MIKEY-TICKET

DRAFT-MATTSSON-MIKEY-TICKET-00

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ORIGIN OF IDEA FOR MIKEY-TICKET

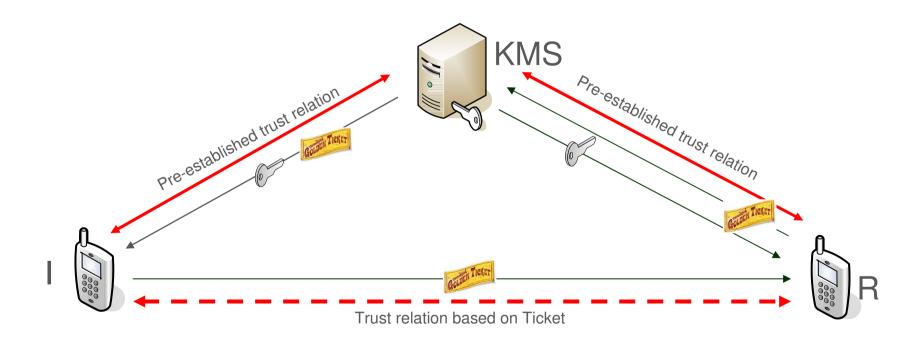
IMS media plane security (3GPP) for users requiring high security, e.g. enterprises and national security and public safety organizations.

Requirements

- Anchoring of trust/security external to IMS/SIP possible, i.e. trust/security is independent of trust/security in operator.
- Central policy control.
- Scalable and efficient.
- Group key management.
- Secure forking late binding of keys to users.
- Remote-end user identity assurance.
- Pre-distribution of tickets.
- Support of deferred delivery.

SOLUTION BY TICKET BASED KEY MANAGEMENT

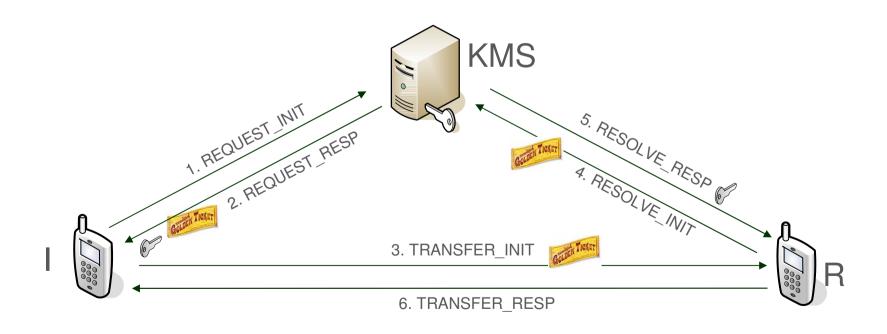
Making the solution an extension to MIKEY should be credited to Lakshminath Dondeti who suggested this during the 3GPP work by.



The KMS resolves the ticket and modifies the keys, making them cryptographically unique for each responder targeted by forking.

MIKEY-TICKET

- > Extends MIKEY with a set of new modes that uses a trusted KMS (Key Management Service) and a ticket concept, similar to Kerberos.
- MIKEY-TICKET consists of up to three roundtrips
 - Ticket Request, Ticket Transfer, Ticket Resolve



TICKET REQUEST

- This exchange is used by the Initiator to request keys and a ticket from a trusted Key Management Service, with which the Initiator have a preestablished trust relation.
- The initiation message REQUEST_INIT comes in two variants corresponding to the pre-shared key (PSK) and public-key encryption (PK) methods of [RFC3830].
- TP is the desired ticket policy.

TICKET TRANSFER

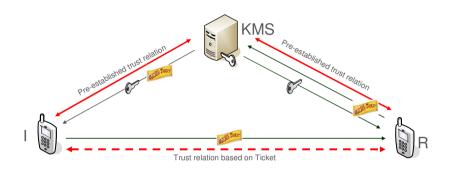
- This exchange is used to transfer the ticket as well as session information from the Initiator to a Responder.
- Similar to MIKEY-PSK but
 - TICKET instead of KEMAC
 - RANDi and RANDr gives mutual key freshness guarantee
 - IDRr, RANDkms is used to modify keys, making them cryptographically unique for each responder targeted by the forking.

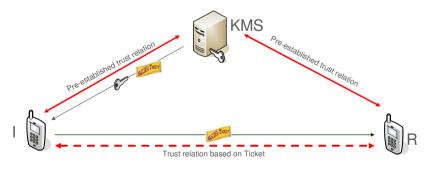
TICKET RESOLVE

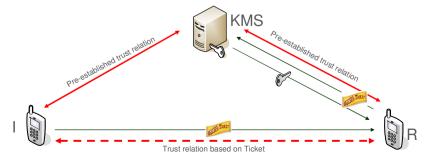
- This exchange is used by the Responder to request the KMS to return the keys encoded in a ticket.
- The initiation message RESOLVE_INIT comes in two variants corresponding to the pre-shared key (PSK) and public-key encryption (PK) methods of [RFC3830].

```
Responder
                                       KMS
RESOLVE INIT PSK =
HDR, T, RANDr, [IDRr],
   [IDRkms], TICKET,
                               <---- RESOLVE RESP
   [IDRpsk], V
                                       HDR, T, [IDRkms], KEMAC,
                                             [IDRr], [RANDkms], V
RESOLVE INIT PK =
HDR, T, RANDr, [IDRr], {CERTr},
   [IDRkms], TICKET,
                               <---- RESOLVE RESP
   [CHASH], PKE, SIGNr
                                      HDR, T, [IDRkms], KEMAC,
                                             [IDRr], [RANDkms], V
```

3 MAIN MODES OF OPERATION







Full three roundtrip

- Ticket generated by KMS
- Policy enforcement when ticket is generated
- Allows pre-distribution of tickets

Kerberos like

- Ticket generated by KMS
- Policy enforcement when ticket is generated
- Cannot handle secure forking/retargeting

Otway-Rees like

- Ticket generated by Initiator
- Policy enforcement when ticket is resolved

KEY FEATURES

- Trust is anchored in KMS, which can be independent of network operator.
- Allowed message exchanges and ticket options are determined by policy and can be adapted for different deployment scenarios: Policy enforcement is performed by the KMS.
- > By defining groups of recipients, group key management becomes simple; 3GPP allows wildcarding for user group definitions. KMS verifies group membership.
- Secure forking and assurance of remote-end user identity
- "Reusable" tickets can be pre-distributed.

3GPP STATUS

- Specification is being finalized.
- MIKEY-TICKET based key management for high security applications.
- > SDES based key management for other applications.

