A Multiplexing Extension for WebSockets
draft-tamplin-hybi-google-mux-03

Abstract

The WebSocket Protocol [RFC6455] requires a new transport connection for every WebSocket connection. This presents a scalability problem when many clients connect to the same server, and is made worse by having multiple clients running in different tabs of the same user agent. This extension provides a way for separate logical WebSocket connections to share an underlying transport connection.

Please send feedback to the hybi@ietf.org mailing list.

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Table of Contents

1. Overview ..................................................... 3
   1.1. Physical and Logical Channel ............................ 3
2. Conformance Requirements ...................................... 4
3. Interaction with other Extensions / Framing Mechanisms .... 5
   3.1. Choosing the point to apply an extension ............. 6
4. Logical Channels ............................................... 7
5. Flow Control .................................................. 8
6. Framing ....................................................... 9
7. Multiplex Control Frames ..................................... 11
   7.1. Multiplex Control Opcodes .............................. 12
8. Examples ........................................................ 16
9. Client Behavior ............................................... 17
10. Buffering .................................................... 18
11. Fairness ..................................................... 19
12. Proxies ...................................................... 20
13. Nesting ....................................................... 21
14. Timeout ...................................................... 22
15. Close the Logical Channel ................................... 23
16. Fail the Logical Channel .................................... 24
17. Fail the Physical Channel .................................. 25
18. Handling Operations On Logical Channel ................... 26
19. Security Considerations ..................................... 27
20. IANA Considerations ......................................... 28
21. Normative References ....................................... 29
Authors’ Addresses ............................................... 30
1. Overview

This document describes a multiplexing extension for the WebSocket Protocol. A client that supports this extension will advertise support for it in the client's opening handshake using the "Sec-WebSocket-Extensions" header. If the server supports this extension and supports parameters compatible with the client’s request, it accepts the use of this extension by the "Sec-WebSocket-Extensions" header in the server’s opening handshake.

1.1. Physical and Logical Channel

Under this extension, one transport connection is shared by multiple application-level instances. The WebSocket connection which lies directly on the transport connection and negotiated this multiplexing extension is called "physical channel". Virtually established WebSocket connections for each WebSocket appplication instances are called "logical channels".

Data for different logical channels are distinguished by the channel ID allocated in the "Extension data" portion of each frame.
2. Conformance Requirements

All diagrams, examples, and notes in this specification are non-normative, as are all sections explicitly marked non-normative. Everything else in this specification is normative.

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC2119. [RFC2119]

Requirements phrased in the imperative as part of algorithms (such as "strip any leading space characters" or "return false and abort these steps") are to be interpreted with the meaning of the key word ("must", "should", "may", etc) used in introducing the algorithm.

Conformance requirements phrased as algorithms or specific steps MAY be implemented in any manner, so long as the end result is equivalent. (In particular, the algorithms defined in this specification are intended to be easy to follow, and not intended to be performant.)
3. Interaction with other Extensions / Framing Mechanisms

If WebSocket payload data is masked by a per-frame key, such masking is applied to frames for each logical channel separately.

If any extension (e.g. compression) is placed before this extension in the "Sec-WebSocket-Extensions" header of the physical channel, that extension is applied to logical channels unless otherwise noted in the extension’s spec.

If such an extension define fields in the "Extension data", they come after this multiplexing extension’s field.

If any extension is placed after this extension in the "Sec-WebSocket-Extensions" header of the physical channel, that extension is applied to frames after multiplexing on the sender side, and before demultiplexing on the receiver side unless otherwise noted in the extension’s spec.

If such an extension define fields in the "Extension data", they come before this multiplexing extension’s field.

A client MAY request such an extension for both the physical channel and the logical channels by placing extension entries before and after this multiplexing extension. In this case, the server SHOULD reject at least either of them if it’s useless to apply the same extension twice.

For example, if we have a compression extension called foo-compress, the client sends

```
Sec-WebSocket-Extensions: foo-compress, mux, foo-compress
```

in the client’s opening handshake of the physical channel to request use of the compression for both physical and logical channels. Then, the server would send back

```
Sec-WebSocket-Extensions: mux, foo-compress
```

to apply compression after multiplexing, or

```
Sec-WebSocket-Extensions: foo-compress, mux
```

to apply compression to logical channels.
3.1. Choosing the point to apply an extension

Where to apply a compression extension makes difference to resource consumption and flexibility. Compression algorithms often use some memory to keep its context. Some of compression extensions may keep using the same context for all the frames on the same connection.

If such an extension is applied to the physical channel, intermediaries that want to demultiplex or multiplex the connection need to decompress (before demultiplexing) and recompress (before multiplexing again) all the frames.

If such an extension is applied to each logical channel, we can control to which channel we apply the compression, so we can avoid applying compression to channels transferring incompressible data. Intermediaries that want to demultiplex can forward Application data field leaving it untouched. However, compressing each logical channel is expensive in terms of memory consumption.
4. Logical Channels

The multiplexing extension maintains separate logical channels, each of which is fully the logical equivalent of an independent WebSocket connection, including separate handshake headers. If the multiplexing extension is successfully negotiated, the headers on the opening handshake of the physical channel are automatically taken to mean one for the logical channel 1, which is implicitly opened by completing the handshake. New channels are added by the client issuing the AddChannel request (note that only the client may initiate new WebSocket connections), including any handshake headers which do not have the same value as the client’s opening handshake of the physical channel. The server’s AddChannel response likewise includes any handshake headers which are different from the server’s opening handshake of the physical channel (the details of this are TBD, but a simple suggestion for a delta encoding is given below). Channel 0 (control channel) is reserved for multiplex control frames and does not contain payload data from any logical channel. In interpreting "Sec-WebSocket-Extensions" header for a logical channel, the entry for this multiplexing extension is ignored but is used to adjust parameters for the logical channel. A client which attempts to add a channel to an existing connection that is not accepted by the server SHOULD attempt to open a new WebSocket connection.

If any inconsistency is found between the "Sec-WebSocket-Extensions" header for the physical channel and one for a logical channel (after decoding header compression), the server MUST reject the AddChannel request.

Once the multiplexing extension is negotiated on a connection, all frames must be prefixed with a channel ID number in the "Extension data". Control frames with a channel ID 0 refer to the physical channel, other control frames MUST be delivered on the logical channel in order with data frames for that logical channel. Control frames SHOULD be sent only on channel 0 where possible, though control frames for other extensions in particular may need to apply to individual logical channels.

A receiver MUST _Fail the Physical Channel_ if any of these rules are violated by the sender.
5. Flow Control

Each logical channel, including the implicitly created channel 1, is initially given a quota of bytes that may be transmitted in each direction without acknowledgement. It is illegal to send more bytes than the remaining send quota, and the receiver MUST _Fail the Logical Channel_ for any sender that does so. This send quota is replenished via control frames as the receiver processes the data.

The initial send quota is specified with the "quota" extension parameter, and defaults to 64k (TBD) if it is not specified. The client and server each may specify a "quota" parameter and these are unrelated -- each specifies how many bytes the other side may send without acknowledgement. The quota values in the opening handshakes of the physical channel apply to the implicitly opened channel 1.
6. Framing

If the extension is successfully negotiated during the opening handshake, all frames have a channel ID in the "Extension data". The channel ID is encoded as a variable number of bytes, as follows:

```
0 1 2 3 4 5 6 7
+-------------+
|0|Channel ID(7)|
+-------------+
```

```
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5
+---------------------------+
|1|0| Channel ID (14) |
+---------------------------+
```

```
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3
+---------------------------+
|1|1|0| Channel ID (21) |
+---------------------------+
```

```
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
+---------------------------+
|1|1|1| Channel ID (29) |
+---------------------------+
```

The base spec requires that a sequence of frames on the wire be a sequence of valid fragments (or one of valid unfragmented frames). The multiplexing extension relaxes this requirement to be for just frames of one logical channel, and that frames of other logical channels may be interleaved arbitrarily.

All frames with a non-zero channel ID must be delivered to the specified logical channel in the order they are received, though fragmentation may be changed if appropriate. Control frames with a non-zero channel ID may also trigger additional processing by the multiplexing extension.

Control frames with a channel ID of 0 refer to the physical connection, and may also trigger additional processing - for example, a close frame on the physical channel will close all logical channels as well (details TBD).

If a frame doesn’t contain valid channel ID, _Fail the Physical Channel_. The cases where it’s considered that the channel ID is
invalid are:

- The "Payload data" portion doesn’t contain a complete channel ID.
- No channel has been opened for the channel ID.
- The channel has been closed and not reopened.
7. Multiplex Control Frames

Binary frames with a channel ID of 0 are multiplex control frames. Unless another negotiated extension defines a meaning for them, any data frames on channel 0 with an opcode other than "binary frame" MUST _Fail the Physical Channel_ "Payload data" of a multiplex control frames consists of a zero or more multiplex control blocks, each defined as follows:

```
0 1 2 3 4 5 6 7
+----------------+
| Objective      |
+---------------+
| : channel ID   :
+---------------+
| (8-32)        |
+----------------+
| Opc | Opcdata     |
+----------------+
| Additional     |
+---------------+
| : data         :
+---------------+
```

Objective channel ID

The channel ID of the logical channel objective to this operation. Encoding is the same as that in the extension data (designated as control channel)

opc

A multiplex control opcode as defined in Section 7.1.

opcd
data

Data interpreted according to that opcode

Additional data

Zero or more bytes defined by that opcode

If any incomplete multiplex control block is found, _Fail the Physical Channel_.

7.1. Multiplex Control Opcodes

0 - AddChannel request (only from client)

Create a new logical channel, exactly as if a new connection were received on a separate transport connection, except for the encoding of the headers. opcdata is interpreted as follows:

```
  3 4 5 6 7
 +-----+
 |R|Enc|Len|
 +-----+
```

R is reserved for future use.

Len is the number of bytes used to represent the length of following handshake data minus 1.

Enc is an encoding scheme type:

0 - uncompressed

The handshake data that follow are uncompressed, and constitute the complete set of a Request-Line and headers that would have been sent on a WebSocket opening handshake.

1 - delta-encoded

The handshake data that follow are delta-encoded, where any header that is not given is assumed to have the same value as that given on the client’s opening handshake of the physical connection. The only exceptions are the Request-Line and the "Sec-WebSocket-Extensions" header. The Request-Line MUST be sent even if it’s the same as one in the client’s opening handshake for the physical channel. If the "Sec-WebSocket-Extensions" header is not given, its value is assumed to be the extension entry for this multiplexing extension and ones following that in the client’s opening handshake of the physical channel. A header with an empty value means that header is not inherited from the initial connection. (TBD: this means that valueless headers cannot be encoded with this scheme).

2-3 - reserved

Reserved for future use.
The following n bytes, where n is the value of len inside opcdata plus 1, are an 8-32 bit length of the client’s opening handshake for the new logical channel that follows, in network byte order. It’s encoded as defined by the enc value in opcdata.

The initial quota for the new logical channel is 0, so the client may not send any data for this connection until the AddChannel response is received.

The server always responds with an AddChannel response, described below.

1 - AddChannel response (only from server)

opcdata is defined as follows:

```
  3 4 5 6 7
+-----+---+
|F|Enc|Len|
+-----+---+
```

F is true if this response indicates a rejection of AddChannel request.

Len is the number of bytes used to represent the length of following handshake data minus 1.

Enc is an encoding scheme type defined as in the AddChannel request (but replacing Request-Line with Response-Line).

If F is set, then the server has rejected the AddChannel request and this SHOULD be treated exactly the same as if a separate connection was attempted and the opening handshake failed. Enc is ignored in this case, and the following n bytes, where n is the value of len inside opcdata plus 1, are an 8-32 bit length of the server’s opening handshake for this logical channel that follows, in network byte order. It SHOULD be treated as the response to an HTTP Upgrade request for the request made by the AddChannel request, For example:

```
HTTP/1.1 404 Not found

404 message body...
```

If F is not set, then the server has accepted the AddChannel request.

The following n bytes, where n is the value of len inside opcdata
plus 1, are an 8-32 bit length of the server’s opening handshake for this logical channel that follows, in network byte order. It’s encoded according to \texttt{enc} as defined in the AddChannel request, and the complete set of a Response-Line and headers after decoding is treated exactly as if it was received in response to a client’s opening handshake on a separate connection. If the server’s opening handshake is validated, the client MUST take this as _The WebSocket Connection is Established_.

2 - FlowControl

\texttt{opcdata} is defined as follows:

\begin{verbatim}
  3 4 5 6 7
  +-----+-
  \ | RSV | Len |
  +-----+-
\end{verbatim}

RSV is reserved for future use.

Len is the number of bytes used to represent the number of bytes to be added to the quota minus 1.

The following \texttt{n} bytes, treated as an unsigned integer in network byte order, is added to the quota of the number of bytes the receiver can have outstanding towards the sender of the FlowControl message. (TBD: is it worth having some non-linear encoding to reduce the average bits required to represent these values?)

3 - DropChannel

DropChannel is used to close a logical channel for both error cases and normal cases.

\begin{verbatim}
  3 4 5 6 7
  +-----+-
  \ | R | RSV | Len |
  +-----+-
\end{verbatim}

If \texttt{R} is set, it means that this DropChannel control block was sent due to _Fail the Logical Channel_. If \texttt{R} is unset, it means that this DropChannel control block was sent due to _Close the Logical Channel_.

RSV is reserved for future use.
Len is the number of bytes used to represent the length of following reason data minus 1.

The following n bytes, where n is the value of len inside opcdata plus 1, are an 8-32 bit length of the DropChannel reason string in network byte order.

When an endpoint received DropChannel, the endpoint MUST remove the logical channel and the application instance that used the logical channel MUST treat this as closure of underlying transport.

4-7 - reserved

Reserved for future use (TBD: do we need some support for quiescence?)
8. Examples

_This section is non-normative._

The examples below assume the handshake has already completed and the x-google-mux extension was negotiated.

```
01 06 01 "Hello" 81 04 02 "bye" 80 07 01 " world"
```

This is a fragmented text message of "Hello world" on channel 1 interleaved with a text message of "bye" on channel 2. Note that the sequence of opcodes/FIN bits cannot be understood without considering the channel ID of each frame.
9. Client Behavior

When a client is asked to _Establish a WebSocket Connection_ by some WebSocket application instance, it MAY choose to reuse an existing WebSocket connection if all of the following are true:

- the multiplexing extension was successfully negotiated on that connection
- the scheme portions of the URIs match exactly
- the host portions of the URIs either match exactly or resolve to the same IP address (TBD: consider DNS rebind attacks)
- the port portions of the URIs (either explicit or implied by the scheme) match exactly
- the connection has an available logical channel ID

If the client chooses to reuse an existing multiplexed connection, it sends an AddChannel request as described above. If the AddChannel request is accepted, WebSocket frames may be sent over that channel as normal. If the server rejects the AddChannel, the client SHOULD attempt to open a new physical WebSocket connection (for example, in a shared hosting environment a server may not be prepared to multiplex connections from different customers despite having a single IP address for them).
10. Buffering

There will be lots of small frames sent in this protocol (particularly replenishing send quotas), so a sender SHOULD attempt to aggregate multiplex control blocks into larger WebSocket frames. For data frames, a sender also SHOULD attempt to aggregate fragments into one packet of the underlying transport. However, care must be taken to avoid introducing excessive latency - the exact heuristics for delaying in order to aggregate blocks is TBD.
11.  Fairness

A multiplexing implementation MUST ensure reasonable fairness among the logical channels. This is accomplished in several ways:

Receiver side

- The receiver MAY limit the other peer’s send quota of a logical channel by not replenishing the send quota to make sure that any logical channel cannot dominate its buffer space on the sender.

- Send quota for one logical channel SHOULD be determined considering the processing capacity (buffer size, processing power, throughput, etc.) of that logical channel. For example, when a logical channel with excess load cannot drain data from the connection smoothly, the other logical channels get stuck even when they have room of processing capacity. Unless there’s special need to give such a big quota for the channel, such condition just makes overall performance low.

Sender side

- The sender MUST use a fair mechanism for selecting which logical channel’s data to send in the next WebSocket frame. Simple implementations may choose a round-robin scheduler, while more advanced implementations may adjust priority based on the amount or frequency of data sent by each logical channel.

- The sender MUST fragment a message into smaller frames when it’s too big so that that logical channel will occupy the connection and the other logical channels get stuck for long time.

- Logical channel frames that are sent SHOULD be limited in size (such as by refragmenting) when there is contention for the physical channel to minimize head-of-line blocking.
12. Proxies

Proxies which do not multiplex/demultiplex are not affected by the presence of this extension -- they simply process WebSocket frames as usual. Proxies which filter or monitor WebSocket traffic will need to understand the multiplexing extension in order to extract the data from logical connections or to terminate individual logical connections when policy is violated. Proxies which actively multiplex connections or demultiplex them (for example, a mobile network might have a proxy which aggregates WebSocket connections at a single cell to conserve bandwidth to the main gateway) will require additional configuration (perhaps including the client) that is outside the scope of this document.
13. Nesting

TBD: Should we allow nesting of multiplexed channels, or should we require that an intermediary multiplexing channels flatten it? The advantage of nesting is it is conceptually cleaner and less work for an intermediary, while the disadvantage is that flow control messages will get amplified by nesting and the ultimate server’s job is a bit more complicated to keep a tree of channel mappings.
14. Timeout

When all the logical channels are closed, each endpoint MAY "Start the WebSocket Closing Handshake" on the physical connection. Such "Start the WebSocket Closing Handshake" operation SHOULD be delayed assuming the physical channel may be reused after some idle period.
15. Close the Logical Channel

To _Close the Logical Channel_, an endpoint MUST send a DropChannel multiplex control block with R bit unset. The endpoint MAY provide the reason of closure in the DropChannel block.
16. Fail the Logical Channel

To _Fail the Logical Channel_, an endpoint MUST send a DropChannel multiplex control block with R bit set. The endpoint MAY provide the reason of failure in the DropChannel block.
17. Fail the Physical Channel

To _Fail the Physical Channel_, an endpoint MUST send a DropChannel multiplex control block with objective channel ID of 0, and then _Fail the WebSocket Connection_ on the physical channel with status code of 1002 (TBD).
18. Handling Operations On Logical Channel

When an endpoint is asked to perform any operation defined in the WebSocket Protocol except for _Close the WebSocket Connection_ by some application instance, it MUST perform it on the corresponding logical channel.

Any event on a logical channel except for _The WebSocket Connection is Closed_, MUST be taken as one for the corresponding application instance.

When an endpoint is asked to do _Close the WebSocket Connection_ by some application instance, it MUST perform _Close the Logical Channel_ on the corresponding logical channel.

When a DropChannel is received and the logical channel hasn’t yet received DropChannel before that, it MUST be taken as _The WebSocket Connection is Closed_ event for the corresponding application instance.
19. Security Considerations

To protect a server from denial-of-service attack, implementation SHOULD have a way to limit the number of concurrent logical channels.

TBD
20. IANA Considerations

This specification is registering a value of the Sec-WebSocket-Extension header field in accordance with Section 11.4 of the WebSocket protocol [RFC6455] as follows:

Extension Identifier

mux

Extension Common Name

Multiplexing Extension for WebSockets

Extension Definition

This document [draft-tamplin-hybi-google-mux] defines the mux extension.

Known Incompatible Extensions

None
21. Normative References


Authors’ Addresses

John A. Tamplin
Google, Inc.

Email: jat@google.com

Takeshi Yoshino
Google, Inc.

Email: tyoshino@google.com
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Draft Authors:
Martin Thomson<martin.thomson@andrew.com>
Salvatore Loreto<Salvatore.Loreto@ericsson.com>
Greg Wilkins<gregw@intalio.com>
Abstract

This specification defines a general per-frame compression scheme for the WebSocket Protocol and one specific compression extension using DEFLATE. This scheme compresses the "Application data" part of WebSocket data frames using specified compression algorithm.

Please send feedback to the hybi@ietf.org mailing list.

Status of this Memo

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Table of Contents

1. Introduction 3
2. Conformance Requirements 4
3. General Per-frame Compression Scheme 5
  3.1. Sending 5
  3.2. Receiving 5
4. Per-frame DEFLATE Extension 6
  4.1. Extension Negotiation 6
  4.2. Application Data Transformation 6
    4.2.1. Compression 7
    4.2.2. Decompression 7
    4.2.3. Examples 8
  4.3. Implementation Note 8
5. Security Considerations 9
6. IANA Considerations 10
  6.1. Registration of the "deflate-frame" WebSocket Extension Name 10
  6.2. Registration of the "Per-frame compressed" WebSocket Framing Header Bit 10
7. Acknowledgements 11
8. References 12
  8.1. Normative References 12
  8.2. Informative References 12
Author’s Address 13
1. Introduction

_This section is non-normative._

As well as other protocols, octets exchanged over the WebSocket Protocol [RFC6455] can benefit from compression technology. This specification introduces a scheme to apply compression algorithms to the WebSocket Protocol, and then adds DEFLATE [RFC1951] based compression functionality to the WebSocket Protocol using its extension framework.

The per-frame compression scheme applies the specified compression algorithm to the octets in the "Application data" part of data frames. It specifies the use of the RSV1 bit of the WebSocket framing to indicate whether any compression is applied to the frame or not, so that we can choose to skip compression for frames with incompressible contents. By specifying extension negotiation and how to transform octets in "Application data", we can define per-frame compression extensions for various compression algorithms based on this scheme.

The specific extension we introduce in this specification by applying DEFLATE to the scheme is called "Per-frame DEFLATE extension". To align the end of compressed data to octet boundary, the extension uses the algorithm described in the Section 2.1 of the PPP Deflate Protocol [RFC1979]. Endpoints can take over the LZ77 sliding window [LZ77] used to build previous frames to get better compression ratio.

The simplest "Sec-WebSocket-Extensions" header in the client’s opening handshake to request per-frame DEFLATE extension is the following:

```
Sec-WebSocket-Extensions: deflate-frame
```

The simplest header from the server to accept this extension is the same.
2. Conformance Requirements

Everything in this specification except for sections explicitly marked non-normative is normative.

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].
3. General Per-frame Compression Scheme

This section describes a general scheme to apply a compression algorithm to the contents of WebSocket frames. Any other compression extension specifications MAY apply their compression algorithm to this scheme to define extensions.

This scheme allocates one bit field called "Per-frame compressed" at the RSV1 bit. This bit indicates whether any kind of per-frame compression is applied to the frame or not.

This scheme operates only on data frames, and only on the "Application data" therein (it does not affect the "Extension data" portion of the "Payload data").

3.1. Sending

To send a frame with the compression applied, an endpoint MUST use the following algorithm.

1. Apply the compression to the "Application data" portion of the frame.

2. Build a frame by putting the resulting octets in the "Application data" portion instead of the original octets. The payload length field of the frame MUST be the sum of the size of the "Extension data" portion and one of these resulting octets. "Per-frame compressed" bit MUST be set to 1.

To send a frame with the compression not applied, an endpoint MUST set "Per-frame compressed" bit of the frame to 0.

3.2. Receiving

To receive a frame with "Per-frame compressed" bit set to 1, an endpoint MUST decompress the octets in the "Application data" portion based on the compression algorithm.

An endpoint MUST receive a frame with "Per-frame compressed" bit set to 0 as-is without any compression processing.
4. Per-frame DEFLATE Extension

4.1. Extension Negotiation

The registered extension token for this extension is "deflate-frame".

To request use of per-frame DEFLATE extension, a client MUST include the "deflate-frame" extension token in the "Sec-WebSocket-Extensions" header in its opening handshake.

To accept use of per-frame DEFLATE extension requested by the client, a server MUST include the "deflate-frame" extension token in the "Sec-WebSocket-Extensions" header in its opening handshake.

An endpoint MAY attach one or more extension parameters as defined below to the extension token.

Maximum LZ77 sliding window size
An endpoint MAY attach "max_window_bits" extension parameter to limit the LZ77 sliding window size that the other peer uses to build frames. This parameter MUST have an integer value in the range between 8 to 15 indicating the base-2 logarithm of the LZ77 sliding window size. An endpoint that received this parameter MUST NOT use LZ77 sliding window size greater than this value to build frames.

Disallow compression context takeover
An endpoint MAY attach "no_context_takeover" extension parameter to disallow the other peer to take over the LZ77 sliding window used to build previous frames. This parameter has no value. An endpoint that received this parameter MUST use an empty LZ77 sliding window to build every frame.

A server MUST ignore any unknown extension parameter attached to "deflate-frame" extension token in the client’s opening handshake.

A client MUST _Fail the WebSocket Connection_ if any unknown extension parameter is attached to "deflate-frame" extension token in the server’s opening handshake.

Once per-frame DEFLATE extension is accepted, both endpoints MUST use the algorithm described in Section 4.2 to exchange frames.

4.2. Application Data Transformation

This extension transforms the "Application data" portion by using the scheme described in Section 3 with DEFLATE as follows.
4.2.1. Compression

An endpoint MUST use the following algorithm to compress the "Application data" portion.

1. Apply DEFLATE [RFC1951] to all the octets. Multiple blocks MAY be used. Any type of block MAY be used. Both block with "BFINAL" set to 0 and 1 MAY be used.

2. If the resulting data does not end with an empty block with no compression ("BTYPE" set to 0), append an empty block with no compression to the tail.

3. Remove 4 octets (that are 0x00 0x00 0xff 0xff) from the tail.

An endpoint MUST NOT use LZ77 sliding window size greater than 32,768 to build frames.

If an endpoint received the "max_window_bits" extension parameter on opening handshake, it MUST NOT use LZ77 sliding window size greater than the "max_window_bits"-th power of 2 to build frames.

Unless it’s prohibited by the other peer by the "no_context_takeover" extension parameter on opening handshake, an endpoint MAY take over the LZ77 sliding window used to build the last frame with DEFLATE applied.

4.2.2. Decompression

An endpoint MUST use the following algorithm to decompress the "Application data" portion.

1. Append 4 octets of 0x00 0x00 0xff 0xff to the tail.

2. Decompress the resulting octets using DEFLATE.

Unless an endpoint sent the "max_window_bits" extension parameter on opening handshake, the endpoint MUST use 32,768 byte LZ77 sliding window to decode frames.

If an endpoint sent the "max_window_bits" extension parameter on opening handshake, it MAY reduce LZ77 sliding window size down to the "max_window_bits"-th power of 2 to decode frames.

Unless the endpoint sent the "no_context_takeover" extension parameter on opening handshake, an endpoint MUST take over the LZ77 sliding window used to decode the last frame with DEFLATE applied.
4.2.3. Examples

_This section is non-normative._

These are examples of resulting data after applying the algorithm above.

- "Hello" in one compressed block
  * 0xf2 0x48 0xcd 0xc9 0xc9 0x07 0x00
  "Hello" in one compressed block in the next frame
  * 0xf2 0x00 0x11 0x00 0x00
- "Hello" in one block with no compression
  * 0x00 0x05 0x00 0xfa 0xff 0x48 0x65 0x6c 0x6f 0x00
- "Hello" in one block with "BFINAL" set to 1
  * 0xf3 0x48 0xcd 0xc9 0xc9 0x07 0x00 0x00
- "He" and "llo" in separate blocks
  * 0xf2 0x48 0x05 0x00 0x00 0xff 0x0f 0xcb 0xc9 0xc9 0x07 0x00

4.3. Implementation Note

_This section is non-normative._

On common software development platforms, the operation of aligning compressed data to octet boundary using an empty block with no compression is available as library. For example, Zlib [Zlib] does this when "Z_SYNC_FLUSH" is passed to deflate function.

To get sufficient compression ratio, LZ77 sliding window size of 1,024 or more is recommended.
5. Security Considerations

There’s no security concern for now.
6. IANA Considerations

6.1. Registration of the "deflate-frame" WebSocket Extension Name

This section describes a WebSocket extension name registration in the WebSocket Extension Name Registry. [RFC6455].

Extension Identifier
deflate-frame

Extension Common Name
WebSocket Per-frame DEFLATE

Extension Definition
Section 4.1 and Section 4.2 of this document.

Known Incompatible Extensions
None

The "deflate-frame" token is used in the "Sec-WebSocket-Extensions" header in the WebSocket opening handshake to negotiate use of per-frame DEFLATE compression extension.

6.2. Registration of the "Per-frame compressed" WebSocket Framing Header Bit

This section describes a WebSocket framing header bit registration in the WebSocket Framing Header Bits Registry. [RFC6455]

Header Bit
RSV1

Common Name
Per-frame compressed

Meaning
Compression is applied to the frame or not.

Reference
Section 3 of this document.

The "Per-frame compressed" framing header bit is used to indicate whether any negotiated per-frame compression extension applied compression to the "Application data" portion of the frame or not.
7. Acknowledgements

Special thanks to Patrick McManus who wrote up the initial specification of DEFLATE based compression extension for the WebSocket Protocol which I referred to write this specification.
8. References

8.1. Normative References


8.2. Informative References


Author’s Address

Takeshi Yoshino
Google, Inc.

Email: tyoshino@google.com