### RMCAT architectural overview

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## Disclaimer

- I'll be talking about "sender" and "receiver" here, just to differentiate roles
  - Yes we're dealing with bidirectional traffic, but so is TCP, and talking about "sender" and "receiver" roles never was a problem there

### A framework

- Delay-based congestion control has many, many issues
  - Unlikely that we solve them all straight away
- Charter: "Determine if extensions to RTP/RTCP are needed for carrying congestion control feedback, using DCCP as a model. If so, provide the requirements for such extensions to the AVTCORE working group for standardization there."
  - This sounds like "define fields", but DCCP had to do much more to become a framework
  - It may also be overloaded...
- The framework involves some general design decisions that would affect all cc. mechanisms we standardize
  - Better get them right from day 1

# TCP, for example

#### Feedback

- Unreliable ACKs; can lead to misinterpretation of backward loss as forward congestion
- not a big deal because information in ACKs redundant, and lots of ACKs are sent
- to detect backward congestion (and do ACK cc.), sender must know receiver's ACK ratio

#### Reaction to ECN

 MUST be similar to reaction to loss, for compatibility with non-ECN-capable TCP flows

### Pluggable congestion control

- even without standardization, sender-side change, no need to even inform the receiver
- possible because of rather "dumb" (better: "generic") receiver

## Lessons learned from TCP

#### Feedback

 To avoid misinterpreting feedback loss as forward congestion and/ or do backwards congestion control: consider making ACKs reliable (see DCCP)

#### Reaction to ECN

- We're designing stuff from scratch here, could make all RMCAT flows ECN-capable => no need to protect non-ECN-capable RMCAT flows (?)
- Flows that don't get ECN marks from the same bottleneck: not likely (?)

### Pluggable congestion control

- Probably desirable
- Requires one side to be generic or [Randell]: both sides generic,
  might be possible to exchange either one of them

# Devil's advocate: Consider RRTCC (draft-alvestrand-rtcweb-

congestion-03), for example

### Receiver:

- Look at changes in inter-packet delay, apply some maths (Kalman filter)
- If the sender should stop increasing (or a feedback timer expires), tell it "your new rate is X"

#### Sender:

- calculate TFRC equation; rate is max (result of TFRC equation, X)
- In the absence of a feedback, increase rate (but: no feedback for a long time => timeout)

#### Can we agree that:

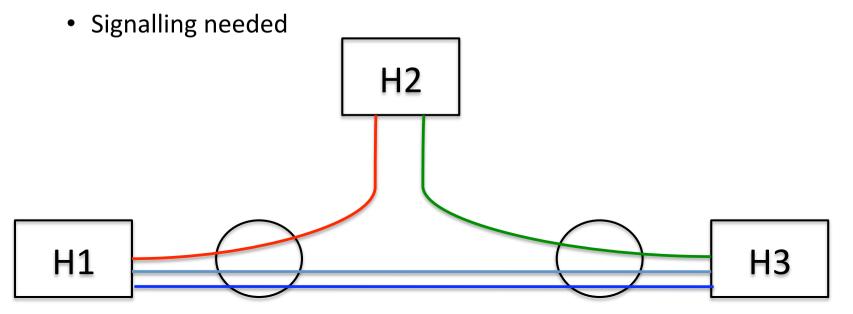
- this receiver behavior is ok for <u>all</u> possible future senders?
- this sender behavior is ok for <u>all</u> possible future receivers?
- Thinking of TCP again: the simpler one side, the more flexible the other becomes

## Sender- vs. Receiver-based

- Perhaps we should decide now which side to make simple?
- Many sides to sender- vs. receiver-based CC... Key question: always minimize feedback or not?
  - might make the control unnecessarily fragile
  - + less traffic is less traffic... also: simpler than feedback-CC, and e.g. smaller chance of collision on wireless
- Many more pro's and con's... e.g., "interactions with applications": importance of packets in send buffer could play a role for congestion control decision
  - affects signalling in case of receivers-side cc.

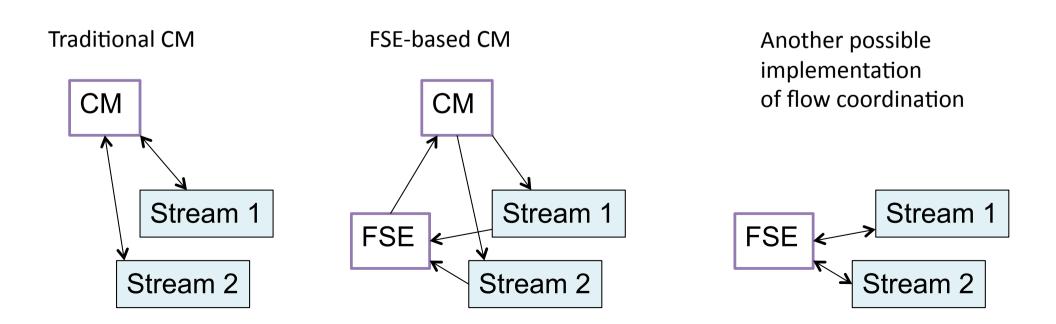
# Coupled congestion control

- Only makes sense for flows that share a bottleneck
  - Obvious for some flows in case of WebRTC (same 5-tuple = same bottleneck)
  - Less so in the general case, but there are working methods
  - Coordinate streams between multiple hosts: need to detect shared bottlenecks on both sides



# "Flow State Exchange" (FSE)

- The result of searching for minimum-necessarystandardization: only define what goes in / out, how data are maintained
  - Could reside in a single app (e.g. browser) and/or in the OS



Thank you!

Questions?