

# Opportunistic Encryption revisited

(We're getting the cypherpunks band back together)

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## Terminology was discussed yesterday at perpass

- Opportunistic Encryption means many different things
- I use the old FreeS/WAN definition
- Encrypt between two endpoints without specific setup
- Can be "anonymous" (no authentication)

# draft-wouters-dane-openpgp-01

Publish OpenPGP key in DNSSEC to allow mail clients, MUA and MTA's to encrypt on the file.

```
echo "python -c 'import base64; print base64.b32encode("paul")'" \
  "_openpgpkey.nohats.ca IN TYPE65280 \#\ \
  (" 'gpg --export --export-options export-minimal \
  paul@nohats.ca | base64 | wc -c' ; gpg --export \
  --export-options export-minimal paul@nohats.ca | \
  base64; echo ")"
```

```
dig +dnssec -t TYPE65280 obqk3a=._openpgpkey.nohats.ca
```

TODO: Write postfix/sendmail milter proof of concept

## draft-wouters-edns-tcp-chain-query and draft-wouters-edns-tcp-keepalive

Both drafts are meant to speed up DNSSEC on high latency links  
(read: phone)

- Improve client - server communication to keep TCP 53 open
- Get all DNSSEC data for validation of IPSECKEY record in one round trip

## old FreeS/WAN OE

- 1 Startup needs to confirm its own identify and public key (often fails)
- 2 Startup of FreeS/WAN causes 30 seconds of DNS misses and packetloss
- 3 Application sends packet to remote host (eg www.nohats.ca)
- 4 Kernel intercepts packet, sends to IKE daemon
- 5 IKE daemon tries to find IPSECKEY/TXT record in reverse DNS
- 6 (meanwhile application retries initial packet, or fails loudly)
- 7 IKE daemon sets up IPsec tunnel
- 8 Application, if not given up, send packet through tunnel
- 9 (LOTS of 'failures' to remember in SPD/SAD on both sides)

## The FreeS/WAN era problems and mistakes

- Ahead of its time (technically, politically)
- Only mutual authenticated IPsec - enduser must publish IPSECKEY (too hard)
- Key distribution via reverse DNS(SEC) which has been abandoned
- Hoped IPv6 would obsolete NATs
- Intercept packet, then find identity
- The common "no OE" fallback to plaintext took too long
- (later, supported initiator-OE and NAT-T)

# What has changed?

- Pervasive monitoring - anonymous IPsec better than plaintext
- Deployed DNSSEC - possible to put validator on every device
- Devices powerful enough to do lots of IPsec and DNSSEC

## What has NOT changed?

- Users are still not able to configure IPSECKEY in DNS
- Still no IPv6 host-to-host, but NATed client-to-server
- Reverse DNS still unusable

## Opportunistic Encryption with IPsec

- 1 Application sends DNS request for A record (eg `www.nohats.ca`) to local DNSSEC resolver
- 2 DNS server attempts to find A record as well as IPSECKEY record for `www.nohats.ca`
- 3 If IPSECKEY record found, send IP address, FQDN and IPSECKEY to IKE daemon
- 4 IKE daemon sets up IPsec tunnel using (internet wide) PSK for its own authentication, RSA for remote auth.
- 5 DNS server returns A record to application
- 6 Application send data, automatically goes over IPsec tunnel

## Anonymous OE with IPsec

- 1 Application sends DNS request for A record (eg www.nohats.ca) to local DNSSEC resolver
- 2 DNS server attempts to find A record as well as IPSECKEY record for www.nohats.ca
- 3 No IPSECKEY record found, send IP address, FQDN to IKE daemon
- 4 IKE daemon sends blind IKE attempt to remote IP with universal PSK. Works of fails.
- 5 DNS server returns A record to application
- 6 Application send data, automatically goes over IPsec tunnel if other end supported IKE.

## Comparison with BTNS

- public keys / CERT exchanged inline only - no server authentication
- Connection latching, channel binding, upgrades, IKE policies, new SAD/SPD flags
- Requires async DNS lookups (or other async auths) in IKE daemon
- Does not use asymmetrical IKEv2 authentication (PSK  $\neq$  RSA) thus privacy leak
- Needs to generate ephemeral RSA keys which is bad for low entropy embedded gateways
- IKE daemon becomes complex - needs modification instead of configuration ( $\approx$  3 lines of code)

## Implementation details

Planned for Libreswan 3.7, tentatively in Fedora 21 and enabled per default in Fedora 22

- Deal with overlapping NAT using server-side NAT to 127.\* addresses
- If both have published IPSECKEY, how to handle role reversal when starting as anonymous
- Reluctantly accept client-server versus host-host
- Combine OE with "static configurations" by simply using multiple policies (no modification)
- On Linux/BSD set OE IPsec "priority" field to be always lower than non-OE IPsec
- (What to do with IKEv1 old OE code - remove?)
- Keep kinds of OE connection to a minimum - no leap of faith IPsec