TSO, fair queuing, pacing: three's a charm

Eric Dumazet <edumazet@google.com>
Yuchung Cheng <ycheng@google.com>

A world of TCP bursts

- TCP is window-based and ack-clocked
 - Sudden cwnd increase
 - cwnd: stretch ACKs (e.g., LRO)
 - rwnd: receiver buffering
 - Idling between ACK and data
- TSO deferral bugs
- Switches local aggregation
- Modern structured traffic

Burst losses are bad signals to CC as network is often not 100% utilized

Oversize bursts are bad

- Bad throughput
 - Losses and ECN flags
 - Network is often not 100% utilized
 - False congestion alarms to TCP
- Vicious cycle
 - Bursts losses recovery rwin surge bursts ...

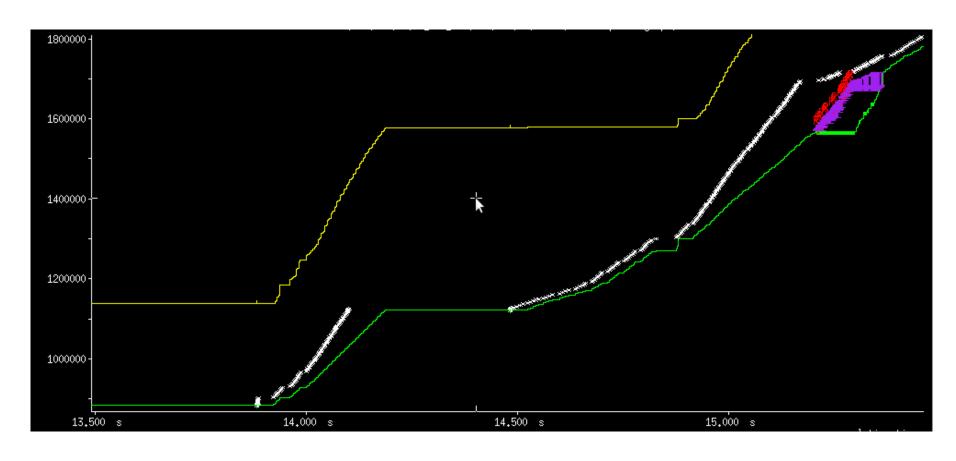
Many solutions exist

- Limit the bursts
 - moderate cwnd = inflight + 3 (linux)
 - send <=max_burst pkts per ACK (bsd)
- Tweak cwnd when idle
 - Reduce cwnd by Y% after X time
 - e.g., cwnd = min(cwnd, IW) after RTO
- Disable TSO/LRO
- Throttle large flows to improve fairness in qdisc

fq/pacing

- TCP cwnd controls the amount to send per RTT
 - Clocked by ACKs
 - No more TCP tweaks for burst
- qdisc layer handles sub-RTT scheduling
 - Pace at cwnd/RTT after idling
 - Flow fair queuing to improve mixing and fairness
- Break large burst into microburst
 - Dynamic TSO sizing base on the pacing rate
 - High performance host I/O

fq/pacing to reduce video burst



fq/pacing in Linux 3.11-rc7

- High performance: allows millions of concurrent flows per Qdisc
 - Small memory footprint: 8K per Qdisc, and 104 bytes per flow
 - Single high resolution timer to pace flows
 - One RB tree to link throttled flows.
 - fast flow match (not stochastic hash like SFQ/FQ codel)
- Uses the new_flow/old_flow separation from FQ_codel
- Special FIFO queue for high prio packets (no need for PRIO + FQ)

Example usage:

tc qdisc add dev \$ETH root fq

