

Kerberos Security Model for SNMPv3



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- ❑ Introduction: Kerberos Security Model for SNMPv3
- ❑ Why we need a new security model
- ❑ Use cases driving this proposal
- ❑ New requirements for security model
- ❑ Proposed security model
- ❑ Elements of Procedure
- ❑ Next steps



□ About the authors

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□ Previous Submission: draft-hornstein-snmpv3-ksm-00

- Ken Hornstein & Wes Hardaker, June 25, 1999.

❑ Untrusted Managed Devices

- Examples: Modems, Set Top Boxes, Home Routers.
- They can be tampered with because they are physically located in customer's homes.
 - ✓ It may be possible for an attacker to replace and spoof one of these devices.
- Any globally sensitive data sent to them may be compromised.
 - ✓ Example: SNMP administrator's SSH's username and password.

❑ Low end Managed Devices

- Examples: Modems, Set Top Boxes, Home Routers.
- They may not have the math processing capabilities to do PK operations quickly.
- They may not be able to maintain session state due to memory limitations.

❑ Large numbers of Managed Devices

- Examples: there are millions such devices deployed in North America.
- Devices will be periodically queried to retrieve device health & traffic load values.
- Automated Managers will poll multiple devices per second.
- Human administrators will access multiple devices while troubleshooting.

- ❑ USM has its own local table of users.

- ❑ RFC5592 + RFC5608 requires:
 - Use SSH to establish a secure session between Network Management Application to the SNMP Engine/RADIUS Client.
 - SSH may outsource the validation of a user's password via a local RADIUS client to a RADIUS server.
 - Upon successful authentication, SNMP stack may receive the *groupName*.
 - This model requires the Network Management Application and SNMP Engine to form a SSH session.

❑ Centralized Security Administration.

- For authentication of Kerberos users (device administrators)
 - ✓ Authentication is handled without interaction with the managed device
- For authorization of SNMP users (device administrators)
 - ✓ Addressed the same way as the I-D.ietf-isms-radius-vacm draft

❑ Strong Authentication (using two factor mechanisms.)

- Enterprises typically require this for accessing sensitive Managed Devices.
- Hardware security tokens sometimes require additional interactions with the user.
 - ✓ Not explicitly addressed by RFC5608, but could be extended.

❑ Convenience

- Each subsequent device does not require user re-authentication.

❑ Efficiency

- Does not require Managed Devices to save state between SNMP requests.
- Does not require Managed Devices to perform excessive computations.
- Minimizes the setup overhead before sending request.

□ Architectural placement of KSM

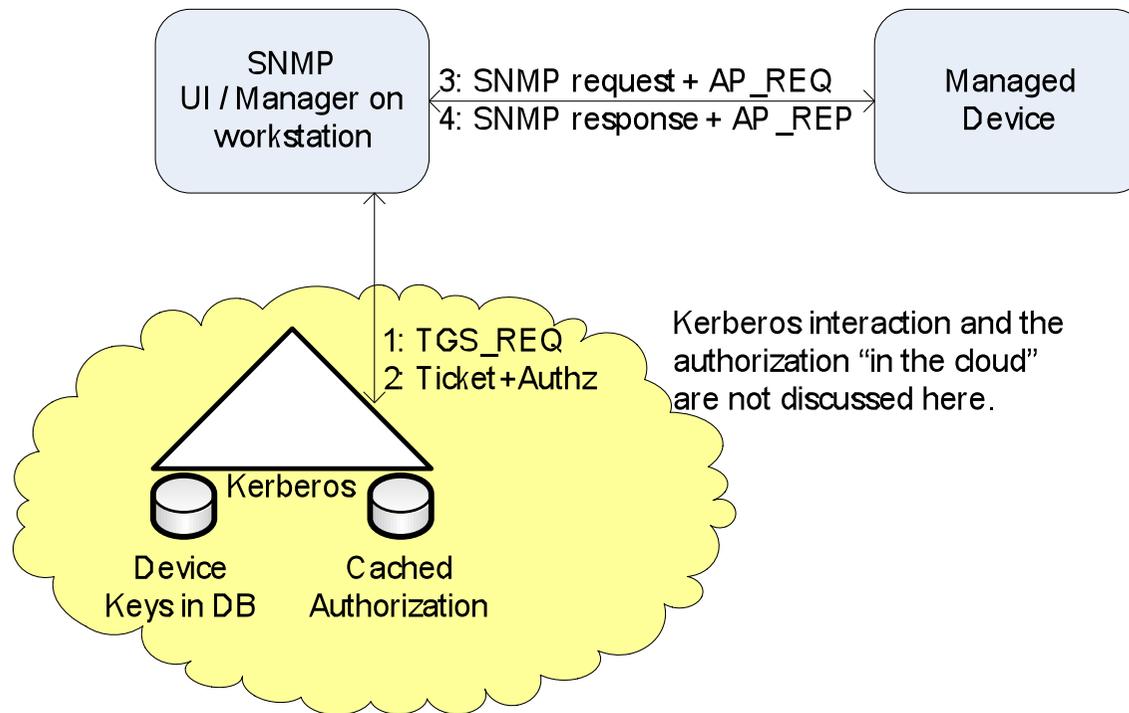
- This model is a peer to USM in the SNMP architecture.
- It uses VACM, and does not require any modifications to it.
- It uses VACM just like I-D.ietf-isms-radius-vacm does.
- It does not use or rely on any transport models.

□ Dependencies

- This model requires a Kerberos KDC server.
- It uses an Authorization Database for centralized authorization mappings.
 - ✓ Specifically, it maps securityName → groupName.
 - ✓ For example:
 - Jack → ConfigurationMgr; Jill → Auditor; Joe → Assistant;
 - Jack gets write access; Jill gets read access; Joe gets nothing.
 - ✓ The groupName may also be thought of as a role, permissions, ...
 - ✓ This value of groupName must be recognized by the Command Responder.
- The KDC and Authorization Database will not be discussed here.

❑ SNMPv3 Headers:

- *securityModel* must contain a new value indicating KSM.
- *securityParameters* must contain *ksmSecurityParameters*.
 - ✓ *ksmSecurityParameters* must contain Kerberos AP_REQ or AP_REP.
- *securityLevel* must contain *noAuthNoPriv*, *authNoPriv*, or *authPriv*.



❑ ksmSecurityParameters

```
ksmSecurityParameters ::= SEQUENCE {  
    -- The Kerberos 5 checksum type used to checksum this message  
    ksmChecksumType  INTEGER(0..2147483647),  
    -- The actual keyed checksum data returned by Kerberos  
    ksmChecksum      OCTET STRING,  
    -- The Kerberos 5 message (AP_REQ or AP_REP)  
    ksmKerberosMsg   OCTET STRING  
}
```

- Message is encrypted when the securityLevel is *authPriv*
 - ✓ *scopedPDU* is encrypted, resulting in a KRB_PRIV message.
- All messages are Integrity protected, except for *noAuthNoPriv* messages.
 - ✓ The entire message, including the SNMPv3 header, is protected.
 - ✓ Kerberos 'checksums' are actually keyed hashes, described in RFC 3961.

❑ KSM notes

- Timeliness & replay detection are addressed by KRB_PRIV methods.
- securityNames for users and devices must be Kerberos Principal names.
 - ✓ Example: joe@example.com
- Each request and response must carry a Kerberos message (AP_REQ/P).

□ Procedure for Outgoing Requests

- Command Generator contacts the KDC server to retrieve the Kerberos ticket. The ticket contains the *groupName* and *securityName*.
- Command Generator hashes the SNMP's PDU and creates the *ksmChecksum*.
- Command Generator creates the *ksmSecurityParameters* and sends the request to the Command Responder.

□ Procedure for Incoming Requests

- Command Responder extracts the kerberos ticket, decrypts the PDU and extracts the *groupName* and *securityName* from the ticket.
- Command Responder creates an entry in *vacmSecurityToGroupTable*:
 - ✓ *vacmSecurityModel* is *KSM*
 - ✓ *vacmSecurityName* is the extracted principle and realm (joe@example.com)
 - ✓ *vacmGroupName* is the extracted value
 - ✓ *vacmSecurityToGroupStorageType* is "volatile"
 - ✓ *vacmSecurityToGroupStatus* is "active"

- Status
 - draft-pejaver-isms-kerberos-01 was published.
 - ✓ It needs more work.
 - Issues are open for discussion.

- Demo of sample implementation.

- Adopt KSM as a ISMS Working Group item

Discussion

