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Transport Layer Security (TLS) Cached Information Extension
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Abstract

Transport Layer Security (TLS) handshakes often include fairly static information, such as the server certificate and a list of trusted Certification Authorities (CAs). This information can be of considerable size, particularly if the server certificate is bundled with a complete certificate path (including all intermediary certificates up to the trust anchor public key).

This document defines an extension that omits the exchange of already available information. The TLS client informs a server of cached information, for example from a previous TLS handshake, allowing the server to omit the already available information.

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1. Introduction

Transport Layer Security (TLS) handshakes often include fairly static information, such as the server certificate and a list of trusted Certification Authorities (CAs). This information can be of considerable size, particularly if the server certificate is bundled with a complete certificate path (including all intermediary certificates up to the trust anchor public key).

Optimizing the exchange of information to a minimum helps to improve performance in environments where devices are connected to a network with characteristics like low bandwidth, high latency and high loss rate. These types of networks exist, for example, when smart objects are connected using a low power IEEE 802.15.4 radio. For more information about the challenges with smart object deployments please see [RFC6574].

This specification defines a TLS extension that allows a client and a server to exclude transmission of cached information from the TLS handshake.

A typical example exchange may therefore look as follows. First, the TLS exchange executes the usual TLS handshake. It may decide to store the certificate provided by the server for a future exchange. When the TLS client then connects to the TLS server some time in the future, without using session resumption, it then attaches the `cached_information` extension defined in this document to the client hello message to indicate that it had cached the certificate, and it provides the fingerprint of it. If the server's certificate had not changed then the TLS server does not need to send the full certificate to the client again. In case the information had changed, the certificate payload is transmitted to the client to allow the client to update it's state information.

2. Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "MUST", "MUST NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

3. Cached Information Extension

This document defines a new extension type (`cached_information(TBD)`), which is used in client hello and server hello messages. The extension type is specified as follows.

```
enum {
    cached_information(TBD), (65535)
} ExtensionType;
```

The `extension_data` field of this extension, when included in the client hello, MUST contain the `CachedInformation` structure.

```
enum {
    certificate_chain(1), trusted_cas(2) (255)
} CachedInformationType;

struct {
    CachedInformationType type;
    HashAlgorithm hash;
    opaque hash_value<1..255>;
} CachedObject;

struct {
    CachedObject cached_info<1..2^16-1>;
} CachedInformation;
```

When the `CachedInformationType` identifies a `certificate_chain`, then the `hash_value` field MUST include the hash calculated over the `certificate_list` element of the `Certificate` payload provided by the TLS server in an earlier exchange, excluding the three length bytes of the `certificate_list` vector.

When the `CachedInformationType` identifies a `trusted_cas`, then the `hash_value` MUST include a hash calculated over the `certificate_authorities` element of the `CertificateRequest` payload provided by the TLS server in an earlier exchange, excluding the two length bytes of the `certificate_authorities` vector.

The hash algorithm used to calculate hash values is conveyed in the 'hash' field of the `CachedObject` element. The list of registered hash algorithms can be found in the TLS HashAlgorithm Registry, which was created by RFC 5246 [RFC5246]. The value zero (0) for 'none' is not an allowed choice for a hash algorithm and MUST NOT be used.

This document establishes a registry for `CachedInformationType` types and additional values can be added following the policy described in Section 7.

4. Exchange Specification

Clients supporting this extension MAY include the "cached_information" extension in the (extended) client hello, which MAY contain zero or more CachedObject attributes.

Server supporting this extension MAY include the "cached_information" extension in the (extended) server hello, which MAY contain one or more CachedObject attributes. By returning the "cached_information" extension the server indicates that it supports caching of each present CachedObject that matches the specified hash value. The server MAY support other cached objects that are not present in the extension.

Note: Clients may need the ability to cache different values depending on other information in the Client Hello that modify what values the server uses, in particular the Server Name Indication [RFC6066] value.

Following a successful exchange of "cached_information" extensions, the server MAY send fingerprints of the cached information in the handshake exchange as a replacement for the exchange of the full data. Section 4.1 and Section 4.2 defines the syntax of the fingerprinted information.

The handshake protocol MUST proceed using the information as if it was provided in the handshake protocol. The Finished message MUST be calculated over the actual data exchanged in the handshake protocol. That is, the Finished message will be calculated over the hash values of cached information objects and not over the cached information that were omitted from transmission.

The server MUST NOT include more than one fingerprint for a single information element, i.e., at maximum only one CachedObject structure per replaced information is provided.

4.1. Fingerprint of the Certificate Chain

When an object of type 'certificate_chain' is provided in the client hello, the server MAY send a fingerprint instead of the complete certificate chain as shown below.

The original handshake message syntax is defined in RFC 5246 [RFC5246] and has the following structure:

```
opaque ASN.1Cert<1..2^24-1>;

struct {
    ASN.1Cert certificate_list<0..2^24-1>;
} Certificate;
```

By using the extension defined in this document the following information is sent:

```
struct {
    CachedObject ASN.1Cert<1..2^24-1>;
} Certificate;
```

The opaque ASN.1Cert structure is replaced with the CachedObject structure defined in this document.

Note: [I-D.ietf-tls-oob-pubkey] allows a PKIX certificate containing only the SubjectPublicKeyInfo instead of the full information typically found in a certificate. Hence, when this specification is used in combination with [I-D.ietf-tls-oob-pubkey] and the negotiated certificate type is a raw public key then the TLS server sends the hashed Certificate payload that contains a ASN.1Cert structure of the SubjectPublicKeyInfo.

4.2. Fingerprint for Trusted CAs

When a hash for an object of type 'trusted_cas' is provided in the client hello, the server MAY send a fingerprint instead of the complete certificate authorities information as shown below.

The original handshake message syntax is defined in RFC 5246 [RFC5246] and has the following structure:

```
opaque DistinguishedName<1..2^16-1>;

struct {
    ClientCertificateType certificate_types<1..2^8-1>;
    SignatureAndHashAlgorithm
        supported_signature_algorithms<2^16-1>;
    DistinguishedName certificate_authorities<0..2^16-1>;
} CertificateRequest;
```

By using the extension defined in this document the following information is sent:

```
struct {
    ClientCertificateType certificate_types<1..2^8-1>;
    SignatureAndHashAlgorithm
        supported_signature_algorithms<2^16-1>;
    CachedObject DistinguishedName<1..2^16-1>;
} CertificateRequest;
```

The opaque DistinguishedName structure is replaced with the CachedObject structure defined in this document.

5. Example

Figure 1 illustrates an example exchange using the TLS cached info extension. In the normal TLS handshake exchange shown in flow (A) the TLS server provides its certificate in the Certificate payload to the client, see step [1]. This allows the client to store the certificate for future use. After some time the TLS client again interacts with the same TLS server and makes use of the TLS cached info extension, as shown in flow (B). The TLS client indicates support for this specification via the `cached_information` extension, see [2], and indicates that it has stored the `certificate_chain` from the earlier exchange. With [3] the TLS server indicates that it also supports this specification and informs the client that it also supports caching of other objects beyond the `'certificate_chain'`, namely `'trusted_cas'` (also defined in this document), and the `'foo-bar'` extension (i.e., an imaginary extension that yet needs to be defined). With [4] the TLS server provides the fingerprint of the certificate chain as described in Section 4.1.

(A) Initial (full) Exchange

```

client_hello ->
                <- server_hello,
                   certificate, // [1]
                   server_key_exchange,
                   server_hello_done

client_key_exchange,
change_cipher_spec,
finished ->
                <- change_cipher_spec,
                   finished

Application Data <-----> Application Data

```

(B) TLS Cached Extension Usage

```

client_hello,
cached_information=(certificate_chain) -> // [2]
                <- server_hello,
                   cached_information= // [3]
                   (certificate_chain, trusted_cas, foo-bar)
                   certificate, // [4]
                   server_key_exchange,
                   server_hello_done

client_key_exchange,
change_cipher_spec,
finished ->
                <- change_cipher_spec,
                   finished

Application Data <-----> Application Data

```

Figure 1: Example Message Exchange

6. Security Considerations

This specification defines a mechanism to reference stored state using a fingerprint. The hash algorithm used in this specification is required to have reasonable random properties in order to provide reasonably unique identifiers. There is no requirement that this hash algorithm must have strong collision resistance.

Caching information in an encrypted handshake (such as a renegotiated handshake) and sending a hash of that cached information in an unencrypted handshake might introduce integrity or data disclosure issues as it enables an attacker to identify if a known object (such as a known server certificate) has been used in previous encrypted handshakes. Information object types defined in this specification, such as server certificates, are public objects and usually not sensitive in this regard, but implementers should be aware if any cached information are subject to such security concerns and in such case SHOULD NOT send a hash over encrypted data in unencrypted handshake.

7. IANA Considerations

7.1. New Entry to the TLS ExtensionType Registry

IANA is requested to add an entry to the existing TLS ExtensionType registry, defined in RFC 5246 [RFC5246], for `cached_information(TBD)` defined in this document.

7.2. New Registry for CachedInformationType

IANA is requested to establish a registry for TLS CachedInformationType values. The first entries in the registry are

- o `certificate_chain(1)`
- o `trusted_cas(2)`

The policy for adding new values to this registry, following the terminology defined in RFC 5226 [RFC5226], is as follows:

- o 0-63 (decimal): Standards Action
- o 64-223 (decimal): Specification Required
- o 224-255 (decimal): reserved for Private Use

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9. References

9.1. Normative References

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9.2. Informative References

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