

CONCERTO and BRAVO

Experiences with practical, full implementations of Network Coding systems

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Projects in Networking and Communications

DARPA CBMANET/ CONCERTO



Network Coding system for MANETs

ONR Adaptive Networks/ BRAVO BRAVO NG



Network Coding using local info, high mobility

DARPA IAMANET/ PIANO



Secure Network Coding

DARPA SAFER/ SONATA



Anonymous Comms through Net Coding



CBMANET/ CONCERTO Motivation

Challenges in Wireless MANETs

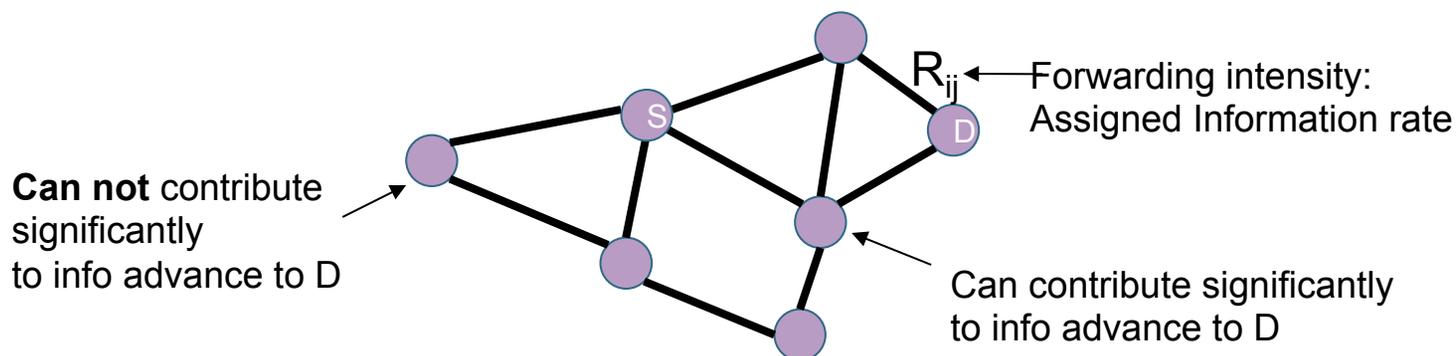
- Wireless communication
 - MANET Broadcast transmissions – use Rx is better than ignore
 - Reception is prone to errors – needs as much help as can get
- Mobility removes certainty from traditional routing

CONCERTO Approach

- Clean slate network and transport layers
- Network Coding Transport: fluid model
- Routing on Subgraphs
 - Enabled by fluid model
 - Robust to link errors and topo change

Subgraph Computation

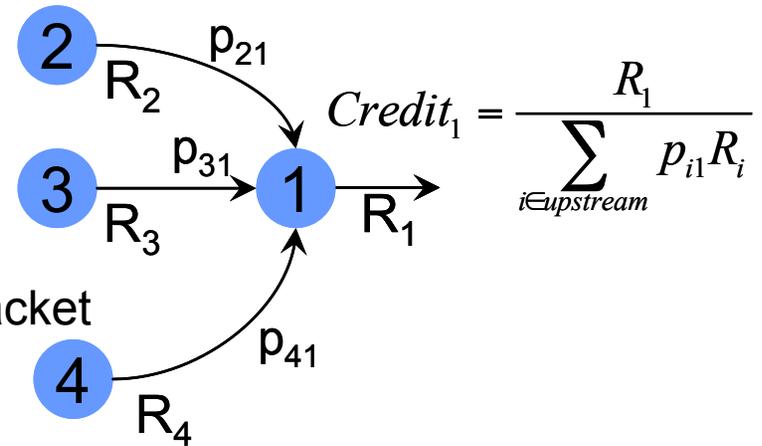
- Optimal Subgraph computation is complex (exponential in no. neighbors)
- Simplified approach: similar to MORE by Katabi et al
 - Requires Link State information
 - Computational complexity is $O(n^2)$ (n nodes in network)
 - Computes rate at which nodes should participate in forwarding mixtures
- Basic Concept
 - Forwarding nodes are chosen from those that are most probable to advance the information propagation to each destination
 - Forwarding intensity is related to the probability of contributing to info advance



Packet Forwarding

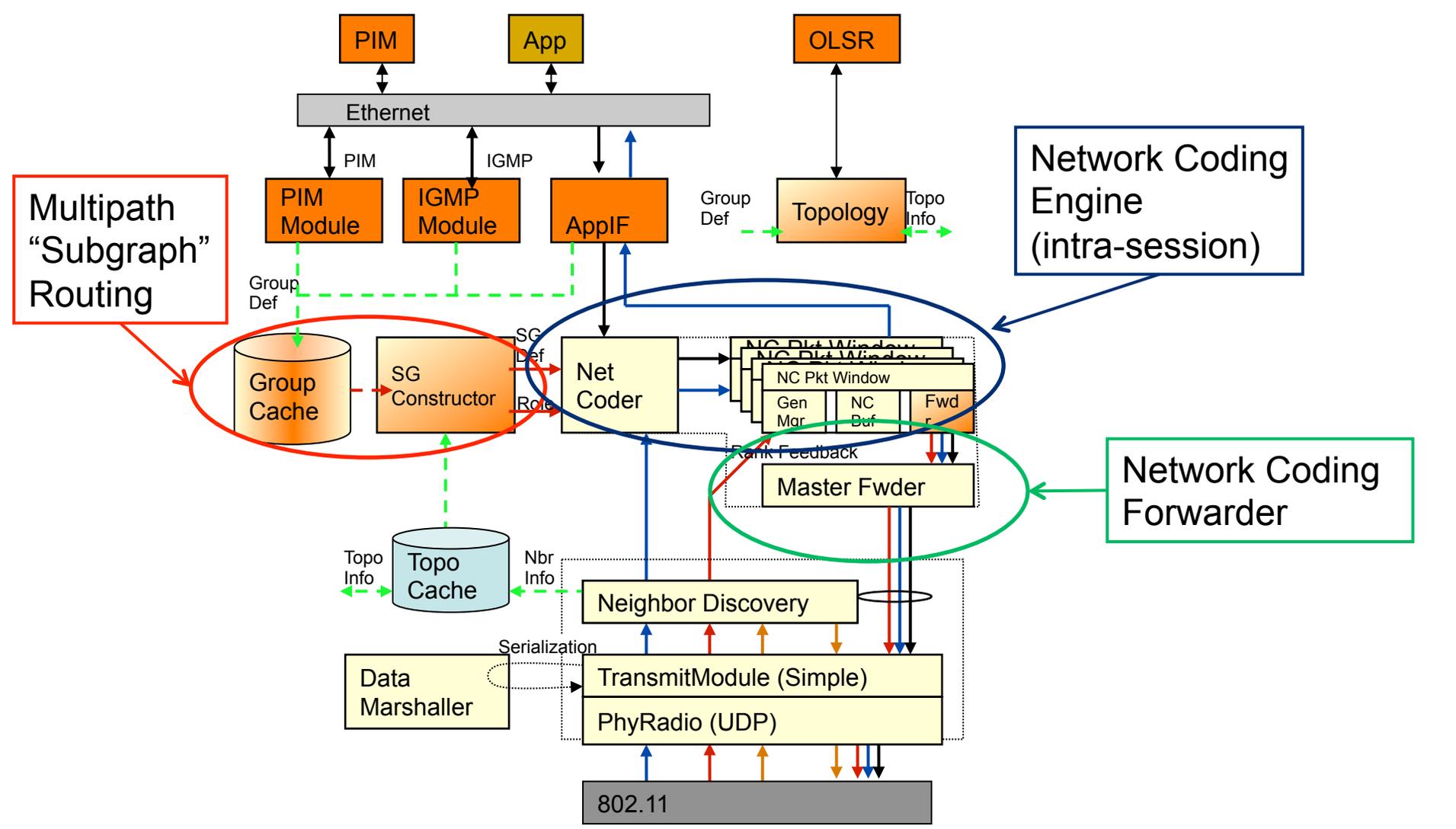
- Packet Forwarding

- Uses transmit rates computed by Subgraph computation
- Computes “receive credit” as ratio of transmit rate to expected receive rate
- Node earns credit when it receives innov. packet
- Spends credit when it transmits packets
- Provides automatic scaling to source rate



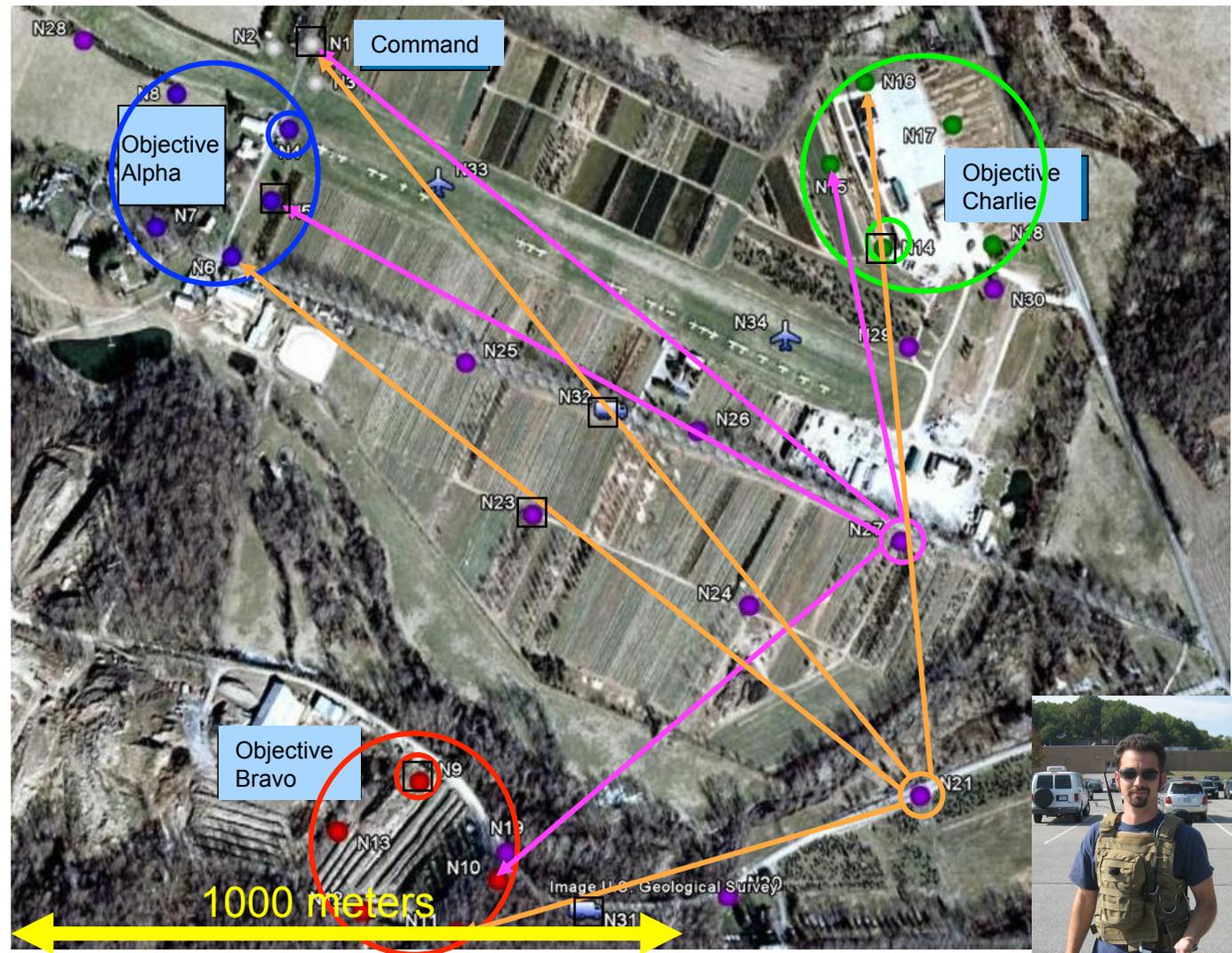
- Repair process, semi-reliable protocol (media streaming)
 - Nodes request **locally** extra transmissions if necessary
 - Requests are piggy-backed on forwarded packets when possible
- Fully reliable protocol (file transfer)
 - Additional algorithm for propagation of repair requests, detection of missing generations, beginning/end of files, late join, etc.

CONCERTO Architecture

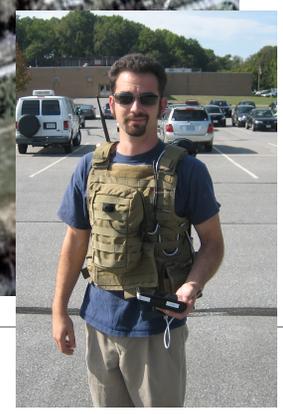


Layerless Modular Architecture: module info sharing → Adaptation to wireless dynamics

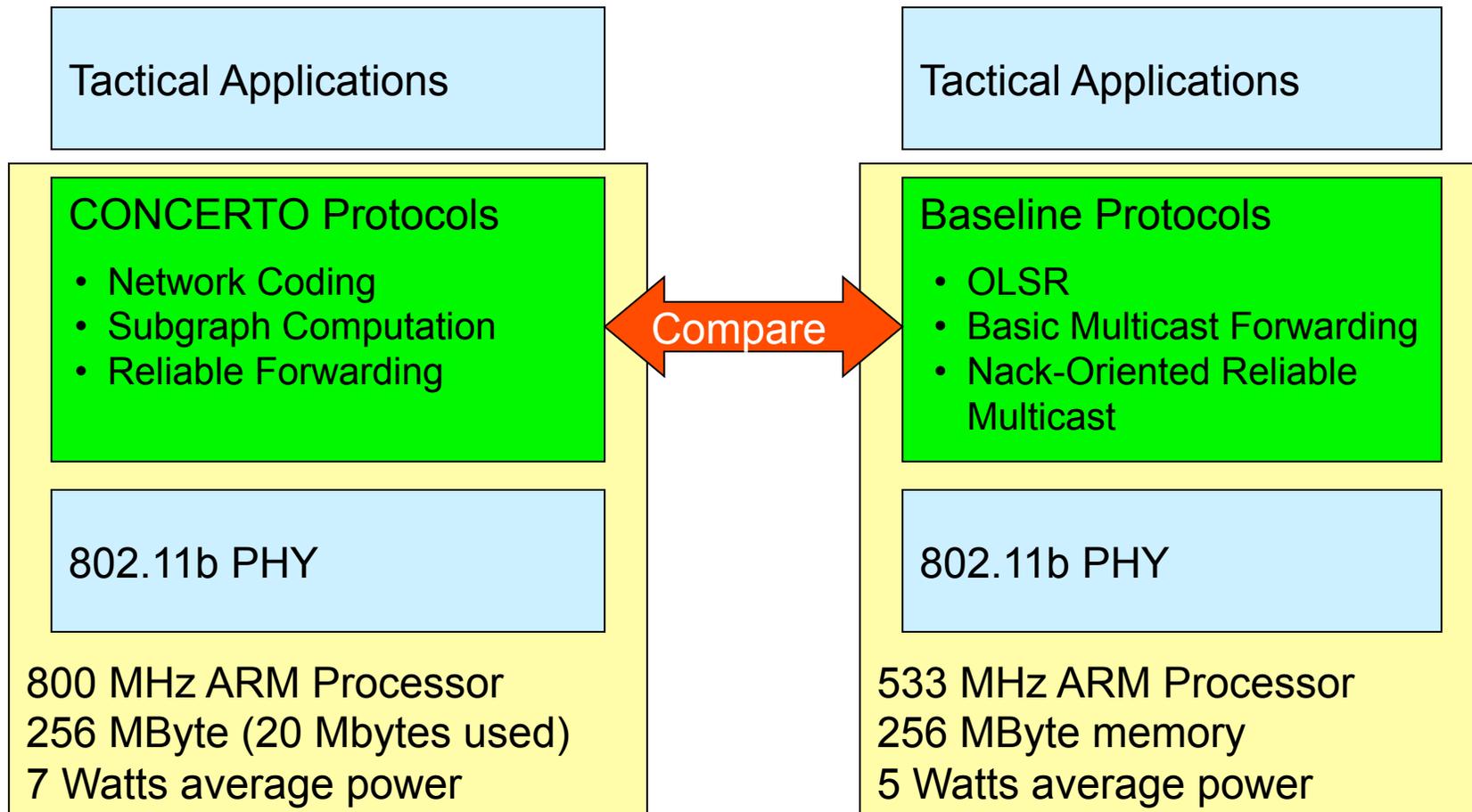
Scenario



- 31 mobile nodes:
- 3 moving groups of 5 people each
- 2 Vehicles traveling on access roads
- Tested ***with and without*** 2 aircraft
- Applications
 - Chat
 - Video
 - File Exchange

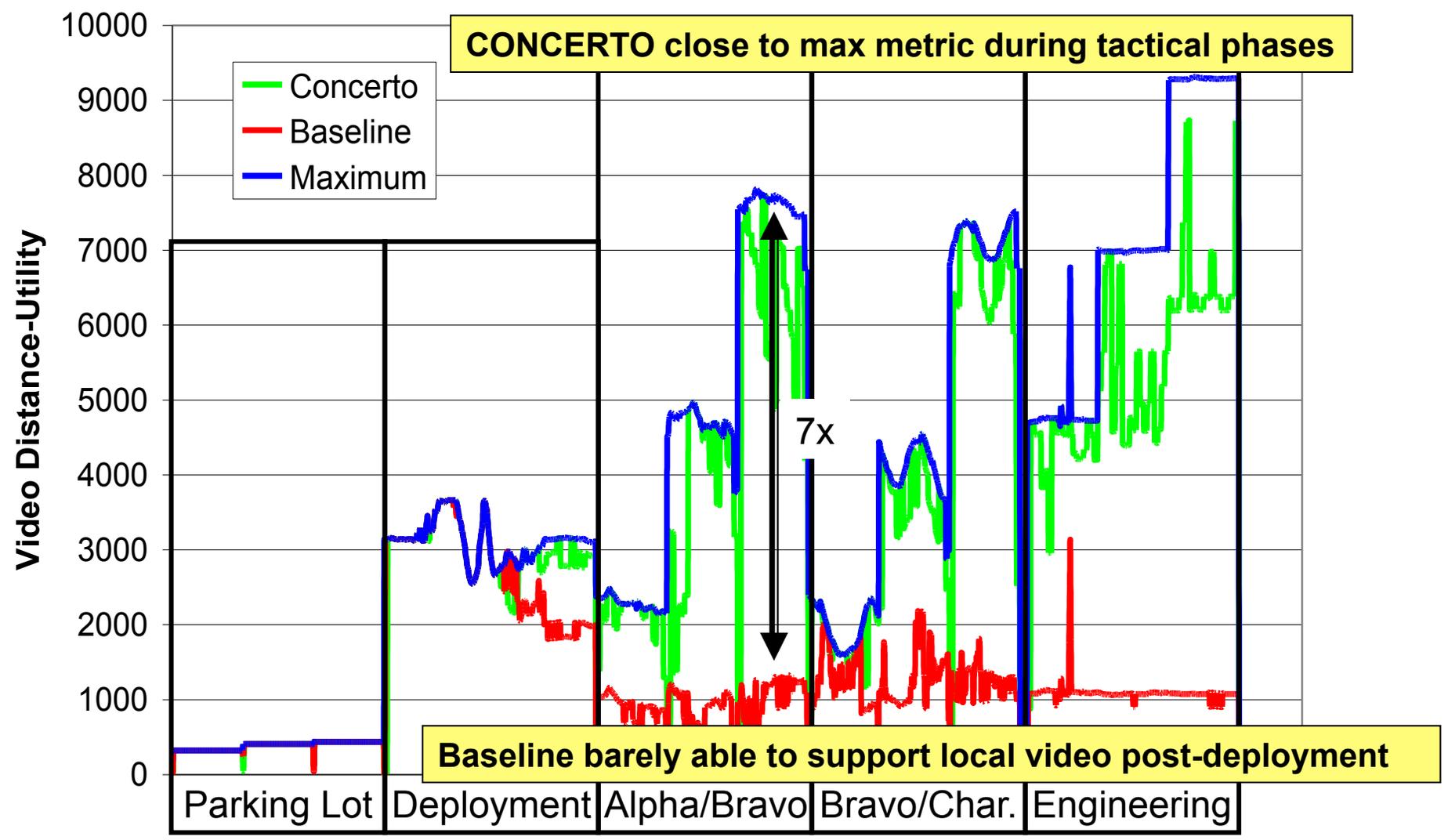


CONCERTO-Baseline Comparison



Common applications and PHY allow “apples-to-apples” comparison

Air Scenario – Distance Utility Metric



CONCERTO provides 7x higher performance than baseline in Air Scenario

Challenges In Subgraph Construction

- Requires Link State from all network
- Inaccurate LS info results in inefficient or defective subgraph

Solution: subgraph construction based on local info

- Gradient-based Routing
- Constructs field and currents for information flow
- Made possible by network coding fluid flow
- Provides natural properties of (Electric) Potential Fields
 - Efficient flow allocation
 - Locality of perturbation, Stability
 - No local minima: data never stalls in the middle

Gradient-based Routing

Information Flow Modeled on Electric Network



Objective: Minimize total transmissions: $G = \sum \{all\ links\ ab\} |I_{ab}| / Q_{ab}$

- I_{ab} = Information rate from a to b; Q_{ab} = P[success Tx from a to b]

Solution: Data network – equivalent to electric network:

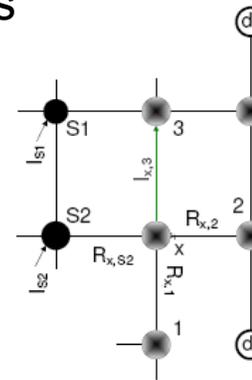
- Approximation: minimize $G = \sum (I_{ab})^2 / Q_{ab} \rightarrow$ minimize $E = \sum (I_{ab})^2 R_{ab}$
- Information rate from a to b \rightarrow intensity of current I_{ab}
- Link quality from a to b = $Q_{ab} \rightarrow 1/Resistance = 1/R_{ab}$
- Minimize Data Transmission: reduced to minimizing total electric energy
–Classic problem of solving electric circuits

• Network Level

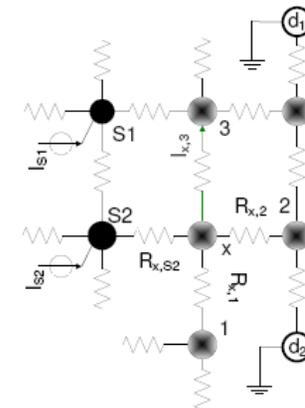
$$E = \frac{1}{2} \sum_{x,y} I_{x,y}^2 R_{x,y} = \frac{1}{2} \sum_{x,y} g_{x,y} (V_x - V_y)^2$$

Flow Level

$$I_{x,y} = \frac{(V_x - V_y)}{R_{x,y}}$$



(a) Communication Network



(b) Equivalent Electric Network

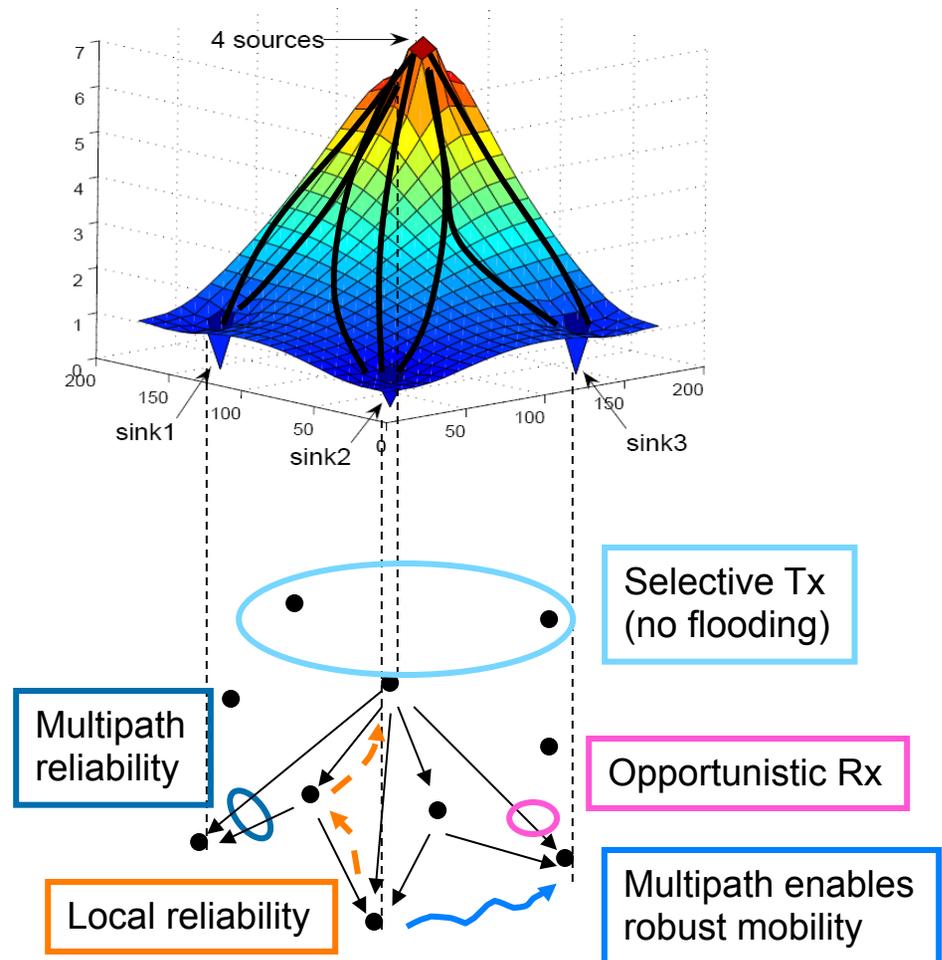
BRAVO: Network Coding over Gradient-based Subgraphs

Routing Challenges

- Reduced overhead through local info
- Multipath increases stability

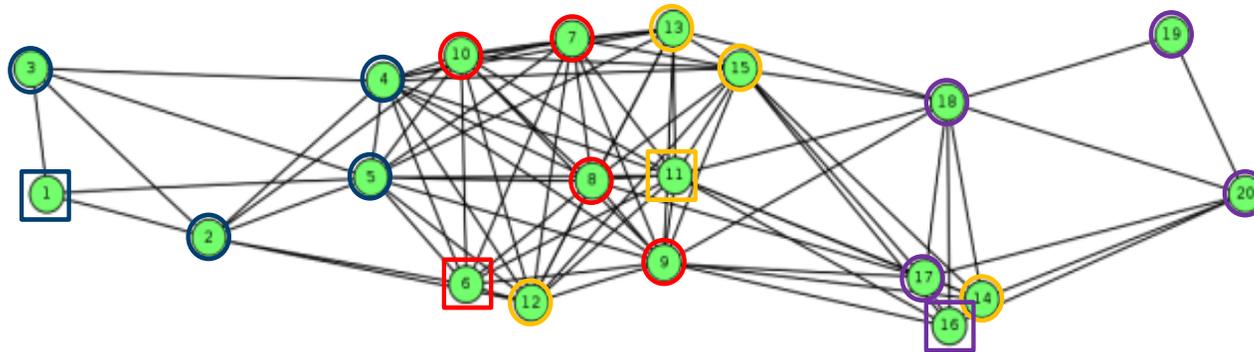
Data Transport Challenges

- Selective Tx is efficient
- Opportunistic Rx increases Tx efficiency
- Local & multi-path reliability increases end-end goodput effectiveness



Netcoded data over gradient routes gains from multipath, opportunistic network usage

BRAVO Performance



The Gauntlet Scenario

- 20 nodes, 2 stationary, 18 moving in random in regions

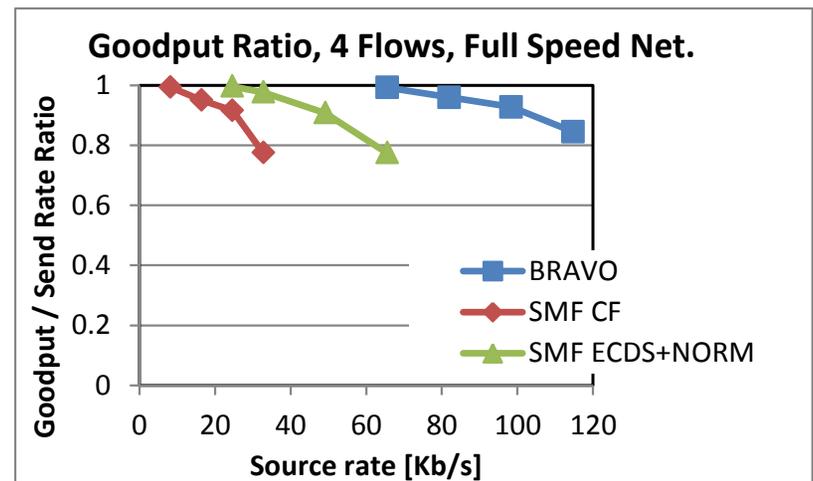
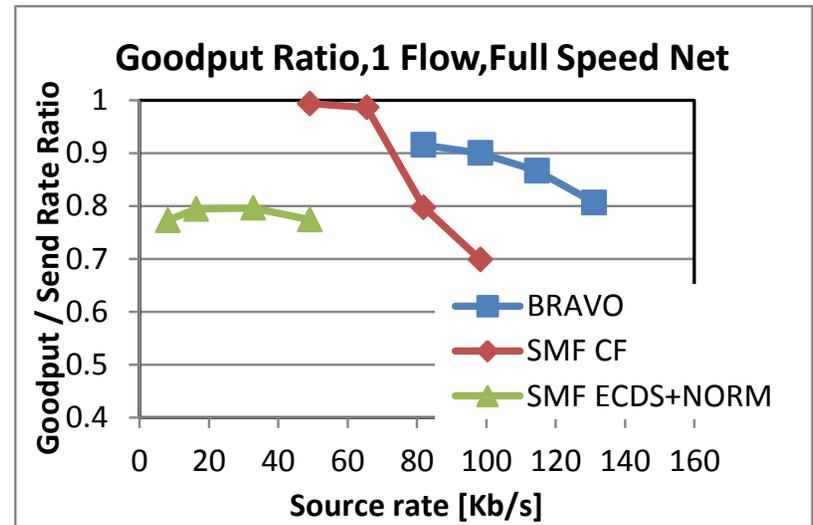
Scenario	Mean duration [s]	Std dev [s]
Full speed	42.275	54.530
Half speed	80.338	88.238
Quarter speed	138.518	143.131

BRAVO Outperforms SMF + NORM

- BRAVO robust at high speeds

Why:

- Gradient-based subgraph adjusts locally to topo changes
- Multi-path routing compensates the temporary lack of a link
- Network Coding adds redundancy: help deliver data, no matter where they come from



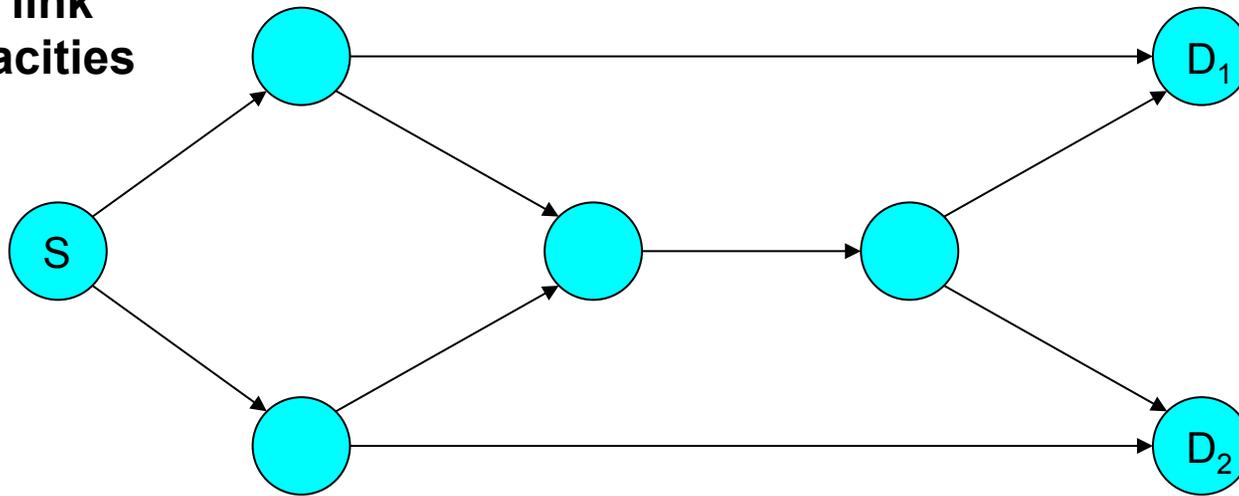
Conclusions

- Gains achieved besides information theoretical
- Network Coding enables multi-path reliability
- Network Coding enables hop-by-hop reliability
- Gradient-based subgraph routing
 - More efficient and scalable multipath routing

BACKUP

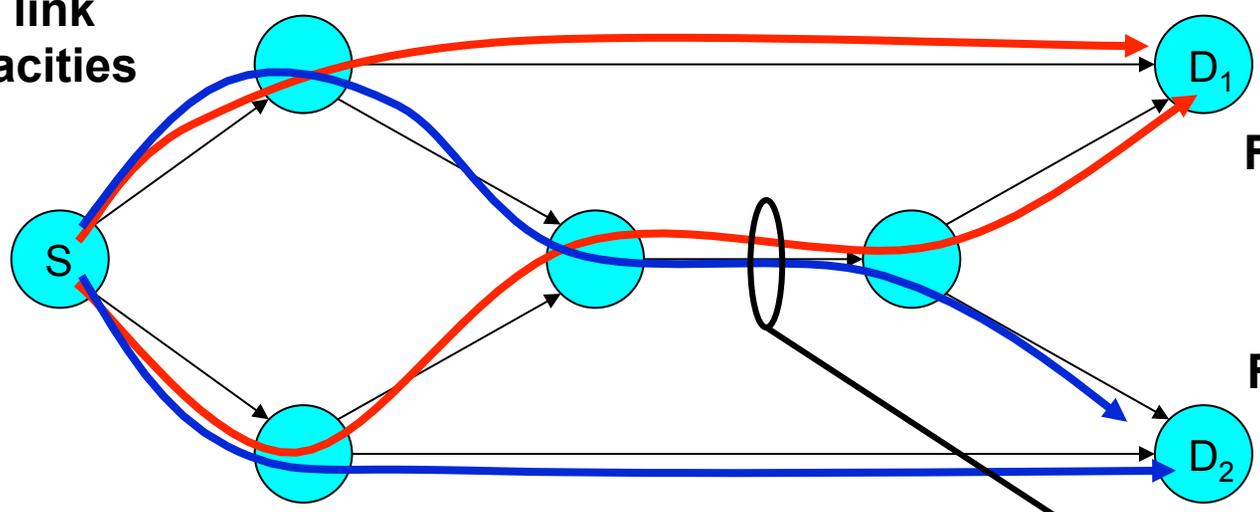
Intra-Session Coding: Wireline Butterfly

Unit link capacities



Wireline Butterfly with Multicast Routing

Unit link capacities

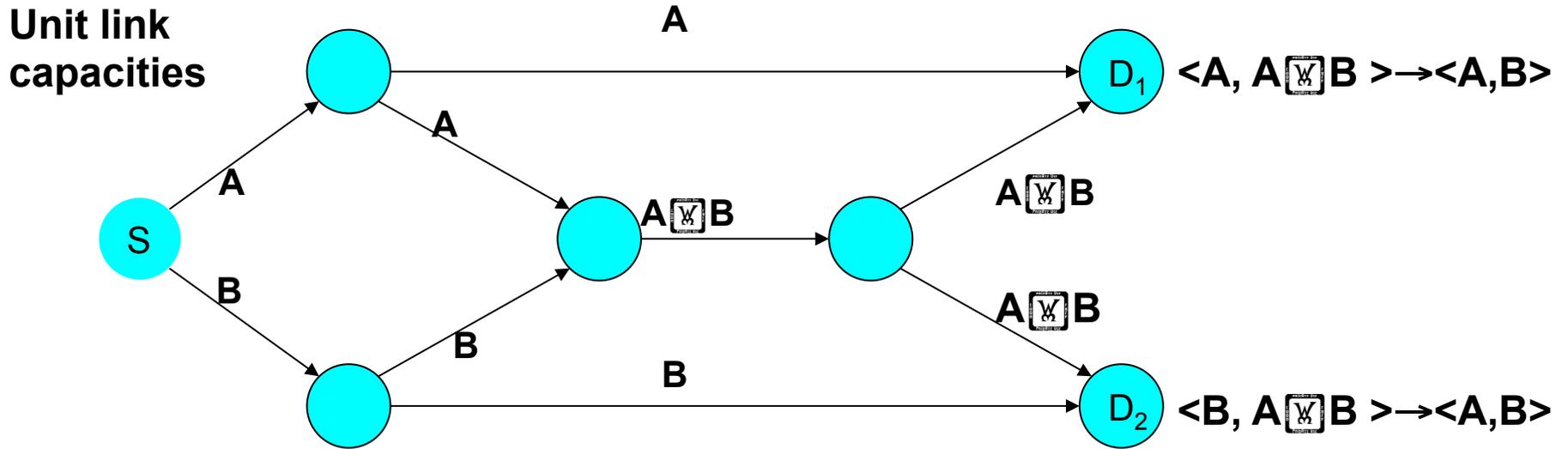


Flow of rate 2

Flow of rate 2

Multicast to D_1 and D_2 :
Can we support two flows of rate 2 with unit link constraint?

Wireline Butterfly with Intra-Session Network Coding:



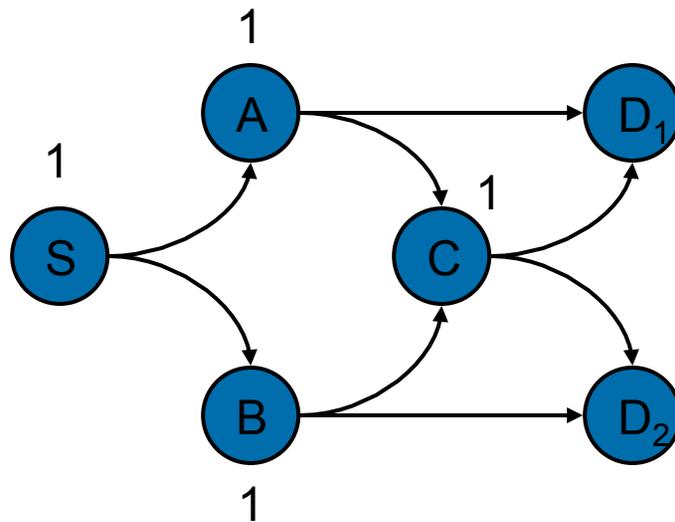
We can support “full rate” multicast to both destinations by XORing packets

Routing: Sum of “per destination” flows on a link must be less than link capacity

Network Coding: Maximum over “per destination” flows on a link must be less than link capacity

Wireless Butterfly with Multicast Routing

- Lossless links

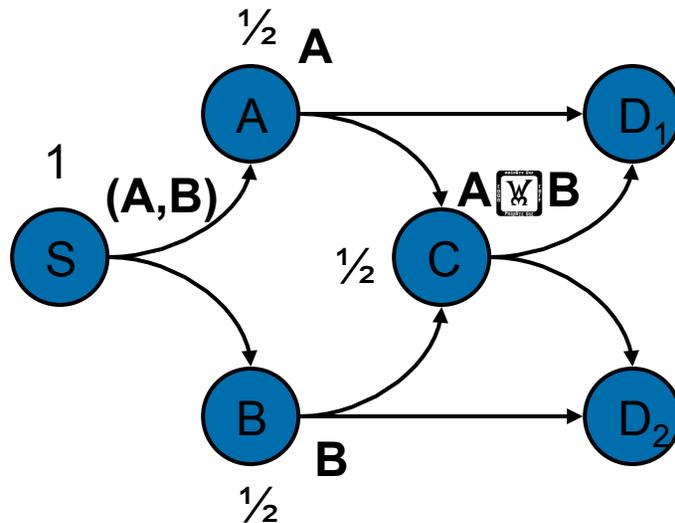


Routing: Sum of Tx Rates / symbol = 4

Contention-free schedule: S, A, B, C

Wireless Butterfly with Intra-Session Network Coding

- Lossless links



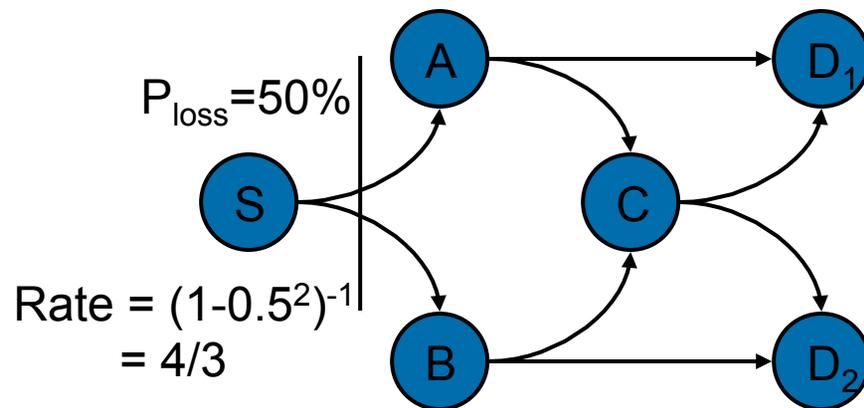
(Routing: Sum of Rates = 4)

NC: Sum of Tx Rates / symbol = 2.5

Contention-free schedule: S, S, A, B, C

Wireless Butterfly with Intra-Session Network Coding

- Lossy links



- When links are lossy, need only send at rate which transfers info to *some* next node on the other side of a cut set...

Intra-Session Network Coding

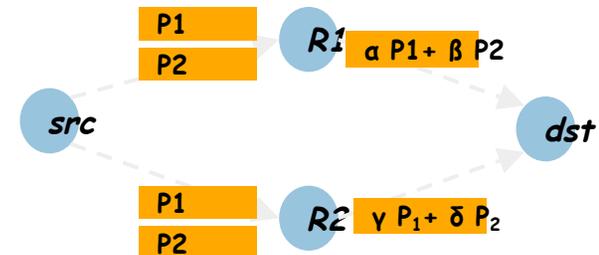
- Ahlswede et al. proved that for a single multicast session, network coding achieves the maximum possible rate allowed in the network
- Problem Decomposition
 - Computing Mixtures
 - How to combine packets into mixtures that “work”
 - Subgraph Construction
 - Which nodes forward mixtures and how much do they participate

Computing Mixtures: Random Linear Coding

- Ho et al.: **Random Linear Coding**

- Encoding:
$$Q_j = \sum_{i=1,n} A_{ij} P_i \quad \text{or} \quad \bar{Q} = A\bar{P}$$

- Identity of packets received – does not matter
- Only matters: Quantity of mixed packets = group size



- Practical Network Coding (Chou et al.)

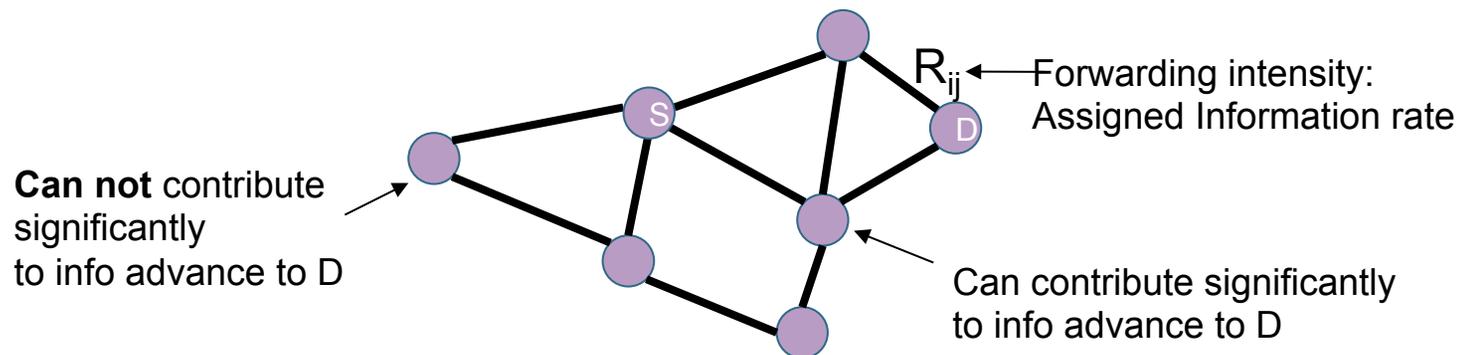
- Break file into N packets
- Collect packets into groups of G packets (generations)
- Source/intermediate nodes transmit random mixtures from generations
- Innovative (linearly independent) packets are stored for future mixtures
- Packet headers collect coefficients used in random coding
- Destination collects G linearly independent coded packets
- Inverts the matrix of random coefficients to recover original packets

$$\bar{P} = A^{-1}\bar{Q}$$

- Intra-session coding is basically matrix inversion

Subgraph Construction

- Not all nodes in a network should participate in each multicast flow
 - Want to maximize network capacity by minimizing transmissions
- Subgraph
 - The nodes which participate in random linear coding for a multicast flow
 - The rate at which these nodes forward random linear combinations



Subgraph Computation: Optimization Problem

$$\min \sum_{(i,j) \in A} f_{ij}(z_{ij}, R) \leftarrow \text{Convex cost function can capture goals of minimum energy, congestion, cost or maximum throughput}$$

$$\text{subject to } z_{ij} \geq x_{ij}^{(t)} \quad \forall (i, j) \in A, t \in T,$$

netcoded transmissions on link (i, j)
greater than *maximum per dest flow*

$$\sum_{(j|(i,j) \in A)} x_{ij}^{(t)} - \sum_{(j|(j,i) \in A)} x_{ji}^{(t)} = \begin{cases} R & \text{if } i = s, \\ -R & \text{if } i = t, \quad \forall i \in N, t \in T, \\ 0 & \text{otherwise} \end{cases}$$

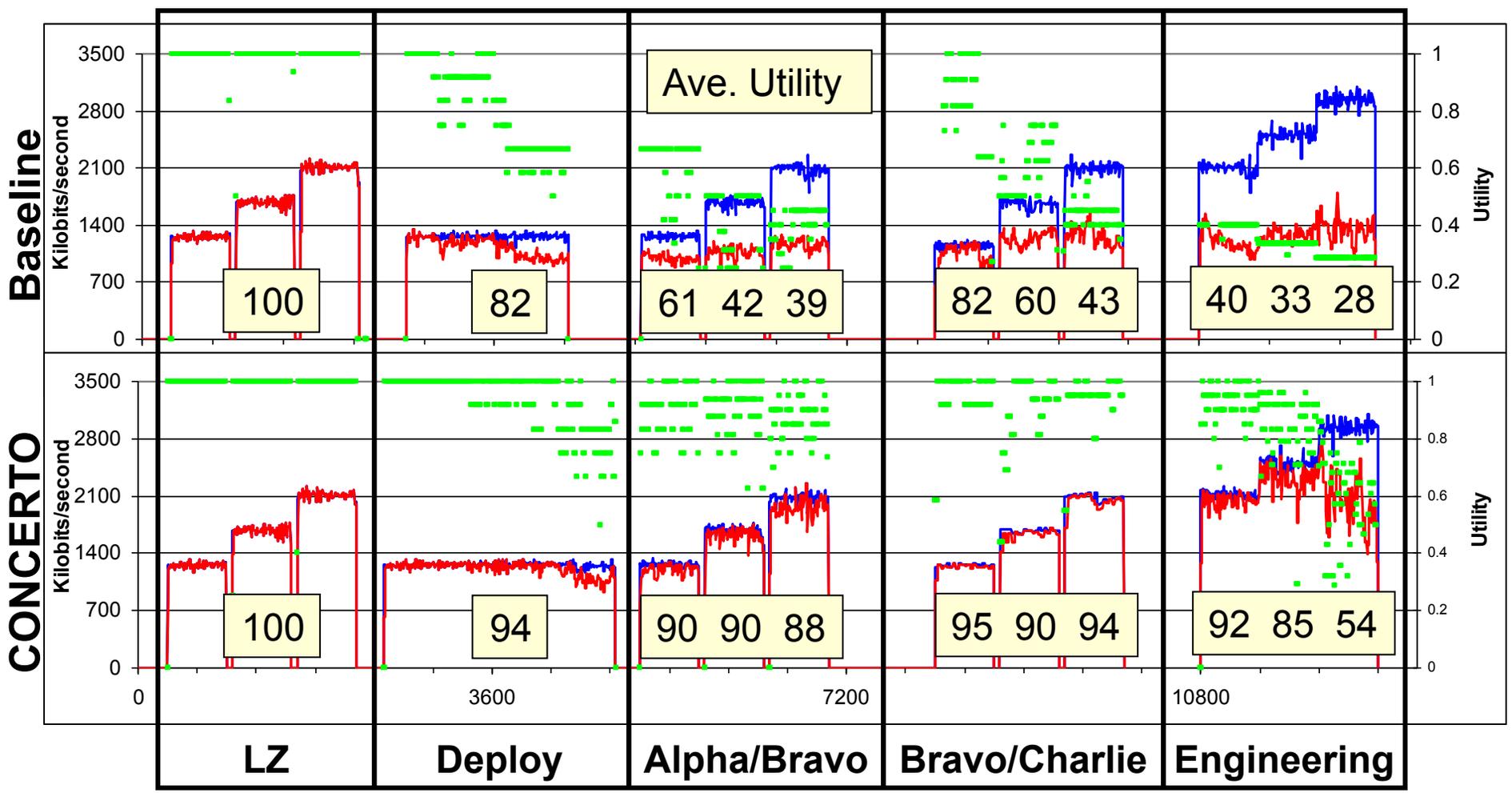
conservation of flow to destination t at node i

$$c_{ij} \geq x_{ij}^{(t)} \geq 0, \quad \forall (i, j) \in A, t \in T$$

link capacity constraints

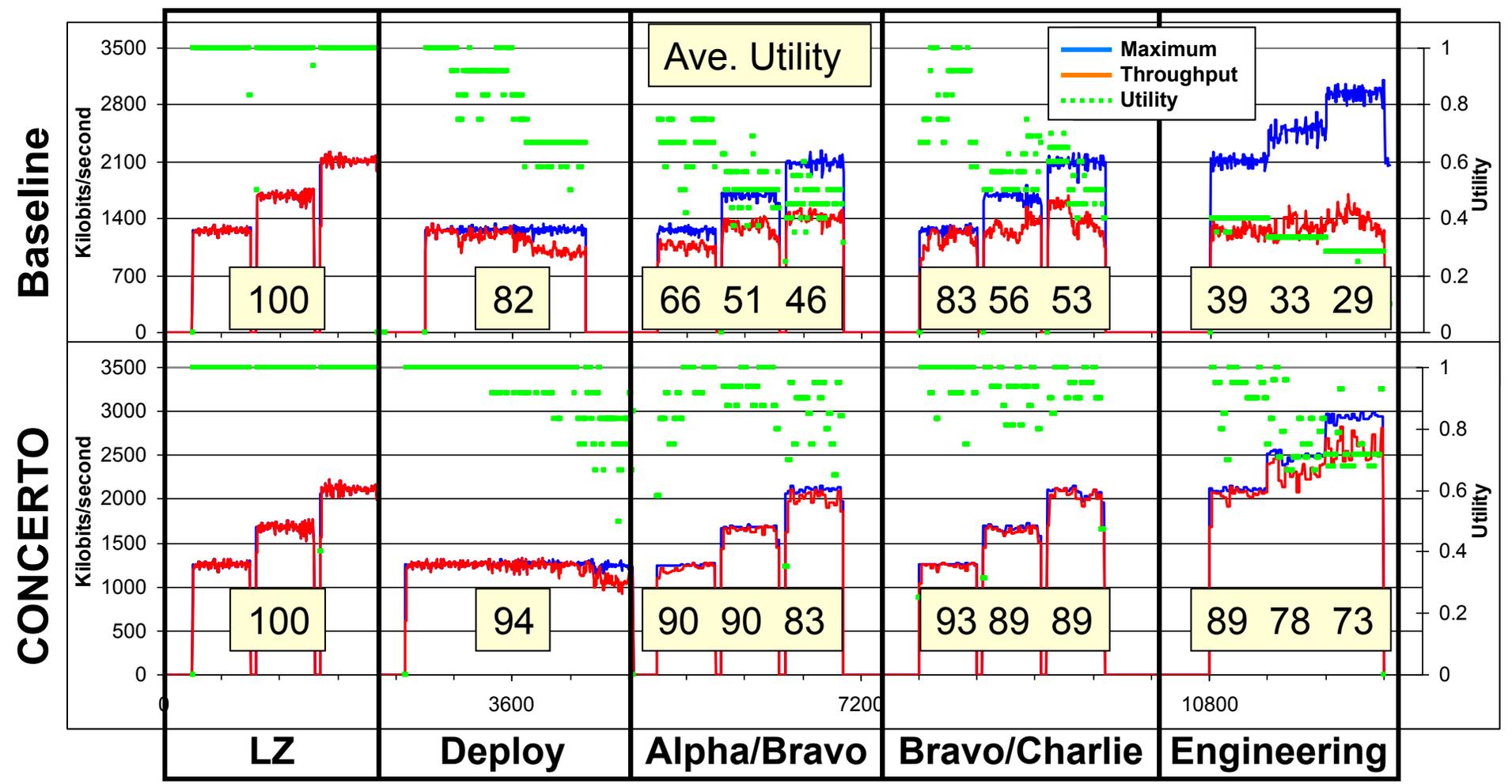
Optimal Subgraph computation is complex (exponential in no. neighbors)

Video Utility - Ground



Baseline barely usable at lowest load; CONCERTO works well at highest tactical loads

Video Utility - Air



Baseline barely usable at lowest load; CONCERTO works well at highest tactical loads