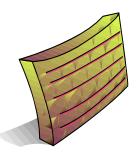
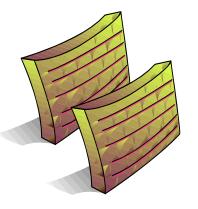


Where are we?

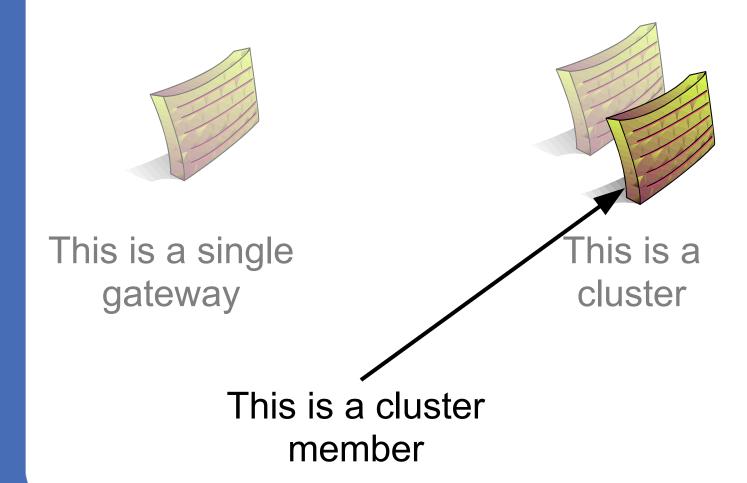
- draft-ietf-ipsecme-ipsec-ha accepted as a WG item early this year
- Now at version -09
- Draft was approved by the IESG on 15-Jul
- Now in RFC Editor's queue
- Let's go over the issues and the terminology



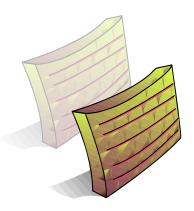
This is a single gateway



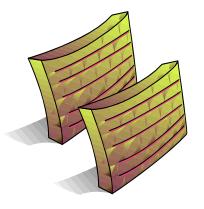
This is a cluster



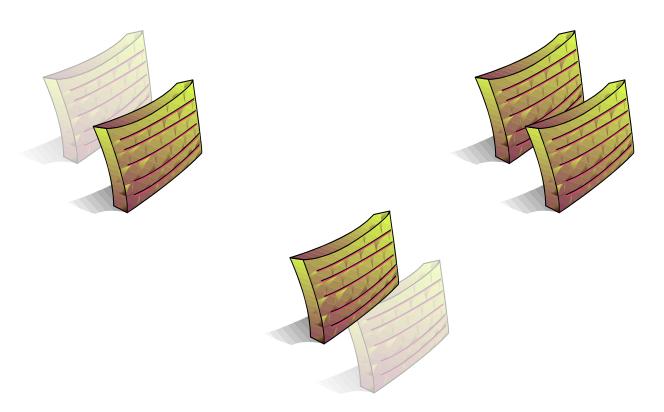
- Availability portion of the time a system can do its work. Expressed as percentage or "nines"
- High Availability the property of a system where the down time is low.
- Fault Tolerance a property of a system where functionality is maintained even following a specified set of fault condition.



This is a hot standby cluster



This is a loadsharing cluster



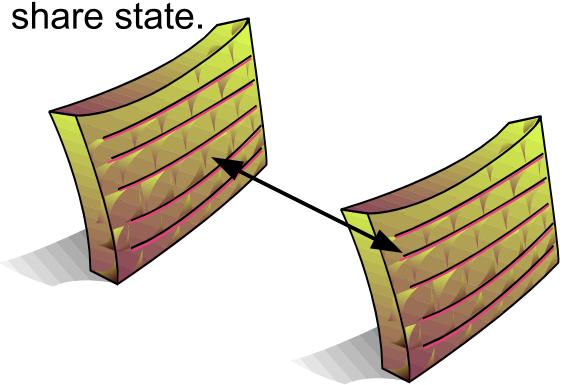
This is a hot standby cluster following a failover

- Failover is when a part of the load goes from one cluster member to another.
- In HS cluster a standby member becomes the active member, and the formerly active member either becomes a standby member, or is out of commission.
- In LS the decision function changes.
 - So the handling of a certain peer, SA, or selctor migrates.
 - Or one of the members is out of commission.

- Tight Cluster a cluster where all the members share an IP address.
- Loose Cluster a cluster where the members don't share an IP address.
 - They may share a DNS name
 - They may use RFC 5685 redirect to send traffic to the correct gateway.

- Tight Cluster a cluster where all the members share an IP address.
- Loose Cluster a cluster where the members don't share an IP address.
 - They may share a DNS name
 - They may use RFC 5685 redirect to send traffic to the correct gateway.
- We don't care about loose clusters.
 - They're out of scope.

 Synch channel is the means by which cluster members communicate in order to share state



- Out of scope:
 - How the synch channel work
 - Synchronizing policy
- In scope:
 - Any behavior of a cluster, that may appear different from that of a single gateway.
 - Any altered behavior following a failover.

- Lots of state (section 3.2):
 - IKE SAs
 - Keys
 - Authentication Information
 - IPsec SAs
 - With replay counters
 - SPD Cache entries

- IKE Counters (section 3.3)
 - An implementation MUST keep careful track of Message Ids, both inbound and outgoing.
- Synch after every IKE exchange?

- Outbound SA counter (section 3.4):
 - MUST NOT reuse a replay counter value.
- Synch after every IPSec packet?
 - Not feasible!
- Synch occasionally?
 - State will mismatch with peer after failover.
- Does the peer actually enforce this?

- Reminder:
 - IKE Message Counters
 - MUST NOT repeat
 - MUST NOT skip
 - MUST process in order
 - IPSec Replay Counter
 - MUST NOT repeat
 - May skip as much as you want
 - Enforcement is OPTIONAL.
 - If you enforce, MUST NOT process outside window

- Inbound SA Counter (section 3.5):
 - Like the previous problem, only causes a security vulnerability
 - Should not accept a packet with an old replay number.
- Synch after every packet?
 - Not practical and you might still miss.
- Live with it, assuming an attacker can't both replay and cause/detect a failover?
 - After all, enforcement is OPTIONAL.

- Missing Synch Messages (section 3.6):
 - No transport is 100% reliable.
 - If failover happened, there's a good chance some synch messages are missing.
- We have to assume that our state is mismatched with the peer's.
 - Maybe there's an SA we don't know about.
 - Maybe an SA was deleted.

- Simultaneous use of IKE or IPSec SA by more than one member (section 3.7):
 - Relevant for LS cluster
 - Replay counters cannot synch.
- Solutions fall into two broad categories:
 - "Sticky" only one member handles a particular class of traffic, so no shared SA.
 - "Duplicate" Similar SAs, one for each member with the same peer.
- Also a problem choosing distinct IVs.

- Overloading the load balancer (section 3.8)
 - We'd like the IPsec SA to directly to the member, bypassing the load balancer.
- draft-arora-ipsecme-ikev2-alt-tunneladdresses addresses this.
- Later on, we'll talk about whether this is interesting for the WG.

- Allocation of SPIs (section 3.9):
 - SPIs for inbound SAs MUST be distinct.
 - Members MUST NOT create two SAs with the same SPI, at least not with the same peer.
 - Do we really need a protocol extension to solve this?
 - We think not

- That's it for the problems.
- We'll come back to these slides when we discuss the solutions later.