Transport Architectures for an Evolving Internet

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March 5, 2014



Joint work with Anirudh Sivaraman, Pratiksha Thaker, and Hari Balakrishnan.

The Internet evolves

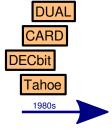
In 20 years, computer networks have seen dramatic change:

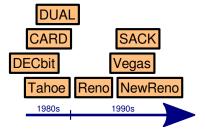
- Wi-Fi
- Cellular networks
- Datacenters
- ▶ 10 GigE
- Transoceanic links
- Ubiquitous mobility
- Huge amounts of streaming video

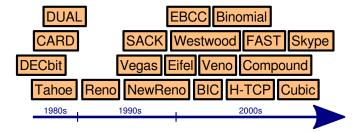
Coping with change

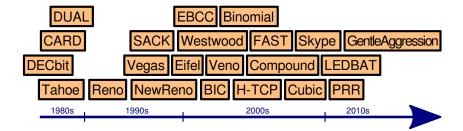
How should users deal with an evolving network?

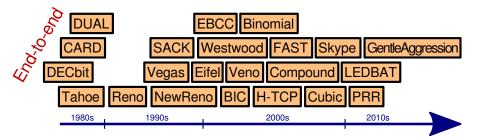
One approach: design new protocols.

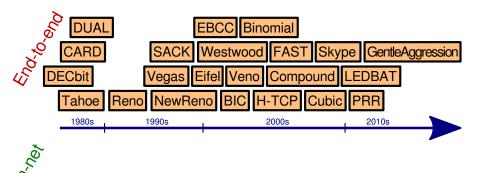


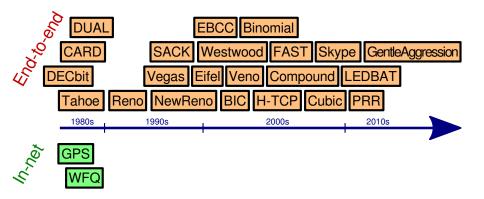


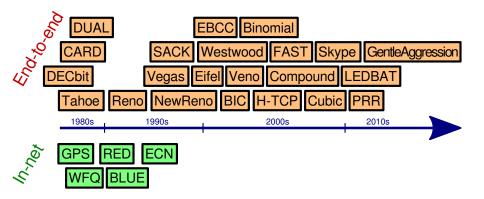


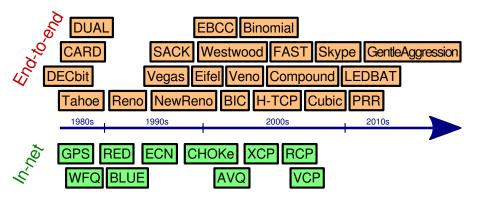


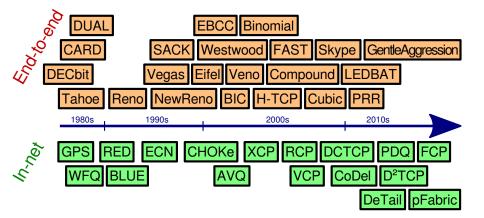












Declarative design

Systems with a **model** and a **mission**.

Explicitness in systems design

Model: explicit statement of assumptions about the problem

Mission: objective that the application wants

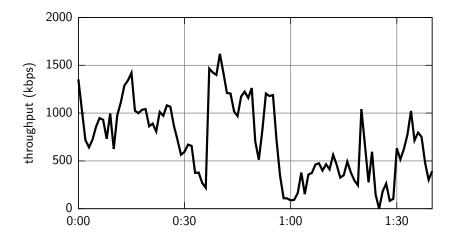
Explicit design considerations \rightarrow **freedom to make changes**

Sprout: a transport protocol designed for variability

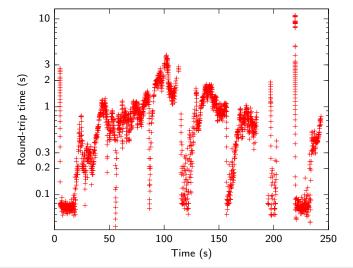
Observation:

Videoconferences perform poorly over cellular networks.

Verizon LTE uplink throughput



Verizon LTE ping delay during one TCP download

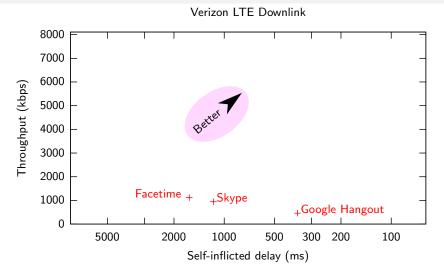


Interactive apps work poorly

- We measured cellular networks while driving:
 - Verizon LTE
 - ► Verizon 3G (1xEV-DO)
 - ► AT&T LTE
 - ► T-Mobile 3G (UMTS)
- Then ran apps across replayed network trace:
 - Skype (Windows 7)
 - Google Hangouts (Chrome on Windows 7)
 - Apple Facetime (OS X)

Skype's performance

Performance summary



What's wrong?

- Existing schemes react to congestion signals.
 - ▶ Packet loss.
 - Increase in round-trip time.
- Feedback comes too late.
- ► The killer: **self-inflicted queueing delay**.

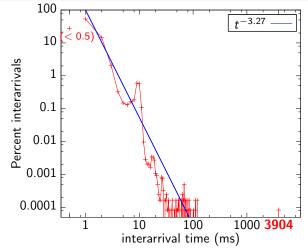
Sprout's mission

- Most throughput
- ightharpoonup Bounded risk of delay $> 100~{
 m ms}$

Bounded risk of delay

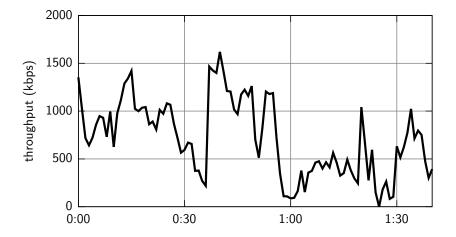
- Model variation in link speed
- ► Infer current link speed
- Predict future link speed
 - Don't wait for congestion
- ► **Control:** Send as much as possible, but require:
 - ▶ 95% chance all packets arrive within 100 ms

Model: packet deliveries looks like flicker noise

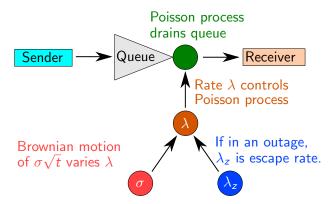


(Verizon LTE, phone stationary.)

Model: average rate looks like random walk



Sprout's model



Sprout's model parameters

Volatility σ : fixed 0 200 $\frac{\text{pkts/s}}{\sqrt{s}}$ Expected outage time $1/\lambda_z$: 1 s

Expected outage time $1/\lambda_z$: 1 s Tick length (τ) : 20 ms

Forecast length: 160 ms

Delay target: 100 ms

Risk tolerance: 5%

All source code was frozen before data collection began.

Infer: current link speed

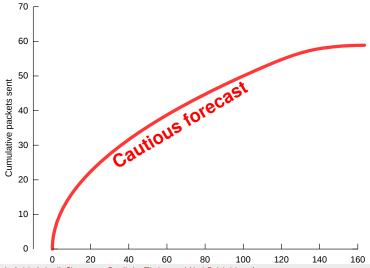
▶ Observe packets received every τ

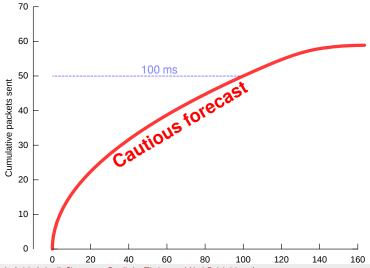
▶ **Update** $P(\lambda)$

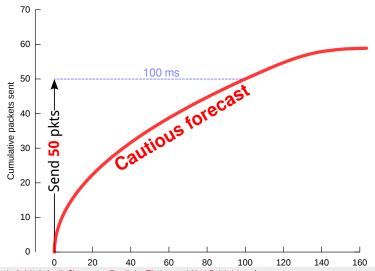
Predict: future link speed

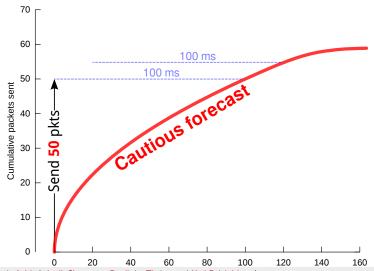
Evolve model forward

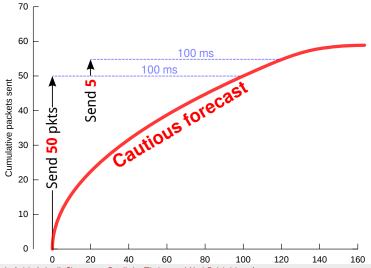
Forecast 5th percentile cumulative packets

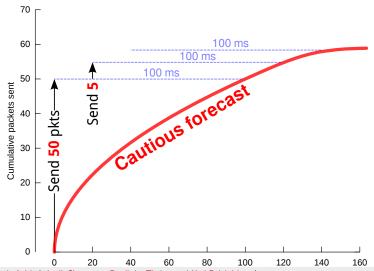




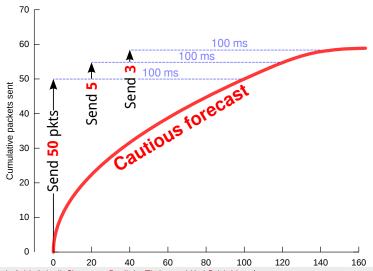






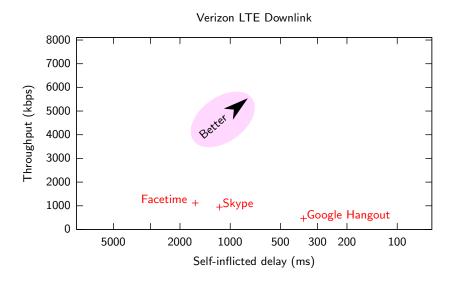


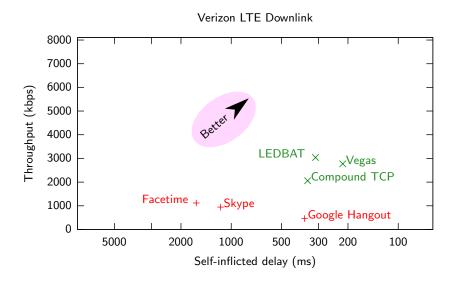
Control: fill up 100 ms forecast window

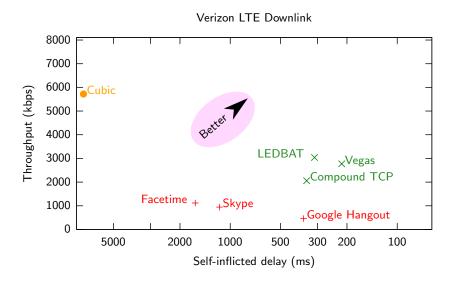


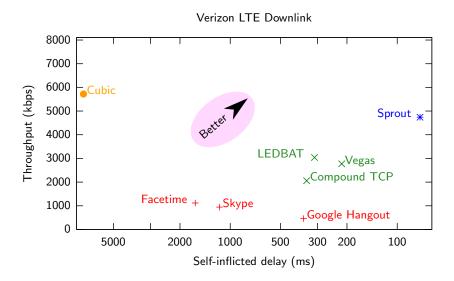
Keith Winstein (with Anirudh Sivaraman, Pratiksha Thaker, and Hari Balakrishnan)

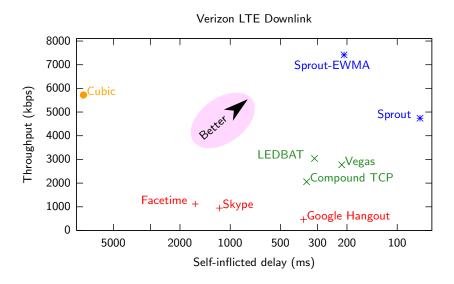
Sprout's results

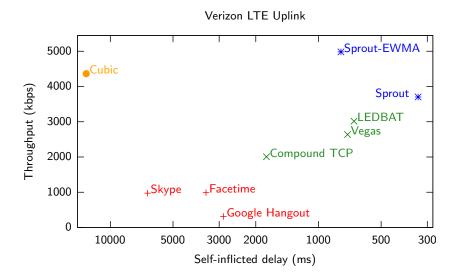










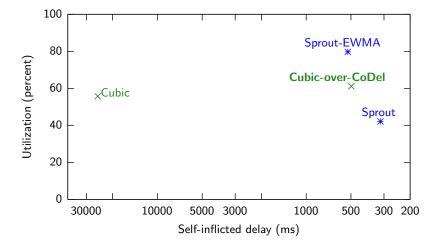


Overall results on 8 links

Verizon 3G/LTE, AT&T LTE, T-Mobile 3G uplink and downlink:

Sprout vs.	Avg. speedup	Delay reduction
Skype	2.2×	7.9×
Hangout	4.4×	7.2×
Facetime	1.9×	8.7×
Compound	1.3×	4.8×
TCP Vegas	1.1 imes	2.1×
LEDBAT	Same	2.8×
Cubic	0.91×	79×

Sprout is end-to-end, but comparable to in-net control



M.I.T. 6.829 contest (March-April 2013)

- Turnkey network emulator, evaluation
- Sender, receiver run in Linux containers
- Mission: maximize throughput/delay
- 4th prize: \$20
- 3rd prize: \$30
- 2nd prize: \$40
- ► (If beat Sprout) 1st prize:

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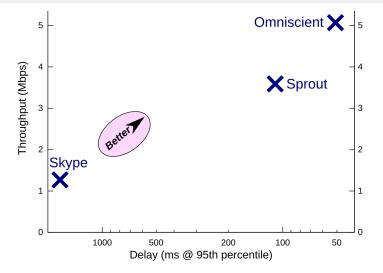
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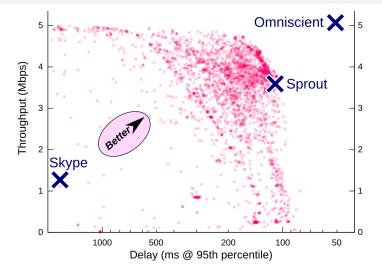
Anirudh Sivaraman, KW, Pauline Varley, Somak Das, Joshua Ma, Ameesh Goyal, João Batalha, and Hari Balakrishnan, **Protocol Design Contests**, *in submission*

Baseline

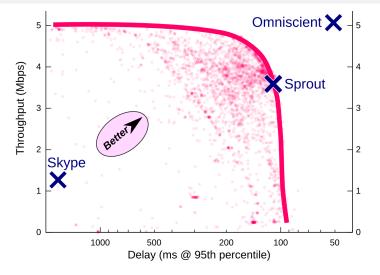


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Land of 3,000 student protocols



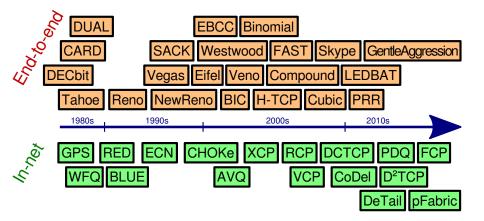
Sprout was on the frontier



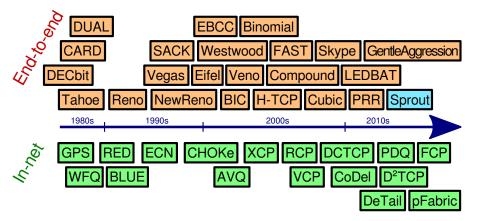
Limitations

- Sprout wants to control all of the traffic on a queue.
 - Cells generally have per-user queues...
 - ...but Wi-Fi and wired networks usually don't.
- ▶ What if cell link *isn't* the bottleneck?
- Assumption: application always has data to send

Sprout's mark



Sprout's mark



Now that we have 40+ algorithms...

- Sprout for cellular networks?
- Wireless-TCP for Wi-Fi?
- High-BDP-TCP for transoceanic links?
- Datacenter-TCP for datacenters?
- CoDel for cable modems?
- ▶ TBA-TCP for tomorrow's networks?

Rational choice of scheme is challenging



- Different missions?
- Different assumptions about network?
- One scheme just plain better?

Networks constrained by a fuzzy idea of TCP's assumptions

- Mask stochastic loss
- Bufferbloat
- Mask out-of-order delivery
- No parallel/multipath routing

Advice for Internet Subnetwork Designers (RFC 3819) is 21,000 words!

- Open lots of flows
- Goose slow start
- Add pacing
- Give up and do it yourself

Open lots of flows











- Goose slow start
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- Open lots of flows
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Google MICR@SOFT



- Open lots of flows
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- Give up and do it yourself









Google MICR@SOFT

You Tube

Chrome (QUIC) BitTorrent (μ TP) Mosh (SSP) IBM Aspera (fasp)

Idea: computer-generated protocols

Transport layer should adapt to whatever:

- network does
- application wants

Idea: computer-generated protocols

Transport layer should adapt to **whatever**:

- network does (model)
- application wants (mission)

What we built

Remy: a program that generates congestion-control schemes offline

Input:

- Assumptions about network and workload (model)
- Application's objective (mission)

Output: CC algorithm for a TCP sender (RemyCC)

Time: hours to days

The basic question of congestion control

At this moment, do I:

- send a packet
- not send a packet?

Maximize

 $\sum_{i} \log [\text{throughput}_{i}] \quad (\text{proportionally fair throughput})$

Maximize

 $\sum_{i} \log [\text{throughput}_{i}] \quad (\text{proportionally fair throughput})$

$$\sum_{i} \log \left[\frac{\text{throughput}_{i}}{\text{delay}_{i}} \right]$$
 (proportionally fair throughput/delay)

Maximize

 $\sum_{i} \log [\text{throughput}_{i}] \quad (\text{proportionally fair throughput})$

$$\sum_{i} \log \left[\frac{\text{throughput}_{i}}{\left(\text{delay}_{i}\right)^{\delta}} \right] \text{ (proportionally fair throughput/delay)}$$

Maximize

- $\sum_{i} \log [\text{throughput}_{i}] \quad (\text{proportionally fair throughput})$
- $\sum_{i} \log \left[\frac{\mathsf{throughput}_{i}}{\left(\mathsf{delay}_{i}\right)^{\delta}} \right] \text{ (proportionally fair throughput/delay)}$
- min_i throughput_i (max-min throughput)

Minimize

- mean flow completion time
- page load time

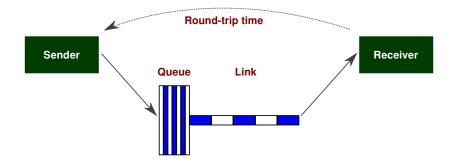
Prevent

- pathological behavior
- congestion collapse

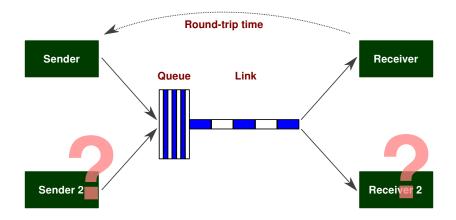
Encoding the designer's prior assumptions

- Model of network uncertainty
 - Link speed distribution
 - Delay distribution
 - Topology distribution
- Model of workload
 - Web browsing
 - MapReduce
 - videoconferencing
 - streaming video (YouTube/Netflix)

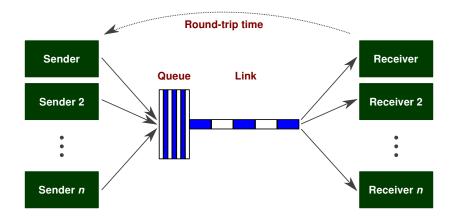
Dumbbell network



Dumbbell network



Dumbbell network



Superrational congestion control

At this moment,* do I:

- ▶ send a packet
- not send a packet?

Superrational congestion control

At this moment,* do I:

- send a packet
- not send a packet?

* Assuming every node is running the same algorithm.

Remy: tractable search for best policy

- Best decision given all history: not tractable
- Instead, summarize the history

A RemyCC tracks four congestion signals

- $r_{-ewma_{\alpha}}$: **short-term** moving average of interval between acks "How fast are packets arriving (now)?"
- $r_{-ewma_{\beta}}$: **long-term** moving average of same "How fast are packets arriving (smoothed)?"
 - s_ewma: moving average of interval between acked timestamps "How fast was I sending?"
- rtt_ratio: ratio of last RTT to smallest RTT so far
 "How long is the queue?"

Why these four features?

We can measure the benefit of each!

- Removing any one hurts
 - ▶ losing $r_{-ewma_{\alpha}}$ hurts the most

- More signals increase search time
- ... but others might help on some networks

A RemyCC maps each state to an action

$$\operatorname{REMYCC}(r_ewma_{\alpha\beta}, s_ewma, rtt_ratio) \rightarrow \langle m, b, \tau \rangle$$

- *m* Multiple to congestion window
- b Increment to congestion window
- au Minimum interval between two outgoing packets

Runtime for a RemyCC

On ack:

- $ightharpoonup \langle m, b, \tau \rangle \leftarrow \text{RemyCC}(r_\text{ewma}_{\alpha\beta}, s_\text{ewma}, rtt_ratio)$
- ▶ cwnd $\leftarrow m \cdot \text{cwnd} + b$

Send packet if:

- cwnd > FlightSize, and
- ▶ last packet sent $> \tau$ ago

Remy's job

Find piecewise-continuous REMYCC() that optimizes expected value of objective function

Remy example: 2D state space

On ack:

```
\langle m, b, \tau \rangle \leftarrow \text{RemyCC}(s\_ewma, r\_ewma_{\alpha}, r\_ewma_{\beta}, rtt\_ratio)
```

Remy example: 2D state space

On ack:

$$\langle m, b, \tau \rangle \leftarrow \text{RemyCC}(s_ewma, r_ewma_{\alpha},$$

Remy example: model

Quantity	Distribution	Units
Link speed	Uniform(10, 20)	Mbps
RTT	Uniform(100, 200)	ms
n	Uniform(1, 16)	
"On" process "Off" process	$\exp[\mu=5]$ same	seconds

Remy example: mission

$$\sum_{i} \log \left[\frac{\mathsf{throughput}_{i}}{\mathsf{delay}_{i}} \right]$$

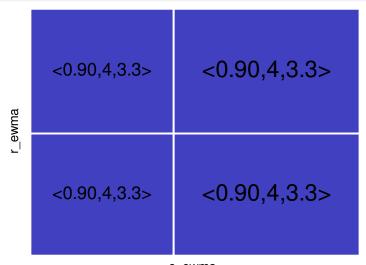
One action for all states. Find the best value.

r_ewma <?,?,?>

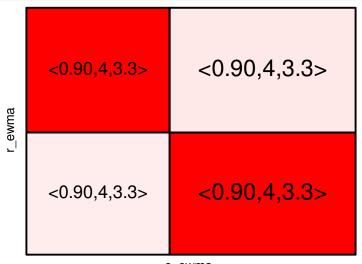
The best (single) action. Now split it on median.

r_ewma <0.90,4,3.3>

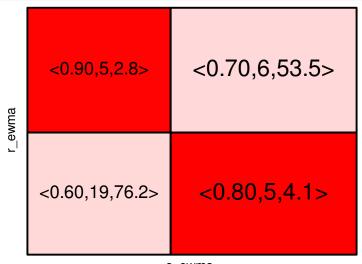
Simulate



Optimize each of the new actions



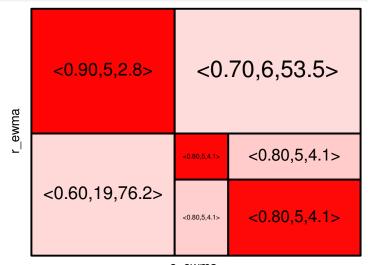
Now split the most-used rule



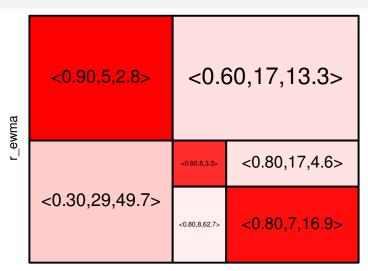
Simulate



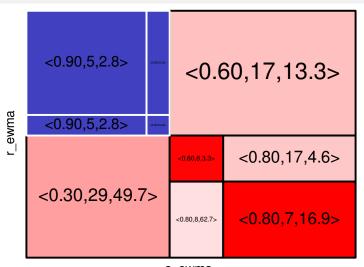
Optimize



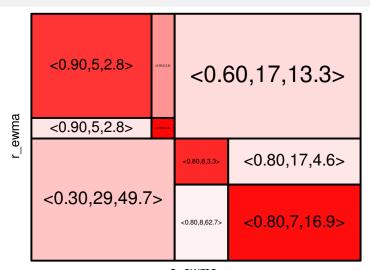
Split



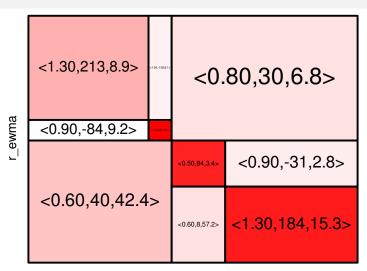
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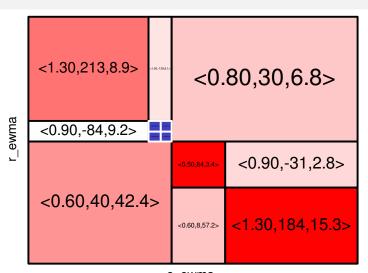
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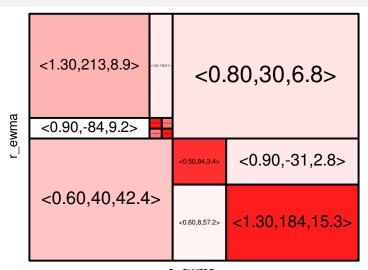
Split



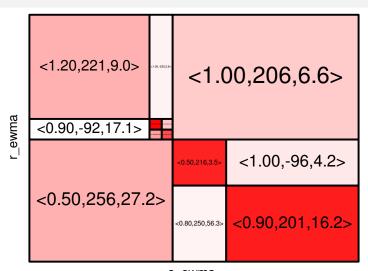
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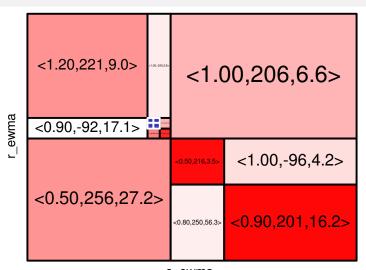
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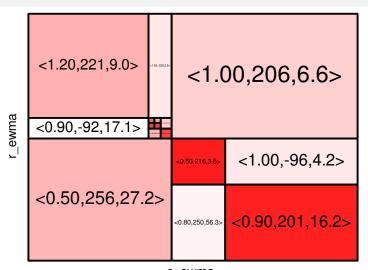
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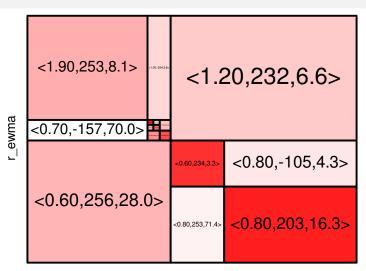
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Optimize



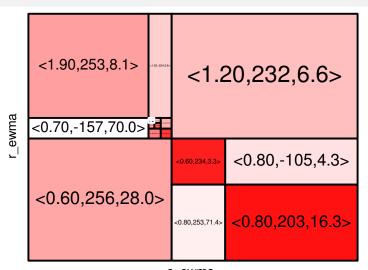
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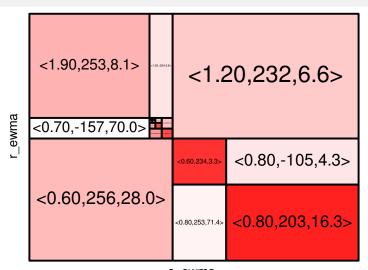
s_ewma

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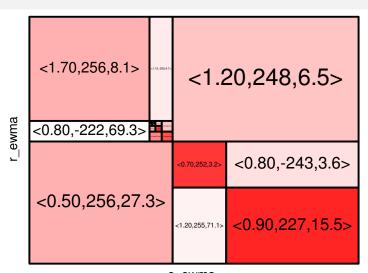
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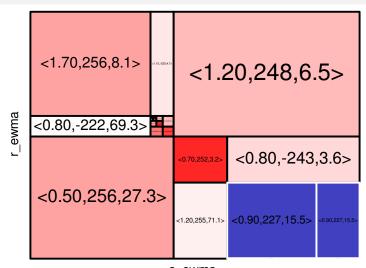
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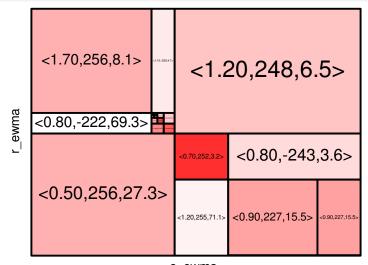


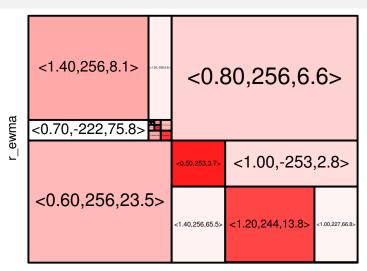
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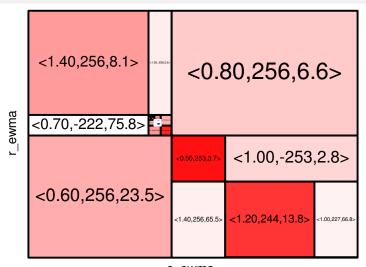


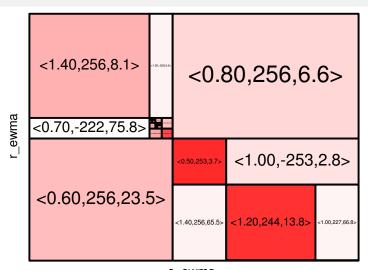
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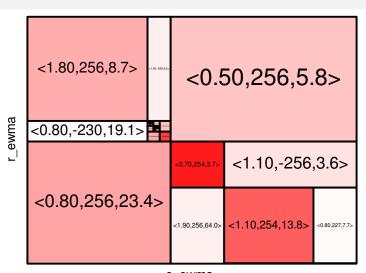


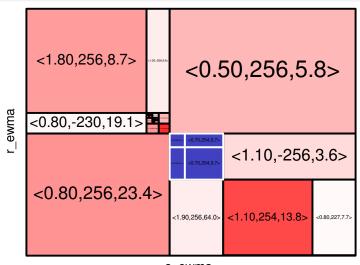


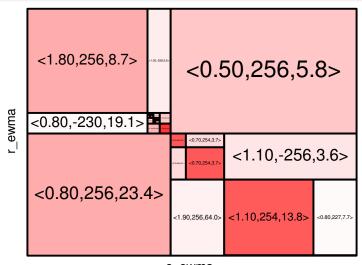


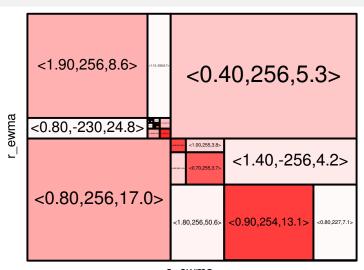


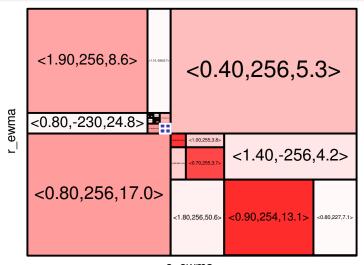


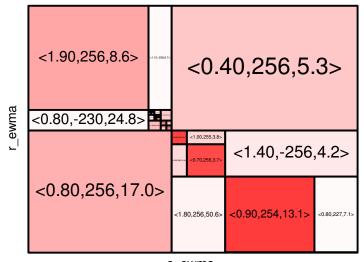


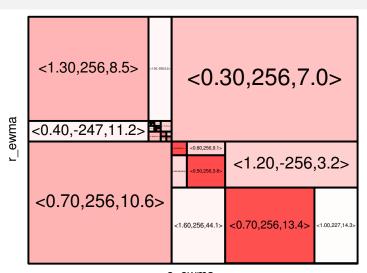


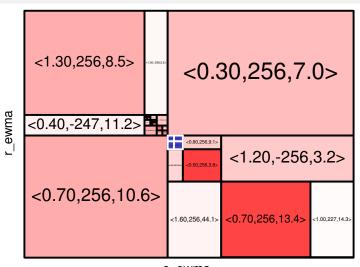


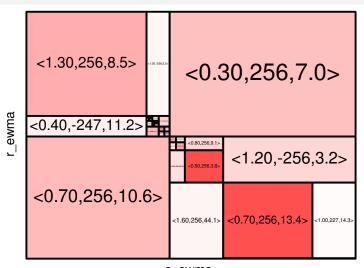


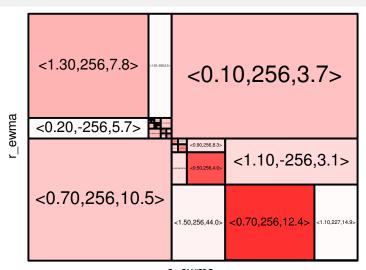


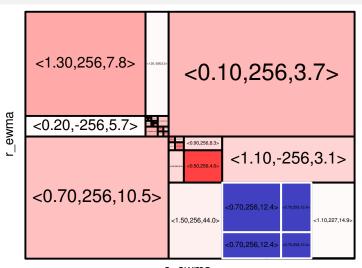


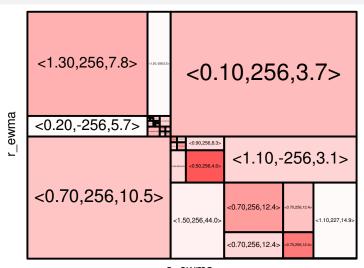


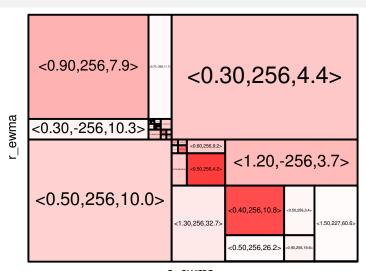


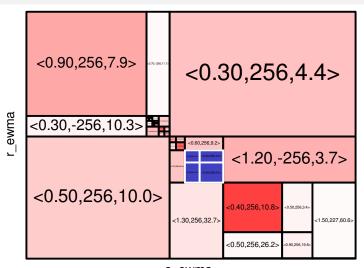


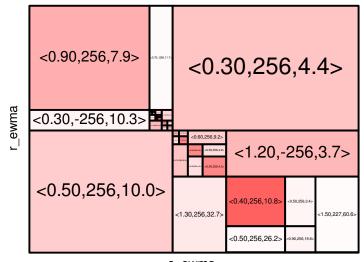


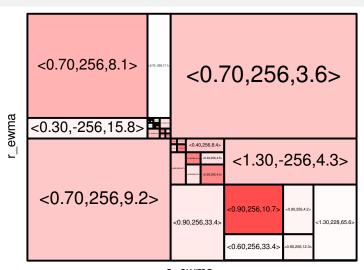


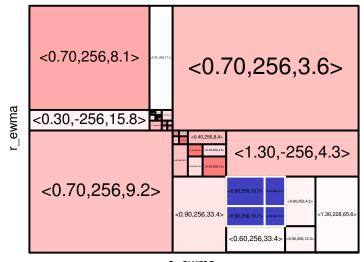






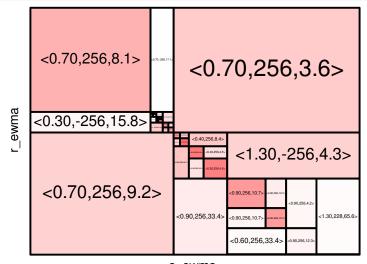


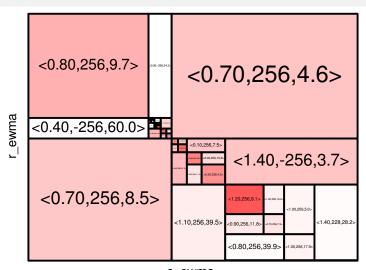


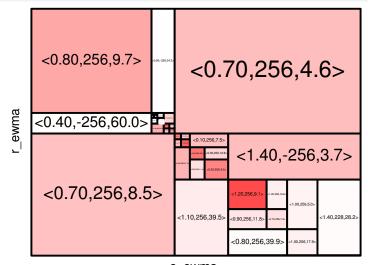


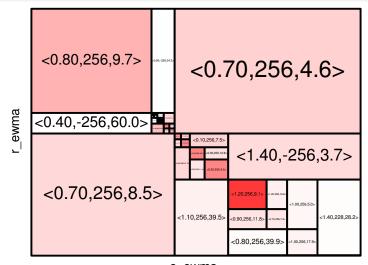
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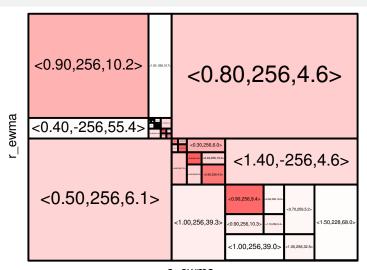
Keith Winstein (with Anirudh Sivaraman, Pratiksha Thaker, and Hari Balakrishnan)

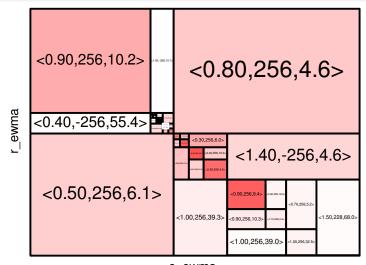


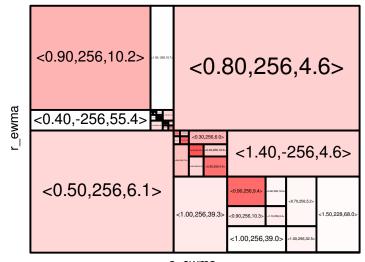


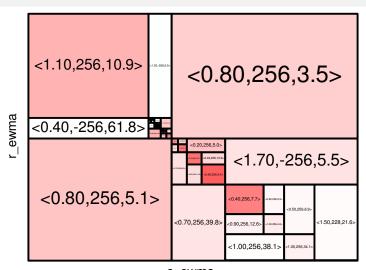


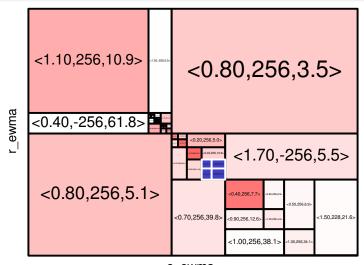


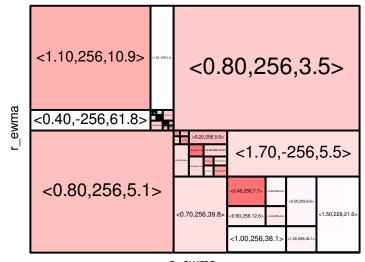


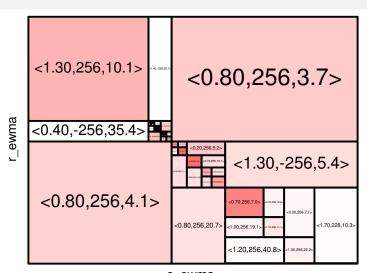


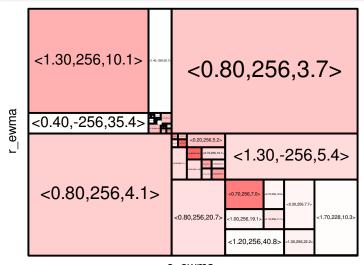


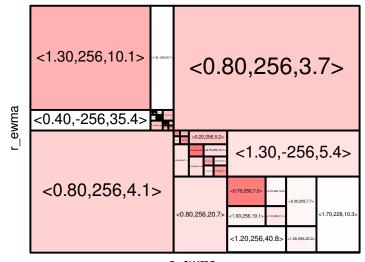


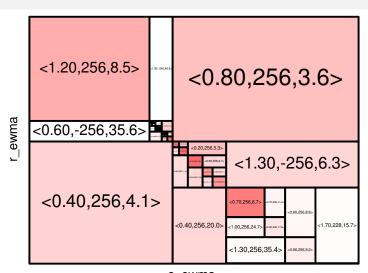


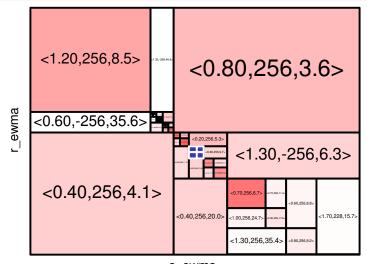




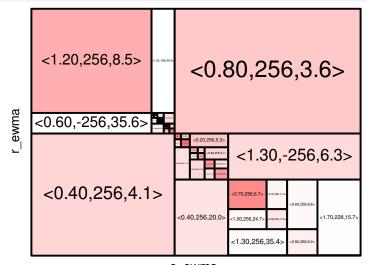




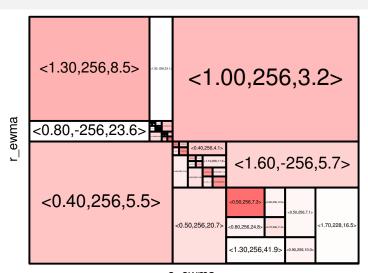




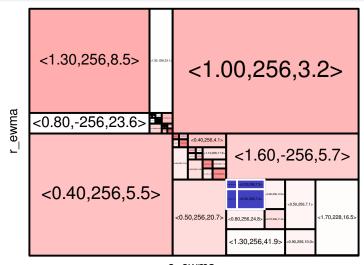
Optimize



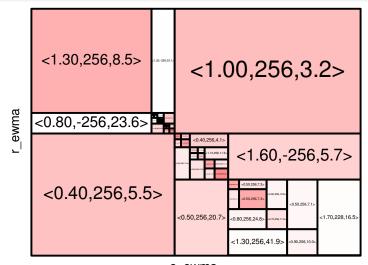
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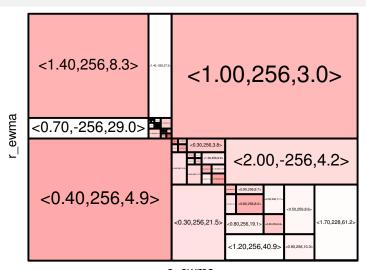
Simulate



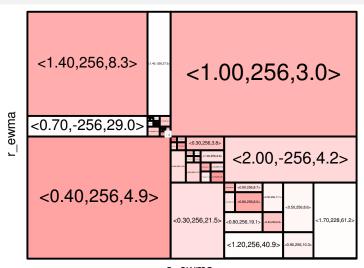
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Split



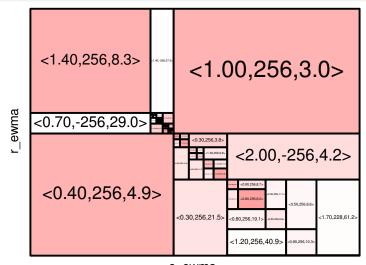
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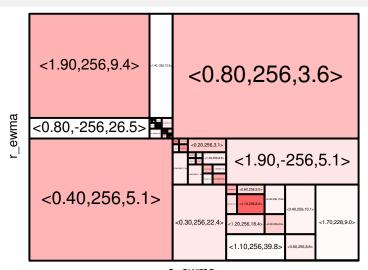
s_ewma

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Optimize



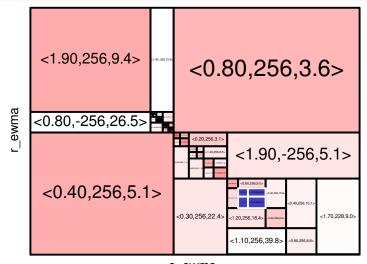
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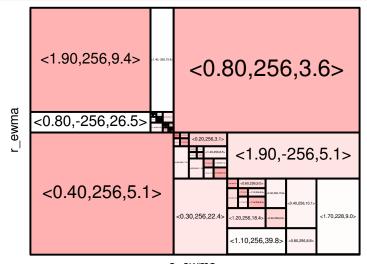
s_ewma

Keith Winstein (with Anirudh Sivaraman, Pratiksha Thaker, and Hari Balakrishnan)

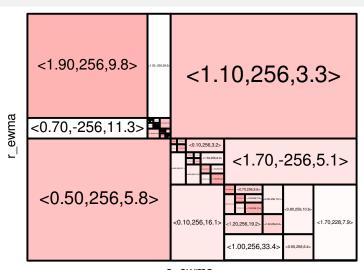
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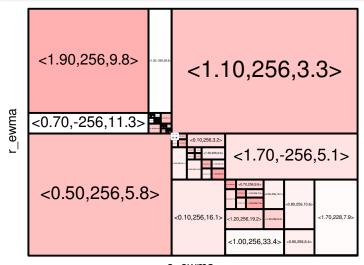
Optimize



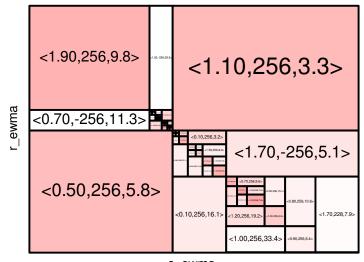
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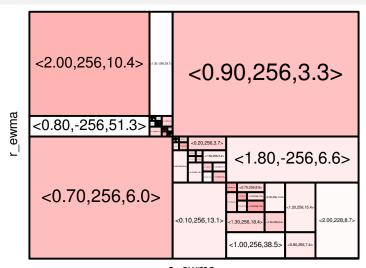
Simulate



Optimize



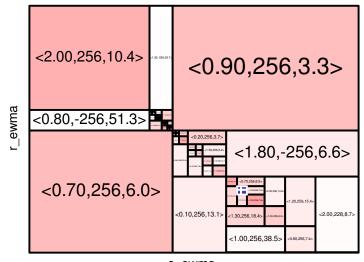
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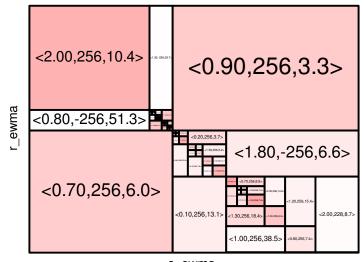
s_ewma

Keith Winstein (with Anirudh Sivaraman, Pratiksha Thaker, and Hari Balakrishnan)

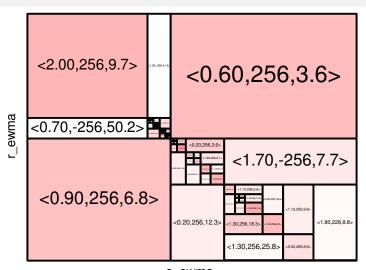
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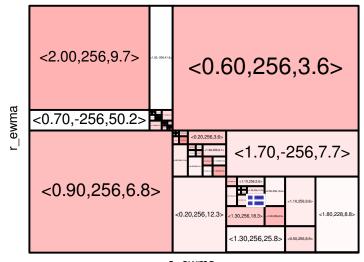
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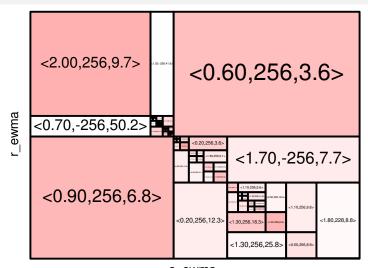


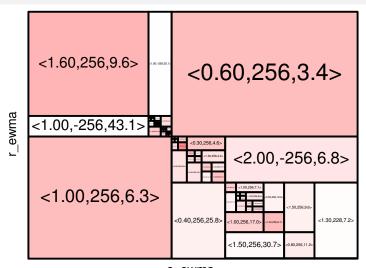
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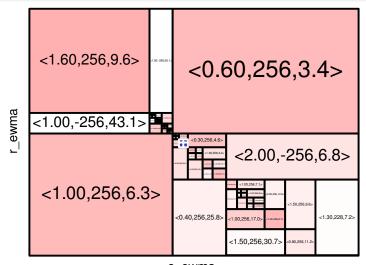


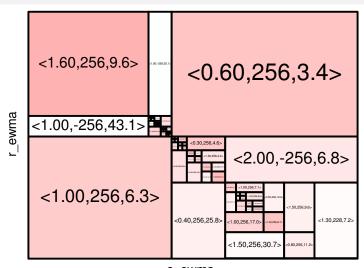
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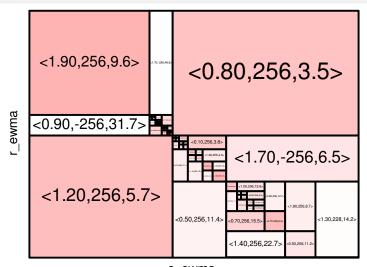


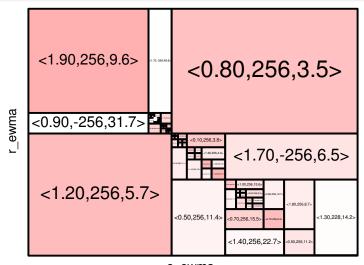


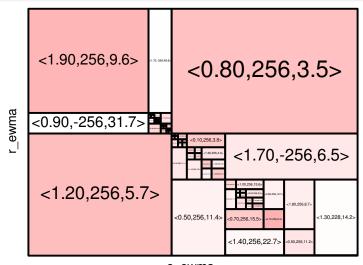


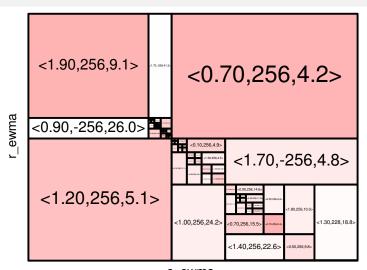


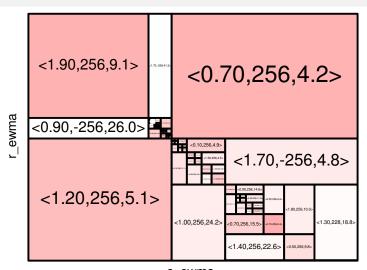


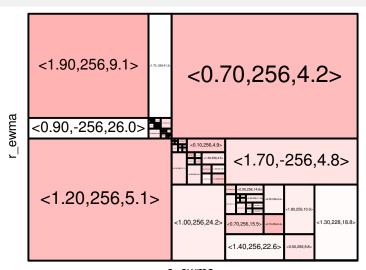


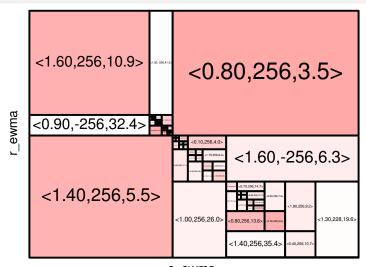


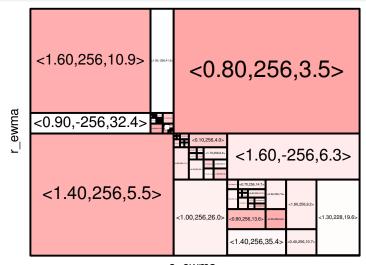


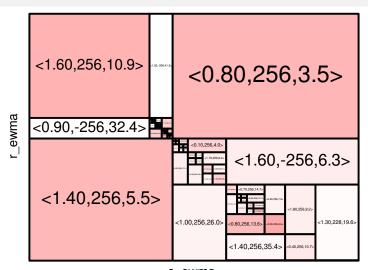


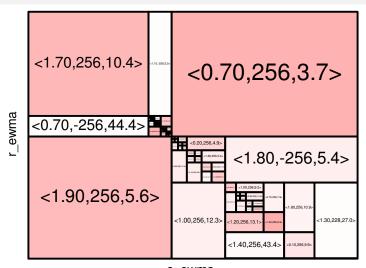










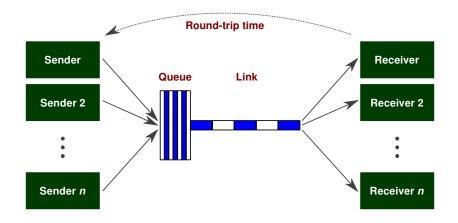


Evaluation in ns-2

- ► End-to-end comparators: NewReno, Cubic, Compound, Vegas
- In-net comparators: Cubic-over-sfqCoDel, XCP
- Simulation setup published for replication



Scenario 1: fixed-rate network, homogenous senders



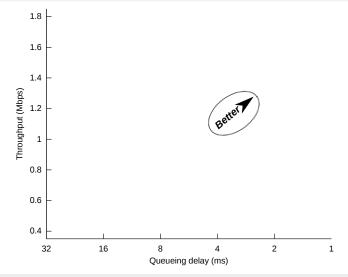
Scenario 1: details

Quantity	Simulation parameter	Remy assumptions
Link speed	15 Mbps	Uniform(10, 20) Mbps
RTT	150 ms	Uniform(100, 200) ms
n	8	Uniform $(1, 16)$
"On" process	$\exp[\mu=100] extsf{kB}$	$\exp[\mu=5]\mathbf{s}$
"Off" process	$\exp\left[\mu = \frac{1}{2}\right]$ s	$\exp[\mu={f 5}]{\sf s}$

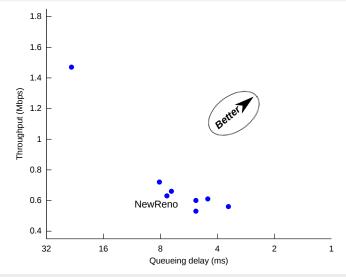
Remy objective:

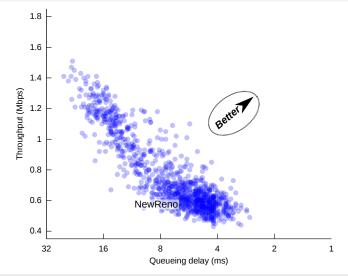
$$\sum_{i} \log \left[rac{\mathsf{throughput}_{i}}{\left(\mathsf{delay}_{i}
ight)^{\delta}}
ight]$$
 $oldsymbol{\delta} \in \left\{ rac{1}{10}, 1, 10
ight\}$

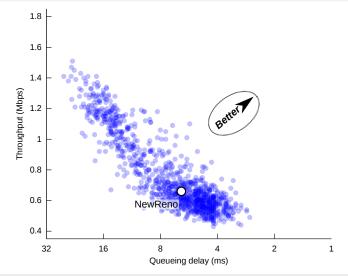
Scenario 1: throughput-delay plot

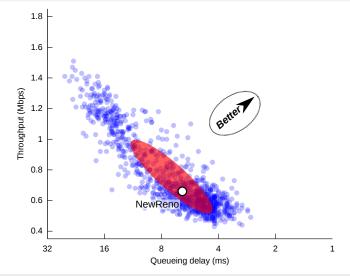


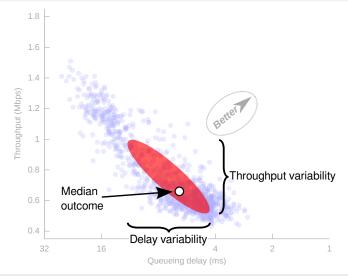
Keith Winstein (with Anirudh Sivaraman, Pratiksha Thaker, and Hari Balakrishnan)

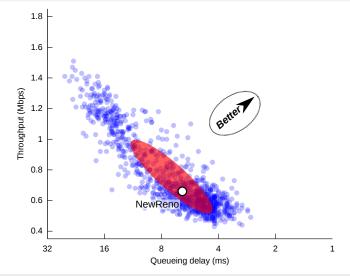


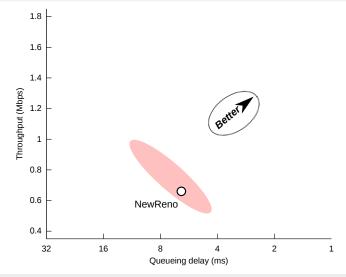


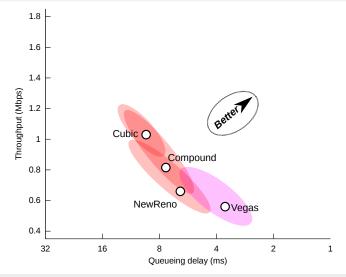


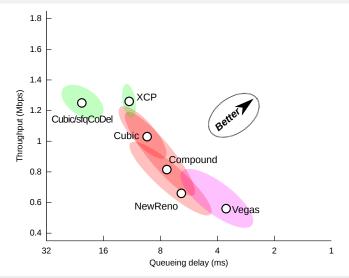


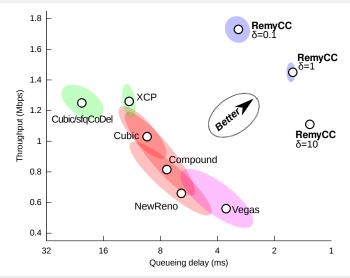


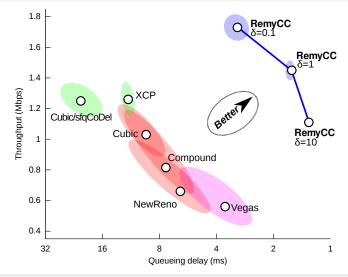




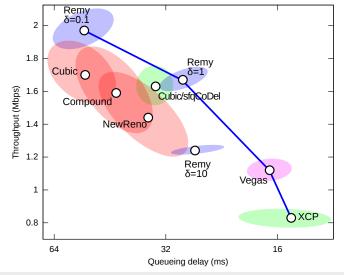








Scenario 2: Verizon LTE, n = 8

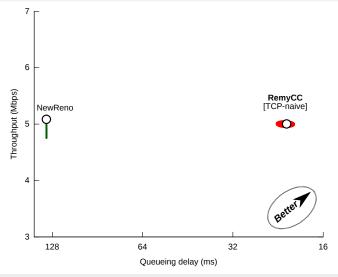


Remy as an instrument to study network science

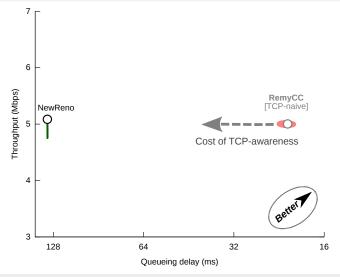
From the perspective of an endpoint, what does it help to know about the network?

How difficult is it to learn a good protocol, given an **imperfect** model of the network?

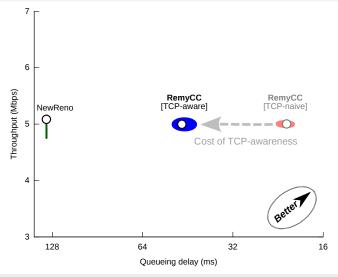
RemyCC competing against itself



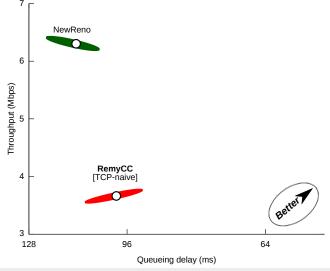
RemyCC competing against itself



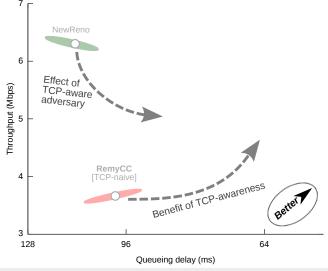
RemyCC competing against itself



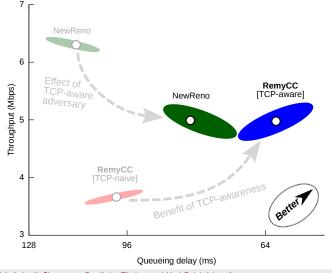
RemyCC competing against TCP NewReno

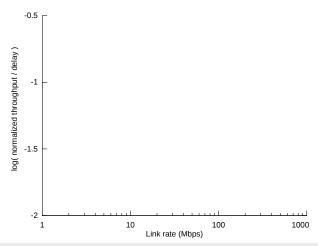


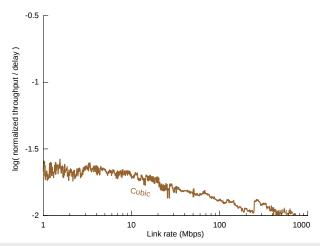
RemyCC competing against TCP NewReno

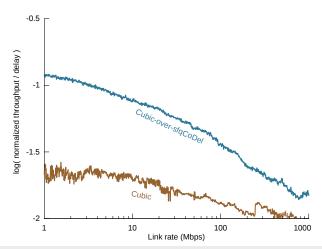


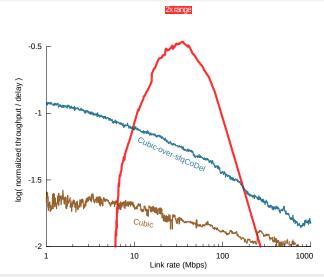
RemyCC competing against TCP NewReno

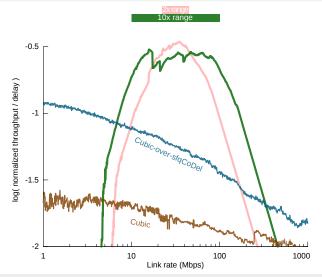


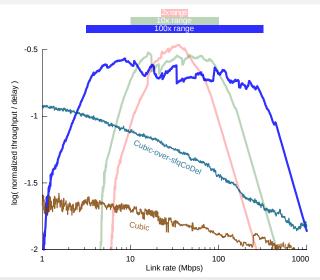


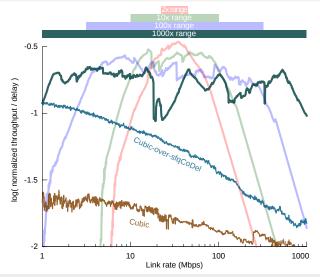












Systems ex Machina

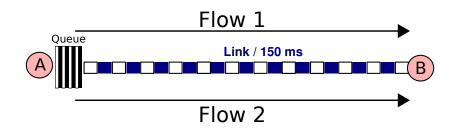
- ► Explicit design considerations → freedom to make changes
- "If this system is the answer, what's the question?"

Sprout $2\text{--}4\times$ the throughput and $7\text{--}9\times$ less delay than Skype, etc. Remy computer-generated protocol design

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http://mit.edu/keithw

One bottleneck



Two bottlenecks

