CoRE Working Group   A. Castellani
Internet-Draft   University of Padova
Intended status: Informational   S. Loreto
Expires: August 29, 2013   Ericsson
                                                               A. Rahman
                                         InterDigital Communications, LLC
                                                              T. Fossati
                                                               KoanLogic
                                                                E. Dijk
                                                        Philips Research
                                                      February 25, 2013

            Best Practices for HTTP-CoAP Mapping Implementation
                  draft-castellani-core-http-mapping-07

Abstract

   This draft provides reference information for HTTP-CoAP protocol
   translation proxy implementors, focusing primarily on the reverse
   proxy case. It details deployment options, discusses possible
   approaches for URI mapping, and provides a set of guidelines and
   considerations related to protocol translation.

Status of this Memo

   This Internet-Draft is submitted in full conformance with the
   provisions of BCP 78 and BCP 79.

   Internet-Drafts are working documents of the Internet Engineering
   Task Force (IETF). Note that other groups may also distribute
   working documents as Internet-Drafts. The list of current Internet-
   Drafts is at http://datatracker.ietf.org/drafts/current/.

   Internet-Drafts are draft documents valid for a maximum of six months
   and may be updated, replaced, or obsoleted by other documents at any
   time. It is inappropriate to use Internet-Drafts as reference
   material or to cite them other than as "work in progress."

   This Internet-Draft will expire on August 29, 2013.

Copyright Notice

Table of Contents

1. Introduction

   CoAP [I-D.ietf-core-coap] has been designed with the twofold aim to
   be an application protocol specialized for constrained environments
   and to be easily used in REST architectures such as the Web. The
   latter goal has led to define CoAP to easily interoperate with HTTP
   [RFC2616] through an intermediary proxy which performs cross-protocol
   conversion.

   Section 10 of [I-D.ietf-core-coap] describes the fundamentals of the
   CoAP-HTTP (and vice-versa) cross-protocol mapping process. However,
   implementing such a cross-protocol proxy can be complex, and many
   details regarding its internal procedures and design choices require
   further elaboration. Therefore a first goal of this document is to
   provide more detailed information to proxy designers and
   implementers, to help implement proxies that correctly inter-work
   with other CoAP and HTTP client/server implementations that adhere to
   the HTTP and CoAP specifications.

   The second goal of this informational document is to define a
   consistent set of guidelines that a HTTP-to-CoAP proxy implementation
   MAY adhere to. The main reason of adhering to such guidelines is to
   reduce variation between proxy implementations, thereby increasing
   interoperability. (As an example use case, a proxy conforming to
   these guidelines made by vendor A can be easily replaced by a proxy
   from vendor B that also conforms to the guidelines.)

   This draft is organized as follows:

   o Section 2 describes terminology to identify proxy types, mapping
      approaches and proxy deployments;

   o Section 3 discusses how URIs refer to resources independent of
      access protocols;

   o Section 5 analyzes the mapping that allows HTTP clients to contact
      CoAP servers;

   o Section 7 discusses possible security impact related to HTTP/CoAP
      cross-protocol mapping.

2. Terminology

   This document assumes readers are familiar with the terms Reverse
   Proxy as defined in [I-D.ietf-httpbis-p1-messaging] and Interception
   Proxy as defined in [RFC3040]. In addition, the following terms are
   defined:

Cross-Protocol Proxy (or Cross Proxy): is a proxy performing a cross-protocol mapping, in the context of this document a HTTP-CoAP (HC) mapping. A Cross-Protocol Proxy can behave as a Forward Proxy, Reverse Proxy or Interception Proxy. Note: In this document we focus on the Reverse Proxy mode of the Cross-Protocol Proxy.

Forward Proxy: a message forwarding agent that is selected by the client, usually via local configuration rules, to receive requests for some type(s) of absolute URI and to attempt to satisfy those requests via translation to the protocol indicated by the absolute URI. The user decides (is willing to) use the proxy as the forwarding/dereferencing agent for a predefined subset of the URI space.

Reverse Proxy: a receiving agent that acts as a layer above some other server(s) and translates the received requests to the underlying server's protocol. It behaves as an origin (HTTP) server on its connection towards the (HTTP) client and as a (CoAP) client on its connection towards the (CoAP) origin server. The (HTTP) client uses the "origin-form" [I-D.ietf-httpbis-p1-messaging] as a request-target URI.

Reverse and Forward proxies are technically very similar, with main differences being that the former appears to a client as an origin server while the latter does not, and that clients may be unaware they are communicating with a proxy.

Placement terms: a server-side (SS) proxy is placed in the same network domain as the server; conversely a client-side (CS) proxy is in the same network domain as the client. In any other case than SS or CS, the proxy is said to be External (E).

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

3. Cross-Protocol Usage of URIs

   A Uniform Resource Identifier (URI) provides a simple and extensible method for identifying a resource. It enables uniform identification of resources via a separately defined extensible set of naming schemes [RFC3986].

   URIs are formed of at least three components: scheme, authority and path. The scheme often corresponds to the protocol used to access the resource. However, as noted in Section 1.2.2 of [RFC3986] the

scheme does not imply that a particular protocol is used to access
the resource. So, we can define the same resource to be accessible
by different protocols i.e. the resource can have cross-protocol URIs
referring to it.

HTTP clients typically only support 'http' and 'https' schemes.
Therefore, they cannot directly access CoAP servers (which support
'coap' and/or 'coaps'). In this situation, communication is enabled
by a Cross-Protocol Proxy, as shown in Figure 1, supporting URI
mapping features. Such features are discussed in the following
section.

4. HTTP to CoAP URI Mapping

Assume that a HTTP client wants to access a CoAP resource and
indicates a target resource of "http://node.something.net/foobar" to
a Forward cross proxy. A possible URI mapping done by the proxy
could result in "coap://node.coap.something.net/foo".

As shown in the above example, in a cross-protocol URI the scheme,
authority and path parts of the URI may all change. The process of
providing cross-protocol URIs may be complex, since a mechanism to
statically or dynamically (e.g., discovery) map the URI is needed.

Two simple static URI mapping solutions are proposed in the following
subsections. Note that other mapping approaches are possible as
well.

4.1. Embedded Mapping

In an embedded mapping approach, the HTTP URI has embedded inside it
the authority and path part of the CoAP URI.

Example: The CoAP resource "//node.coap.something.net/foo" can be
accessed by an HTTP client by inserting in the request
"http://hc-proxy.something.net/coap/node.coap.something.net/foo".
The Cross-Protocol Proxy then maps the URI to
"coap://node.coap.something.net/foo"

4.2. Homogeneous Mapping

In a homogeneous mapping approach, only the scheme portion of the URI
needs to be mapped. The rest of the URI (i.e. authority, path, etc.)
remains unchanged.

Example: The CoAP resource "coap://node.coap.something.net/foo" can
be accessed by an HTTP client by requesting

"http://node.coap.something.net/foo". The Cross-Protocol Proxy
receiving the request is responsible to map the URI to
"coap://node.coap.something.net/foo"

Background info: The assumption in this case is that the HTTP client
would be able to successfully resolve "node.coap.something.net" using
DNS infrastructure to return the IP address of the HC proxy. Most
likely this would be through a two step DNS lookup where the first
DNS lookup would resolve "something.net" using public DNS
infrastructure. Then the second DNS lookup on the subdomain "coap"
and the host "node" would typically be resolved by a DNS server
operated by the owner of domain "something.net". So this domain
owner can manage its own internal node names and subdomain allocation
which would correspond to the CoAP namespace

4.3. Scheme Security Mapping

In general, regardless of the URI mapping scheme used in the Cross-
Protocol Proxy, an "https" request SHOULD be translated to a "coaps"
request. The exception case being cases where security on the CoAP
side is not needed because the network is well enough protected
already by other means (e.g. strong link-layer security, or the CoAP
network runs inside a firewalled network, etc.).

5. HTTP-CoAP Reverse Proxy

A HTTP-CoAP Reverse Cross-Protocol Proxy is accessed by web clients
only supporting HTTP, and handles their requests by mapping these to
CoAP requests, which are forwarded to CoAP servers; and mapping back
the received CoAP responses to HTTP. This mechanism is transparent
to the client, which may assume that it is communicating with the
intended target HTTP server. In other words, the client accesses the
proxy as an origin server using the "origin-form"
[I-D.ietf-httpbis-p1-messaging] as a Request Target.

Normative requirements on the translation of HTTP requests to CoAP
and of the CoAP responses back to HTTP responses are defined in
Section 10.2 of [I-D.ietf-core-coap]. However, that section only
considers the case of a HTTP-CoAP Forward Cross-Protocol Proxy in
which a client explicitly indicates it targets a request to a CoAP
server, and does not cover all aspects of proxy implementation in
detail. The present section provides guidelines and more details for
the implementation of a Reverse Cross-Protocol Proxy, which MAY be
followed in addition to the normative requirements.

Translation of unicast HTTP requests into multicast CoAP requests is
currently out of scope since in a reverse proxy scenario a HTTP

client typically expects to receive a single response, not multiple.
However a Cross-Protocol Proxy MAY include custom application-
specific functions to generate a multicast CoAP request based on a
unicast HTTP request and aggregate multiple CoAP responses into a
single HTTP response.

Note that the guidelines in this section also apply to an HTTP-CoAP
Intercepting Cross-Protocol Proxy.

5.1. Proxy Placement

Typically, a Cross-Protocol Proxy is located at the edge of the
constrained network. See Figure 1. The arguments supporting server-
side (SS) placement are the following:

Caching: Efficient caching requires that all request traffic to a
   CoAP server is handled by the same proxy which receives HTTP
   requests from multiple source locations. This maximally reduces
   the load on (constrained) CoAP servers.

Multicast: To support CoAPs use of local-multicast functionalities
   available in a constrained network, the Cross-Protocol Proxy
   requires a network interface directly attached to the constrained
   network.

TCP/UDP: Translation between HTTP and CoAP requires also TCP/UDP
   translation; TCP may be the preferred way for communicating with
   the constrained network due to its reliability or due to
   intermediate gateways configured to block UDP traffic.

Arguments against SS placement, in favor of client-side (CS), are:

Scalability: A solution where a single SS proxy has to manage
   numerous open TCP/IP connections to a large number of HTTP clients
   is not scalable. (Unless multiple SS proxies are employed with a
   load-balancing mechanism, which adds complexity.)

```
                         +------+
                         |      |
                         | DNS  |
                         |      |
                         +------+        Constrained Network
                                        --------------------
                                       /                    \
                                      /  /-----\    /-----\ \
                                     /   CoAP      CoAP    \
                                    /   server    server    \
                                   ||   \-----/    \-----/   ||
    +------+     HTTP Request      +----------+               ||
    |HTTP  |--------------------->| HTTP/CoAP| Req  /-----\   ||
    |Client|                      | Cross-   |------->| CoAP   ||
    |      |<---------------------| Proxy    |<-------|server  ||
    +------+     HTTP Response     +----------+ Resp  \-----/   ||
                                   ||                           ||
                                   ||  /-----\                  ||
                                   ||  CoAP                     ||
                                    \  server                   /
                                     \ \-----/                 /
                                      \      /-----\          /
                                       \     CoAP            /
                                        \    server        /
                                         \   \-----/      /
                                          ----------------
```
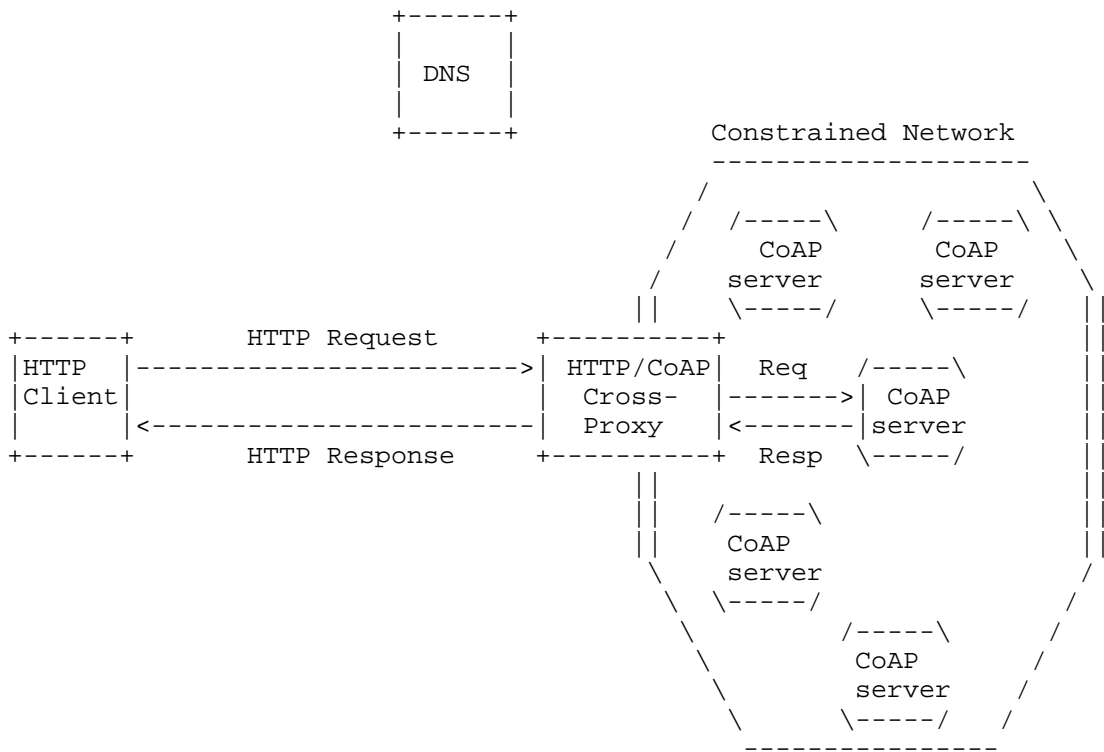
           Figure 1: Reverse Cross-Protocol Proxy Deployment Scenario

5.2. Response Code Translations

   Table 1 defines all possible CoAP responses along with the HTTP
   response to which each CoAP response SHOULD be translated. This
   table complies with the Section 10.2 requirements of
   [I-D.ietf-core-coap] and is intended to cover all possible cases.
   Multiple appearances of a HTTP status code in the second column
   indicates multiple equivalent HTTP responses are possible, depending
   on the conditions cited in the Notes (third column).

```
+----------------------------+----------------------------+-------+
| CoAP Response Code         | HTTP Status Code           | Notes |
+----------------------------+----------------------------+-------+
| 2.01 Created               | 201 Created                | 1     |
| 2.02 Deleted               | 200 OK                     | 2     |
|                            | 204 No Content             | 2     |
| 2.03 Valid                 | 304 Not Modified           | 3     |
|                            | 200 OK                     | 4     |
| 2.04 Changed               | 200 OK                     | 2     |
|                            | 204 No Content             | 2     |
| 2.05 Content               | 200 OK                     |       |
| 4.00 Bad Request           | 400 Bad Request            |       |
| 4.01 Unauthorized          | 400 Bad Request            | 5     |
| 4.02 Bad Option            | 400 Bad Request            | 6     |
| 4.03 Forbidden             | 403 Forbidden              |       |
| 4.04 Not Found             | 404 Not Found              |       |
| 4.05 Method Not Allowed    | 400 Bad Request            | 7     |
| 4.06 Not Acceptable        | 406 Not Acceptable         |       |
| 4.12 Precondition Failed   | 412 Precondition Failed    |       |
| 4.13 Request Entity Too    | 413 Request Repr. Too Large|       |
| Large                      |                            |       |
| 4.15 Unsupported Media Type| 415 Unsupported Media Type |       |
| 5.00 Internal Server Error | 500 Internal Server Error  |       |
| 5.01 Not Implemented       | 501 Not Implemented        |       |
| 5.02 Bad Gateway           | 502 Bad Gateway            |       |
| 5.03 Service Unavailable   | 503 Service Unavailable    | 8     |
| 5.04 Gateway Timeout       | 504 Gateway Timeout        |       |
| 5.05 Proxying Not Supported| 502 Bad Gateway            | 9     |
+----------------------------+----------------------------+-------+
```

Table 1: HTTP-CoAP Response Mapping

Notes:

1. A CoAP server may return an arbitrary format payload along with
    this response. This payload SHOULD be returned as entity in the
    HTTP 201 response. Section 7.3.2 of
    [I-D.ietf-httpbis-p2-semantics] does not put any requirement on
    the format of the payload. (In the past, [RFC2616] did.)

2. The HTTP code is 200 or 204 respectively for the case that a CoAP
    server returns a payload or not. [I-D.ietf-httpbis-p2-semantics]
    Section 5.3 requires code 200 in case a representation of the
    action result is returned for DELETE, POST and PUT and code 204
    if not. Hence, a proxy SHOULD transfer any CoAP payload
    contained in a 2.02 response to the HTTP client in a 200 OK
    response.

3. A CoAP 2.03 (Valid) response only (1) confirms that the request
    ETag is valid and (2) provides a new Max-Age value. HTTP 304
    (Not Modified) also updates some header fields of a stored
    response. A non-caching proxy may not have enough information to
    fill in the required values in the HTTP 304 (Not Modified)
    response, so it may not be advisable for a non-caching proxy to
    provoke the 2.03 (Valid) response by forwarding an ETag. A
    caching proxy will fill the information out of the cache.

4. A 200 response to a CoAP 2.03 occurs only when the proxy is
    caching and translated a HTTP request (without validation
    request) to a CoAP request that includes validation, for
    efficiency. The proxy receiving 2.03 updates the freshness of
    the cached representation and returns the entire representation
    to the HTTP client.

5. The HTTP code 401 Unauthorized MUST NOT be used, as long as in
    CoAP there is no equivalent defined of the required WWW-
    Authenticate header (Section 3.1 of [I-D.ietf-httpbis-p7-auth]).

6. In some cases a proxy receiving 4.02 may retry the request with
    less CoAP Options in the hope that the server will understand the
    newly formulated request. For example, if the proxy tried using
    a Block Option which was not recognized by the CoAP server it may
    retry without that Block Option.

7. The HTTP code "405 Method Not Allowed" MUST NOT be used since
    CoAP does not provide enough information to determine a value for
    the required "Allow" response-header field.

8. The value of the HTTP "Retry-After" response-header field is
    taken from the value of the CoAP Max-Age Option, if present.

9. This CoAP response can only happen if the proxy itself is
    configured to use a CoAP Forward Proxy to execute some, or all,
    of its CoAP requests.

5.3. Media Type Translations

   A Cross-Protocol Proxy translates a media type string, carried in a
   HTTP Content-Type header in a request, to a CoAP Content-Format
   Option with the equivalent numeric value. The media types supported
   by CoAP are defined in the CoAP Content-Format Registry. Any HTTP
   request with a Content-Type for which the proxy does not know an
   equivalent CoAP Content-Format number, MUST lead to HTTP response 415
   (Unsupported Media Type).

   Also, a CoAP Content-Format value in a response is translated back to

the equivalent HTTP Content-Type. If a proxy receives a CoAP
Content-Format value that it does not recognize (e.g. because the
value is IANA-registered after the proxy software was deployed), and
is unable to look up the equivalent HTTP Content-Type on the fly, the
proxy SHOULD return an HTTP entity (payload) without Content-Type
header (complying to Section 3.1.1.5 of
[I-D.ietf-httpbis-p2-semantics]).

5.4. Caching and Congestion Control

   A Cross-Protocol Proxy SHOULD limit the number of requests to CoAP
   servers by responding, where applicable, with a cached representation
   of the resource.

   Duplicate idempotent pending requests by a Cross-Protocol Proxy to
   the same CoAP resource SHOULD in general be avoided, by duplexing the
   response to the requesting HTTP clients without duplicating the CoAP
   request.

   If the HTTP client times out and drops the HTTP session to the Cross-
   Protocol Proxy (closing the TCP connection) after the HTTP request
   was made, a Cross-Protocol Proxy SHOULD wait for the associated CoAP
   response and cache it if possible. Further requests to the Cross-
   Protocol Proxy for the same resource can use the result present in
   cache, or, if a response has still to come, the HTTP requests will
   wait on the open CoAP session.

   According to [I-D.ietf-core-coap], a proxy MUST limit the number of
   outstanding interactions to a given CoAP server to NSTART. To limit
   the amount of aggregate traffic to a constrained network, the Cross-
   Protocol Proxy SHOULD also pose a limit to the number of concurrent
   CoAP requests pending on the same constrained network; further
   incoming requests MAY either be queued or dropped (returning 503
   Service Unavailable). This limit and the proxy queueing/dropping
   behavior SHOULD be configurable. In order to efficiently apply this
   congestion control, the Cross-Protocol Proxy SHOULD be SS placed.

   Resources experiencing a high access rate coupled with high
   volatility MAY be observed [I-D.ietf-core-observe] by the Cross-
   Protocol Proxy to keep their cached representation fresh while
   minimizing the number CoAP messages. See Section 5.5.

5.5. Cache Refresh via Observe

   There are cases where using the CoAP observe protocol
   [I-D.ietf-core-observe] to handle proxy cache refresh is preferable
   to the validation mechanism based on ETag as defined in
   [I-D.ietf-core-coap]. Such scenarios include, but are not limited

to, sleepy nodes -- with possibly high variance in requests'
distribution -- which would greatly benefit from a server driven
cache update mechanism. Ideal candidates would also be crowded or
very low throughput networks, where reduction of the total number of
exchanged messages is an important requirement.

This subsection aims at providing a practical evaluation method to
decide whether the refresh of a cached resource R is more efficiently
handled via ETag validation or by establishing an observation on R.

Let T_R be the mean time between two client requests to resource R,
let F_R be the freshness lifetime of R representation, and let M_R be
the total number of messages exchanged towards resource R. If we
assume that the initial cost for establishing the observation is
negligible, an observation on R reduces M_R iff $T\_R < 2*F\_R$ with
respect to using ETag validation, that is iff the mean arrival time
of requests for resource R is greater than half the refresh rate of
R.

When using observations M_R is always upper bounded by 2*F_R: in the
constrained network no more than 2*F_R messages will be generated
towards resource R.

5.6. Use of CoAP Blockwise Transfer

A Cross-Protocol Proxy SHOULD support CoAP blockwise transfers
[I-D.ietf-core-block] to allow transport of large CoAP payloads while
avoiding excessive link-layer fragmentation in LLNs, and to cope with
small datagram buffers in CoAP end-points as described in
[I-D.ietf-core-coap] Section 4.6.

A Cross-Protocol Proxy SHOULD attempt to retry a payload-carrying
CoAP PUT or POST request with blockwise transfer if the destination
CoAP server responded with 4.13 (Request Entity Too Large) to the
original request. A Cross-Protocol Proxy SHOULD attempt to use
blockwise transfer when sending a CoAP PUT or POST request message
that is larger than a value BLOCKWISE_THRESHOLD. The value of
BLOCKWISE_THRESHOLD MAY be implementation-specific, for example
calculated based on a known or typical UDP datagram buffer size for
CoAP end-points, or set to N times the size of a link-layer frame
where e.g. N=5, or preset to a known IP MTU value, or set to a known
Path MTU value. The value BLOCKWISE_THRESHOLD or parameters from
which it is calculated SHOULD be configurable in a proxy
implementation.

The Cross-Protocol Proxy SHOULD detect CoAP end-points not supporting
blockwise transfers by checking for a 4.02 (Bad Option) response
returned by an end-point in response to a CoAP request with a Block*

Option. This allows the Cross-Protocol Proxy to be more efficient,
not attempting repeated blockwise transfers to CoAP servers that do
not support it. However if a request payload is too large to be sent
as a single CoAP request and blockwise transfer would be unavoidable,
the proxy still SHOULD attempt blockwise transfer on such an end-
point before returning 413 (Request Entity Too Large) to the HTTP
client.

For improved latency a cross proxy MAY initiate a blockwise CoAP
request triggered by an incoming HTTP request even when the HTTP
request message has not yet been fully received, but enough data has
been received to send one or more data blocks to a CoAP server
already. This is particularly useful on slow client-to-proxy
connections.

5.7. Security Translation

A HC proxy SHOULD implement explicit rules for security context
translations. A translation may involve e.g. applying a rule that
any "https" request is translated to a "coaps" request, or e.g.
applying a rule that a "https" request is translated to an unsecured
"coap" request. Another rule could specify the security policy and
parameters used for DTLS connections. Such rules will largely depend
on the application and network context in which a proxy is applied.
To enable widest possible use of a proxy implementation, these rules
SHOULD be configurable in a HC proxy.

5.8. Other guidelines

For long delays of a CoAP server, the HTTP client or any other proxy
in between MAY timeout. Further discussion of timeouts in HTTP is
available in Section 6.2.4 of [I-D.ietf-httpbis-p1-messaging].

A cross proxy MUST define an internal timeout for each pending CoAP
request, because the CoAP server may silently die before completing
the request. The timeout value SHOULD be approximately less than or
equal to MAX_RTT defined in [I-D.ietf-core-coap].

When the DNS protocol is not used between CoAP nodes in a constrained
network, defining valid FQDN (i.e., DNS entries) for constrained CoAP
servers, where possible, MAY help HTTP clients to access the
resources offered by these servers via a HC proxy.

HTTP connection pipelining (section 6.2.2.1 of
[I-D.ietf-httpbis-p1-messaging]) MAY be supported by the proxy and is
transparent to the CoAP network: the HC cross proxy will sequentially
serve the pipelined requests by issuing different CoAP requests.

6. IANA Considerations

   This memo includes no request to IANA.

7. Security Considerations

   The security concerns raised in Section 15.7 of [RFC2616] also apply
   to the cross proxy scenario. In fact, the cross proxy is a trusted
   (not rarely a transparently trusted) component in the network path.

   The trustworthiness assumption on the cross proxy cannot be dropped.
   Even if we had a blind, bi-directional, end-to-end, tunneling
   facility like the one provided by the CONNECT method in HTTP, and
   also assuming the existence of a DTLS-TLS transparent mapping, the
   two tunneled ends should be speaking the same application protocol,
   which is not the case. Basically, the protocol translation function
   is a core duty of the cross proxy that can't be removed, and makes it
   a necessarily trusted, impossible to bypass, component in the
   communication path.

   A reverse proxy deployed at the boundary of a constrained network is
   an easy single point of failure for reducing availability. As such,
   a special care should be taken in designing, developing and operating
   it, keeping in mind that, in most cases, it could have fewer
   limitations than the constrained devices it is serving.

   The following sub paragraphs categorize and argue about a set of
   specific security issues related to the translation, caching and
   forwarding functionality exposed by a cross proxy module.

7.1. Traffic overflow

   Due to the typically constrained nature of CoAP nodes, particular
   attention SHOULD be posed in the implementation of traffic reduction
   mechanisms (see Section 5.4), because inefficient implementations can
   be targeted by unconstrained Internet attackers. Bandwidth or
   complexity involved in such attacks is very low.

   An amplification attack to the constrained network may be triggered
   by a multicast request generated by a single HTTP request mapped to a
   CoAP multicast resource, as considered in Section TBD of
   [I-D.ietf-core-coap].

   The impact of this amplification technique is higher than an
   amplification attack carried out by a malicious constrained device
   (e.g. ICMPv6 flooding, like Packet Too Big, or Parameter Problem on
   a multicast destination [RFC4732]), since it does not require direct

access to the constrained network.

The feasibility of this attack, disruptive in terms of CoAP server availability, can be limited by access controlling the exposed HTTP multicast resource, so that only known/authorized users access such URIs.

7.2. Handling Secured Exchanges

It is possible that the request from the client to the cross proxy is sent over a secured connection. However, there may or may not exist a secure connection mapping to the other protocol. For example, a secure distribution method for multicast traffic is complex and MAY not be implemented (see [I-D.ietf-core-groupcomm]).

By default, a cross proxy SHOULD reject any secured client request if there is no configured security policy mapping. This recommendation MAY be relaxed in case the destination network is believed to be secured by other, complementary, means. E.g.: assumed that CoAP nodes are isolated behind a firewall (e.g. as the SS cross proxy deployment shown in Figure 1), the cross proxy may be configured to translate the incoming HTTPS request using plain CoAP (i.e. NoSec mode.)

The HC URI mapping MUST NOT map to HTTP (see Section 4) a CoAP resource intended to be accessed only using HTTPS.

A secured connection that is terminated at the cross proxy, i.e. the proxy decrypts secured data locally, raises an ambiguity about the cacheability of the requested resource. The cross proxy SHOULD NOT cache any secured content to avoid any leak of secured information. However in some specific scenario, a security/efficiency trade-off could motivate caching secured information; in that case the caching behavior MAY be tuned to some extent on a per-resource basis.

8. Acknowledgements

An initial version of the table found in Section 5.2 has been provided in revision -05 of [I-D.ietf-core-coap]. Special thanks to Peter van der Stok for countless comments and discussions on this document, that contributed to its current structure and text.

Thanks to Carsten Bormann, Zach Shelby, Michele Rossi, Nicola Bui, Michele Zorzi, Klaus Hartke, Cullen Jennings, Kepeng Li, Brian Frank, Peter Saint-Andre, Kerry Lynn, Linyi Tian, Dorothy Gellert, Francesco Corazza for helpful comments and discussions that have shaped the document.

9. References

9.1. Normative References

   [I-D.ietf-core-block]
              Bormann, C. and Z. Shelby, "Blockwise transfers in CoAP",
              draft-ietf-core-block-10 (work in progress), October 2012.

   [I-D.ietf-core-coap]
              Shelby, Z., Hartke, K., Bormann, C., and B. Frank,
              "Constrained Application Protocol (CoAP)",
              draft-ietf-core-coap-13 (work in progress), December 2012.

   [I-D.ietf-core-groupcomm]
              Rahman, A. and E. Dijk, "Group Communication for CoAP",
              draft-ietf-core-groupcomm-05 (work in progress),
              February 2013.

   [I-D.ietf-core-observe]
              Hartke, K., "Observing Resources in CoAP",
              draft-ietf-core-observe-07 (work in progress),
              October 2012.

   [I-D.ietf-httpbis-p1-messaging]
              Fielding, R. and J. Reschke, "Hypertext Transfer Protocol
              (HTTP/1.1): Message Syntax and Routing",
              draft-ietf-httpbis-p1-messaging-22 (work in progress),
              February 2013.

   [I-D.ietf-httpbis-p2-semantics]
              Fielding, R. and J. Reschke, "Hypertext Transfer Protocol
              (HTTP/1.1): Semantics and Content",
              draft-ietf-httpbis-p2-semantics-22 (work in progress),
              February 2013.

   [I-D.ietf-httpbis-p7-auth]
              Fielding, R. and J. Reschke, "Hypertext Transfer Protocol
              (HTTP/1.1): Authentication", draft-ietf-httpbis-p7-auth-22
              (work in progress), February 2013.

   [RFC2119] Bradner, S., "Key words for use in RFCs to Indicate
              Requirement Levels", BCP 14, RFC 2119, March 1997.

   [RFC2616] Fielding, R., Gettys, J., Mogul, J., Frystyk, H.,
              Masinter, L., Leach, P., and T. Berners-Lee, "Hypertext
              Transfer Protocol -- HTTP/1.1", RFC 2616, June 1999.

   [RFC3986] Berners-Lee, T., Fielding, R., and L. Masinter, "Uniform
              Resource Identifier (URI): Generic Syntax", STD 66,
              RFC 3986, January 2005.

9.2. Informative References

   [I-D.bormann-core-simple-server-discovery]
              Bormann, C., "CoRE Simple Server Discovery",
              draft-bormann-core-simple-server-discovery-01 (work in
              progress), March 2012.

   [I-D.shelby-core-resource-directory]
              Shelby, Z., Krco, S., and C. Bormann, "CoRE Resource
              Directory", draft-shelby-core-resource-directory-04 (work
              in progress), July 2012.

   [RFC3040] Cooper, I., Melve, I., and G. Tomlinson, "Internet Web
              Replication and Caching Taxonomy", RFC 3040, January 2001.

   [RFC4732] Handley, M., Rescorla, E., and IAB, "Internet Denial-of-
              Service Considerations", RFC 4732, December 2006.

Authors' Addresses

   Angelo P. Castellani
   University of Padova
   Via Gradenigo 6/B
   Padova 35131
   Italy

   Email: angelo@castellani.net

   Salvatore Loreto
   Ericsson
   Hirsalantie 11
   Jorvas 02420
   Finland

   Email: salvatore.loreto@ericsson.com

Akbar Rahman
InterDigital Communications, LLC
1000 Sherbrooke Street West
Montreal H3A 3G4
Canada

Phone: +1 514 585 0761
Email: Akbar.Rahman@InterDigital.com

Thomas Fossati
KoanLogic
Via di Sabbiuno 11/5
Bologna 40136
Italy

Phone: +39 051 644 82 68
Email: tho@koanlogic.com

Esko Dijk
Philips Research
High Tech Campus 34
Eindhoven 5656 AE
The Netherlands

Email: esko.dijk@philips.com