

TRILL Working Group
Internet Draft
Intended status: Standards Track
Expires: April 2014

T. Mizrahi
Marvell
T. Senevirathne
S. Salam
D. Kumar
Cisco
D. Eastlake 3rd
Huawei
October 10, 2013

Loss and Delay Measurement in
Transparent Interconnection of Lots of Links (TRILL)
<draft-ietf-trill-loss-delay-00.txt>

Abstract

Performance Monitoring (PM) is a key aspect of Operations, Administration and Maintenance (OAM). It allows network operators to verify the Service Level Agreement (SLA) provided to customers, and to detect network anomalies. This document specifies mechanisms for Loss Measurement and Delay Measurement in TRILL networks.

Status of this Memo

This Internet-Draft is submitted to IETF in full conformance with the provisions of BCP 78 and BCP 79.

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF), its areas, and its working groups. Note that other groups may also distribute working documents as Internet-Drafts.

Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as "work in progress."

The list of current Internet-Drafts can be accessed at
<http://www.ietf.org/ietf/lid-abstracts.txt>.

The list of Internet-Draft Shadow Directories can be accessed at
<http://www.ietf.org/shadow.html>.

This Internet-Draft will expire on April 10, 2014.

Copyright Notice

Copyright (c) 2013 IETF Trust and the persons identified as the document authors. All rights reserved.

This document is subject to BCP 78 and the IETF Trust's Legal Provisions Relating to IETF Documents (<http://trustee.ietf.org/license-info>) in effect on the date of publication of this document. Please review these documents carefully, as they describe your rights and restrictions with respect to this document. Code Components extracted from this document must include Simplified BSD License text as described in Section 4.e of the Trust Legal Provisions and are provided without warranty as described in the Simplified BSD License.

Table of Contents

1. Introduction	3
2. Conventions Used in this Document	4
2.1. Keywords	4
2.2. Definitions	4
2.3. Abbreviations	5
3. Loss and Delay Measurement in the TRILL Architecture	5
3.1. Performance Monitoring Granularity	6
3.2. One-Way vs. Two-Way Performance Monitoring	6
3.2.1. One-Way Performance Monitoring	7
3.2.2. Two-Way Performance Monitoring	7
3.3. Point-to-point vs. Point-to-multipoint Performance Monitoring	8
4. Loss Measurement	8
4.1. One-Way Loss Measurement	8
4.1.1. 1SL Message Transmission	9
4.1.2. 1SL Message Reception	10
4.2. Two-Way Loss Measurement	11
4.2.1. SLM Message Transmission	12
4.2.2. SLM Message Reception	12
4.2.3. SLR Message Reception	13
5. Delay Measurement	14
5.1. One-Way Delay Measurement	14
5.1.1. 1DM Message Transmission	15
5.1.2. 1DM Message Reception	16
5.2. Two-Way Delay Measurement	16
5.2.1. DMM Message Transmission	17
5.2.2. DMM Message Reception	17
5.2.3. DMR Message Reception	18
6. Packet Formats	19
6.1. TRILL OAM Encapsulation	19

6.2. Loss Measurement Packet Formats	21
6.2.1. Counter Format	21
6.2.2. ISL Packet Format	22
6.2.3. SLM Packet Format	23
6.2.4. SLR Packet Format	24
6.3. Delay Measurement Packet Formats	25
6.3.1. Timestamp Format	25
6.3.2. LDM Packet Format	25
6.3.3. DMM Packet Format	26
6.3.4. DMR Packet Format	27
7. Performance Monitoring Process	28
8. Security Considerations	29
9. IANA Considerations	29
9.1. OpCode Values	29
10. Acknowledgments	29
11. References	30
11.1. Normative References	30
11.2. Informative References	30

1. Introduction

TRILL [RFC2111] is a protocol for transparent least cost routing, where Rbridges route traffic to their destination based on least cost, using a TRILL encapsulation header with a hop count.

Operations, Administration and Maintenance (OAM) [OAM] is a set of tools for detecting, isolating and reporting connection failures and performance degradation. Performance Monitoring (PM) is a key aspect of OAM. PM allows network operators to detect and debug network anomalies and incorrect behavior. PM consists of two main building blocks - Loss Measurement and Delay Measurement. PM may also include other derived metrics such as Packet Delivery Rate, and Inter-Frame Delay Variation.

The requirements of OAM in TRILL networks are defined in [OAM-REQ], and the TRILL OAM framework is described in [OAM-FRAMEWK]. These two documents also highlight the main requirements in terms of performance monitoring.

This document defines protocols for loss measurement and for delay measurement in TRILL networks. These protocols are somewhat based on the mechanisms defined in ITU-T G.8013/Y.1731 [Y.1731].

- o Loss Measurement: the Loss Measurement protocol measures packet loss between two RBridges. The measurement is performed by sending a set of synthetic packets, and counting the number of packets transmitted and received during the test. The frame loss is calculated by comparing the numbers of transmitted and received packets.
This provides a statistical estimate of the packet loss between the involved RBridges, with a margin of error that can be controlled by varying the number of transmitted synthetic packets. This document does not define procedures for packet loss computation based on counting user data. For further details see [OAM-FRAMEWK].
- o Delay Measurement: the Delay Measurement protocol measures the packet delay and packet delay variation between two RBridges. The measurement is performed using timestamped OAM messages.

2. Conventions Used in this Document

2.1. Keywords

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [KEYWORDS].

The requirement level of PM in [OAM-REQ] is 'SHOULD'. Nevertheless, this memo uses the entire range of requirement levels, including 'MUST'; the requirements in this memo are to be read as 'A MEP that implements TRILL PM MUST/SHOULD/MAY/...'.

2.2. Definitions

- o One-way packet delay - (as defined in [OAM-REQ]) the time elapsed from the start of transmission of the first bit of a packet by an RBridge until the reception of the last bit of the packet by the remote RBridge.
- o Two-way packet delay - (as defined in [OAM-REQ]) the time elapsed from the start of transmission of the first bit of a packet from the local RBridge, receipt of the packet at the remote RBridge, the remote RBridge sending a response packet back to the local RBridge and the local RBridge receiving the last bit of that response packet.
- o Packet loss - the number of packets lost in a specific probe instance, and a specific observation period.

- o Far-end packet loss - the number of packets lost on the path from the local RBridge to the remote RBridge in a specific probe instance, and a specific observation period.
- o Near-end packet loss - the number of packets lost on the path from the remote RBridge to the local RBridge in a specific probe instance, and a specific observation period.

2.3. Abbreviations

1DM	One-way Delay Measurement message
1SL	One-way Synthetic Loss Measurement message
DMM	Delay Measurement Message
DMR	Delay Measurement Reply
MD	Maintenance Domain
MD-L	Maintenance Domain Level
MEP	Maintenance End Point
MIP	Maintenance Intermediate Point
MP	Maintenance Point
OAM	Operations, Administration and Maintenance
PM	Performance Monitoring
SLM	Synthetic Loss Measurement Message
SLR	Synthetic Loss Measurement Reply
TLV	Type, Length and Value
TRILL	Transparent Interconnection of Lots of Links

3. Loss and Delay Measurement in the TRILL Architecture

As described in [OAM-FRAMEWK], OAM protocols in a TRILL campus operate over two types of Maintenance Points (MPs): Maintenance End Points (MEPs) and Maintenance Intermediate Points (MIPs).

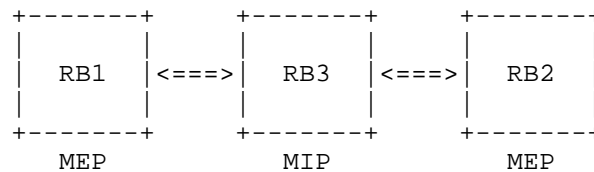


Figure 1 Maintenance Points in a TRILL Campus

Performance Monitoring (PM) allows a MEP to perform loss and delay measurements to any other MEP in the campus. Performance Monitoring is performed in the context of a specific Maintenance Domain (MD).

The PM functionality defined in this document is not applicable to MIPs.

3.1. Performance Monitoring Granularity

As defined in [OAM-FRAMEWK], PM can be applied at three levels of granularity: 'Network', 'Service' and 'Flow'.

- o Network-level PM: the PM protocol is run over a dedicated test VLAN or FGL.
- o Service-level PM: the PM protocol is used to perform measurements of actual user VLANs or FGL.
- o Flow-level PM: the PM protocol is used to perform measurements on a per-flow basis. A flow, as defined in [OAM-REQ], is a set of packets that share the same path and per-hop behavior (such as priority).
As defined in [OAM-FRAMEWK], flow-based monitoring uses a Flow Entropy field that resides at the beginning of the OAM packet header (see Section 6.1.), and mimics the forwarding behavior of the monitored flow.

3.2. One-Way vs. Two-Way Performance Monitoring

Paths in a TRILL network are not necessarily symmetric, i.e., a packet sent from RB1 to RB2 does not necessarily traverse the same set of RBridges or links as a packet sent from RB2 to RB1. Even within a given flow, packets from RB1 to RB2 do not necessarily traverse the same path as packets from RB2 to RB1. Therefore, this document provides tools for one-way performance monitoring and for two-way performance monitoring.

3.2.1. One-Way Performance Monitoring

In one-way PM, RB1 sends PM messages to RB2, allowing RB2 to monitor the performance on the path from RB1 to RB2.

A MEP that implements TRILL PM SHOULD support one-way performance monitoring. A MEP that implements TRILL PM SHOULD support both the functionality of the sender, RB1, and the functionality of the receiver, RB2.

One-way PM can be applied either proactively or on-demand, although the more typical scenario is the proactive mode, where RB1 and RB2 periodically transmit PM messages to each other, allowing each of them to monitor the performance on the incoming path from the peer MEP.

3.2.2. Two-Way Performance Monitoring

In two-way PM, a sender, RB1, sends PM messages to a reflector, RB2, and RB2 responds to these messages, allowing RB1 to monitor the performance of:

- o The path from RB1 to RB2.
- o The path from RB2 to RB1.
- o The two-way path from RB1 to RB2, and back to RB1.

Note that in some cases it may be interesting for RB1 to monitor only the path from RB1 to RB2. Two-way PM allows the sender, RB1, to monitor the path from RB1 to RB2, as opposed to one-way PM (Section 3.2.1.), which allows the receiver, RB2, to monitor this path.

A MEP that implements TRILL PM MUST support two-way PM. A MEP that implements TRILL PM MUST support both the sender and the reflector functionality.

As described in Section 3.1. , flow-based PM uses the Flow Entropy field as one of the parameters that identify a flow. In two-way PM, the Flow Entropy of the path from RB1 to RB2 is typically different from the Flow Entropy of the path from RB2 to RB1. This document uses the Reflector Entropy TLV [TRILL-FM,], which allows the sender to specify the Flow Entropy value to be used in the response message.

Two-way PM can be applied either proactively or on-demand.

3.3. Point-to-point vs. Point-to-multipoint Performance Monitoring

PM can be applied either as a point-to-point measurement protocol, or as a point-to-multi-point measurement protocol.

The point-to-point approach measures the performance between two RBridges using unicast PM messages.

In the point-to-multipoint approach, an RBridge RB1 sends PM messages to multiple RBridges using multicast messages. The reflectors (in two-way PM) respond to RB1 using unicast messages. To protect against reply storms, the reflectors MUST send the response messages after a random delay in the range of 0 to 2 seconds. This ensures that the responses are staggered in time, and that the initiating RBridge is not overwhelmed with responses. Moreover, a scope TLV [TRILL-FM] can be used to limit the set of RBridges from which a response is expected, thus reducing the impact of potential response bursts.

4. Loss Measurement

The Loss Measurement protocol has two flavors, one-way Loss Measurement, and two-way Loss Measurement.

Note: The terms 'one-way' and 'two-way' Loss Measurement should not be confused with the terms 'single-ended' and 'dual-ended' Loss Measurement used in [Y.1731]. As defined in Section 3.2. , the terms 'one-way' and 'two-way' specify whether the protocol monitors performance on one direction, or on both directions. The terms 'single-ended' and 'dual-ended', on the other hand, describe whether the protocol is asymmetric or symmetric, respectively.

4.1. One-Way Loss Measurement

One-way Loss Measurement measures the one-way packet loss from one MEP to another. The loss ratio is measured using a set of One-way Synthetic Loss Measurement (1SL) messages. The packet format of the 1SL message is specified in Section 6.2.2. Figure 2 illustrates a one-way Loss Measurement message exchange.

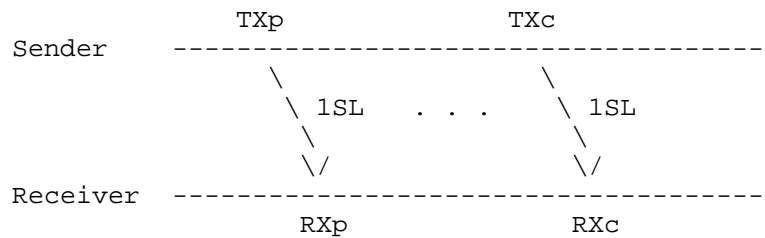


Figure 2 One-Way Loss Measurement

The one-way Loss Measurement procedure uses a set of 1SL messages to measure the packet loss. The figure shows two non-consecutive messages from the set.

The sender maintains a counter of transmitted 1SL messages, and includes the value of this counter, TX, in each 1SL message it transmits. The receiver maintains a counter of received 1SL messages, RX, and can calculate the loss by comparing its counter values to the counter values received in the 1SL messages.

In Figure 2, the subscript 'c' is an abbreviation for current, and 'p' is an abbreviation for previous.

4.1.1.1. 1SL Message Transmission

One-way Loss Measurement can be applied either proactively or on-demand, although as mentioned in Section 3.2.1. , it is more likely to be applied proactively.

The term 'on-demand' in the context of one-way Loss Measurement implies that the sender transmits a fixed set of 1SL messages, allowing the receiver to perform the measurement based on this set.

A MEP that supports one-way Loss Measurement **MUST** support unicast transmission of 1SL messages.

A MEP that supports one-way Loss Measurement **MAY** support multicast transmission of 1SL messages.

The sender **MUST** maintain a packet counter for each peer MEP and probe instance (test ID). Every time the sender transmits a 1SL packet, it

increments the corresponding counter, and then integrates the value of the counter into the <Counter TX> field of the 1SL packet.

The 1SL message MAY be sent with a variable size Data TLV, allowing loss measurement for various packet sizes.

4.1.2. 1SL Message Reception

The receiver MUST maintain a reception counter for each peer MEP and probe instance (test ID). Upon receiving a 1SL packet, the receiver MUST verify that:

- o The 1SL packet is destined to the current MEP.
- o The packet's MD level matches the MEP's MD level.

If both conditions are satisfied, the receiver increments the corresponding receive packet counter, and records the new value of the counter, RX1.

A MEP that supports one-way Loss Measurement MUST support reception of both unicast and multicast 1SL messages.

The receiver computes the one-way packet loss with respect to a probe instance measurement interval. A probe instance measurement interval includes a sequence of 1SL messages with the same test ID. The one-way packet loss is computed by comparing the counter values TXp and RXp at the beginning of the measurement interval, and the counter values TXc and RXc at the end of the measurement interval (Figure 2):

$$\text{one-way packet loss} = (\text{TXc} - \text{TXp}) - (\text{RXc} - \text{RXp}) \quad (1)$$

The calculation in Equation (1) is based on counter value differences, implying that the sender's counter, TX, and the receiver's counter, RX, are not required to be synchronized with respect to a common initial value.

It is noted that if the sender or receiver resets one of the counters, TX or RX, the calculation in Equation (1) produces a false measurement result. Hence the sender and receiver SHOULD NOT clear the TX and RX counters during a measurement interval.

When the receiver calculates the packet loss per Equation (1) it MUST perform a wraparound check. If the receiver detects that one of the counters has wrapped around, the receiver adjusts the result of Equation (1) accordingly.

A 1SL receiver MUST support reception of 1SL messages with a Data TLV.

Since synthetic one-way Loss Measurement is performed using 1SL messages, obviously some 1SL messages may be dropped during a measurement interval. Thus, when the receiver does not receive a 1SL, the receiver cannot perform the calculations in Equation (1) for that specific 1SL message.

4.2. Two-Way Loss Measurement

Two-way Loss Measurement allows a MEP to measure the packet loss on the paths to and from a peer MEP. Two-way Loss Measurement uses a set of Synthetic loss Measurement Messages (SLM) to compute the packet loss. Each SLM is answered with a Synthetic loss Measurement Reply (SLR). The packet formats of the SLM and SLR packets are specified in Sections 6.2.3. and 6.2.4. , respectively. Figure 2 illustrates a two-way Loss Measurement message exchange.

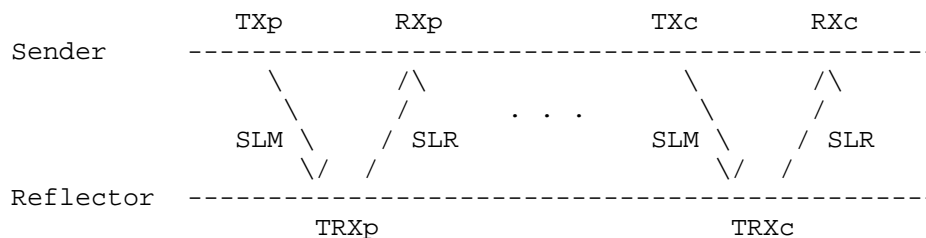


Figure 3 Two-Way Loss Measurement

The two-way Loss Measurement procedure uses a set of SLM-SLR handshakes. The figure shows two non-consecutive handshakes from the set.

The sender maintains a counter of transmitted SLM messages, and includes the value of this counter, TX, in each transmitted SLM message. The reflector maintains a counter of received SLM messages, TRX. The reflector generates an SLR, and incorporates TRX into the SLR packet. The sender maintains a counter of received SLR messages,

RX. Upon receiving an SLR message, the sender can calculate the loss by comparing the local counter values to the counter values received in the SLR messages.

The subscript 'c' is an abbreviation for current, and 'p' is an abbreviation for previous.

4.2.1. SLM Message Transmission

Two-way Loss Measurement can be applied either proactively or on-demand.

A MEP that supports two-way Loss Measurement MUST support unicast transmission of SLM messages.

A MEP that supports two-way Loss Measurement MAY support multicast transmission of SLM messages.

The sender MUST maintain a counter of transmitted SLM packets for each peer MEP and probe instance (test ID). Every time the sender transmits an SLM packet it increments the corresponding counter, and then integrates the value of the counter into the <Counter TX> field of the SLM packet.

A sender MAY include a Reflector Entropy TLV in an SLM message. The Reflector Entropy TLV format is specified in [TRILL-FM].

An SLM message MAY be sent with a Data TLV, allowing loss measurement for various packet sizes.

4.2.2. SLM Message Reception

The reflector MUST maintain a reception counter, TRX, for each peer MEP and probe instance (test ID).

Upon receiving an SLM packet, the reflector MUST verify that:

- o The SLM packet is destined to the current MEP.
- o The packet's MD level matches the MEP's MD level.

If both conditions are satisfied, the reflector increments the corresponding packet counter, and records the value of the new counter, TRX. The reflector then generates an SLR message that is identical to the received SLM, except for the following modifications:

- o The reflector incorporates TRX into the <Counter TRX> field of the SLR.
- o The <OpCode> field in the OAM header is set to the SLR OpCode.
- o The reflector assigns its MEP ID in the <Reflector MEP ID> field.
- o If the received SLM includes a Reflector Entropy TLV [TRILL-FM], the reflector copies the value of the Flow Entropy from the TLV into the <Flow Entropy> field of the SLR message. The outgoing SLR message does not include a Reflector Entropy TLV.
- o The TRILL header and transport header are modified to reflect the source and destination of the SLR packet. The SLR is always a unicast message.

A MEP that supports two-way Loss Measurement MUST support reception of both unicast and multicast SLM messages.

A reflector MUST support reception of SLM packets with a Data TLV. When receiving an SLM with a Data TLV, the reflector includes the unmodified TLV in the SLR.

4.2.3. SLR Message Reception

The sender MUST maintain a reception counter, RX, for each peer MEP and probe instance (test ID).

Upon receiving an SLR message, the sender MUST verify that:

- o The SLR packet is destined to the current MEP.
- o The <Sender MEP ID> field in the SLR packet matches the current MEP.
- o The packet's MD level matches the MEP's MD level.

If the conditions above are met, the sender increments the corresponding reception counter, and records the new value, RX.

The sender computes the packet loss with respect to a probe instance measurement interval. A probe instance measurement interval includes a sequence of SLM messages, and their corresponding SLR messages, all with the same test ID. The packet loss is computed by comparing the counters at the beginning of the measurement interval, denoted with a subscript 'p', and the counters at the end of the measurement interval, denoted with a subscript 'c' (as illustrated in Figure 3).

$$\text{far-end packet loss} = (\text{TXc-TXp}) - (\text{TRXc-TRXp}) \quad (2)$$

$$\text{near-end packet loss} = (\text{TRXc-TRXp}) - (\text{RXc-RXp}) \quad (3)$$

Note: total two-way packet loss is the sum of the far and near end packet losses, that is $(\text{TXc-TXp}) - (\text{RXc-RXp})$.

The calculations in the two equations above are based on counter value differences, implying that the sender's counters, TX and RX, and the reflector's counter, TRX, are not required to be synchronized with respect to a common initial value.

It is noted that if the sender or reflector resets one of the counters, TX, TRX or RX, the calculation in Equations (2) and (3) produces a false measurement result. Hence the sender and reflector SHOULD NOT clear the TX, TRX and RX counters during a measurement interval.

When the sender calculates the packet loss per Equations (2) and (3) it MUST perform a wraparound check. If the reflector detects that one of the counters has wrapped around, the reflector adjusts the result of Equations (2) and (3) accordingly.

Since synthetic two-way Loss Measurement is performed using SLM and SLR messages, obviously some SLM and SLR messages may be dropped during a measurement interval. When an SLM or an SLR is dropped, the corresponding two-way handshake (Figure 3) is not completed successfully, and thus the reflector does not perform the calculations in Equations (2) and (3) for that specific message exchange.

A sender MAY choose to monitor only the far-end packet loss, i.e., perform the computation in Equation (2), and ignore the computation in Equation (3). Note that, in this case, the sender can run flow-based PM of the path TO the peer MEP without using the Reflector Entropy TLV.

5. Delay Measurement

The Delay Measurement protocol has two flavors, One-Way Delay Measurement, and Two-Way Delay Measurement.

5.1. One-Way Delay Measurement

One-way Delay Measurement is used for computing the one-way packet delay from one MEP to another. The packet format used in one-way Delay Measurement is referred to as 1DM, and is specified in Section

6.3.2. The one-way Delay Measurement message exchange is illustrated in Figure 4.

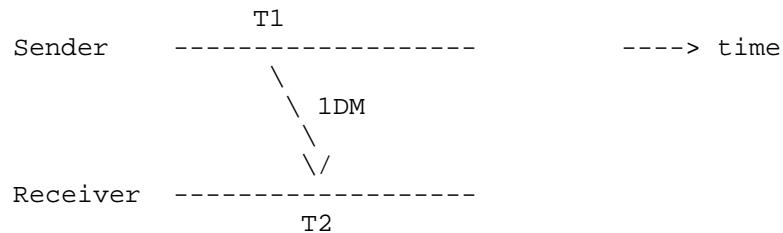


Figure 4 One-Way Delay Measurement

The sender transmits a 1DM message incorporating its time of transmission, T1. The receiver then receives the message at time T2, and calculates the one-way delay as:

$$\text{one-way delay} = T2 - T1 \quad (4)$$

Equation (4) implies that T2 and T1 are measured with respect to a common reference time. Hence, two MEPs running an one-way Delay Measurement protocol MUST be time-synchronized. The method used for synchronizing the clocks associated with the two MEPs is outside the scope of this document.

5.1.1.1. 1DM Message Transmission

1DM packets can be transmitted proactively or on-demand, although as mentioned in Section 3.2.1. , they are typically transmitted proactively.

A MEP that supports one-way Delay Measurement MUST support unicast transmission of 1DM messages.

A MEP that supports one-way Delay Measurement MAY support multicast transmission of 1DM messages.

A 1DM message MAY be sent with a variable size Data TLV, allowing packet delay measurement for various packet sizes.

The sender incorporates the 1DM packet's time of transmission into the <Timestamp T1> field.

5.1.2. 1DM Message Reception

Upon receiving a 1DM packet, the receiver records its time of reception, T2. The receiver MUST verify two conditions:

- o The 1DM packet is destined to the current MEP.
- o The packet's MD level matches the MEP's MD level.

If both conditions are satisfied, the receiver terminates the packet and calculates the one-way delay as specified in Equation (4).

A MEP that supports one-way Delay Measurement MUST support reception of both unicast and multicast 1DM messages.

A 1DM receiver MUST support reception of 1DM messages with a Data TLV.

When one-way Delay Measurement packets are received periodically, the receiver MAY compute the packet delay variation based on multiple measurements. Note that packet delay variation can be computed even when the two peer MEPs are not time synchronized.

5.2. Two-Way Delay Measurement

Two-way Delay Measurement uses a two-way handshake for computing the two-way packet delay between two MEPs. The handshake includes two packets, a Delay Measurement Message (DMM) and a Delay Measurement Reply (DMR). The DMM and DMR packet formats are specified in Section 6.3.3. and 6.3.4. , respectively.

The two-way Delay Measurement message exchange is illustrated in Figure 5.

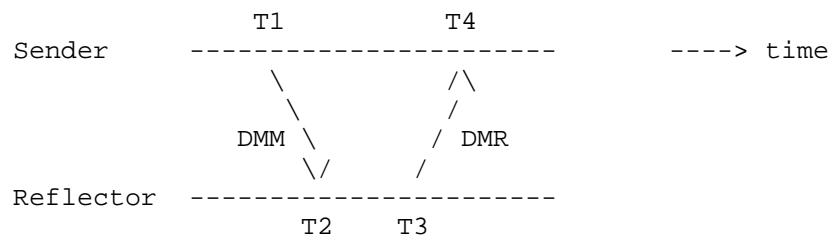


Figure 5 Two-Way Delay Measurement

The sender generates a DMM message incorporating its time of transmission, T1. The reflector receives the DMM message and records its time of reception, T2. The reflector then generates a DMR message, incorporating T1, T2 and the DMR's transmission time, T3. The sender receives the DMR message at T4, and using the 4 timestamps it calculates the two-way packet delay.

5.2.1. DMM Message Transmission

DMM packets can be transmitted periodically or on-demand.

A MEP that supports two-way Delay Measurement **MUST** support unicast transmission of DMM messages.

A MEP that supports two-way Delay Measurement **MAY** support multicast transmission of DMM messages.

A sender **MAY** include a Reflector Entropy TLV in a DMM message. The Reflector Entropy TLV format is specified in [TRILL-FM].

A DMM **MAY** be sent with a variable size Data TLV, allowing packet delay measurement for various packet sizes.

The sender incorporates the DMM packet's time of transmission into the <Timestamp T1> field.

5.2.2. DMM Message Reception

Upon receiving a DMM packet, the reflector records its time of reception, T2. The reflector **MUST** verify two conditions:

- o The DMM packet is destined to the current MEP.
- o The packet's MD level matches the MEP's MD level.

If both conditions are satisfied, the reflector terminates the packet, and generates a DMR packet. The DMR is identical to the received DMM, except for the following modifications:

- o The reflector incorporates T2 into the <Timestamp T2> field of the DMR.
- o The reflector incorporates the DMR's transmission time, T3, into the <Timestamp T3> field of the DMR.

- o The <OpCode> field in the OAM header is set to the DMR OpCode.
- o If the received DMM includes a Reflector Entropy TLV [TRILL-FM], the reflector copies the value of the Flow Entropy from the TLV into the <Flow Entropy> field of the DMR message. The outgoing DMR message does not include a Reflector Entropy TLV.
- o The TRILL header and transport header are modified to reflect the source and destination of the DMR packet. The DMR is always a unicast message.

A MEP that supports two-way Delay Measurement MUST support reception of both unicast and multicast DMM messages.

A reflector MUST support reception of DMM packets with a Data TLV. When receiving a DMM with a Data TLV, the reflector includes the unmodified TLV in the DMR.

5.2.3. DMR Message Reception

Upon receiving the DMR message, the sender records its time of reception, T4. The sender MUST verify:

- o The DMR packet is destined to the current MEP.
- o The packet's MD level matches the MEP's MD level.

If both conditions above are met, the sender uses the 4 timestamps to compute the two-way delay:

$$\text{two-way delay} = (T4 - T1) - (T3 - T2) \quad (5)$$

Note that two-way delay can be computed even when the two peer MEPs are not time synchronized. One-way Delay Measurement, on the other hand, requires the two MEPs to be synchronized.

Two MEPs running a two-way Delay Measurement protocol MAY be time-synchronized. If two-way Delay Measurement is run between two time-synchronized MEPs, the sender MAY compute the one-way delays:

$$\text{one-way delay \{sender->reflector\}} = T2 - T1 \quad (6)$$

$$\text{one-way delay \{reflector->sender\}} = T4 - T3 \quad (7)$$

When two-way Delay Measurement is run periodically, the sender MAY also compute the delay variation based on multiple measurements.

A sender MAY choose to monitor only the sender->reflector delay, i.e., perform the computation in Equation (6), and ignore the computations in (5) and (7). Note that in this case the sender can run flow-based PM of the path TO the peer MEP without using the Reflector Entropy TLV.

6. Packet Formats

6.1. TRILL OAM Encapsulation

The TRILL OAM encapsulation is defined in [OAM-FRAMEWK], and is quoted in this document for clarity. For further details see [OAM-FRAMEWK].

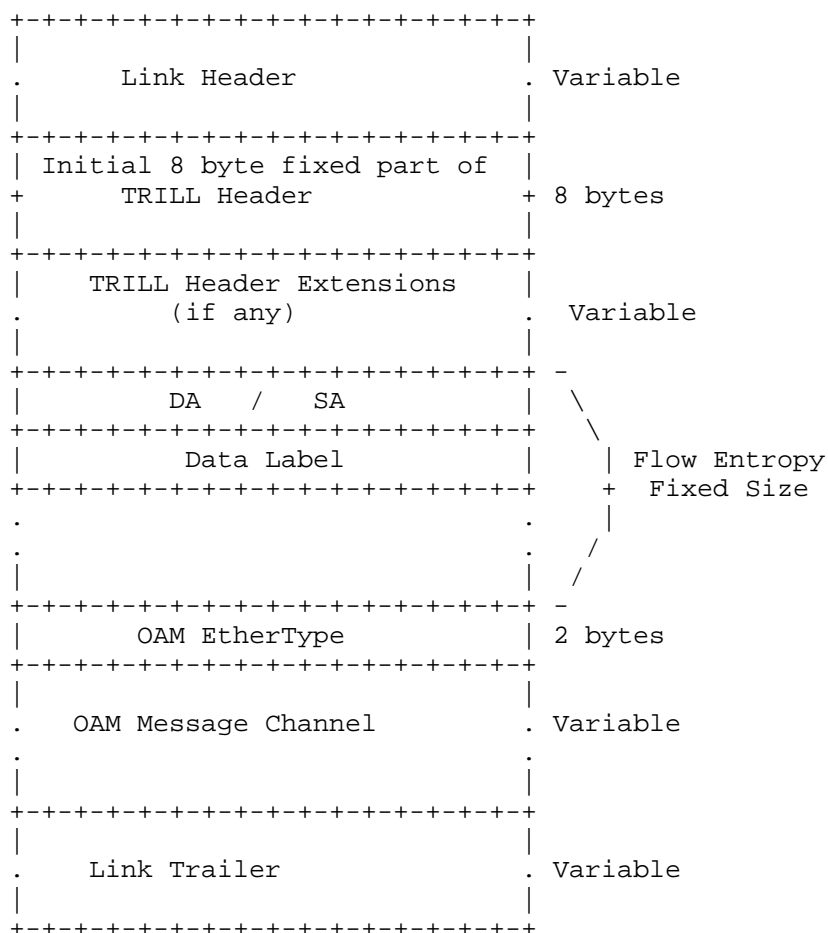


Figure 6 TRILL OAM Encapsulation

The OAM Message Channel used in this document is defined in [TRILL-FM], and has the following structure:

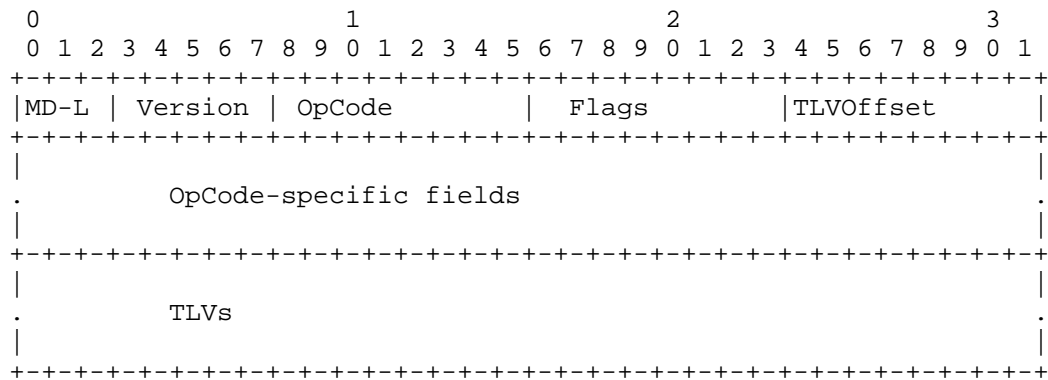


Figure 7 OAM Packet Format

The first 4 octets of the OAM Message Channel are common to all OpCodes, whereas the rest is OpCode-specific. Below is a brief summary of the fields in the first 4 octets:

- o MD-L : Maintenance Domain Level.
- o Version: indicates the version of this protocol. Always zero in the context of this document.
- o Flags: always zero in the context of this document.
- o FirstTLVOffset: defines the location of the first TLV, in octets, starting from the end of the FirstTLVOffset field.

For further details about the OAM packet format, see [TRILL-FM].

6.2. Loss Measurement Packet Formats

6.2.1. Counter Format

Loss Measurement packets use a 32-bit packet counter field. When a counter is incremented beyond its maximal value, 0xFFFFFFFF, it wraps around back to 0.

6.2.2. 1SL Packet Format

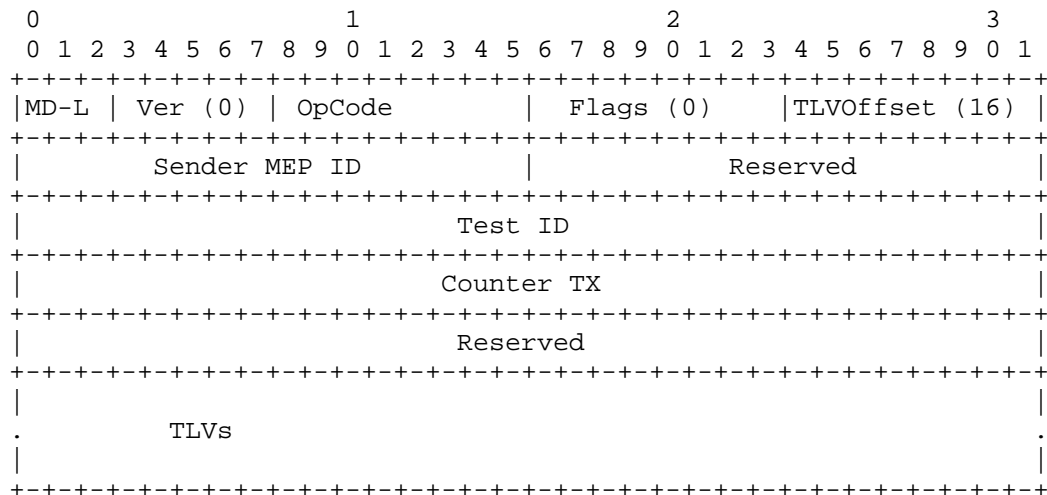


Figure 8 1SL Packet Format

- o Sender MEP ID: the MEP ID of the MEP that initiated the 1SL.
- o Reserved: always 0.
- o Test ID: a 32-bit unique test identifier.
- o Counter TX: the value of the sender's transmission counter, including this packet, at the time of transmission.

6.2.3. SLM Packet Format

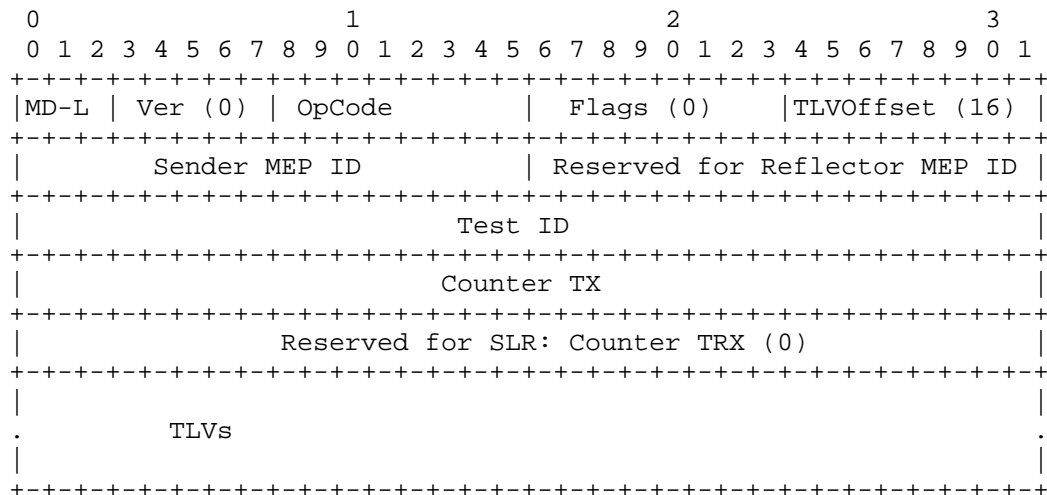


Figure 9 SLM Packet Format

- o Sender MEP ID: the MEP ID of the MEP that initiated this packet.
- o Reserved: this field is reserved for the reflector's MEP ID, to be added in the SLR.
- o Test ID: a 32-bit unique test identifier.
- o Counter TX: the value of the sender's transmission counter, including this packet, at the time of transmission.
- o Reserved: this field is reserved for the SLR corresponding to this packet. The reflector uses this field in the SLR for carrying TRX, the value of its reception counter.

6.2.4. SLR Packet Format

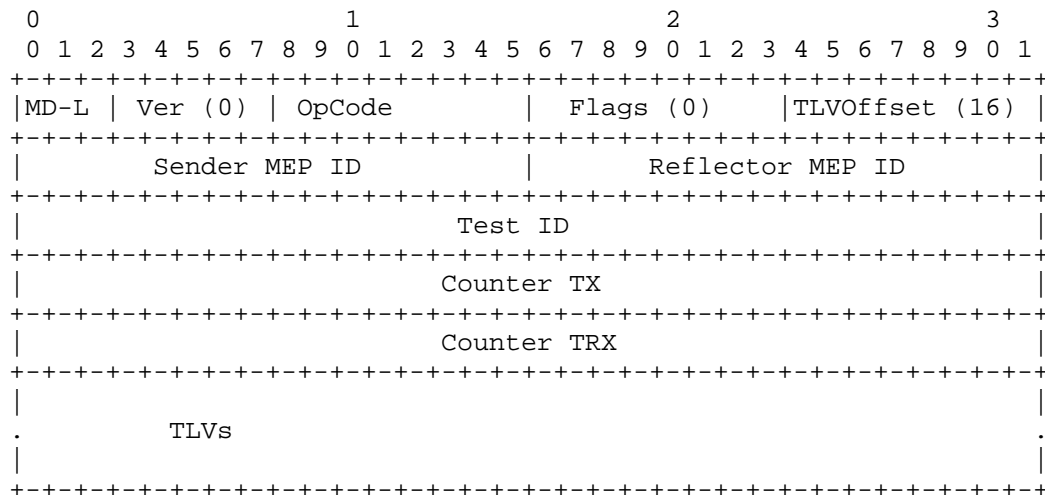


Figure 10 SLR Packet Format

- o Sender MEP ID: the MEP ID of the MEP that initiated the SLM that this SLR replies to.
- o Reflector MEP ID: the MEP ID of the MEP that transmits this SLR message.
- o Test ID: a 32-bit unique test identifier, copied from the corresponding SLM message.
- o Counter TX: the value of the sender's transmission counter at the time of the SLM transmission.
- o Counter TRX: the value of the reflector's reception counter, including this packet, at the time of reception of the corresponding SLM packet.

6.3. Delay Measurement Packet Formats

6.3.1. Timestamp Format

The timestamps used in Delay Measurement packets are 64 bits long. These timestamps use the 64 least significant bits of the IEEE 1588-2008 (1588v2) Precision Time Protocol timestamp format [IEEE1588].

This truncated format consists of a 32-bit seconds field followed by a 32-bit nanoseconds field. This truncated format is also used in IEEE 1588v1, in [Y.1731], and in [MPLS-LM-DM].

6.3.2. 1DM Packet Format

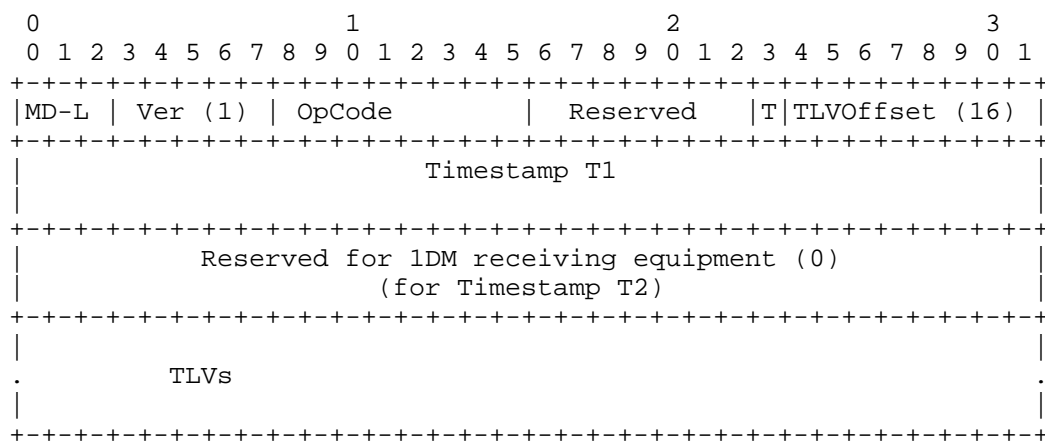


Figure 11 1DM Packet Format

- o T: Type flag. When this flag is set it indicates proactive operation, and when cleared it indicates on-demand mode.
- o Timestamp T1: specifies the time of transmission of this packet.
- o Reserved: this field is reserved for internal usage of the 1DM receiver. The receiver can use this field for carrying T2, the time of reception of this packet.

6.3.3. DMM Packet Format

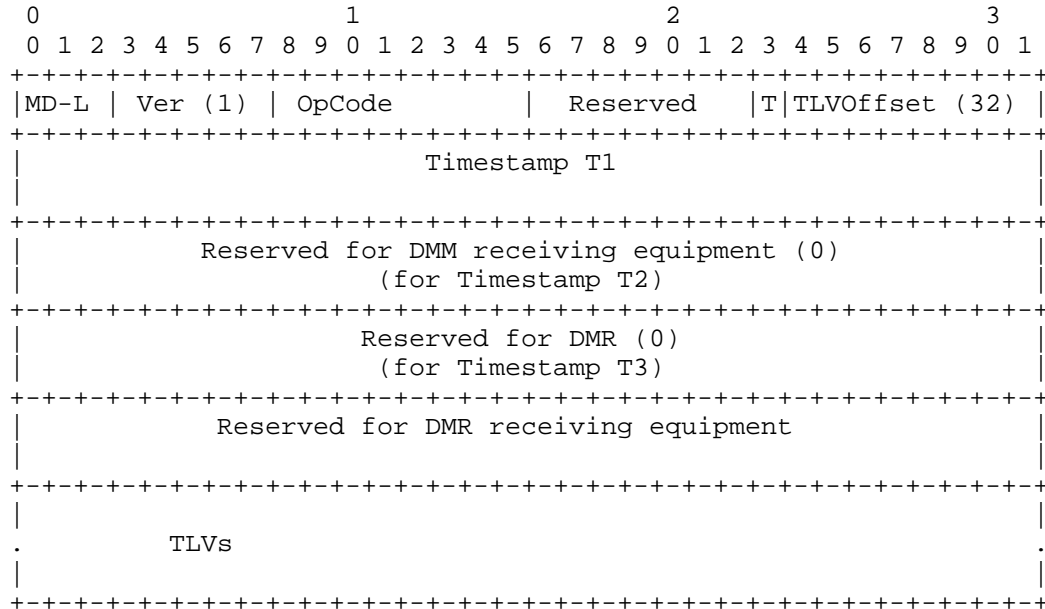


Figure 12 DMM Packet Format

- o T: Type flag. When this flag is set it indicates proactive operation, and when cleared it indicates on-demand mode.
- o Timestamp T1: specifies the time of transmission of this packet.
- o Reserved: this field is reserved for internal usage of the MEP that receives the DMM (the reflector). The reflector can use this field for carrying T2, the time of reception of this packet.
- o Reserved for DMR: two timestamp fields are reserved for the DMR message. One timestamp field is reserved for T3, the DMR transmission time, and the other field is reserved for internal usage of the MEP that receives the DMR.

6.3.4. DMR Packet Format

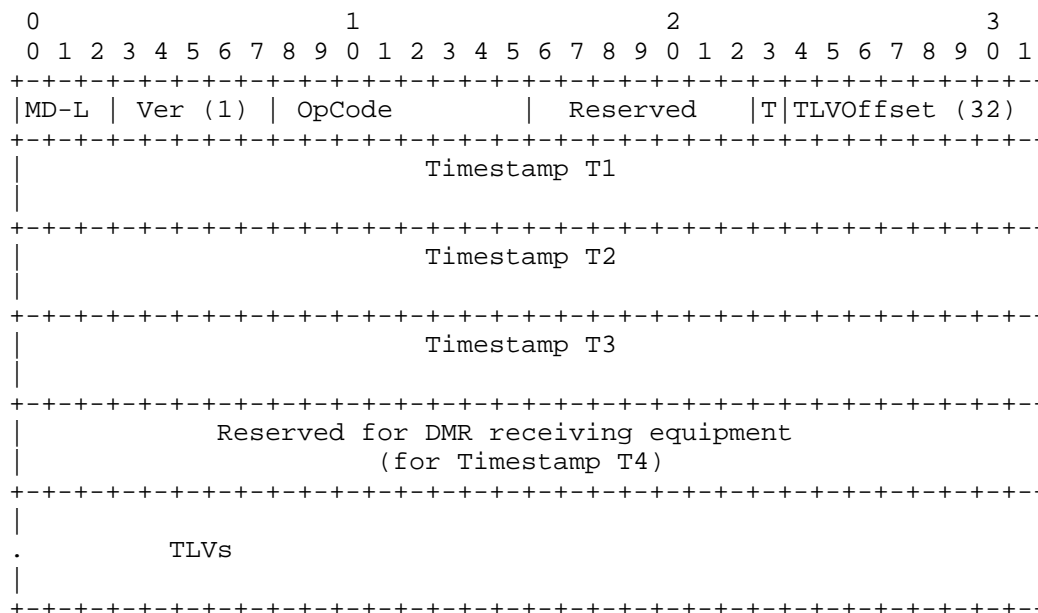


Figure 13 DMR Packet Format

- o T: Type flag. When this flag is set it indicates proactive operation, and when cleared it indicates on-demand mode.
- o Timestamp T1: specifies the time of transmission of the DMM packet that this DMR replies to.
- o Timestamp T2: specifies the time of reception of the DMM packet that this DMR replies to.
- o Timestamp T3: specifies the time of transmission of this DMR packet.
- o Reserved: this field is reserved for internal usage of the MEP that receives the DMR (the sender). The sender can use this field for carrying T4, the time of reception of this packet.

7. Performance Monitoring Process

The Performance Monitoring process is made up of a number of Performance Monitoring instances, known as PM Sessions. A PM session can be initiated between two MEPs on a specific flow and be defined as either a Loss Measurement session or Delay Measurement session.

The Loss Measurement session can be used to determine the performance metrics Frame Loss Ratio, availability, and resiliency. The Delay Measurement session can be used to determine the performance metrics Frame Delay, Inter-Frame Delay Variation, Frame Delay Range, and Mean Frame Delay.

The PM session is defined by the specific PM function (PM tool) being run, and also by the Start Time, Stop time, Message Period, Measurement Interval, and Repetition Time. These terms are defined as follows:

- o The Start Time is the time that the PM session begins.
- o The Stop Time is the time that the measurement ends.
- o The Message Period is the message transmission frequency (the time between message transmissions).
- o The Measurement Interval is the time period over which measurements are gathered and then summarized. The Measurement Interval can align with the PM Session duration, but it doesn't need to. PM messages are only transmitted during a PM Session.
- o The Repetition Time is the time between start times of the Measurement Intervals.

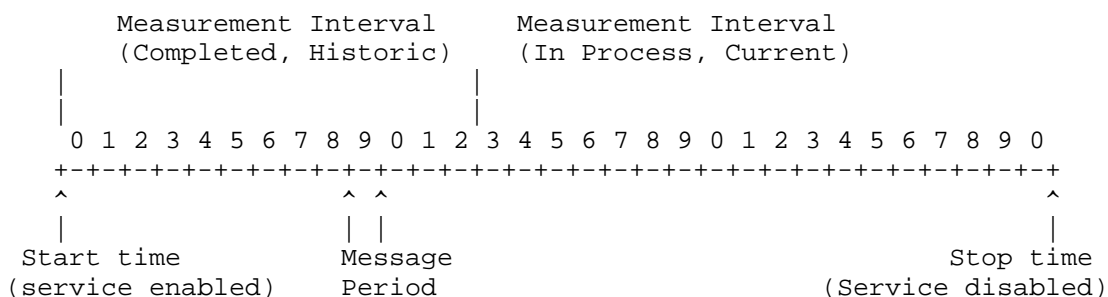


Figure 14 Relationship Between Different Timing Parameters

8. Security Considerations

The security considerations of TRILL OAM are discussed in [OAM-REQ] and in [OAM-FRAMEWK] and in [TRILL-FM]. General TRILL security considerations are discussed in [RFCTRILL]. This document does not inflict further security considerations.

9. IANA Considerations

9.1. OpCode Values

IANA is requested to assign TRILL OAM OpCode values to the packet types defined in this document. The suggested OpCode values are identical to the ones defined in [Y.1731]:

TBD-45 : 1DM

TBD-46 : DMR

TBD-47 : DMM

TBD-53 : 1SL

TBD-54 : SLR

TBD-55 : SLM

10. Acknowledgments

This document was prepared using 2-Word-v2.0.template.dot.

11. References

11.1. Normative References

- [KEYWORDS] Bradner, S., "Key words for use in RFCs to Indicate Requirement Levels", BCP 14, RFC 2119, March 1997.
- [RFCTRILL] Perlman, R., Eastlake, D., Dutt, D., Gai, S., Ghanwani, A., "Routing Bridges (RBridges): Base Protocol Specification", RFC 6325, July 2011.
- [OAM-FRAMEWK] Salam, S., Senevirathne, T., Aldrin, S., Eastlake, D., "TRILL OAM Framework", draft-ietf-trill-oam-framework (work in progress), September 2013.
- [TRILL-FM] Senevirathne, T., Finn, N., Salam, S., Kumar, D., Eastlake, D., Aldrin, S., Li, Y., "TRILL Fault Management", draft-ietf-trill-oam-fm (work in progress), July 2013.

11.2. Informative References

- [OAM-REQ] Senevirathne, T., Bond, D., Aldrin, S., Li, Y., Watve, R., "Requirements for Operations, Administration and Maintenance (OAM) in Transparent Interconnection of Lots of Links (TRILL)", RFC 6905, March 2013.
- [Y.1731] ITU-T Recommendation G.8013/Y.1731, "OAM Functions and Mechanisms for Ethernet-based Networks", July 2011.
- [802.1Q] "IEEE Standard for Local and metropolitan area networks - Media Access Control (MAC) Bridges and Virtual Bridged Local Area Networks", IEEE Std 802.1Q(tm), 2012 Edition, October 2012.
- [IEEE1588] IEEE TC 9 Instrumentation and Measurement Society, "1588 IEEE Standard for a Precision Clock Synchronization Protocol for Networked Measurement and Control Systems Version 2", IEEE Standard, 2008.
- [MPLS-LM-DM] Frost, D., Bryant, S., "Packet Loss and Delay Measurement for MPLS Networks", RFC 6374, September 2011.
- [OAM] Andersson, L., Van Helvoort, H., Bonica, R., Romascanu, D., Mansfield, S., "Guidelines for the use of the OAM acronym in the IETF ", RFC 6291, June 2011.

Authors' Addresses

Tal Mizrahi
Marvell
6 Hamada St.
Yokneam, 20692 Israel

Email: talmi@marvell.com

Tissa Senevirathne
Cisco
375 East Tasman Drive
San Jose, CA 95134, USA

Email: tsenevir@cisco.com

Samer Salam
Cisco
595 Burrard Street, Suite 2123
Vancouver, BC V7X 1J1, Canada

Email: ssalam@cisco.com

Deepak Kumar
Cisco
510 McCarthy Blvd,
Milpitas, CA 95035, USA

Phone : +1 408-853-9760
Email: dekumar@cisco.com

Donald Eastlake 3rd
Huawei USA R&D
155 Beaver Street
Milford, MA 01757 USA
Phone: +1-508-333-2270
Email: d3e3e3@gmail.com

