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Transporting PTP messages (1588) over MPLS Networks  
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Abstract

This document defines the method for transporting PTP messages (PDUs) over an MPLS network to enable a proper handling of these packets (e.g. implementation of Transparent Clocks (TC)) in LSRs.

The basic idea is to transport PTP messages inside dedicated MPLS LSPs. These LSPs only carry PTP messages and possibly Control and Management packets, but they do not carry customer traffic.

Two methods for transporting 1588 over MPLS are defined. The first method is to transport PTP messages directly over the dedicated MPLS LSP via UDP/IP encapsulation, which is suitable for IP/MPLS networks. The second method is to transport PTP messages inside a PW via Ethernet encapsulation, which is more suitable for MPLS-TP networks.

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The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC2119 [RFC2119].

When used in lower case, these words convey their typical use in common language, and are not to be interpreted as described in RFC2119 [RFC2119].

## 1. Introduction

The objective of Precision Time Protocol (PTP) is to synchronize independent clocks running on separate nodes of a distributed system. [IEEE] defines PTP messages for clock and time synchronization. The PTP messages include PTP PDUs over UDP/IP (Annex D & E of [IEEE]) and PTP PDUs over Ethernet (Annex F of [IEEE]). This document defines mapping and transport of the PTP messages defined in [IEEE] over MPLS networks.

PTP defines intermediate clock functions (called transparent clocks) between the source of time (Master) and the Slave clocks. Boundary Clocks (BC) form Master-Slave hierarchy with the Master clock as root. The messages related to synchronization, establishing the Master-Slave hierarchy, and signaling, terminate in the protocol engine of a boundary clock and are not forwarded. Management messages however, are forwarded to other ports on the boundary clock.

Transparent clocks modify a "correction field" (CF) within the synchronization messages to compensate for residence and propagation delays. Transparent clocks do not terminate synchronization, Master-Slave hierarchy control messages or signaling messages.

There is a need to transport PTP messages over MPLS networks. The MPLS network could be a transit network between 1588 Masters and Slaves. The accuracy of the recovered clock improves and the Slave logic simplifies when intermediate nodes (e.g. LSRs) properly handle PTP messages (e.g. perform TC), otherwise the jitter at the 1588 Slave may be excessive and therefore the Slave may not be able to properly recover the clock and time of day.

This document defines a "1588-aware LSR" that is able to identify 1588 timing flows carried over MPLS.

Transparent Clock (TC) function requires a 1588-aware LSR in the middle of an LSP to identify the PTP messages and perform proper update of the CF, via a 1-step or 2-step process.

More generally this document requires that an LSR should be able to properly handle the PTP messages. For instance for those cases when the TC function is not viable (e.g. due to layer violation) as an alternative it should be possible to instead control the delay for these messages on both directions across the node.

In the above cases it is beneficial that PTP packets can be easily identified when carried over MPLS.

This document provides two methods for transporting PTP messages over

MPLS. The main objectives are for LSRs to be able to deterministically detect and identify the PTP messages.

## 2. Terminology

1588: The timing and synchronization as defined by IEEE 1588

PTP: The timing and synchronization protocol used by 1588

Master: The Source of 1588 Timing and clock

Slave: The Destination of 1588 Timing and clock that tries to follow the Master clock

OC: Ordinary Clock

TC: Transparent Clock, a time stamping method applied by intermediate nodes between Master and Slave

BC: Boundary Clock, is a node that recovers the Master clock via a Slave function and uses that clock as the Master for other Slaves

PTP LSP: An LSP dedicated to carry PTP messages

PTP PW: A PW within a PTP LSP that is dedicated to carry PTP messages.

CW: Pseudowire Control Word

LAG: Link Aggregation

ECMP: Equal Cost Multipath

CF: Correction Field, a field inside certain PTP messages (message type 0-3) that holds the accumulative transit time inside intermediate switches



### 3. Problem Statement

When PTP messages are transported over MPLS networks, there is a need for intermediate LSRs to detect such messages and perform proper processing (e.g. Transparent Clock (TC)). Note the TC processing could be in the form of 1-Step or 2-Step time stamping.

PTP messages over Ethernet or IP can always be tunneled over MPLS. However the 1588 over MPLS mapping defined in this document is applicable whenever MPLS LSRs are 1588-aware and the intention is for those LSRs to perform proper processing on these packets.

When 1588-awareness is needed, PTP messages should not be transported over LSPs or PWs that are carrying customer traffic because LSRs perform Label switching based on the top label in the stack. To detect PTP messages inside such LSPs require special Hardware (HW) to do deep packet inspection at line rate. Even if one assumes a deep packet inspection HW at line rate exists, the payload can't be deterministically identified by LSRs because the payload type is a context of the PW label and the PW label and its context are only known to the Edge routers (PEs) and LSRs don't know what is a PW's payload (Ethernet, ATM, FR, CES, etc). Even if one assumes only Ethernet PWs are permitted in an LSP, the LSRs don't have the knowledge of whether PW Control Word (CW) is present or not and therefore can't deterministically identify the payload.

Therefore a generic method is defined in this document that does not require deep packet inspection at line rate, and can deterministically identify PTP messages. The defined method is applicable to both MPLS and MPLS-TP networks.

#### 4. Dedicated LSPs for PTP messages

Many methods were considered for identifying the 1588 messages when they are encapsulated in MPLS such as by using GAL/ACH or a new reserved label. These methods were not attractive since they either required deep packet inspection and snooping at line rate or they required use of scarce new reserved label. Also one of the goals was to reuse existing OAM and protection mechanisms.

The method defined in this document can be used by LSRs to identify PTP messages in MPLS tunnels by using dedicated LSPs to carry PTP messages.

Compliant implementations MUST use dedicated LSPs to carry PTP messages over MPLS. Let's call these LSPs as the "PTP LSPs" and the labels associated with these LSPs as "PTP labels". These LSPs could be P2P or P2MP LSPs. The PTP LSP between Master and Slaves MAY be P2MP or P2P LSP while the PTP LSP between each Slave and Master SHOULD be P2P LSP. The PTP LSP between a Master and a Slave and the PTP LSP between the same Slave and Master MUST be co-routed. Alternatively, a single bidirectional co-routed LSP can be used. The PTP LSP MAY be MPLS LSP or MPLS-TP LSP.

The PTP LSPs could be configured or signaled via RSVP-TE/GMPLS. New RSVP-TE/GMPLS TLVs and objects are defined in this document to indicate that these LSPs are PTP LSPs.

We should be selective about the kind of traffic that flows over PTP LSPs as these will be handled as a special case by the LSR. The only LSP user plane traffic MUST be PTP, but the LSP MAY also carry essential MPLS/MPLS-TP control plane traffic such as BFD and LSP-Ping.

## 5. 1588 over MPLS Encapsulation

This document defines two methods for carrying PTP messages over MPLS. The first method is carrying IP encapsulated PTP messages over PTP LSPs and the second method is to carry PTP messages over dedicated Ethernet PWs (called PTP PWs) inside PTP LSPs.

### 5.1. 1588 over LSP Encapsulation

The simplest method of transporting PTP messages over MPLS is to encapsulate PTP PDUs in UDP/IP and then encapsulate them in PTP LSP. The 1588 over LSP format is shown in Figure 1.

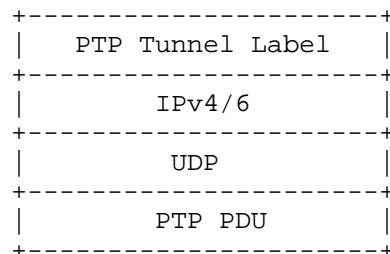


Figure 1 - 1588 over LSP Encapsulation

This encapsulation is very simple and is useful when the networks between 1588 Master and Slave are IP/MPLS networks.

In order for an LSR to process PTP messages, the PTP Label must be the top label of the label stack.

The UDP/IP encapsulation of PTP MUST follow Annex D and E of [IEEE].

### 5.2. 1588 over PW Encapsulation

Another method of transporting 1588 over MPLS networks is by encapsulating PTP PDUs in Ethernet and then transporting them over Ethernet PW (PTP PW) as defined in [RFC4448], which in turn is transported over PTP LSPs. Alternatively PTP PDUs MAY be encapsulated in UDP/IP/Ethernet and then transported over Ethernet PW.

Both Raw and Tagged modes for Ethernet PW are permitted. The 1588 over PW format is shown in Figure 2.

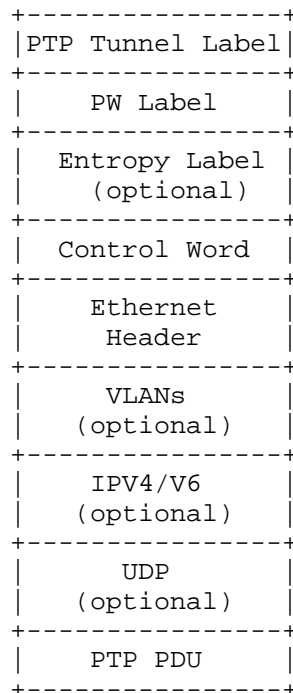


Figure 2 - 1588 over PW Encapsulation

The Control Word (CW) as specified in [RFC4448] SHOULD be used to ensure a more robust detection of PTP messages inside the MPLS packet. If CW is used, the use of Sequence number is optional.

The use of VLAN and UDP/IP are optional. Note that 1 or 2 VLANs MAY exist in the PW payload.

In order for an LSR to process PTP messages, the top label of the label stack (the Tunnel Label) MUST be from PTP label range. However in some applications the PW label may be the top label in the stack, such as cases where there is only one-hop between PEs or in case of PHP. In such cases, the PW label SHOULD be chosen from the PTP Label range.

An Entropy label [I-D.ietf-pwe3-fat-pw] MAY be present at the bottom of stack.

The Ethernet encapsulation of PTP MUST follow Annex F of [IEEE] and the UDP/IP encapsulation of PTP MUST follow Annex D and E of [IEEE].

For 1588 over MPLS encapsulations that are PW based, there are some

cases in which the PTP LSP label may not be present:

- o When PHP is applied to the PTP LSP, and the packet is received without PTP LSP label at PW termination point .
- o When the PW is established between two routers directly connected to each other and no PTP LSP is needed.

In such cases it is required for a router to identify these packets as PTP packets. This would require the PW label to also be a label that is distributed specifically for carrying PTP traffic (aka PTP PW label). Therefore there is a need to add extension to LDP/BGP PW label distribution protocol to indicate that a PW label is a PTP PW labels.

### 5.3. 1588 over pure MPLS mode

Editor Note: The encapsulation is general enough and can support transporting 1588 in a pure MPLS mode (i.e., without any IP/UDP or Ethernet headers). Should the WG pursue this?

## 6. 1588 Message Transport

1588 protocol comprises of the following message types:

- o Announce
- o SYNC
- o FOLLOW UP
- o DELAY REQ (Delay Request)
- o DELAY RESP (Delay Response)
- o PDELAY REQ (Peer Delay Request)
- o PDELAY RESP (Peer Delay Response)
- o PDELAY RESP FOLLOW UP (Peer Delay Response Follow up)
- o Management
- o Signaling

A subset of PTP message types that require TC processing are called Event messages:

- o SYNC
- o DELAY REQ (Delay Request)
- o PDELAY REQ (Peer Delay Request)
- o PDELAY RESP (Peer Delay Response)

SYNC and DELAY\_REQ are exchanged between Master and Slave and MUST be transported over PTP LSPs. PDELAY\_REQ and PDELAY\_RESP are exchanged between adjacent routers and MAY be transported over single hop PTP LSPs. If Two Step Transparent clocks are present, then the FOLLOW\_UP and DELAY\_RESP messages must also be transported over the PTP LSPs.

For a given instance of 1588 protocol, SYNC and DELAY\_REQ MUST be transported over two PTP LSPs that are in opposite directions. These PTP LSPs, which are in opposite directions MUST be congruent and co-routed. Alternatively, a single bidirectional co-routed LSP can be used.

Except as indicated above for the two-step Transparent clocks, Non-

Event PTP message types don't need to be processed by intermediate routers. These message types MAY be carried in PTP Tunnel LSPs.

## 7. Protection and Redundancy

In order to ensure continuous uninterrupted operation of 1588 Slaves, usually as a general practice, Redundant Masters are tracked by each Slave. It is the responsibility of the network operator to ensure that physically disjoint PTP tunnels that don't share any link are used between the redundant Masters and a Slave.

When redundant Masters are tracked by a Slave, any PTP LSP or PTP PW failure will trigger the slave to switch to the Redundant Master. However LSP/PW protection such as Linear Protection Switching (1:1,1+1), Ring protection switching or MPLS Fast Reroute (FRR) SHOULD still be used to ensure the LSP/PW is ready for a future failure.

Note that any protection or reroute mechanism that adds additional label to the label stack, such as Facility Backup Fast Reroute, MUST ensure that the pushed label is a PTP Label to ensure proper processing of PTP messages by LSRs in the backup path.



## 8. ECMP

To ensure the proper operation of 1588 Slaves, the physical path for PTP messages from Master to Slave and vice versa must be the same for all PTP messages listed in section 7 and must not change even in the presence of ECMP in the MPLS network.

To ensure the forward and reverse paths are the same PTP LSPs and PWs MUST not be subject to ECMP.

## 9. OAM, Control and Management

In order to manage PTP LSPs and PTP PWs, they MAY carry OAM, Control and Management messages. These control and management messages can be differentiated from PTP messages via already defined IETF methods.

In particular BFD [RFC5880], [RFC5884] and LSP-Ping [RFC4389] MAY run over PTP LSPs via UDP/IP encapsulation or via GAL/G-ACH. These Management protocols are easily identified by the UDP Destination Port number or by GAL/ACH respectively.

Also BFD, LSP-Ping and other Management messages MAY run over PTP PW via one of the defined VCCVs (Type 1, 2 or 3) [RFC5085]. In this case G-ACH, Router Alert Label (RAL), or PW label (TTL=1) are used to identify such management messages.

## 10. QoS Considerations

The PTP messages are time critical and must be treated with the highest priority. Therefore 1588 over MPLS messages must be treated with the highest priority in the routers. This can be achieved by proper setup of PTP tunnels. It is recommended that the PTP LSPs are setup and marked properly to indicate EF-PHB for the CoS and Green for drop eligibility.

11. FCS Recalculation

Ethernet FCS MUST be recalculated at every LSR that performs the TC processing and FCS retention described in [RFC4720] MUST not be used.

## 12. UDP Checksum Correction

For UDP/IP encapsulation mode of 1588 over MPLS, the UDP checksum is optional when used for IPv4 encapsulation and mandatory in case of IPv6. When IPv4/v6 UDP checksum is used each 1588-aware LSR must either incrementally update the UDP checksum after the CF update or should verify the UDP checksum on reception from upstream and recalculate the checksum completely on transmission after CF update to downstream node.

13. Routing extensions for 1588aware LSRs

MPLS-TE routing relies on extensions to OSPF [RFC2328] [RFC5340] and IS-IS [ISO] [RFC1195] in order to advertise Traffic Engineering (TE) link information used for constraint-based routing.

Indeed, it is useful to advertise data plane TE node capabilities, such as the capability for a router to be 1588-aware. This capability MUST then be taken into account during path computation to prefer nodes that advertise themselves as 1588-aware, so that the PTP LSPs can be properly handled.

For this purpose, the following sections specify extensions to OSPF and IS-IS in order to advertise 1588 aware capabilities of a node.

Editor Note: There is an open issue on whether we must consider LSRs that may not want to support PTP on all ports. An example could be an LSR where a few blades have been upgraded to support PTP timestamping in silicon. In such cases, routers must explicitly indicate the ports that are 1588-aware. If the WG agrees about this then we will need to change the subsequent OSPF and IS-IS sections to advertise the 1588-aware capability on per port/interface basis, rather than per node as is current described.

13.1. 1588aware Node Capability for OSPF

This extension makes use of the Router Information (RI) Opaque LSA defined in [RFC4970]for both OSPFv2 and OSPFv3, by defining a new OSPF Router Information (RI) TLV - The 1588-aware Capability TLV.

The 1588-aware Capability TLV is OPTIONAL and is defined as follows:

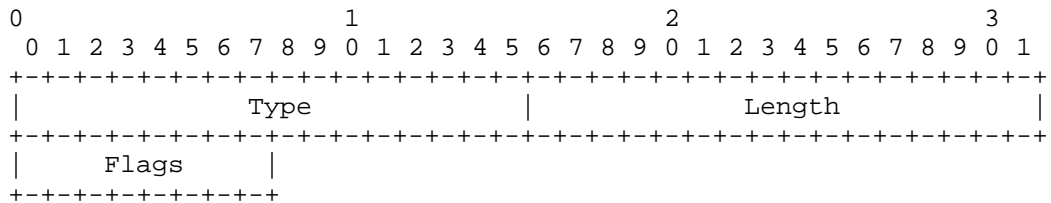


Figure 3: 1588-aware Capability TLV

Where:

Type, 16 bits: 1588-aware Capability TLV where the value is TBD

Length, 16 bits: Gives the length of the flags field in octets, and

is currently set to 1

Flags, 8 bits: The bits are defined least-significant-bit (LSB) first, so bit 7 is the least significant bit of the flags octet.

```

  0 1 2 3 4 5 6 7
+-----+
|   Reserved   |C|
+-----+
```

Figure 4: Flags Format

Correction (C) field Update field, 1 bit: Setting the C bit to 1 indicates that the node is capable of recognizing the PTP event packets and can compensate for residence time by updating the PTP packet Correction Field. When this is set to 0, it means that this node cannot perform the residence time correction but is capable of performing MPLS frame forwarding of the frames with PTP labels using a method that support the end to end delivery of accurate timing. The exact method is not defined herein.

Reserved, 7 bits: Reserved for future use. The reserved bits must be ignored by the receiver.

The 1588-aware Capability TLV is applicable to both OSPFv2 and OSPFv3.

The 1588-aware Capability TLV MAY be advertised within an area-local or autonomous system (AS) scope Router Information (RI) LSA. But the 1588-aware Capability TLV SHOULD NOT be advertised into an area in more than one RI LSA irrespective of the scope of the LSA.

The flooding scope is controlled by the Opaque LSA type in OSPFv2 and by the S1 and S2 bits in OSPFv3. For area scope, the 1588-aware Capability TLV MUST be carried within an OSPFv2 Type 10 RI LSA or an OSPFv3 RI LSA with the S1 bit set and S2 bit clear. If the flooding scope is the entire routing domain (AS scope), the 1588-aware Capability TLV MUST be carried within an OSPFv2 Type 11 RI LSA or OSPFv3 RI LSA with the S1 bit clear and the S2 bit set.

### 13.2. 1588aware Node Capability for IS-IS

Generic capability advertisement mechanisms for IS-IS are defined in [RFC4971]. These allow a router to advertise its capabilities within an IS-IS area or an entire IS-IS routing domain. This document defines a new sub-TLV (named the 1588-aware Capability) to be carried

within the IS-IS Router Capability TLV.

The IS-IS extensions defined in this document allow for discovering 1588-aware nodes within an IS-IS routing domain. Solutions for 1588-aware nodes discovery across AS boundaries are beyond the scope of this document, and are for further study.

The format of the IS-IS 1588-aware sub-TLV is identical to the TLV format used by the Traffic Engineering Extensions to IS-IS [RFC3784]. That is, the TLV is comprised of 1 octet for the type, 1 octet specifying the TLV length, and a value field. The Length field defines the length of the value portion in octets.

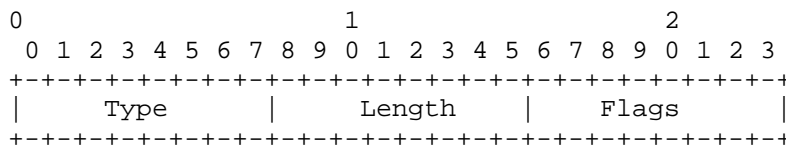


Figure 5: 1588-aware Capability sub-TLV

Where:

Type, 8 bits: 1588-aware Capability sub-TLV where the value is TBD

Length, 8 bits: Gives the length of the flags field in octets, and is currently set to 1

Flags, 8 bits: The bits are defined least-significant-bit (LSB) first, so bit 7 is the least significant bit of the flags octet.

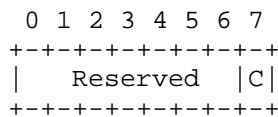


Figure 6: Flags Format

Correction (C) field Update field, 1 bit: Setting the C bit to 1 indicates that the node is capable of recognizing the PTP event packets and can compensate for residence time by updating the PTP packet Correction Field. When this is set to 0, it means that this node cannot perform the residence time correction but is capable of performing MPLS frame forwarding of the frames with PTP labels using a method that support the end to end delivery of accurate timing. The exact method is not defined herein.



Reserved, 7 bits: Reserved for future use. The reserved bits must be ignored by the receiver.

The 1588-aware sub-TLV is optional and is carried within an IS-IS Capability TLV [RFC4971] to facilitate selection of 1588-aware nodes.

The flooding scope for 1588-aware node information advertised through IS-IS can be a single L1 area, an L1 area and the L2 sub-domain, or the entire IS-IS routing domain.

## 14. RSVP-TE Extensions for support of 1588

RSVP-TE signaling MAY be used to setup the PTP LSPs. A new RSVP object is defined to signal that this is a PTP LSP. The OFFSET to the start of the PTP message header MAY also be signaled. Implementations can trivially locate the correctionField (CF) location given this information. The OFFSET points to the start of the PTP header as a node may want to check the PTP messageType before it touches the correctionField (CF).

The LSRs that receive and process the RSVP-TE/GMPLS messages MAY use the OFFSET to locate the start of the PTP message header.

Note that the new object/TLV Must be ignored by LSRs that are not compliant to this specification.

The new RSVP 1588\_PTP\_LSP object should be included in signaling PTP LSPs and is defined as follows:

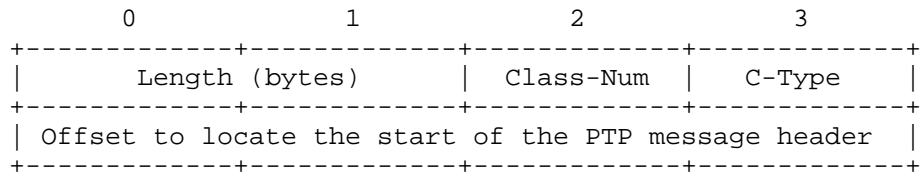


Figure 7: RSVP 1588\_PTP\_LSP object

The ingress LSR MUST include this object in the RSVP PATH Message. It is just a normal RSVP path that is exclusively set up for PTP messages

15. Distributing PW labels

15.1. LDP extensions for distributing PW labels

TBD

15.2. BGP extensions for distributing PW labels

TBD

## 16. Behavior of LER/LSR

### 16.1. Behavior of 1588-aware LER

A 1588-aware LER advertises its 1588-awareness via the OSPF procedure explained in earlier section of this specification. The 1588-aware LER then signals PTP LSPs by including the 1588\_PTP\_LSP object in the RSVP-TE signaling.

When a 1588 message is received from a non-MPLS interface, the LER MUST redirect them to a previously established PTP LSP. When a 1588 over MPLS message is received from an MPLS interface, the processing is similar to 1588-aware LSR processing.

### 16.2. Behavior of 1588-aware LSR

1588-aware LSRs are LSRs that understand the 1588\_PTP\_LSP RSVP object and can perform 1588 processing (e.g. TC processing).

A 1588-aware LSR advertises its 1588-awareness via the OSPF procedure explained in earlier section of this specification.

When a 1588-aware LSR distributes a label for PTP LSP, it maintains this information. When the 1588-aware LSR receives an MPLS packet, it performs a label lookup and if the label lookup indicates it is a PTP label then further parsing must be done to positively identify that the payload is 1588 and not OAM, BFD or control and management. Ruling out non-1588 messages can easily be done when parsing indicates the presence of GAL, ACH or VCCV (Type 1, 2, 3) or when the UDP port number does not match one of the 1588 UDP port numbers.

After a 1588 message is positively identified in a PTP LSP, the PTP message type indicates what type of processing (TC) if any is required. After 1588 processing the packet is forwarded as a normal MPLS packet to downstream node.

### 16.3. Behavior of non-1588-aware LSR

It is most beneficial that all LSRs in the path of a PTP LSP be 1588-aware LSRs. This would ensure the highest quality time and clock synchronization by 1588 Slaves. However, this specification does not mandate that all LSRs in path of a PTP LSP be 1588-aware.

Non-1588-aware LSRs are LSRs that either don't have the capability to process 1588 packets (e.g. TC processing) or don't understand the 1588\_PTP\_LSP RSVP object.

Non-1588-aware LSRs ignore the RSVP 1588\_PTP\_LSP object and just

switch the MPLS packets carrying 1588 messages as data packets and don't perform any TC processing. However as explained in QoS section the 1588 over MPLS packets MUST be still be treated with the highest priority.

## 17. Other considerations

The use of Explicit Null (Label= 0 or 2) is acceptable as long as either the Explicit Null label is the bottom of stack label (applicable only to UDP/IP encapsulation) or the label below the Explicit Null label is a PTP label.

The use of Penultimate Hop Pop (PHP) is acceptable as long as either the PHP label is the bottom of stack label (applicable only to UDP/IP encapsulation) or the label below the PHP label is a PTP label.

## 18. Security Considerations

MPLS PW security considerations in general are discussed in [RFC3985] and [RFC4447], and those considerations also apply to this document.

An experimental security protocol is defined in [IEEE]. The PTP security extension and protocol provides group source authentication, message integrity, and replay attack protection for PTP messages.

19. IANA Considerations

19.1. IANA Considerations for OSPF

IANA has defined a registry for TLVs carried in the Router Information LSA defined in [RFC4970]. IANA is requested to assign a new TLV code-point for the PCED TLV carried within the Router Information LSA.

Value	Sub-TLV	References
TBD	1588aware node sub-TLV	(this document)

19.2. IANA Considerations for IS-IS

IANA has defined a registry for the sub-TLVs carried in the IS-IS Router Capability sub-TLVs defined in [RFC4971]. IANA is requested to assign a new sub-TLV code-point for the 1588aware node sub-TLV carried within the Router Capability sub-TLV.

Value	Sub-TLV	References
TBD	1588aware node sub-TLV	(this document)

19.3. IANA Considerations for RSVP

IANA is requested to assign a new Class Number for 1588 PTP LSP object that is used to signal PTP LSPs.

1588 PTP LSP Object

Class-Num of type 11bbbbbb

Suggested value TBD

Defined CType: 1 (1588 PTP LSP)



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## Abstract

This draft describes network time synchronization mechanisms which may enable increased accuracy, beyond that possible with the current Network Time Protocol version 4 (NTPv4) standard, to time and frequency of computer clocks. The mechanisms considered are those that will provide improved estimates as to when a packet is put on the network, transferred across a network, or taken from the network. Potential standardization actions will be considered for the mechanisms considered, though no such actions are recommended at this time.

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## 1. Introduction

The IETF Timing over IP Connection and Transfer of Clock (TICTOC) Working Group was formed to investigate emerging needs to distribute highly accurate time and frequency information over Internet Protocol (IP) and Multiprotocol Label Switching (MPLS) Packet Switched Networks (PSNs). In this draft, new mechanisms beyond those identified in the NTPv4 standard (i.e., Request For Comment (RFC)

5905) are considered to provide increased time synchronization accuracy for computer (i.e., operating system) clocks' time and frequency. The mechanisms considered are those that will provide improved estimates as to when a packet is put on the network, transferred across a network, or taken from the network. This draft identifies a set of mechanisms that are candidates for experimentation. Standardization considerations will be described for the mechanisms identified.

The purpose of this draft is to examine methods for improving NTPv4 time synchronization performance. The authors are requesting comments and contributions on the mechanisms described and on additional mechanisms that should be considered. It is hoped that discussions within the IETF TICTOC Working Group will motivate experimentation that will lead to standardization actions to enable better accuracy to those utilizing a future Network Time Protocol (NTP) specification.

### 1.1. Motivation for Increased Performance

There are two reasons to improve upon the time synchronization performance that is currently available from NTP. Not only is the increased performance needed for existing product designs that would make use of the added performance if it were available, but it is expected that new uses will be identified that were not even possible until performance is improved. This is similar to how network speeds are increased every several years, and the uses for the increased network bandwidth soon follow.

The current methods for achieving an increase in time synchronization performance involve use of a technology separate from the existing computer network (e.g., Inter-Range Instrumentation Group [IRIG] technology) or use of a technology like Precision Time Protocol (PTP), which is defined in the Institute of Electrical and Electronics Engineers (IEEE) 1588 standard. With PTP, the computer applications must interface with the installed PTP hardware in order to read time from its oscillator. There is a lot of resiliency built into NTP, which does not exist in the PTP hardware oscillator. It is unknown what happens to the time provided by the PTP hardware when a network switch in the network path to the time source is temporarily unavailable (i.e., the network switch gets rebooted). It would be beneficial to have the resiliency of the NTP algorithms be paired with the highly accurate PTP hardware-based time distribution.

## 1.2. NTP/PTP Commonality and Differences

NTP and PTP are both packet-based protocols for exchanging time with a time server over a computer network. Both protocols are used to determine the offset between two independent clocks. Both implement a hierarchical tree structure for obtaining time from a master time source through intermediary time sources to time clients. Both assume network paths are symmetric, and both have their own methods for addressing network delays that are not symmetric. NTP uses its algorithms to determine which of several consecutive time measurements are most accurate and uses that measurement. PTP makes use of hardware means of measuring delays as packets traverse intermediate network devices and corrects its received time information based upon those measured delays.

Because PTP has the ability to measure actual packet delays and to correct for them, PTP can provide the most accurate measurement of clock offset between two clocks. PTP does not define the method for synchronizing that clock once the highly accurate time measurements have been obtained. PTP is normally used to synchronize a hardware clock located on an interface card and does not synchronize the operating system clock. NTP, on the other hand, possesses the ability to synchronize the operating system clock based upon received clock offset measurements. The NTP algorithms have a lot of resiliency so that operating system clocks stay stable despite the conditions on the network.

## 2. Use Case Targeted

The use case considered by this internet draft is a dense concentration of computing elements connected by a network. A satellite-based time source (e.g., Global Positioning System [GPS]) is used for synchronizing primary time servers. Secondary time servers and leaf computing elements are synchronized to the primary time servers via the network. In this use case, there are approximately 150 or so total computers where there are 3 to 4 levels of time servers. These time servers may have to communicate to each other through layer 2 and layer 3 network switches (could be 10 to 20 different layer 2 subnetworks). All of the computers are connected together through gigabit or faster network connections. In this environment, there will be some groups of computers that will need to synchronize to each other to within a microsecond, while other groups of computers only have to be synchronized to each other to within a millisecond. In this use case, there is one interconnected time synchronization scheme where NTP, PTP, or a combination of both is used to meet all time synchronization needs.



### 2.1. Emerging Need for NTP and PTP Commonality

Due to its accuracy capabilities, PTP is beginning to replace NTP as the base protocol for time clients within dense computing sites. This results in an implementation in which some hosts use PTP while others use NTP within the same building and sometimes within the same room. Over time, hosts are being changed from NTP to PTP. This leads to an emerging need to provide similar approaches for basic time service functions for operational ease of managing the time distribution assets. Examples of functions where commonality is considered to be an emerging need include providing synchronization of the computer clock, providing management to time clients, and configuring timeservers. A standard means of synchronizing the computer clock for both protocols is in particular of interest because there appears to be no value in using different methods when the hosts that both protocols are supporting are often working within the same system.

In addition, there are highly accurate PTP time clients that could serve as highly accurate secondary timeservers for the NTP time clients if this capability were supported in vendor products.

## 3. Approach

The approach taken by the authors of this draft includes determining what the current accuracy capabilities are with NTPv4 and investigating additional mechanisms that may provide improvements in accuracy. Through experiments of those additional mechanisms, estimations of improvements can be calculated. Depending on the standardization difficulty and potential benefits offered, more than one standardization action may be recommended in the future.

## 4. Mechanisms Considered

### 4.1. NTP Interleave

NTP Interleave is an extension of the NTPv4 protocol, which is included in the current NTP distribution [1]. It is designed to be backward compatible (i.e., not affecting NTP implementations that do not use the interleave extensions). It also utilizes the same NTP packet format as the current standard NTPv4. NTP Interleave uses an IEEE 1588 PTP-like feature that provides a follow-up packet with a better estimate of when a previous NTP packet was sent on the network and a message exchange sequence to determine network mean path delay.

This mechanism could be used by some of the primary time servers for synchronizing secondary (i.e., lower stratum) time servers and leaf computing elements, which have very accurate time synchronization requirements.

Future experimentation may identify what gains are possible with this mechanism.

#### 4.1.1. Standardization Considerations

If the NTP Interleave investigation generates promising results, this may initiate a standardization proposal to add this capability to the NTPv4 standard.

#### 4.2. Use of IEEE 1588 PTP and 802.1AS Mechanisms in the Underlying Network Service (e.g., Network Interface Controller [NIC])

The purpose of investigating this mechanism is to determine if using special capabilities in the underlying network service can improve the timestamp estimates when NTP packets are put on the network, transferred across networks, or taken from the network.

##### 4.2.1. Standardization Considerations

If the investigation of these mechanisms generates promising results, this may initiate a standardization proposal for additions to the NTPv4 standard to make use of these capabilities. Alternatively, a modification to the NTPv4 standard may be proposed, which would enable hardware assists to be incorporated into future NICs.

#### 4.3. Use of IEEE 1588 PTP to Synchronize Computer Clocks

The purpose of investigating this mechanism is to determine the viability of using the IEEE 1588 to synchronize computing elements, which require very accurate synchronization. This mechanism considers bringing the 1588 synchronization all the way to the computer clock through a standardized clock discipline algorithm. Computing elements synchronized by 1588 are candidates to be time servers (by the use of NTP) for computing elements not synchronized by 1588. Based on their respective strengths, the natural way to merge NTP and PTP would be to use PTP as the means of obtaining extremely accurate time information from across the network and to let the NTP algorithms use that time to keep local clocks synchronized.

#### 4.3.1. Standardization Considerations

If the investigation of this mechanism generates promising results, this may initiate a standardization proposal to specify a 1588 profile for use by NTP. It may be possible to replace the current NTP clock coordination services without affecting the NTP time management services or the clock access mechanisms used by each operating system. A variety of studies will be needed if this approach is pursued, including a study to determine if there are any issues for secondary time servers to run both NTP and PTP. If such issues are identified, standards activities may be needed in the IETF or in IEEE 1588.

#### 5. Early Experimentation

Some preliminary experiments were performed that tested the new Interleave mode available in NTP version 4.2.6. This mode mimics the operation of PTP defined by IEEE 1588 where an additional follow-up message is sent so that a more accurate transmission time can be used. In this experiment, seven workstations running Red Hat Enterprise Linux (RHEL) 5 were used. These workstations are 2-years old and make use of two dual-core processors. Since the GPS-based (i.e., stratum 1) time server does not currently have version 4.2.6 of NTP available, which supports the interleave mode, one of the seven workstations was synchronized to the stratum 1 time server; then, served as the stratum 2 time server to the other six stratum 3 workstations. The stratum 2 server and four of the six other workstations were upgraded to use NTP version 4.2.6, while the remaining two workstations were left running version 4.2.2 of NTP that was included with RHEL 5.

Interleave mode was achieved by using the "xleave" option when either the broadcast mode or peer mode was used under NTP. The stratum 2 server was configured as a broadcast server, making use of the standard multicast address and using the "xleave" option. Two of the workstations running NTP 4.2.6 were configured as multicast clients so that the Interleave mode was utilized. The other two NTP 4.2.6 workstations were configured to synchronize to the stratum 2 server using standard client/server (unicast) mode. The two workstations running 4.2.2 of NTP were configured to be broadcast clients; however, they did not use Interleave mode since NTP 4.2.2 does not include Interleave support. All NTP polling intervals were configured to 16 seconds.

Offset measurements were obtained between the six clients and the stratum 2 server using the ntpdate command with the "-q" option. Measurements were taken every minute over a period of approximately

four days. These workstations were not running any other major tasks, and NTP ran over a network with no discernable network load. All of the workstations were connected through the same Virtual Local Area Network on the same network switch (i.e., no routers involved). All of the network connections were 100 Mbit/sec Ethernet.

Some results were obtained from the experiments where the average and the standard deviation of the absolute value of clock offset were measured. The worst behaved NTP Interleave client was able to stay synchronized with an average clock offset of 9 microseconds with a standard deviation of 8 microseconds. The worst behaved computer that synchronized using client/server mode was able to maintain an average clock offset of 11 microseconds with a standard deviation of 10 microseconds. The worst behaved broadband (without NTP Interleave) client stayed synchronized with an average clock offset of 49 microseconds and with a standard deviation of 58 microseconds.

These results illustrate that broadcast with NTP Interleave does provide results that are better than having every client poll the server via unicast. However, the result is not significantly better (e.g., not an order of magnitude better). Preliminary experiments with hardware-based PTP have been performed in the past where the average offsets between PTP NICs and the PTP Grandmaster clock are in the hundreds of nanoseconds with standard deviations in the tens of nanoseconds.

#### 5.1. Future Experimentation

The results so far should be treated as preliminary, and further experiments are needed to draw conclusions. One concern is that the clock offsets are in the microsecond range. The use of ntpdate may not be a valid way to accurately measure clock offsets at this level since ntpdate makes measurements across the network and is susceptible to errors caused by variations in network delay. Further work is needed to ensure valid offset measurements. An out-of-band measurement technique, which is not affected by variations in network delay, needs to be investigated for use in future experiments.

Another concern is that no load has been applied either to the processor or to the network carrying the NTP packets. Future experiments need to be performed to measure the resiliency of the NTP Interleave while under network and processor load and then make comparisons to the other time synchronization methods. This experiment was performed in a configuration in which the test

workstations were connected to the same network switch. Future experiments should be performed to determine how performance is affected when a more complex network configuration is used.

Developing a standard to define time synchronization performance metrics would be beneficial by allowing different experimental efforts to be performed in a way that the results are comparable.

## 6. Analysis of Results

In 2008, it was reported [2] that the synchronization of computer clocks using NTPv4 over a Local Area Network (LAN) can approach values on the order of ~145us (mean plus 3 sigma). This has been demonstrated in the laboratory and was achieved across a single stratum (e.g., stratum 1 to stratum 2). Recently, using newer hardware and a newer version of NTP, time synchronization values were measured on the order of ~40us (mean plus 3 sigma). The variation in the time synchronization comprises the majority of the 40us value. The results need to be analyzed in detail, but in a little over two years, the time synchronization of computer clocks over a LAN has improved; most likely due to hardware improvements and minor software enhancements.

If this 40us synchronization can be maintained from stratum to stratum for all subsequent tiers in a well-engineered network with 3 to 5 stratum, a maximum theoretical time synchronization offset on the order of 120 to 200us could be achieved. A 50 percent improvement in the variability of clock synchronization from stratum to stratum would reduce this number to less than 125us. This does not imply that stratum-to-stratum accuracy should not be improved. On the contrary, by working on the accuracy and variability together all users of time synchronization will benefit. This needs to be accomplished without overloading the computer or the network.

A 50 percent improvement appears to be a reasonable goal; however, if the stratum-to-stratum synchronization variability could be improved by an order of magnitude, it is reasonable to anticipate maximum theoretical time synchronization offsets of 50us or less in a well-behaved LAN and potentially in the hundreds of microseconds for a Wide Area Network. Follow-on discussions within the IETF TICTOC Working Group are needed to define what level of improvement in accuracy would warrant a standardization action to add an Interleave option to the NTP standard.

## 7. Security Considerations

Security aspects of the aforementioned options will need to be considered in more detail.

## 8. Internet Assigned Numbers Authority (IANA) Considerations

No IANA actions are required as a result of the publication of this document.

## 9. Conclusions

Preliminary results tentatively indicate that an Interleave capability can provide a modest improvement to the accuracy achieved with NTPv4 and, as a result of these preliminary findings, the authors recommend conducting further testing to determine whether this should be added as an option to the NTPv4 standard. Preliminary results also tentatively indicate that Interleave will not provide accuracies in the range that PTP with hardware assists can; thus, the other two mechanisms described in this paper should be pursued in parallel.

Assuming any of the proposed mechanisms in this draft or other mechanisms proposed through IETF TICTOC discussions provide a significant improvement in computer clock time synchronization, it is recommended that the IETF TICTOC Working Group initiate work on a standards-track targeted working group Internet-Draft.

It would be of benefit to the IETF TICTOC Working Group to standardize the performance metrics used in describing the behavior of clock synchronization. Such a standard would enable better comparisons between the mechanisms considered. In addition to the same definitions of the metrics used, better agreement should be obtained in experiments being performed by different organizations.

The authors are interested in contributions on the mechanisms described in this draft as well as on additional mechanisms that may improve the accuracy of computer clocks synchronized via a network. The authors request that other experimental results on mechanisms that can improve NTPv4 accuracy be shared. It is hoped that discussions on this topic within the IETF TICTOC Working Group will lead to standardization NTP actions to enable better accuracy to those utilizing a future NTP specification.

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Abstract

The TICTOC charter specifies that the working group is concerned with highly accurate time and frequency distribution over native IP and MPLS-enabled IP Packet Switched Networks (PSNs). We discuss here issues specific to MPLS PSNs.

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## 1. Motivation

The TICTOC charter specifies that the working group is concerned with highly accurate time and frequency distribution over native IP and MPLS-enabled IP Packet Switched Networks (PSNs). To date, discussions have focused on NTPv4 and 1588-2008 timing distribution using UDP/IP encapsulations. The present document discusses transport of timing packets over MPLS PSNs, and is based on material presented and discussed in previous TICTOC meetings.

We first must address the question as to why special treatment is needed at all for transport of timing packets over MPLS. Timing packets in IP format can certainly be transported over MPLS without the LSRs along the path being aware of them. However, there are advantages to being able to recognize, and potentially manipulate, timing packets.

For highly accurate timing distribution, timing packets are required to travel through the PSN with minimal, symmetric, and quasistationary delay, as well as minimal and uncorrelated packet delay variation. Thus, prioritization and symmetric routing of timing packets are minimal requirements. For the most demanding applications, timing distribution mechanisms avail themselves of "on-path support", such as Transparent Clock (TC) overwriting of header fields. None of these can be selectively applied to timing packets unless they can be recognized by the LSR.

## 2. Encapsulation Options

MPLS as a server layer presently permits three types of clients:

1. MPLS (via label stacking)
2. IP (either IPv4 or IPv6)
3. pseudowires (PWs)

although we can not rule out defining a new client for timing distribution. Taking into account the two defined associated channels that carry non-user traffic, namely the VCCV associated channel for PWs [RFC5085] and the GAL G-ACh defined for MPLS-TP [RFC5586] any proposal for carrying timing packets over MPLS will need to put the timing information in one of the following six formats:

1. a new MPLS client type
2. IP
3. an Ethernet PW

4. a new "timing" pseudowire type
  5. a new VCCV channel type
  6. a new G-ACh channel type.
- We will discuss each of these options in turn.

### 2.1. Using IP

Since the two main timing distribution protocols have UDP/IP encapsulations, arguably the simplest method of transporting timing packets is in this format. There are several methods for intermediate LSRs to recognize the timing packets :

- o Deep Packet Inspection (i.e., peeking under the MPLS stack and identifying well-known port numbers)
- o using an arbitrary configured or signaled MPLS label
- o using a new reserved MPLS label
- o using a specific MPLS Traffic Class (ex-EXP).

It should be noted that there are very few available reserved labels, and that traffic class usage is not standardized.

If an arbitrary label is required to be signaled, then an extension to the signaling protocol will be required. Such an extension will identify the FEC as belonging to a timing flow.

### 2.2. Using an Ethernet PW

The UDP/IP packets of the timing distribution protocols of the previous subsection may be contained in Ethernet frames, that can be transported over an MPLS network in an Ethernet PW. In addition, IEEE 1588-2008 has a native Ethernet (non-IP) encapsulation.

Methods for identifying timing packets inside an Ethernet PW are similar to those of the previous subsection.

### 2.3. Defining a New MPLS Client or a new PW Type

IEEE 1588-2008 allows for different transport layers to be defined, with annexes to the standard defining UDP/IPv4, UDP/IPv6, Ethernet, and several other transport networks. It is possible to define a new transport mechanism for MPLS, which entails specifying how to encapsulate a PTP PDU in an MPLS packet. This can be done in the context of the PWE3 architecture (this was proposed in draft-ronc-ntp-mpls-00.txt, now expired) or without reference to that architecture. If the PWE3 architecture is used, then PWE3 features, such as the control word and the control protocol, may be used.

Whether the MPLS encapsulation is a PW or not, the timing packet will be recognized by virtue of its bottom-of-stack label. When additional labels are present, the LSR will need to search for the

label with S=1. This technique is technically a violation of the MPLS architecture, but is presently performed for other purposes, e.g., ECMP avoidance [RFC4928].

The particular label(s) used to signify timing packets may be distributed by manual configuration or a signaling protocol. If the latter is employed, a new FEC type will need to be defined. If a PW is used, a new MPLS Pseudowire Type [RFC4446] will be needed for use in the PWE3 control protocol [RFC4447].

#### 2.4. Using VCCV or the G-ACh

Timing is often distributed in order to enable proper functioning of applications that already define PWs between the end points of the timing flow. In this case, the timing information may be placed in the VCCV associated channel of the application's PW. The associated channel is typically identified by having the same PW (bottom-of-stack) label as the application, but a control word with 0001 in its initial nibble. The timing packets may then be identified by a new channel type, or may use the IP channel type and then would be identified by the well-known port numbers. It is recognized that this requires the LSR to perform relatively deep packet inspection.

If there is no application PW between the timing end points, then we may still use the Generic Associated CHannel (G-ACh) defined in [RFC5586] in a similar manner. The G-ACh is identified by the G-ACh Label (GAL), which is a reserved MPLS label of value 13, at its bottom-of-stack. The format after the GAL is the same as that of VCCV, and thus the considerations of the previous paragraph apply.

#### 3. IANA Considerations

This document requires no IANA actions.

#### 4. References

- [RFC4446] Martini, L., "IANA Allocations for Pseudowire Edge to Edge Emulation (PWE3)", BCP 116, RFC 4446, April 2006.
- [RFC4447] Martini, L., Rosen, E., El-Aawar, N., Smith, T., and G. Heron, "Pseudowire Setup and Maintenance Using the Label Distribution Protocol (LDP)", RFC 4447, April 2006.
- [RFC4928] Swallow, G., Bryant, S., and L. Andersson, "Avoiding Equal Cost Multipath Treatment in MPLS Networks", BCP 128, RFC 4928, June 2007.

[RFC5085] Nadeau, T. and C. Pignataro, "Pseudowire Virtual Circuit Connectivity Verification (VCCV): A Control Channel for Pseudowires", RFC 5085, December 2007.

[RFC5586] Bocci, M., Vigoureux, M., and S. Bryant, "MPLS Generic Associated Channel", RFC 5586, June 2009.

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Precision Time Protocol Version 2 (PTPv2)  
Management Information Base  
draft-vinay-tictoc-ntp-mib-00.txt

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## Abstract

This memo defines a portion of the Management Information Base (MIB) for use with network management protocols in TCP/IP-based internets. In particular, it defines objects for managing networks using Precision Time Protocol.

This memo specifies a MIB module in a manner that is both compliant to the SNMPv2 SMI, and semantically identical to the peer SNMPv1 definitions.

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## 1. Introduction

This memo defines a portion of the Management Information Base (MIB) for use with network management protocols in the Internet Community. In particular, it describes managed objects used for managing PTP devices including the ordinary clock, transparent clock, boundary clocks. It is envisioned this MIB will complement other managed objects defined to monitor, measure the performance of the PTP devices and telecom clocks. Those objects are considered out of scope for the current draft.

## 1.1. Change Log

This section tracks changes made to the revisions of the Internet Drafts of this document. It will be *\*deleted\** when the document is published as an RFC. This section tracks changes made to the revisions of the Internet Drafts of this document. It will be *\*deleted\** when the document is published as an RFC.

## 2. The SNMP Management Framework

The SNMP Management Framework presently consists of five major components:

- o An overall architecture, described in [RFC 2271].
- o Mechanisms for describing and naming objects and events for the purpose of management. The first version of this Structure of Management Information (SMI) is called SMIV1 and described in [RFC 1155], [RFC 1212] and [RFC 1215]. The second version, called SMIV2, is described in [RFC 2578], [RFC 2579] and [RFC 2580].
- o Message protocols for transferring management information. The first version of the SNMP message protocol is called SNMPv1 and described in [RFC 1157]. A second version of the SNMP message protocol, which is not an Internet standards track protocol, is called SNMPv2c and described in [RFC 1901] and [RFC 1905]. The third version of the message protocol is called SNMPv3 and described in [RFC 1906], [RFC 2572] and [RFC 2574].
- o Protocol operations for accessing management information. The first set of protocol operations and associated PDU formats is described in [RFC 1157]. A second set of protocol operations and associated PDU formats is described in [RFC 1905].
- o A set of fundamental applications described in [RFC 2573] and the view-based access control mechanism described in [RFC 2575].

Managed objects are accessed via a virtual information store, termed the Management Information Base or MIB. Objects in the MIB are defined using the mechanisms defined in the SMI.

This memo specifies a MIB module that is compliant to the SMIV2. A MIB conforming to the SMIV1 can be produced through the appropriate translations. The resulting translated MIB must be semantically equivalent, except where objects or events are omitted because no translation is possible (e.g., use of Counter64). Some machine readable information in SMIV2 will be converted into textual descriptions in SMIV1 during the translation process. However, this

loss of machine readable information is not considered to change the semantics of the MIB.

### 3. Overview

The objects defined in this MIB are to be used when describing Precision Time Protocol (PTPv2).

### 4. Structure of MIB Module

```
IETF-PTP-MIB DEFINITIONS ::= BEGIN
```

```
ietfPtpMIB MODULE-IDENTITY
```

```
  LAST-UPDATED      "201101280000Z"  
  ORGANIZATION      "TICTOC Working Group"  
  CONTACT-INFO
```

```
    "WG Email: tictoc@ietf.org
```

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```

```
    Greg Dowd,  
    Symmetricom Inc.,  
    Email: gdowd@symmetricom.com"
```

```
DESCRIPTION
```

```
"The MIB module for PTPv2 (IEEE (TM) 1588 - 2008)
```

```
Overview of PTPv2 (IEEE (TM) 1588-2008)
```

```
[IEEE1588-2008] defines a protocol enabling precise  
synchronization of clocks in measurement and control systems  
implemented with packet-based networks, the Precision Time  
protocol version 2 (PTPv2). This MIB does not address the  
standard IEEE(TM) 1588 (2002) and PTPv1. The protocol is  
applicable to network element communicating using IP. The  
protocol enables heterogeneous systems that include clocks  
of various inherent precision, resolution, and stability to
```

synchronize to a grandmaster clock.

The protocol supports system-wide synchronization accuracy in the sub-microsecond range with minimal network and local clock computing resources. [IEEE1588-2008] uses UDP/IP or Ethernet and can be adapted to other mapping. It includes formal mechanisms for message extensions, higher sampling rates, correction for asymmetry, a clock type to reduce error accumulation in large topologies, and specifications on how to incorporate the resulting additional data into the synchronization protocol. The [IEEE1588-2008] defines conformance and management capability also.

#### MIB description

This MIB is to support the Precision Time Protocol version 2 (PTPv2, hereafter designated as PTP) features of network Element System devices.

#### Acronyms:

ARB	arbitrary
BMC	best master clock
CAN	Controller Area Network
CP	Communication Profile (according to IEC 61784-1:200710)
CPF	Communication Profile Family (according to IEC 61784-1:2007)
DS	Differentiated Service
E2E	End-to-End
E2ETC	End-to-End Transparent Clock
EUI	Extended Unique Identifier.
FFO	Fractional Frequency Offset
GPS	Global Positioning System
IANA	Internet Assigned Numbers Authority
ICV	Integrity Check Value
ID	Identification
IPv4	Internet Protocol version 4
IPv6	Internet Protocol version 6
JD	Julian Date
JDN	Julian Day Number
MAC	Media Access Control (according to IEEE Std 802.3-2005)
MJD	Modified Julian Day
NIST	National Institute of Standards and Technology (see <a href="http://www.nist.gov">http://www.nist.gov</a> )
NTP	Network Time Protocol (see [RFC 1305])
OUI	Organizational Unique Identifier (allocated by the IEEE)
P2P	Peer-to-Peer

P2PTC	Peer-To-Peer Transparent Clock
PHY	physical layer (according to IEEE Std 802.3 2005)
POSIX	Portable Operating System Interface (see ISO/IEC 9945:2003)
PPS	Pulse per Second
PTP	Precision Time Protocol
SA	Security Associations
SNTP	Simple Network Time Protocol
SOF	Start of Frame
TAI	International Atomic Time
TC	Traffic Class
TC	Transparent Clock
TLV	Type, Length, Value (according to IEEE Std 02.1AB)
ToD	Time of Day Synchronization
ToS	Type of Service
UCMM	UnConnect Message Manager
UDP/IP	User Datagram Protocol
UTC	Coordinated Universal Time

#### References:

[1] Precision clock synchronization protocol for networked measurement and control systems - IEC 61588 IEEE 1588(tm) Edition 2.0 2009-02

As defined in [IEEE1588-2008]:

#### Accuracy:

The mean of the time or frequency error between the clock under test and a perfect reference clock, over an ensemble of measurements. Stability is a measure of how the mean varies with respect to variables such as time, temperature, and so on.

The precision is a measure of the deviation of the error from the mean.

#### Atomic process:

A process is atomic if the values of all inputs to the process are not permitted to change until all of the results of the process are instantiated, and the outputs of the process are not visible to other processes until the processing of each output is complete.

#### Boundary clock:

A clock that has multiple Precision Time Protocol (PTP) ports in a domain and maintains the timescale used in the domain. It may serve as the source of time, i.e., be a master clock,

and may synchronize to another clock, i.e., be a slave clock.

**Clock:**

A node participating in the Precision Time Protocol (PTP) that is capable of providing a measurement of the passage of time since a defined epoch.

**Domain:**

A logical grouping of clocks that synchronize to each other using the protocol, but that are not necessarily synchronized to clocks in another domain.

**End-to-end transparent clock:**

A transparent clock that supports the use of the end-to-end delay measurement mechanism between slave clocks and the master clock. Each node must measure the residence time of PTP event messages and accumulate it in Correction Field.

**Epoch:**

The origin of a timescale.

**Event:**

An abstraction of the mechanism by which signals or conditions are generated and represented.

**Foreign master:**

An ordinary or boundary clock sending Announce messages to another clock that is not the current master recognized by the other clock.

**Grandmaster clock:**

Within a domain, a clock that is the ultimate source of time for clock synchronization using the protocol.

**Holdover:**

A clock previously synchronized/syntonized to another clock (normally a primary reference or a master clock) but now free-running based on its own internal oscillator, whose frequency is being adjusted using data acquired while it had been synchronized/syntonized to the other clock. It is said to be in holdover or in the holdover mode, as long as it is within its accuracy requirements.

**Link:**

A network segment between two Precision Time Protocol ports supporting the peer delay mechanism of this standard. The peer delay mechanism is designed to measure the propagation time over such a link.

**Management node:**

A device that configures and monitors clocks.

**Master clock:**

In the context of a single Precision Time Protocol communication path, a clock that is the source of time to which all other clocks on that path synchronize.

**Message timestamp point:**

A point within a Precision Time Protocol event message serving as a reference point in the message. A timestamp is defined by the instant a message timestamp point passes the reference plane of a clock.

**Multicast communication:**

A communication model in which each Precision Time Protocol message sent from any PTP port is capable of being received and processed by all PTP ports on the same PTP communication path.

**Node:**

A device that can issue or receive Precision Time Protocol communications on a network.

**One-step clock:**

A clock that provides time information using a single event message.

**On-pass support:**

Indicates that each node in the synchronization chain from master to slave can support IEEE-1588.

**Ordinary clock:**

A clock that has a single Precision Time Protocol port in a domain and maintains the timescale used in the domain. It may serve as a source of time, i.e., be a master clock, or may synchronize to another clock, i.e., be a slave clock.

**Parent clock:**

The master clock to which a clock is synchronized.

**Peer-to-peer transparent clock:**

A transparent clock that, in addition to providing Precision Time Protocol event transit time information, also provides corrections for the propagation delay of the link connected to the port receiving the PTP event message. In the presence of peer-to-peer transparent clocks, delay measurements between slave clocks and the master clock are



performed using the peer-to-peer delay measurement mechanism.

**Phase change rate:**

The observed rate of change in the measured time with respect to the reference time. The phase change rate is equal to the fractional frequency offset between the measured frequency and the reference frequency.

**PortNumber:**

An index identifying a specific Precision Time Protocol port on a PTP node.

**Primary reference:**

A source of time and or frequency that is traceable to international standards.

**Profile:**

The set of allowed Precision Time Protocol features applicable to a device.

**Precision Time Protocol communication:**

Information used in the operation of the protocol, transmitted in a PTP message over a PTP communication path.

**Precision Time Protocol communication path:**

The signaling path portion of a particular network enabling direct communication among ordinary and boundary clocks.

**Precision Time Protocol node:**

PTP ordinary, boundary, or transparent clock or a device that generates or parses PTP messages.

**Precision Time Protocol port:**

A logical access point of a clock for PTP communications to the communications network.

**Recognized standard time source:**

A recognized standard time source is a source external to Precision Time Protocol that provides time and/or frequency as appropriate that is traceable to the international standards laboratories maintaining clocks that form the basis for the International Atomic Time and Universal Coordinated Time timescales. Examples of these are Global Positioning System, NTP, and National Institute of Standards and Technology (NIST) timeservers.

**Requestor:**

The port implementing the peer-to-peer delay mechanism that

initiates the mechanism by sending a Pdelay\_Req message.

**Responder:**

The port responding to the receipt of a Pdelay\_Req message as part of the operation of the peer-to-peer delay mechanism.

**Synchronized clocks:**

Two clocks are synchronized to a specified uncertainty if they have the same epoch and their measurements of the time of a single event at an arbitrary time differ by no more than that uncertainty.

**Syntonized clocks:**

Two clocks are syntonized if the duration of the second is the same on both, which means the time as measured by each advances at the same rate. They may or may not share the same epoch.

**Timeout:**

A mechanism for terminating requested activity that, at least from the requester's perspective, does not complete within the specified time.

**Timescale:**

A linear measure of time from an epoch.

**Traceability:**

A property of the result of a measurement or the value of a standard whereby it can be related to stated references, usually national or international standards, through an unbroken chain of comparisons all having stated uncertainties.

**Translation device:**

A boundary clock or, in some cases, a transparent clock that translates the protocol messages between regions implementing different transport and messaging protocols, between different versions of IEEE Std 1588-2008/IEC 61588:2009, or different Precision Time Protocol profiles.

**Transparent clock:**

A device that measures the time taken for a Precision Time Protocol event message to transit the device and provides this information to clocks receiving this PTP event message.

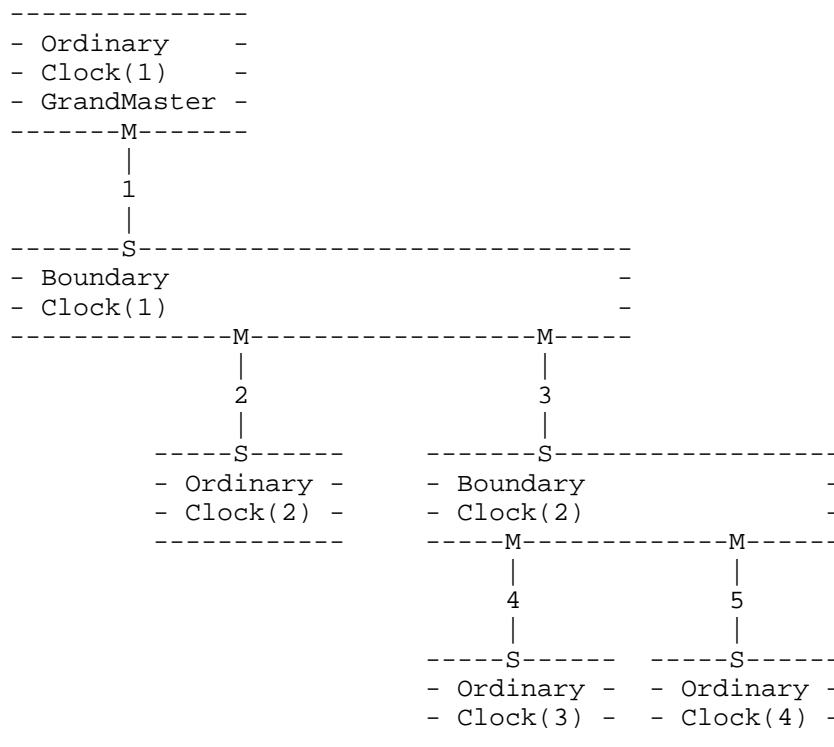
**Two-step clock:**

A clock that provides time information using the combination of an event message and a subsequent general message.

The below table specifies the object formats of the various textual conventions used.

Data type mapping	Textual Convention	SYNTAX
5.3.2 TimeInterval	ClockTimeInterval	OCTET STRING (SIZE(1..255))
5.3.3 Timestamp	ClockTimestamp	OCTET STRING(SIZE(6))
5.3.4 ClockIdentity	ClockIdentity	OCTET STRING (SIZE(1..255))
5.3.5 PortIdentity	ClockPortNumber	INTEGER(1..65535)
5.3.7 ClockQuality	ClockQualityClassType	

Simple master-slave hierarchy [1] section 6.6.2.4



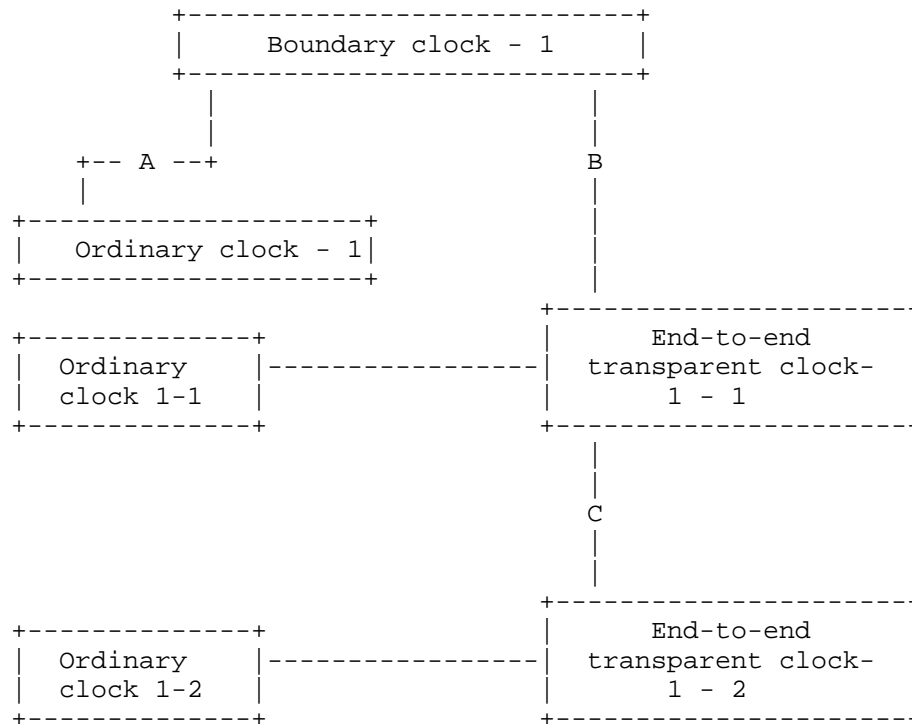
Grandmaster

Boundary Clocks(0-N)    Ordinary Clocks(0-N)  
 Ordinary Clocks(0-N)

Relationship cardinality

PTP system 1 : N PTP Clock  
 PTP Clock 1 : 1 Domain  
 PTP Clock 1 : N PTP Ports  
 PTP Port N : N Physical Port (interface in IF-MIB)

Transparent clock diagram from section 6.7.1.3 of [1]



The MIB refers to the sections of [IEEE1588-2008]."

#### 4.1. Textual Conventions

```

ClockDomainType ::= TEXTUAL-CONVENTION
    DISPLAY-HINT    "d"
    STATUS          current
    DESCRIPTION
        "The Domain is identified by an integer, the domainNumber,
        in the range of 0 to 255. An integer value that is used to
        assign each PTP device to a particular domain. The following
        values define the valid domains. [1] Section 7.1 Domains
        Table 2
  
```

Value definition.

```

-----
0          Default domain
1          Alternate domain 1
2          Alternate domain 2
3          Alternate domain 3
4 - 127   User-defined domains
128 - 255 Reserved"

```

```

REFERENCE      "Section 7.1 Domains and Table 2 of [1]"
SYNTAX         Unsigned32 (0..255)

```

ClockIdentity ::= TEXTUAL-CONVENTION

STATUS current

DESCRIPTION

"The clock Identity is an 8-octet array and will be presented in the form of a character array. The value of the ClockIdentity should be taken from the IEEE EUI-64 individual assigned numbers as indicated in Section 7.5.2.2.2 of [1]. The EUI-64 address is divided into the following fields.

```

OUI byte[0-2]
Extension identifier bytes:[3-7]

```

The clock identifier can be constructed from existing EUI-48 assignments and here is an abbreviated example extracted from section 7.5.2.2.2 [1].

```

Company EUI-48 = 0xACDE4823456716
EUI-64 = ACDE48FFFE23456716

```

It is important to note the IEEE Registration Authority has deprecated the use of MAC-48 in any new design."

```

REFERENCE      "Section 7.5.2.2.1 from [1]"
SYNTAX         OCTET STRING (SIZE (1..255))

```

ClockIntervalBase2 ::= TEXTUAL-CONVENTION

DISPLAY-HINT "d"

STATUS current

DESCRIPTION

"The interval included in message types Announce, Sync, Delay\_Req, and Pdelay\_Req as indicated in section 7.7.2.1 of [1].

The mean time interval between successive messages shall be represented as the logarithm to the base 2 of this time interval measured in seconds on the local clock of the device sending the message. The values of these logarithmic

attributes shall be selected from integers in the range -128 to 127 subject to further limits established in an applicable PTP profile."

REFERENCE

"Section 7.7.2.1 General interval specification of [1]"

SYNTAX Integer32 (-128..127)

ClockMechanismType ::= TEXTUAL-CONVENTION

STATUS current

DESCRIPTION

"The clock type based on whether End to End or peer to peer mechanisms are used. The mechanism used to calculate the Mean Path Delay as indicated in Table 9 of [IEEE1588-2008].

Delay mechanism	Value(hex)	Specification
E2E	01	The port is configured to use the delay request-response mechanism.
P2P	02	The port is configured to use the peer delay mechanism.
DISABLED	FE	The port does not implement the delay mechanism."

REFERENCE "Sections 8.2.5.4.4, 6.6.4 and 7.4.2 of [1]."

SYNTAX INTEGER { e2e(1), p2p(2), disabled(254) }

ClockInstanceType ::= TEXTUAL-CONVENTION

DISPLAY-HINT "d"

STATUS current

DESCRIPTION

"The instance of the Clock of a given clock type in a given domain."

SYNTAX Unsigned32 (0..255)

ClockPortNumber ::= TEXTUAL-CONVENTION

DISPLAY-HINT "d"

STATUS current

DESCRIPTION

"An index identifying a specific Precision Time Protocol (PTP) port on a PTP node."

REFERENCE "Section 7.5.2.3 Port Number and 5.3.5 of [1]"

SYNTAX Unsigned32 (0..65535)

ClockPortState ::= TEXTUAL-CONVENTION

STATUS current

DESCRIPTION

"This is the value of the current state of the protocol engine associated with this port.

Port state	Value	Description
initializing	1	In this state a port initializes its data sets, hardware, and communication facilities.
faulty	2	The fault state of the protocol.
disabled	3	The port shall not place any messages on its communication path.
listening	4	The port is waiting for the announceReceiptTimeout to expire or to receive an Announce message from a master.
preMaster	5	The port shall behave in all respects as though it were in the MASTER state except that it shall not place any messages on its communication path except for Pdelay_Req, Pdelay_Resp, Pdelay_Resp_Follow_Up, signaling, or management messages.
master	6	The port is behaving as a master port.
passive	7	The port shall not place any messages on its communication path except for Pdelay_Req, Pdelay_Resp, Pdelay_Resp_Follow_Up, or signaling messages, or management messages that are a required response to another management message
uncalibrated	8	The local port is preparing to synchronize to the master port.
slave	9	The port is synchronizing to the selected master port."

-----

REFERENCE "Section 8.2.5.3.1 portState and 9.2.5 of [1]"

SYNTAX INTEGER {  
initializing(1),  
faulty(2),  
disabled(3),  
listening(4),  
preMaster(5),  
master(6),  
passive(7),  
uncalibrated(8),

```

        slave(9)
    }

```

ClockProfileType ::= TEXTUAL-CONVENTION

STATUS current

DESCRIPTION

"Clock Profile used. From [1] section 3.1.30, Profile is the set of allowed Precision Time Protocol (PTP) features applicable to a device."

REFERENCE "Section 3.1.30 and 19.3 PTP profiles of [1]"

SYNTAX INTEGER {  
     default(1),  
     telecom(2),  
     vendorspecific(3)  
 }

ClockQualityAccuracyType ::= TEXTUAL-CONVENTION

STATUS current

DESCRIPTION

"The ClockQuality as specified in section 5.3.7, 7.6.2.5 and Table 6 of [1].

The following values are not represented in the enumerated values.

```

0x01-0x1F Reserved
0x32-0x7F Reserved

```

It is important to note that section 7.1.1 [RFC 2578] allows for gaps and enumerate values to start with zero when indicated by the protocol."

REFERENCE "Section 5.3.7, 7.6.2.5 and Table 6 of [1]"

SYNTAX INTEGER {  
     reserved00(1),            -- 0  
     nanoSecond25(32),        -- 0x20  
     nanoSecond100(33),       -- 0x21  
     nanoSecond250(34),       -- 0x22  
     microSec1(35),            -- 0x23  
     microSec2dot5(36),       -- 0x24  
     microSec10(37),           -- 0x25  
     microSec25(38),           -- 0x26  
     microSec100(39),         -- 0x27  
     microSec250(40),         -- 0x28  
     milliSec1(41),            -- 0x29  
     milliSec2dot5(42),        -- 0x2A  
     milliSec10(43),           -- 0x2B  
     milliSec25(44),           -- 0x2C



```

        milliSec100(45),      -- 0x2D
        milliSec250(46),     -- 0x2E
        second1(47),         -- 0x2F
        second10(48),        -- 0x30
        secondGreater10(49), -- 0x31
        unknown(254),        -- 0xFE
        reserved255(255)     -- 0xFF
    }

```

ClockQualityClassType ::= TEXTUAL-CONVENTION

```

    DISPLAY-HINT    "d"
    STATUS          current
    DESCRIPTION

```

"The ClockQuality as specified in section 5.3.7, 7.6.2.4 and Table 5 of [1].

Value	Description
0	Reserved to enable compatibility with future versions.
1-5	Reserved
6	Shall designate a clock that is synchronized to a primary reference time source. The timescale distributed shall be PTP. A clockClass 6 clock shall not be a slave to another clock in the domain.
7	Shall designate a clock that has previously been designated as clockClass 6 but that has lost the ability to synchronize to a primary reference time source and is in holdover mode and within holdover specifications. The timescale distributed shall be PTP. A clockClass 7 clock shall not be a slave to another clock in the domain.
8	Reserved.
9-10	Reserved to enable compatibility with future versions.
11-12	Reserved.
13	Shall designate a clock that is synchronized to an application-specific source of time. The timescale distributed shall be ARB. A clockClass 13 clock shall not be a slave to another clock in the domain.
14	Shall designate a clock that has previously been designated as clockClass 13 but that has lost the ability to synchronize to an application-specific source of time and is in holdover mode and within holdover specifications. The timescale distributed

-----

shall be ARB. A clockClass 14 clock shall not be a slave to another clock in the domain.

15-51 Reserved.

52 Degradation alternative A for a clock of clockClass 7 that is not within holdover specification. A clock of clockClass 52 shall not be a slave to another clock in the domain.

53-57 Reserved.

58 Degradation alternative A for a clock of clockClass 14 that is not within holdover specification. A clock of clockClass 58 shall not be a slave to another clock in the domain.

59-67 Reserved.

68-122 For use by alternate PTP profiles.

123-127 Reserved.

128-132 Reserved.

133-170 For use by alternate PTP profiles.

171-186 Reserved.

187 Degradation alternative B for a clock of clockClass 7 that is not within holdover specification. A clock of clockClass 187 may be a slave to another clock in the domain.

188-192 Reserved.

193 Degradation alternative B for a clock of clockClass 14 that is not within holdover specification. A clock of clockClass 193 may be a slave to another clock in the domain.

194-215 Reserved.

216-232 For use by alternate PTP profiles.

233-247 Reserved.

248 Default. This clockClass shall be used if none of the other clockClass definitions apply.

249-250 Reserved.

251 Reserved for ver. 1 compatibility; see Clause 18.

252-254 Reserved.

255 Shall be the clockClass of a slave-only clock; see 9.2.2."

REFERENCE "section 5.3.7, 7.6.2.4 and Table 5 of [1]."  
 SYNTAX Unsigned32 (0..255)

ClockRoleType ::= TEXTUAL-CONVENTION

STATUS current

DESCRIPTION

"The Clock Role. The protocol generates a Master Slave relationship among the clocks in the system.

Clock Role	Value	Description
------------	-------	-------------

```

-----
Master clock      1      A clock that is the source of
                    time to which all other clocks on
                    that path synchronize.

Slave clock      2      A clock which synchronizes to
                    another clock (master).

SYNTAX           INTEGER {
                    master(1),
                    slave(2)
                  }

```

ClockStateType ::= TEXTUAL-CONVENTION

STATUS current

DESCRIPTION

"The clock state returned by PTP engine.

Clock State	Value	Description
Freerun state	1	Applies to a slave device that is not locked to a master. This is the initial state a slave starts out with when it is not getting any PTP packets from the master or because of some other input error (erroneous packets, etc).
Holdover state	2	In this state the slave device is locked to a master but communication with the master is lost or the timestamps in the ptp packets are incorrect. But since the slave was locked to the master, it can run with the same accuracy for some time. The slave can continue to operate in this state for some time. If communication with the master is not restored for a while, the device is moved to the FREERUN state.
Acquiring state	3	The slave device is receiving packets from a master and is trying to acquire a clock.
Freq_locked state	4	The slave device is locked to the master with respect to frequency, but not phase aligned
Phase_aligned state	5	The slave device is locked to the

master with respect to frequency and phase."

```
SYNTAX          INTEGER {
                    freerun(1),
                    holdover(2),
                    acquiring(3),
                    frequencyLocked(4),
                    phaseAligned(5)
                }
```

ClockTimeSourceType ::= TEXTUAL-CONVENTION

STATUS current

DESCRIPTION

"The ClockQuality as specified in section 5.3.7, 7.6.2.6 and Table 7 of [1].

The following values are not represented in the enumerated values.

0xF0-0xFE For use by alternate PTP profiles  
0xFF Reserved

It is important to note that section 7.1.1 [RFC 2578] allows for gaps and enumerate values to start with zero when indicated by the protocol."

REFERENCE "section 5.3.7, 7.6.2.6 and Table 7 of [1]."

```
SYNTAX          INTEGER {
                    atomicClock(16), -- 0x10
                    gps(32), -- 0x20
                    terrestrialRadio(48), -- 0x22
                    ptp(64), -- 0x40
                    ntp(80), -- 0x50
                    handSet(96), -- 0x60
                    other(144), -- 0x90
                    internalOillator(160) -- 0xA0
                }
```

ClockTimeInterval ::= TEXTUAL-CONVENTION

STATUS current

DESCRIPTION

"This textual convention corresponds to the TimeInterval structure indicated in section 5.3.2 of [1]. It will be presented in the form of a character array.

The TimeInterval type represents time intervals.

```
struct TimeInterval
```

```
{
    Integer64 scaledNanoseconds;
};
```

The scaledNanoseconds member is the time interval expressed in units of nanoseconds and multiplied by 2\*\*16.

Positive or negative time intervals outside the maximum range of this data type shall be encoded as the largest positive and negative values of the data type, respectively.

For example, 2.5 ns is expressed as 0000 0000 0002 8000 in Base16."

REFERENCE

"Section 5.3.2 and section 7.7.2.1 Timer interval specification of [1]"

SYNTAX OCTET STRING (SIZE (1..255))

ClockTxModeType ::= TEXTUAL-CONVENTION

STATUS current

DESCRIPTION

"Transmission mode.

unicast. Using unicast communication channel.

multicast. Using Multicast communication channel.

multicast-mix. Using multicast-unicast comm. channel"

```
SYNTAX INTEGER {
    unicast(1),
    multicast(2),
    multicastmix(3)
}
```

ClockType ::= TEXTUAL-CONVENTION

STATUS current

DESCRIPTION

"The clock types as defined in the MIB module description."

REFERENCE "section 6.5.1 of [1]."

```
SYNTAX INTEGER {
    ordinaryClock(1),
    boundaryClock(2),
    transparentClock(3),
    boundaryNode(4)
}
```

#### 4.2. The Notifications Subtree

```
ietfPtpMIB Notifs OBJECT IDENTIFIER
 ::= { ietfPtpMIB 0 }

ietfPtpMIB Objects OBJECT IDENTIFIER
 ::= { ietfPtpMIB 1 }

ietfPtpMIB Conformance OBJECT IDENTIFIER
 ::= { ietfPtpMIB 2 }

ietfPtpMIB SystemInfo OBJECT IDENTIFIER
 ::= { ietfPtpMIB Objects 1 }

-- Conformance Information Definition

ietfPtpMIB Compliances OBJECT IDENTIFIER
 ::= { ietfPtpMIB Conformance 1 }

ietfPtpMIB Groups OBJECT IDENTIFIER
 ::= { ietfPtpMIB Conformance 2 }

ietfPtpMIB Compliances1 MODULE-COMPLIANCE
STATUS          current
DESCRIPTION
  "Compliance statement for agents that provide read-only
  support for IETF-PTP-MIB. Such devices can only be monitored
  using this MIB module.

  The Module is implemented with support for read-only. In
  other words, only monitoring is available by implementing
  this MODULE-COMPLIANCE."
MODULE          -- this module
MANDATORY-GROUPS { ietfPtpMIB SystemInfoGroup }
 ::= { ietfPtpMIB Compliances 1 }

ietfPtpMIB Compliances2 MODULE-COMPLIANCE
STATUS          current
DESCRIPTION
  "Compliance statement for agents that provide read-only
  support for IETF-PTP-MIB. Such devices can only be monitored
  using this MIB module.

  The Module is implemented with support for read-only. In
  other words, only monitoring is available by implementing
  this MODULE-COMPLIANCE."
MODULE          -- this module
MANDATORY-GROUPS {
    ietfPtpMIB ClockCurrentDSGroup,
    ietfPtpMIB ClockParentDSGroup,
    ietfPtpMIB ClockDefaultDSGroup,
```

```

        ietfPtpMIB ClockRunningGroup,
        ietfPtpMIB ClockTimepropertiesGroup
    }
 ::= { ietfPtpMIB Compliances 2 }

ietfPtpMIB Compliances3 MODULE-COMPLIANCE
STATUS          current
DESCRIPTION
    "Compliance statement for agents that provide read-only
    support for IETF-PTP-MIB. Such devices can only be monitored
    using this MIB module.

    The Module is implemented with support for read-only. In
    other words, only monitoring is available by implementing
    this MODULE-COMPLIANCE."
MODULE          -- this module
MANDATORY-GROUPS {
    ietfPtpMIB ClockPortGroup,
    ietfPtpMIB ClockPortDSGroup,
    ietfPtpMIB ClockPortRunningGroup,
    ietfPtpMIB ClockPortAssociateGroup
}
 ::= { ietfPtpMIB Compliances 3 }

ietfPtpMIB Compliances4 MODULE-COMPLIANCE
STATUS          current
DESCRIPTION
    "Compliance statement for agents that provide read-only
    support for IETF-PTP-MIB. Such devices can only be monitored
    using this MIB module.

    The Module is implemented with support for read-only. In
    other words, only monitoring is available by implementing
    this MODULE-COMPLIANCE."
MODULE          -- this module
MANDATORY-GROUPS {
    ietfPtpMIB ClockTranparentDSGroup,
    ietfPtpMIB ClockPortTransDSGroup
}
 ::= { ietfPtpMIB Compliances 4 }

ietfPtpMIB SystemInfoGroup OBJECT-GROUP
OBJECTS        {
    ietfptpSystemDomainTotals,
    ietfptpDomainClockPortsTotal,
    ietfptpSystemProfile
}
STATUS          current
```

```
DESCRIPTION
    "Group which aggregates objects describing system-wide
    information"
 ::= { ietfPtpMIB Groups 1 }

ietfPtpMIB ClockCurrentDSGroup OBJECT-GROUP
OBJECTS
    {
        ietfntpClockCurrentDSStepsRemoved,
        ietfntpClockCurrentDSOffsetFromMaster,
        ietfntpClockCurrentDSMeanPathDelay
    }
STATUS
    current
DESCRIPTION
    "Group which aggregates objects describing PTP Current
    Dataset information"
 ::= { ietfPtpMIB Groups 2 }

ietfPtpMIB ClockParentDSGroup OBJECT-GROUP
OBJECTS
    {
        ietfntpClockParentDSParentPortIdentity,
        ietfntpClockParentDSParentStats,
        ietfntpClockParentDSOffset,
        ietfntpClockParentDSClockPhChRate,
        ietfntpClockParentDSGMClockIdentity,
        ietfntpClockParentDSGMClockPriority1,
        ietfntpClockParentDSGMClockPriority2,
        ietfntpClockParentDSGMClockQualityClass,
        ietfntpClockParentDSGMClockQualityAccuracy,
        ietfntpClockParentDSGMClockQualityOffset
    }
STATUS
    current
DESCRIPTION
    "Group which aggregates objects describing PTP Parent
    Dataset information"
 ::= { ietfPtpMIB Groups 3 }

ietfPtpMIB ClockDefaultDSGroup OBJECT-GROUP
OBJECTS
    {
        ietfntpClockDefaultDSTwoStepFlag,
        ietfntpClockDefaultDSClockIdentity,
        ietfntpClockDefaultDSPriority1,
        ietfntpClockDefaultDSPriority2,
        ietfntpClockDefaultDSSlaveOnly,
        ietfntpClockDefaultDSQualityClass,
        ietfntpClockDefaultDSQualityAccuracy,
        ietfntpClockDefaultDSQualityOffset
    }
STATUS
    current
DESCRIPTION
```



```
        "Group which aggregates objects describing PTP Default
        Dataset information"
 ::= { ietfPtpMIB Groups 4 }

ietfPtpMIB ClockRunningGroup OBJECT-GROUP
  OBJECTS      {
                ietfntpClockRunningState,
                ietfntpClockRunningPacketsSent,
                ietfntpClockRunningPacketsReceived
              }
  STATUS       current
  DESCRIPTION  "Group which aggregates objects describing PTP running state
  information"
 ::= { ietfPtpMIB Groups 5 }

ietfPtpMIB ClockTimepropertiesGroup OBJECT-GROUP
  OBJECTS      {
                ietfntpClockTimePropertiesDSCurrentUTCOffsetValid,
                ietfntpClockTimePropertiesDSCurrentUTCOffset,
                ietfntpClockTimePropertiesDSLeap59,
                ietfntpClockTimePropertiesDSLeap61,
                ietfntpClockTimePropertiesDSTimeTraceable,
                ietfntpClockTimePropertiesDSFreqTraceable,
                ietfntpClockTimePropertiesDSPTPTimescale,
                ietfntpClockTimePropertiesDSSource
              }
  STATUS       current
  DESCRIPTION  "Group which aggregates objects describing PTP Time
  Properties information"
 ::= { ietfPtpMIB Groups 6 }

ietfPtpMIB ClockTransparentDSGroup OBJECT-GROUP
  OBJECTS      {
                ietfntpClockTransDefaultDSClockIdentity,
                ietfntpClockTransDefaultDSNumOfPorts,
                ietfntpClockTransDefaultDSDelay,
                ietfntpClockTransDefaultDSPrimaryDomain
              }
  STATUS       current
  DESCRIPTION  "Group which aggregates objects describing PTP Transparent
  Dataset information"
 ::= { ietfPtpMIB Groups 7 }

ietfPtpMIB ClockPortGroup OBJECT-GROUP
```

```
OBJECTS          {
    ietfntpClockPortName,
    ietfntpClockPortSyncOneStep,
    ietfntpClockPortCurrentPeerAddress,
    ietfntpClockPortNumOfAssociatedPorts,
    ietfntpClockPortCurrentPeerAddressType,
    ietfntpClockPortRole
}
STATUS           current
DESCRIPTION      "Group which aggregates objects describing information for a
                 given PTP Port."
::= { ietfPtpMIB Groups 8 }

ietfPtpMIB ClockPortDSGroup OBJECT-GROUP
OBJECTS          {
    ietfntpClockPortDSName,
    ietfntpClockPortDSPortIdentity,
    ietfntpClockPortDSAnnouncementInterval,
    ietfntpClockPortDSAnnounceRctTimeout,
    ietfntpClockPortDSSyncInterval,
    ietfntpClockPortDSMinDelayReqInterval,
    ietfntpClockPortDSPeerDelayReqInterval,
    ietfntpClockPortDSDelayMech,
    ietfntpClockPortDSPeerMeanPathDelay,
    ietfntpClockPortDSGrantDuration,
    ietfntpClockPortDSPTPVersion
}
STATUS           current
DESCRIPTION      "Group which aggregates objects describing PTP Port Dataset
                 information"
::= { ietfPtpMIB Groups 9 }

ietfPtpMIB ClockPortRunningGroup OBJECT-GROUP
OBJECTS          {
    ietfntpClockPortRunningName,
    ietfntpClockPortRunningState,
    ietfntpClockPortRunningRole,
    ietfntpClockPortRunningInterfaceIndex,
    ietfntpClockPortRunningIPVersion,
    ietfntpClockPortRunningEncapsulationType,
    ietfntpClockPortRunningTxMode,
    ietfntpClockPortRunningRxMode,
    ietfntpClockPortRunningPacketsReceived,
    ietfntpClockPortRunningPacketsSent
}
STATUS           current
DESCRIPTION
```

```

        "Group which aggregates objects describing PTP running
        interface information"
 ::= { ietfPtpMIB Groups 10 }

ietfPtpMIB ClockPortTransDSGroup OBJECT-GROUP
  OBJECTS      {
                ietfntpClockPortTransDSPortIdentity,
                ietfntpClockPortTransDSlogMinPdelayReqInt,
                ietfntpClockPortTransDSFaultyFlag,
                ietfntpClockPortTransDSPeerMeanPathDelay
              }
  STATUS       current
  DESCRIPTION  "Group which aggregates objects describing PTP TransparentDS
                Dataset information"
 ::= { ietfPtpMIB Groups 11 }

ietfPtpMIB ClockPortAssociateGroup OBJECT-GROUP
  OBJECTS      {
                ietfntpClockPortAssociatePacketsSent,
                ietfntpClockPortAssociatePacketsReceived,
                ietfntpClockPortAssociateAddress,
                ietfntpClockPortAssociateAddressType,
                ietfntpClockPortAssociateInErrors,
                ietfntpClockPortAssociateOutErrors
              }
  STATUS       current
  DESCRIPTION  "Group which aggregates objects describing information on
                peer PTP ports for a given PTP clock-port."
 ::= { ietfPtpMIB Groups 12 }

ietfPtpMIB ClockInfo OBJECT IDENTIFIER
 ::= { ietfPtpMIB Objects 2 }

```

#### 4.3. The Table Structures

```

ietfntpSystemTable OBJECT-TYPE
  SYNTAX      SEQUENCE OF ietfntpSystemEntry
  MAX-ACCESS  not-accessible
  STATUS      current
  DESCRIPTION "Table of count information about the PTP system for all
                domains."
 ::= { ietfPtpMIB SystemInfo 1 }

ietfntpSystemEntry OBJECT-TYPE
  SYNTAX      ietfntpSystemEntry
  MAX-ACCESS  not-accessible

```

```

STATUS          current
DESCRIPTION
    "An entry in the table, containing count information about a
    single domain. New row entries are added when the PTP clock
    for this domain is configured, while the unconfiguration of
    the PTP clock removes it."
INDEX           {
                ietfntpDomainIndex,
                ietfntpInstanceIndex
            }
 ::= { ietfntpSystemTable 1 }

ietfntpSystemEntry ::= SEQUENCE {
    ietfntpDomainIndex
ClockDomainType,
    ietfntpInstanceIndex
ClockInstanceType,
    ietfntpDomainClockPortsTotal          Gauge32
}

ietfntpDomainIndex OBJECT-TYPE
SYNTAX          ClockDomainType
MAX-ACCESS     not-accessible
STATUS         current
DESCRIPTION
    "This object specifies the domain number used to create
    logical group of PTP devices. The Clock Domain is a logical
    group of clocks and devices that synchronize with each other
    using the PTP protocol.

    0          Default domain
    1          Alternate domain 1
    2          Alternate domain 2
    3          Alternate domain 3
    4 - 127    User-defined domains
    128 - 255 Reserved"
 ::= { ietfntpSystemEntry 1 }

ietfntpInstanceIndex OBJECT-TYPE
SYNTAX          ClockInstanceType (0..255)
MAX-ACCESS     not-accessible
STATUS         current
DESCRIPTION
    "This object specifies the instance of the Clock for this
    domain."
 ::= { ietfntpSystemEntry 2 }

ietfntpDomainClockPortsTotal OBJECT-TYPE
SYNTAX          Gauge32

```

```
UNITS          "ntp ports"
MAX-ACCESS     read-only
STATUS         current
DESCRIPTION    "This object specifies the total number of clock ports
               configured within a domain."
 ::= { ietfntpSystemEntry 3 }

ietfntpSystemDomainTable OBJECT-TYPE
SYNTAX        SEQUENCE OF ietfntpSystemDomainEntry
MAX-ACCESS     not-accessible
STATUS         current
DESCRIPTION    "Table of information about the PTP system for all clock
               modes -- ordinary, boundary or transparent."
 ::= { ietfNtpMIB SystemInfo 2 }

ietfntpSystemDomainEntry OBJECT-TYPE
SYNTAX        ietfntpSystemDomainEntry
MAX-ACCESS     not-accessible
STATUS         current
DESCRIPTION    "An entry in the table, containing information about a
               single clock mode for the PTP system. A row entry gets added
               when PTP clocks are configured on the router."
INDEX         { ietfntpSystemDomainClockTypeIndex }
 ::= { ietfntpSystemDomainTable 1 }

ietfntpSystemDomainEntry ::= SEQUENCE {
    ietfntpSystemDomainClockTypeIndex ClockType,
    ietfntpSystemDomainTotals         Gauge32
}

ietfntpSystemDomainClockTypeIndex OBJECT-TYPE
SYNTAX        ClockType
MAX-ACCESS     not-accessible
STATUS         current
DESCRIPTION    "This object specifies the clock type as defined in the
               Textual convention description."
 ::= { ietfntpSystemDomainEntry 1 }

ietfntpSystemDomainTotals OBJECT-TYPE
SYNTAX        Gauge32
UNITS         "domains"
MAX-ACCESS     read-only
STATUS         current
DESCRIPTION    "This object specifies the total number of PTP domains for
```

```
        this particular clock type configured in this node."
 ::= { ietfntpSystemDomainEntry 2 }

ietfntpClockNodeTable OBJECT-TYPE
    SYNTAX          SEQUENCE OF ietfntpClockNodeEntry
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "Table of information about the PTP system for a given
        domain."
 ::= { ietfNtpMIB SystemInfo 3 }

ietfntpClockNodeEntry OBJECT-TYPE
    SYNTAX          ietfntpClockNodeEntry
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "An entry in the table, containing information about a
        single domain. A entry is added when a new PTP clock domain
        is configured on the router."
    INDEX           {
                    ietfntpClockDomainIndex,
                    ietfntpClockTypeIndex,
                    ietfntpClockInstanceIndex
                    }
 ::= { ietfntpClockNodeTable 1 }

ietfntpClockNodeEntry ::= SEQUENCE {
    ietfntpClockDomainIndex      ClockDomainType,
    ietfntpClockTypeIndex        ClockType,
    ietfntpClockInstanceIndex    ClockInstanceType,
    ietfntpSystemProfile          ClockProfileType
}

ietfntpClockDomainIndex OBJECT-TYPE
    SYNTAX          ClockDomainType
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the domain number used to create
        logical group of PTP devices."
 ::= { ietfntpClockNodeEntry 1 }

ietfntpClockTypeIndex OBJECT-TYPE
    SYNTAX          ClockType
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the clock type as defined in the
```

```
        Textual convention description."
 ::= { ietfntpClockNodeEntry 2 }

ietfntpClockInstanceIndex OBJECT-TYPE
    SYNTAX          ClockInstanceType (0..255)
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the instance of the Clock for this
        clock type for the given domain."
 ::= { ietfntpClockNodeEntry 3 }

ietfntpSystemProfile OBJECT-TYPE
    SYNTAX          ClockProfileType
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the NTP Profile implemented on the
        system."
    REFERENCE       "Section 19.3 NTP profiles of [1]"
 ::= { ietfNtpMIB SystemInfo 1 }

ietfntpClockCurrentDSTable OBJECT-TYPE
    SYNTAX          SEQUENCE OF ietfntpClockCurrentDSEntry
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "Table of information about the NTP clock Current Datasets
        for all domains."
 ::= { ietfNtpMIB ClockInfo 1 }

ietfntpClockCurrentDSEntry OBJECT-TYPE
    SYNTAX          ietfntpClockCurrentDSEntry
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "An entry in the table, containing information about a
        single NTP clock Current Datasets for a domain."
    REFERENCE       "[IEEE1588-2008] Section 8.2.2 currentDS data set member
        specifications of [1]"
    INDEX           {
                    ietfntpClockCurrentDSDomainIndex,
                    ietfntpClockCurrentDSClockTypeIndex,
                    ietfntpClockCurrentDSInstanceIndex
                    }
 ::= { ietfntpClockCurrentDSTable 1 }
```

```

ietfntpClockCurrentDSEntry ::= SEQUENCE {
    ietfntpClockCurrentDSDomainIndex      ClockDomainType,
    ietfntpClockCurrentDSClockTypeIndex   ClockType,
    ietfntpClockCurrentDSInstanceIndex    ClockInstanceType,
    ietfntpClockCurrentDSStepsRemoved     Counter32,
    ietfntpClockCurrentDSOffsetFromMaster ClockTimeInterval,
    ietfntpClockCurrentDSMeanPathDelay    ClockTimeInterval
}

```

```

ietfntpClockCurrentDSDomainIndex OBJECT-TYPE
    SYNTAX      ClockDomainType
    MAX-ACCESS  not-accessible
    STATUS      current
    DESCRIPTION
        "This object specifies the domain number used to create
        logical group of PTP devices."
    ::= { ietfntpClockCurrentDSEntry 1 }

```

```

ietfntpClockCurrentDSClockTypeIndex OBJECT-TYPE
    SYNTAX      ClockType
    MAX-ACCESS  not-accessible
    STATUS      current
    DESCRIPTION
        "This object specifies the clock type as defined in the
        Textual convention description."
    ::= { ietfntpClockCurrentDSEntry 2 }

```

```

ietfntpClockCurrentDSInstanceIndex OBJECT-TYPE
    SYNTAX      ClockInstanceType (0..255)
    MAX-ACCESS  not-accessible
    STATUS      current
    DESCRIPTION
        "This object specifies the instance of the clock for this
        clock type in the given domain."
    ::= { ietfntpClockCurrentDSEntry 3 }

```

```

ietfntpClockCurrentDSStepsRemoved OBJECT-TYPE
    SYNTAX      Counter32
    UNITS       "steps"
    MAX-ACCESS  read-only
    STATUS      current
    DESCRIPTION
        "The current clock dataset StepsRemoved value.

        This object specifies the distance measured by the number of
        boundary clocks between the local clock and the Foreign
        master as indicated in the stepsRemoved field of Announce
        messages."
    REFERENCE   "[IEEE1588-2008] Section 8.2.2.2 stepsRemoved"

```



```
 ::= { ietfntpClockCurrentDSEntry 4 }

ietfntpClockCurrentDSOffsetFromMaster OBJECT-TYPE
    SYNTAX          ClockTimeInterval
    UNITS           "Time Interval"
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the current clock dataset ClockOffset
        value. The value of the computation of the offset in time
        between
        a slave and a master clock."
    REFERENCE       "[IEEE1588-2008] Section 8.2.2.3 of [1]"
 ::= { ietfntpClockCurrentDSEntry 5 }

ietfntpClockCurrentDSMeanPathDelay OBJECT-TYPE
    SYNTAX          ClockTimeInterval
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the current clock dataset
        MeanPathDelay value.

        The mean path delay between a pair of ports as measure by
        the delay request-response mechanism."
    REFERENCE       "[IEEE1588-2008] Section 8.2.2.4 mean path
delay"
 ::= { ietfntpClockCurrentDSEntry 6 }

ietfntpClockParentDSTable OBJECT-TYPE
    SYNTAX          SEQUENCE OF ietfntpClockParentDSEntry
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "Table of information about the PTP clock Parent Datasets
        for all domains."
 ::= { ietfPtpMIB ClockInfo 2 }

ietfntpClockParentDSEntry OBJECT-TYPE
    SYNTAX          ietfntpClockParentDSEntry
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "An entry in the table, containing information about a
        single PTP clock Parent Datasets for a domain."
    REFERENCE       "Section 8.2.3 parentDS data set member specifications of
[1]"
    INDEX          {
```

```

        ietfntpClockParentDSDomainIndex,
        ietfntpClockParentDSClockTypeIndex,
        ietfntpClockParentDSInstanceIndex
    }
 ::= { ietfntpClockParentDSTable 1 }

ietfntpClockParentDSEntry ::= SEQUENCE {
    ietfntpClockParentDSDomainIndex      ClockDomainType,
    ietfntpClockParentDSClockTypeIndex   ClockType,
    ietfntpClockParentDSInstanceIndex    ClockInstanceType,
    ietfntpClockParentDSParentPortIdentity OCTET STRING,
    ietfntpClockParentDSParentStats      TruthValue,
    ietfntpClockParentDSOffset           ClockIntervalBase2,
    ietfntpClockParentDSClockPhChRate    Integer32,
    ietfntpClockParentDSGMClockIdentity  ClockIdentity,
    ietfntpClockParentDSGMClockPriority1  Integer32,
    ietfntpClockParentDSGMClockPriority2  Integer32,
    ietfntpClockParentDSGMClockQualityClass
ClockQualityClassType,
    ietfntpClockParentDSGMClockQualityAccuracy
ClockQualityAccuracyType,
    ietfntpClockParentDSGMClockQualityOffset Unsigned32
}

ietfntpClockParentDSDomainIndex OBJECT-TYPE
    SYNTAX      ClockDomainType
    MAX-ACCESS  not-accessible
    STATUS      current
    DESCRIPTION
        "This object specifies the domain number used to create
        logical group of NTP devices."
 ::= { ietfntpClockParentDSEntry 1 }

ietfntpClockParentDSClockTypeIndex OBJECT-TYPE
    SYNTAX      ClockType
    MAX-ACCESS  not-accessible
    STATUS      current
    DESCRIPTION
        "This object specifies the clock type as defined in the
        Textual convention description."
 ::= { ietfntpClockParentDSEntry 2 }

ietfntpClockParentDSInstanceIndex OBJECT-TYPE
    SYNTAX      ClockInstanceType (0..255)
    MAX-ACCESS  not-accessible
    STATUS      current
    DESCRIPTION
        "This object specifies the instance of the clock for this
        clock type in the given domain."

```

```
 ::= { ietfntpClockParentDSEntry 3 }

ietfntpClockParentDSParentPortIdentity OBJECT-TYPE
    SYNTAX          OCTET STRING
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the value of portIdentity of the port
        on the master that issues the Sync messages used in
        synchronizing this clock."
    REFERENCE
        "section 8.2.3.2 parentDS.parentPortIdentity of [1]"
    ::= { ietfntpClockParentDSEntry 4 }

ietfntpClockParentDSParentStats OBJECT-TYPE
    SYNTAX          TruthValue
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the Parent Dataset ParentStats value.

        This value indicates whether the values of ParentDSOffset
        and ParentDSClockPhChRate have been measured and are valid.
        A TRUE value shall indicate valid data."
    REFERENCE
        "section 8.2.3.3 parentDS.parentStats of [1]"
    ::= { ietfntpClockParentDSEntry 5 }

ietfntpClockParentDSOffset OBJECT-TYPE
    SYNTAX          ClockIntervalBase2 (-128..127)
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the Parent Dataset
        ParentOffsetScaledLogVariance value.

        This value is the variance of the parent clocks phase as
        measured by the local clock."
    REFERENCE
        "section 8.2.3.4
        parentDS.observedParentOffsetScaledLogVariance [1]"
    ::= { ietfntpClockParentDSEntry 6 }

ietfntpClockParentDSClockPhChRate OBJECT-TYPE
    SYNTAX          Integer32
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the clock's parent dataset
        ParentClockPhaseChangeRate value."
```

This value is an estimate of the parent clocks phase change rate as measured by the slave clock."

## REFERENCE

"section 8.2.3.5

parentDS.observedParentClockPhaseChangeRate of [1]"

::= { ietfntpClockParentDSEntry 7 }

## ietfntpClockParentDSGMClockIdentity OBJECT-TYPE

SYNTAX ClockIdentity

MAX-ACCESS read-only

STATUS current

## DESCRIPTION

"This object specifies the parent dataset Grandmaster clock identity."

## REFERENCE

"section 8.2.3.6 parentDS.grandmasterIdentity of [1]"

::= { ietfntpClockParentDSEntry 8 }

## ietfntpClockParentDSGMClockPriority1 OBJECT-TYPE

SYNTAX Integer32

MAX-ACCESS read-only

STATUS current

## DESCRIPTION

"This object specifies the parent dataset Grandmaster clock priority1."

## REFERENCE

"section 8.2.3.8 parentDS.grandmasterPriority1 of [1]"

::= { ietfntpClockParentDSEntry 9 }

## ietfntpClockParentDSGMClockPriority2 OBJECT-TYPE

SYNTAX Integer32

MAX-ACCESS read-only

STATUS current

## DESCRIPTION

"This object specifies the parent dataset grandmaster clock priority2."

## REFERENCE

"section 8.2.3.9 parentDS.grandmasterPriority2 of [1]"

::= { ietfntpClockParentDSEntry 10 }

## ietfntpClockParentDSGMClockQualityClass OBJECT-TYPE

SYNTAX ClockQualityClassType (0..255)

MAX-ACCESS read-only

STATUS current

## DESCRIPTION

"This object specifies the parent dataset grandmaster clock quality class."

## REFERENCE

```

    "section 8.2.3.7 parentDS.grandmasterClockQuality of [1]"
    ::= { ietfntpClockParentDSEntry 11 }

ietfntpClockParentDSGMCClockQualityAccuracy OBJECT-TYPE
    SYNTAX          ClockQualityAccuracyType
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the parent dataset grandmaster clock
        quality accuracy."
    REFERENCE
        "section 8.2.3.7 parentDS.grandmasterClockQuality of [1]"
    ::= { ietfntpClockParentDSEntry 12 }

ietfntpClockParentDSGMCClockQualityOffset OBJECT-TYPE
    SYNTAX          Unsigned32
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the parent dataset grandmaster clock
        quality offset."
    REFERENCE
        "section 8.2.3.7 parentDS.grandmasterClockQuality of [1]"
    ::= { ietfntpClockParentDSEntry 13 }

ietfntpClockDefaultDSTable OBJECT-TYPE
    SYNTAX          SEQUENCE OF ietfntpClockDefaultDSEntry
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "Table of information about the PTP clock Default Datasets
        for all domains."
    ::= { ietfNtpMIB ClockInfo 3 }

ietfntpClockDefaultDSEntry OBJECT-TYPE
    SYNTAX          ietfntpClockDefaultDSEntry
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "An entry in the table, containing information about a
        single PTP clock Default Datasets for a domain."
    INDEX
        {
            ietfntpClockDefaultDSDomainIndex,
            ietfntpClockDefaultDSClockTypeIndex,
            ietfntpClockDefaultDSInstanceIndex
        }
    ::= { ietfntpClockDefaultDSTable 1 }

ietfntpClockDefaultDSEntry ::= SEQUENCE {
```

```
    ietfntpClockDefaultDSDomainIndex      ClockDomainType,
    ietfntpClockDefaultDSClockTypeIndex  ClockType,
    ietfntpClockDefaultDSInstanceIndex   ClockInstanceType,
    ietfntpClockDefaultDSTwoStepFlag     TruthValue,
    ietfntpClockDefaultDSClockIdentity   ClockIdentity,
    ietfntpClockDefaultDSPriority1      Integer32,
    ietfntpClockDefaultDSPriority2      Integer32,
    ietfntpClockDefaultDSSlaveOnly      TruthValue,
    ietfntpClockDefaultDSQualityClass    ClockQualityClassType,
    ietfntpClockDefaultDSQualityAccuracy ClockQualityAccuracyType,
    ietfntpClockDefaultDSQualityOffset   Integer32
}
```

ietfntpClockDefaultDSDomainIndex OBJECT-TYPE

```
SYNTAX          ClockDomainType
MAX-ACCESS      not-accessible
STATUS          current
DESCRIPTION
    "This object specifies the domain number used to create
    logical group of NTP devices."
 ::= { ietfntpClockDefaultDSEntry 1 }
```

ietfntpClockDefaultDSClockTypeIndex OBJECT-TYPE

```
SYNTAX          ClockType
MAX-ACCESS      not-accessible
STATUS          current
DESCRIPTION
    "This object specifies the clock type as defined in the
    Textual convention description."
 ::= { ietfntpClockDefaultDSEntry 2 }
```

ietfntpClockDefaultDSInstanceIndex OBJECT-TYPE

```
SYNTAX          ClockInstanceType (0..255)
MAX-ACCESS      not-accessible
STATUS          current
DESCRIPTION
    "This object specifies the instance of the clock for this
    clock type in the given domain."
 ::= { ietfntpClockDefaultDSEntry 3 }
```

ietfntpClockDefaultDSTwoStepFlag OBJECT-TYPE

```
SYNTAX          TruthValue
MAX-ACCESS      read-only
STATUS          current
DESCRIPTION
    "This object specifies whether the Two Step process is
    used."
 ::= { ietfntpClockDefaultDSEntry 4 }
```

```
ietfntpClockDefaultDSClockIdentity OBJECT-TYPE
    SYNTAX          ClockIdentity
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the default Datasets clock identity."
    ::= { ietfntpClockDefaultDSEntry 5 }

ietfntpClockDefaultDSPriority1 OBJECT-TYPE
    SYNTAX          Integer32
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the default Datasets clock
        Priority1."
    ::= { ietfntpClockDefaultDSEntry 6 }

ietfntpClockDefaultDSPriority2 OBJECT-TYPE
    SYNTAX          Integer32
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the default Datasets clock
        Priority2."
    ::= { ietfntpClockDefaultDSEntry 7 }

ietfntpClockDefaultDSSlaveOnly OBJECT-TYPE
    SYNTAX          TruthValue
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "Whether the SlaveOnly flag is set."
    ::= { ietfntpClockDefaultDSEntry 8 }

ietfntpClockDefaultDSQualityClass OBJECT-TYPE
    SYNTAX          ClockQualityClassType (0..255)
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the default dataset Quality Class."
    ::= { ietfntpClockDefaultDSEntry 9 }

ietfntpClockDefaultDSQualityAccuracy OBJECT-TYPE
    SYNTAX          ClockQualityAccuracyType
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the default dataset Quality
        Accuracy."
```

```

 ::= { ietfntpClockDefaultDSEntry 10 }

ietfntpClockDefaultDSQualityOffset OBJECT-TYPE
    SYNTAX          Integer32
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the default dataset Quality offset."
 ::= { ietfntpClockDefaultDSEntry 11 }

ietfntpClockRunningTable OBJECT-TYPE
    SYNTAX          SEQUENCE OF ietfntpClockRunningEntry
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "Table of information about the PTP clock Running Datasets
         for all domains."
 ::= { ietfNtpMIB ClockInfo 4 }

ietfntpClockRunningEntry OBJECT-TYPE
    SYNTAX          ietfntpClockRunningEntry
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "An entry in the table, containing information about a
         single PTP clock running Datasets for a domain."
    INDEX           {
                    ietfntpClockRunningDomainIndex,
                    ietfntpClockRunningClockTypeIndex,
                    ietfntpClockRunningInstanceIndex
                    }
 ::= { ietfntpClockRunningTable 1 }

ietfntpClockRunningEntry ::= SEQUENCE {
    ietfntpClockRunningDomainIndex      ClockDomainType,
    ietfntpClockRunningClockTypeIndex   ClockType,
    ietfntpClockRunningInstanceIndex     ClockInstanceType,
    ietfntpClockRunningState             ClockStateType,
    ietfntpClockRunningPacketsSent       Counter64,
    ietfntpClockRunningPacketsReceived   Counter64
}

ietfntpClockRunningDomainIndex OBJECT-TYPE
    SYNTAX          ClockDomainType
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the domain number used to create
         logical group of PTP devices."

```



```
 ::= { ietfntpClockRunningEntry 1 }

ietfntpClockRunningClockTypeIndex OBJECT-TYPE
    SYNTAX          ClockType
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the clock type as defined in the
        Textual convention description."
    ::= { ietfntpClockRunningEntry 2 }

ietfntpClockRunningInstanceIndex OBJECT-TYPE
    SYNTAX          ClockInstanceType (0..255)
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the instance of the clock for this
        clock type in the given domain."
    ::= { ietfntpClockRunningEntry 3 }

ietfntpClockRunningState OBJECT-TYPE
    SYNTAX          ClockStateType
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the Clock state returned by PTP
        engine which was described earlier.

        Freerun state. Applies to a slave device that is not locked
        to a master. This is the initial state a slave starts out
        with when
        it is not getting any PTP packets from the master or because
        of some other input error (erroneous packets, etc).

        Holdover state. In this state the slave device is locked to
        a master but communication with the master is lost or the
        timestamps in the ntp packets are incorrect. But since the
        slave was locked to the master, it can run with the same
        accuracy for sometime. The slave can continue to operate in
        this state for some time. If communication with the master
        is not restored for a while, the device is moved to the
        FREERUN state.

        Acquiring state. The slave device is receiving packets from
        a master and is trying to acquire a lock.

        Freq_locked state. The slave device is locked to the master
        with respect to frequency, but is not phase aligned.
```

Phase\_aligned state. The slave device is locked to the master with respect to frequency and phase."  
 ::= { ietfntpClockRunningEntry 4 }

ietfntpClockRunningPacketsSent OBJECT-TYPE  
SYNTAX Counter64  
MAX-ACCESS read-only  
STATUS current  
DESCRIPTION  
"This object specifies the total number of all packet Unicast and multicast that have been sent out for this clock in this domain for this type."  
 ::= { ietfntpClockRunningEntry 5 }

ietfntpClockRunningPacketsReceived OBJECT-TYPE  
SYNTAX Counter64  
MAX-ACCESS read-only  
STATUS current  
DESCRIPTION  
"This object specifies the total number of all packet Unicast and multicast that have been received for this clock in this domain for this type."  
 ::= { ietfntpClockRunningEntry 6 }

ietfntpClockTimePropertiesDSTable OBJECT-TYPE  
SYNTAX SEQUENCE OF ietfntpClockTimePropertiesDSEntry  
MAX-ACCESS not-accessible  
STATUS current  
DESCRIPTION  
"Table of information about the NTP clock Timeproperties Datasets for all domains."  
 ::= { ietfNtpMIB ClockInfo 5 }

ietfntpClockTimePropertiesDSEntry OBJECT-TYPE  
SYNTAX ietfntpClockTimePropertiesDSEntry  
MAX-ACCESS not-accessible  
STATUS current  
DESCRIPTION  
"An entry in the table, containing information about a single NTP clock timeproperties Datasets for a domain."  
REFERENCE "Section 8.2.4 of [1]"  
INDEX {  
 ietfntpClockTimePropertiesDSDomainIndex,  
 ietfntpClockTimePropertiesDSClockTypeIndex,  
 ietfntpClockTimePropertiesDSInstanceIndex  
 }  
 ::= { ietfntpClockTimePropertiesDSTable 1 }

ietfntpClockTimePropertiesDSEntry ::= SEQUENCE {

```

    ietfntpClockTimePropertiesDSDomainIndex
ClockDomainType,
    ietfntpClockTimePropertiesDSClockTypeIndex      ClockType,
    ietfntpClockTimePropertiesDSInstanceIndex
ClockInstanceType,
    ietfntpClockTimePropertiesDSCurrentUTCOffsetValid TruthValue,
    ietfntpClockTimePropertiesDSCurrentUTCOffset      Integer32,
    ietfntpClockTimePropertiesDSLeap59                 TruthValue,
    ietfntpClockTimePropertiesDSLeap61                 TruthValue,
    ietfntpClockTimePropertiesDSTimeTraceable          TruthValue,
    ietfntpClockTimePropertiesDSFreqTraceable          TruthValue,
    ietfntpClockTimePropertiesDSPTPTimescale           TruthValue,
    ietfntpClockTimePropertiesDSSource
ClockTimeSourceType
}

```

```

ietfntpClockTimePropertiesDSDomainIndex OBJECT-TYPE

```

```

    SYNTAX          ClockDomainType

```

```

    MAX-ACCESS      not-accessible

```

```

    STATUS          current

```

```

    DESCRIPTION

```

```

        "This object specifies the domain number used to create
        logical group of PTP devices."

```

```

    ::= { ietfntpClockTimePropertiesDSEntry 1 }

```

```

ietfntpClockTimePropertiesDSClockTypeIndex OBJECT-TYPE

```

```

    SYNTAX          ClockType

```

```

    MAX-ACCESS      not-accessible

```

```

    STATUS          current

```

```

    DESCRIPTION

```

```

        "This object specifies the clock type as defined in the
        Textual convention description."

```

```

    ::= { ietfntpClockTimePropertiesDSEntry 2 }

```

```

ietfntpClockTimePropertiesDSInstanceIndex OBJECT-TYPE

```

```

    SYNTAX          ClockInstanceType (0..255)

```

```

    MAX-ACCESS      not-accessible

```

```

    STATUS          current

```

```

    DESCRIPTION

```

```

        "This object specifies the instance of the clock for this
        clock type in the given domain."

```

```

    ::= { ietfntpClockTimePropertiesDSEntry 3 }

```

```

ietfntpClockTimePropertiesDSCurrentUTCOffsetValid OBJECT-TYPE

```

```

    SYNTAX          TruthValue

```

```

    MAX-ACCESS      read-only

```

```

    STATUS          current

```

```

    DESCRIPTION

```

```

        "This object specifies the timeproperties dataset value of

```

whether current UTC offset is valid."  
REFERENCE "Section 8.2.4.2 of [1]"  
::= { ietfntpClockTimePropertiesDSEntry 4 }

ietfntpClockTimePropertiesDSCurrentUTCOffset OBJECT-TYPE  
SYNTAX Integer32  
MAX-ACCESS read-only  
STATUS current  
DESCRIPTION  
"This object specifies the timeproperties dataset value of  
current UTC offset.  
  
In PTP systems whose epoch is the PTP epoch, the value of  
timePropertiesDS.currentUtcOffset is the offset  
between TAI and UTC; otherwise the value has no meaning. The  
value shall be in units of seconds.  
The initialization value shall be selected as follows:  
a) If the timePropertiesDS.ptpTimescale (see 8.2.4.8) is  
TRUE, the value is the value obtained from a  
primary reference if the value is known at the time of  
initialization, else.  
b) The value shall be the current number of leap seconds  
(7.2.3) when the node is designed."  
REFERENCE "Section 8.2.4.3 of [1]"  
::= { ietfntpClockTimePropertiesDSEntry 5 }

ietfntpClockTimePropertiesDSLeap59 OBJECT-TYPE  
SYNTAX TruthValue  
MAX-ACCESS read-only  
STATUS current  
DESCRIPTION  
"This object specifies the Leap59 value in the clock Current  
Dataset."  
REFERENCE "Section 8.2.4.4 of [1]"  
::= { ietfntpClockTimePropertiesDSEntry 6 }

ietfntpClockTimePropertiesDSLeap61 OBJECT-TYPE  
SYNTAX TruthValue  
MAX-ACCESS read-only  
STATUS current  
DESCRIPTION  
"This object specifies the Leap61 value in the clock Current  
Dataset."  
REFERENCE "Section 8.2.4.5 of [1]"  
::= { ietfntpClockTimePropertiesDSEntry 7 }

ietfntpClockTimePropertiesDSTimeTraceable OBJECT-TYPE  
SYNTAX TruthValue  
MAX-ACCESS read-only

```
STATUS          current
DESCRIPTION
  "This object specifies the Timetraceable value in the clock
  Current Dataset."
REFERENCE       "Section 8.2.4.6 of [1]"
::= { ietfntpClockTimePropertiesDSEntry 8 }

ietfntpClockTimePropertiesDSFreqTraceable OBJECT-TYPE
SYNTAX          TruthValue
MAX-ACCESS      read-only
STATUS          current
DESCRIPTION
  "This object specifies the Frequency Traceable value in the
  clock Current Dataset."
REFERENCE       "Section 8.2.4.7 of [1]"
::= { ietfntpClockTimePropertiesDSEntry 9 }

ietfntpClockTimePropertiesDSPTPTimescale OBJECT-TYPE
SYNTAX          TruthValue
MAX-ACCESS      read-only
STATUS          current
DESCRIPTION
  "This object specifies the PTP Timescale value in the clock
  Current Dataset."
REFERENCE       "Section 8.2.4.8 of [1]"
::= { ietfntpClockTimePropertiesDSEntry 10 }

ietfntpClockTimePropertiesDSSource OBJECT-TYPE
SYNTAX          ClockTimeSourceType
MAX-ACCESS      read-only
STATUS          current
DESCRIPTION
  "This object specifies the Timesource value in the clock
  Current Dataset."
REFERENCE       "Section 8.2.4.9 of [1]"
::= { ietfntpClockTimePropertiesDSEntry 11 }

ietfntpClockTransDefaultDSTable OBJECT-TYPE
SYNTAX          SEQUENCE OF ietfntpClockTransDefaultDSEntry
MAX-ACCESS      not-accessible
STATUS          current
DESCRIPTION
  "Table of information about the PTP transparent clock
  Default Datasets for all domains."
::= { ietfPtpMIB ClockInfo 6 }

ietfntpClockTransDefaultDSEntry OBJECT-TYPE
SYNTAX          ietfntpClockTransDefaultDSEntry
MAX-ACCESS      not-accessible
```

```

STATUS          current
DESCRIPTION
    "An entry in the table, containing information about a
    single PTP transparent clock Default Datasets for a domain."
REFERENCE       "Section 8.3.2 of [1]"
INDEX          {
                ietfntpClockTransDefaultDSDomainIndex,
                ietfntpClockTransDefaultDSInstanceIndex
            }
 ::= { ietfntpClockTransDefaultDSTable 1 }

ietfntpClockTransDefaultDSEntry ::= SEQUENCE {
    ietfntpClockTransDefaultDSDomainIndex  ClockDomainType,
    ietfntpClockTransDefaultDSInstanceIndex ClockInstanceType,
    ietfntpClockTransDefaultDSClockIdentity ClockIdentity,
    ietfntpClockTransDefaultDSNumOfPorts   Counter32,
    ietfntpClockTransDefaultDSDelay       ClockMechanismType,
    ietfntpClockTransDefaultDSPrimaryDomain Integer32
}

ietfntpClockTransDefaultDSDomainIndex OBJECT-TYPE
SYNTAX          ClockDomainType
MAX-ACCESS     not-accessible
STATUS         current
DESCRIPTION
    "This object specifies the domain number used to create
    logical group of PTP devices."
 ::= { ietfntpClockTransDefaultDSEntry 1 }

ietfntpClockTransDefaultDSInstanceIndex OBJECT-TYPE
SYNTAX          ClockInstanceType (0..255)
MAX-ACCESS     not-accessible
STATUS         current
DESCRIPTION
    "This object specifies the instance of the clock for this
    clock type in the given domain."
 ::= { ietfntpClockTransDefaultDSEntry 2 }

ietfntpClockTransDefaultDSClockIdentity OBJECT-TYPE
SYNTAX          ClockIdentity
MAX-ACCESS     read-only
STATUS         current
DESCRIPTION
    "This object specifies the value of the clockIdentity
    attribute of the local clock."
REFERENCE       "Section 8.3.2.2.1 of [1]"
 ::= { ietfntpClockTransDefaultDSEntry 3 }

ietfntpClockTransDefaultDSNumOfPorts OBJECT-TYPE

```

```
SYNTAX          Counter32
MAX-ACCESS      read-only
STATUS          current
DESCRIPTION     "This object specifies the number of PTP ports of the
                device."
REFERENCE       "Section 8.3.2.2.2 of [1]"
::= { ietfntpClockTransDefaultDSEntry 4 }

ietfntpClockTransDefaultDSDelay OBJECT-TYPE
SYNTAX          ClockMechanismType
MAX-ACCESS      read-only
STATUS          current
DESCRIPTION     "This object, if the transparent clock is an end-to-end
                transparent clock, has the value shall be E2E; If the
                transparent clock is a peer-to-peer transparent clock, the
                value shall be P2P."
REFERENCE       "Section 8.3.2.3.1 of [1]"
::= { ietfntpClockTransDefaultDSEntry 5 }

ietfntpClockTransDefaultDSPrimaryDomain OBJECT-TYPE
SYNTAX          Integer32
MAX-ACCESS      read-only
STATUS          current
DESCRIPTION     "This object specifies the value of the primary
                syntonization domain. The initialization value shall be 0."
REFERENCE       "Section 8.3.2.3.2 of [1]"
::= { ietfntpClockTransDefaultDSEntry 6 }

ietfntpClockPortTable OBJECT-TYPE
SYNTAX          SEQUENCE OF ietfntpClockPortEntry
MAX-ACCESS      not-accessible
STATUS          current
DESCRIPTION     "Table of information about the clock ports for a particular
                domain."
::= { ietfPtpMIB ClockInfo 7 }

ietfntpClockPortEntry OBJECT-TYPE
SYNTAX          ietfntpClockPortEntry
MAX-ACCESS      not-accessible
STATUS          current
DESCRIPTION     "An entry in the table, containing information about a
                single clock port."
INDEX           {
                ietfntpClockPortDomainIndex,
```

```

        ietfntpClockPortClockTypeIndex,
        ietfntpClockPortClockInstanceIndex,
        ietfntpClockPortTablePortNumberIndex
    }
 ::= { ietfntpClockPortTable 1 }

ietfntpClockPortEntry ::= SEQUENCE {
    ietfntpClockPortDomainIndex          ClockDomainType,
    ietfntpClockPortClockTypeIndex       ClockType,
    ietfntpClockPortClockInstanceIndex   ClockInstanceType,
    ietfntpClockPortTablePortNumberIndex ClockPortNumber,
    ietfntpClockPortName                  DisplayString,
    ietfntpClockPortRole                   ClockRoleType,
    ietfntpClockPortSyncOneStep           TruthValue,
    ietfntpClockPortCurrentPeerAddressType InetAddressType,
    ietfntpClockPortCurrentPeerAddress    InetAddress,
    ietfntpClockPortNumOfAssociatedPorts  Gauge32
}

ietfntpClockPortDomainIndex OBJECT-TYPE
    SYNTAX          ClockDomainType
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the domain number used to create
        logical group of PTP devices."
 ::= { ietfntpClockPortEntry 1 }

ietfntpClockPortClockTypeIndex OBJECT-TYPE
    SYNTAX          ClockType
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the clock type as defined in the
        Textual convention description."
 ::= { ietfntpClockPortEntry 2 }

ietfntpClockPortClockInstanceIndex OBJECT-TYPE
    SYNTAX          ClockInstanceType (0..255)
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the instance of the clock for this
        clock type in the given domain."
 ::= { ietfntpClockPortEntry 3 }

ietfntpClockPortTablePortNumberIndex OBJECT-TYPE
    SYNTAX          ClockPortNumber (1..65535)
    MAX-ACCESS      not-accessible

```



```

STATUS          current
DESCRIPTION
  "This object specifies the PTP Portnumber for this port."
 ::= { ietfntpClockPortEntry 4 }

ietfntpClockPortName OBJECT-TYPE
SYNTAX          DisplayString (SIZE (1..64))
MAX-ACCESS      read-only
STATUS          current
DESCRIPTION
  "This object specifies the PTP clock port name configured on
  the router."
 ::= { ietfntpClockPortEntry 5 }

ietfntpClockPortRole OBJECT-TYPE
SYNTAX          ClockRoleType
MAX-ACCESS      read-only
STATUS          current
DESCRIPTION
  "This object describes the current role (slave/master) of
  the port."
 ::= { ietfntpClockPortEntry 6 }

ietfntpClockPortSyncOneStep OBJECT-TYPE
SYNTAX          TruthValue
MAX-ACCESS      read-only
STATUS          current
DESCRIPTION
  "This object specifies that one-step clock operation between
  the PTP master and slave device is enabled."
 ::= { ietfntpClockPortEntry 7 }

ietfntpClockPortCurrentPeerAddressType OBJECT-TYPE
SYNTAX          InetAddressType
MAX-ACCESS      read-only
STATUS          current
DESCRIPTION
  "This object specifies the current peer's network address
  used for PTP communication. Based on the scenario and the
  setup involved, the values might look like these -
  Scenario                Value
  -----                -
  Single Master           master port
  Multiple Masters       selected master port
  Single Slave            slave port
  Multiple Slaves        <empty>

  (In relevant setups, information on available
  slaves and available masters will be available through

```

```

        ietfntpClockPortAssociateTable)"
 ::= { ietfntpClockPortEntry 8 }

ietfntpClockPortCurrentPeerAddress OBJECT-TYPE
    SYNTAX          InetAddress
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the current peer's network address
        used for NTP communication. Based on the scenario and the
        setup involved, the values might look like these -
        Scenario          Value
        -----          -
        Single Master      master port
        Multiple Masters   selected master port
        Single Slave       slave port
        Multiple Slaves    <empty>

        (In relevant setups, information on available
        slaves and available masters will be available through
        ietfntpClockPortAssociateTable)"
 ::= { ietfntpClockPortEntry 9 }

ietfntpClockPortNumOfAssociatedPorts OBJECT-TYPE
    SYNTAX          Gauge32
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies -
        For a master port - the number of NTP slave sessions (peers)
        associated with this NTP port.
        For a slave port - the number of masters available to this
        slave port (might or might not be peered)."
 ::= { ietfntpClockPortEntry 10 }

ietfntpClockPortDSTable OBJECT-TYPE
    SYNTAX          SEQUENCE OF ietfntpClockPortDSEntry
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "Table of information about the clock ports dataset for a
        particular domain."
 ::= { ietfPtpMIB ClockInfo 8 }

ietfntpClockPortDSEntry OBJECT-TYPE
    SYNTAX          ietfntpClockPortDSEntry
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION

```

"An entry in the table, containing port dataset information for a single clock port."

```

INDEX      {
            ietfntpClockPortDSDomainIndex,
            ietfntpClockPortDSClockTypeIndex,
            ietfntpClockPortDSClockInstanceIndex,
            ietfntpClockPortDSPortNumberIndex
          }
 ::= { ietfntpClockPortDSTable 1 }

ietfntpClockPortDSEntry ::= SEQUENCE {
    ietfntpClockPortDSDomainIndex      ClockDomainType,
    ietfntpClockPortDSClockTypeIndex   ClockType,
    ietfntpClockPortDSClockInstanceIndex ClockInstanceType,
    ietfntpClockPortDSPortNumberIndex  ClockPortNumber,
    ietfntpClockPortDSName              DisplayString,
    ietfntpClockPortDSPortIdentity      OCTET STRING,
    ietfntpClockPortDSAnnouncementInterval Integer32,
    ietfntpClockPortDSAnnounceRctTimeout Integer32,
    ietfntpClockPortDSSyncInterval      Integer32,
    ietfntpClockPortDSMinDelayReqInterval Integer32,
    ietfntpClockPortDSPeerDelayReqInterval Integer32,
    ietfntpClockPortDSDelayMech         ClockMechanismType,
    ietfntpClockPortDSPeerMeanPathDelay ClockTimeInterval,
    ietfntpClockPortDSGrantDuration     Unsigned32,
    ietfntpClockPortDSPTPVersion        Integer32
}

ietfntpClockPortDSDomainIndex OBJECT-TYPE
    SYNTAX      ClockDomainType
    MAX-ACCESS  not-accessible
    STATUS      current
    DESCRIPTION
        "This object specifies the domain number used to create
        logical group of PTP devices."
    ::= { ietfntpClockPortDSEntry 1 }

ietfntpClockPortDSClockTypeIndex OBJECT-TYPE
    SYNTAX      ClockType
    MAX-ACCESS  not-accessible
    STATUS      current
    DESCRIPTION
        "This object specifies the clock type as defined in the
        Textual convention description."
    ::= { ietfntpClockPortDSEntry 2 }

ietfntpClockPortDSClockInstanceIndex OBJECT-TYPE
    SYNTAX      ClockInstanceType (0..255)
    MAX-ACCESS  not-accessible

```

```
STATUS          current
DESCRIPTION
  "This object specifies the instance of the clock for this
  clock type in the given domain."
 ::= { ietfntpClockPortDSEntry 3 }

ietfntpClockPortDSPortNumberIndex OBJECT-TYPE
SYNTAX          ClockPortNumber (1..65535)
MAX-ACCESS      not-accessible
STATUS          current
DESCRIPTION
  "This object specifies the PTP portnumber associated with
  this PTP port."
 ::= { ietfntpClockPortDSEntry 4 }

ietfntpClockPortDSName OBJECT-TYPE
SYNTAX          DisplayString (SIZE (1..64))
MAX-ACCESS      read-only
STATUS          current
DESCRIPTION
  "This object specifies the PTP clock port name."
 ::= { ietfntpClockPortDSEntry 5 }

ietfntpClockPortDSPortIdentity OBJECT-TYPE
SYNTAX          OCTET STRING
MAX-ACCESS      read-only
STATUS          current
DESCRIPTION
  "This object specifies the PTP clock port Identity."
 ::= { ietfntpClockPortDSEntry 6 }

ietfntpClockPortDSAnnouncementInterval OBJECT-TYPE
SYNTAX          Integer32
MAX-ACCESS      read-only
STATUS          current
DESCRIPTION
  "This object specifies the Announce message transmission
  interval associated with this clock port."
 ::= { ietfntpClockPortDSEntry 7 }

ietfntpClockPortDSAnnounceRctTimeout OBJECT-TYPE
SYNTAX          Integer32
MAX-ACCESS      read-only
STATUS          current
DESCRIPTION
  "This object specifies the Announce receipt timeout
  associated with this clock port."
 ::= { ietfntpClockPortDSEntry 8 }
```

```
ietfntpClockPortDSSyncInterval OBJECT-TYPE
    SYNTAX          Integer32
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the Sync message transmission
        interval."
    ::= { ietfntpClockPortDSEntry 9 }

ietfntpClockPortDSMinDelayReqInterval OBJECT-TYPE
    SYNTAX          Integer32
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the Delay_Req message transmission
        interval."
    ::= { ietfntpClockPortDSEntry 10 }

ietfntpClockPortDSPeerDelayReqInterval OBJECT-TYPE
    SYNTAX          Integer32
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the Pdelay_Req message transmission
        interval."
    ::= { ietfntpClockPortDSEntry 11 }

ietfntpClockPortDSDelayMech OBJECT-TYPE
    SYNTAX          ClockMechanismType
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the delay mechanism used. If the
        clock is an end-to-end clock, the value of the is e2e, else
        if the clock is a peer to-peer clock, the value shall be
        p2p."
    ::= { ietfntpClockPortDSEntry 12 }

ietfntpClockPortDSPeerMeanPathDelay OBJECT-TYPE
    SYNTAX          ClockTimeInterval
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the peer meanPathDelay."
    ::= { ietfntpClockPortDSEntry 13 }

ietfntpClockPortDSGrantDuration OBJECT-TYPE
    SYNTAX          Unsigned32
    MAX-ACCESS      read-only
```

```

STATUS          current
DESCRIPTION
    "This object specifies the grant duration allocated by the
    master."
 ::= { ietfntpClockPortDSEntry 14 }

ietfntpClockPortDSPTPVersion OBJECT-TYPE
SYNTAX          Integer32
MAX-ACCESS     read-only
STATUS         current
DESCRIPTION
    "This object specifies the PTP version being used."
 ::= { ietfntpClockPortDSEntry 15 }

ietfntpClockPortRunningTable OBJECT-TYPE
SYNTAX          SEQUENCE OF ietfntpClockPortRunningEntry
MAX-ACCESS     not-accessible
STATUS         current
DESCRIPTION
    "Table of information about the clock ports running dataset
    for a particular domain."
 ::= { ietfNtpMIB ClockInfo 9 }

ietfntpClockPortRunningEntry OBJECT-TYPE
SYNTAX          ietfntpClockPortRunningEntry
MAX-ACCESS     not-accessible
STATUS         current
DESCRIPTION
    "An entry in the table, containing running dataset
    information about a single clock port."
INDEX          {
                ietfntpClockPortRunningDomainIndex,
                ietfntpClockPortRunningClockTypeIndex,
                ietfntpClockPortRunningClockInstanceIndex,
                ietfntpClockPortRunningPortNumberIndex
            }
 ::= { ietfntpClockPortRunningTable 1 }

ietfntpClockPortRunningEntry ::= SEQUENCE {
    ietfntpClockPortRunningDomainIndex      ClockDomainType,
    ietfntpClockPortRunningClockTypeIndex   ClockType,
    ietfntpClockPortRunningClockInstanceIndex ClockInstanceType,
    ietfntpClockPortRunningPortNumberIndex  ClockPortNumber,
    ietfntpClockPortRunningName            DisplayString,
    ietfntpClockPortRunningState           ClockPortState,
    ietfntpClockPortRunningRole            ClockRoleType,

```

```
        ietfntpClockPortRunningInterfaceIndex
InterfaceIndexOrZero,
        ietfntpClockPortRunningIPversion          Integer32,
        ietfntpClockPortRunningEncapsulationType  Integer32,
        ietfntpClockPortRunningTxMode            ClockTxModeType,
        ietfntpClockPortRunningRxMode            ClockTxModeType,
        ietfntpClockPortRunningPacketsReceived    Counter64,
        ietfntpClockPortRunningPacketsSent        Counter64
}

```

ietfntpClockPortRunningDomainIndex OBJECT-TYPE

```
SYNTAX          ClockDomainType
MAX-ACCESS      not-accessible
STATUS          current
DESCRIPTION
    "This object specifies the domain number used to create
    logical group of NTP devices."
 ::= { ietfntpClockPortRunningEntry 1 }

```

ietfntpClockPortRunningClockTypeIndex OBJECT-TYPE

```
SYNTAX          ClockType
MAX-ACCESS      not-accessible
STATUS          current
DESCRIPTION
    "This object specifies the clock type as defined in the
    Textual convention description."
 ::= { ietfntpClockPortRunningEntry 2 }

```

ietfntpClockPortRunningClockInstanceIndex OBJECT-TYPE

```
SYNTAX          ClockInstanceType (0..255)
MAX-ACCESS      not-accessible
STATUS          current
DESCRIPTION
    "This object specifies the instance of the clock for this
    clock type in the given domain."
 ::= { ietfntpClockPortRunningEntry 3 }

```

ietfntpClockPortRunningPortNumberIndex OBJECT-TYPE

```
SYNTAX          ClockPortNumber (1..65535)
MAX-ACCESS      not-accessible
STATUS          current
DESCRIPTION
    "This object specifies the NTP portnumber associated with
    this clock port."
 ::= { ietfntpClockPortRunningEntry 4 }

```

ietfntpClockPortRunningName OBJECT-TYPE

```
SYNTAX          DisplayString (SIZE (1..64))
MAX-ACCESS      read-only

```

```
STATUS          current
DESCRIPTION
  "This object specifies the PTP clock port name."
 ::= { ietfntpClockPortRunningEntry 5 }

ietfntpClockPortRunningState OBJECT-TYPE
SYNTAX          ClockPortState
MAX-ACCESS      read-only
STATUS          current
DESCRIPTION
  "This object specifies the port state returned by PTP
  engine.

  initializing - In this state a port initializes
                 its data sets, hardware, and
                 communication facilities.
  faulty       - The fault state of the protocol.
  disabled     - The port shall not place any
                 messages on its communication path.
  listening    - The port is waiting for the
                 announceReceiptTimeout to expire or
                 to receive an Announce message from
                 a master.
  preMaster    - The port shall behave in all respects
                 as though it were in the MASTER state
                 except that it shall not place any
                 messages on its communication path
                 except for Pdelay_Req, Pdelay_Resp,
                 Pdelay_Resp_Follow_Up, signaling, or
                 management messages.
  master       - The port is behaving as a master port.
  passive      - The port shall not place any
                 messages on its communication path
                 except for Pdelay_Req, Pdelay_Resp,
                 Pdelay_Resp_Follow_Up, or signaling
                 messages, or management messages
                 that are a required response to
                 another management message
  uncalibrated - The local port is preparing to
                 synchronize to the master port.
  slave        - The port is synchronizing to the
                 selected master port."
 ::= { ietfntpClockPortRunningEntry 6 }

ietfntpClockPortRunningRole OBJECT-TYPE
SYNTAX          ClockRoleType
MAX-ACCESS      read-only
STATUS          current
DESCRIPTION
```



```
        "This object specifies the Clock Role."
 ::= { ietfntpClockPortRunningEntry 7 }

ietfntpClockPortRunningInterfaceIndex OBJECT-TYPE
    SYNTAX      InterfaceIndexOrZero
    MAX-ACCESS  read-only
    STATUS      current
    DESCRIPTION
        "This object specifies the interface on the router being
        used by the PTP Clock for PTP communication."
 ::= { ietfntpClockPortRunningEntry 8 }

ietfntpClockPortRunningIPversion OBJECT-TYPE
    SYNTAX      Integer32
    MAX-ACCESS  read-only
    STATUS      current
    DESCRIPTION
        "This object specifies the IP version being used for PTP
        communication (the mapping used)."
 ::= { ietfntpClockPortRunningEntry 9 }

ietfntpClockPortRunningEncapsulationType OBJECT-TYPE
    SYNTAX      Integer32
    MAX-ACCESS  read-only
    STATUS      current
    DESCRIPTION
        "This object specifies the type of encapsulation if the
        interface is adding extra layers (eg. VLAN, Pseudowire
        encapsulation...) for the PTP messages."
 ::= { ietfntpClockPortRunningEntry 10 }

ietfntpClockPortRunningTxMode OBJECT-TYPE
    SYNTAX      ClockTxModeType
    MAX-ACCESS  read-only
    STATUS      current
    DESCRIPTION
        "This object specifies the clock transmission mode as

        unicast:      Using unicast communication channel.
        multicast:    Using Multicast communication channel.
        multicast-mix: Using multicast-unicast comm. channel"
 ::= { ietfntpClockPortRunningEntry 11 }

ietfntpClockPortRunningRxMode OBJECT-TYPE
    SYNTAX      ClockTxModeType
    MAX-ACCESS  read-only
    STATUS      current
    DESCRIPTION
        "This object specifies the clock receive mode as
```

```
unicast:      Using unicast communication channel.
multicast:    Using Multicast communication channel.
multicast-mix: Using multicast-unicast comm. channel"
 ::= { ietfntpClockPortRunningEntry 12 }
```

ietfntpClockPortRunningPacketsReceived OBJECT-TYPE

```
SYNTAX      Counter64
UNITS       "packets"
MAX-ACCESS  read-only
STATUS      current
DESCRIPTION
    "This object specifies the packets received on the clock
    port (cumulative)."
```

```
 ::= { ietfntpClockPortRunningEntry 13 }
```

ietfntpClockPortRunningPacketsSent OBJECT-TYPE

```
SYNTAX      Counter64
UNITS       "packets"
MAX-ACCESS  read-only
STATUS      current
DESCRIPTION
    "This object specifies the packets sent on the clock port
    (cumulative)."
```

```
 ::= { ietfntpClockPortRunningEntry 14 }
```

ietfntpClockPortTransDSTable OBJECT-TYPE

```
SYNTAX      SEQUENCE OF ietfntpClockPortTransDSEntry
MAX-ACCESS  not-accessible
STATUS      current
DESCRIPTION
    "Table of information about the transparent clock ports
    running dataset for a particular domain."
```

```
 ::= { ietfNtpMIB ClockInfo 10 }
```

ietfntpClockPortTransDSEntry OBJECT-TYPE

```
SYNTAX      ietfntpClockPortTransDSEntry
MAX-ACCESS  not-accessible
STATUS      current
DESCRIPTION
    "An entry in the table, containing clock port transparent
    dataset information about a single clock port"
```

```
INDEX      {
            ietfntpClockPortTransDSDomainIndex,
            ietfntpClockPortTransDSInstanceIndex,
            ietfntpClockPortTransDSPortNumberIndex
            }
```

```
 ::= { ietfntpClockPortTransDSTable 1 }
```

```
ietfptpClockPortTransDSEntry ::= SEQUENCE {
    ietfptpClockPortTransDSDomainIndex      ClockDomainType,
    ietfptpClockPortTransDSInstanceIndex    ClockInstanceType,
    ietfptpClockPortTransDSPortNumberIndex  ClockPortNumber,
    ietfptpClockPortTransDSPortIdentity     ClockIdentity,
    ietfptpClockPortTransDSlogMinPdelayReqInt Integer32,
    ietfptpClockPortTransDSFaultyFlag      TruthValue,
    ietfptpClockPortTransDSPeerMeanPathDelay ClockTimeInterval
}
```

```
ietfptpClockPortTransDSDomainIndex OBJECT-TYPE
    SYNTAX          ClockDomainType
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the domain number used to create
        logical group of PTP devices."
    ::= { ietfptpClockPortTransDSEntry 1 }
```

```
ietfptpClockPortTransDSInstanceIndex OBJECT-TYPE
    SYNTAX          ClockInstanceType (0..255)
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the instance of the clock for this
        clock type in the given domain."
    ::= { ietfptpClockPortTransDSEntry 2 }
```

```
ietfptpClockPortTransDSPortNumberIndex OBJECT-TYPE
    SYNTAX          ClockPortNumber (1..65535)
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the PTP port number associated with
        this port."
    REFERENCE       "Section 7.5.2 Port Identity of [1]"
    ::= { ietfptpClockPortTransDSEntry 3 }
```

```
ietfptpClockPortTransDSPortIdentity OBJECT-TYPE
    SYNTAX          ClockIdentity
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the value of the PortIdentity
        attribute of the local port."
    REFERENCE       "Section 8.3.3.2.1 of [1]"
    ::= { ietfptpClockPortTransDSEntry 4 }
```

```
ietfptpClockPortTransDSlogMinPdelayReqInt OBJECT-TYPE
```

SYNTAX Integer32  
MAX-ACCESS read-only  
STATUS current  
DESCRIPTION  
    "This object specifies the value of the logarithm to the  
    base 2 of the minPdelayReqInterval."  
REFERENCE "Section 8.3.3.3.1 of [1]"  
::= { ietfptpClockPortTransDSEntry 5 }

ietfptpClockPortTransDSFaultyFlag OBJECT-TYPE

SYNTAX TruthValue  
MAX-ACCESS read-only  
STATUS current  
DESCRIPTION  
    "This object specifies the value TRUE if the port is faulty  
    and FALSE if the port is operating normally."  
REFERENCE "Section 8.3.3.3.2 of [1]"  
::= { ietfptpClockPortTransDSEntry 6 }

ietfptpClockPortTransDSPeerMeanPathDelay OBJECT-TYPE

SYNTAX ClockTimeInterval  
MAX-ACCESS read-only  
STATUS current  
DESCRIPTION  
    "This object specifies, (if the delayMechanism used is P2P)  
    the value is the estimate of the current one-way propagation  
    delay, i.e. <meanPathDelay> on the link attached to this  
    port computed using the peer delay mechanism. If the value  
    of the delayMechanism used is E2E, then the value will be  
    zero."  
REFERENCE "Section 8.3.3.3.3 of [1]"  
::= { ietfptpClockPortTransDSEntry 7 }

ietfptpClockPortAssociateTable OBJECT-TYPE

SYNTAX SEQUENCE OF ietfptpClockPortAssociateEntry  
MAX-ACCESS not-accessible  
STATUS current  
DESCRIPTION  
    "Table of information about a given port's associated ports.  
  
    For a master port - multiple slave ports which have  
    established sessions with the current master port.  
    For a slave port - the list of master clocks available for a  
    given slave port.  
  
    Session information (pkts, errors) to be displayed based on  
    availability and scenario."  
::= { ietfPtpMIB ClockInfo 11 }

```
ietfntpClockPortAssociateEntry OBJECT-TYPE
    SYNTAX          ietfntpClockPortAssociateEntry
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION     "An entry in the table, containing information about a
                    single associated port for the given clockport."
    INDEX           {
                    ietfntpClockPortCurrentDomainIndex,
                    ietfntpClockPortCurrentClockTypeIndex,
                    ietfntpClockPortCurrentClockInstanceIndex,
                    ietfntpClockPortCurrentPortNumberIndex,
                    ietfntpClockPortAssociatePortIndex
                    }
    ::= { ietfntpClockPortAssociateTable 1 }

ietfntpClockPortAssociateEntry ::= SEQUENCE {
    ietfntpClockPortCurrentDomainIndex      ClockDomainType,
    ietfntpClockPortCurrentClockTypeIndex   ClockType,
    ietfntpClockPortCurrentClockInstanceIndex ClockInstanceType,
    ietfntpClockPortCurrentPortNumberIndex  ClockPortNumber,
    ietfntpClockPortAssociatePortIndex      Unsigned32,
    ietfntpClockPortAssociateAddressType    InetAddressType,
    ietfntpClockPortAssociateAddress        InetAddress,
    ietfntpClockPortAssociatePacketsSent    Counter64,
    ietfntpClockPortAssociatePacketsReceived Counter64,
    ietfntpClockPortAssociateInErrors       Counter64,
    ietfntpClockPortAssociateOutErrors      Counter64
}

ietfntpClockPortCurrentDomainIndex OBJECT-TYPE
    SYNTAX          ClockDomainType
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION     "This object specifies the given port's domain number."
    ::= { ietfntpClockPortAssociateEntry 1 }

ietfntpClockPortCurrentClockTypeIndex OBJECT-TYPE
    SYNTAX          ClockType
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION     "This object specifies the given port's clock type."
    ::= { ietfntpClockPortAssociateEntry 2 }

ietfntpClockPortCurrentClockInstanceIndex OBJECT-TYPE
    SYNTAX          ClockInstanceType
    MAX-ACCESS      not-accessible
```

```
STATUS          current
DESCRIPTION
  "This object specifies the instance of the clock for this
  clock type in the given domain."
 ::= { ietfntpClockPortAssociateEntry 3 }

ietfntpClockPortCurrentPortNumberIndex OBJECT-TYPE
SYNTAX          ClockPortNumber
MAX-ACCESS     not-accessible
STATUS         current
DESCRIPTION
  "This object specifies the NTP Port Number for the given
  port."
 ::= { ietfntpClockPortAssociateEntry 4 }

ietfntpClockPortAssociatePortIndex OBJECT-TYPE
SYNTAX          Unsigned32 (1..65535)
MAX-ACCESS     not-accessible
STATUS         current
DESCRIPTION
  "This object specifies the associated port's serial number
  in the current port's context."
 ::= { ietfntpClockPortAssociateEntry 5 }

ietfntpClockPortAssociateAddressType OBJECT-TYPE
SYNTAX          InetAddressType
MAX-ACCESS     read-only
STATUS         current
DESCRIPTION
  "This object specifies the peer port's network address type
  used for NTP communication."
 ::= { ietfntpClockPortAssociateEntry 6 }

ietfntpClockPortAssociateAddress OBJECT-TYPE
SYNTAX          InetAddress
MAX-ACCESS     read-only
STATUS         current
DESCRIPTION
  "This object specifies the peer port's network address used
  for NTP communication."
 ::= { ietfntpClockPortAssociateEntry 7 }

ietfntpClockPortAssociatePacketsSent OBJECT-TYPE
SYNTAX          Counter64
UNITS          "packets"
MAX-ACCESS     read-only
STATUS         current
DESCRIPTION
  "The number of packets sent to this peer port from the
```

```
        current port."
 ::= { ietfntpClockPortAssociateEntry 8 }

ietfntpClockPortAssociatePacketsReceived OBJECT-TYPE
    SYNTAX          Counter64
    UNITS           "packets"
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "The number of packets received from this peer port by the
        current port."
 ::= { ietfntpClockPortAssociateEntry 9 }

ietfntpClockPortAssociateInErrors OBJECT-TYPE
    SYNTAX          Counter64
    UNITS           "packets"
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the input errors associated with the
        peer port."
 ::= { ietfntpClockPortAssociateEntry 10 }

ietfntpClockPortAssociateOutErrors OBJECT-TYPE
    SYNTAX          Counter64
    UNITS           "packets"
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the output errors associated with the
        peer port."
 ::= { ietfntpClockPortAssociateEntry 11 }

END
```

## 5. Relationship to the Interface MIB

This section clarifies the relationship of this MIB to the Interfaces MIB [RFC 2233]. Several areas of correlation are addressed in the following subsections. The implementor is referred to the Interfaces MIB document in order to understand the general intent of these areas.

## 6. MIB modules required for IMPORTS

```
MODULE-IDENTITY,  
OBJECT-TYPE,  
Integer32,  
Gauge32,  
Unsigned32,  
Counter32,  
Counter64  
    FROM SNMPv2-SMI  
OBJECT-GROUP,  
MODULE-COMPLIANCE  
    FROM SNMPv2-CONF  
TEXTUAL-CONVENTION,  
TruthValue,  
DisplayString  
    FROM SNMPv2-TC  
InetAddressType,  
InetAddress  
    FROM INET-ADDRESS-MIB  
InterfaceIndexOrZero  
    FROM IETF-TC
```

## 7. Security Considerations

This MIB contains readable objects whose values provide information related to PTP objects. While unauthorized access to the readable objects is relatively innocuous, unauthorized access to the writeable objects could cause a denial of service, or could cause unauthorized creation and/or manipulation of tunnels. Hence, the support for SET operations in a non-secure environment without proper protection can have a negative effect on network operations.

SNMPv1 by itself is such an insecure environment. Even if the network itself is secure (for example by using IPSec), even then, there is no control as to who on the secure network is allowed to access and SET (change/create/delete) the objects in this MIB.

It is recommended that the implementers consider the security features as provided by the SNMPv3 framework. Specifically, the use of the User-based Security Model [RFC 2574] and the View-based Access Control Model [RFC 2575] is recommended.

It is then a customer/user responsibility to ensure that the SNMP entity giving access to this MIB, is properly configured to give access to those objects only to those principals (users) that have legitimate rights to access them.



## 8. IANA Considerations

To be added.

## 9. References

### 9.1. Normative References

[IEEE1588-2008] "Standard for A Precision Clock Synchronization Protocol for Networked Measurement and Control Systems", IEEE1588-2008.

### 9.2. Informative References

[RFC 1155] Rose, M., and K. McCloghrie, "Structure and Identification of Management Information for TCP/IP-based Internets", RFC 1155, Performance Systems International, Hughes LAN Systems, May 1990

[RFC 1157] Case, J., Fedor, M., Schoffstall, M., and J. Davin, "Simple Network Management Protocol", RFC 1157, SNMP Research, Performance Systems International, Performance Systems International, MIT Laboratory for Computer Science, May 1990.

[RFC 1212] Rose, M., and K. McCloghrie, "Concise MIB Definitions", RFC 1212, Performance Systems International, Hughes LAN Systems, March 1991

[RFC 1215] M. Rose, "A Convention for Defining Traps for use with the SNMP", RFC 1215, Performance Systems International, March 1991

[RFC 1305] David L. Mills, "Network Time Protocol (Version 3) - Specification, Implementation and Analysis", RFC 1305, University of Delaware, March 1992.

[RFC 1901] SNMPv2 Working Group, Case, J., McCloghrie, K., Rose, M., and S. Waldbusser, "Introduction to Community-based SNMPv2", RFC 1901, SNMP Research, Inc., Cisco Systems, Inc., Dover Beach Consulting, Inc., International Network Services, January 1996.

[RFC 1905] SNMPv2 Working Group, Case, J., McCloghrie, K., Rose, M., and S. Waldbusser, "Protocol Operations for Version 2 of the Simple Network Management Protocol (SNMPv2)", RFC 1905, SNMP Research, Inc., Cisco Systems, Inc., Dover Beach Consulting, Inc., International Network Services, January 1996.

[RFC 1906] SNMPv2 Working Group, Case, J., McCloghrie, K., Rose, M., and S. Waldbusser, "Transport Mappings for Version 2 of the Simple Network Management Protocol (SNMPv2)", RFC 1906, SNMP Research,

Inc., Cisco Systems, Inc., Dover Beach Consulting, Inc.,  
International Network Services, January 1996.

[RFC 2233] "The Interfaces Group MIB using SMIV2", K. McCloghrie, F.  
Kastenholz, November 1997, RFC 2233

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#### Abstract

Cellular networks often use Internet standard technologies to handle synchronization. This document analyses the need for security methods for synchronization messages distributed over the Internet. This document also gives a solution on how to mark the synchronization message when IPSec is implemented in end to end frequency synchronization."

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## 1. Introduction

When transferring timing in internet, a shared infrastructure is used, and hence the path is no longer physically deterministic. It leaves open the possibility to disrupt, corrupt or even spoof the timing flow, where a timing signal purports to come from a higher quality clock than it actually does. In the extreme, this may be used to attack the integrity of the network, to disrupt the synchronization flow, or cause authentication failures. On the other hand, it may be possible for unauthorized users to request service from a clock server. This may overload a clock server and compromise its ability to deliver timing to authorized users.

For the cellular backhaul applications, two kinds of synchronization is needed, one is the recovery of an accurate and stable frequency synchronization signal as a reference for the radio signal (e.g. GSM, UMTS FDD, LTE FDD). In addition to frequency synchronization, phase/time synchronization are also needed in Mobile technologies, This is the case for the TDD technologies such as UMTS TDD, LTE TDD.

Frequency synchronization is normally implemented in an end-to-end scenario where none of the intermediate nodes in the network have to recognize and process the synchronization packets. However In phase/time synchronization, a hop-by-hop scenario will request intermediate nodes to process the synchronization packets If very accurate phase/time is needed (e.g. sub-microsecond accuracy).

Femtocell is the typical cellular backhaul application that requires time synchronization. A Femtocell is defined as a wireless base station for deployment in residential environments and is typically connected to the mobile core network via a public broadband connection (eg., DSL modem, cable modem). Femtocell improves cellular network coverage and saves cost for operators. Just like a typical macrocell (larger base station), Femtocells (residential base stations) require a certain level of synchronization (frequency or phase/time) on the air interface, predominantly frequency requirements.

The [3GPP.33.320] specification defines some of the high-level network architecture aspects of a Home NodeB (3G UMTS) and a Home eNodeB (4G LTE). In addition, the Femto Forum organization also provides a network reference model very similar to 3GPP. Both architectures have commonalities as illustrated in Figure 1.



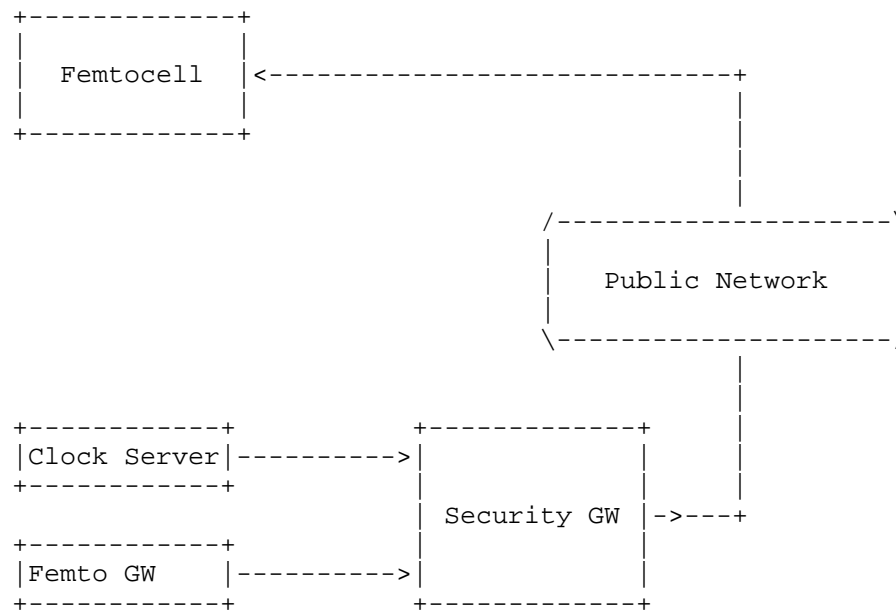


Figure 1. Typical Architecture of a Femtocell Network

The network architecture shows that a public network is used to establish connectivity between Femtocell and core network elements (e.g., Security Gateway, Femto Gateway, Clock server, etc.). With respect to synchronization process, Femtocell will therefore see synchronization messages exchanged over the public network (e.g, Internet). This presents a set of unique challenges for mobile operators.

One challenge involves the security aspects of such the Femto architecture. In both reference models, the communication between Femtocell and Femto Gateway is secured by a mandatory Security Gateway function. The Security Gateway is mandatory since the Femto Gateway and Clock server communicate to Femtocell via a public backhaul broadband connection (also known as the 3GPP iuh interface or Femto Forum Fa interface). The [3GPP.33.320] specification requires that the Femtocell SHALL support receiving time synchronization messages over the secure backhaul link between Femtocell and the Security Gateway, and Femtocell SHALL use IKEv2 protocol to set up at least one IPsec tunnel to protect the traffic with Security Gateway.

This document provides analysis on security requirements for packet-based synchronization and proposes IPsec security solution for end to end frequency synchronization.

## 2. Terminology used in this document

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

## 3. Security requirements for synchronization

The ITUT [G.8265] specification provides general consideration on synchronization security. Because packet-based timing streams may be observed at different points in the network, there may be cases where timing packets flow across multiple network domains which may introduce specific security requirements. There may also be aspects of security that may be related to both the network (e.g. authentication and/or authorization) and to the synchronization protocol itself. ITUT [G.8265] specification recommends to use existing, standards-based security techniques to help ensure the integrity of the synchronization. Examples may include encryption and/or authentication techniques, or network techniques for separating traffic, such as VLANs or LSPs. Specifically for the performance issue, it may not be possible to implement some security requirements without actually degrading the overall level of timing or system performance. From above analysis, following synchronizations requirements are listed:

1. synchronization client SHOULD be prevented from connecting to rogue clock servers
2. clock servers SHOULD be prevented from providing service to unauthorized synchronization client
3. Security mechanisms to achieve synchronization SHOULD minimize any degradation in performance and this side effect SHOULD be controlled to meet specific synchronization requirements(e.g., Femtocell synchronization)

## 4. Security mechanism for synchronization

There are mainly two kinds of security mechanism used in current synchronization: authentication-based and encryption-based.

For the authentication-based security mechanism, a shared secret key between the synchronization client and the clock servers is used to compute an authentication code (known as an "Integrity Check Value",

ICV) over the entire message datagram. [IEEE1588] contains an experimental security annex defining an authentication-based approach. This approach also implements a challenge-response mechanism to confirm the creation of any security association (SA) between a clock servers and a synchronization client. A limitation of the process is that no method of sharing the key is proposed in [IEEE1588]. This MUST be handled by other means.

For the encryption-based security mechanism, a shared-key approach is also used. Instead of creating an ICV, the shared key is used to encrypt the contents of the packet completely. The encryption might be performed in the synchronization device itself, or it might be performed in a separate device, e.g. a secure gateway. An example might be where the timing packets have to pass through an encrypted tunnel (e.g. an IPSec tunnel). Full encryption might be required for various reasons. The contents of the packet may be considered secret, such as might be the case where accuracy of the time distribution is being sold as a service. Alternatively, it may be because other traffic from a device is considered secret, and hence it is easier to encrypt all traffic.

IPsec, as a popular security mechanism, is being considered in some mobile applications, especially in case of unsecure backhaul links (e.g. Femtocells, [3GPP.33.320]) being involved. IPsec can provide data source authentication, confidentiality, integrity that is suitable to end to end synchronization without intermediate nodes. For example, if only frequency synchronization is needed, an end-to-end scenario where none of the intermediate nodes in the network have to recognise and process the synchronization packets might be suitable to use IPsec security mechanism. In this case, the synchronization packets will be encrypted if the packet is transported in the IPSec tunnel.

IPsec can meet synchronization requirement 1 and 2 in section 3, however IPsec still need some enhancement to meet requirement 3. Normally, device will decrypt IPSec message in IP layer, but in order to improve the synchronization accuracy, some synchronization protocol (e.g. [IEEE1588]) requests to process the synchronization message in hardware, therefore the synchronization device may need to identify synchronization messages in physical layer before the message is decrypted. How to identify the synchronization messages in IPsec becomes the most important issue to keep the synchronization accuracy in IPsec synchronization scenario.

#### 5. ESP format enhancement

As discussed above section, it has advantage to identify whether the

tunnel packets received by synchronization client are the special timing packets or not. This section proposes a solution to identify the timing packets When using IPsec to protect the whole time synchronization message. The main thought is to use time packet identifier which is included in a new defined flexible ESP format to identify whether the received data packet is a timing packet or not.

5.1. Existing ESP format

ESP provides confidentiality, data integrity, access control, and data source authentication to IP datagrams as specified in [RFC4303]. The ESP contains several parts (Figure 2): Security Parameters Index(SPI) and Sequence Number(SQN),ESP Payload,ESP trailer and the ICV. SPI and SQN are used to identity a SA and replay protection respectively. ESP trailer is comprised of Padding, Pad length, Next Header. The integrity scope is from SPI to Next Header. The encryption protection is provided for the Payload Data and ESP trailer. For SPI and SQN, only the authentication of data integrity is provided.

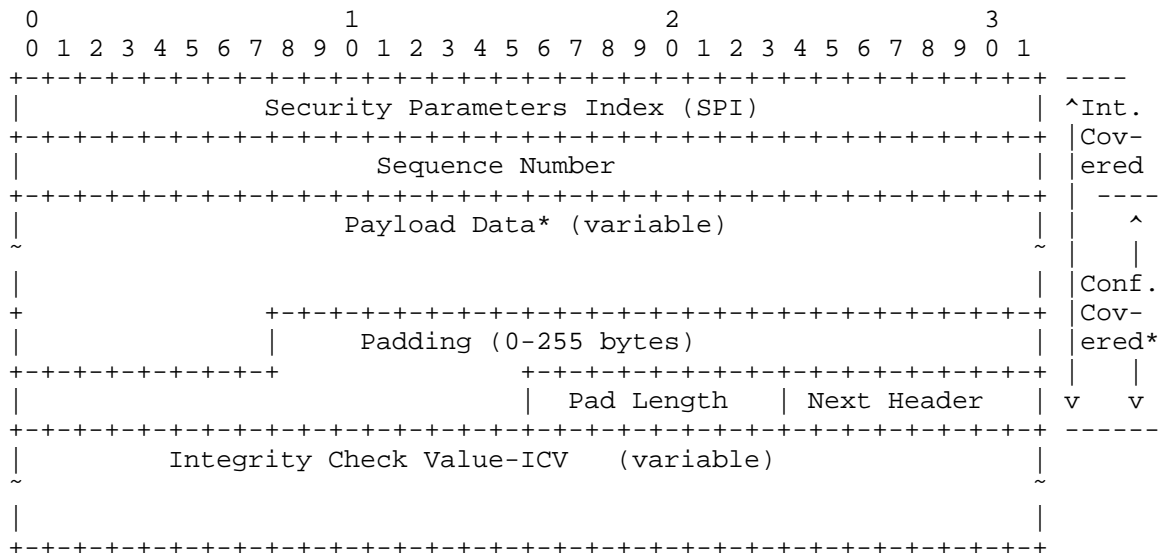


Figure 2. Top-Level Format of an ESP Packet

Except for the fields discussed above, there are no other reserved bits in ESP. However, In the protection of time packets over IPsec scenario, the time packet is encrypted in Payload Data, the receiver could not identify whether it is the time packet or not.

5.2. Flexible ESP format

This documents proposes to define a new flexible ESP format. The new extended ESP format not only contains the fields described in [RFC4303], but also has additional authentication information. The additional authentication information is comprised of ESP special usage flags(one octet zeros),extended data type, extended data length, and Authentication Payload (Figure 3). In the extended ESP additional authentication information, it includes a data type to identify the time packets, and could also identify whether the time packet is the event message or not by additional time-packet information in Authentication Payload. In addition, the authentication of data integrity for the whole extended Data is provided. The figure of the proposed flexible ESP format is as following:

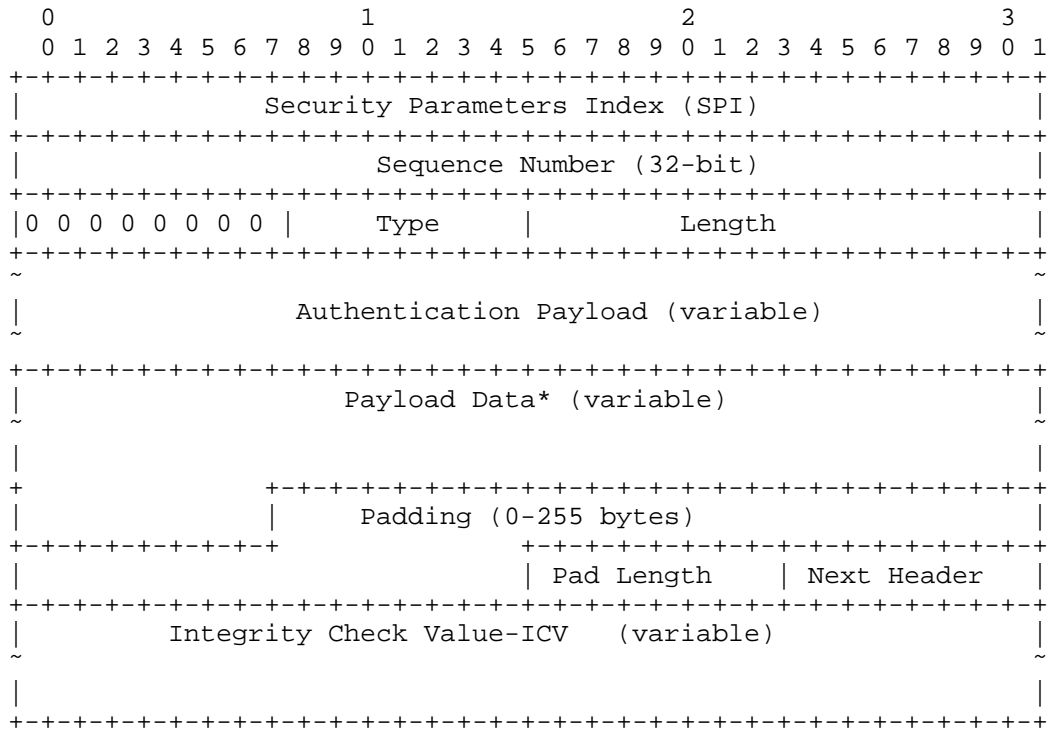


Figure 3. New defined ESP format with 32-bit Sequence Number

o Security Parameters Index(32-bit)-Defined in [RFC4303].

- o Sequence Number (32-bit or the extended 64-bit)-Defined in [RFC4303].
- o One octet zero bits - The inspection bits, used to distinguish from the existing ESP.
- o ED-Type(Extended Data Type (8-bit))- The message type flag in extended additional authentication information which indicates the message type in encrypted Payload Data.
- o ED-Length(Extended Data Length (16-bit))- The length of extended additional authentication data contains the whole extended additional authentication information.
- o Authentication Payload- It contains additional message information, and also contains the information of integrity for extended data, such as, integrity algorithm, and the extended data integrity check value. it is an optional part to provide more information of encrypted messages in Payload Data and also provide authentication of data integrity for extended data, which includes One octet zero bits, Extended Data Type, Extended Data Length and Authentication Payload data.
- o Other fields- Defined in [RFC4303].

In Femtocell scenario, as the link between Security Gateway and clock server is normally security path, the message transmitted between them are in plain text. When Security Gateway receives the message, it identifies the time packet at first, then put appropriate value to Data type field to identify the message type in Payload Data, after that, it could put more packet information into Authentication Payload, such as UDP port number or timestamps, then Extended Data Length, Algorithm ID, Extended Data integrity Check value (Figure 4), could also be filled consequently. following figure illustrates on how to use this new flexible ESP format to identify time packet.

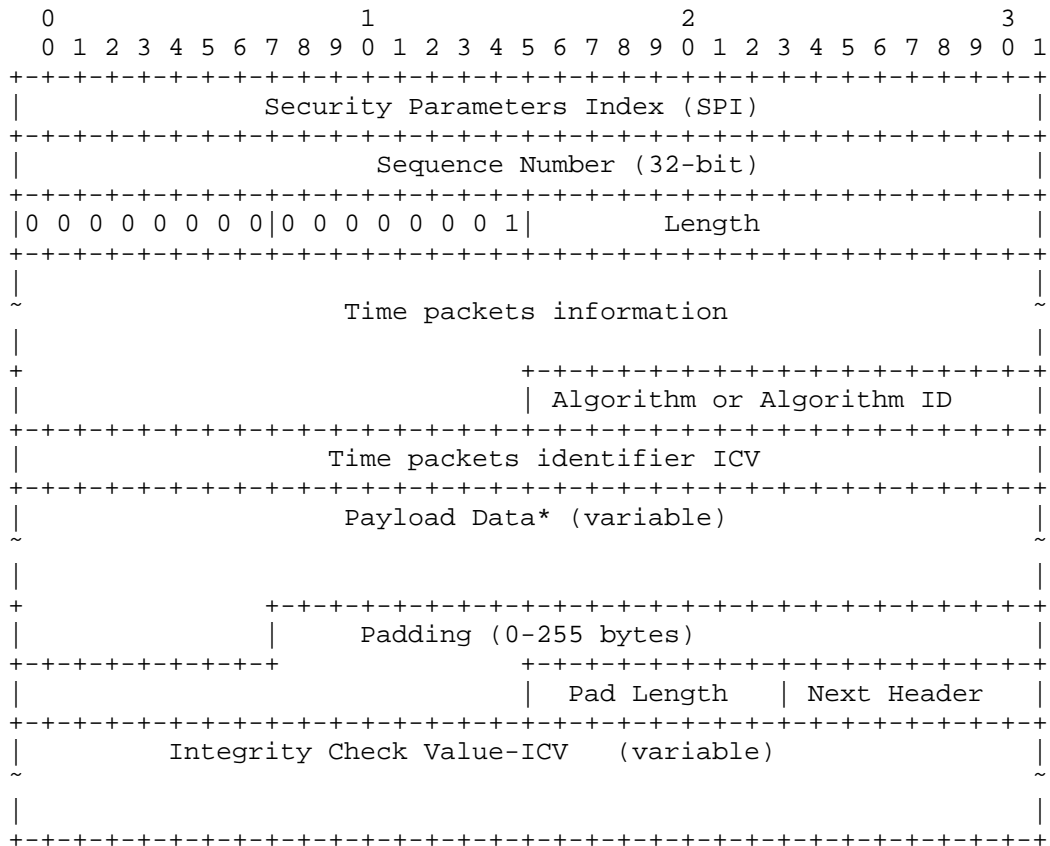


Figure 4. New defined ESP format with 32-bit Sequence Number for time-packet

- o Extended Data Type (8-bit) - The value 0x1 here indicates that the extended context is time packet.
- o Extended Data Length (16-bit)- The length of whole extended additional authentication data
- o Time packets information(variable)- the addintional message information, such as UDP port number or timestamps. It is a part of Authentication payload.
- o Algorithm or Algorithm ID- It indicates which algorithm could be used to generate the extended data ICV. It is a part of Authentication payload.The integrity algorithm negotiated during IKEv2 could be used, also Algorithm ID field in the extended additional authentication data could be marked to indicate the integrity algorithm, such as HMAC- SHA1, HMAC-256, or others. It

is a part of Authentication payload.

- o Extended Data integrity Check value (variable) - Time packets identifier integrity Check value. It is a part of Authentication payload.

Time packets information, Algorithm or Algorithm ID and Extended Data integrity Check value form the Authentication Payload, it is an optional field and used to guarantee the integrity of transmission. As the integrity protection is only for the Extended Data but not for the whole ESP packet, the time delay of calculation can be decreased. In addition, if the integrity protection is not necessary, this part of security validation could be ignored.

6. Example

In this section, the procedure to identify time packet in Security Gateway scenario is depicted.

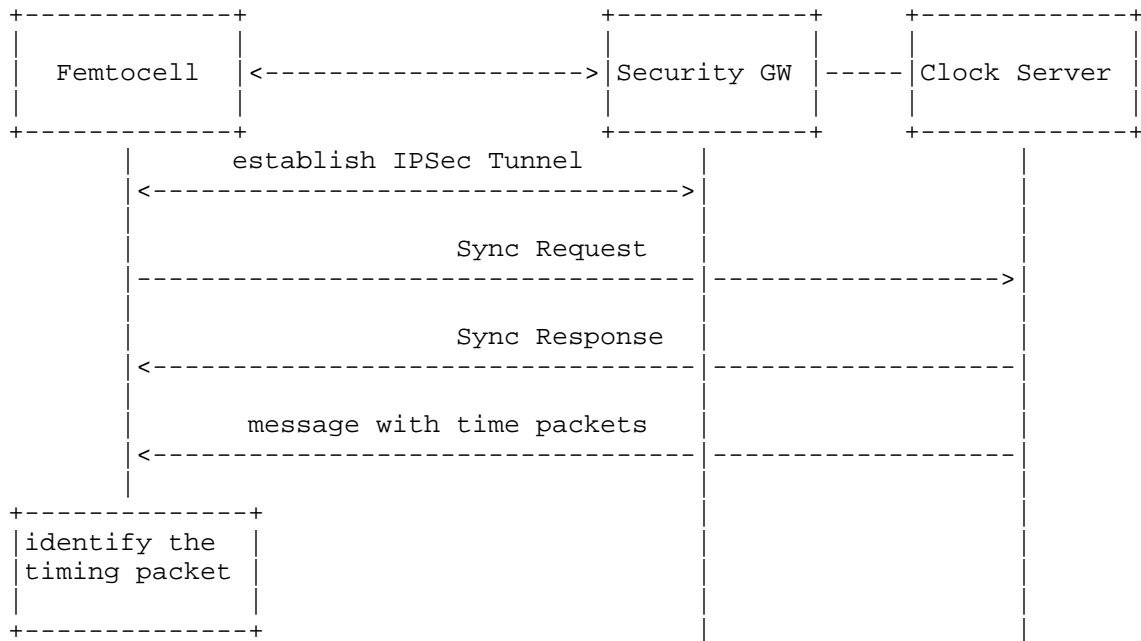


Figure 5. example procedure

In the Security Gateway scenario, The IPsec with tunnel mode is established between Femtocell and Security Gateway. After Femtocell



and Clock server exchange the Sync Request and Sync Response, the clock server will send the time packets to Femtocell to implement frequency synchronization with the protection of IPsec tunnel. When Femtocell receives the message, it can identify whether it is time packet, and can also identify whether the time packet is the event message by the time packet information in the unencrypted field as defined in the new ESP format. If the message is time packet and identifies that it is the event message, Femtocell will do special process for the event message, such as recording the message receiving time. On the server side, When Security Gateway receives the message, it identifies the time packet at first, then put appropriate value to Data type field to identify the message type in Payload Data, after that, it could put more packet information into Authentication Payload, such as UDP port number or timestamps, then Extended Data Length, Algorithm ID, Extended Data integrity Check value, could also be filled consequently.

#### 7. IPv4/v6 consideration for IPsec based synchronization

IPsec is a security mechanism used both for IPv4 and IPv6, and ESP-based solution has no impact on the IPv4 header and makes the transition/migration from IPv4 to IPv6 seamless.

#### 8. Security Considerations

This protocol variation inherits all the security properties of regular ESP as described in [RFC4303].

This document defines a new flexible ESP format, which contains Extended Data Type Extended Data Length and additional authentication payload. The authentication of data integrity for the extended data is provided, and the data type is carried in the unencryption part. Therefore the receiver could identify whether the receiving messages are time packets or not.

#### 9. IANA Considerations

There have been no IANA considerations so far in this document.

#### 10. Acknowledgments

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