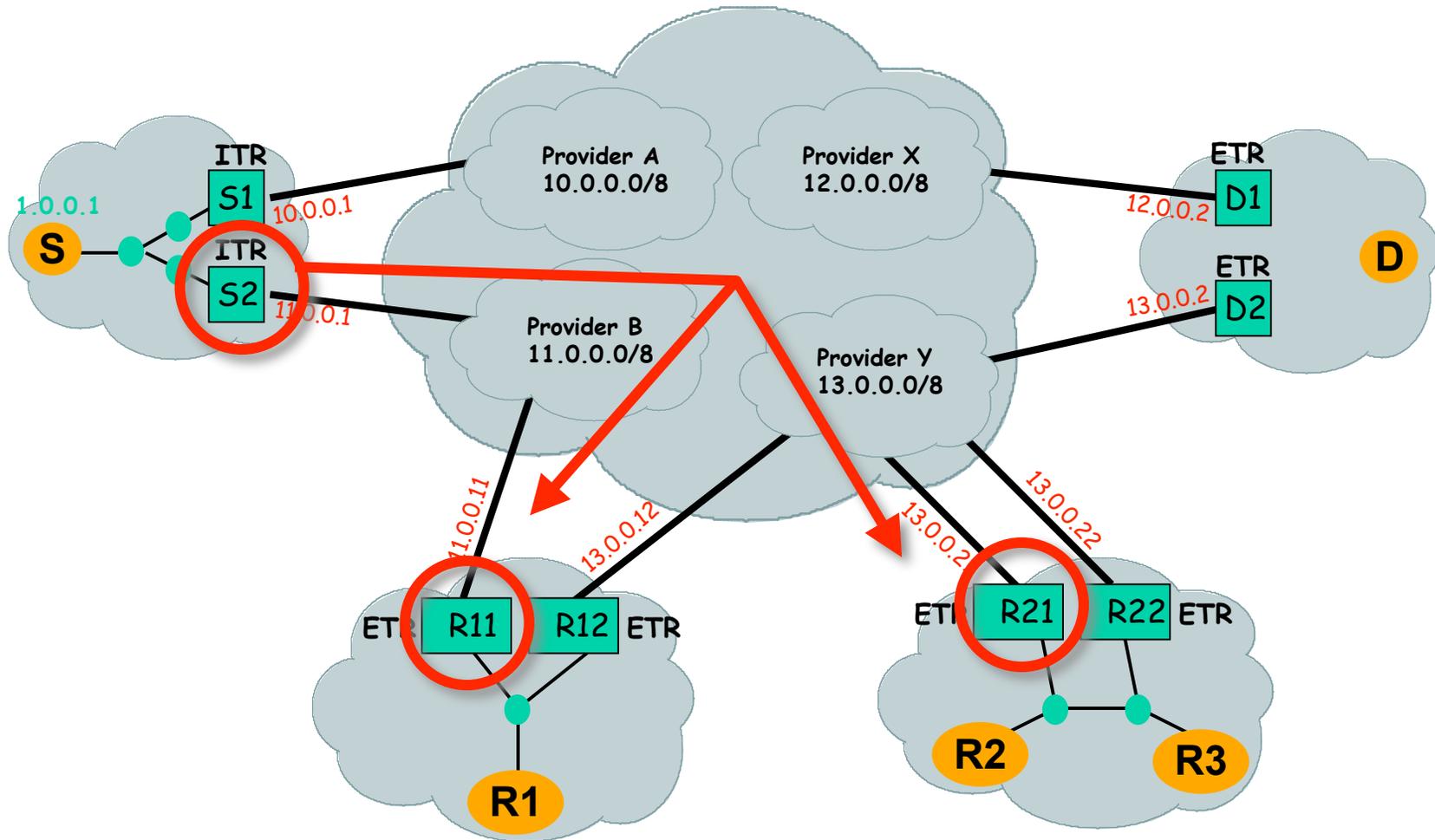
Multicast Packet Forwarding



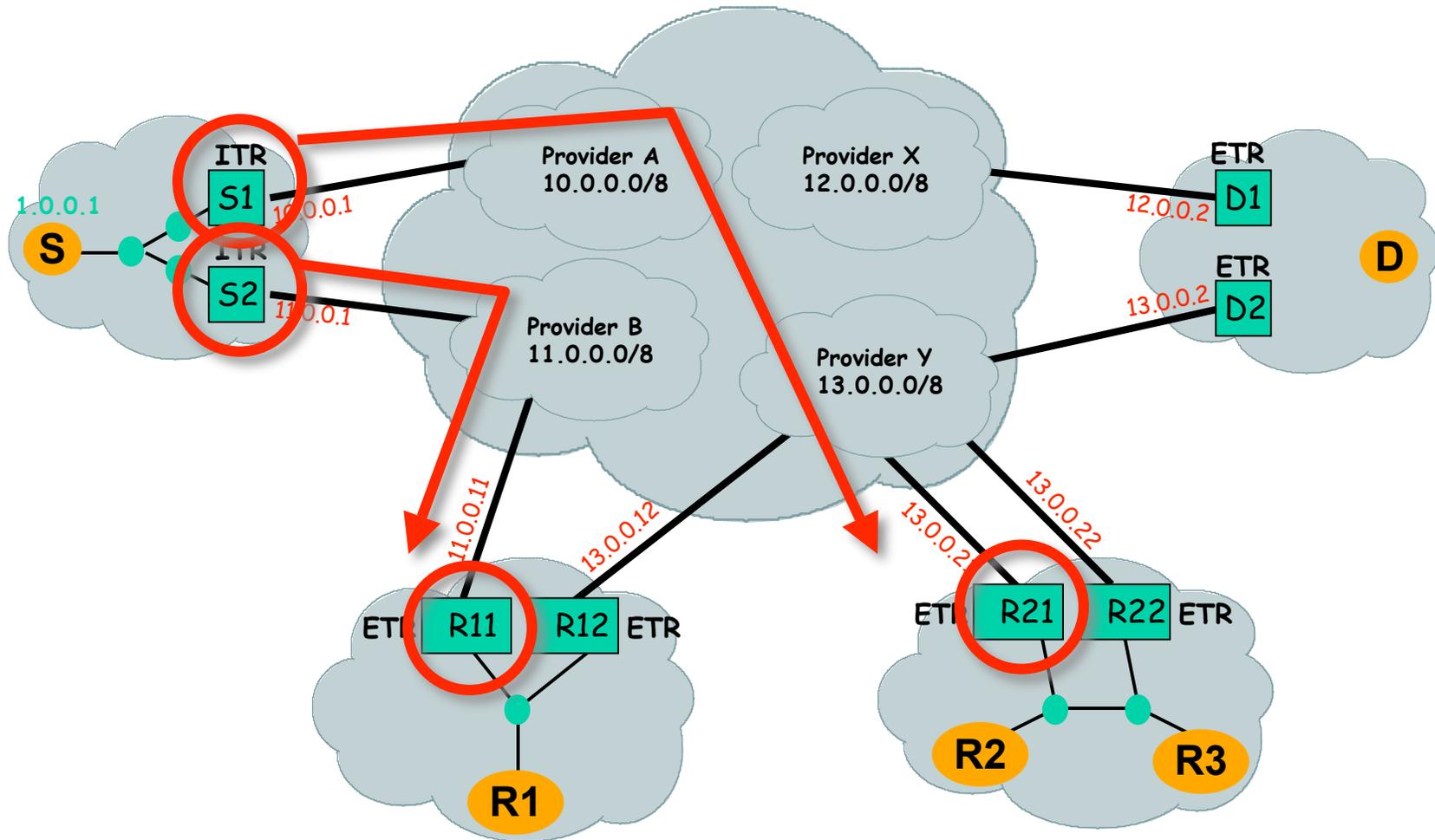
Multicast Packet Forwarding



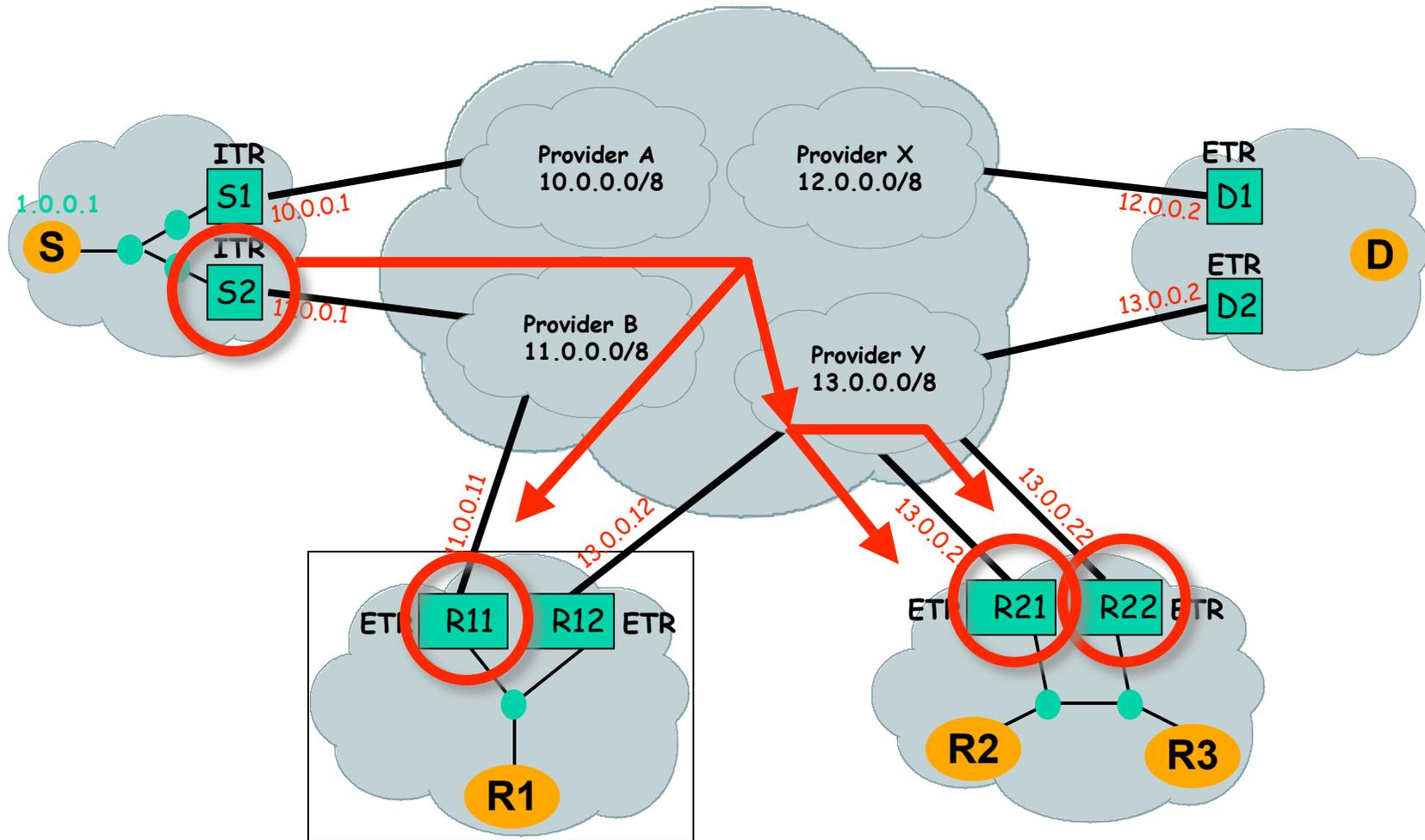
Join Policy - 1 Exit, 1 Entrance



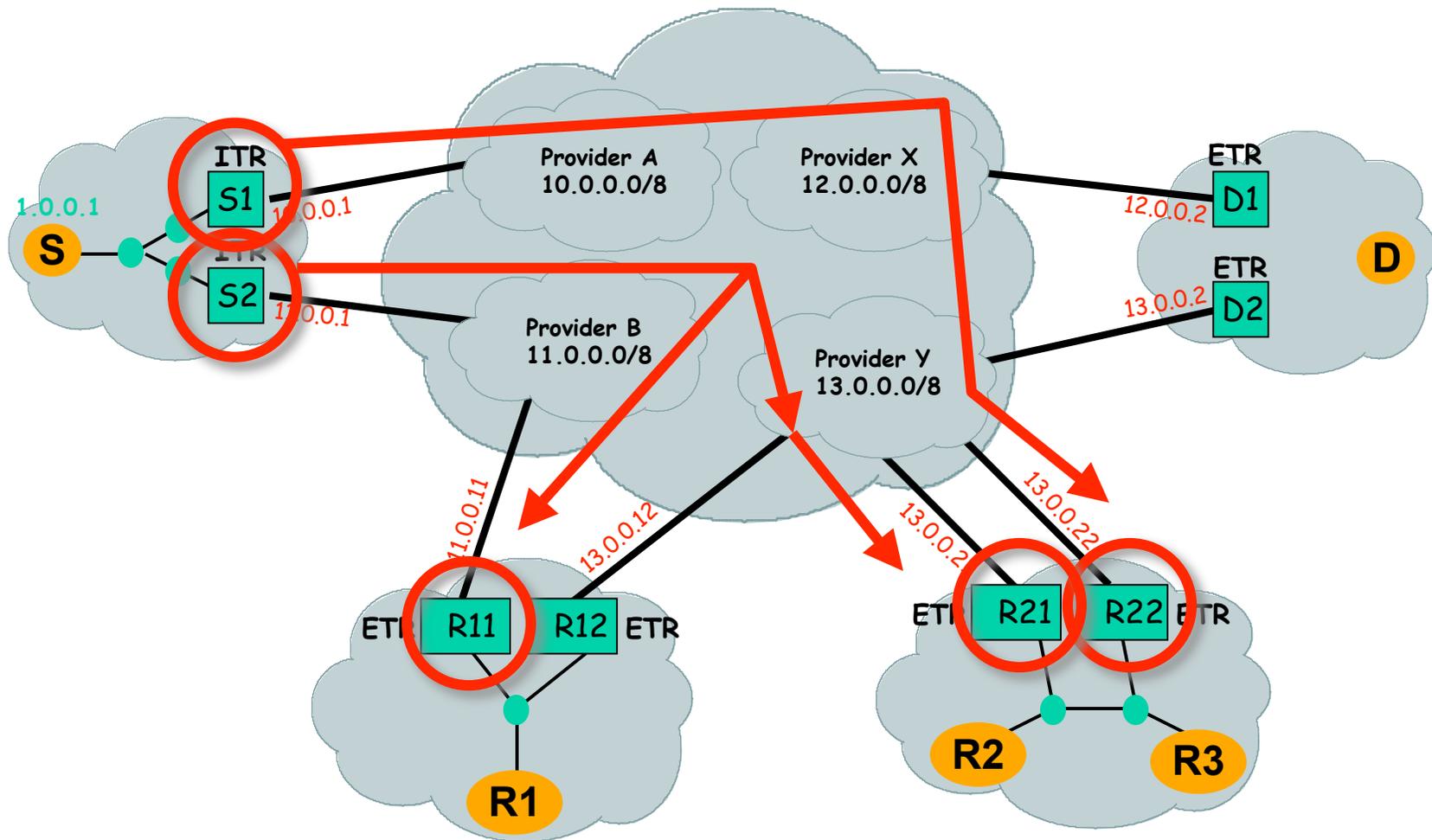
Join Policy - 2 Exits, 1 Entrance



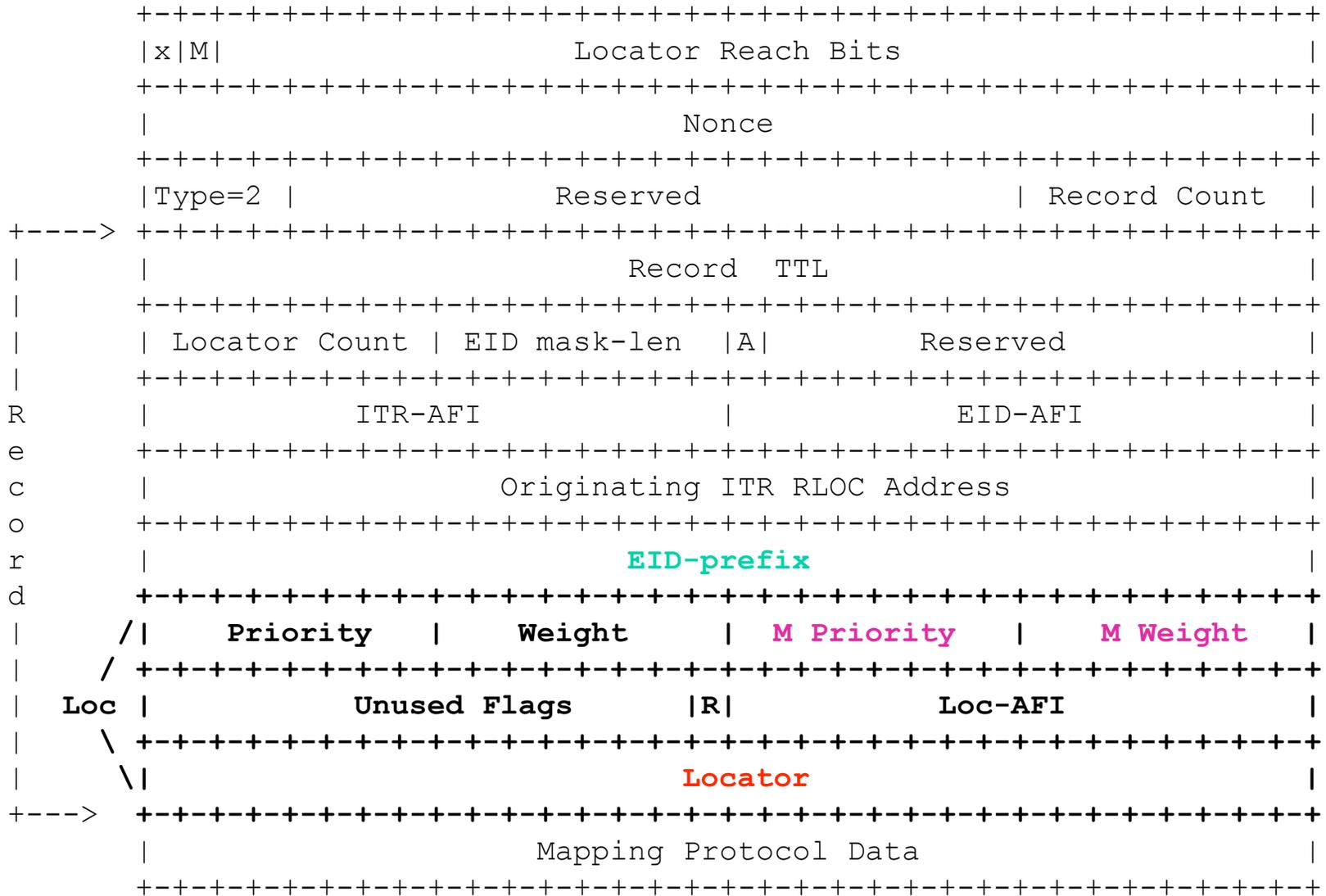
Join Policy - 1 Exit, 2 Entrances



Join Policy - 2 Exits, 2 Entrances



6.1.4. Map-Reply Message Format



Locator Reachability

- Loc-reach-bits
 - Used for both unicast and multicast
 - Mapping entries contain all locators
 - You can turn locators on for unicast and off for multicast by setting priority to 255
 - Multicast and Unicast fate-share each other
 - A mapping obtained for a unicast packet can be used for a multicast join and vice-versa

LISP-Multicast Interworking

- Non-LISP sites
 - Unicast only
 - Unicast and multicast
- uLISP sites
 - Unicast-LISP
 - Multicast-Traditional
 - No Multicast at all
- Multicast-LISP sites
 - Both Unicast-LISP and Multicast-LISP
- Simplification
 - Treat Non-LISP-multicast and uLISP-multicast-traditional the same

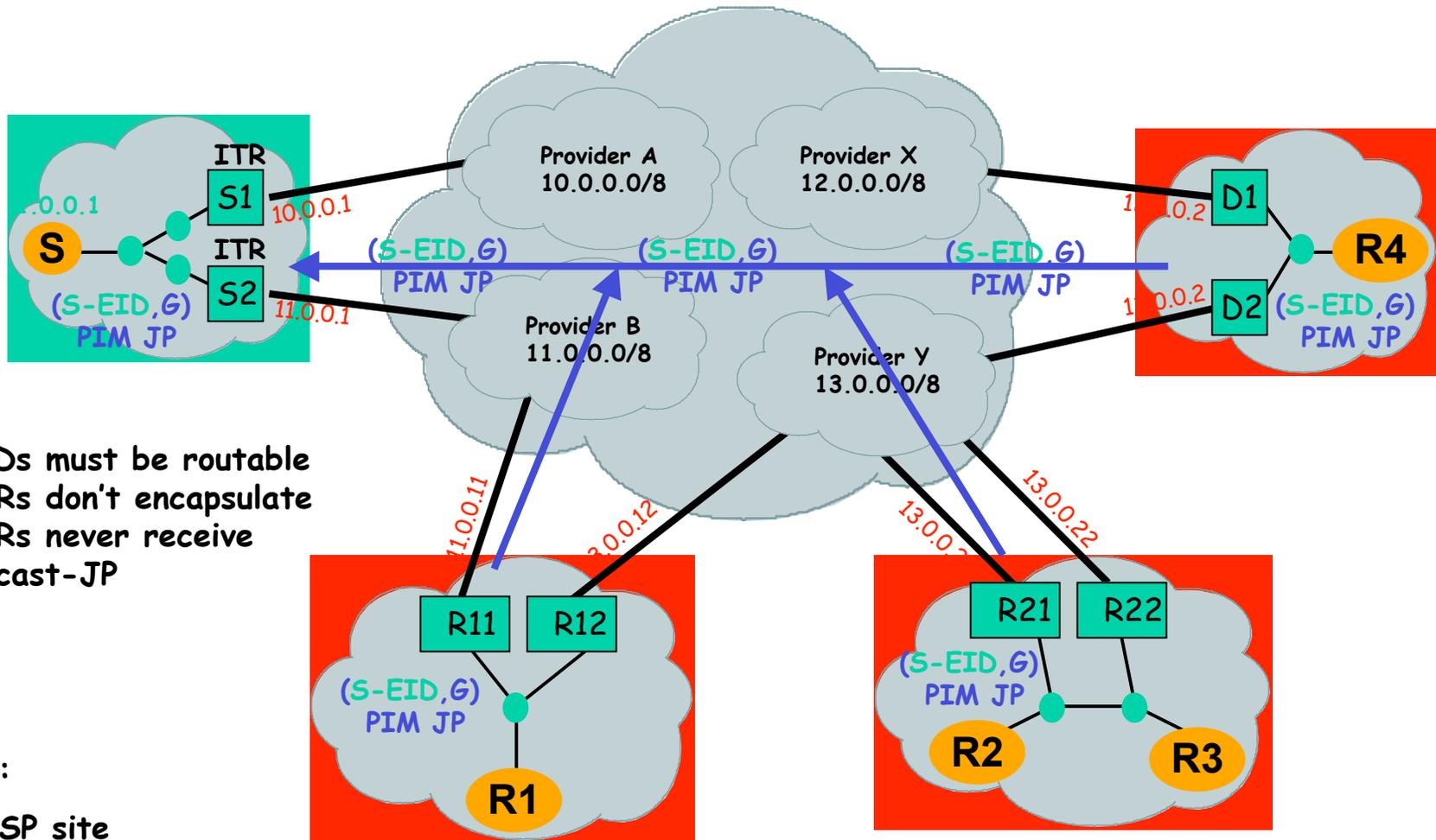
Interworking Solutions

- Unicast Interworking
 - Translation at source site
 - Proxy Tunnel Routers (PTR)
- Multicast Interworking
 - Translation at receiver site
 - Would require another mapping table
 - Receiver sites must translate to same routable (S-EID)
 - Operationally a non-starter
 - Multicast Proxy Tunnel Routers (mPTRs)
 - Reduce where (S-EID) state resides in network
 - (S-EID) state exists between non-LISP site and mPTR

Interworking Case Studies

- Non-LISP multicast source site
 - All receiver sites are LISP
 - Mix of LISP and non-LISP receiver sites
- uLISP source site -- same as above
- LISP multicast source site
 - All receiver sites are non-LISP
 - All receiver sites are LISP
 - Mix of LISP and non-LISP

LISP-Multicast Interworking

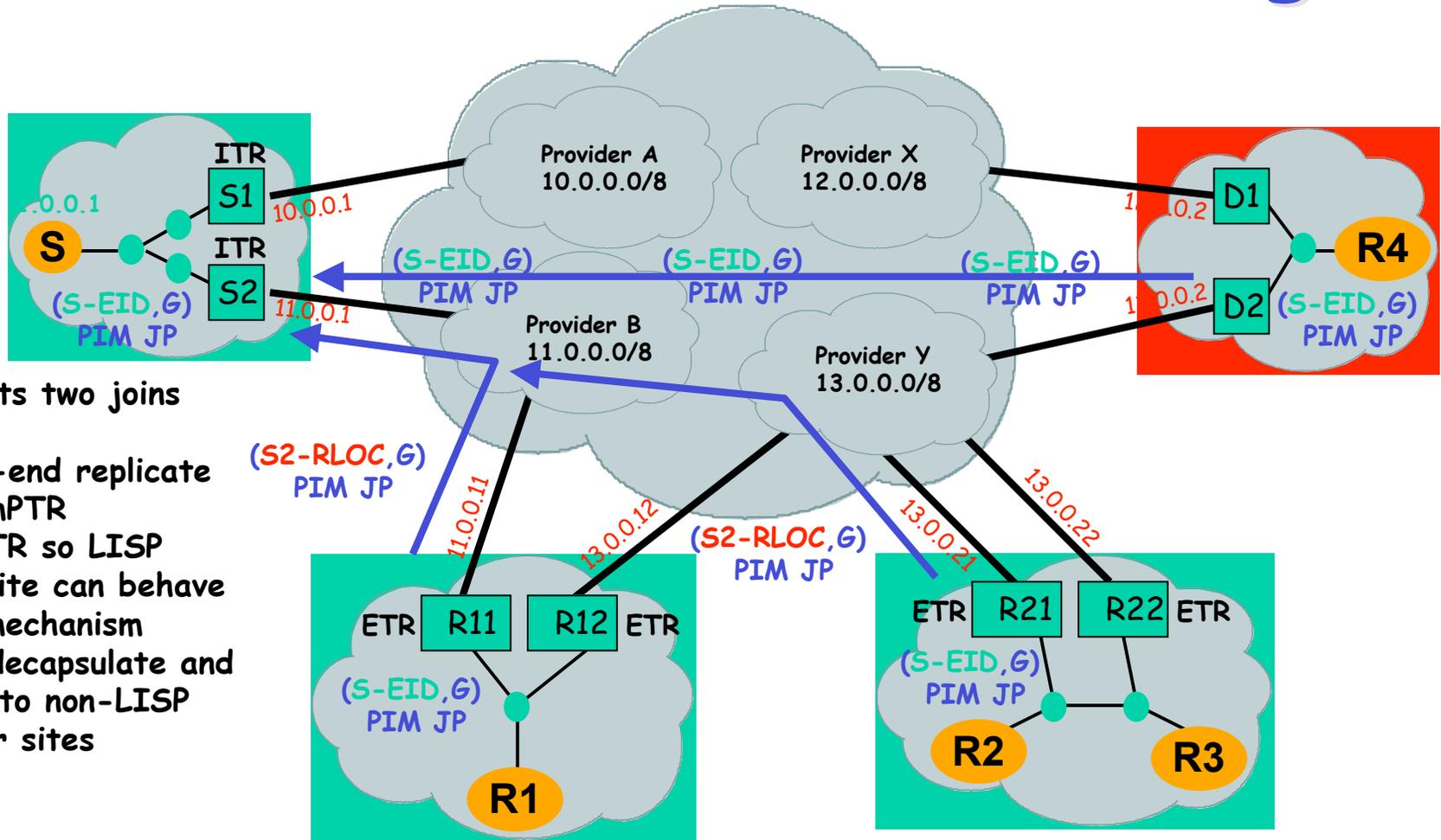


- 1) EIDs must be routable
- 2) ITRs don't encapsulate
- 3) ITRs never receive unicast-JP

Legend:

-  LISP site
-  non-LISP site

LISP-Multicast Interworking



- 1) ITRs gets two joins
Options:
Head-end replicate
Use mPTR
- 2) Use mPTR so LISP source site can behave with 1 mechanism
- 3) mPTRs decapsulate and deliver to non-LISP receiver sites

Legend:

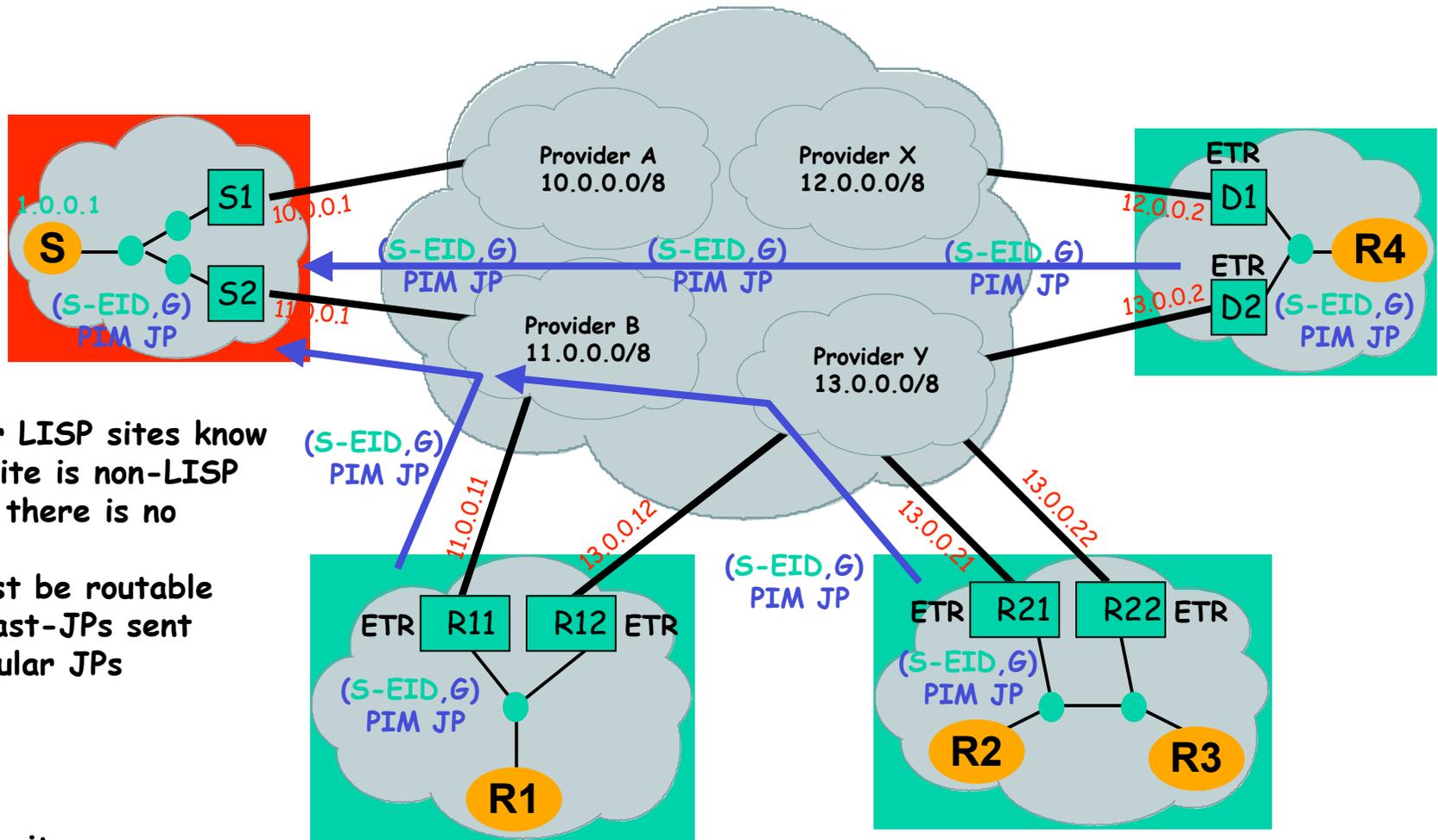
-  LISP site
-  non-LISP site

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LISP-Multicast Interworking



- 1) Receiver LISP sites know source site is non-LISP because there is no mapping
- 2) EID must be routable
- 3) No unicast-JPs sent just regular JPs

Legend:

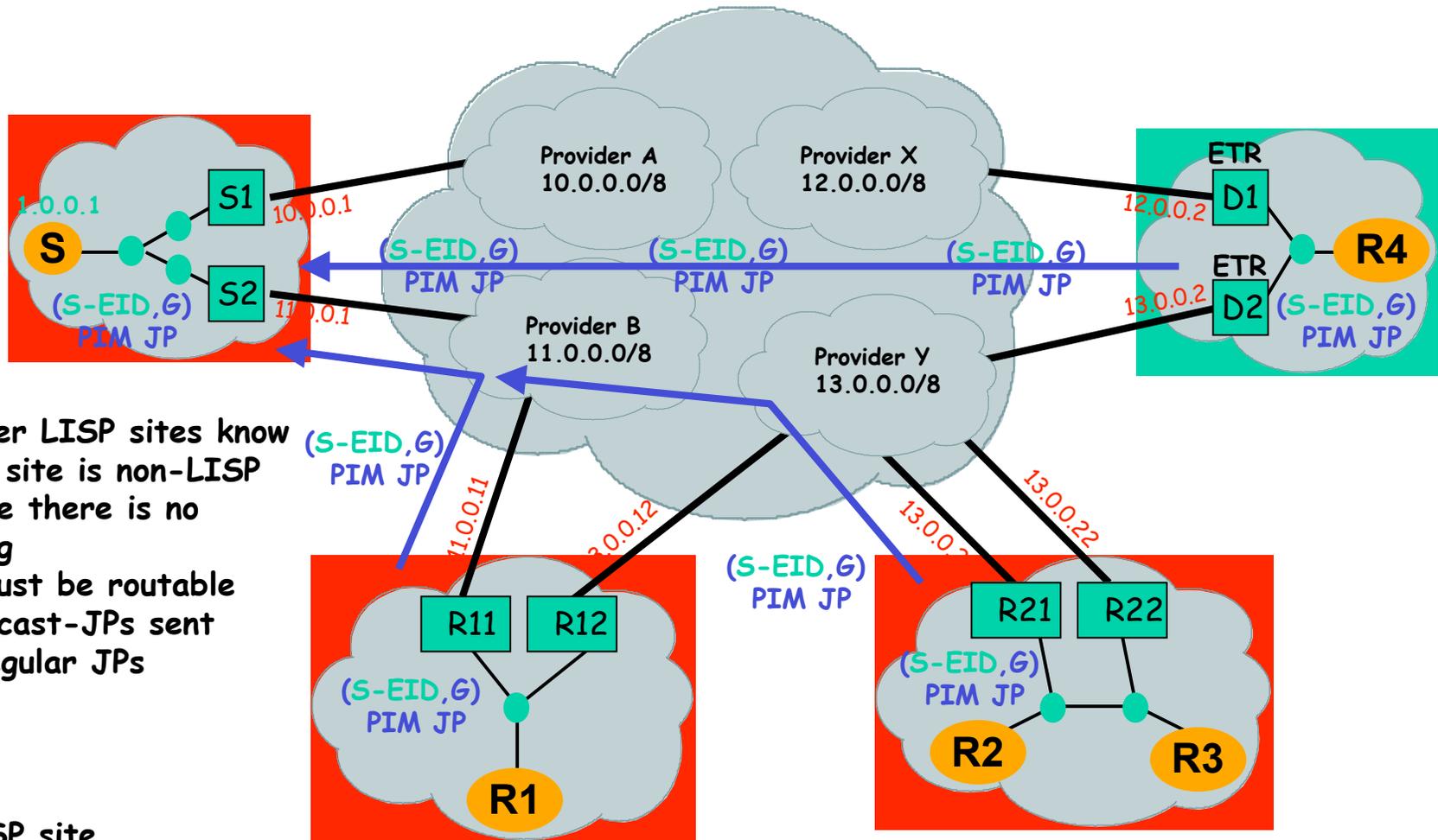
-  LISP site
-  non-LISP site

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LISP-Multicast Interworking



- 1) Receiver LISP sites know source site is non-LISP because there is no mapping
- 2) EID must be routable
- 3) No unicast-JPs sent just regular JPs

Legend:



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Multicast PTRs

- Unicast PTRs
 - They proxy for ITR functionality
 - They advertise coarse prefixes to attract and encapsulate packets
- Multicast PTRs
 - They proxy for ETR functionality
 - They advertise coarse prefixes to attract Joins and therefore attract and decapsulate packets
- No unicast EID state in core from PTRs to ETRs
- No multicast (S-EID, G) state in core from ITRs to mPTRs

AF Case Studies

1. LISP-Multicast IPv4 Site
2. LISP-Multicast IPv6 Site
3. LISP-Multicast Dual-Stack Site
4. uLISP IPv4 Site
5. uLISP IPv6 Site
6. uLISP Dual-Stack Site
7. non-LISP IPv4 Site
8. non-LISP IPv6 Site
9. non-LISP Dual-Stack Site

Case Studies

- A Combinatorial Explosion
 - $(9\ 1) + (9\ 2) + \dots + (9\ 9)$
 - 502 combinations!
- Rationalize by forcing multicast tree be AF homogenous
 - Avoid head-end replication and translation
 - Source site mapping entry priority decides which AF
 - When source site is non-LISP, use application-level directory to determine AF
- Further simplification
 - Make all sites that deploy LISP be unicast & multicast capable from the day-one

LISP Protocol Mechanisms

- Changes only to ITRs, ETRs, & PTRs
- Simple change to PIM
 - Sending/receiving unicast JPs
 - Translation of JP from EID to RLOC
- Simple change to MSDP
 - Translation for RPF-peering for RP
- Simplification
 - Only support SSM for inter-domain
- Use mPTRs to reduce (S-EID, G) state
 - Using (S-RLOC, G) reduces multicast routing table size we have today!

Best Possible Outcome

