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PCEP Extensions for MPLS-TE LSP Automatic Bandwidth Adjustment with
stateful PCE
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Abstract

The Path Computation Element Communication Protocol (PCEP) provides mechanisms for Path Computation Elements (PCEs) to perform path computations in response to Path Computation Clients (PCCs) requests. The stateful PCE extensions provide stateful control of Multiprotocol Label Switching (MPLS) Traffic Engineering Label Switched Paths (TE LSP) via PCEP, for a model where the PCC delegates control over one or more locally configured LSPs to the PCE.

This document describes the automatic bandwidth adjustment of such LSPs under the Active Stateful PCE model.

Status of This Memo

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1. Introduction

[RFC5440] describes the Path Computation Element Protocol (PCEP) as the communication between a Path Computation Client (PCC) and a Path Control Element (PCE), or between PCE and PCE, enabling computation

of Multiprotocol Label Switching (MPLS) for Traffic Engineering Label Switched Path (TE LSP).

[I-D.ietf-pce-stateful-pce] specifies extensions to PCEP to enable stateful control of MPLS TE LSPs. In this document focus is on Active Stateful PCE where LSPs are configured on the PCC and control over them is delegated to the PCE.

Over time, based on the varying traffic pattern, an LSP established with certain bandwidth may require to adjust the reserved bandwidth over time automatically. Ingress Label Switch Router (LSR) samples the traffic rate at each sample-interval (BwSample) to determine the traffic information as Maximum Average Bandwidth (MaxAvgBw). Further adjustment to the reserved bandwidth should be made at every adjustment-interval automatically.

Enabling Auto-Bandwidth on a LSP results in the LSP automatically adjusting its bandwidth based on the actual traffic flowing through the LSP. A LSP can therefore be setup with some arbitrary (or zero) bandwidth value such that the LSP automatically monitors the traffic flow and adjusts its bandwidth every adjustment-interval period. The bandwidth adjustment uses the make-before-break signaling method so that there is no interruption to traffic flow. This is described in detail in Section 4.1. [I-D.ietf-pce-stateful-pce-app] describes the usecase for auto-bandwidth adjustment for passive and active stateful PCE.

There are two approaches to automatic bandwidth adjustments in case of active stateful PCE -

- o PCE to decide adjusted bandwidth:
 - * Active stateful PCE can use other information such as historical trending data, application-specific information about expected demands and central policy information along with realtime actual flow volumes to make smarter bandwidth adjustment to delegated LSPs. Since LSP has delegated control to the PCE, it is inherently suited that it should be stateful PCE that decides the bandwidth adjustments. But this requires PCC to report the realtime bandwidth usage as well as the configuration knobs etc.
- o PCC to decide adjusted bandwidth:
 - * This approach would be similar to passive stateful PCE model, where the headend (PCC) monitor and calculate the new adjusted bandwidth and request the computed adjusted bandwidth to be updated. The passive stateful PCE would use path request/reply

mechanism where as in active stateful PCE report/update mechanism is used to adjust the bandwidth. This approach only require PCC to report the calculated bandwidth to be adjusted. But this approach does not utilize the optimization advantages offered by the active stateful PCE.

This document defines extensions needed to support Auto-Bandwidth feature along with mechanism to provide traffic information of the LSPs in a stateful PCE model using PCEP.

This document does not exclude use of any other mechanism employed by stateful PCE to learn real time traffic information etc. But at the same time, using the same protocol (PCEP in this case) for updating and reporting the LSP parameters as well as to support automatic bandwidth adjustment is operationally beneficial.

1.1. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

2. Terminology

The following terminology is used in this document.

Active Stateful PCE: PCE that uses tunnel state information learned from PCCs to optimize path computations. Additionally, it actively updates tunnel parameters in those PCCs that delegated control over their tunnels to the PCE.

Delegation: :An operation to grant a PCE temporary rights to modify a subset of tunnel parameters on one or more PCC's tunnels. Tunnels are delegated from a PCC to a PCE.

PCC: Path Computation Client: any client application requesting a path computation to be performed by a Path Computation Element.

PCE: Path Computation Element. An entity (component, application, or network node) that is capable of computing a network path or route based on a network graph and applying computational constraints.

TE LSP: Traffic Engineering Label Switched Path.

Note the additional terms defined in Section 4.1.

3. Motivation

An active stateful PCE can update the bandwidth for a delegated LSP via mechanisms described in [I-D.ietf-pce-stateful-pce]. Note that further extension are needed because of following reasons:

1. To identify the LSPs that would like to use this feature. Not all LSPs in some deployments would like their bandwidth to be dependent on the live traffic but be constant as set by the operator. In case of PCC initiated LSP, they would be configured at PCC and PCEP should support a mechanism to identify the LSP with auto bandwidth feature enabled at the PCE. Where as for PCE initiated LSP, PCEP should support mechanisms to request PCC to provide live traffic information.
2. Further for LSP with auto bandwidth feature enabled, operator should be able to specify the knobs to control this feature like the bandwidth-range etc and PCEP should support their encoding.
3. PCC would need to report the live traffic information using the same protocol (PCEP in this case) making the network operations easier.

Extensions as specified in this document is one of the way for PCE to learn this information. But at the same time a stateful PCE MAY choose to learn this information from other means like management, performance tools.

4. Architectural Overview

4.1. Auto-Bandwidth Overview

Auto-Bandwidth feature allows an LSP to automatically and dynamically adjust its reserved bandwidth over time, i.e. without network operator intervention. The bandwidth adjustment uses the make-before-break adaptive signaling method so that there is no interruption to traffic flow.

The new bandwidth reservation is determined by sampling the actual traffic flowing through the LSP. If the traffic flowing through the LSP is lower than the configured or current bandwidth of the LSP, the extra bandwidth is being reserved needlessly and being wasted. Conversely, if the actual traffic flowing through the LSP is higher than the configured or current bandwidth of the LSP, it can potentially cause congestion or packet loss. With Auto-Bandwidth feature, the LSP bandwidth can be set to some arbitrary value (even zero) during initial setup time, and it will be periodically adjusted over time based on the actual bandwidth requirement.

Note the following terms:

Maximum Average Bandwidth (MaxAvgBw): The maximum average bandwidth is the unit to measure the current traffic demand between a time interval. This is the maximum value of the averaged traffic pattern in a particular time interval.

Sample-Interval: The time interval in which the traffic rate is collected as a sample.

Adjustment-Interval: The time interval in which the bandwidth adjustment should be made based on the MaxAvgBw.

Minimum Bandwidth: The minimum bandwidth that should be reserved for the LSP.

Maximum Bandwidth: The maximum bandwidth that can be reserved for the LSP.

Report-Threshold: This value indicates when the current live traffic bandwidth sample (BwSample) must be reported to stateful PCE via PCRpt message. Only if the percentage difference between the current BwSample and the last BwSample is greater than or equal to the threshold percentage the LSP bandwidth is reported to PCE.

Adjust-Threshold: This value indicates when the bandwidth must be adjusted. Only if the percentage difference between the current MaxAvgBw and the current bandwidth allocation is greater than or equal to the threshold percentage the LSP bandwidth is adjusted to the current bandwidth demand.

4.2. Deploying Auto-Bandwidth Feature

The traffic rate is repeatedly sampled at each sample-interval (which can be configured by the user and the default value as 5 minutes). The sampled traffic rates are accumulated over the adjustment-interval period (which can be configured by the user and the default value as 24 hours).

The ingress LSR reports the live traffic information to the stateful PCE via the PCRpt message, to avoid multiple reports, the Report-Threshold percentage is used. Only if the percentage difference between the current BwSample and the last BwSample is greater than or equal to the threshold percentage the LSP bandwidth is reported to PCE.

Stateful PCE will adjust the bandwidth of the LSP to the highest sampled traffic rate amongst the set of samples taken over the

adjustment-interval. Note that the highest sampled traffic rate could be higher or lower than the current LSP bandwidth. Only if the current MaxAvgBw and the current bandwidth allocation is greater than or equal to the Adjust-Threshold percentage the LSP bandwidth is adjusted to the current bandwidth demand.

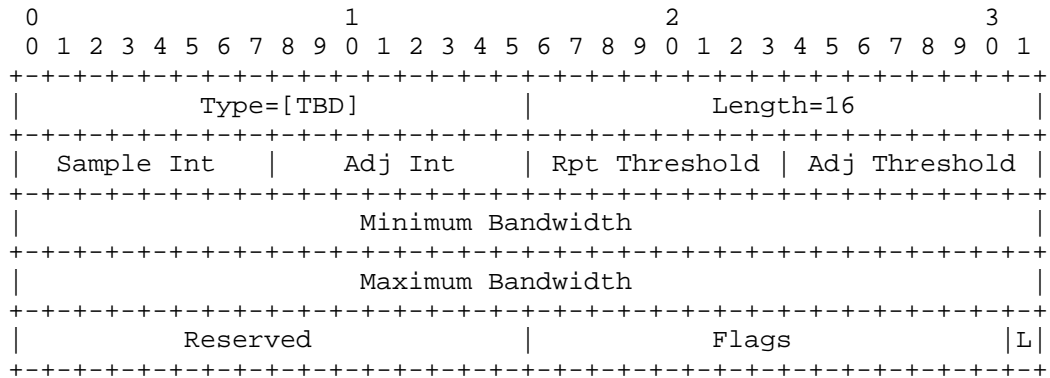
Also to avoid multiple LSP re-signaling, sometimes operator set up longer adjustment intervals. However long adjustment-interval can also result in an undesirable effect of masking sudden changes in traffic patterns. To avoid this, the stateful PCE MAY pre-maturely expire the adjustment-interval to accommodate sudden bursts in traffic.

5. Extensions to the PCEP

5.1. AUTO-BANDWIDTH-ATTRIBUTE TLV

The AUTO-BANDWIDTH-ATTRIBUTE TLV can be included as an optional TLV in the LSP object as described in [I-D.ietf-pce-stateful-pce]. Whenever the LSP with Auto-Bandwidth feature enabled is delegated, AUTO-BANDWIDTH-ATTRIBUTE TLV is carried in PCRpt message. The TLV provides PCE with the 'local configurable knobs' of this feature. In case of PCE Initiated LSP ([I-D.ietf-pce-pce-initiated-lsp]) with this feature enabled, this TLV is included in LSP object with PCInitiate message.

The format of the AUTO-BANDWIDTH-ATTRIBUTE TLV is shown in the following figure:



AUTO-BANDWIDTH-ATTRIBUTE TLV format

The type of the TLV is [TBD] and it has a fixed length of 16 octets.

The value contains the following fields:

Sample Int (8 bits): The Sample-Interval, time interval in which the traffic rate is collected at the PCC.

Adj Int (8 bits): The Adjustment-Interval, time interval in which the bandwidth adjustment should be made.

Rpt Threshold (8 bits): The Report-Threshold value is encoded in percentage. Only if the percentage difference between the between the current BwSample and the last BwSample is greater than or equal to the threshold percentage the real time bandwidth sample is reported to PCE.

Adj Threshold (8 bits): The Adjust-Threshold value is encoded in percentage. Only if the percentage difference between the current MaxAvgBw and the current bandwidth allocation is greater than or equal to the threshold percentage the LSP bandwidth is adjusted to the current bandwidth demand.

Minimum Bandwidth (32 bits): The minimum bandwidth allowed is encoded in IEEE floating point format (see [IEEE.754.1985]), expressed in bytes per second. Refer to Section 3.1.2 of [RFC3471] for a table of commonly used values.

Maximum Bandwidth (32 bits): The maximum bandwidth allowed is encoded in IEEE floating point format (see [IEEE.754.1985]), expressed in bytes per second. Refer to Section 3.1.2 of [RFC3471] for a table of commonly used values.

Flags (16 bits): One flag is currently defined:

- * L (Live-Traffic - 1 bit): If set, PCC SHOULD report the live traffic information flowing on the LSP as per the Report-Threshold set. Otherwise PCC only reports the calculated bandwidth to be adjusted to the PCE.

Unassigned flags MUST be set to zero on transmission and MUST be ignored on receipt.

Reserved (16 bits): This field MUST be set to zero on transmission and MUST be ignored on receipt.

If the above parameters are not specified by the user, based on the local policy at Ingress (PCC) the default value can be encoded.

If no default value is specified at Ingress, value 'zero' can be encoded for the particular field. The stateful PCE can then apply its own default value based on the local policy.

5.2. BANDWIDTH Object

As per [RFC5440], the BANDWIDTH object is defined with two Object-Type values:

- o Requested Bandwidth: BANDWIDTH Object-Type is 1.
- o Re-optimization Bandwidth: Bandwidth of an existing TE LSP for which a reoptimization is requested. BANDWIDTH Object-Type is 2.

The new BANDWIDTH object type 3 [TBD] is used to specify the BwSample determined from the existing TE LSP Traffic flow at every sample-interval when L bit is set in AUTO-BANDWIDTH-ATTRIBUTE TLV. The Report-Threshold percentage is used to determine if there is a need to report the current BwSample.

If Live-Traffic (L-Bit) is not set, PCC only reports the calculated bandwidth to be adjusted (MaxAvgBw) to the PCE. This is done via the existing 'Requested Bandwidth with BANDWIDTH Object-Type as 1'.

5.3. The PCRpt Message

When the delegated LSP is enabled with the Auto-Bandwidth adjustment feature with Live-Traffic (L-Bit) set, PCC SHOULD include the BANDWIDTH object of type 3 [TBD] in the PCRpt message. The definition of the PCRpt message (see [I-D.ietf-pce-stateful-pce]) is unchanged.

When LSP is delegated to a PCE for the very first time, BANDWIDTH object of type 1 is used to specify the requested bandwidth in the PCRpt message. To report the live traffic flow information (as the BwSample) the BANDWIDTH object of type 3 [TBD] is encoded in further PCRpt message.

If Live-Traffic (L-Bit) is not set, PCC SHOULD include the BANDWIDTH object of type 1 to specify the he calculated bandwidth to be adjusted to the PCE.

5.4. The PCInitiate Message

For PCE Initiated LSP ([I-D.ietf-pce-pce-initiated-lsp]) with Auto-Bandwidth feature enabled, AUTO-BANDWIDTH-ATTRIBUTE TLV is included in LSP object with the PCInitiate message. The rest of the processing remains unchanged.

6. Security Considerations

This document defines a new BANDWIDTH type and AUTO-BANDWIDTH-ATTRIBUTE TLV which does not add any new security concerns beyond those discussed in [RFC5440] and [I-D.ietf-pce-stateful-pce] in itself. Some deployments may find the live traffic bandwidth information as extra sensitive and thus should employ suitable PCEP security mechanisms like TCP-AO or [I-D.ietf-pce-pceps].

7. Manageability Considerations

7.1. Control of Function and Policy

The Auto-Bandwidth feature MUST BE controlled per tunnel at Ingress (PCC), the values for parameters like sample-interval, adjustment-interval, minimum-bandwidth, maximum-bandwidth, report-threshold, adjust-threshold, Live-Traffic (L-Bit) SHOULD BE configurable by an operator.

7.2. Information and Data Models

[RFC7420] describes the PCEP MIB, there are no new MIB Objects for this document.

7.3. Liveness Detection and Monitoring

Mechanisms defined in this document do not imply any new liveness detection and monitoring requirements in addition to those already listed in [RFC5440].

7.4. Verify Correct Operations

Mechanisms defined in this document do not imply any new operation verification requirements in addition to those already listed in [RFC5440].

7.5. Requirements On Other Protocols

Mechanisms defined in this document do not imply any new requirements on other protocols.

7.6. Impact On Network Operations

Mechanisms defined in this document do not have any impact on network operations in addition to those already listed in [RFC5440].

8. IANA Considerations

8.1. PCEP TLV Type Indicators

This document defines the following new PCEP TLVs; IANA is requested to make the following allocations from this registry.

Value	Meaning	Reference
TBD	AUTO-BANDWIDTH-ATTRIBUTE	[This I.D.]

8.2. AUTO-BANDWIDTH-ATTRIBUTE

This document requests that a registry is created to manage the Flags field in the AUTO-BANDWIDTH-ATTRIBUTE TLV in the LSP object. New values are to be assigned by Standards Action [RFC5226]. Each bit should be tracked with the following qualities:

- o Bit number (counting from bit 0 as the most significant bit)
- o Capability description
- o Defining RFC

The following values are defined in this document:

Bit	Description	Reference
31	Live-Traffic (L-Bit)	[This I.D.]

8.3. BANDWIDTH Object

This document defines new object type for the BANDWIDTH object; IANA is requested to make the following allocations from this registry.

Object-Class Value	Name	Reference
5	BANDWIDTH Object-Type	[This I.D.]
	3: MaxAvgBw determined from the existing TE LSP Traffic flow.	

9. Acknowledgments

We would like to thank Venugopal Reddy, Reeja Paul, Sandeep Boina and Avantika for their useful comments and suggestions.

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Appendix A. Contributor Addresses

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PCEP Extensions for MPLS-TE LSP Automatic Bandwidth Adjustment with
Stateful PCE
draft-dhody-pce-stateful-pce-auto-bandwidth-09

Abstract

The Path Computation Element Communication Protocol (PCEP) provides mechanisms for Path Computation Elements (PCEs) to perform path computations in response to Path Computation Clients (PCCs) requests. The stateful PCE extensions allow stateful control of Multi-Protocol Label Switching (MPLS) Traffic Engineering Label Switched Paths (TE LSPs) using PCEP.

This document describes PCEP extensions for automatic bandwidth adjustment when employing an Active Stateful PCE for both PCE-initiated and PCC-initiated LSPs.

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1. Introduction

[RFC5440] describes the Path Computation Element Protocol (PCEP) as a communication mechanism between a Path Computation Client (PCC) and a Path Control Element (PCE), or between PCE and PCE, that enables computation of Multi-Protocol Label Switching (MPLS) Traffic Engineering Label Switched Paths (TE LSPs).

[I-D.ietf-pce-stateful-pce] specifies extensions to PCEP to enable stateful control of MPLS TE LSPs. It describes two mode of operations - Passive stateful PCE and Active stateful PCE. In this document, the focus is on Active stateful PCE where LSPs are configured at the PCC and control over them is delegated to the PCE. Further [I-D.ietf-pce-pce-initiated-lsp] describes the setup, maintenance and teardown of PCE-initiated LSPs for the stateful PCE model.

Over time, based on the varying traffic pattern, an LSP established with certain bandwidth may require to adjust the bandwidth, reserved in the network automatically. Ingress Label Switch Router (LSR) collects the traffic rate at each sample interval to determine the bandwidth demand of the LSP. This bandwidth information is then used to adjust the LSP bandwidth periodically. This feature is commonly referred to as Auto-Bandwidth.

Enabling Auto-Bandwidth feature on an LSP results in the LSP automatically adjusting its bandwidth reservation based on the actual traffic flowing through the LSP. The initial LSP bandwidth can be set to an arbitrary value (including zero), in practice, it can be operator expected value based on design and planning. Once the LSP is set-up, the LSP monitors the traffic flow and adjusts its bandwidth every adjustment-interval period. The bandwidth adjustment uses the make-before-break signaling method so that there is no interruption to the traffic flow. The Auto-Bandwidth is described in detail in Section 4.1. [I-D.ietf-pce-stateful-pce-app] describes the use-case for Auto-Bandwidth adjustment for passive and active stateful PCE.

- o The PCC (head-end of the LSP) monitors and calculates the new adjusted bandwidth. The PCC reports the calculated bandwidth to be adjusted to the PCE.
- o This approach would be similar to passive stateful PCE model, while the passive stateful PCE uses path request/reply mechanism, the active stateful PCE uses report/update mechanism to adjust the LSP bandwidth.
- o For PCE-initiated LSP, the PCC is requested during the LSP initiation to monitor and calculate the new adjusted bandwidth.

This document defines extensions needed to support Auto-Bandwidth feature on the LSPs in a active stateful PCE model using PCEP.

2. Conventions Used in This Document

2.1. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

2.2. Terminology

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Active Stateful PCE: PCE that uses tunnel state information learned from PCCs to optimize path computations. Additionally, it actively updates tunnel parameters in those PCCs that delegated control over their tunnels to the PCE.

Delegation: An operation to grant a PCE temporary rights to modify a subset of tunnel parameters on one or more PCC's tunnels. Tunnels

are delegated from a PCC to a PCE.

PCC: Path Computation Client. Any client application requesting a path computation to be performed by a Path Computation Element.

PCE: Path Computation Element. An entity (component, application, or network node) that is capable of computing a network path or route based on a network graph and applying computational constraints.

TE LSP: Traffic Engineering Label Switched Path.

Note the Auto-Bandwidth feature specific terms defined in Section 4.1.

3. Requirements for PCEP Extensions

The PCEP speaker supporting this document MUST have a mechanism to advertise the automatic bandwidth adjustment capability.

PCEP extensions required are summarized in the following table.

PCC Initiated	PCE Initiated
PCC monitors the traffic and reports the calculated bandwidth to be adjusted to the PCE.	At the time of initiation, PCE request PCC to monitor the traffic and report the calculated bandwidth to be adjusted to the PCE.
No new extensions are needed.	Extension is needed for PCE to pass on the adjustment parameters at the time of Initiation.

Table 1: Auto-Bandwidth PCEP extensions

Further Auto-Bandwidth deployment considerations are summarized below:

- o It is required to identify and inform the PCEP peer, the LSP that are enabled with Auto-Bandwidth feature. Not all LSPs in some

deployments would like their bandwidth to be dependent on the real-time bandwidth usage but be constant as set by the operator.

- o Further for the LSP with Auto-Bandwidth feature enabled, an operator should be able to specify the adjustment parameters (i.e. configuration knobs) to control this feature (e.g. minimum/maximum bandwidth range) and PCEP peer should be informed.

4. Architectural Overview

4.1. Auto-Bandwidth Overview

Auto-Bandwidth feature allows an LSP to automatically and dynamically adjust its reserved bandwidth over time, i.e. without network operator intervention. The bandwidth adjustment uses the make-before-break signaling method so that there is no interruption to the traffic flow.

The new bandwidth reservation is determined by sampling the actual traffic flowing through the LSP. If the traffic flowing through the LSP is lower than the configured or current bandwidth of the LSP, the extra bandwidth is being reserved needlessly and being wasted. Conversely, if the actual traffic flowing through the LSP is higher than the configured or current bandwidth of the LSP, it can potentially cause congestion or packet loss in the network. With Auto-Bandwidth feature, the LSP bandwidth can be set to some arbitrary value (including zero) during initial setup time, and it will be periodically adjusted over time based on the actual bandwidth requirement.

Note the following definitions of the Auto-Bandwidth terms:

Maximum Average Bandwidth (MaxAvgBw): The maximum average bandwidth represents the current traffic bandwidth demand during a time interval. This is the maximum value of the averaged traffic bandwidth rate in a given adjustment-interval.

Adjusted Bandwidth: This is the Auto-Bandwidth computed bandwidth that needs to be adjusted for the LSP.

Sample-Interval: The periodic time interval at which the traffic rate is collected as a sample.

Bandwidth-Sample (BwSample): The bandwidth sample collected at every sample interval to measure the traffic rate.

Adjustment-Interval: The periodic time interval at which the

bandwidth adjustment should be made using the MaxAvgBw.

Maximum-Bandwidth: The maximum bandwidth that can be reserved for the LSP.

Minimum-Bandwidth: The minimum bandwidth that can be reserved for the LSP.

Adjustment-Threshold: This value is used to decide when the bandwidth should be adjusted. If the percentage or absolute difference between the current MaxAvgBw and the current bandwidth reservation is greater than or equal to the threshold value, the LSP bandwidth is adjusted to the current bandwidth demand (Adjusted Bandwidth) at the adjustment-interval expiry.

Overflow-Count: This value is used to decide when the bandwidth should be adjusted when there is a sudden increase in traffic demand. This value indicates how many times consecutively, the percentage or absolute difference between the current MaxAvgBw and the current bandwidth reservation is greater than or equal to the Overflow-Threshold value.

Overflow-Threshold: This value is used to decide when the bandwidth should be adjusted when there is a sudden increase in traffic demand. If the percentage or absolute difference between the current MaxAvgBw and the current bandwidth reservation is greater than or equal to the threshold value, the overflow-condition is set to be met. The LSP bandwidth is adjusted to the current bandwidth demand bypassing the adjustment-interval if the overflow-condition is met consecutively for the Overflow-Count.

Underflow-Count: This value is used to decide when the bandwidth should be adjusted when there is a sudden decrease in traffic demand. This value indicates how many times consecutively, the percentage or absolute difference between the current MaxAvgBw and the current bandwidth reservation is greater than or equal to the Underflow-Threshold value.

Underflow-Threshold: This value is used to decide when the bandwidth should be adjusted when there is a sudden decrease in traffic demand. If the percentage or absolute difference between the current MaxAvgBw and the current bandwidth reservation is greater than or equal to the threshold value, the underflow-condition is set to be met. The LSP bandwidth is adjusted to the current bandwidth demand bypassing the adjustment-interval if the underflow-condition is met consecutively for the Underflow-Count.

4.2. Auto-bandwidth Theory of Operation

The traffic rate is periodically sampled at each sample-interval (which can be configured by the user and the default value as 5 minutes) by the head-end node of the LSP. The sampled traffic rates are accumulated over the adjustment-interval period (which can be configured by the user and the default value as 24 hours). The PCEP peer which is in-charge of calculating the bandwidth to be adjusted, will adjust the bandwidth of the LSP to the highest sampled traffic rate (MaxAvgBw) amongst the set of bandwidth samples collected over the adjustment-interval.

Note that the highest sampled traffic rate could be higher or lower than the current LSP bandwidth. Only if the difference between the current bandwidth demand (MaxAvgBw) and the current bandwidth reservation is greater than or equal to the Adjustment-Threshold (percentage or absolute value), the LSP bandwidth is adjusted to the current bandwidth demand (MaxAvgBw). Some LSPs are less eventful while other LSPs may encounter a lot of changes in the traffic pattern. PCE sets the intervals for adjustment based on the traffic pattern of the LSP.

In order to avoid frequent re-signaling, an operator may set a longer adjustment-interval value. However, longer adjustment-interval can result in an undesirable effect of masking sudden changes in traffic demands of an LSP. To avoid this, the Auto-Bandwidth feature may pre-maturely expire the adjustment-interval and adjust the LSP bandwidth to accommodate the sudden bursts of increase in traffic demand as an overflow condition or decrease in traffic demand as an underflow condition.

All thresholds in this document could be represented in both absolute value and percentage, and could be used together.

4.3. Scaling Considerations

It should be noted that any bandwidth change would require re-signaling of an LSP in a make-before-break fashion, which can further trigger preemption of lower priority LSPs in the network. When deployed under scale, this can lead to a signaling churn in the network. The Auto-bandwidth application algorithm is thus advised to take this into consideration before adjusting the LSP bandwidth. Operators are advised to set the values of various auto-bandwidth adjustment parameters appropriate for the deployed LSP scale.

If a PCE gets overwhelmed, it can notify the PCC to temporarily suspend its auto-bandwidth reporting (see Section 5.6). Similarly if

a PCC gets overwhelmed due to signaling churn, it can notify the PCE to temporarily suspend the LSP bandwidth adjustment.

5. Extensions to the PCEP

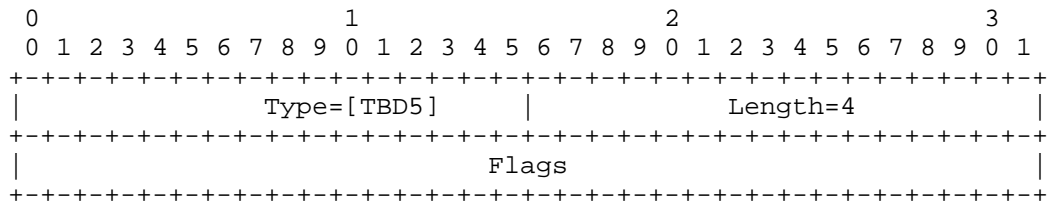
5.1. Capability Advertisement

During PCEP Initialization Phase, PCEP Speakers (PCE or PCC) advertise their support of Automatic Bandwidth Adjustment. A PCEP Speaker includes the "Auto-Bandwidth Capability" TLV, in the OPEN Object to advertise its support for PCEP Auto-Bandwidth extensions. The presence of the "Auto-Bandwidth Capability" TLV in the OPEN Object indicates that the Automatic Bandwidth Adjustment is supported as described in this document.

The PCEP protocol extensions for Auto-Bandwidth adjustments MUST NOT be used if one or both PCEP Speakers have not included the "Auto-Bandwidth Capability" TLV in their respective OPEN message. If the PCEP speaker that supports the extensions of this draft but did not advertise this capability, then upon receipt of AUTO-BANDWIDTH-ATTRIBUTE TLV in LSPA object, it SHOULD generate a PCerr with error-type 19 (Invalid Operation), error-value TBD4 (Auto-Bandwidth capability was not advertised) and it will terminate the PCEP session.

5.1.1 AUTO-BANDWIDTH-CAPABILITY TLV

The AUTO-BANDWIDTH-CAPABILITY TLV is an optional TLV for use in the OPEN Object for Automatic Bandwidth Adjustment via PCEP capability advertisement. Its format is shown in the following figure:



AUTO-BANDWIDTH-CAPABILITY TLV format

The type of the TLV is (TBD5) and it has a fixed length of 4 octets.

The value comprises a single field - Flags (32 bits). Currently no flags are defined for this TLV.

Unassigned bits are considered reserved. They MUST be set to 0 on transmission and MUST be ignored on receipt.

Advertisement of the Auto-Bandwidth capability TLV implies support of auto-bandwidth adjustment, as well as the objects, TLVs and procedures defined in this document.

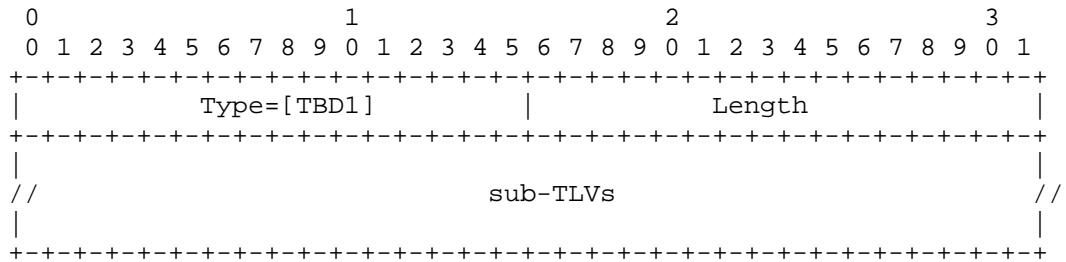
5.2. AUTO-BANDWIDTH-ATTRIBUTE TLV

The AUTO-BANDWIDTH-ATTRIBUTE TLV provides the 'configurable knobs' of the feature and it can be included as an optional TLV in the LSPA Object (as described in [RFC5440]).

For PCE-Initiated LSP ([I-D.ietf-pce-pce-initiated-lsp]), this TLV is included in the LSPA Object with PCInitiate message. For delegated LSPs, this TLV is carried in PCRpt message in LSPA Object.

The TLV is encoded in all PCEP messages for the LSP till the auto bandwidth adjustment feature is enabled, the absence of the TLV indicate the PCEP speaker wish to disable the feature.

The format of the AUTO-BANDWIDTH-ATTRIBUTE TLV is shown in the following figure:



AUTO-BANDWIDTH-ATTRIBUTE TLV format

Type: TBD1

Length: Variable

Value: This comprises one or more sub-TLVs.

Following sub-TLVs are defined in this document:

Type	Len	Name
1	4	Sample-Interval sub-TLV

- 2 4 Adjustment-Interval sub-TLV
- 3 4 Adjustment-Threshold sub-TLV
- 4 4 Adjustment-Threshold-Percentage sub-TLV
- 5 4 Minimum-Bandwidth sub-TLV
- 6 4 Maximum-Bandwidth sub-TLV
- 7 8 Overflow-Threshold sub-TLV
- 8 4 Overflow-Threshold-Percentage sub-TLV
- 9 8 Underflow-Threshold sub-TLV
- 10 4 Underflow-Threshold-Percentage sub-TLV

Future specification can define additional sub-TLVs.

The presence of AUTO-BANDWIDTH-ATTRIBUTE TLV in LSPA Object means that the automatic bandwidth adjustment feature is enabled. All sub-TLVs are optional and any unrecognized sub-TLV MUST be silently ignored. If a sub-TLV of same type appears more than once, only the first occurrence is processed and all others MUST be ignored.

The AUTO-BANDWIDTH-ATTRIBUTE TLV can also be carried in PCUpd message in LSPA Object in order to make updates to auto-bandwidth attributes such as Adjustment-Interval.

If sub-TLVs are not present, the default values based on the local policy are assumed.

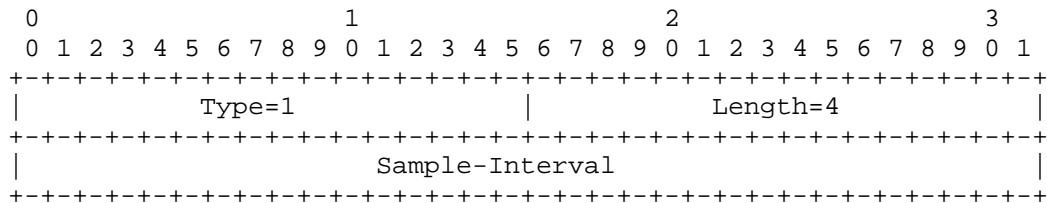
The sub-TLVs are encoded to inform the PCEP peer the various sampling and adjustment parameters.

The following sub-sections describe the sub-TLVs which are currently defined to be carried within the AUTO-BANDWIDTH-ATTRIBUTE TLV.

5.2.1. Sample-Interval sub-TLV

The Sample-Interval sub-TLV specifies a time interval in seconds at which traffic samples are collected at the PCC.

The Type is 1, Length is 4, and the value comprises of 4-octet time interval, the valid range is from 1 to 604800, in seconds. The default value is 300 seconds.

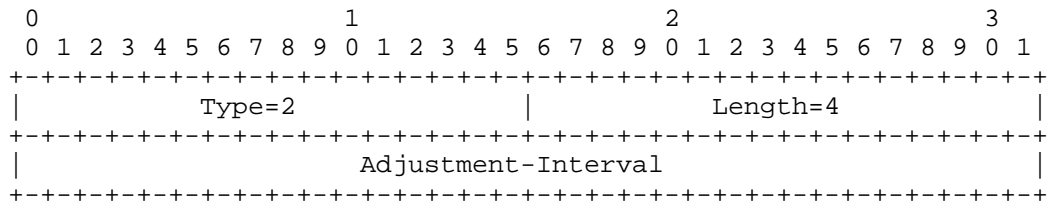


Sample-Interval sub-TLV format

5.2.2. Adjustment-Interval sub-TLV

The Adjustment-Interval sub-TLV specifies a time interval in seconds at which bandwidth adjustment should be made.

The Type is 2, Length is 4, and the value comprises of 4-octet time interval, the valid range is from 1 to 604800, in seconds. The default value is 300 seconds.



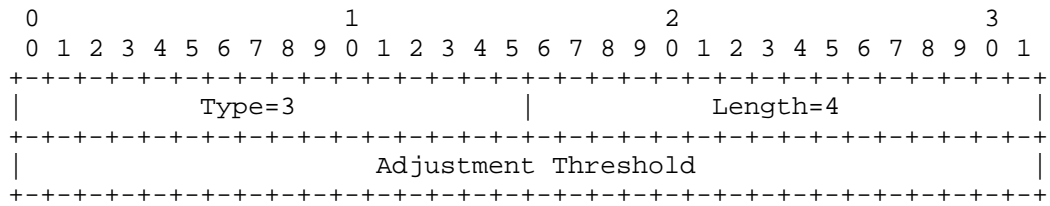
Adjustment-Interval sub-TLV format

5.2.3. Adjustment Threshold

The sub-TLVs in this section are encoded to inform the PCEP peer the adjustment threshold parameters. An implementation MAY include both sub-TLVs for the absolute value and the percentage, in which case the bandwidth is adjusted when either of the adjustment threshold conditions are met.

5.2.3.1. Adjustment-Threshold sub-TLV

The Adjustment-Threshold sub-TLV is used to decide when the LSP bandwidth should be adjusted.



Adjustment-Threshold sub-TLV format

The Type is 3, Length is 4, and the value comprises of -

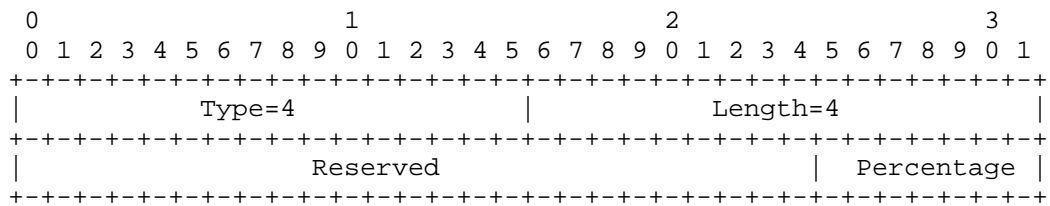
- o Adjustment Threshold: The absolute Adjustment-Threshold bandwidth value, encoded in IEEE floating point format (see

[IEEE.754.1985]), expressed in bytes per second. Refer to Section 3.1.2 of [RFC3471] for a table of commonly used values.

If the difference between the current MaxAvgBw and the current bandwidth reservation is greater than or equal to the threshold value, the LSP bandwidth is adjusted to the current bandwidth demand.

5.2.3.2. Adjustment-Threshold-Percentage sub-TLV

The Adjustment-Threshold-Percentage sub-TLV is used to decide when the LSP bandwidth should be adjusted.



Adjustment-Threshold-Percentage sub-TLV format

The Type is 4, Length is 4, and the value comprises of -

- o Reserved: SHOULD be set to zero on transmission and MUST be ignored on receipt.
- o Percentage: The Adjustment-Threshold value, encoded in percentage (an integer from 0 to 100). If the percentage difference between the current MaxAvgBw and the current bandwidth reservation is greater than or equal to the threshold percentage, the LSP bandwidth is adjusted to the current bandwidth demand.

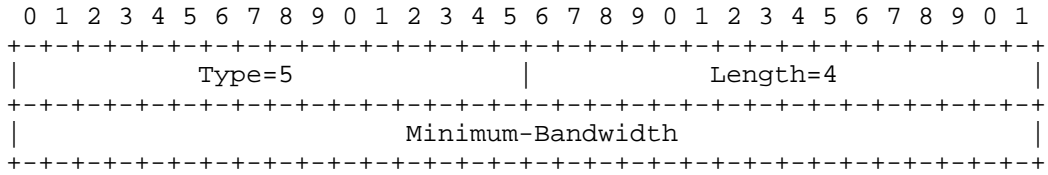
5.2.4. Minimum and Maximum Bandwidth Values

5.2.4.1. Minimum-Bandwidth sub-TLV

The Minimum-Bandwidth sub-TLV specify the minimum bandwidth allowed for the LSP, and is expressed in bytes per second. The LSP bandwidth cannot be adjusted below the minimum bandwidth value.

The Type is 5, Length is 4, and the value comprises of 4-octet bandwidth value encoded in IEEE floating point format (see [IEEE.754.1985]), expressed in bytes per second. Refer to Section 3.1.2 of [RFC3471] for a table of commonly used values.



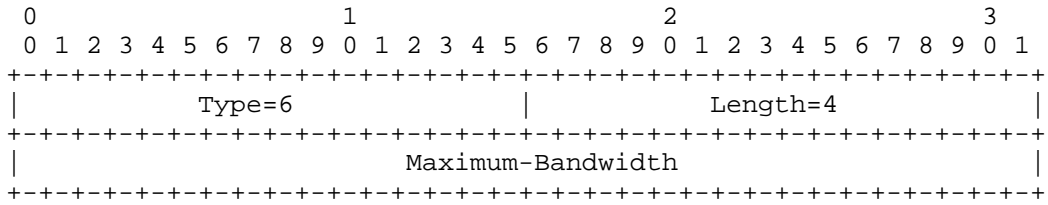


Minimum-Bandwidth sub-TLV format

5.2.4.2. Maximum-Bandwidth sub-TLV

The Maximum-Bandwidth sub-TLV specify the maximum bandwidth allowed for the LSP, and is expressed in bytes per second. The LSP bandwidth cannot be adjusted above the maximum bandwidth value.

The Type is 6, Length is 4, and the value comprises of 4-octet bandwidth value encoded in IEEE floating point format (see [IEEE.754.1985]), expressed in bytes per second. Refer to Section 3.1.2 of [RFC3471] for a table of commonly used values.



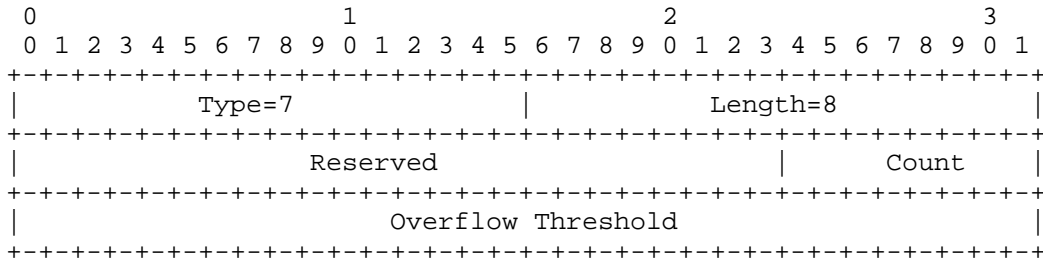
Maximum-Bandwidth sub-TLV format

5.2.5. Overflow and Underflow Conditions

The sub-TLVs in this section are encoded to inform the PCEP peer the overflow and underflow threshold parameters. An implementation MAY include sub-TLVs for the absolute value and the percentage for the threshold, in which case the bandwidth is immediately adjusted when either of the adjustment threshold conditions are met consecutively for the given count.

5.2.5.1. Overflow-Threshold sub-TLV

The Overflow-Threshold sub-TLV is used to decide if the bandwidth should be adjusted immediately.



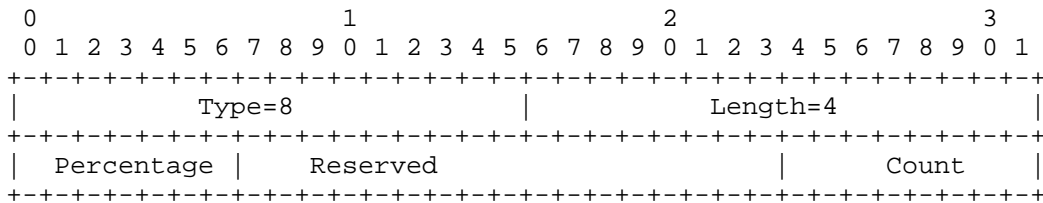
Overflow-Threshold sub-TLV format

The Type is 7, Length is 8, and the value comprises of -

- o Reserved: SHOULD be set to zero on transmission and MUST be ignored on receipt.
- o Count: The Overflow-Count value, encoded in integer. The value 0 is considered to be invalid. The number of consecutive samples for which the overflow condition MUST be met for the LSP bandwidth to be immediately adjusted to the current bandwidth demand, bypassing the adjustment-interval.
- o Overflow Threshold: The absolute Overflow-Threshold bandwidth value, encoded in IEEE floating point format (see [IEEE.754.1985]), expressed in bytes per second. Refer to Section 3.1.2 of [RFC3471] for a table of commonly used values. If the increase of the current MaxAvgBw from the current bandwidth reservation is greater than or equal to the threshold value, the overflow condition is met.

5.2.5.2. Overflow-Threshold-Percentage sub-TLV

The Overflow-Threshold-Percentage sub-TLV is used to decide if the bandwidth should be adjusted immediately.



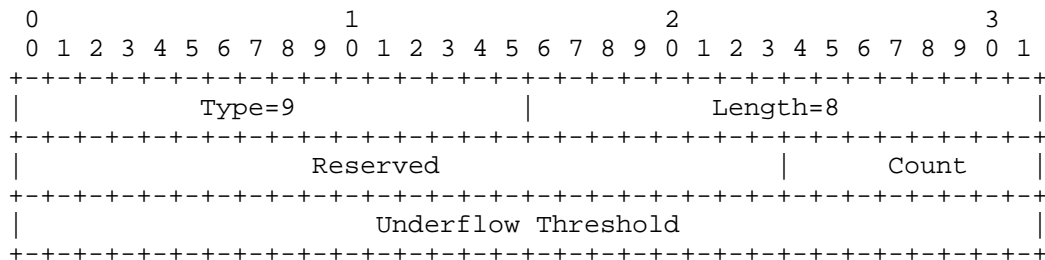
Overflow-Threshold-Percentage sub-TLV format

The Type is 8, Length is 4, and the value comprises of -

- o Percentage: The Overflow-Threshold value, encoded in percentage (an integer from 0 to 100). If the percentage increase of the current MaxAvgBw from the current bandwidth reservation is greater than or equal to the threshold percentage, the overflow condition is met.
- o Reserved: SHOULD be set to zero on transmission and MUST be ignored on receipt.
- o Count: The Overflow-Count value, encoded in integer. The value 0 is considered to be invalid. The number of consecutive samples for which the overflow condition MUST be met for the LSP bandwidth to be immediately adjusted to the current bandwidth demand, bypassing the adjustment-interval.

5.2.5.3. Underflow-Threshold sub-TLV

The Underflow-Threshold sub-TLV is used to decide if the bandwidth should be adjusted immediately.



Underflow-Threshold sub-TLV format

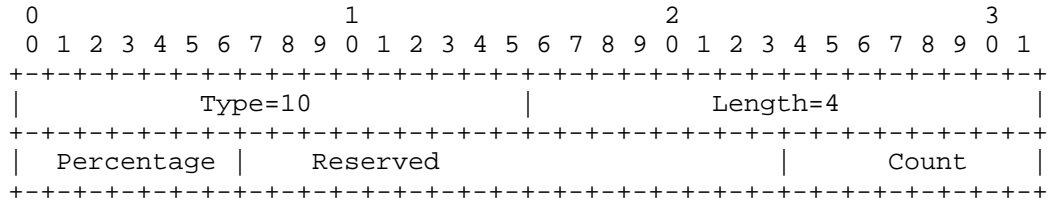
The Type is 9, Length is 8, and the value comprises of -

- o Reserved: SHOULD be set to zero on transmission and MUST be ignored on receipt.
- o Count: The Underflow-Count value, encoded in integer. The value 0 is considered to be invalid. The number of consecutive samples for which the underflow condition MUST be met for the LSP bandwidth to be immediately adjusted to the current bandwidth demand, bypassing the adjustment-interval.
- o Underflow Threshold: The absolute Underflow-Threshold bandwidth value, encoded in IEEE floating point format (see [IEEE.754.1985]), expressed in bytes per second. Refer to Section 3.1.2 of [RFC3471] for a table of commonly used values. If the decrease of the current MaxAvgBw from the current bandwidth

reservation is greater than or equal to the threshold value, the underflow condition is met.

5.2.5.4. Underflow-Threshold-Percentage sub-TLV

The Underflow-Threshold-Percentage sub-TLV is used to decide if the bandwidth should be adjusted immediately.



Underflow-Threshold-Percentage sub-TLV format

The Type is 10, Length is 4, and the value comprises of -

- o Percentage: The Underflow-Threshold value, encoded in percentage (an integer from 0 to 100). If the percentage decrease of the current MaxAvgBw from the current bandwidth reservation is greater than or equal to the threshold percentage, the underflow condition is met.
- o Reserved: SHOULD be set to zero on transmission and MUST be ignored on receipt.
- o Count: The Underflow-Count value, encoded in integer. The value 0 is considered to be invalid. The number of consecutive samples for which the underflow condition MUST be met for the LSP bandwidth to be immediately adjusted to the current bandwidth demand, bypassing the adjustment-interval.

5.3. BANDWIDTH Object

As per [RFC5440], the BANDWIDTH object (Object-Class value 5) is defined with two Object-Type values as following:

- o Requested Bandwidth: BANDWIDTH Object-Type value is 1.
- o Re-optimization Bandwidth: Bandwidth of an existing TE LSP for which a re-optimization is requested. BANDWIDTH Object-Type value is 2.

PCC reports the calculated bandwidth to be adjusted (MaxAvgBw) to the

PCE using existing 'Requested Bandwidth with BANDWIDTH Object-Type as 1.

5.4. The PCInitiate Message

A PCInitiate message is a PCEP message sent by a PCE to a PCC to trigger LSP instantiation or deletion [I.D.ietf-pce-pce-initiated-lsp].

For the PCE-initiated LSP [I.D.ietf-pce-pce-initiated-lsp] with Auto-Bandwidth feature enabled, AUTO-BANDWIDTH-ATTRIBUTE TLV MUST be included in the LSPA object with the PCInitiate message. The rest of the processing remains unchanged.

5.5. The PCRpt Message

As specified in [I.D.ietf-pce-pce-initiated-lsp], the PCC creates the LSP using the attributes communicated by the PCE, and local values for the unspecified parameters. After the successful instantiation of the LSP, PCC automatically delegates the LSP to the PCE and generates an LSP State Report (PCRpt) for the LSP.

When LSP is delegated to a PCE for the very first time, BANDWIDTH object of type 1 is used to specify the requested bandwidth in the PCRpt message.

When the LSP is enabled with the Auto-Bandwidth feature, PCC SHOULD include the BANDWIDTH object of type 1 to specify the calculated bandwidth to be adjusted to the PCE in the PCRpt message.

The definition of the PCRpt message (see [I-D.ietf-pce-stateful-pce]) is unchanged by this document.

5.6. The PCNtf Message

As per [RFC5440], the PCEP Notification message (PCNtf) can be sent by a PCEP speaker to notify its peer of a specific event. As described in Section 4.3 of this document, a PCEP speaker SHOULD notify its PCEP peer that it is overwhelmed, and on receipt of such notification the peer SHOULD NOT send any PCEP messages related to auto-bandwidth adjustment. If a PCEP message related to auto-bandwidth adjustment is received, it MUST be silently ignored.

When a PCEP speaker is overwhelmed, it SHOULD notify its peer by sending a PCNtf message with Notification Type = TBD6 (Auto-bandwidth

Overwhelm State) and Notification Value = 1 (Entering auto-bandwidth overwhelm state). Optionally, OVERLOADED-DURATION TLV [RFC5440] MAY be included that specifies the time period during which no further PCEP messages related to auto-bandwidth adjustment should be sent. When the PCEP speaker is no longer in the overwhelm state and is available to process the auto-bandwidth adjustment, it SHOULD notify its peer by sending a PCNtf message with Notification Type = TBD6 (Auto-bandwidth Overwhelm State) and Notification Value = 2 (Clearing auto-bandwidth overwhelm state).

When Auto-Bandwidth feature is deployed, a PCE can send this notification to PCC when a PCC is reporting frequent auto-bandwidth adjustments. If a PCC is overwhelmed with re-signaling/re-routing, it can also notify the PCE to not adjust the LSP bandwidth while in overwhelm state.

6. Security Considerations

This document defines AUTO-BANDWIDTH-CAPABILITY TLV, AUTO-BANDWIDTH-ATTRIBUTE TLV which do not add any new security concerns beyond those discussed in [RFC5440] and [I-D.ietf-pce-stateful-pce].

Some deployments may find the reporting of the auto-bandwidth information as extra sensitive and thus SHOULD employ suitable PCEP security mechanisms like TCP-AO or [I-D.ietf-pce-pceps].

7. Manageability Considerations

7.1. Control of Function and Policy

The Auto-Bandwidth feature SHOULD be controlled per tunnel (at ingress (PCC) or PCE), the values for parameters like sample-interval, adjustment-interval, minimum-bandwidth, maximum-bandwidth, adjustment-threshold SHOULD be configurable by an operator.

7.2. Information and Data Models

[RFC7420] describes the PCEP MIB, there are no new MIB Objects defined in this document.

7.3. Liveness Detection and Monitoring

Mechanisms defined in this document do not imply any new liveness detection and monitoring requirements in addition to those already listed in [RFC5440].

7.4. Verify Correct Operations

Mechanisms defined in this document do not imply any new operation verification requirements in addition to those already listed in [RFC5440].

7.5. Requirements On Other Protocols

Mechanisms defined in this document do not add any new requirements on other protocols.

7.6. Impact On Network Operations

Mechanisms defined in this document do not have any impact on network operations in addition to those already listed in [RFC5440].

8. IANA Considerations

8.1. PCEP TLV Type Indicators

This document defines the following new PCEP TLVs; IANA is requested to make the following allocations from this registry.

<<http://www.iana.org/assignments/pcep/pcep.xhtml#pcep-tlv-type-indicators>>.

Value	Name	Reference
TBD5	AUTO-BANDWIDTH-CAPABILITY	[This I.D.]
TBD1	AUTO-BANDWIDTH-ATTRIBUTE	[This I.D.]

8.2. AUTO-BANDWIDTH-CAPABILITY TLV Flag Field

IANA is requested to create a registry to manage the Flag field of the AUTO-BANDWIDTH-CAPABILITY TLV.

New bit numbers may be allocated only by an IETF Consensus action. Each bit should be tracked with the following qualities:

- o Bit number (counting from bit 0 as the most significant bit)
- o Capability description
- o Defining RFC

No bit is defined for the AUTO-BANDWIDTH-CAPABILITY TLV Object flag field in this document.

8.3. AUTO-BANDWIDTH-ATTRIBUTE Sub-TLV

This document specifies the AUTO-BANDWIDTH-ATTRIBUTE Sub-TLVs. IANA is requested to create an "AUTO-BANDWIDTH-ATTRIBUTE Sub-TLV Types" sub-registry in the "PCEP TLV Type Indicators" for the sub-TLVs carried in the AUTO-BANDWIDTH-ATTRIBUTE TLV. This document defines the following types:

Type	Name	Reference
0	Reserved	[This I.D.]
1	Sample-Interval sub-TLV	[This I.D.]
2	Adjustment-Interval sub-TLV	[This I.D.]
3	Adjustment-Threshold sub-TLV	[This I.D.]
4	Adjustment-Threshold-Percentage sub-TLV	[This I.D.]

5	Minimum-Bandwidth sub-TLV	[This I.D.]
6	Maximum-Bandwidth sub-TLV	[This I.D.]
7	Overflow-Threshold sub-TLV	[This I.D.]
8	Overflow-Threshold-Percentage sub-TLV	[This I.D.]
9	Underflow-Threshold sub-TLV	[This I.D.]
10	Underflow-Threshold-Percentage sub-TLV	[This I.D.]
11-	Unassigned	[This I.D.]
65535		

8.4. Error Object

This document defines a new Error-Value for PCErr message of type 19 (Invalid Operation) [I-D.ietf-pce-stateful-pce]); IANA is requested to make the following allocation from this registry.
[<http://www.iana.org/assignments/pcep/pcep.xhtml#pcep-error-object>](http://www.iana.org/assignments/pcep/pcep.xhtml#pcep-error-object)

Error-Value	Meaning	Reference
TBD4	Auto-Bandwidth Capability was not Advertised	[This I.D.]

8.5. Notification Object

IANA is requested to allocate new Notification Types and Notification Values within the "Notification Object" sub-registry of the PCEP Numbers registry, as follows:

Type	Meaning	Reference
TBD6	Auto-Bandwidth Overwhelm State	[This I.D.]
	Notification-value=1: Entering Auto-Bandwidth overwhelm state	
	Notification-value=2: Clearing Auto-Bandwidth overwhelm state	

9. References

9.1. Normative References

- [RFC2119] Bradner, S., "Key words for use in RFCs to Indicate Requirement Levels", BCP 14, RFC 2119, March 1997.
- [RFC5440] Vasseur, JP. and JL. Le Roux, "Path Computation Element (PCE) Communication Protocol (PCEP)", RFC 5440, March 2009.

[I-D.ietf-pce-stateful-pce] Crabbe, E., Minei, I., Medved, J., and R. Varga, "PCEP Extensions for Stateful PCE", draft-ietf-pce-stateful-pce (work in progress).

[I-D.ietf-pce-pce-initiated-lsp] Crabbe, E., Minei, I., Sivabalan, S., and R. Varga, "PCEP Extensions for PCE-initiated LSP Setup in a Stateful PCE Model", draft-ietf-pce-pce-initiated-lsp (work in progress).

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Conveying path setup type in PCEP messages
draft-ietf-pce-lsp-setup-type-10

Abstract

A Path Computation Element (PCE) can compute Traffic Engineering (TE) paths through a network that are subject to various constraints. Currently, TE paths are Label Switched Paths (LSPs) which are set up using the RSVP-TE signaling protocol. However, other TE path setup methods are possible within the PCE architecture. This document proposes an extension to the PCE communication protocol (PCEP) to allow support for different path setup methods over a given PCEP session.

Status of This Memo

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1. Introduction

[RFC5440] describes the Path Computation Element communication Protocol (PCEP) for communication between a Path Computation Client (PCC) and a Path Computation Element (PCE), or between a PCE and a PCE. A PCC requests a path subject to various constraints and optimization criteria from a PCE. The PCE responds to the PCC with a hop-by-hop path in an Explicit Route Object (ERO). The PCC uses the ERO to set up the path in the network.

[RFC8231] specifies extensions to PCEP that allow a PCC to delegate its LSPs to a PCE. The PCE can then update the state of LSPs

delegated to it. In particular, the PCE may modify the path of an LSP by sending a new ERO. The PCC uses this ERO to re-route the LSP in a make-before-break fashion. [RFC8281] specifies a mechanism allowing a PCE to dynamically instantiate an LSP on a PCC by sending the ERO and the characteristics of the LSP. The PCC creates the LSP using the ERO and other attributes sent by the PCE.

So far, PCEP and its extensions have assumed that the TE paths are label switched and are established via the RSVP-TE protocol. However, other methods of LSP setup are possible in the PCE architecture (see [RFC4655] and [RFC4657]). This document generalizes PCEP to allow other LSP setup methods to be used. It defines two new TLVs and specifies the base procedures to facilitate this, as follows.

- o The PATH-SETUP-TYPE-CAPABILITY TLV, which allows a PCEP speaker to announce which LSP setup methods it supports when the PCEP session is established.
- o The PATH-SETUP-TYPE TLV, which allows a PCEP speaker to specify which setup method should be used for a given LSP. When multiple path setup types are deployed in a network, a given PCEP session may have to simultaneously support more than one path setup type. A PCEP speaker uses the PATH-SETUP-TYPE TLV to explicitly indicate the intended path setup type in the appropriate PCEP messages, unless the path setup type is RSVP-TE (which is assumed to be the path setup type if no other setup type is indicated). This is so that both the PCC and the PCE can take the necessary steps to set up the path.

This document defines a path setup type code for RSVP-TE. When a new path setup type (other than RSVP-TE) is introduced for setting up a path, a path setup type code and, optionally, a sub-TLV pertaining to the new path setup type will be defined by the document that specifies the new path setup type.

1.1. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

2. Terminology

The following terminologies are used in this document:

- ERO: Explicit Route Object.
- LSR: Label Switching Router.
- PCC: Path Computation Client.
- PCE: Path Computation Element.
- PCEP: Path Computation Element Protocol.
- PST: Path Setup Type.
- TLV: Type, Length, and Value.

3. Path Setup Type Capability TLV

A PCEP speaker indicates which PSTs it supports during the PCEP initialization phase, as follows. When the PCEP session is created, it sends an Open message with an OPEN object containing the PATH-SETUP-TYPE-CAPABILITY TLV. The format of this TLV is as follows.

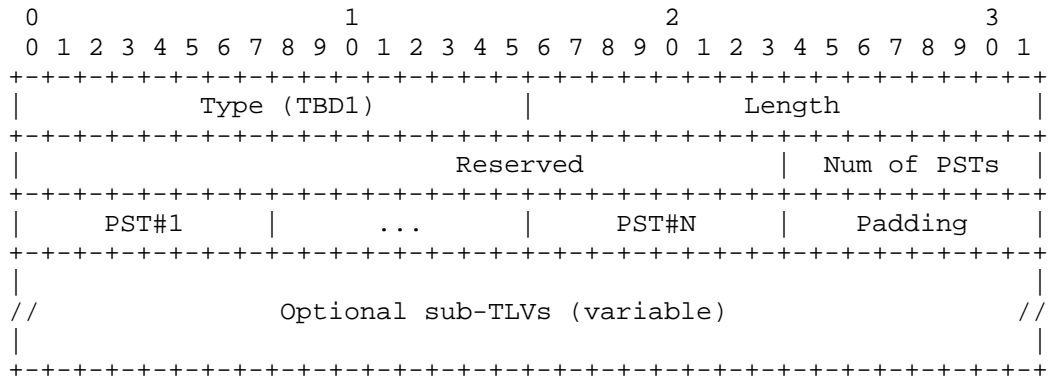


Figure 1: PATH-SETUP-TYPE-CAPABILITY TLV

The TLV type is TBD1 (to be assigned by IANA). Its reserved field MUST be set to zero by the sender and MUST be ignored by the receiver. The other fields in the TLV are as follows.

Length: The total length in bytes of the remainder of the TLV, that is, excluding the Type and Length fields.

Number of PSTs: The number of PSTs in the following list, excluding padding.

List of PSTs: A list of the PSTs that the PCEP speaker supports. Each PST is a single byte in length. Duplicate entries in this list MUST be ignored. The PCEP speaker MUST pad the list with zeros so that it is a multiple of four bytes in length. This document defines the following PST value:

* PST = 0: Path is setup using the RSVP-TE signaling protocol.

Optional sub-TLVs: A list of sub-TLVs associated with the supported PSTs. Each PST has zero or one sub-TLVs associated with it, and each sub-TLV is associated with exactly one PST. Each sub-TLV MUST obey the rules for TLV formatting defined in ([RFC5440]). That is, each sub-TLV is padded to a four byte alignment, and the length field of each sub-TLV does not include the padding bytes. This document does not define any sub-TLVs; an example can be found in [I-D.ietf-pce-segment-routing].

A PCEP speaker MUST check that this TLV is correctly formatted, as follows.

- o If there are no sub-TLVs, then the TLV length field MUST be equal to four bytes plus the size of the PST list, excluding any padding bytes.
- o If there are sub-TLVs then the TLV Length field MUST be equal to four bytes plus the size of the PST list (rounded up to the nearest multiple of four) plus the size of the appended sub-TLVs excluding any padding bytes in the final sub-TLV.
- o The Number of PSTs field MUST be greater than zero.

If a PCEP speaker receives a PATH-SETUP-TYPE-CAPABILITY TLV which violates these rules, then the PCEP speaker MUST send a PCerr message with Error-Type = 10 (Reception of an invalid object) and Error-Value = 11 (Malformed object) and MUST close the PCEP session. The PCEP speaker MAY include the malformed OPEN object in the PCerr message as well.

If a PCEP speaker receives an OPEN object with more than one PATH-SETUP-TYPE-CAPABILITY TLV then it MUST ignore all but the first instance of this TLV.

The absence of the PATH-SETUP-TYPE-CAPABILITY TLV from the OPEN object is equivalent to a PATH-SETUP-TYPE-CAPABILITY TLV containing a single PST of 0 (RSVP-TE signaling protocol) and no sub-TLVs. A PCEP

speaker MAY omit the PATH-SETUP-TYPE-CAPABILITY TLV if the only PST it supports is RSVP-TE. If a PCEP speaker supports other PSTs besides RSVP-TE, then it SHOULD include the PATH-SETUP-TYPE-CAPABILITY TLV in its OPEN object.

If a PCEP speaker does not recognize the PATH-SETUP-TYPE-CAPABILITY TLV, it will ignore the TLV in accordance with [RFC5440].

4. Path Setup Type TLV

When a PCEP session is used to set up TE paths using different methods, the corresponding PCE and PCC must be aware of the path setup method used. That means, a PCE must be able to specify paths in the correct format and a PCC must be able to take control plane and forwarding plane actions appropriate to the PST.

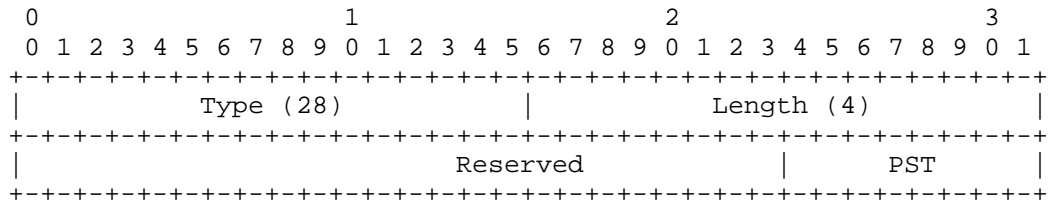


Figure 2: PATH-SETUP-TYPE TLV

PATH-SETUP-TYPE TLV is an optional TLV associated with the RP ([RFC5440]) and the SRP ([RFC8231]) objects. Its format is shown in the above figure. The TLV type is 28. Its reserved field MUST be set to zero. The one byte value contains the PST as defined for the PATH-SETUP-TYPE-CAPABILITY TLV.

The absence of the PATH-SETUP-TYPE TLV is equivalent to a PATH-SETUP-TYPE TLV with a PST value of 0 (RSVP-TE). A PCEP speaker MAY omit the TLV if the PST is RSVP-TE. If the RP or SRP object contains more than one PATH-SETUP-TYPE TLV, only the first TLV MUST be processed and the rest MUST be ignored.

If a PCEP speaker does not recognize the PATH-SETUP-TYPE TLV, it will ignore the TLV in accordance with [RFC5440], and will use RSVP-TE to set up the path.

5. Operation

During the PCEP initialization phase, if a PCEP speaker receives a PATH-SETUP-TYPE-CAPABILITY TLV from its peer, it MUST assume that the peer supports only the PSTs listed in the TLV. If the PCEP speaker and its peer have no PSTs in common, then the PCEP speaker MUST send

a PCErr message with Error-Type = 21 (Invalid traffic engineering path setup type) and Error-Value = 2 (Mismatched path setup type) and close the PCEP session.

If the peer has sent no PATH-SETUP-TYPE-CAPABILITY TLV, then the PCEP speaker MUST infer that the peer supports path setup using at least RSVP-TE. The PCEP speaker MAY also infer that the peer supports other path setup types, but the means of inference are outside the scope of this document.

When a PCC sends a PCReq message to a PCE ([RFC5440]), it MUST include the PATH-SETUP-TYPE TLV in the RP object, unless the intended PST is RSVP-TE, in which case it MAY omit the PATH-SETUP-TYPE TLV. If the PCE is capable of expressing the path in a format appropriate to the intended PST, it MUST use the appropriate ERO format in the PCRep message.

When a PCE sends a PCRep message to a PCC ([RFC5440]), it MUST include the PATH-SETUP-TYPE TLV in the RP object, unless the PST is RSVP-TE, in which case it MAY omit the PATH-SETUP-TYPE TLV. If the PCE does not support the intended PST, it MUST send a PCErr message with Error-Type = 21 (Invalid traffic engineering path setup type) and Error-Value = 1 (Unsupported path setup type) and close the PCEP session. If the PSTs corresponding to the PCReq and PCRep messages do not match, the PCC MUST send a PCErr message with Error-Type = 21 (Invalid traffic engineering path setup type) and Error-Value = 2 (Mismatched path setup type) and close the PCEP session.

When a stateful PCE sends a PCUpd message ([RFC8231]) or a PCInitiate message ([RFC8281]) to a PCC, it MUST include the PATH-SETUP-TYPE TLV in the SRP object, unless the intended PST is RSVP-TE, in which case it MAY omit the PATH-SETUP-TYPE TLV. If the PCC does not support the PST associated with the PCUpd or PCInitiate message, it MUST send a PCErr message with Error-Type = 21 (Invalid traffic engineering path setup type) and Error-Value = 1 (Unsupported path setup type) and close the PCEP session.

When a PCC sends a PCRpt message to a stateful PCE ([RFC8231]), it MUST include the PATH-SETUP-TYPE TLV in the SRP object, unless the PST is RSVP-TE, in which case it MAY omit the PATH-SETUP-TYPE TLV. The PCC MUST include the SRP object in the PCRpt message if the PST is not RSVP-TE, even when the SRP-ID-number is the reserved value of 0x00000000. If the PCRpt message is triggered by a PCUpd or PCInitiate message, then the PST that the PCC indicates in the PCRpt MUST match the PST that the stateful PCE intended in the PCUpd or PCInitiate. If it does not, then the PCE MUST send a PCErr message with Error-Type = 21 (Invalid traffic engineering path setup type)

and Error-Value = 2 (Mismatched path setup type) and close the PCEP session.

6. Manageability Considerations

This document generalises PCEP to allow path setup methods other than RSVP-TE to be used by the network (but does not define any new path setup types, besides RSVP-TE). It is possible that, in a given network, multiple path setup methods will be used. It is also possible that not all devices will support the same set of path setup methods. Managing networks that combine multiple path setup methods may therefore raise some challenges from a configuration and observability point of view.

Each document that defines a new Path Setup Type in the Path Setup Type Registry (Section 8.2) must include a manageability section. The manageability section must explain how operators can manage PCEP with the new path setup type. It must address the following questions, which are generally applicable when working with multiple path setup types in PCEP.

- o What are the criteria for when devices will use the new path setup type in PCEP, and how can the operator control this?
- o How can the network be migrated to the new path setup type, and are there any backwards compatibility issues that operators need to be aware of?
- o Are paths set up using the new path setup type intended to coexist with other paths over the long term and, if so, how is this situation managed with PCEP?
- o How can operators verify the correct operation of PCEP in the network with respect to the new path setup type? Which fault conditions must be reported to the operators?
- o Are there any existing management interfaces (such as YANG models) that must be extended to model the operation of PCEP in the network with respect to the new path setup type?

See [RFC5706] for further guidance on how to write manageability sections in standards-track documents.

7. Security Considerations

The security considerations described in [RFC5440] and [RFC8281] are applicable to this specification. No additional security measure is required.

Note that, if the security mechanisms of [RFC5440] and [RFC8281] are not used, then the protocol described by this draft could be attacked in the following new way. An attacker, using a TCP man-in-the-middle attack, could inject error messages into the PCEP session when a particular PST is (or is not) used. By doing so, the attacker could potentially force the use of a specific PST, which may allow them to subsequently attack a weakness in that PST.

8. IANA Considerations

8.1. PCEP TLV Type Indicators

IANA is requested to confirm the early allocation of the following code point in the PCEP TLV Type Indicators registry.

Value	Description	Reference
28	PATH-SETUP-TYPE	This document

IANA is requested to allocate a new code point for the following TLV in the PCEP TLV Type Indicators registry.

Value	Description	Reference
TBD1	PATH-SETUP-TYPE-CAPABILITY	This document

Note to IANA: The above TLV type was not part of the early code point allocation that was done for this draft. It was added to the draft after the early code point allocation had taken place. Please assign a code point from the indicated registry and replace each instance of "TBD1" in this document with the allocated code point.

8.2. New Path Setup Type Registry

IANA is requested to create a new sub-registry within the "Path Computation Element Protocol (PCEP) Numbers" registry called "PCEP Path Setup Types". The allocation policy for this new registry should be by IETF Review. The new registry should contain the following value:

Value	Description	Reference
0	Path is setup using the RSVP-TE signaling protocol.	This document

8.3. PCEP-Error Object

IANA is requested to confirm the early allocation of the following code-points in the PCEP-ERROR Object Error Types and Values registry.

Error-Type	Meaning
10	Reception of an invalid object
	Error-value=11: Malformed object
Error-Type	Meaning
21	Invalid traffic engineering path setup type
	Error-value=0: Unassigned
	Error-value=1: Unsupported path setup type
	Error-value=2: Mismatched path setup type

Note to IANA: the early allocation for Error-Type=10, Error-value=11 was originally done by draft-ietf-pce-segment-routing. However, we have since moved its definition into this document. Therefore, please update the reference for this Error-value in the indicated registry to point to RFC.ietf-pce-lsp-setup-type.

9. Contributors

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10. Acknowledgements

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PCEP Extensions for Segment Routing
draft-ietf-pce-segment-routing-16

Abstract

Segment Routing (SR) enables any head-end node to select any path without relying on a hop-by-hop signaling technique (e.g., LDP or RSVP-TE). It depends only on "segments" that are advertised by link-state Interior Gateway Protocols (IGPs). A Segment Routing Path can be derived from a variety of mechanisms, including an IGP Shortest Path Tree (SPT), explicit configuration, or a Path Computation Element (PCE). This document specifies extensions to the Path Computation Element Communication Protocol (PCEP) that allow a stateful PCE to compute and initiate Traffic Engineering (TE) paths, as well as a PCC to request a path subject to certain constraints and optimization criteria in SR networks.

Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

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1. Introduction

Segment Routing (SR) leverages the source routing paradigm. Using SR, a source node steers a packet through a path without relying on hop-by-hop signaling protocols such as LDP or RSVP-TE. Each path is specified as an ordered list of instructions called "segments". Each segment is an instruction to route the packet to a specific place in the network, or to perform a function on the packet. A database of segments can be distributed through the network using a routing protocol (such as IS-IS or OSPF) or by any other means. Several types of segment are defined. A node segment uniquely identifies a specific node in the SR domain. Each router in the SR domain associates a node segment with an ECMP-aware shortest path to the node that it identifies. An adjacency segment represents a unidirectional adjacency. An adjacency segment is local to the node which advertises it. Both node segments and adjacency segments can be used for SR.

[RFC8402] describes the SR architecture. The corresponding IS-IS and OSPF extensions are specified in [I-D.ietf-isis-segment-routing-extensions] and [I-D.ietf-ospf-segment-routing-extensions], respectively.

The SR architecture can be implemented using either an MPLS forwarding plane [I-D.ietf-spring-segment-routing-mpls] or an IPv6 forwarding plane [I-D.ietf-6man-segment-routing-header]. The MPLS forwarding plane can be applied to SR without any change, in which case an SR path corresponds to an MPLS Label Switching Path (LSP). This document is relevant to the MPLS forwarding plane only. In this

document, "Node-SID" and "Adjacency-SID" denote Node Segment Identifier and Adjacency Segment Identifier respectively.

A Segment Routing path (SR path) can be derived from an IGP Shortest Path Tree (SPT). SR-TE paths may not follow an IGP SPT. Such paths may be chosen by a suitable network planning tool and provisioned on the ingress node of the SR-TE path.

[RFC5440] describes the Path Computation Element Communication Protocol (PCEP) for communication between a Path Computation Client (PCC) and a Path Computation Element (PCE) or between a pair of PCEs. A PCE computes paths for MPLS Traffic Engineering LSPs (MPLS-TE LSPs) based on various constraints and optimization criteria. [RFC8231] specifies extensions to PCEP that allow a stateful PCE to compute and recommend network paths in compliance with [RFC4657] and defines objects and TLVs for MPLS-TE LSPs. Stateful PCEP extensions provide synchronization of LSP state between a PCC and a PCE or between a pair of PCEs, delegation of LSP control, reporting of LSP state from a PCC to a PCE, controlling the setup and path routing of an LSP from a PCE to a PCC. Stateful PCEP extensions are intended for an operational model in which LSPs are configured on the PCC, and control over them is delegated to the PCE.

A mechanism to dynamically initiate LSPs on a PCC based on the requests from a stateful PCE or a controller using stateful PCE is specified in [RFC8281]. This mechanism is useful in Software Defined Networking (SDN) applications, such as on-demand engineering, or bandwidth calendaring [RFC8413].

It is possible to use a stateful PCE for computing one or more SR-TE paths taking into account various constraints and objective functions. Once a path is chosen, the stateful PCE can initiate an SR-TE path on a PCC using PCEP extensions specified in [RFC8281] using the SR specific PCEP extensions specified in this document. Additionally, using procedures described in this document, a PCC can request an SR path from either a stateful or a stateless PCE.

This specification relies on the procedures specified in [RFC8408] to exchange the segment routing capability and to specify that the path setup type of an LSP is segment routing. This specification also updates [RFC8408] to clarify the use of sub-TLVs in the PATH-SETUP-TYPE-CAPABILITY TLV. See Section 4.1.1 for details.

This specification provides a mechanism for a network controller (acting as a PCE) to instantiate candidate paths for an SR Policy onto a head-end node (acting as a PCC) using PCEP. For more information on the SR Policy Architecture, see [I-D.ietf-spring-segment-routing-policy].

2. Terminology

The following terminologies are used in this document:

ERO: Explicit Route Object

IGP: Interior Gateway Protocol

IS-IS: Intermediate System to Intermediate System

LSR: Label Switching Router

MSD: Base MPLS Imposition Maximum SID Depth, as defined in [RFC8491]

NAI: Node or Adjacency Identifier

OSPF: Open Shortest Path First

PCC: Path Computation Client

PCE: Path Computation Element

PCEP: Path Computation Element Communication Protocol

RRO: Record Route Object

SID: Segment Identifier

SR: Segment Routing

SR-DB: Segment Routing Database: the collection of SRGBs, SRLBs and SIDs and the objects they map to, advertised by a link state IGP

SRGB: Segment Routing Global Block

SRLB: Segment Routing Local Block

SR-TE: Segment Routing Traffic Engineering

3. Overview of PCEP Operation in SR Networks

In an SR network, the ingress node of an SR path prepends an SR header to all outgoing packets. The SR header consists of a list of SIDs (or MPLS labels in the context of this document). The header has all necessary information so that, in combination with the information distributed by the IGP, the packets can be guided from the ingress node to the egress node of the path; hence, there is no need for any signaling protocol.

In PCEP messages, LSP route information is carried in the Explicit Route Object (ERO), which consists of a sequence of subobjects. SR-TE paths computed by a PCE can be represented in an ERO in one of the following forms:

- o An ordered set of IP addresses representing network nodes/links.
- o An ordered set of SIDs, with or without the corresponding IP addresses.
- o An ordered set of MPLS labels, with or without corresponding IP address.

The PCC converts these into an MPLS label stack and next hop, as described in Section 5.2.2.

This document defines a new ERO subobject denoted by "SR-ERO subobject" capable of carrying a SID as well as the identity of the node/adjacency represented by the SID. SR-capable PCEP speakers should be able to generate and/or process such ERO subobject. An ERO containing SR-ERO subobjects can be included in the PCEP Path Computation Reply (PCRep) message defined in [RFC5440], the PCEP LSP Initiate Request message (PCInitiate) defined in [RFC8281], as well as in the PCEP LSP Update Request (PCUpd) and PCEP LSP State Report (PCRpt) messages defined in [RFC8231].

When a PCEP session between a PCC and a PCE is established, both PCEP speakers exchange their capabilities to indicate their ability to support SR-specific functionality.

A PCE can update an LSP that is initially established via RSVP-TE signaling to use an SR-TE path, by sending a PCUpd to the PCC that delegated the LSP to it ([RFC8231]). A PCC can update an undelegated LSP that is initially established via RSVP-TE signaling to use an SR-TE path as follows. First, it requests an SR-TE Path from a PCE by sending a PCReq message. If it receives a suitable path, it establishes the path in the data plane, and then tears down the original RSVP-TE path. If the PCE is stateful, then the PCC sends PCRpt messages indicating that the new path is set up and the old path is torn down, per [RFC8231].

Similarly, a PCE or PCC can update an LSP initially created with an SR-TE path to use RSVP-TE signaling, if necessary. This capability is useful for rolling back a change when a network is migrated from RSVP-TE to SR-TE technology.

A PCC MAY include an RRO containing the recorded LSP in PCReq and PCRpt messages as specified in [RFC5440] and [RFC8231], respectively.

This document defines a new RRO subobject for SR networks. The methods used by a PCC to record the SR-TE LSP are outside the scope of this document.

In summary, this document:

- o Defines a new ERO subobject, a new RRO subobject and new PCEP error codes.
- o Specifies how two PCEP speakers can establish a PCEP session that can carry information about SR-TE paths.
- o Specifies processing rules for the ERO subobject.
- o Defines a new path setup type to be used in the PATH-SETUP-TYPE and PATH-SETUP-TYPE-CAPABILITY TLVs ([RFC8408]).
- o Defines a new sub-TLV for the PATH-SETUP-TYPE-CAPABILITY TLV.

The extensions specified in this document complement the existing PCEP specifications to support SR-TE paths. As such, the PCEP messages (e.g., Path Computation Request, Path Computation Reply, Path Computation Report, Path Computation Update, Path Computation Initiate, etc.,) are formatted according to [RFC5440], [RFC8231], [RFC8281], and any other applicable PCEP specifications.

4. Object Formats

4.1. The OPEN Object

4.1.1. The Path Setup Type Capability TLV

[RFC8408] defines the PATH-SETUP-TYPE-CAPABILITY TLV for use in the OPEN object. The PATH-SETUP-TYPE-CAPABILITY TLV contains an optional list of sub-TLVs which are intended to convey parameters that are associated with the path setup types supported by a PCEP speaker.

This specification updates [RFC8408], as follows. It creates a new registry which defines the valid type indicators of the sub-TLVs of the PATH-SETUP-TYPE-CAPABILITY TLV (see Section 8.6). A PCEP speaker MUST NOT include a sub-TLV in the PATH-SETUP-TYPE-CAPABILITY TLV unless it appears in this registry. If a PCEP speaker receives a sub-TLV whose type indicator does not match one of those from the registry, or else is not recognised by the speaker, then the speaker MUST ignore the sub-TLV.

4.1.2. The SR PCE Capability sub-TLV

This document defines a new Path Setup Type (PST) for SR, as follows:

- o PST = 1: Path is setup using Segment Routing Traffic Engineering.

A PCEP speaker SHOULD indicate its support of the function described in this document by sending a PATH-SETUP-TYPE-CAPABILITY TLV in the OPEN object with this new PST included in the PST list.

This document also defines the SR-PCE-CAPABILITY sub-TLV. PCEP speakers use this sub-TLV to exchange information about their SR capability. If a PCEP speaker includes PST=1 in the PST List of the PATH-SETUP-TYPE-CAPABILITY TLV then it MUST also include the SR-PCE-CAPABILITY sub-TLV inside the PATH-SETUP-TYPE-CAPABILITY TLV.

The format of the SR-PCE-CAPABILITY sub-TLV is shown in the following figure:

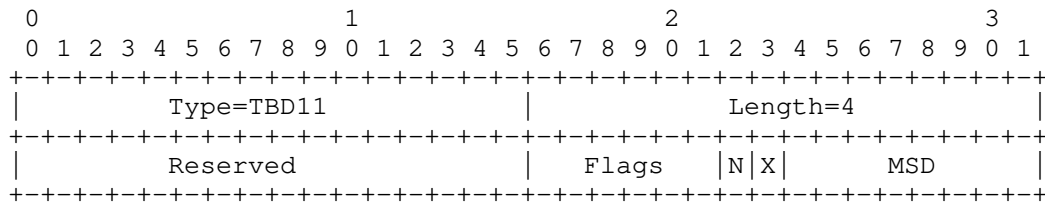


Figure 1: SR-PCE-CAPABILITY sub-TLV format

The code point for the TLV type is TBD11. The TLV length is 4 octets.

The 32-bit value is formatted as follows.

Reserved: MUST be set to zero by the sender and MUST be ignored by the receiver.

Flags: This document defines the following flag bits. The other bits MUST be set to zero by the sender and MUST be ignored by the receiver.

- * N: A PCC sets this flag bit to 1 to indicate that it is capable of resolving a Node or Adjacency Identifier (NAI) to a SID.
- * X: A PCC sets this flag bit to 1 to indicate that it does not impose any limit on the MSD.

Maximum SID Depth (MSD): specifies the maximum number of SIDs (MPLS label stack depth in the context of this document) that a PCC is capable of imposing on a packet. Section 5.1 explains the relationship between this field and the X flag.

4.2. The RP/SRP Object

To set up an SR-TE LSP using SR, the RP (Request Parameters) or SRP (Stateful PCE Request Parameters) object MUST include the PATH-SETUP-TYPE TLV, specified in [RFC8408], with the PST set to 1 (path setup using SR-TE).

The LSP-IDENTIFIERS TLV MAY be present for the above PST type.

4.3. ERO

An SR-TE path consists of one or more SIDs where each SID MAY be associated with the identifier that represents the node or adjacency corresponding to the SID. This identifier is referred to as the 'Node or Adjacency Identifier' (NAI). As described later, a NAI can be represented in various formats (e.g., IPv4 address, IPv6 address, etc). Furthermore, a NAI is used for troubleshooting purposes and, if necessary, to derive SID value as described below.

The ERO specified in [RFC5440] is used to carry SR-TE path information. In order to carry SID and/or NAI, this document defines a new ERO subobject referred to as "SR-ERO subobject" whose format is specified in the following section. An ERO carrying an SR-TE path consists of one or more ERO subobjects, and MUST carry only SR-ERO subobjects. Note that an SR-ERO subobject does not need to have both SID and NAI. However, at least one of them MUST be present.

When building the MPLS label stack from ERO, a PCC MUST assume that SR-ERO subobjects are organized as a last-in-first-out stack. The first subobject relative to the beginning of ERO contains the information about the topmost label. The last subobject contains information about the bottommost label.

4.3.1. SR-ERO Subobject

An SR-ERO subobject is formatted as shown in the following diagram.

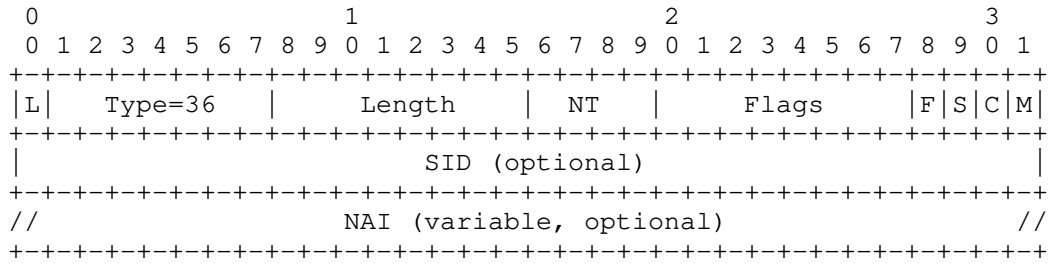


Figure 2: SR-ERO subobject format

The fields in the SR-ERO Subobject are as follows:

The 'L' Flag: Indicates whether the subobject represents a loose-hop in the LSP [RFC3209]. If this flag is set to zero, a PCC MUST NOT overwrite the SID value present in the SR-ERO subobject. Otherwise, a PCC MAY expand or replace one or more SID values in the received SR-ERO based on its local policy.

Type: Set to 36.

Length: Contains the total length of the subobject in octets. The Length MUST be at least 8, and MUST be a multiple of 4. An SR-ERO subobject MUST contain at least one of a SID or an NAI. The flags described below indicate whether the SID or NAI fields are absent.

NAI Type (NT): Indicates the type and format of the NAI contained in the object body, if any is present. If the F bit is set to zero (see below) then the NT field has no meaning and MUST be ignored by the receiver. This document describes the following NT values:

- NT=0 The NAI is absent.
- NT=1 The NAI is an IPv4 node ID.
- NT=2 The NAI is an IPv6 node ID.
- NT=3 The NAI is an IPv4 adjacency.
- NT=4 The NAI is an IPv6 adjacency with global IPv6 addresses.
- NT=5 The NAI is an unnumbered adjacency with IPv4 node IDs.
- NT=6 The NAI is an IPv6 adjacency with link-local IPv6 addresses.

Flags: Used to carry additional information pertaining to the SID. This document defines the following flag bits. The other bits MUST be set to zero by the sender and MUST be ignored by the receiver.

- * M: If this bit is set to 1, the SID value represents an MPLS label stack entry as specified in [RFC3032]. Otherwise, the SID value is an administratively configured value which represents an index into an MPLS label space (either SRGB or SRLB) per [RFC8402].
- * C: If the M bit and the C bit are both set to 1, then the TC, S, and TTL fields in the MPLS label stack entry are specified by the PCE. However, a PCC MAY choose to override these values according its local policy and MPLS forwarding rules. If the M bit is set to 1 but the C bit is set to zero, then the TC, S, and TTL fields MUST be ignored by the PCC. The PCC MUST set these fields according to its local policy and MPLS forwarding rules. If the M bit is set to zero then the C bit MUST be set to zero.
- * S: When this bit is set to 1, the SID value in the subobject body is absent. In this case, the PCC is responsible for choosing the SID value, e.g., by looking up in the SR-DB using the NAI which, in this case, MUST be present in the subobject. If the S bit is set to 1 then the M and C bits MUST be set to zero.
- * F: When this bit is set to 1, the NAI value in the subobject body is absent. The F bit MUST be set to 1 if NT=0, and otherwise MUST be set to zero. The S and F bits MUST NOT both be set to 1.

SID: The Segment Identifier. Depending on the M bit, it contains either:

- * A 4 octet index defining the offset into an MPLS label space per [RFC8402].
- * A 4 octet MPLS Label Stack Entry, where the 20 most significant bits encode the label value per [RFC3032].

NAI: The NAI associated with the SID. The NAI's format depends on the value in the NT field, and is described in the following section.

At least one of the SID and the NAI MUST be included in the SR-ERO subobject, and both MAY be included.

4.3.2. NAI Associated with SID

This document defines the following NAIs:

'IPv4 Node ID' is specified as an IPv4 address. In this case, the NT value is 1 and the NAI field length is 4 octets.

'IPv6 Node ID' is specified as an IPv6 address. In this case, the NT value is 2 and the NAI field length is 16 octets.

'IPv4 Adjacency' is specified as a pair of IPv4 addresses. In this case, the NT value is 3 and the NAI field length is 8 octets. The format of the NAI is shown in the following figure:

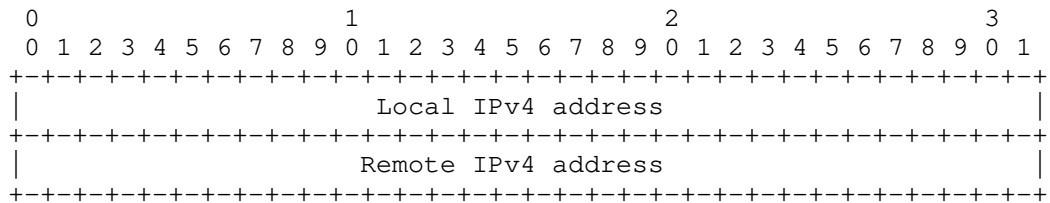


Figure 3: NAI for IPv4 adjacency

'IPv6 Global Adjacency' is specified as a pair of global IPv6 addresses. It is used to describe an IPv6 adjacency for a link that uses global IPv6 addresses. Each global IPv6 address is configured on a specific router interface, so together they identify an adjacency between a pair of routers. In this case, the NT value is 4 and the NAI field length is 32 octets. The format of the NAI is shown in the following figure:

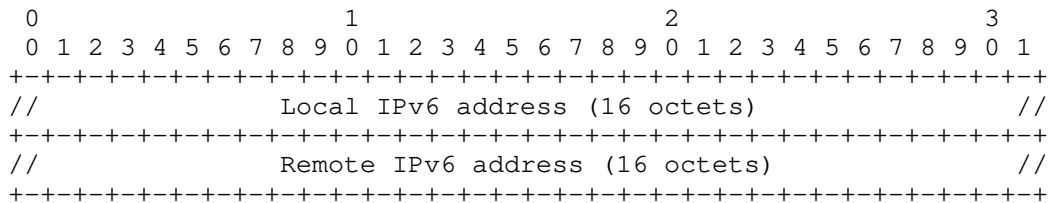


Figure 4: NAI for IPv6 global adjacency

'Unnumbered Adjacency with IPv4 NodeIDs' is specified as a pair of (node ID, interface ID) tuples. In this case, the NT value is 5 and the NAI field length is 16 octets. The format of the NAI is shown in the following figure:

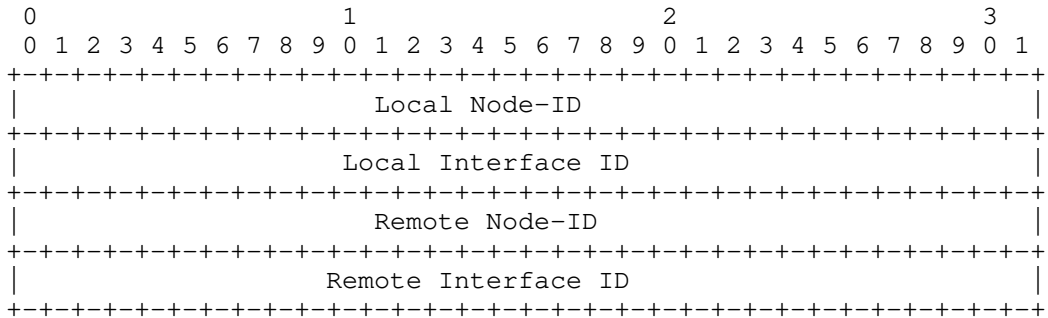


Figure 5: NAI for Unnumbered adjacency with IPv4 Node IDs

'IPv6 Link-Local Adjacency' is specified as a pair of (global IPv6 address, interface ID) tuples. It is used to describe an IPv6 adjacency for a link that uses only link local IPv6 addresses. Each global IPv6 address is configured on a specific router, so together they identify a pair of adjacent routers. The interface IDs identify the link that the adjacency is formed over. In this case, the NT value is 6 and the NAI field length is 40 octets. The format of the NAI is shown in the following figure:

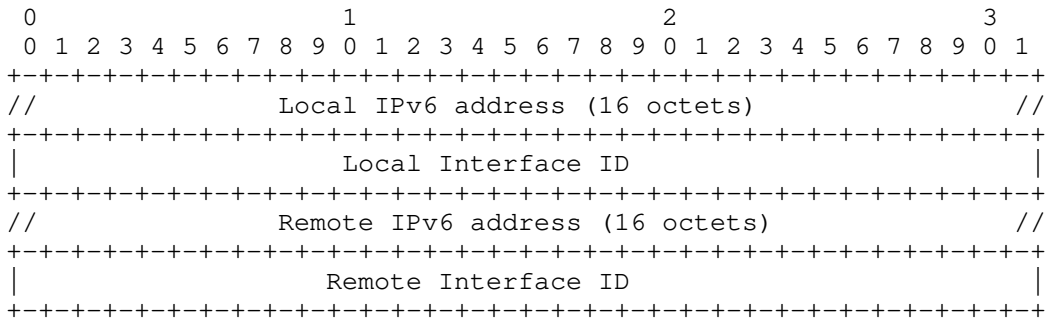


Figure 6: NAI for IPv6 link-local adjacency

4.4. RRO

A PCC reports an SR-TE LSP to a PCE by sending a PCRpt message, per [RFC8231]. The RRO on this message represents the SID list that was applied by the PCC, that is, the actual path taken by the LSP. The procedures of [RFC8231] with respect to the RRO apply equally to this specification without change.

An RRO contains one or more subobjects called "SR-RRO subobjects" whose format is shown below:

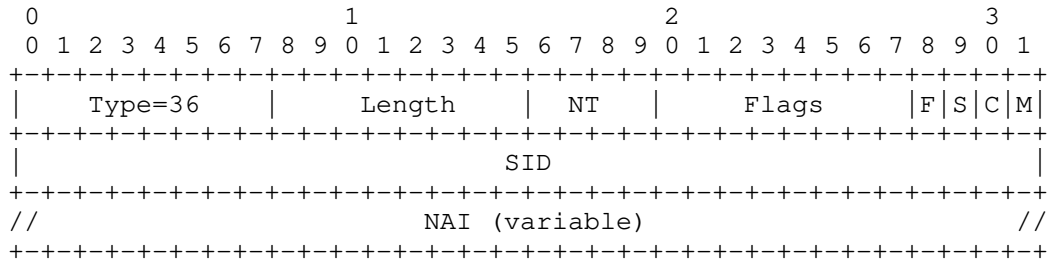


Figure 7: SR-RRO Subobject format

The format of the SR-RRO subobject is the same as that of the SR-ERO subobject, but without the L flag.

A PCC MUST order the SR-RRO subobjects such that the first subobject relative to the beginning of the RRO identifies the first segment visited by the SR-TE LSP, and the last subobject identifies the final segment of the SR-TE LSP, that is, its endpoint.

4.5. METRIC Object

A PCC MAY request that PCE optimizes an individual path computation request to minimize the SID depth of the computed path by using the METRIC object defined in [RFC5440]. This document defines a new type for the METRIC object to be used for this purpose, as follows:

- o T = 11: Maximum SID Depth of the requested path.

If the PCC includes a METRIC object of this type on a path computation request, then the PCE minimizes the SID depth of the computed path. If the B (bound) bit is set to to 1 in the METRIC object, then the PCE MUST NOT return a path whose SID depth exceeds the given metric-value. If the PCC did not set the X flag in its SR-PCE-CAPABILITY TLV, then it MUST set the B bit to 1. If the PCC set the X flag in its SR-PCE-CAPABILITY TLV, then it MAY set the B bit to 1 or zero.

If a PCEP session is established with a non-zero default MSD value, then the PCC MUST NOT send an MSD METRIC object with an MSD greater than the session's default MSD. If the PCE receives a path computation request with an MSD METRIC object on such a session that is greater than the session's default MSD, then it MUST consider the request invalid and send a PCerr with Error-Type = 10 ("Reception of an invalid object") and Error-Value 9 ("MSD exceeds the default for the PCEP session").

5. Procedures

5.1. Exchanging the SR PCE Capability

A PCC indicates that it is capable of supporting the head-end functions for SR-TE LSP by including the SR-PCE-CAPABILITY sub-TLV in the Open message that it sends to a PCE. A PCE indicates that it is capable of computing SR-TE paths by including the SR-PCE-CAPABILITY sub-TLV in the Open message that it sends to a PCC.

If a PCEP speaker receives a PATH-SETUP-TYPE-CAPABILITY TLV with a PST list containing PST=1, and supports that path setup type, then it checks for the presence of the SR-PCE-CAPABILITY sub-TLV. If that sub-TLV is absent, then the PCEP speaker MUST send a PCErr message with Error-Type 10 (Reception of an invalid object) and Error-Value TBD1 (Missing PCE-SR-CAPABILITY sub-TLV) and MUST then close the PCEP session. If a PCEP speaker receives a PATH-SETUP-TYPE-CAPABILITY TLV with a SR-PCE-CAPABILITY sub-TLV, but the PST list does not contain PST=1, then the PCEP speaker MUST ignore the SR-PCE-CAPABILITY sub-TLV.

If a PCC sets the N flag to 1, then the PCE MAY send an SR-ERO subobject containing NAI and no SID (see Section 5.2). Otherwise, the PCE MUST NOT send an SR-ERO subobject containing NAI and no SID.

The number of SIDs that can be imposed on a packet depends on the PCC's data plane's capability. If a PCC sets the X flag to 1 then the MSD is not used and MUST be set to zero. If a PCE receives an SR-PCE-CAPABILITY sub-TLV with the X flag set to 1 then it MUST ignore the MSD field and assumes that the sender can impose a SID stack of any depth. If a PCC sets the X flag to zero, then it sets the MSD field to the maximum number of SIDs that it can impose on a packet. In this case, the PCC MUST set the MSD to a number greater than zero. If a PCE receives an SR-PCE-CAPABILITY sub-TLV with the X flag and MSD both set to zero then it MUST send a PCErr message with Error-Type 10 (Reception of an invalid object) and Error-Value TBD10 (Maximum SID depth must be nonzero) and MUST then close the PCEP session.

Note that the MSD value exchanged via the SR-PCE-CAPABILITY sub-TLV indicates the SID/label imposition limit for the PCC node. It is anticipated that, in many deployments, the PCCs will have network interfaces that are homogeneous with respect to MSD (that is, each interface has the same MSD). In such cases, having a per-node MSD on the PCEP session is sufficient; the PCE SHOULD interpret this to mean that all network interfaces on the PCC have the given MSD. However, the PCE MAY also learn a per-node MSD and a per-interface MSD from the routing protocols, as specified in: [RFC8491]; [RFC8476];

[I-D.ietf-idr-bgp-ls-segment-routing-msd]. If the PCE learns the per-node MSD of a PCC from a routing protocol, then it MUST ignore the per-node MSD value in the SR-PCE-CAPABILITY sub-TLV and use the per-node MSD learned from the routing protocol instead. If the PCE learns the MSD of a network interface on a PCC from a routing protocol, then it MUST use the per-interface MSD instead of the MSD value in the SR-PCE-CAPABILITY sub-TLV when it computes a path that uses that interface.

Once an SR-capable PCEP session is established with a non-zero MSD value, the corresponding PCE MUST NOT send SR-TE paths with a number of SIDs exceeding that MSD value. If a PCC needs to modify the MSD value, it MUST close the PCEP session and re-establish it with the new MSD value. If a PCEP session is established with a non-zero MSD value, and the PCC receives an SR-TE path containing more SIDs than specified in the MSD value, the PCC MUST send a PCERR message with Error-Type 10 (Reception of an invalid object) and Error-Value 3 (Unsupported number of Segment ERO subobjects). If a PCEP session is established with an MSD value of zero, then the PCC MAY specify an MSD for each path computation request that it sends to the PCE, by including a "maximum SID depth" metric object on the request, as defined in Section 4.5.

The N flag, X flag and MSD value inside the SR-PCE-CAPABILITY sub-TLV are meaningful only in the Open message sent from a PCC to a PCE. As such, a PCE MUST set the N flag to zero, the X flag to 1 and MSD value to zero in an outbound message to a PCC. Similarly, a PCC MUST ignore any MSD value received from a PCE. If a PCE receives multiple SR-PCE-CAPABILITY sub-TLVs in an Open message, it processes only the first sub-TLV received.

5.2. ERO Processing

5.2.1. SR-ERO Validation

If a PCC does not support the SR PCE Capability and thus cannot recognize the SR-ERO or SR-RRR subobjects, it will respond according to the rules for a malformed object per [RFC5440].

On receiving an SR-ERO, a PCC MUST validate that the Length field, the S bit, the F bit and the NT field are consistent, as follows.

- o If NT=0, the F bit MUST be 1, the S bit MUST be zero and the Length MUST be 8.
- o If NT=1, the F bit MUST be zero. If the S bit is 1, the Length MUST be 8, otherwise the Length MUST be 12.

- o If NT=2, the F bit MUST be zero. If the S bit is 1, the Length MUST be 20, otherwise the Length MUST be 24.
- o If NT=3, the F bit MUST be zero. If the S bit is 1, the Length MUST be 12, otherwise the Length MUST be 16.
- o If NT=4, the F bit MUST be zero. If the S bit is 1, the Length MUST be 36, otherwise the Length MUST be 40.
- o If NT=5, the F bit MUST be zero. If the S bit is 1, the Length MUST be 20, otherwise the Length MUST be 24.
- o If NT=6, the F bit MUST be zero. If the S bit is 1, the Length MUST be 44, otherwise the Length MUST be 48.

If a PCC finds that the NT field, Length field, S bit and F bit are not consistent, it MUST consider the entire ERO invalid and MUST send a PCerr message with Error-Type = 10 ("Reception of an invalid object") and Error-Value = 11 ("Malformed object").

If a PCC does not recognise or support the value in the NT field, it MUST consider the entire ERO invalid and MUST send a PCerr message with Error-Type = 10 ("Reception of an invalid object") and Error-Value = TBD2 ("Unsupported NAI Type in Segment ERO subobject").

If a PCC receives an SR-ERO subobject in which the S and F bits are both set to 1 (that is, both the SID and NAI are absent), it MUST consider the entire ERO invalid and send a PCerr message with Error-Type = 10 ("Reception of an invalid object") and Error-Value = 6 ("Both SID and NAI are absent in SR-ERO subobject").

If a PCC receives an SR-ERO subobject in which the S bit is set to 1 and the F bit is set to zero (that is, the SID is absent and the NAI is present), but the PCC does not support NAI resolution, it MUST consider the entire ERO invalid and send a PCerr message with Error-Type = 4 ("Not supported object") and Error-Value = 4 ("Unsupported parameter").

If a PCC receives an SR-ERO subobject in which the S bit is set to 1 and either or both of the M or C bits is set to 1, it MUST consider the entire ERO invalid and send a PCerr message with Error-Type = 10 ("Reception of an invalid object") and Error-Value = 11 ("Malformed object").

If a PCC receives an SR-ERO subobject in which the S bit is set to zero and the M bit is set to 1, then the subobject contains an MPLS label. The PCC MAY choose not to accept a label provided by the PCE, based on its local policy. The PCC MUST NOT accept MPLS label value 3

(Implicit NULL), but it MAY accept other special purpose MPLS label values. If the PCC decides not to accept an MPLS label value, it MUST send a PCErr message with Error-Type = 10 ("Reception of an invalid object") and Error Value = 2 ("Bad label value").

If both M and C bits of an SR-ERO subobject are set to 1, and if a PCC finds erroneous setting in one or more of TC, S, and TTL fields, it MAY overwrite those fields with values chosen according to its own policy. If the PCC does not overwrite them, it MUST send a PCErr message with Error-Type = 10 ("Reception of an invalid object") and Error-Value = 4 ("Bad label format").

If the M bit of an SR-ERO subobject is set to zero but the C bit is set to 1, then the PCC MUST consider the entire ERO invalid and MUST send a PCErr message with Error-Type = 10 ("Reception of an invalid object") and Error-Value = 11 ("Malformed object").

If a PCC receives an SR-ERO subobject in which the S bit is set to zero and the M bit is set to zero, then the subobject contains a SID index value. If the SID is an Adjacency-SID then the L flag MUST NOT be set. If the L flag is set for an Adjacency-SID then the PCC MUST send a PCErr message with Error-Type = 10 ("Reception of an invalid object") and Error-Value = 11 ("Malformed object").

If a PCC detects that the subobjects of an ERO are a mixture of SR-ERO subobjects and subobjects of other types, then it MUST send a PCErr message with Error-Type = 10 ("Reception of an invalid object") and Error-Value = 5 ("ERO mixes SR-ERO subobjects with other subobject types").

The SR-ERO subobjects can be classified according to whether they contain a SID representing an MPLS label value, a SID representing an index value, or no SID. If a PCC detects that the SR-ERO subobjects are a mixture of more than one of these types, then it MUST send a PCErr message with Error-Type = 10 ("Reception of an invalid object") and Error-Value = TBD9 ("Inconsistent SIDs in SR-ERO / SR-RO subobjects").

If an ERO specifies a new SR-TE path for an existing LSP and the PCC determines that the ERO contains SR-ERO subobjects that are not valid, then the PCC MUST NOT update the LSP.

5.2.2. Interpreting the SR-ERO

The SR-ERO contains a sequence of subobjects. Each SR-ERO subobject in the sequence identifies a segment that the traffic will be directed to, in the order given. That is, the first subobject

identifies the first segment the traffic will be directed to, the second subobject represents the second segment, and so on.

The PCC interprets the SR-ERO by converting it to an MPLS label stack plus a next hop. The PCC sends packets along the segment routed path by prepending the MPLS label stack onto the packets and sending the resulting, modified packet to the next hop.

The PCC uses a different procedure to do this conversion, depending on the information that the PCE has provided in the subobjects.

- o If the subobjects contain SID index values, then the PCC converts them into the corresponding MPLS labels by following the procedure defined in [I-D.ietf-spring-segment-routing-mpls].
- o If the subobjects contain NAI only, the PCC first converts each NAI into a SID index value and then proceeds as above. To convert an NAI to a SID index, the PCC looks for a fully-specified prefix or adjacency matching the fields in the NAI. If the PCC finds a matching prefix/adjacency, and the matching prefix/adjacency has a SID associated with it, then the PCC uses that SID. If the PCC cannot find a matching prefix/adjacency, or if the matching prefix/adjacency has no SID associated with it, the PCC behaves as specified in Section 5.2.2.1.
- o If the subobjects contain MPLS labels, then the PCC looks up the offset of the first subobject's label in its SRGB or SRLB. This gives the first SID. The PCC pushes the labels in any remaining subobjects onto the packet (with the final subobject specifying the bottom-of-stack label).

For all cases above, after the PCC has imposed the label stack on the packet, it sends the packet to the segment identified by the first SID.

5.2.2.1. Handling Errors During SR-ERO Conversion

There are several errors that can occur during the process of converting an SR-ERO sequence to an MPLS label stack and a next hop. The PCC deals with them as follows.

- o If the PCC cannot find a SID index in the SR-DB, it MUST send a PCErr message with Error-Type = 10 ("Reception of an invalid object") and Error-Value = TBD3 ("Unknown SID").
- o If the PCC cannot find an NAI in the SR-DB, it MUST send a PCErr message with Error-Type = 10 ("Reception of an invalid object") and Error-Value = TBD4 ("NAI cannot be resolved to a SID").

- o If the PCC needs to convert a SID into an MPLS label value but cannot find the corresponding router's SRGB in the SR-DB, it MUST send a PCErr message with Error-Type = 10 ("Reception of an invalid object") and Error-Value = TBD5 ("Could not find SRGB").
- o If the PCC finds that a router's SRGB is not large enough for a SID index value, it MUST send a PCErr message with Error-Type = 10 ("Reception of an invalid object") and Error-Value = TBD6 ("SID index exceeds SRGB size").
- o If the PCC needs to convert a SID into an MPLS label value but cannot find the corresponding router's SRLB in the SR-DB, it MUST send a PCErr message with Error-Type = 10 ("Reception of an invalid object") and Error-Value = TBD7 ("Could not find SRLB").
- o If the PCC finds that a router's SRLB is not large enough for a SID index value, it MUST send a PCErr message with Error-Type = 10 ("Reception of an invalid object") and Error-Value = TBD8 ("SID index exceeds SRLB size").
- o If the number of labels in the computed label stack exceeds the maximum number of SIDs that the PCC can impose on the packet, it MUST send a PCErr message with Error-Type = 10 ("Reception of an invalid object") and Error-Value = 3 ("Unsupported number of Segment ERO subobjects").

If an ERO specifies a new SR-TE path for an existing LSP and the PCC encounters an error while processing the ERO, then the PCC MUST NOT update the LSP.

5.3. RRO Processing

The syntax checking rules that apply to the SR-RRO subobject are identical to those of the SR-ERO subobject, except as noted below.

If a PCEP speaker receives an SR-RRO subobject in which both SID and NAI are absent, it MUST consider the entire RRO invalid and send a PCErr message with Error-Type = 10 ("Reception of an invalid object") and Error-Value = 7 ("Both SID and NAI are absent in SR-RRO subobject").

If a PCE detects that the subobjects of an RRO are a mixture of SR-RRO subobjects and subobjects of other types, then it MUST send a PCErr message with Error-Type = 10 ("Reception of an invalid object") and Error-Value = 10 ("RRO mixes SR-RRO subobjects with other subobject types").

The SR-RRO subobjects can be classified according to whether they contain a SID representing an MPLS label value or a SID representing an index value, or no SID. If a PCE detects that the SR-RRO subobjects are a mixture of more than one of these types, then it MUST send a PCerr message with Error-Type = 10 ("Reception of an invalid object") and Error-Value = TBD9 ("Inconsistent SIDs in SR-ERO / SR-RRO subobjects").

6. Management Considerations

This document adds a new path setup type to PCEP to allow LSPs to be set up using segment routing techniques. This path setup type may be used with PCEP alongside other path setup types, such as RSVP-TE, or it may be used exclusively.

6.1. Controlling the Path Setup Type

The following factors control which path setup type is used for a given LSP.

- o The available path setup types are constrained to those that are supported by, or enabled on, the PCEP speakers. The PATH-SETUP-TYPE-CAPABILITY TLV indicates which path setup types a PCEP speaker supports. To use segment routing as a path setup type, it is a prerequisite that the PCC and PCE both include PST=1 in the list of supported path setup types in this TLV, and also include the SR-PCE-CAPABILITY sub-TLV.
- o When a PCE initiates an LSP, it proposes which path setup type to use by including it in the PATH-SETUP-TYPE TLV in the SRP object of the PCInitiate message. The PCE chooses the path setup type based on the capabilities of the network nodes on the path and on its local policy. The PCC MAY choose to accept the proposed path setup type, or to reject the PCInitiate request, based on its local policy.
- o When a PCC requests a path for an LSP, it can nominate a preferred path setup type by including it in the PATH-SETUP-TYPE TLV in the RP object of the PCReq message. The PCE MAY choose to reply with a path of the requested type, or to reply with a path of a different type, or to reject the request, based on the capabilities of the network nodes on the path and on its local policy.

The operator can influence the path setup type as follows.

- o Implementations MUST allow the operator to enable and disable the segment routing path setup type on a PCEP-speaking device.

Implementations MAY also allow the operator to enable and disable the RSVP-TE path setup type.

- o PCE implementations MUST allow the operator to specify that an LSP should be instantiated using segment routing or RSVP-TE as the proposed path setup type.
- o PCE implementations MAY allow the operator to configure a preference for the PCE to propose paths using segment routing or RSVP-TE in the absence of a specified path setup type.
- o PCC implementations MUST allow the operator to specify that a path requested for an LSP nominates segment routing or RSVP-TE as the path setup type.
- o PCC implementations MAY allow the operator to configure a preference for the PCC to nominate segment routing or RSVP-TE as the path setup type if none is specified for an LSP.
- o PCC implementations SHOULD allow the operator to configure a PCC to refuse to set up an LSP using an undesired path setup type.

6.2. Migrating a Network to Use PCEP Segment Routed Paths

This section discusses the steps that the operator takes when migrating a network to enable PCEP to set up paths using segment routing as the path setup type.

- o The operator enables the segment routing PST on the PCE servers.
- o The operator enables the segment routing PST on the PCCs.
- o The operator resets each PCEP session. The PCEP sessions come back up with segment routing enabled.
- o If the operator detects a problem, they can roll the network back to its initial state by disabling the segment routing PST on the PCEP speakers and resetting the PCEP sessions.

Note that the data plane is unaffected if a PCEP session is reset. Any LSPs that were set up before the session reset will remain in place and will still be present after the session comes back up.

An implementation SHOULD allow the operator to manually trigger a PCEP session to be reset.

An implementation MAY automatically reset a PCEP session when an operator reconfigures the PCEP speaker's capabilities. However, note

that if the capabilities at both ends of the PCEP session are not reconfigured simultaneously, then the session could be reset twice, which could lead to unnecessary network traffic. Therefore, such implementations SHOULD allow the operator to override this behaviour and wait instead for a manual reset.

Once segment routing is enabled on a PCEP session, it can be used as the path setup type for future LSPs.

User traffic is not automatically migrated from existing LSPs onto segment routed LSPs just by enabling the segment routing PST in PCEP. The migration of user traffic from existing LSPs onto segment routing LSPs is beyond the scope of this document.

6.3. Verification of Network Operation

The operator needs the following information to verify that PCEP is operating correctly with respect to the segment routing path setup type.

- o An implementation SHOULD allow the operator to view whether the PCEP speaker sent the segment routing PST capability to its peer. If the PCEP speaker is a PCC, then the implementation SHOULD also allow the operator to view the values of the L and N flags that were sent, and the value of the MSD field that was sent.
- o An implementation SHOULD allow the operator to view whether the peer sent the segment routing PST capability. If the peer is a PCC, then the implementation SHOULD also allow the operator to view the values of the L and N flags and MSD fields that the peer sent.
- o An implementation SHOULD allow the operator to view whether the segment routing PST is enabled on the PCEP session.
- o If one PCEP speaker advertises the segment routing PST capability, but the other does not, then the implementation SHOULD create a log to inform the operator of the capability mismatch.
- o An implementation SHOULD allow the operator to view the PST that was proposed, or requested, for an LSP, and the PST that was actually used.
- o If a PCEP speaker decides to use a different PST to the one that was proposed, or requested, for an LSP, then the implementation SHOULD create a log to inform the operator that the expected PST has not been used. The log SHOULD give the reason for this choice (local policy, equipment capability etc.)

- o If a PCEP speaker rejects a segment routing path, then it SHOULD create a log to inform the operator, giving the reason for the decision (local policy, MSD exceeded etc.)

6.4. Relationship to Existing Management Models

The PCEP YANG module is defined in [I-D.ietf-pce-pcep-yang]. In future, this YANG module should be extended or augmented to provide the following additional information relating to segment routing:

- o The advertised PST capabilities and MSD per PCEP session.
- o The PST configured for, and used by, each LSP.

The PCEP MIB [RFC7420] could also be updated to include this information.

7. Security Considerations

The security considerations described in [RFC5440], [RFC8231], [RFC8281] and [RFC8408] are applicable to this specification. No additional security measure is required.

Note that this specification enables a network controller to instantiate a path in the network without the use of a hop-by-hop signaling protocol (such as RSVP-TE). This creates an additional vulnerability if the security mechanisms of [RFC5440], [RFC8231] and [RFC8281] are not used. If there is no integrity protection on the session, then an attacker could create a path which is not subjected to the further verification checks that would be performed by the signaling protocol.

Note that this specification adds the MSD field to the OPEN message (see Section 4.1.2) which discloses how many MPLS labels the sender can push onto packets that it forwards into the network. If the security mechanisms of [RFC8231] and [RFC8281] are not used with strong encryption, then an attacker could use this new field to gain intelligence about the capabilities of the edge devices in the network.

8. IANA Considerations

8.1. PCEP ERO and RRO subobjects

This document defines a new subobject type for the PCEP explicit route object (ERO), and a new subobject type for the PCEP record route object (RRO). The code points for subobject types of these objects is maintained in the RSVP parameters registry, under the

EXPLICIT_ROUTE and ROUTE_RECORD objects. IANA is requested to confirm the early allocation of the following code points in the RSVP Parameters registry for each of the new subobject types defined in this document.

Object	Subobject	Subobject Type
EXPLICIT_ROUTE	SR-ERO (PCEP-specific)	36
ROUTE_RECORD	SR-RRO (PCEP-specific)	36

8.2. New NAI Type Registry

IANA is requested to create a new sub-registry within the "Path Computation Element Protocol (PCEP) Numbers" registry called "PCEP SR-ERO NAI Types". The allocation policy for this new registry should be by IETF Review. The new registry should contain the following values:

Value	Description	Reference
0	NAI is absent.	This document
1	NAI is an IPv4 node ID.	This document
2	NAI is an IPv6 node ID.	This document
3	NAI is an IPv4 adjacency.	This document
4	NAI is an IPv6 adjacency with global IPv6 addresses.	This document
5	NAI is an unnumbered adjacency with IPv4 node IDs.	This document
6	NAI is an IPv6 adjacency with link-local IPv6 addresses.	This document

8.3. New SR-ERO Flag Registry

IANA is requested to create a new sub-registry, named "SR-ERO Flag Field", within the "Path Computation Element Protocol (PCEP) Numbers" registry to manage the Flag field of the SR-ERO subobject. New values are to be assigned by Standards Action [RFC8126]. Each bit should be tracked with the following qualities:

- o Bit number (counting from bit 0 as the most significant bit)
- o Capability description
- o Defining RFC

The following values are defined in this document:

Bit	Description	Reference
0-7	Unassigned	
8	NAI is absent (F)	This document
9	SID is absent (S)	This document
10	SID specifies TC, S and TTL in addition to an MPLS label (C)	This document
11	SID specifies an MPLS label (M)	This document

8.4. PCEP-Error Object

IANA is requested to confirm the early allocation of the code-points in the PCEP-ERROR Object Error Types and Values registry for the following new error-values:

Error-Type	Meaning
-----	-----
10	Reception of an invalid object.
	Error-value = 2: Bad label value Error-value = 3: Unsupported number of SR-ERO subobjects Error-value = 4: Bad label format Error-value = 5: ERO mixes SR-ERO subobjects with other subobject types Error-value = 6: Both SID and NAI are absent in SR-ERO subobject Error-value = 7: Both SID and NAI are absent in SR-RRO subobject Error-value = 9: MSD exceeds the default for the PCEP session Error-value = 10: RRO mixes SR-RRO subobjects with other subobject types Error-value = TBD1: Missing PCE-SR-CAPABILITY sub-TLV

Error-value = TBD2:	Unsupported NAI Type in SR-ERO subobject
Error-value = TBD3:	Unknown SID
Error-value = TBD4:	NAI cannot be resolved to a SID
Error-value = TBD5:	Could not find SRGB
Error-value = TBD6:	SID index exceeds SRGB size
Error-value = TBD7:	Could not find SRLB
Error-value = TBD8:	SID index exceeds SRLB size
Error-value = TBD9:	Inconsistent SIDs in SR-ERO / SR-RRO subobjects
Error-value = TBD10:	MSD must be nonzero

Note to IANA: this draft originally had an early allocation for Error-value=11 (Malformed object) in the above list. However, we have since moved the definition of that code point to RFC8408.

Note to IANA: some Error-values in the above list were defined after the early allocation took place, and so do not currently have a code point assigned. Please assign code points from the indicated registry and replace each instance of "TBD1", "TBD2" etc. in this document with the respective code points.

Note to IANA: some of the Error-value descriptive strings above have changed since the early allocation. Please refresh the registry.

8.5. PCEP TLV Type Indicators

IANA is requested to confirm the early allocation of the following code point in the PCEP TLV Type Indicators registry. Note that this TLV type indicator is deprecated but retained in the registry to ensure compatibility with early implementations of this specification. See Appendix A for details.

Value	Meaning	Reference
26	SR-PCE-CAPABILITY (deprecated)	This document

8.6. PATH-SETUP-TYPE-CAPABILITY Sub-TLV Type Indicators

IANA is requested to create a new sub-registry, named "PATH-SETUP-TYPE-CAPABILITY Sub-TLV Type Indicators", within the "Path Computation Element Protocol (PCEP) Numbers" registry to manage the

type indicator space for sub-TLVs of the PATH-SETUP-TYPE-CAPABILITY TLV. New values are to be assigned by Standards Action [RFC8126]. The valid range of values in the registry is 0-65535. IANA is requested to initialize the registry with the following values. All other values in the registry should be marked as "Unassigned".

Value	Meaning	Reference
0	Reserved	This document
TBD11 (recommended 26)	SR-PCE-CAPABILITY	This document

Note to IANA: Please replace each instance of "TBD11" in this document with the allocated code point. We have recommended that value 26 be used for consistency with the deprecated value in the PCEP TLV Type Indicators registry.

8.7. New Path Setup Type

[RFC8408] created a sub-registry within the "Path Computation Element Protocol (PCEP) Numbers" registry called "PCEP Path Setup Types". IANA is requested to allocate a new code point within this registry, as follows:

Value	Description	Reference
1	Traffic engineering path is setup using Segment Routing.	This document

8.8. New Metric Type

IANA is requested to confirm the early allocation of the following code point in the PCEP METRIC object T field registry:

Value	Description	Reference
11	Segment-ID (SID) Depth.	This document

8.9. SR PCE Capability Flags

IANA is requested to create a new sub-registry, named "SR Capability Flag Field", within the "Path Computation Element Protocol (PCEP) Numbers" registry to manage the Flag field of the SR-PCE-CAPABILITY TLV. New values are to be assigned by Standards Action [RFC8126]. Each bit should be tracked with the following qualities:

- o Bit number (counting from bit 0 as the most significant bit)
- o Capability description
- o Defining RFC

The following values are defined in this document:

Bit	Description	Reference
0-5	Unassigned	
6	Node or Adjacency Identifier (NAI) is supported (N)	This document
7	Unlimited Maximum SID Depth (X)	This document

Note to IANA: The name of bit 7 has changed from "Unlimited Maximum SID Depth (L)" to "Unlimited Maximum SID Depth (X)".

9. Contributors

The following people contributed to this document:

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Appendix A. Compatibility with Early Implementations

An early implementation of this specification will send the SR-CAPABILITY-TLV as a top-level TLV in the OPEN object instead of sending the PATH-SETUP-TYPE-CAPABILITY TLV in the OPEN object. Implementations that wish to interoperate with such early implementations should also send the SR-CAPABILITY-TLV as a top-level TLV in their OPEN object and should interpret receiving this top-level TLV as though the sender had sent a PATH-SETUP-TYPE-CAPABILITY TLV with a PST list of (0, 1) (that is, both RSVP-TE and SR-TE PSTs are supported) with the SR-CAPABILITY-TLV as a sub-TLV. If a PCEP speaker receives an OPEN object in which both the SR-CAPABILITY-TLV and PATH-SETUP-TYPE-CAPABILITY TLV appear as top-level TLVs, then it should ignore the top-level SR-CAPABILITY-TLV and process only the PATH-SETUP-TYPE-CAPABILITY TLV.

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A YANG Data Model for Path Computation Element Communications Protocol
(PCEP)
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Abstract

This document defines a YANG data model for the management of Path Computation Element communications Protocol (PCEP) for communications between a Path Computation Client (PCC) and a Path Computation Element (PCE), or between two PCEs. The data model includes configuration data and state data (status information and counters for the collection of statistics).

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1. Introduction

The Path Computation Element (PCE) defined in [RFC4655] is an entity that is capable of computing a network path or route based on a network graph, and applying computational constraints. A Path Computation Client (PCC) may make requests to a PCE for paths to be computed.

PCEP is the communication protocol between a PCC and PCE and is defined in [RFC5440]. PCEP interactions include path computation requests and path computation replies as well as notifications of specific states related to the use of a PCE in the context of

Multiprotocol Label Switching (MPLS) and Generalized MPLS (GMPLS) Traffic Engineering (TE).

This document defines a YANG [RFC6020] data model for the management of PCEP speakers. It is important to establish a common data model for how PCEP speakers are identified, configured, and monitored. The data model includes configuration data and state data (status information and counters for the collection of statistics).

This document contains a specification of the base PCEP YANG module, "ietf-pcep" which provides the basic PCEP [RFC5440] data model.

[Editor's Note: Further modules augmenting the data model with advanced features maybe handled in a future revision or a separate document.]

2. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

3. Terminology and Notation

This document uses the terminology defined in [RFC4655] and [RFC5440]. In particular, it uses the following acronyms.

- o Path Computation Request message (PCReq).
- o Path Computation Reply message (PCRep).
- o Notification message (PCNtf).
- o Error message (PCErr).
- o Request Parameters object (RP).
- o Synchronization Vector object (SVEC).
- o Explicit Route object (ERO).

This document also uses the following terms defined in [RFC7420]:

- o PCEP entity: a local PCEP speaker.
- o PCEP peer: to refer to a remote PCEP speaker.

- o PCEP speaker: where it is not necessary to distinguish between local and remote.

3.1. Tree Diagrams

A graphical representation of the complete data tree is presented in Section 5. The meaning of the symbols in these diagrams is as follows and as per [I-D.ietf-netmod-rfc6087bis]:

- o Brackets "[" and "]" enclose list keys.
- o Curly braces "{" and "}" contain names of optional features that make the corresponding node conditional.
- o Abbreviations before data node names: "rw" means configuration (read-write), and "ro" state data (read-only).
- o Symbols after data node names: "?" means an optional node and "*" denotes a "list" or "leaf-list".
- o Parentheses enclose choice and case nodes, and case nodes are also marked with a colon (":").
- o Ellipsis ("...") stands for contents of subtrees that are not shown.

3.2. Prefixes in Data Node Names

In this document, names of data nodes and other data model objects are often used without a prefix, as long as it is clear from the context in which YANG module each name is defined. Otherwise, names are prefixed using the standard prefix associated with the corresponding YANG module, as shown in Table 1.

Prefix	YANG module	Reference
yang	ietf-yang-types	[RFC6991]
inet	ietf-inet-types	[RFC6991]

Table 1: Prefixes and corresponding YANG modules

4. Objectives

This section describes some of the design objectives for the model:

- o In case of existing implementations, it needs to map the data model defined in this document to their proprietary native data model. To facilitate such mappings, the data model should be simple.
- o The data model should be suitable for new implementations to use as is.
- o Mapping to the PCEP MIB Module should be clear.
- o The data model should allow for static configurations of peers.
- o The data model should include read-only counters in order to gather statistics for sent and received PCEP messages, received messages with errors, and messages that could not be sent due to errors.
- o It should be fairly straightforward to augment the base data model for advanced PCE features.

5. The Design of PCEP Data Model

The module, "ietf-pcep", defines the basic components of a PCE speaker.

```

module: ietf-pcep
  +--rw pcep
  |   +--rw entity* [addr]
  |   |   +--rw addr                inet:ip-address
  |   |   +--rw enabled?           boolean
  |   |   +--rw role                pcep-role
  |   |   +--rw description?       string
  |   |   +--rw pce-info
  |   |   |   +--rw scope
  |   |   |   |   +--rw intra-area-scope?    boolean
  |   |   |   |   +--rw intra-area-pref?    uint8
  |   |   |   |   +--rw inter-area-scope?   boolean
  |   |   |   |   +--rw inter-area-scope-default? boolean
  |   |   |   |   +--rw inter-area-pref?   uint8
  |   |   |   |   +--rw inter-as-scope?    boolean
  |   |   |   |   +--rw inter-as-scope-default? boolean
  |   |   |   |   +--rw inter-as-pref?    uint8
  |   |   |   |   +--rw inter-layer-scope?  boolean
  |   |   |   |   +--rw inter-layer-pref?  uint8
  |   |   |   +--rw domain
  |   |   |   |   +--rw domain-type?    pce-domain-type
  |   |   |   |   +--rw domain?        pce-domain
  |   |   |   +--rw neigh-domains
  |   |
  |   +--rw domain
  |   |   +--rw domain-type?    pce-domain-type
  |   |   +--rw domain?        pce-domain
  |   +--rw neigh-domains

```

```

| | | +--rw domain* [domain-type domain]
| | | | +--rw domain-type    pce-domain-type
| | | | +--rw domain        pce-domain
+--rw capability
+--rw  gmpls?                boolean
+--rw  bi-dir?               boolean
+--rw  diverse?              boolean
+--rw  load-balance?         boolean
+--rw  synchronize?         boolean
+--rw  objective-function?   boolean
+--rw  add-path-constraint?  boolean
+--rw  prioritization?       boolean
+--rw  multi-request?        boolean
+--rw  gco?                  boolean
+--rw  p2mp?                 boolean
+--rw  connect-timer?        uint32
+--rw  connect-max-retry?    uint32
+--rw  init-backoff-timer?   uint32
+--rw  max-backoff-timer?    uint32
+--rw  open-wait-timer?      uint32
+--rw  keep-wait-timer?      uint32
+--rw  keep-alive-timer?     uint32
+--rw  dead-timer?           uint32
+--rw  allow-negotiation?    boolean
+--rw  max-keep-alive-timer? uint32
+--rw  max-dead-timer?       uint32
+--rw  min-keep-alive-timer? uint32
+--rw  min-dead-timer?       uint32
+--rw  sync-timer?           uint32
+--rw  request-timer?        uint32
+--rw  max-sessions?         uint32
+--rw  max-unknown-reqs?     uint32
+--rw  max-unknown-msgs?     uint32
+--rw  pcep-notification-max-rate uint32
+--rw  peers
+--rw  peer* [addr]
+--rw  | addr                inet:ip-address
+--rw  | description?        string
+--rw  | pce-info
+--rw  | | scope
+--rw  | | | intra-area-scope?    boolean
+--rw  | | | intra-area-pref?     uint8
+--rw  | | | inter-area-scope?    boolean
+--rw  | | | inter-area-scope-default? boolean
+--rw  | | | inter-area-pref?     uint8
+--rw  | | | inter-as-scope?      boolean
+--rw  | | | inter-as-scope-default? boolean
+--rw  | | | inter-as-pref?       uint8

```

```

|         |   +-rw inter-layer-scope?           boolean
|         |   +-rw inter-layer-pref?          uint8
+-rw domain
|         |   +-rw domain-type?    pce-domain-type
|         |   +-rw domain?         pce-domain
+-rw neigh-domains
|         |   +-rw domain* [domain-type domain]
|         |       +-rw domain-type    pce-domain-type
|         |       +-rw domain        pce-domain
+-rw capability
|         |   +-rw gmpls?            boolean
|         |   +-rw bi-dir?          boolean
|         |   +-rw diverse?        boolean
|         |   +-rw load-balance?    boolean
|         |   +-rw synchronize?    boolean
|         |   +-rw objective-function? boolean
|         |   +-rw add-path-constraint? boolean
|         |   +-rw prioritization?  boolean
|         |   +-rw multi-request?   boolean
|         |   +-rw gco?             boolean
|         |   +-rw p2mp?           boolean
+-ro pcep-state
  +-ro entity* [addr]
    +-ro index?                    uint32
    +-ro addr                       inet:ip-address
    +-ro admin-status?              pcep-admin-status
    +-ro oper-status?              pcep-admin-status
    +-ro role?                      pcep-role
    +-ro pce-info
      +-ro scope
        +-ro intra-area-scope?     boolean
        +-ro intra-area-pref?     uint8
        +-ro inter-area-scope?     boolean
        +-ro inter-area-scope-default? boolean
        +-ro inter-area-pref?     uint8
        +-ro inter-as-scope?       boolean
        +-ro inter-as-scope-default? boolean
        +-ro inter-as-pref?       uint8
        +-ro inter-layer-scope?    boolean
        +-ro inter-layer-pref?    uint8
      +-ro domain
        +-ro domain-type?    pce-domain-type
        +-ro domain?         pce-domain
      +-ro neigh-domains
        +-ro domain* [domain-type domain]
          +-ro domain-type    pce-domain-type
          +-ro domain        pce-domain
      +-ro capability

```

```

|      +--ro gmpls?                boolean
|      +--ro bi-dir?              boolean
|      +--ro diverse?             boolean
|      +--ro load-balance?        boolean
|      +--ro synchronize?         boolean
|      +--ro objective-function?   boolean
|      +--ro add-path-constraint?  boolean
|      +--ro prioritization?       boolean
|      +--ro multi-request?        boolean
|      +--ro gco?                  boolean
|      +--ro p2mp?                 boolean
+--ro connect-timer?              uint32
+--ro connect-max-retry?          uint32
+--ro init-backoff-timer?         uint32
+--ro max-backoff-timer?          uint32
+--ro open-wait-timer?            uint32
+--ro keep-wait-timer?            uint32
+--ro keep-alive-timer?           uint32
+--ro dead-timer?                 uint32
+--ro allow-negotiation?          boolean
+--ro max-keep-alive-timer?       uint32
+--ro max-dead-timer?             uint32
+--ro min-keep-alive-timer?       uint32
+--ro min-dead-timer?             uint32
+--ro sync-timer?                 uint32
+--ro request-timer?              uint32
+--ro max-sessions?               uint32
+--ro max-unknown-reqs?           uint32
+--ro max-unknown-msgs?           uint32
+--ro peers
  +--ro peer* [addr]
    +--ro addr                     inet:ip-address
    +--ro role?                     pcep-role
    +--ro pce-info
      +--ro scope
        +--ro intra-area-scope?     boolean
        +--ro intra-area-pref?      uint8
        +--ro inter-area-scope?     boolean
        +--ro inter-area-scope-default? boolean
        +--ro inter-area-pref?      uint8
        +--ro inter-as-scope?       boolean
        +--ro inter-as-scope-default? boolean
        +--ro inter-as-pref?        uint8
        +--ro inter-layer-scope?    boolean
        +--ro inter-layer-pref?     uint8
      +--ro domain
        +--ro domain-type?          pce-domain-type
        +--ro domain?                pce-domain

```

```

+--ro neigh-domains
|   +--ro domain* [domain-type domain]
|       +--ro domain-type      pce-domain-type
|       +--ro domain           pce-domain
+--ro capability
    +--ro gmpls?                boolean
    +--ro bi-dir?              boolean
    +--ro diverse?             boolean
    +--ro load-balance?        boolean
    +--ro synchronize?         boolean
    +--ro objective-function?  boolean
    +--ro add-path-constraint? boolean
    +--ro prioritization?      boolean
    +--ro multi-request?       boolean
    +--ro gco?                 boolean
    +--ro p2mp?                boolean
+--ro discontinuity-time?      yang:timestamp
+--ro initiate-session?       boolean
+--ro session-exists?         boolean
+--ro num-sess-setup-ok?      yang:counter32
+--ro num-sess-setup-fail?    yang:counter32
+--ro session-up-time?        yang:timestamp
+--ro session-fail-time?      yang:timestamp
+--ro session-fail-up-time?   yang:timestamp
+--ro avg-rsp-time?           uint32
+--ro lwm-rsp-time?           uint32
+--ro hwm-rsp-time?           uint32
+--ro num-pcreq-sent?         yang:counter32
+--ro num-pcreq-rcvd?        yang:counter32
+--ro num-pcrep-sent?        yang:counter32
+--ro num-pcrep-rcvd?        yang:counter32
+--ro num-pcerr-sent?        yang:counter32
+--ro num-pcerr-rcvd?        yang:counter32
+--ro num-pcntf-sent?        yang:counter32
+--ro num-pcntf-rcvd?        yang:counter32
+--ro num-keepalive-sent?    yang:counter32
+--ro num-keepalive-rcvd?    yang:counter32
+--ro num-unknown-rcvd?     yang:counter32
+--ro num-corrupt-rcvd?     yang:counter32
+--ro num-req-sent?          yang:counter32
+--ro num-svec-sent?         yang:counter32
+--ro num-svec-req-sent?     yang:counter32
+--ro num-req-sent-pend-rep? yang:counter32
+--ro num-req-sent-ero-rcvd? yang:counter32
+--ro num-req-sent-nopath-rcvd? yang:counter32
+--ro num-req-sent-cancel-rcvd? yang:counter32
+--ro num-req-sent-error-rcvd? yang:counter32
+--ro num-req-sent-timeout?  yang:counter32

```

```

+--ro num-req-sent-cancel-sent? yang:counter32
+--ro num-req-rcvd? yang:counter32
+--ro num-svec-rcvd? yang:counter32
+--ro num-svec-req-rcvd? yang:counter32
+--ro num-req-rcvd-pend-rep? yang:counter32
+--ro num-req-rcvd-ero-sent? yang:counter32
+--ro num-req-rcvd-nopath-sent? yang:counter32
+--ro num-req-rcvd-cancel-sent? yang:counter32
+--ro num-req-rcvd-error-sent? yang:counter32
+--ro num-req-rcvd-cancel-rcvd? yang:counter32
+--ro num-rep-rcvd-unknown? yang:counter32
+--ro num-req-rcvd-unknown? yang:counter32
+--ro num-req-sent-closed? yang:counter32
+--ro num-req-rcvd-closed? yang:counter32
+--ro sessions
  +--ro session* [initiator]
    +--ro initiator pcep-initiator
    +--ro state-last-change? yang:timestamp
    +--ro state? pcep-sess-state
    +--ro connect-retry? yang:counter32
    +--ro local-id? uint32
    +--ro remote-id? uint32
    +--ro keepalive-timer? uint32
    +--ro peer-keepalive-timer? uint32
    +--ro dead-timer? uint32
    +--ro peer-dead-timer? uint32
    +--ro ka-hold-time-rem? uint32
    +--ro overloaded? boolean
    +--ro overload-time? uint32
    +--ro peer-overloaded? boolean
    +--ro peer-overload-time? uint32
    +--ro discontinuity-time? yang:timestamp
    +--ro avg-rsp-time? uint32
    +--ro lwm-rsp-time? uint32
    +--ro hwm-rsp-time? uint32
    +--ro num-pcreq-sent? yang:counter32
    +--ro num-pcreq-rcvd? yang:counter32
    +--ro num-pcrep-sent? yang:counter32
    +--ro num-pcrep-rcvd? yang:counter32
    +--ro num-pcerr-sent? yang:counter32
    +--ro num-pcerr-rcvd? yang:counter32
    +--ro num-pcntf-sent? yang:counter32
    +--ro num-pcntf-rcvd? yang:counter32
    +--ro num-keepalive-sent? yang:counter32
    +--ro num-keepalive-rcvd? yang:counter32
    +--ro num-unknown-rcvd? yang:counter32
    +--ro num-corrupt-rcvd? yang:counter32
    +--ro num-req-sent? yang:counter32

```

```

        +--ro num-svec-sent?          yang:counter32
        +--ro num-svec-req-sent?     yang:counter32
        +--ro num-req-sent-pend-rep? yang:counter32
        +--ro num-req-sent-ero-rcvd? yang:counter32
        +--ro num-req-sent-nopath-rcvd? yang:counter32
        +--ro num-req-sent-cancel-rcvd? yang:counter32
        +--ro num-req-sent-error-rcvd? yang:counter32
        +--ro num-req-sent-timeout?   yang:counter32
        +--ro num-req-sent-cancel-sent? yang:counter32
        +--ro num-req-rcvd?           yang:counter32
        +--ro num-svec-rcvd?         yang:counter32
        +--ro num-svec-req-rcvd?     yang:counter32
        +--ro num-req-rcvd-pend-rep?  yang:counter32
        +--ro num-req-rcvd-ero-sent?  yang:counter32
        +--ro num-req-rcvd-nopath-sent? yang:counter32
        +--ro num-req-rcvd-cancel-sent? yang:counter32
        +--ro num-req-rcvd-error-sent? yang:counter32
        +--ro num-req-rcvd-cancel-rcvd? yang:counter32
        +--ro num-rep-rcvd-unknown?   yang:counter32
        +--ro num-req-rcvd-unknown?   yang:counter32
notifications:
  +---n pcep-session-up
  |   +--ro entity-addr?          leafref
  |   +--ro peer-addr?           leafref
  |   +--ro session-initiator?   leafref
  |   +--ro state-last-change?   yang:timestamp
  |   +--ro state?               pcep-sess-state
  +---n pcep-session-down
  |   +--ro entity-addr?          leafref
  |   +--ro peer-addr?           leafref
  |   +--ro session-initiator?   pcep-initiator
  |   +--ro state-last-change?   yang:timestamp
  |   +--ro state?               pcep-sess-state
  +---n pcep-session-local-overload
  |   +--ro entity-addr?          leafref
  |   +--ro peer-addr?           leafref
  |   +--ro session-initiator?   leafref
  |   +--ro overloaded?          boolean
  |   +--ro overload-time?       uint32
  +---n pcep-session-local-overload-clear
  |   +--ro entity-addr?          leafref
  |   +--ro peer-addr?           leafref
  |   +--ro overloaded?          boolean
  +---n pcep-session-peer-overload
  |   +--ro entity-addr?          leafref
  |   +--ro peer-addr?           leafref
  |   +--ro session-initiator?   leafref
  |   +--ro peer-overloaded?     boolean

```



```
|  +--ro peer-overload-time?  uint32
+---n pcep-session-peer-overload-clear
  +--ro entity-addr?         leafref
  +--ro peer-addr?          leafref
  +--ro peer-overloaded?    boolean
```

5.1. The Entity Lists

The PCEP yang module may contain status information for multiple logical local PCEP entities. There are several scenarios in which there may be more than one local PCEP entity, it is listed in [RFC7420].

The data model for PCEP presented in this document uses a flat list of entities. Each entity in the list is identified by its IP address (addr-type, addr). Furthermore, each entity has a mandatory "role" leaf (the local entity PCEP role). The ietf-inet-types [RFC6991] is used.

There is one list for configuration ("/pcep/entity"), and a separate list for the operational state of all entities ("/pcep-state/entity").

The PCEP MIB module [RFC7420] uses a system generated entity index as a primary index to the read only entity table. If the device implements the PCEP MIB, the "index" leaf MUST contain the value of the corresponding pcePcepEntityIndex.

5.2. The Peer Lists

The peer list contains peer(s) that the local PCEP entity knows about. A PCEP speaker is identified by its IP address. If there is a PCEP speaker in the network that uses multiple IP addresses then it looks like multiple distinct peers to the other PCEP speakers in the network.

Since PCEP sessions can be ephemeral, the peer list tracks a peer even when no PCEP session currently exists to that peer. The statistics contained are an aggregate of the statistics for all successive sessions to that peer.

To limit the quantity of information that is stored, an implementation MAY choose to discard this information if and only if no PCEP session exists to the corresponding peer.

The data model for PCEP peer presented in this document uses a flat list of peers. Each peer in the list is identified by its IP address (addr-type, addr).

There is one list for static peer configuration ("/pcep/entity/peers"), and a separate list for the operational state of all peers (i.e. static as well as discovered)("/pcep-state/entity/peers").

5.3. The Session Lists

The session list contains PCEP session that the PCEP entity (PCE or PCC) is currently participating in. The statistics in session are semantically different from those in peer since the former applies to the current session only, whereas the latter is the aggregate for all sessions that have existed to that peer.

Although [RFC5440] forbids more than one active PCEP session between a given pair of PCEP entities at any given time, there is a window during session establishment where two sessions may exist for a given pair, one representing a session initiated by the local PCEP entity and the other representing a session initiated by the peer. If either of these sessions reaches active state first, then the other is discarded.

The data model for PCEP session presented in this document uses a flat list of sessions. Each session in the list is identified by its initiator. This index allows two sessions to exist transiently for a given peer, as discussed above.

There is only one list for the operational state of all sessions ("/pcep-state/entity/peers/peer/sessions/session").

5.4. Notifications

This YANG model defines a list of notifications to inform client of important events detected during the protocol operation. The notifications defined cover the PCEP MIB notifications.

6. Advanced PCE Features

This document contains a specification of the base PCEP YANG module, "ietf-pcep" which provides the basic PCEP [RFC5440] data model.

A means and procedure to handle to following PCE features needs to be decided:

- o Capability and Scope

- o Domain information (local/neighbour)
- o Path-Key
- o OF
- o GCO
- o P2MP
- o GMPLS
- o Inter-Layer
- o Stateful PCE

7. PCEP YANG Module

RFC Ed.: In this section, replace all occurrences of 'XXXX' with the actual RFC number and all occurrences of the revision date below with the date of RFC publication (and remove this note).

<CODE BEGINS> file "ietf-pcep@2015-02-xx.yang"

```
module ietf-pcep {
  namespace "urn:ietf:params:xml:ns:yang:ietf-pcep";
  prefix pcep;

  import ietf-inet-types {
    prefix inet;
  }

  import ietf-yang-types {
    prefix yang;
  }

  organization
    "IETF PCE (Path Computation Element) Working Group";

  contact
    "WG Web: <http://tools.ietf.org/wg/pce/>
    WG List: <mailto:pcep@ietf.org>
```

WG Chair: JP Vasseur
<mailto:jpv@cisco.com>

WG Chair: Julien Meuric
<mailto:julien.meuric@orange.com>

Editor: Dhruv Dhody
<mailto:dhruv.ietf@gmail.com>;

```
description
  "The YANG module defines a generic configuration and
  operational model for PCEP common across all of the
  vendor implementations.";

revision 2015-02-23 {
  description "Initial revision.";
  reference
    "RFC XXXX: A YANG Data Model for Path Computation
    Element Communications Protocol
    (PCEP)";
}

/*
 * Identities
 */

identity pcep {
  description "Identity for the PCEP protocol.";
}

/*
 * Typedefs
 */
typedef pcep-role {
  type enumeration {
    enum unknown {
      value "0";
      description
        "An unknown role";
    }
    enum pcc {
      value "1";
      description
        "The role of a Path Computation Client";
    }
    enum pce {
      value "2";
    }
  }
}
```

```
        description
            "The role of Path Computation Element";
    }
    enum pcc-and-pce {
        value "3";
        description
            "The role of both Path Computation Client and
            Path Computation Element";
    }
}

description
    "The role of a PCEP speaker.
    Takes one of the following values
    - unknown(0): the role is not known.
    - pcc(1): the role is of a Path Computation
      Client (PCC).
    - pce(2): the role is of a Path Computation
      Server (PCE).
    - pccAndPce(3): the role is of both a PCC and
      a PCE.";
}

typedef pcep-admin-status {
    type enumeration {
        enum admin-status-up {
            value "1";
            description
                "Admin Status is Up";
        }
        enum admin-status-down {
            value "2";
            description
                "Admin Status is Down";
        }
    }
}

description
    "The Admin Status of the PCEP entity.
    Takes one of the following values
    - admin-status-up(1): Admin Status is Up.
    - admin-status-down(2): Admin Status is Down";
}

typedef pcep-oper-status {
    type enumeration {
        enum oper-status-up {
```

```
        value "1";
        description
            "The PCEP entity is active";
    }
    enum oper-status-down {
        value "2";
        description
            "The PCEP entity is inactive";
    }
    enum oper-status-going-up {
        value "3";
        description
            "The PCEP entity is activating";
    }
    enum oper-status-going-down {
        value "4";
        description
            "The PCEP entity is deactivating";
    }
    enum oper-status-failed {
        value "5";
        description
            "The PCEP entity has failed and will recover
            when possible.";
    }
    enum oper-status-failed-perm {
        value "6";
        description
            "The PCEP entity has failed and will not recover
            without operator intervention";
    }
}
description
    "The operational status of the PCEP entity.
    Takes one of the following values
    - oper-status-up(1): Active
    - oper-status-down(2): Inactive
    - oper-status-going-up(3): Activating
    - oper-status-going-down(4): Deactivating
    - oper-status-failed(5): Failed
    - oper-status-failed-perm(6): Failed Permanantly";
}

typedef pcep-initiator {
    type enumeration {
        enum local {
            value "1";
            description
```

```
        "The local PCEP entity initiated the session";
    }

    enum remote {
        value "2";
        description
            "The remote PCEP peer initiated the session";
    }
}
description
    "The initiator of the session, that is, whether the TCP
    connection was initiated by the local PCEP entity or
    the remote peer.
    Takes one of the following values
    - local(1): Initiated locally
    - remote(2): Initiated remotely";
}

typedef pcep-sess-state {
    type enumeration {
        enum tcp-pending {
            value "1";
            description
                "The tcp-pending state of PCEP session.";
        }

        enum open-wait {
            value "2";
            description
                "The open-wait state of PCEP session.";
        }

        enum keep-wait {
            value "3";
            description
                "The keep-wait state of PCEP session.";
        }

        enum session-up {
            value "4";
            description
                "The session-up state of PCEP session.";
        }
    }
}
description
    "The current state of the session.

    The set of possible states excludes the idle state
```

```
        since entries do not exist in the idle state.
        Takes one of the following values
        - tcp-pending(1): PCEP TCP Pending state
        - open-wait(2): PCEP Open Wait state
        - keep-wait(3): PCEP Keep Wait state
        - session-up(4): PCEP Session Up state";
    }

typedef pce-domain-type {
    type enumeration {
        enum ospf-area {
            value "1";
            description
                "The OSPF area.";
        }
        enum isis-area {
            value "2";
            description
                "The IS-IS area.";
        }
        enum as {
            value "3";
            description
                "The Autonomous System (AS).";
        }
    }
    description
        "The PCE Domain Type";
}

typedef domain-ospf-area {
    type union {
        type uint32;
        type yang:dotted-quad;
    }
    description
        "OSPF Area ID.";
}

typedef domain-isis-area {
    type string {
        pattern '[0-9A-Fa-f]{2}\.([0-9A-Fa-f]{4}\.){0,3}';
    }
    description
        "IS-IS Area ID.";
}

typedef domain-as {
```



```
    type uint32;
    description
        "Autonomous System number.";
}

typedef pce-domain {
    type union {
        type domain-ospf-area;
        type domain-isis-area;
        type domain-as;
    }
    description
        "The PCE Domain";
}

/*
 * Features - TBD
 */

/*
 * Groupings
 */

grouping pcep-entity-info{
    description
        "This grouping defines the attributes for PCEP entity.";
    leaf connect-timer {
        type uint32 {
            range "1..65535";
        }
        units "seconds";
        default 60;
        description
            "The time in seconds that the PCEP entity will wait
            to establish a TCP connection with a peer.  If a
            TCP connection is not established within this time
            then PCEP aborts the session setup attempt.";
        reference
            "RFC 5440: Path Computation Element (PCE)
            Communication Protocol (PCEP)";
    }

    leaf connect-max-retry {
```

```
    type uint32;
    default 5;
    description
        "The maximum number of times the system tries to
        establish a TCP connection to a peer before the
        session with the peer transitions to the idle
        state.";
    reference
        "RFC 5440: Path Computation Element (PCE)
        Communication Protocol (PCEP)";
}

leaf init-backoff-timer {
    type uint32 {
        range "1..65535";
    }
    units "seconds";
    description
        "The initial back-off time in seconds for retrying
        a failed session setup attempt to a peer.

        The back-off time increases for each failed
        session setup attempt, until a maximum back-off
        time is reached. The maximum back-off time is
        max-backoff-timer.";
}

leaf max-backoff-timer {
    type uint32;
    units "seconds";
    description
        "The maximum back-off time in seconds for retrying
        a failed session setup attempt to a peer.

        The back-off time increases for each failed session
        setup attempt, until this maximum value is reached.
        Session setup attempts then repeat periodically
        without any further increase in back-off time.";
}

leaf open-wait-timer {
    type uint32 {
        range "1..65535";
    }
    units "seconds";
    default 60;
    description
        "The time in seconds that the PCEP entity will wait
```

to receive an Open message from a peer after the TCP connection has come up.

If no Open message is received within this time then PCEP terminates the TCP connection and deletes the associated sessions.";

```
reference
  "RFC 5440: Path Computation Element (PCE)
    Communication Protocol (PCEP)";
}

leaf keep-wait-timer {
  type uint32 {
    range "1..65535";
  }
  units "seconds";
  default 60;
  description
    "The time in seconds that the PCEP entity will wait
    to receive a Keepalive or PCErr message from a peer
    during session initialization after receiving an
    Open message.  If no Keepalive or PCErr message is
    received within this time then PCEP terminates the
    TCP connection and deletes the associated
    sessions.";
  reference
    "RFC 5440: Path Computation Element (PCE)
      Communication Protocol (PCEP)";
}

leaf keep-alive-timer {
  type uint32 {
    range "0..255";
  }
  units "seconds";
  default 30;
  description
    "The keep alive transmission timer that this PCEP
    entity will propose in the initial OPEN message of
    each session it is involved in.  This is the
    maximum time between two consecutive messages sent
    to a peer.  Zero means that the PCEP entity prefers
    not to send Keepalives at all.

    Note that the actual Keepalive transmission
    intervals, in either direction of an active PCEP
    session, are determined by negotiation between the
    peers as specified by RFC 5440, and so may differ
```

```
        from this configured value.";
    reference
        "RFC 5440: Path Computation Element (PCE)
        Communication Protocol (PCEP)";
}

leaf dead-timer {
    type uint32 {
        range "0..255";
    }
    units "seconds";
    must ". >= ../keep-alive-timer" {
        error-message "The dead timer must be "
            + "larger than the keep alive timer";
        description
            "This value MUST be greater than
            keep-alive-timer.";
    }
    default 120;
    description
        "The dead timer that this PCEP entity will propose
        in the initial OPEN message of each session it is
        involved in. This is the time after which a peer
        should declare a session down if it does not
        receive any PCEP messages. Zero suggests that the
        peer does not run a dead timer at all." ;
    reference
        "RFC 5440: Path Computation Element (PCE)
        Communication Protocol (PCEP)";
}

leaf allow-negotiation{
    type boolean;
    description
        "Whether the PCEP entity will permit negotiation of
        session parameters.";
}

leaf max-keep-alive-timer{
    type uint32 {
        range "0..255";
    }
    units "seconds";
    description
        "In PCEP session parameter negotiation in seconds,
        the maximum value that this PCEP entity will
        accept from a peer for the interval between
```

```
        Keepalive transmissions. Zero means that the PCEP
        entity will allow no Keepalive transmission at
        all." ;
    }

    leaf max-dead-timer{
        type uint32 {
            range "0..255";
        }
        units "seconds";
        description
            "In PCEP session parameter negotiation in seconds,
            the maximum value that this PCEP entity will accept
            from a peer for the Dead timer. Zero means that
            the PCEP entity will allow not running a Dead
            timer.";
    }

    leaf min-keep-alive-timer{
        type uint32 {
            range "0..255";
        }
        units "seconds";
        description
            "In PCEP session parameter negotiation in seconds,
            the minimum value that this PCEP entity will
            accept for the interval between Keepalive
            transmissions. Zero means that the PCEP entity
            insists on no Keepalive transmission at all.";
    }

    leaf min-dead-timer{
        type uint32 {
            range "0..255";
        }
        units "seconds";
        description
            "In PCEP session parameter negotiation in
            seconds, the minimum value that this PCEP entity
            will accept for the Dead timer. Zero means that
            the PCEP entity insists on not running a Dead
            timer.";
    }

    leaf sync-timer{
        type uint32 {
            range "0..65535";
        }
    }
}
```

```
    units "seconds";
    default 60;
    description
        "The value of SyncTimer in seconds is used in the
        case of synchronized path computation request
        using the SVEC object. Consider the case where a
        PCReq message is received by a PCE that contains
        the SVEC object referring to M synchronized path
        computation requests. If after the expiration of
        the SyncTimer all the M path computation requests
        have not been, received a protocol error is
        triggered and the PCE MUST cancel the whole set
        of path computation requests.

        The aim of the SyncTimer is to avoid the storage
        of unused synchronized requests should one of
        them get lost for some reasons (for example, a
        misbehaving PCC).

        Zero means that the PCEP entity does not use the
        SyncTimer.";
    reference
        "RFC 5440: Path Computation Element (PCE)
        Communication Protocol (PCEP)";
}

leaf request-timer{
    type uint32 {
        range "1..65535";
    }
    units "seconds";
    description
        "The maximum time that the PCEP entity will wait
        for a response to a PCReq message.";
}

leaf max-sessions{
    type uint32;
    description
        "Maximum number of sessions involving this PCEP
        entity that can exist at any time.";
}

leaf max-unknown-reqs{
    type uint32;
    default 5;
    description
        "The maximum number of unrecognized requests and
```

replies that any session on this PCEP entity is willing to accept per minute before terminating the session.

A PCRep message contains an unrecognized reply if it contains an RP object whose request ID does not correspond to any in-progress request sent by this PCEP entity.

A PCReq message contains an unrecognized request if it contains an RP object whose request ID is zero.";

```
reference
  "RFC 5440: Path Computation Element (PCE)
    Communication Protocol (PCE)";
}

leaf max-unknown-msgs{
  type uint32;
  default 5;
  description
    "The maximum number of unknown messages that any
    session on this PCEP entity is willing to accept
    per minute before terminating the session.";
  reference
    "RFC 5440: Path Computation Element (PCE)
      Communication Protocol (PCE)";
}
} //pcep-entity-info

grouping pce-scope{
  description
    "This grouping defines PCE path computation scope
    information which maybe relevant to PCE selection.
    This information corresponds to PCE auto-discovery
    information.";
  reference
    "RFC 5088: OSPF Protocol Extensions for Path
      Computation Element (PCE)
      Discovery
    RFC 5089: IS-IS Protocol Extensions for Path
      Computation Element (PCE)
      Discovery";
  leaf intra-area-scope{
    type boolean;
    default true;
    description
      "PCE can compute intra-area paths.";
  }
}
```

```
    }
    leaf intra-area-pref{
      type uint8{
        range "0..7";
      }
      description
        "The PCE's preference for intra-area TE LSP
        computation.";
    }
    leaf inter-area-scope{
      type boolean;
      default false;
      description
        "PCE can compute inter-area paths.";
    }
    leaf inter-area-scope-default{
      type boolean;
      default false;
      description
        "PCE can act as a default PCE for inter-area
        path computation.";
    }
    leaf inter-area-pref{
      type uint8{
        range "0..7";
      }
      description
        "The PCE's preference for inter-area TE LSP
        computation.";
    }
    leaf inter-as-scope{
      type boolean;
      default false;
      description
        "PCE can compute inter-AS paths.";
    }
    leaf inter-as-scope-default{
      type boolean;
      default false;
      description
        "PCE can act as a default PCE for inter-AS
        path computation.";
    }
    leaf inter-as-pref{
      type uint8{
        range "0..7";
      }
      description
```



```
        "The PCE's preference for inter-AS TE LSP
        computation.";
    }
    leaf inter-layer-scope{
        type boolean;
        default false;
        description
            "PCE can compute inter-layer paths.";
    }
    leaf inter-layer-pref{
        type uint8{
            range "0..7";
        }
        description
            "The PCE's preference for inter-layer TE LSP
            computation.";
    }
} //pce-scope

grouping pce-domain{
    description
        "This grouping specifies a PCE-Domain where the
        PCE has topology visibility and through which
        the PCE can compute paths.";
    leaf domain-type{
        type pce-domain-type;
        description
            "The PCE domain type.";
    }
    leaf domain{
        type pce-domain;
        description
            "The PCE domain.";
    }
} //pce-domain

grouping pce-capability{
    description
        "This grouping specifies a PCE-capability
        information which maybe relevant to PCE selection.
        This information corresponds to PCE auto-discovery
        information.";
    reference
        "RFC 5088: OSPF Protocol Extensions for Path
        Computation Element (PCE)
        Discovery
        RFC 5089: IS-IS Protocol Extensions for Path
        Computation Element (PCE)"
}
```

```
                Discovery";
leaf gmpls{
    type boolean;
    description
        "Path computation with GMPLS link
        constraints.";
}
leaf bi-dir{
    type boolean;
    description
        "Bidirectional path computation.";
}
leaf diverse{
    type boolean;
    description
        "Diverse path computation.";
}
leaf load-balance{
    type boolean;
    description
        "Load-balanced path computation.";
}
leaf synchronize{
    type boolean;
    description
        "Synchronized paths computation.";
}
leaf objective-function{
    type boolean;
    description
        "Support for multiple objective functions.";
}
leaf add-path-constraint{
    type boolean;
    description
        "Support for additive path constraints (max
        hop count, etc.).";
}
leaf prioritization{
    type boolean;
    description
        "Support for request prioritization.";
}
leaf multi-request{
    type boolean;
    description
        "Support for multiple requests per message.";
}
```

```
    leaf gco{
      type boolean;
      description
        "Support for Global Concurrent Optimization
        (GCO).";
    }
    leaf p2mp{
      type boolean;
      description
        "Support for P2MP path computation.";
    }
  }//pce-capability

grouping pce-info{
  description
    "This grouping specifies all PCE information
    which maybe relevant to the PCE selection.
    This information corresponds to PCE auto-discovery
    information.";
  container scope{
    description
      "The path computation scope";
    uses pce-scope;
  }
  container domain{
    description
      "The PCE domain";
    uses pce-domain{
      description
        "The PCE local domain.";
    }
  }
  container neigh-domains{
    description
      "The list of neighbour PCE-Domain
      toward which a PCE can compute
      paths";
    list domain{
      key "domain-type domain";

      description
        "The neighbour domain.";
      uses pce-domain{
        description
          "The PCE neighbour domain.";
      }
    }
  }
}
```

```
    container capability{
      description
        "The PCE capability";
      uses pce-capability{
        description
          "The PCE entity supported
          capabilities.";
      }
    }
  }//pce-info

grouping pcep-stats{
  description
    "This grouping defines statistics for PCEP. It is used
    for both peer and current session.";
  leaf avg-rsp-time{
    type uint32;
    units "milliseconds";
    must "(/pcep-state/entity/peers/peer/role != 'pcc'" +
      " or " +
      "(/pcep-state/entity/peers/peer/role = 'pcc'" +
      " and avg-rsp-time = 0))" {
      error-message
        "Invalid average response time";
      description
        "If role is pcc then this leaf is meaningless
        and is set to zero.";
    }
    description
      "The average response time.

      If an average response time has not been
      calculated then this leaf has the value zero.";
  }

  leaf lwm-rsp-time{
    type uint32;
    units "milliseconds";
    must "(/pcep-state/entity/peers/peer/role != 'pcc'" +
      " or " +
      "(/pcep-state/entity/peers/peer/role = 'pcc'" +
      " and lwm-rsp-time = 0))" {
      error-message
        "Invalid smallest (low-water mark)
        response time";
      description
        "If role is pcc then this leaf is meaningless
        and is set to zero.";
    }
  }
}
```

```
    }
    description
      "The smallest (low-water mark) response time seen.

      If no responses have been received then this
      leaf has the value zero.";
  }

  leaf hwm-rsp-time{
    type uint32;
    units "milliseconds";
    must "(/pcep-state/entity/peers/peer/role != 'pcc'" +
      " or " +
      "(/pcep-state/entity/peers/peer/role = 'pcc'" +
      " and hwm-rsp-time = 0))" {
      error-message
        "Invalid greatest (high-water mark)
        response time seen";
      description
        "If role is pcc then this field is
        meaningless and is set to zero.";
    }
    description
      "The greatest (high-water mark) response time seen.

      If no responses have been received then this object
      has the value zero.";
  }

  leaf num-pcreq-sent{
    type yang:counter32;
    description
      "The number of PCReq messages sent.";
  }

  leaf num-pcreq-rcvd{
    type yang:counter32;
    description
      "The number of PCReq messages received.";
  }

  leaf num-pcrep-sent{
    type yang:counter32;
    description
      "The number of PCRep messages sent.";
  }

  leaf num-pcrep-rcvd{
```

```
        type yang:counter32;
        description
            "The number of PCRep messages received.";
    }

    leaf num-pcerr-sent{
        type yang:counter32;
        description
            "The number of PCErr messages sent.";
    }

    leaf num-pcerr-rcvd{
        type yang:counter32;
        description
            "The number of PCErr messages received.";
    }

    leaf num-pcntf-sent{
        type yang:counter32;
        description
            "The number of PCNtf messages sent.";
    }

    leaf num-pcntf-rcvd{
        type yang:counter32;
        description
            "The number of PCNtf messages received.";
    }

    leaf num-keepalive-sent{
        type yang:counter32;
        description
            "The number of Keepalive messages sent.";
    }

    leaf num-keepalive-rcvd{
        type yang:counter32;
        description
            "The number of Keepalive messages received.";
    }

    leaf num-unknown-rcvd{
        type yang:counter32;
        description
            "The number of unknown messages received.";
    }

    leaf num-corrupt-rcvd{
```

```
    type yang:counter32;
    description
      "The number of corrupted PCEP message received.";
  }

  leaf num-req-sent{
    type yang:counter32;
    description
      "The number of requests sent. A request corresponds
      1:1 with an RP object in a PCReq message. This might
      be greater than num-pcreq-sent because multiple
      requests can be batched into a single PCReq
      message.";
  }

  leaf num-svec-sent{
    type yang:counter32;
    description
      "The number of SVEC objects sent in PCReq messages.
      An SVEC object represents a set of synchronized
      requests.";
  }

  leaf num-svec-req-sent{
    type yang:counter32;
    description
      "The number of requests sent that appeared in one
      or more SVEC objects.";
  }

  leaf num-req-sent-pend-rep{
    type yang:counter32;
    description
      "The number of requests that have been sent for
      which a response is still pending.";
  }

  leaf num-req-sent-ero-rcvd{
    type yang:counter32;
    description
      "The number of requests that have been sent for
      which a response with an ERO object was received.
      Such responses indicate that a path was
      successfully computed by the peer.";
  }

  leaf num-req-sent-nopath-rcvd{
    type yang:counter32;
```

```
description
  "The number of requests that have been sent for
  which a response with a NO-PATH object was
  received. Such responses indicate that the peer
  could not find a path to satisfy the
  request.";
}

leaf num-req-sent-cancel-rcvd{
  type yang:counter32;
  description
    "The number of requests that were cancelled with
    a PCNtf message.

    This might be different than num-pcntf-rcvd because
    not all PCNtf messages are used to cancel requests,
    and a single PCNtf message can cancel multiple
    requests.";
}

leaf num-req-sent-error-rcvd{
  type yang:counter32;
  description
    "The number of requests that were rejected with a
    PCErr message.

    This might be different than num-pcerr-rcvd because
    not all PCErr messages are used to reject requests,
    and a single PCErr message can reject multiple
    requests.";
}

leaf num-req-sent-timeout{
  type yang:counter32;
  description
    "The number of requests that have been sent to a peer
    and have been abandoned because the peer has taken too
    long to respond to them.";
}

leaf num-req-sent-cancel-sent{
  type yang:counter32;
  description
    "The number of requests that were sent to the peer and
    explicitly cancelled by the local PCEP entity sending
    a PCNtf.";
}
```



```
leaf num-req-rcvd{
  type yang:counter32;
  description
    "The number of requests received. A request
    corresponds 1:1 with an RP object in a PCReq
    message.

    This might be greater than num-pcreq-rcvd because
    multiple requests can be batched into a single
    PCReq message.";
}

leaf num-svec-rcvd{
  type yang:counter32;
  description
    "The number of SVEC objects received in PCReq
    messages. An SVEC object represents a set of
    synchronized requests.";
}

leaf num-svec-req-rcvd{
  type yang:counter32;
  description
    "The number of requests received that appeared in one
    or more SVEC objects.";
}

leaf num-req-rcvd-pend-rep{
  type yang:counter32;
  description
    "The number of requests that have been received for
    which a response is still pending.";
}

leaf num-req-rcvd-ero-sent{
  type yang:counter32;
  description
    "The number of requests that have been received for
    which a response with an ERO object was sent. Such
    responses indicate that a path was successfully
    computed by the local PCEP entity.";
}

leaf num-req-rcvd-nopath-sent{
  type yang:counter32;
  description
    "The number of requests that have been received for
    which a response with a NO-PATH object was sent. Such
```

```
        responses indicate that the local PCEP entity could
        not find a path to satisfy the request.";
    }

    leaf num-req-rcvd-cancel-sent{
        type yang:counter32;
        description
            "The number of requests received that were cancelled
            by the local PCEP entity sending a PCNtf message.

            This might be different than num-pcntf-sent because
            not all PCNtf messages are used to cancel requests,
            and a single PCNtf message can cancel multiple
            requests.";
    }

    leaf num-req-rcvd-error-sent{
        type yang:counter32;
        description
            "The number of requests received that were cancelled
            by the local PCEP entity sending a PCErr message.

            This might be different than num-pcerr-sent because
            not all PCErr messages are used to cancel requests,
            and a single PCErr message can cancel multiple
            requests.";
    }

    leaf num-req-rcvd-cancel-rcvd{
        type yang:counter32;
        description
            "The number of requests that were received from the
            peer and explicitly cancelled by the peer sending
            a PCNtf.";
    }

    leaf num-rep-rcvd-unknown{
        type yang:counter32;
        description
            "The number of responses to unknown requests
            received. A response to an unknown request is a
            response whose RP object does not contain the
            request ID of any request that is currently
            outstanding on the session.";
    }

    leaf num-req-rcvd-unknown{
        type yang:counter32;
```

```
        description
            "The number of unknown requests that have been
            received. An unknown request is a request
            whose RP object contains a request ID of
            zero.";
    }
} // pcep-stats

grouping notification-instance-hdr {
    description
        "This group describes common instance specific data
        for notifications.";

    leaf entity-addr {
        type leafref {
            path "/pcep-state/entity/addr";
        }
        description
            "Reference to local entity address";
    }

    leaf peer-addr {
        type leafref {
            path "/pcep-state/entity/peers/peer/addr";
        }
        description
            "Reference to peer address";
    }
} // notification-instance-hdr

grouping notification-session-hdr {
    description
        "This group describes common session instance specific
        data for notifications.";

    leaf session-initiator {
        type leafref {
            path "/pcep-state/entity/peers/peer/sessions/" +
                "session/initiator";
        }
        description
            "Reference to pcep session initiator leaf";
    }
} // notification-session-hdr

/*
```

```
* Configuration data nodes
*/
container pcep{

  description
    "Parameters for list of configured PCEP entities
    on the device.";

  list entity{
    key "addr";

    description
      "The configured PCEP entity on the device.";

    leaf addr {
      type inet:ip-address;
      description
        "The local Internet address of this PCEP
        entity.

        If operating as a PCE server, the PCEP
        entity listens on this address.

        If operating as a PCC, the PCEP entity
        binds outgoing TCP connections to this
        address.

        It is possible for the PCEP entity to
        operate both as a PCC and a PCE Server, in
        which case it uses this address both to
        listen for incoming TCP connections and to
        bind outgoing TCP connections.";
    }

    leaf enabled {
      type boolean;
      default true;
      description
        "The administrative status of this PCEP
        Entity.";
    }

    leaf role {
      type pcep-role;
      mandatory true;
      description
        "The role that this entity can play.
        Takes one of the following values.
```

```
        - unknown(0): this PCEP Entity role is not
          known.
        - pcc(1): this PCEP Entity is a PCC.
        - pce(2): this PCEP Entity is a PCE.
        - pcc-and-pce(3): this PCEP Entity is both
          a PCC and a PCE.";
    }

    leaf description {
        type string;
        description
            "Description of the PCEP entity configured
             by the user";
    }

    container pce-info {
        must "(../role == 'pce')" +
            " or " +
            "(../role == 'pcc-and-pce')";
        {
            error-message "The PCEP entity must be PCE";
            description
                "When PCEP entity is PCE";
        }
        uses pce-info {
            description
                "Local PCE information";
        }
        description
            "The Local PCE Entity PCE information";
    }

    uses pcep-entity-info {
        description
            "The configuration related to the PCEP
             entity.";
    }

    leaf pcep-notification-max-rate {
        type uint32;
        mandatory true;
        description
            "This variable indicates the maximum number of
             notifications issued per second. If events occur
             more rapidly, the implementation may simply fail
             to emit these notifications during that period,
```

```
        or may queue them until an appropriate time. A
        value of 0 means no notifications are emitted
        and all should be discarded (that is, not
        queued).";
    }

    container peers{
        description
            "The list of configured peers for the
            entity";

        list peer{
            key "addr";

            description
                "The peer configured for the entity.";

            leaf addr {
                type inet:ip-address;
                description
                    "The local Internet address of this PCEP peer.";
            }

            leaf description {
                type string;
                description
                    "Description of the PCEP peer
                    configured by the user";
            }

            container pce-info {
                uses pce-info {
                    description
                        "PCE Peer information";
                }
            }

            description
                "The Local PCE Entity PCE information";
        }
    } //peer
} //peers
} //entity
} //pcep

/*
 * Operational data nodes
 */

container pcep-state{
    config false;
```

```
description
  "The list of operational PCEP entities on the
  device.";

list entity{
  key "addr";
  unique "index";

  description
    "The operational PCEP entity on the device.";

  leaf index{
    type uint32;
    description
      "The index of the operational PCEP
      entity";
  }

  leaf addr {
    type inet:ip-address;
    description
      "The local Internet address of this PCEP
      entity.

      If operating as a PCE server, the PCEP
      entity listens on this address.

      If operating as a PCC, the PCEP entity
      binds outgoing TCP connections to this
      address.

      It is possible for the PCEP entity to
      operate both as a PCC and a PCE Server, in
      which case it uses this address both to
      listen for incoming TCP connections and to
      bind outgoing TCP connections.";
  }

  leaf admin-status {
    type pcep-admin-status;
    description
      "The administrative status of this PCEP Entity.
      This is the desired operational status as
      currently set by an operator or by default in
      the implementation. The value of enabled
      represents the current status of an attempt
      to reach this desired status.";
  }
}
```

```
leaf oper-status {
  type pcep-admin-status;
  description
    "The operational status of the PCEP entity.
    Takes one of the following values.
    - oper-status-up(1): the PCEP entity is
      active.
    - oper-status-down(2): the PCEP entity is
      inactive.
    - oper-status-going-up(3): the PCEP entity is
      activating.
    - oper-status-going-down(4): the PCEP entity is
      deactivating.
    - oper-status-failed(5): the PCEP entity has
      failed and will recover when possible.
    - oper-status-failed-perm(6): the PCEP entity
      has failed and will not recover without
      operator intervention.";
}

leaf role {
  type pcep-role;
  description
    "The role that this entity can play.
    Takes one of the following values.
    - unknown(0): this PCEP entity role is
      not known.
    - pcc(1): this PCEP entity is a PCC.
    - pce(2): this PCEP entity is a PCE.
    - pcc-and-pce(3): this PCEP entity is
      both a PCC and a PCE.";
}

container pce-info {
  when "(../role == 'pce') " +
    " or " +
    "(../role == 'pcc-and-pce')";
  {
    description
      "When PCEP entity is PCE";
  }
  uses pce-info {
    description
      "Local PCE information";
  }
  description
    "The Local PCE Entity PCE information";
}
```



```
uses pcep-entity-info{
  description
    "The operational information related to the
    PCEP entity.";
}

container peers{
  description
    "The list of peers for the entity";

  list peer{
    key "addr";

    description
      "The peer for the entity.";

    leaf addr {
      type inet:ip-address;
      description
        "The local Internet address of this PCEP peer.";
    }

    leaf role {
      type pcep-role;
      description
        "The role of the PCEP Peer.
        Takes one of the following values.
        - unknown(0): this PCEP peer role
          is not known.
        - pcc(1): this PCEP peer is a PCC.
        - pce(2): this PCEP peer is a PCE.
        - pcc-and-pce(3): this PCEP peer
          is both a PCC and a PCE.";
    }

    container pce-info {
      uses pce-info {
        description
          "PCE Peer information";
      }
    }
    description
      "The Local PCE Entity PCE information";
  }

  leaf discontinuity-time {
    type yang:timestamp;
  }
}
```

```
        description
            "The timestamp of the time when the
            information and statistics were
            last reset.";
    }

    leaf initiate-session {
        type boolean;
        description
            "Indicates whether the local PCEP
            entity initiates sessions to this peer,
            or waits for the peer to initiate a
            session.";
    }

    leaf session-exists{
        type boolean;
        description
            "Indicates whether a session with
            this peer currently exists.";
    }

    leaf num-sess-setup-ok{
        type yang:counter32;
        description
            "The number of PCEP sessions successfully
            successfully established with the peer,
            including any current session. This
            counter is incremented each time a
            session with this peer is successfully
            established.";
    }

    leaf num-sess-setup-fail{
        type yang:counter32;
        description
            "The number of PCEP sessions with the peer
            that have been attempted but failed
            before being fully established. This
            counter is incremented each time a
            session retry to this peer fails.";
    }

    leaf session-up-time{
        type yang:timestamp;
        must "(../num-sess-setup-ok != 0 or " +
            "(../num-sess-setup-ok = 0 and " +
            "session-up-time = 0))" {

```

```
        error-message
            "Invalid Session Up timestamp";
        description
            "If num-sess-setup-ok is zero,
             then this leaf contains zero.";
    }
    description
        "The timestamp value of the last time a
         session with this peer was successfully
         established.";
}

leaf session-fail-time{
    type yang:timestamp;
    must "(../num-sess-setup-fail != 0 or " +
        "(../num-sess-setup-fail = 0 and " +
        "session-fail-time = 0))" {
        error-message
            "Invalid Session Fail timestamp";
        description
            "If num-sess-setup-fail is zero,
             then this leaf contains zero.";
    }
    description
        "The timestamp value of the last time a
         session with this peer failed to be
         established.";
}

leaf session-fail-up-time{
    type yang:timestamp;
    must "(../num-sess-setup-ok != 0 or " +
        "(../num-sess-setup-ok = 0 and " +
        "session-fail-up-time = 0))" {
        error-message
            "Invalid Session Fail from
             Up timestamp";
        description
            "If num-sess-setup-ok is zero,
             then this leaf contains zero.";
    }
    description
        "The timestamp value of the last time a
         session with this peer failed from
         active.";
}

uses pcep-stats{
```

```
    description
      "Since PCEP sessions can be ephemeral,
       the peer statistics tracks a peer even
       when no PCEP session currently exists
       to that peer. The statistics contained
       are an aggregate of the statistics for
       all successive sessions to that peer.";
  }

  leaf num-req-sent-closed{
    type yang:counter32;
    description
      "The number of requests that were sent
       to the peer and implicitly cancelled
       when the session they were sent over
       was closed.";
  }

  leaf num-req-rcvd-closed{
    type yang:counter32;
    description
      "The number of requests that were
       received from the peer and implicitly
       cancelled when the session they were
       received over was closed.";
  }

  container sessions {
    description
      "This entry represents a single PCEP
       session in which the local PCEP entity
       participates.

       This entry exists only if the
       corresponding PCEP session has been
       initialized by some event, such as
       manual user configuration, auto-
       discovery of a peer, or an incoming
       TCP connection.";

    list session {
      key "initiator";

      description
        "The list of sessions, note that
         for a time being two sessions
         may exist for a peer";
    }
  }
}
```

```
leaf initiator {
  type pcep-initiator;
  description
    "The initiator of the session,
    that is, whether the TCP
    connection was initiated by
    the local PCEP entity or the
    peer.

    There is a window during
    session initialization where
    two sessions can exist between
    a pair of PCEP speakers, each
    initiated by one of the
    speakers. One of these
    sessions is always discarded
    before it leaves OpenWait state.
    However, before it is discarded,
    two sessions to the given peer
    appear transiently in this MIB
    module. The sessions are
    distinguished by who initiated
    them, and so this field is the
    key.";
}

leaf state-last-change {
  type yang:timestamp;
  description
    "The timestamp value at the
    time this session entered its
    current state as denoted by
    the state leaf.";
}

leaf state {
  type pcep-sess-state;
  description
    "The current state of the
    session.

    The set of possible states
    excludes the idle state since
    entries do not exist in the
    idle state.";
}

leaf connect-retry {
```

```
type yang:counter32;
description
  "The number of times that the
  local PCEP entity has
  attempted to establish a TCP
  connection for this session
  without success. The PCEP
  entity gives up when this
  reaches connect-max-retry.";
}

leaf local-id {
  type uint32 {
    range "0..255";
  }
  description
    "The value of the PCEP session
    ID used by the local PCEP
    entity in the Open message
    for this session.

    If state is tcp-pending then
    this is the session ID that
    will be used in the Open
    message. Otherwise, this is
    the session ID that was sent
    in the Open message.";
}

leaf remote-id {
  type uint32 {
    range "0..255";
  }
  must "((../state != 'tcp-pending' " +
  "and " +
  "../state != 'open-wait' )" +
  "or " +
  "((../state = 'tcp-pending' " +
  " or " +
  "../state = 'open-wait' )" +
  "and remote-id = 0))" {
    error-message
      "Invalid remote-id";
    description
      "If state is tcp-pending
      or open-wait then this
      leaf is not used and
      MUST be set to zero.";
  }
}
```

```
    }
    description
        "The value of the PCEP session
        ID used by the peer in its
        Open message for this
        session.";
}

leaf keepalive-timer {
    type uint32 {
        range "0..255";
    }
    units "seconds";
    must "(../state = 'session-up'" +
        "or " +
        "(../state != 'session-up'" +
        "and keepalive-timer = 0))" {
        error-message
            "Invalid keepalive
            timer";
        description
            "This field is used if
            and only if state is
            session-up. Otherwise,
            it is not used and
            MUST be set to
            zero.";
    }
    description
        "The agreed maximum interval at
        which the local PCEP entity
        transmits PCEP messages on this
        PCEP session. Zero means that
        the local PCEP entity never
        sends Keepalives on this
        session.";
}

leaf peer-keepalive-timer {
    type uint32 {
        range "0..255";
    }
    units "seconds";
    must "(../state = 'session-up'" +
        "or " +
        "(../state != 'session-up'" +
        "and " +
        "peer-keepalive-timer = 0))" {
```

```

        error-message
            "Invalid Peer keepalive
            timer";
        description
            "This field is used if
            and only if state is
            session-up. Otherwise,
            it is not used and MUST
            be set to zero.";
    }
    description
        "The agreed maximum interval at
        which the peer transmits PCEP
        messages on this PCEP session.
        Zero means that the peer never
        sends Keepalives on this
        session.";
}

leaf dead-timer {
    type uint32 {
        range "0..255";
    }
    units "seconds";
    description
        "The dead timer interval for
        this PCEP session.";
}

leaf peer-dead-timer {
    type uint32 {
        range "0..255";
    }
    units "seconds";
    must "((../state != 'tcp-pending' " +
        "and " +
        "../state != 'open-wait' )" +
        "or " +
        "(../state = 'tcp-pending' " +
        " or " +
        "../state = 'open-wait' )" +
        "and " +
        "peer-dead-timer = 0))" {
        error-message
            "Invalid Peer Dead
            timer";
        description
            "If state is tcp-

```



```

        pending or open-wait
        then this leaf is not
        used and MUST be set to
        zero.";
    }
    description
        "The peer's dead-timer interval
        for this PCEP session.";
}

leaf ka-hold-time-rem {
    type uint32 {
        range "0..255";
    }
    units "seconds";
    must "((../state != 'tcp-pending' " +
        "and " +
        "../state != 'open-wait' )" +
        "or " +
        "((../state = 'tcp-pending' " +
        "or " +
        "../state = 'open-wait' )" +
        "and " +
        "ka-hold-time-rem = 0))" {
        error-message
            "Invalid Keepalive hold
            time remaining";
        description
            "If state is tcp-pending
            or open-wait then this
            field is not used and
            MUST be set to zero.";
    }
    description
        "The keep alive hold time
        remaining for this session.";
}

leaf overloaded {
    type boolean;
    description
        "If the local PCEP entity has
        informed the peer that it is
        currently overloaded, then this
        is set to true. Otherwise, it
        is set to false.";
}

```

```
leaf overload-time {
  type uint32;
  units "seconds";
  must "(../overloaded = true or" +
    "(../overloaded != true and" +
    " overload-time = 0))" {
    error-message
      "Invalid overload-time";
    description
      "This field is only used
      if overloaded is set to
      true. Otherwise, it is
      not used and MUST be set
      to zero.";
  }
  description
    "The interval of time that is
    remaining until the local PCEP
    entity will cease to be
    overloaded on this session.";
}

leaf peer-overloaded {
  type boolean;
  description
    "If the peer has informed the
    local PCEP entity that it is
    currently overloaded, then this
    is set to true. Otherwise, it
    is set to false.";
}

leaf peer-overload-time {
  type uint32;
  units "seconds";
  must "(../peer-overloaded = true" +
    " or " +
    "(../peer-overloaded != true" +
    " and " +
    "peer-overload-time = 0))" {
    error-message
      "Invalid peer overload
      time";
    description
      "This field is only used
      if peer-overloaded is
      set to true. Otherwise,
      it is not used and MUST
```

```

        be set to zero.";
    }
    description
        "The interval of time that is
        remaining until the peer will
        cease to be overloaded.  If it
        is not known how long the peer
        will stay in overloaded state,
        this leaf is set to zero.";
}

leaf discontinuity-time {
    type yang:timestamp;
    description
        "The timestamp value of the time
        when the statistics were last
        reset.";
}

uses pcep-stats{
    description
        "The statistics contained are
        for the current sessions to that
        peer.  These are lost when the
        session goes down.";
}
} // session
} // sessions
} // peer
} // peers
} // entity
} // pcep-state

/*
 * Notifications
 */
notification pcep-session-up {
    description
        "This notification is sent when the value of
        '/pcep/pcep-state/peers/peer/sessions/session/state'
        enters the 'session-up' state.";

    uses notification-instance-hdr;

    uses notification-session-hdr;
}

```

```
leaf state-last-change {
  type yang:timestamp;
  description
    "The timestamp value at the time this session entered
    its current state as denoted by the state leaf.";
}

leaf state {
  type pcep-sess-state;
  description
    "The current state of the session.

    The set of possible states excludes the idle state
    since entries do not exist in the idle state.";
}
} //notification

notification pcep-session-down {
  description
    "This notification is sent when the value of
    '/pcep/pcep-state/peers/peer/sessions/session/state'
    leaves the 'session-up' state.";

  uses notification-instance-hdr;

  leaf session-initiator {
    type pcep-initiator;
    description
      "The initiator of the session.";
  }

  leaf state-last-change {
    type yang:timestamp;
    description
      "The timestamp value at the time this session entered
      its current state as denoted by the state leaf.";
  }

  leaf state {
    type pcep-sess-state;
    description
      "The current state of the session.

      The set of possible states excludes the idle state
      since entries do not exist in the idle state.";
  }
} //notification
```

```
notification pcep-session-local-overload {
  description
    "This notification is sent when the local PCEP entity
    enters overload state for a peer.";

  uses notification-instance-hdr;

  uses notification-session-hdr;

  leaf overloaded {
    type boolean;
    description
      "If the local PCEP entity has informed the peer that
      it is currently overloaded, then this is set to
      true. Otherwise, it is set to false.";
  }

  leaf overload-time {
    type uint32;
    units "seconds";
    must "(../overloaded = true or " +
      "(../overloaded != true and " +
      "overload-time = 0))" {
      error-message
        "Invalid overload-time";
      description
        "This field is only used if overloaded is
        set to true. Otherwise, it is not used
        and MUST be set to zero.";
    }
    description
      "The interval of time that is remaining until the
      local PCEP entity will cease to be overloaded on
      this session.";
  }
} //notification

notification pcep-session-local-overload-clear {
  description
    "This notification is sent when the local PCEP entity
    leaves overload state for a peer.";

  uses notification-instance-hdr;

  leaf overloaded {
    type boolean;
    description
      "If the local PCEP entity has informed the peer
```

```
        that it is currently overloaded, then this is set
        to true.  Otherwise, it is set to false.";
    }
} //notification

notification pcep-session-peer-overload {
    description
        "This notification is sent when a peer enters overload
        state.";

    uses notification-instance-hdr;

    uses notification-session-hdr;

    leaf peer-overloaded {
        type boolean;
        description
            "If the peer has informed the local PCEP entity that
            it is currently overloaded, then this is set to true.
            Otherwise, it is set to false.";
    }

    leaf peer-overload-time {
        type uint32;
        units "seconds";
        must "(../peer-overloaded = true or " +
            "(../peer-overloaded != true and " +
            "peer-overload-time = 0))" {
            error-message
                "Invalid peer-overload-time";
            description
                "This field is only used if
                peer-overloaded is set to true.
                Otherwise, it is not used and MUST
                be set to zero.";
        }
        description
            "The interval of time that is remaining until the
            peer will cease to be overloaded.  If it is not known
            how long the peer will stay in overloaded state, this
            leaf is set to zero.";
    }
} //notification

notification pcep-session-peer-overload-clear {
    description
        "This notification is sent when a peer leaves overload
        state.";
```

```
    uses notification-instance-hdr;

    leaf peer-overloaded {
        type boolean;
        description
            "If the peer has informed the local PCEP entity that
             it is currently overloaded, then this is set to true.
             Otherwise, it is set to false.";
    }
} //notification
} //module
```

<CODE ENDS>

8. Security Considerations

The YANG module defined in this memo is designed to be accessed via the NETCONF protocol [RFC6241]. The lowest NETCONF layer is the secure transport layer and the mandatory-to-implement secure transport is SSH [RFC6242]. The NETCONF access control model [RFC6536] provides the means to restrict access for particular NETCONF users to a pre-configured subset of all available NETCONF protocol operations and content.

There are a number of data nodes defined in the YANG module which are writable/creatable/deletable (i.e., config true, which is the default). These data nodes may be considered sensitive or vulnerable in some network environments. Write operations (e.g., <edit-config>) to these data nodes without proper protection can have a negative effect on network operations.

TBD: List specific Subtrees and data nodes and their sensitivity/vulnerability.

9. Manageability Considerations

9.1. Control of Function and Policy

9.2. Information and Data Models

9.3. Liveness Detection and Monitoring

9.4. Verify Correct Operations

9.5. Requirements On Other Protocols

9.6. Impact On Network Operations

10. IANA Considerations

This document registers a URI in the "IETF XML Registry" [RFC3688]. Following the format in RFC 3688, the following registration has been made.

URI: urn:ietf:params:xml:ns:yang:ietf-pcep

Registrant Contact: The PCE WG of the IETF.

XML: N/A; the requested URI is an XML namespace.

This document registers a YANG module in the "YANG Module Names" registry [RFC6020].

Name:	ietf-pcep
Namespace:	urn:ietf:params:xml:ns:yang:ietf-pcep
Prefix:	pcep
Reference:	This I-D

11. Acknowledgements

The initial document is based on the PCEP MIB [RFC7420]. Further this document structure is based on Routing Yang Module [I-D.ietf-netmod-routing-cfg]. We would like to thank the authors of aforementioned documents.

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A YANG Data Model for Path Computation Element Communications Protocol
(PCEP)
draft-pkd-pce-pcep-yang-06

Abstract

This document defines a YANG data model for the management of Path Computation Element communications Protocol (PCEP) for communications between a Path Computation Client (PCC) and a Path Computation Element (PCE), or between two PCEs. The data model includes configuration data and state data (status information and counters for the collection of statistics).

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1. Introduction

The Path Computation Element (PCE) defined in [RFC4655] is an entity that is capable of computing a network path or route based on a network graph, and applying computational constraints. A Path

Computation Client (PCC) may make requests to a PCE for paths to be computed.

PCEP is the communication protocol between a PCC and PCE and is defined in [RFC5440]. PCEP interactions include path computation requests and path computation replies as well as notifications of specific states related to the use of a PCE in the context of Multiprotocol Label Switching (MPLS) and Generalized MPLS (GMPLS) Traffic Engineering (TE). [I-D.ietf-pce-stateful-pce] specifies extensions to PCEP to enable stateful control of MPLS TE LSPs.

This document defines a YANG [RFC6020] data model for the management of PCEP speakers. It is important to establish a common data model for how PCEP speakers are identified, configured, and monitored. The data model includes configuration data and state data (status information and counters for the collection of statistics).

This document contains a specification of the PCEP YANG module, "ietf-pcep" which provides the PCEP [RFC5440] data model.

2. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

3. Terminology and Notation

This document uses the terminology defined in [RFC4655] and [RFC5440]. In particular, it uses the following acronyms.

- o Path Computation Request message (PCReq).
- o Path Computation Reply message (PCRep).
- o Notification message (PCNtf).
- o Error message (PCErr).
- o Request Parameters object (RP).
- o Synchronization Vector object (SVEC).
- o Explicit Route object (ERO).

This document also uses the following terms defined in [RFC7420]:

- o PCEP entity: a local PCEP speaker.

- o PCEP peer: to refer to a remote PCEP speaker.
- o PCEP speaker: where it is not necessary to distinguish between local and remote.

Further, this document also uses the following terms defined in [I-D.ietf-pce-stateful-pce] :

- o Stateful PCE, Passive Stateful PCE, Active Stateful PCE
- o Delegation, Revocation, Redelegation
- o LSP State Report, Path Computation Report message (PCRpt).
- o LSP State Update, Path Computation Update message (PCUpd).

[I-D.ietf-pce-pce-initiated-lsp] :

- o PCE-initiated LSP, Path Computation LSP Initiate Message (PCInitiate).

[I-D.ietf-pce-lsp-setup-type] :

- o Path Setup Type (PST).

[I-D.ietf-pce-segment-routing] :

- o Segment Routing (SR).
- o Segment Identifier (SID).
- o Maximum SID Depth (MSD).

3.1. Tree Diagrams

A graphical representation of the complete data tree is presented in Section 5. The meaning of the symbols in these diagrams is as follows and as per [I-D.ietf-netmod-rfc6087bis]:

- o Brackets "[" and "]" enclose list keys.
- o Curly braces "{" and "}" contain names of optional features that make the corresponding node conditional.
- o Abbreviations before data node names: "rw" means configuration (read-write), and "ro" state data (read-only).

- o Symbols after data node names: "?" means an optional node and "*" denotes a "list" or "leaf-list".
- o Parentheses enclose choice and case nodes, and case nodes are also marked with a colon (":").
- o Ellipsis ("...") stands for contents of subtrees that are not shown.

3.2. Prefixes in Data Node Names

In this document, names of data nodes and other data model objects are often used without a prefix, as long as it is clear from the context in which YANG module each name is defined. Otherwise, names are prefixed using the standard prefix associated with the corresponding YANG module, as shown in Table 1.

Prefix	YANG module	Reference
yang	ietf-yang-types	[RFC6991]
inet	ietf-inet-types	[RFC6991]

Table 1: Prefixes and corresponding YANG modules

4. Objectives

This section describes some of the design objectives for the model:

- o In case of existing implementations, it needs to map the data model defined in this document to their proprietary native data model. To facilitate such mappings, the data model should be simple.
- o The data model should be suitable for new implementations to use as is.
- o Mapping to the PCEP MIB Module should be clear.
- o The data model should allow for static configurations of peers.
- o The data model should include read-only counters in order to gather statistics for sent and received PCEP messages, received messages with errors, and messages that could not be sent due to errors.

- o It should be fairly straightforward to augment the base data model for advanced PCE features.

5. The Design of PCEP Data Model

The module, "ietf-pcep", defines the basic components of a PCE speaker.

```

module: ietf-pcep
+--rw pcep!
|   +--rw entity
|   |   +--rw addr                inet:ip-address
|   |   +--rw enabled?           boolean
|   |   +--rw role                pcep-role
|   |   +--rw description?       string
|   |   +--rw domain
|   |   |   +--rw domain* [domain-type domain]
|   |   |   |   +--rw domain-type  domain-type
|   |   |   |   +--rw domain      domain
|   |   +--rw capability
|   |   |   +--rw gmpls?          boolean {gmpls}?
|   |   |   +--rw bi-dir?        boolean
|   |   |   +--rw diverse?       boolean
|   |   |   +--rw load-balance?  boolean
|   |   |   +--rw synchronize?   boolean {svec}?
|   |   |   +--rw objective-function? boolean {obj-fn}?
|   |   |   +--rw add-path-constraint? boolean
|   |   |   +--rw prioritization? boolean
|   |   |   +--rw multi-request?  boolean
|   |   |   +--rw gco?           boolean {gco}?
|   |   |   +--rw p2mp?         boolean {p2mp}?
|   |   |   +--rw stateful {stateful}?
|   |   |   |   +--rw enabled?    boolean
|   |   |   |   +--rw active?     boolean
|   |   |   |   +--rw pce-initiated? boolean {pce-initiated}?
|   |   +--rw sr {sr}?
|   |   |   +--rw enabled?    boolean
|   |   |   +--rw msd?       uint8
|   +--rw pce-info
|   |   +--rw scope
|   |   |   +--rw intra-area-scope?    boolean
|   |   |   +--rw intra-area-pref?    uint8
|   |   |   +--rw inter-area-scope?   boolean
|   |   |   +--rw inter-area-scope-default? boolean
|   |   |   +--rw inter-area-pref?   uint8
|   |   |   +--rw inter-as-scope?     boolean
|   |   |   +--rw inter-as-scope-default? boolean
|   |   |   +--rw inter-as-pref?     uint8

```

```

| | | +--rw inter-layer-scope?          boolean
| | | +--rw inter-layer-pref?          uint8
+--rw neigh-domains
| | | +--rw domain* [domain-type domain]
| | | | +--rw domain-type      domain-type
| | | | +--rw domain          domain
+--rw (auth-type-selection)?
| | | +---:(auth-key-chain)
| | | | +--rw key-chain?          key-chain:key-chain-ref
+---:(auth-key)
| | | | +--rw key?                string
| | | | +--rw crypto-algorithm
| | | | | +--rw (algorithm)?
| | | | | | +---:(hmac-sha-1-12) {crypto-hmac-sha-1-12}?
| | | | | | | +--rw hmac-sha1-12?      empty
| | | | | | +---:(aes-cmac-prf-128) {aes-cmac-prf-128}?
| | | | | | | +--rw aes-cmac-prf-128?   empty
| | | | | +---:(md5)
| | | | | | +--rw md5?                empty
| | | | | +---:(sha-1)
| | | | | | +--rw sha-1?              empty
| | | | | +---:(hmac-sha-1)
| | | | | | +--rw hmac-sha-1?         empty
| | | | | +---:(hmac-sha-256)
| | | | | | +--rw hmac-sha-256?       empty
| | | | | +---:(hmac-sha-384)
| | | | | | +--rw hmac-sha-384?       empty
| | | | | +---:(hmac-sha-512)
| | | | | | +--rw hmac-sha-512?       empty
| | | | | +---:(clear-text) {clear-text}?
| | | | | | +--rw clear-text?         empty
| | | | | +---:(replay-protection-only) {replay-protection-only}?
| | | | | | +--rw replay-protection-only? empty
+---:(auth-tls) {tls}?
| | | | +--rw tls
+--rw connect-timer?          uint32
+--rw connect-max-retry?     uint32
+--rw init-backoff-timer?    uint32
+--rw max-backoff-timer?     uint32
+--rw open-wait-timer?       uint32
+--rw keep-wait-timer?       uint32
+--rw keep-alive-timer?      uint32
+--rw dead-timer?            uint32
+--rw allow-negotiation?     boolean
+--rw max-keep-alive-timer?  uint32
+--rw max-dead-timer?        uint32
+--rw min-keep-alive-timer?  uint32
+--rw min-dead-timer?        uint32

```

```

+--rw sync-timer?                uint32 {svec}?
+--rw request-timer?            uint32
+--rw max-sessions?             uint32
+--rw max-unknown-reqs?        uint32
+--rw max-unknown-msgs?        uint32
+--rw pcep-notification-max-rate uint32
+--rw stateful-parameter {stateful}?
|   +--rw state-timeout?        uint32
|   +--rw redelegation-timeout? uint32
|   +--rw rpt-non-pcep-lsp?     boolean
+--rw peers
|   +--rw peer* [addr]
|   |   +--rw addr                inet:ip-address
|   |   +--rw description?       string
|   |   +--rw domain
|   |   |   +--rw domain* [domain-type domain]
|   |   |   |   +--rw domain-type  domain-type
|   |   |   |   +--rw domain      domain
|   |   +--rw capability
|   |   |   +--rw gmpls?          boolean {gmpls}?
|   |   |   +--rw bi-dir?        boolean
|   |   |   +--rw diverse?       boolean
|   |   |   +--rw load-balance?   boolean
|   |   |   +--rw synchronize?   boolean {svec}?
|   |   |   +--rw objective-function? boolean {obj-fn}?
|   |   |   +--rw add-path-constraint? boolean
|   |   |   +--rw prioritization? boolean
|   |   |   +--rw multi-request?  boolean
|   |   |   +--rw gco?           boolean {gco}?
|   |   |   +--rw p2mp?          boolean {p2mp}?
|   |   |   +--rw stateful {stateful}?
|   |   |   |   +--rw enabled?    boolean
|   |   |   |   +--rw active?    boolean
|   |   |   |   +--rw pce-initiated? boolean {pce-initiated}?
|   |   +--rw sr {sr}?
|   |   |   +--rw enabled?    boolean
|   |   |   +--rw msd?        uint8
+--rw scope
|   +--rw intra-area-scope?      boolean
|   +--rw intra-area-pref?       uint8
|   +--rw inter-area-scope?      boolean
|   +--rw inter-area-scope-default? boolean
|   +--rw inter-area-pref?       uint8
|   +--rw inter-as-scope?        boolean
|   +--rw inter-as-scope-default? boolean
|   +--rw inter-as-pref?         uint8
|   +--rw inter-layer-scope?     boolean
|   +--rw inter-layer-pref?      uint8

```

```

|      +--rw neigh-domains
|      |      +--rw domain* [domain-type domain]
|      |      |      +--rw domain-type    domain-type
|      |      |      +--rw domain        domain
|      +--rw delegation-pref?    uint8 {stateful}?
|      +--rw (auth-type-selection)?
|      |      +--:(auth-key-chain)
|      |      |      +--rw key-chain?      key-chain:key-chain-ref
|      |      +--:(auth-key)
|      |      |      +--rw key?            string
|      |      +--rw crypto-algorithm
|      |      |      +--rw (algorithm)?
|      |      |      |      +--:(hmac-sha-1-12) {crypto-hmac-sha-1-12}?
|      |      |      |      |      +--rw hmac-sha1-12?          empty
|      |      |      |      +--:(aes-cmac-prf-128) {aes-cmac-prf-128}?
|      |      |      |      |      +--rw aes-cmac-prf-128?      empty
|      |      |      |      +--:(md5)
|      |      |      |      |      +--rw md5?                    empty
|      |      |      |      +--:(sha-1)
|      |      |      |      |      +--rw sha-1?                  empty
|      |      |      |      +--:(hmac-sha-1)
|      |      |      |      |      +--rw hmac-sha-1?              empty
|      |      |      |      +--:(hmac-sha-256)
|      |      |      |      |      +--rw hmac-sha-256?           empty
|      |      |      |      +--:(hmac-sha-384)
|      |      |      |      |      +--rw hmac-sha-384?           empty
|      |      |      |      +--:(hmac-sha-512)
|      |      |      |      |      +--rw hmac-sha-512?           empty
|      |      |      |      +--:(clear-text) {clear-text}?
|      |      |      |      |      +--rw clear-text?              empty
|      |      |      |      +--:(replay-protection-only) {replay-protection-only}
|      |      |      |      |      +--rw replay-protection-only?  empty
|      |      |      +--rw replay-protection-only?  empty
|      |      +--:(auth-tls) {tls}?
|      |      |      +--rw tls
|      +--ro pcep-state
|      |      +--ro entity
|      |      |      +--ro addr?            inet:ip-address
|      |      |      +--ro index?          uint32
|      |      |      +--ro admin-status?    pcep-admin-status
|      |      |      +--ro oper-status?     pcep-admin-status
|      |      |      +--ro role?            pcep-role
|      |      |      +--ro domain
|      |      |      |      +--ro domain* [domain-type domain]
|      |      |      |      |      +--ro domain-type    domain-type
|      |      |      |      |      +--ro domain        domain
|      |      |      +--ro capability
|      |      |      |      +--ro gmpls?      boolean {gmpls}?
|      |      |      |      +--ro bi-dir?     boolean

```

```

|--ro diverse?                boolean
|--ro load-balance?          boolean
|--ro synchronize?          boolean {svec}?
|--ro objective-function?    boolean {obj-fn}?
|--ro add-path-constraint?   boolean
|--ro prioritization?        boolean
|--ro multi-request?         boolean
|--ro gco?                   boolean {gco}?
|--ro p2mp?                  boolean {p2mp}?
|--ro stateful {stateful}?
|   |--ro enabled?           boolean
|   |--ro active?            boolean
|   |--ro pce-initiated?    boolean {pce-initiated}?
|--ro sr {sr}?
|   |--ro enabled?          boolean
|   |--ro msd?              uint8
+--ro pce-info
|   +--ro scope
|   |   |--ro intra-area-scope?        boolean
|   |   |--ro intra-area-pref?        uint8
|   |   |--ro inter-area-scope?       boolean
|   |   |--ro inter-area-scope-default? boolean
|   |   |--ro inter-area-pref?       uint8
|   |   |--ro inter-as-scope?         boolean
|   |   |--ro inter-as-scope-default? boolean
|   |   |--ro inter-as-pref?         uint8
|   |   |--ro inter-layer-scope?     boolean
|   |   |--ro inter-layer-pref?     uint8
|   +--ro neigh-domains
|   |   |--ro domain* [domain-type domain]
|   |   |   |--ro domain-type    domain-type
|   |   |   |--ro domain        domain
|   +--ro (auth-type-selection)?
|   |   +--:(auth-key-chain)
|   |   |   |--ro key-chain?        key-chain:key-chain-ref
|   |   +--:(auth-key)
|   |   |   |--ro key?              string
|   |   |   +--ro crypto-algorithm
|   |   |   |   +--ro (algorithm)?
|   |   |   |   |   +--:(hmac-sha-1-12) {crypto-hmac-sha-1-12}?
|   |   |   |   |   |   |--ro hmac-sha1-12?          empty
|   |   |   |   |   +--:(aes-cmac-prf-128) {aes-cmac-prf-128}?
|   |   |   |   |   |   |--ro aes-cmac-prf-128?      empty
|   |   |   |   +--:(md5)
|   |   |   |   |   |--ro md5?                    empty
|   |   |   |   +--:(sha-1)
|   |   |   |   |   |--ro sha-1?                  empty
|   |   |   +--:(hmac-sha-1)

```

```

|         |         |   +--ro hmac-sha-1?           empty
|         |         |   +---:(hmac-sha-256)
|         |         |   |   +--ro hmac-sha-256?           empty
|         |         |   |   +---:(hmac-sha-384)
|         |         |   |   |   +--ro hmac-sha-384?        empty
|         |         |   |   |   +---:(hmac-sha-512)
|         |         |   |   |   |   +--ro hmac-sha-512?    empty
|         |         |   |   |   |   +---:(clear-text) {clear-text}?
|         |         |   |   |   |   |   +--ro clear-text?    empty
|         |         |   |   |   |   |   +---:(replay-protection-only) {replay-protection-only}?
|         |         |   |   |   |   |   |   +--ro replay-protection-only? empty
|         |         |   |   |   |   |   |   +---:(auth-tls) {tls}?
|         |         |   |   |   |   |   |   |   +--ro tls
|         |         |   |   |   |   |   |   |   +---ro connect-timer?          uint32
|         |         |   |   |   |   |   |   |   +---ro connect-max-retry?       uint32
|         |         |   |   |   |   |   |   |   +---ro init-backoff-timer?        uint32
|         |         |   |   |   |   |   |   |   +---ro max-backoff-timer?         uint32
|         |         |   |   |   |   |   |   |   +---ro open-wait-timer?           uint32
|         |         |   |   |   |   |   |   |   +---ro keep-wait-timer?           uint32
|         |         |   |   |   |   |   |   |   +---ro keep-alive-timer?          uint32
|         |         |   |   |   |   |   |   |   +---ro dead-timer?                uint32
|         |         |   |   |   |   |   |   |   +---ro allow-negotiation?          boolean
|         |         |   |   |   |   |   |   |   +---ro max-keep-alive-timer?       uint32
|         |         |   |   |   |   |   |   |   +---ro max-dead-timer?            uint32
|         |         |   |   |   |   |   |   |   +---ro min-keep-alive-timer?       uint32
|         |         |   |   |   |   |   |   |   +---ro min-dead-timer?            uint32
|         |         |   |   |   |   |   |   |   +---ro sync-timer?                uint32 {svec}?
|         |         |   |   |   |   |   |   |   +---ro request-timer?             uint32
|         |         |   |   |   |   |   |   |   +---ro max-sessions?              uint32
|         |         |   |   |   |   |   |   |   +---ro max-unknown-reqs?           uint32
|         |         |   |   |   |   |   |   |   +---ro max-unknown-msgs?          uint32
|         |         |   |   |   |   |   |   |   +---ro stateful-parameter {stateful}?
|         |         |   |   |   |   |   |   |   |   +--ro state-timeout?            uint32
|         |         |   |   |   |   |   |   |   |   +--ro redelegation-timeout?     uint32
|         |         |   |   |   |   |   |   |   |   +--ro rpt-non-pcep-lsp?        boolean
|         |         |   |   |   |   |   |   |   +---ro lsp-db {stateful}?
|         |         |   |   |   |   |   |   |   |   +--ro association-list* [id source global-source extended-id]
|         |         |   |   |   |   |   |   |   |   |   +--ro type?                assoc-type
|         |         |   |   |   |   |   |   |   |   |   +--ro id                  uint16
|         |         |   |   |   |   |   |   |   |   |   +--ro source               inet:ip-address
|         |         |   |   |   |   |   |   |   |   |   +--ro global-source         uint32
|         |         |   |   |   |   |   |   |   |   |   +--ro extended-id          string
|         |         |   |   |   |   |   |   |   |   |   +--ro lsp* [plsp-id pcc-id]
|         |         |   |   |   |   |   |   |   |   |   |   +--ro plsp-id         -> /pcep-state/entity/lsp-db/lsp/plsp-id
|         |         |   |   |   |   |   |   |   |   |   |   +--ro pcc-id         -> /pcep-state/entity/lsp-db/lsp/pcc-id
|         |         |   |   |   |   |   |   |   |   +---ro lsp* [plsp-id pcc-id]
|         |         |   |   |   |   |   |   |   |   |   +--ro plsp-id             uint32
|         |         |   |   |   |   |   |   |   |   |   +--ro pcc-id             inet:ip-address

```

```

|
|   +--ro lsp-ref
|   |   +--ro source?          -> /te:te/lsp-state/lsp/source
|   |   +--ro destination?     -> /te:te/lsp-state/lsp/destinati
on
|
|   +--ro tunnel-id?          -> /te:te/lsp-state/lsp/tunnel-id
|   +--ro lsp-id?            -> /te:te/lsp-state/lsp/lsp-id
tunnel-id
|   +--ro extended-tunnel-id? -> /te:te/lsp-state/lsp/extended-
|
|   +--ro type?                -> /te:te/lsp-state/lsp/type
+--ro admin-state?            boolean
+--ro operational-state?      operational-state
+--ro delegated
|   +--ro enabled?            boolean
|   +--ro pce?                -> /pcep-state/entity/peers/peer/addr
|   +--ro srp-id?             uint32
+--ro initiation {pce-initiated}?
|   +--ro enabled?            boolean
|   +--ro pce?                -> /pcep-state/entity/peers/peer/addr
+--ro symbolic-path-name?     string
+--ro last-error?             lsp-error
+--ro pst?                    pst
+--ro association-list* [id source global-source extended-id]
n-list/id
|   +--ro id                   -> /pcep-state/entity/lsp-db/associatio
n-list/source
|   +--ro source               -> /pcep-state/entity/lsp-db/associatio
n-list/global-source
|   +--ro global-source        -> /pcep-state/entity/lsp-db/associatio
n-list/extended-id
|   +--ro extended-id         -> /pcep-state/entity/lsp-db/associatio
+--ro peers
+--ro peer* [addr]
|   +--ro addr                 inet:ip-address
|   +--ro role?                pcep-role
+--ro domain
|   +--ro domain* [domain-type domain]
|   |   +--ro domain-type      domain-type
|   |   +--ro domain           domain
+--ro capability
|   +--ro gmpls?               boolean {gmpls}?
|   +--ro bi-dir?              boolean
|   +--ro diverse?             boolean
|   +--ro load-balance?        boolean
|   +--ro synchronize?         boolean {svec}?
|   +--ro objective-function?  boolean {obj-fn}?
|   +--ro add-path-constraint? boolean
|   +--ro prioritization?      boolean
|   +--ro multi-request?       boolean
|   +--ro gco?                 boolean {gco}?
|   +--ro p2mp?                boolean {p2mp}?
+--ro stateful {stateful}?
|   +--ro enabled?             boolean
|   +--ro active?              boolean
|   +--ro pce-initiated?      boolean {pce-initiated}?

```

```

|   +--ro sr {sr}?
|       +--ro enabled?    boolean
|       +--ro msd?       uint8
+--ro pce-info
|   +--ro scope
|       +--ro intra-area-scope?    boolean
|       +--ro intra-area-pref?    uint8
|       +--ro inter-area-scope?    boolean
|       +--ro inter-area-scope-default?    boolean
|       +--ro inter-area-pref?    uint8
|       +--ro inter-as-scope?    boolean
|       +--ro inter-as-scope-default?    boolean
|       +--ro inter-as-pref?    uint8
|       +--ro inter-layer-scope?    boolean
|       +--ro inter-layer-pref?    uint8
+--ro neigh-domains
|   +--ro domain* [domain-type domain]
|       +--ro domain-type    domain-type
|       +--ro domain        domain
+--ro delegation-pref?    uint8 {stateful}?
+--ro (auth-type-selection)?
|   +--:(auth-key-chain)
|   |   +--ro key-chain?    key-chain:key-chain-ref
+--:(auth-key)
|   +--ro key?    string
|   +--ro crypto-algorithm
|       +--ro (algorithm)?
|           +--:(hmac-sha-1-12) {crypto-hmac-sha-1-12}?
|           |   +--ro hmac-sha1-12?    empty
|           +--:(aes-cmac-prf-128) {aes-cmac-prf-128}?
|           |   +--ro aes-cmac-prf-128?    empty
|           +--:(md5)
|           |   +--ro md5?    empty
|           +--:(sha-1)
|           |   +--ro sha-1?    empty
|           +--:(hmac-sha-1)
|           |   +--ro hmac-sha-1?    empty
|           +--:(hmac-sha-256)
|           |   +--ro hmac-sha-256?    empty
|           +--:(hmac-sha-384)
|           |   +--ro hmac-sha-384?    empty
|           +--:(hmac-sha-512)
|           |   +--ro hmac-sha-512?    empty
|           +--:(clear-text) {clear-text}?
|           |   +--ro clear-text?    empty
|           +--:(replay-protection-only) {replay-protection-only}
|           |   +--ro replay-protection-only?    empty
+--:(auth-tls) {tls}?

```

?


```

|      +--ro tls
+--ro discontinuity-time?      yang:timestamp
+--ro initiate-session?      boolean
+--ro session-exists?        boolean
+--ro num-sess-setup-ok?     yang:counter32
+--ro num-sess-setup-fail?   yang:counter32
+--ro session-up-time?       yang:timestamp
+--ro session-fail-time?     yang:timestamp
+--ro session-fail-up-time?  yang:timestamp
+--ro pcep-stats
|   +--ro avg-rsp-time?      uint32
|   +--ro lwm-rsp-time?     uint32
|   +--ro hwm-rsp-time?     uint32
|   +--ro num-pcreq-sent?    yang:counter32
|   +--ro num-pcreq-rcvd?    yang:counter32
|   +--ro num-pcrep-sent?    yang:counter32
|   +--ro num-pcrep-rcvd?    yang:counter32
|   +--ro num-pcerr-sent?    yang:counter32
|   +--ro num-pcerr-rcvd?    yang:counter32
|   +--ro num-pcntf-sent?    yang:counter32
|   +--ro num-pcntf-rcvd?    yang:counter32
|   +--ro num-keepalive-sent? yang:counter32
|   +--ro num-keepalive-rcvd? yang:counter32
|   +--ro num-unknown-rcvd?  yang:counter32
|   +--ro num-corrupt-rcvd?  yang:counter32
|   +--ro num-req-sent?      yang:counter32
|   +--ro num-req-sent-pend-rep? yang:counter32
|   +--ro num-req-sent-ero-rcvd? yang:counter32
|   +--ro num-req-sent-nopath-rcvd? yang:counter32
|   +--ro num-req-sent-cancel-rcvd? yang:counter32
|   +--ro num-req-sent-error-rcvd? yang:counter32
|   +--ro num-req-sent-timeout? yang:counter32
|   +--ro num-req-sent-cancel-sent? yang:counter32
|   +--ro num-req-rcvd?      yang:counter32
|   +--ro num-req-rcvd-pend-rep? yang:counter32
|   +--ro num-req-rcvd-ero-sent? yang:counter32
|   +--ro num-req-rcvd-nopath-sent? yang:counter32
|   +--ro num-req-rcvd-cancel-sent? yang:counter32
|   +--ro num-req-rcvd-error-sent? yang:counter32
|   +--ro num-req-rcvd-cancel-rcvd? yang:counter32
|   +--ro num-rep-rcvd-unknown? yang:counter32
|   +--ro num-req-rcvd-unknown? yang:counter32
|   +--ro svec {svec}?
|   |   +--ro num-svec-sent?      yang:counter32
|   |   +--ro num-svec-req-sent?  yang:counter32
|   |   +--ro num-svec-rcvd?      yang:counter32
|   |   +--ro num-svec-req-rcvd?  yang:counter32
+--ro stateful {stateful}?

```

```

+--ro num-pcrpt-sent?          yang:counter32
+--ro num-pcrpt-rcvd?        yang:counter32
+--ro num-pcupd-sent?        yang:counter32
+--ro num-pcupd-rcvd?        yang:counter32
+--ro num-rpt-sent?          yang:counter32
+--ro num-rpt-rcvd?          yang:counter32
+--ro num-rpt-rcvd-error-sent? yang:counter32
+--ro num-upd-sent?          yang:counter32
+--ro num-upd-rcvd?          yang:counter32
+--ro num-upd-rcvd-unknown?  yang:counter32
+--ro num-upd-rcvd-undelegated? yang:counter32
+--ro num-upd-rcvd-error-sent? yang:counter32
+--ro initiation {pce-initiated}?
  +--ro num-pcinitiate-sent?      yang:counter32
  +--ro num-pcinitiate-rcvd?      yang:counter32
  +--ro num-initiate-sent?        yang:counter32
  +--ro num-initiate-rcvd?        yang:counter32
  +--ro num-initiate-rcvd-error-sent? yang:counter32
+--ro num-req-sent-closed?      yang:counter32
+--ro num-req-rcvd-closed?     yang:counter32
+--ro sessions
  +--ro session* [initiator]
    +--ro initiator              pcep-initiator
    +--ro state-last-change?     yang:timestamp
    +--ro state?                 pcep-sess-state
    +--ro session-creation?      yang:timestamp
    +--ro connect-retry?         yang:counter32
    +--ro local-id?              uint32
    +--ro remote-id?             uint32
    +--ro keepalive-timer?       uint32
    +--ro peer-keepalive-timer?  uint32
    +--ro dead-timer?            uint32
    +--ro peer-dead-timer?       uint32
    +--ro ka-hold-time-rem?      uint32
    +--ro overloaded?            boolean
    +--ro overload-time?         uint32
    +--ro peer-overloaded?       boolean
    +--ro peer-overload-time?    uint32
    +--ro lspdb-sync?            sync-state {stateful}?
    +--ro discontinuity-time?    yang:timestamp
    +--ro pcep-stats
      +--ro avg-rsp-time?        uint32
      +--ro lwm-rsp-time?        uint32
      +--ro hwm-rsp-time?        uint32
      +--ro num-pcreq-sent?      yang:counter32
      +--ro num-pcreq-rcvd?      yang:counter32
      +--ro num-pcrep-sent?      yang:counter32
      +--ro num-pcrep-rcvd?      yang:counter32

```

```

+--ro num-pcerr-sent?          yang:counter32
+--ro num-pcerr-rcvd?         yang:counter32
+--ro num-pcntf-sent?         yang:counter32
+--ro num-pcntf-rcvd?         yang:counter32
+--ro num-keepalive-sent?     yang:counter32
+--ro num-keepalive-rcvd?    yang:counter32
+--ro num-unknown-rcvd?      yang:counter32
+--ro num-corrupt-rcvd?      yang:counter32
+--ro num-req-sent?           yang:counter32
+--ro num-req-sent-pend-rep?  yang:counter32
+--ro num-req-sent-ero-rcvd?  yang:counter32
+--ro num-req-sent-nopath-rcvd? yang:counter32
+--ro num-req-sent-cancel-rcvd? yang:counter32
+--ro num-req-sent-error-rcvd? yang:counter32
+--ro num-req-sent-timeout?   yang:counter32
+--ro num-req-sent-cancel-sent? yang:counter32
+--ro num-req-rcvd?           yang:counter32
+--ro num-req-rcvd-pend-rep?  yang:counter32
+--ro num-req-rcvd-ero-sent?  yang:counter32
+--ro num-req-rcvd-nopath-sent? yang:counter32
+--ro num-req-rcvd-cancel-sent? yang:counter32
+--ro num-req-rcvd-error-sent? yang:counter32
+--ro num-req-rcvd-cancel-rcvd? yang:counter32
+--ro num-rep-rcvd-unknown?   yang:counter32
+--ro num-req-rcvd-unknown?   yang:counter32
+--ro svec {svec}?
|   +--ro num-svec-sent?       yang:counter32
|   +--ro num-svec-req-sent?   yang:counter32
|   +--ro num-svec-rcvd?      yang:counter32
|   +--ro num-svec-req-rcvd?   yang:counter32
+--ro stateful {stateful}?
+--ro num-pcrpt-sent?         yang:counter32
+--ro num-pcrpt-rcvd?         yang:counter32
+--ro num-pcupd-sent?         yang:counter32
+--ro num-pcupd-rcvd?         yang:counter32
+--ro num-rpt-sent?           yang:counter32
+--ro num-rpt-rcvd?           yang:counter32
+--ro num-rpt-rcvd-error-sent? yang:counter32
+--ro num-upd-sent?           yang:counter32
+--ro num-upd-rcvd?           yang:counter32
+--ro num-upd-rcvd-unknown?   yang:counter32
+--ro num-upd-rcvd-undelegated? yang:counter32
+--ro num-upd-rcvd-error-sent? yang:counter32
+--ro initiation {pce-initiated}?
+--ro num-pcinitiate-sent?    yang:counter
32
+--ro num-pcinitiate-rcvd?    yang:counter
32
+--ro num-initiate-sent?      yang:counter
32
+--ro num-initiate-rcvd?      yang:counter
32

```

```

32                                     +--ro num-initiate-rcvd-error-sent?  yang:counter
notifications:
  +---n pcep-session-up
  |   +--ro peer-addr?                -> /pcep-state/entity/peers/peer/addr
  |   +--ro session-initiator?       -> /pcep-state/entity/peers/peer/sessions/sessi
on/initiator
  |   +--ro state-last-change?       yang:timestamp
  |   +--ro state?                   pcep-sess-state
  +---n pcep-session-down
  |   +--ro peer-addr?                -> /pcep-state/entity/peers/peer/addr
  |   +--ro session-initiator?       pcep-initiator
  |   +--ro state-last-change?       yang:timestamp
  |   +--ro state?                   pcep-sess-state
  +---n pcep-session-local-overload
  |   +--ro peer-addr?                -> /pcep-state/entity/peers/peer/addr
  |   +--ro session-initiator?       -> /pcep-state/entity/peers/peer/sessions/sessi
on/initiator
  |   +--ro overloaded?               boolean
  |   +--ro overload-time?           uint32
  +---n pcep-session-local-overload-clear
  |   +--ro peer-addr?                -> /pcep-state/entity/peers/peer/addr
  |   +--ro overloaded?              boolean
  +---n pcep-session-peer-overload
  |   +--ro peer-addr?                -> /pcep-state/entity/peers/peer/addr
  |   +--ro session-initiator?       -> /pcep-state/entity/peers/peer/sessions/sessi
on/initiator
  |   +--ro peer-overloaded?          boolean
  |   +--ro peer-overload-time?      uint32
  +---n pcep-session-peer-overload-clear
  |   +--ro peer-addr?                -> /pcep-state/entity/peers/peer/addr
  |   +--ro peer-overloaded?          boolean

```

5.1. The Entity

The PCEP yang module may contain status information for the local PCEP entity.

The entity has an IP address (using `ietf-inet-types` [RFC6991]) and a "role" leaf (the local entity PCEP role) as mandatory.

Note that, the PCEP MIB module [RFC7420] uses an entity list and a system generated entity index as a primary index to the read only entity table. If the device implements the PCEP MIB, the "index" leaf MUST contain the value of the corresponding `pcepEntityIndex` and only one entity is assumed.

5.2. The Peer Lists

The peer list contains peer(s) that the local PCEP entity knows about. A PCEP speaker is identified by its IP address. If there is a PCEP speaker in the network that uses multiple IP addresses then it looks like multiple distinct peers to the other PCEP speakers in the network.

Since PCEP sessions can be ephemeral, the peer list tracks a peer even when no PCEP session currently exists to that peer. The statistics contained are an aggregate of the statistics for all successive sessions to that peer.

To limit the quantity of information that is stored, an implementation MAY choose to discard this information if and only if no PCEP session exists to the corresponding peer.

The data model for PCEP peer presented in this document uses a flat list of peers. Each peer in the list is identified by its IP address (addr-type, addr).

There is one list for static peer configuration ("/pcep/entity/peers"), and a separate list for the operational state of all peers (i.e. static as well as discovered)("/pcep-state/entity/peers"). The former is used to enable remote PCE configuration at PCC (or PCE) while the latter has the operational state of these peers as well as the remote PCE peer which were discovered and PCC peers that have initiated session.

5.3. The Session Lists

The session list contains PCEP session that the PCEP entity (PCE or PCC) is currently participating in. The statistics in session are semantically different from those in peer since the former applies to the current session only, whereas the latter is the aggregate for all sessions that have existed to that peer.

Although [RFC5440] forbids more than one active PCEP session between a given pair of PCEP entities at any given time, there is a window during session establishment where two sessions may exist for a given pair, one representing a session initiated by the local PCEP entity and the other representing a session initiated by the peer. If either of these sessions reaches active state first, then the other is discarded.

The data model for PCEP session presented in this document uses a flat list of sessions. Each session in the list is identified by its

initiator. This index allows two sessions to exist transiently for a given peer, as discussed above.

There is only one list for the operational state of all sessions ("/pcep-state/entity/peers/peer/sessions/session").

5.4. Notifications

This YANG model defines a list of notifications to inform client of important events detected during the protocol operation. The notifications defined cover the PCEP MIB notifications.

6. Advanced PCE Features

This document contains a specification of the base PCEP YANG module, "ietf-pcep" which provides the basic PCEP [RFC5440] data model.

This document further handles advanced PCE features like -

- o Capability and Scope
- o Domain information (local/neighbour)
- o Path-Key
- o OF
- o GCO
- o P2MP
- o GMPLS
- o Inter-Layer
- o Stateful PCE
- o Segment Routing
- o Authentication including PCEPS (TLS)

[Editor's Note - Some of them would be added in a future revision.]

6.1. Stateful PCE's LSP-DB

In the operational state of PCEP which supports stateful PCE mode, the list of LSP state are maintained in LSP-DB. The key is the PLSP-ID and the PCC IP address.

The PCEP data model contains the operational state of LSPs (/pcep-state/entity/lsp-db/lsp/) with PCEP specific attributes. The generic TE attributes of the LSP are defined in [I-D.ietf-teas-yang-te]. A reference to LSP state in TE model is maintained.

7. Open Issues and Next Step

This section is added so that open issues can be tracked. This section would be removed when the document is ready for publication.

7.1. The PCE-Initiated LSP

The TE Model at [I-D.ietf-teas-yang-te] should support creationg of tunnels at the controller (PCE) and marking them as PCE-Initiated. The LSP-DB in the PCEP Yang (/pcep-state/entity/lsp-db/lsp/initiation) also marks the LSPs which are PCE-initiated.

7.2. PCEP over TLS (PCEPS)

A future version of this document would add TLS related configurations.

8. PCEP YANG Module

RFC Ed.: In this section, replace all occurrences of 'XXXX' with the actual RFC number and all occurrences of the revision date below with the date of RFC publication (and remove this note).

```
<CODE BEGINS> file "ietf-pcep@2016-07-07.yang"
module ietf-pcep {
  namespace "urn:ietf:params:xml:ns:yang:ietf-pcep";
  prefix pcep;

  import ietf-inet-types {
    prefix "inet";
  }

  import ietf-yang-types {
    prefix "yang";
  }

  import ietf-te {
    prefix "te";
  }

  import ietf-key-chain {
    prefix "key-chain";
  }
}
```

```
}

organization
  "IETF PCE (Path Computation Element) Working Group";

contact
  "WG Web:    <http://tools.ietf.org/wg/pce/>
  WG List:    <mailto:pce@ietf.org>
  WG Chair:   JP Vasseur
              <mailto:jpv@cisco.com>
  WG Chair:   Julien Meuric
              <mailto:julien.meuric@orange.com>
  WG Chair:   Jonathan Hardwick
              <mailto:Jonathan.Hardwick@metaswitch.com>
  Editor:     Dhruv Dhody
              <mailto:dhruv.ietf@gmail.com>";

description
  "The YANG module defines a generic configuration and
  operational model for PCEP common across all of the
  vendor implementations.";

revision 2016-07-07 {
  description "Initial revision.";
  reference
    "RFC XXXX:  A YANG Data Model for Path Computation
    Element Communications Protocol
    (PCEP)";
}

/*
 * Identities
 */

identity pcep {
  description "Identity for the PCEP protocol.";
}

/*
 * Typedefs
 */
typedef pcep-role {
  type enumeration {
    enum unknown {
      value "0";
      description
```



```
        "An unknown role";
    }
    enum pcc {
        value "1";
        description
            "The role of a Path Computation Client";
    }
    enum pce {
        value "2";
        description
            "The role of Path Computation Element";
    }
    enum pcc-and-pce {
        value "3";
        description
            "The role of both Path Computation Client and
            Path Computation Element";
    }
}

description
    "The role of a PCEP speaker.
    Takes one of the following values
    - unknown(0): the role is not known.
    - pcc(1): the role is of a Path Computation
      Client (PCC).
    - pce(2): the role is of a Path Computation
      Server (PCE).
    - pccAndPce(3): the role is of both a PCC and
      a PCE.";
}

typedef pcep-admin-status {
    type enumeration {
        enum admin-status-up {
            value "1";
            description
                "Admin Status is Up";
        }
        enum admin-status-down {
            value "2";
            description
                "Admin Status is Down";
        }
    }
}

description
```

```
    "The Admin Status of the PCEP entity.
    Takes one of the following values
      - admin-status-up(1): Admin Status is Up.
      - admin-status-down(2): Admin Status is Down";
  }

typedef pcep-oper-status {
  type enumeration {
    enum oper-status-up {
      value "1";
      description
        "The PCEP entity is active";
    }
    enum oper-status-down {
      value "2";
      description
        "The PCEP entity is inactive";
    }
    enum oper-status-going-up {
      value "3";
      description
        "The PCEP entity is activating";
    }
    enum oper-status-going-down {
      value "4";
      description
        "The PCEP entity is deactivating";
    }
    enum oper-status-failed {
      value "5";
      description
        "The PCEP entity has failed and will recover
        when possible.";
    }
    enum oper-status-failed-perm {
      value "6";
      description
        "The PCEP entity has failed and will not recover
        without operator intervention";
    }
  }
  description
    "The operational status of the PCEP entity.
    Takes one of the following values
      - oper-status-up(1): Active
      - oper-status-down(2): Inactive
      - oper-status-going-up(3): Activating
      - oper-status-going-down(4): Deactivating
```

```
        - oper-status-failed(5): Failed
        - oper-status-failed-perm(6): Failed Permanantly";
    }

typedef pcep-initiator {
    type enumeration {
        enum local {
            value "1";
            description
                "The local PCEP entity initiated the session";
        }

        enum remote {
            value "2";
            description
                "The remote PCEP peer initiated the session";
        }
    }
    description
        "The initiator of the session, that is, whether the TCP
        connection was initiated by the local PCEP entity or
        the remote peer.
        Takes one of the following values
        - local(1): Initiated locally
        - remote(2): Initiated remotely";
}

typedef pcep-sess-state {
    type enumeration {
        enum tcp-pending {
            value "1";
            description
                "The tcp-pending state of PCEP session.";
        }

        enum open-wait {
            value "2";
            description
                "The open-wait state of PCEP session.";
        }

        enum keep-wait {
            value "3";
            description
                "The keep-wait state of PCEP session.";
        }

        enum session-up {
```

```
        value "4";
        description
            "The session-up state of PCEP session.";
    }
}
description
    "The current state of the session.
    The set of possible states excludes the idle state
    since entries do not exist in the idle state.
    Takes one of the following values
    - tcp-pending(1): PCEP TCP Pending state
    - open-wait(2): PCEP Open Wait state
    - keep-wait(3): PCEP Keep Wait state
    - session-up(4): PCEP Session Up state";
}

typedef domain-type {
    type enumeration {
        enum ospf-area {
            value "1";
            description
                "The OSPF area.";
        }
        enum isis-area {
            value "2";
            description
                "The IS-IS area.";
        }
        enum as {
            value "3";
            description
                "The Autonomous System (AS).";
        }
    }
    description
        "The PCE Domain Type";
}

typedef domain-ospf-area {
    type union {
        type uint32;
        type yang:dotted-quad;
    }
    description
        "OSPF Area ID.";
}

typedef domain-isis-area {
```

```
    type string {
      pattern '[0-9A-Fa-f]{2}\.([0-9A-Fa-f]{4}\.){0,3}';
    }
    description
      "IS-IS Area ID.";
  }

  typedef domain-as {
    type uint32;
    description
      "Autonomous System number.";
  }

  typedef domain {
    type union {
      type domain-ospf-area;
      type domain-isis-area;
      type domain-as;
    }
    description
      "The Domain Information";
  }

  typedef operational-state {
    type enumeration {
      enum down {
        value "0";
        description
          "not active.";
      }
      enum up {
        value "1";
        description
          "signalled.";
      }
      enum active {
        value "2";
        description
          "up and carrying traffic.";
      }
      enum going-down {
        value "3";
        description
          "LSP is being torn down, resources are
            being released.";
      }
      enum going-up {
```

```
        value "4";
        description
            "LSP is being signalled.";
    }
}
description
    "The operational status of the LSP";
}

typedef lsp-error {
    type enumeration {
        enum no-error {
            value "0";
            description
                "No error, LSP is fine.";
        }
        enum unknown {
            value "1";
            description
                "Unknown reason.";
        }
        enum limit {
            value "2";
            description
                "Limit reached for PCE-controlled LSPs.";
        }
        enum pending {
            value "3";
            description
                "Too many pending LSP update requests.";
        }
        enum unacceptable {
            value "4";
            description
                "Unacceptable parameters.";
        }
        enum internal {
            value "5";
            description
                "Internal error.";
        }
        enum admin {
            value "6";
            description
                "LSP administratively brought down.";
        }
        enum preempted {
            value "7";
        }
    }
}
```

```

        description
            "LSP preempted.";
    }
    enum rsvp {
        value "8";
        description
            "RSVP signaling error.";
    }
}
description
    "The LSP Error Codes.";
}

typedef sync-state {
    type enumeration {
        enum pending {
            value "0";
            description
                "The state synchronization
                 has not started.";
        }
        enum ongoing {
            value "1";
            description
                "The state synchronization
                 is ongoing.";
        }
        enum finished {
            value "2";
            description
                "The state synchronization
                 is finished.";
        }
    }
}
description
    "The LSP-DB state synchronization operational status.";
}

typedef pst{
    type enumeration{
        enum rsvp-te{
            value "0";
            description
                "RSVP-TE signaling protocol";
        }
        enum sr{
            value "1";
            description

```

```

    }
    }
    description
        "The Path Setup Type";
}

typedef assoc-type{
    type enumeration{
        enum protection{
            value "1";
            description
                "Path Protection Association Type";
        }
    }
    description
        "The PCEP Association Type";
}

/*
 * Features
 */

feature svec {
    description
        "Support synchronized path computation.";
}

feature gmpls {
    description
        "Support GMPLS.";
}

feature obj-fn {
    description
        "Support OF as per RFC 5541.";
}

feature gco {
    description
        "Support GCO as per RFC 5557.";
}

feature pathkey {
    description
        "Support pathkey as per RFC 5520.";
}

```



```
feature p2mp {
  description
    "Support P2MP as per RFC 6006.";
}

feature stateful {
  description
    "Support stateful PCE.";
}

feature pce-initiated {
  description
    "Support PCE-Initiated LSP.";
}

feature tls {
  description
    "Support PCEP over TLS.";
}

feature sr {
  description
    "Support Segement Routing for PCE.";
}

/*
 * Groupings
 */

grouping pcep-entity-info{
  description
    "This grouping defines the attributes for PCEP entity.";
  leaf connect-timer {
    type uint32 {
      range "1..65535";
    }
    units "seconds";
    default 60;
    description
      "The time in seconds that the PCEP entity will wait
      to establish a TCP connection with a peer.  If a
      TCP connection is not established within this time
      then PCEP aborts the session setup attempt.";
  }
  reference
    "RFC 5440: Path Computation Element (PCE)
    Communication Protocol (PCEP)";
}
```

```
    }

    leaf connect-max-retry {
      type uint32;
      default 5;
      description
        "The maximum number of times the system tries to
        establish a TCP connection to a peer before the
        session with the peer transitions to the idle
        state.";
      reference
        "RFC 5440: Path Computation Element (PCE)
        Communication Protocol (PCEP)";
    }

    leaf init-backoff-timer {
      type uint32 {
        range "1..65535";
      }
      units "seconds";
      description
        "The initial back-off time in seconds for retrying
        a failed session setup attempt to a peer.
        The back-off time increases for each failed
        session setup attempt, until a maximum back-off
        time is reached. The maximum back-off time is
        max-backoff-timer.";
    }

    leaf max-backoff-timer {
      type uint32;
      units "seconds";
      description
        "The maximum back-off time in seconds for retrying
        a failed session setup attempt to a peer.
        The back-off time increases for each failed session
        setup attempt, until this maximum value is reached.
        Session setup attempts then repeat periodically
        without any further increase in back-off time.";
    }

    leaf open-wait-timer {
      type uint32 {
        range "1..65535";
      }
      units "seconds";
      default 60;
      description
```

```
        "The time in seconds that the PCEP entity will wait
        to receive an Open message from a peer after the
        TCP connection has come up.
        If no Open message is received within this time then
        PCEP terminates the TCP connection and deletes the
        associated sessions.";
    reference
        "RFC 5440: Path Computation Element (PCE)
        Communication Protocol (PCEP)";
}

leaf keep-wait-timer {
    type uint32 {
        range "1..65535";
    }
    units "seconds";
    default 60;
    description
        "The time in seconds that the PCEP entity will wait
        to receive a Keepalive or PCErr message from a peer
        during session initialization after receiving an
        Open message.  If no Keepalive or PCErr message is
        received within this time then PCEP terminates the
        TCP connection and deletes the associated
        sessions.";
    reference
        "RFC 5440: Path Computation Element (PCE)
        Communication Protocol (PCEP)";
}

leaf keep-alive-timer {
    type uint32 {
        range "0..255";
    }
    units "seconds";
    default 30;
    description
        "The keep alive transmission timer that this PCEP
        entity will propose in the initial OPEN message of
        each session it is involved in.  This is the
        maximum time between two consecutive messages sent
        to a peer.  Zero means that the PCEP entity prefers
        not to send Keepalives at all.
        Note that the actual Keepalive transmission
        intervals, in either direction of an active PCEP
        session, are determined by negotiation between the
        peers as specified by RFC 5440, and so may differ
        from this configured value.";
```

```
        reference
            "RFC 5440: Path Computation Element (PCE)
              Communication Protocol (PCEP)";
    }

    leaf dead-timer {
        type uint32 {
            range "0..255";
        }
        units "seconds";
        must ". >= ../keep-alive-timer" {
            error-message "The dead timer must be "
                + "larger than the keep alive timer";
            description
                "This value MUST be greater than
                 keep-alive-timer.";
        }
        default 120;
        description
            "The dead timer that this PCEP entity will propose
             in the initial OPEN message of each session it is
             involved in. This is the time after which a peer
             should declare a session down if it does not
             receive any PCEP messages. Zero suggests that the
             peer does not run a dead timer at all." ;
        reference
            "RFC 5440: Path Computation Element (PCE)
              Communication Protocol (PCEP)";
    }

    leaf allow-negotiation{
        type boolean;
        description
            "Whether the PCEP entity will permit negotiation of
             session parameters.";
    }

    leaf max-keep-alive-timer{
        type uint32 {
            range "0..255";
        }
        units "seconds";
        description
            "In PCEP session parameter negotiation in seconds,
             the maximum value that this PCEP entity will
             accept from a peer for the interval between
             Keepalive transmissions. Zero means that the PCEP
```

```
        entity will allow no Keepalive transmission at
        all." ;
    }

    leaf max-dead-timer{
        type uint32 {
            range "0..255";
        }
        units "seconds";
        description
            "In PCEP session parameter negotiation in seconds,
            the maximum value that this PCEP entity will accept
            from a peer for the Dead timer.  Zero means that
            the PCEP entity will allow not running a Dead
            timer.";
    }

    leaf min-keep-alive-timer{
        type uint32 {
            range "0..255";
        }
        units "seconds";
        description
            "In PCEP session parameter negotiation in seconds,
            the minimum value that this PCEP entity will
            accept for the interval between Keepalive
            transmissions. Zero means that the PCEP entity
            insists on no Keepalive transmission at all.";
    }

    leaf min-dead-timer{
        type uint32 {
            range "0..255";
        }
        units "seconds";
        description
            "In PCEP session parameter negotiation in
            seconds, the minimum value that this PCEP entity
            will accept for the Dead timer.  Zero means that
            the PCEP entity insists on not running a Dead
            timer.";
    }

    leaf sync-timer{
        if-feature svec;
        type uint32 {
            range "0..65535";
        }
    }
}
```

```
    units "seconds";
    default 60;
    description
        "The value of SyncTimer in seconds is used in the
        case of synchronized path computation request
        using the SVEC object. Consider the case where a
        PCReq message is received by a PCE that contains
        the SVEC object referring to M synchronized path
        computation requests. If after the expiration of
        the SyncTimer all the M path computation requests
        have not been, received a protocol error is
        triggered and the PCE MUST cancel the whole set
        of path computation requests.
        The aim of the SyncTimer is to avoid the storage
        of unused synchronized requests should one of
        them get lost for some reasons (for example, a
        misbehaving PCC).
        Zero means that the PCEP entity does not use the
        SyncTimer.";
    reference
        "RFC 5440: Path Computation Element (PCE)
        Communication Protocol (PCEP)";
}

leaf request-timer{
    type uint32 {
        range "1..65535";
    }
    units "seconds";
    description
        "The maximum time that the PCEP entity will wait
        for a response to a PCReq message.";
}

leaf max-sessions{
    type uint32;
    description
        "Maximum number of sessions involving this PCEP
        entity that can exist at any time.";
}

leaf max-unknown-reqs{
    type uint32;
    default 5;
    description
        "The maximum number of unrecognized requests and
        replies that any session on this PCEP entity is
```

```
    willing to accept per minute before terminating
    the session.
    A PCRep message contains an unrecognized reply
    if it contains an RP object whose request ID
    does not correspond to any in-progress request
    sent by this PCEP entity.
    A PCReq message contains an unrecognized request
    if it contains an RP object whose request ID is
    zero.";
  reference
    "RFC 5440: Path Computation Element (PCE)
    Communication Protocol (PCEP)";
}

leaf max-unknown-msgs{
  type uint32;
  default 5;
  description
    "The maximum number of unknown messages that any
    session on this PCEP entity is willing to accept
    per minute before terminating the session.";
  reference
    "RFC 5440: Path Computation Element (PCE)
    Communication Protocol (PCEP)";
}

} // pcep-entity-info

grouping pce-scope{
  description
    "This grouping defines PCE path computation scope
    information which maybe relevant to PCE selection.
    This information corresponds to PCE auto-discovery
    information.";
  reference
    "RFC 5088: OSPF Protocol Extensions for Path
    Computation Element (PCE)
    Discovery
    RFC 5089: IS-IS Protocol Extensions for Path
    Computation Element (PCE)
    Discovery";
  leaf intra-area-scope{
    type boolean;
    default true;
    description
      "PCE can compute intra-area paths.";
  }
  leaf intra-area-pref{
```

```
    type uint8{
      range "0..7";
    }
    description
      "The PCE's preference for intra-area TE LSP
      computation.";
  }
  leaf inter-area-scope{
    type boolean;
    default false;
    description
      "PCE can compute inter-area paths.";
  }
  leaf inter-area-scope-default{
    type boolean;
    default false;
    description
      "PCE can act as a default PCE for inter-area
      path computation.";
  }
  leaf inter-area-pref{
    type uint8{
      range "0..7";
    }
    description
      "The PCE's preference for inter-area TE LSP
      computation.";
  }
  leaf inter-as-scope{
    type boolean;
    default false;
    description
      "PCE can compute inter-AS paths.";
  }
  leaf inter-as-scope-default{
    type boolean;
    default false;
    description
      "PCE can act as a default PCE for inter-AS
      path computation.";
  }
  leaf inter-as-pref{
    type uint8{
      range "0..7";
    }
    description
      "The PCE's preference for inter-AS TE LSP
      computation.";
```



```
    }
    leaf inter-layer-scope{
        type boolean;
        default false;
        description
            "PCE can compute inter-layer paths.";
    }
    leaf inter-layer-pref{
        type uint8{
            range "0..7";
        }
        description
            "The PCE's preference for inter-layer TE LSP
            computation.";
    }
} //pce-scope

grouping domain{
    description
        "This grouping specifies a Domain where the
        PCEP speaker has topology visibility.";
    leaf domain-type{
        type domain-type;
        description
            "The domain type.";
    }
    leaf domain{
        type domain;
        description
            "The domain Information.";
    }
} //domain

grouping capability{
    description
        "This grouping specifies a capability
        information of local PCEP entity. This maybe
        relevant to PCE selection as well. This
        information corresponds to PCE auto-discovery
        information.";
    reference
        "RFC 5088: OSPF Protocol Extensions for Path
        Computation Element (PCE)
        Discovery
        RFC 5089: IS-IS Protocol Extensions for Path
        Computation Element (PCE)
        Discovery";
    leaf gmpls{
```

```
        if-feature gmpls;
        type boolean;
        description
            "Path computation with GMPLS link
            constraints.";
    }
    leaf bi-dir{
        type boolean;
        description
            "Bidirectional path computation.";
    }
    leaf diverse{
        type boolean;
        description
            "Diverse path computation.";
    }
    leaf load-balance{
        type boolean;
        description
            "Load-balanced path computation.";
    }
    leaf synchronize{
        if-feature svec;
        type boolean;
        description
            "Synchronized paths computation.";
    }
    leaf objective-function{
        if-feature obj-fn;
        type boolean;
        description
            "Support for multiple objective functions.";
    }
    leaf add-path-constraint{
        type boolean;
        description
            "Support for additive path constraints (max
            hop count, etc.).";
    }
    leaf prioritization{
        type boolean;
        description
            "Support for request prioritization.";
    }
    leaf multi-request{
        type boolean;
        description
            "Support for multiple requests per message.";
```

```

    }
    leaf gco{
        if-feature gco;
        type boolean;
        description
            "Support for Global Concurrent Optimization
            (GCO).";
    }
    leaf p2mp{
        if-feature p2mp;
        type boolean;
        description
            "Support for P2MP path computation.";
    }
}

container stateful{
    if-feature stateful;
    description
        "If stateful PCE feature is present";
    leaf enabled{
        type boolean;
        description
            "Enabled or Disabled";
    }
    leaf active{
        type boolean;
        description
            "Support for active stateful PCE.";
    }
    leaf pce-initiated{
        if-feature pce-initiated;
        type boolean;
        description
            "Support for PCE-initiated LSP.";
    }
}

container sr{
    if-feature sr;
    description
        "If segment routing is supported";
    leaf enabled{
        type boolean;
        description
            "Enabled or Disabled";
    }
    leaf msd{ /*should be in MPLS yang model (?)*/
        type uint8;
        must "((../../role == 'pcc')" +

```

```

        " or " +
        "(../../role == 'pcc-and-pce'))))"
    {
        error-message
            "The PCEP entity must be PCC";
        description
            "When PCEP entity is PCC for
            MSD to be applicable";
    }
        description
            "Maximum SID Depth";
    }
}
} //capability

grouping info{
    description
        "This grouping specifies all information which
        maybe relevant to both PCC and PCE.
        This information corresponds to PCE auto-discovery
        information.";
    container domain{
        description
            "The local domain for the PCEP entity";
        list domain{
            key "domain-type domain";
            description
                "The local domain.";
            uses domain{
                description
                    "The local domain for the PCEP entity.";
            }
        }
    }
    container capability{
        description
            "The PCEP entity capability";
        uses capability{
            description
                "The PCEP entity supported
                capabilities.";
        }
    }
} //info

grouping pce-info{
    description
        "This grouping specifies all PCE information

```

```
        which maybe relevant to the PCE selection.
        This information corresponds to PCE auto-discovery
        information.";
    container scope{
        description
            "The path computation scope";
        uses pce-scope;
    }

    container neigh-domains{
        description
            "The list of neighbour PCE-Domain
            toward which a PCE can compute
            paths";
        list domain{
            key "domain-type domain";

            description
                "The neighbour domain.";
            uses domain{
                description
                    "The PCE neighbour domain.";
            }
        }
    }
} //pce-info

grouping pcep-stats{
    description
        "This grouping defines statistics for PCEP. It is used
        for both peer and current session.";
    leaf avg-rsp-time{
        type uint32;
        units "milliseconds";
        must "(/pcep-state/entity/peers/peer/role != 'pcc'" +
            " or " +
            "(/pcep-state/entity/peers/peer/role = 'pcc'" +
            " and avg-rsp-time = 0))" {
            error-message
                "Invalid average response time";
            description
                "If role is pcc then this leaf is meaningless
                and is set to zero.";
        }
    }
    description
        "The average response time.
        If an average response time has not been
        calculated then this leaf has the value zero.";
```

```
    }  
  
    leaf lwm-rsp-time{  
        type uint32;  
        units "milliseconds";  
        must "(/pcep-state/entity/peers/peer/role != 'pcc'" +  
            " or " +  
            "(/pcep-state/entity/peers/peer/role = 'pcc'" +  
            " and lwm-rsp-time = 0))" {  
            error-message  
                "Invalid smallest (low-water mark)  
                response time";  
            description  
                "If role is pcc then this leaf is meaningless  
                and is set to zero.";  
        }  
        description  
            "The smallest (low-water mark) response time seen.  
            If no responses have been received then this  
            leaf has the value zero.";  
    }  
  
    leaf hwm-rsp-time{  
        type uint32;  
        units "milliseconds";  
        must "(/pcep-state/entity/peers/peer/role != 'pcc'" +  
            " or " +  
            "(/pcep-state/entity/peers/peer/role = 'pcc'" +  
            " and hwm-rsp-time = 0))" {  
            error-message  
                "Invalid greatest (high-water mark)  
                response time seen";  
            description  
                "If role is pcc then this field is  
                meaningless and is set to zero.";  
        }  
        description  
            "The greatest (high-water mark) response time seen.  
            If no responses have been received then this object  
            has the value zero.";  
    }  
  
    leaf num-pcreq-sent{  
        type yang:counter32;  
        description  
            "The number of PCReq messages sent.";  
    }  
}
```

```
leaf num-pcreq-rcvd{
  type yang:counter32;
  description
    "The number of PCReq messages received.";
}

leaf num-pcrep-sent{
  type yang:counter32;
  description
    "The number of PCRep messages sent.";
}

leaf num-pcrep-rcvd{
  type yang:counter32;
  description
    "The number of PCRep messages received.";
}

leaf num-pcerr-sent{
  type yang:counter32;
  description
    "The number of PCErr messages sent.";
}

leaf num-pcerr-rcvd{
  type yang:counter32;
  description
    "The number of PCErr messages received.";
}

leaf num-pcntf-sent{
  type yang:counter32;
  description
    "The number of PCNtf messages sent.";
}

leaf num-pcntf-rcvd{
  type yang:counter32;
  description
    "The number of PCNtf messages received.";
}

leaf num-keepalive-sent{
  type yang:counter32;
  description
    "The number of Keepalive messages sent.";
}
```

```
leaf num-keepalive-rcvd{
  type yang:counter32;
  description
    "The number of Keepalive messages received.";
}

leaf num-unknown-rcvd{
  type yang:counter32;
  description
    "The number of unknown messages received.";
}

leaf num-corrupt-rcvd{
  type yang:counter32;
  description
    "The number of corrupted PCEP message received.";
}

leaf num-req-sent{
  type yang:counter32;
  description
    "The number of requests sent. A request corresponds
    1:1 with an RP object in a PCReq message. This might
    be greater than num-pcreq-sent because multiple
    requests can be batched into a single PCReq
    message.";
}

leaf num-req-sent-pend-rep{
  type yang:counter32;
  description
    "The number of requests that have been sent for
    which a response is still pending.";
}

leaf num-req-sent-ero-rcvd{
  type yang:counter32;
  description
    "The number of requests that have been sent for
    which a response with an ERO object was received.
    Such responses indicate that a path was
    successfully computed by the peer.";
}

leaf num-req-sent-nopath-rcvd{
  type yang:counter32;
  description
    "The number of requests that have been sent for
```



```
        which a response with a NO-PATH object was
        received. Such responses indicate that the peer
        could not find a path to satisfy the
        request.";
    }

    leaf num-req-sent-cancel-rcvd{
        type yang:counter32;
        description
            "The number of requests that were cancelled with
            a PCNtf message.
            This might be different than num-pcntf-rcvd because
            not all PCNtf messages are used to cancel requests,
            and a single PCNtf message can cancel multiple
            requests.";
    }

    leaf num-req-sent-error-rcvd{
        type yang:counter32;
        description
            "The number of requests that were rejected with a
            PCErr message.
            This might be different than num-pcerr-rcvd because
            not all PCErr messages are used to reject requests,
            and a single PCErr message can reject multiple
            requests.";
    }

    leaf num-req-sent-timeout{
        type yang:counter32;
        description
            "The number of requests that have been sent to a peer
            and have been abandoned because the peer has taken too
            long to respond to them.";
    }

    leaf num-req-sent-cancel-sent{
        type yang:counter32;
        description
            "The number of requests that were sent to the peer and
            explicitly cancelled by the local PCEP entity sending
            a PCNtf.";
    }

    leaf num-req-rcvd{
        type yang:counter32;
        description
            "The number of requests received. A request
```

```
        corresponds 1:1 with an RP object in a PCReq
        message.
        This might be greater than num-pcreq-rcvd because
        multiple requests can be batched into a single
        PCReq message.";
    }

leaf num-req-rcvd-pend-rep{
    type yang:counter32;
    description
        "The number of requests that have been received for
        which a response is still pending.";
}

leaf num-req-rcvd-ero-sent{
    type yang:counter32;
    description
        "The number of requests that have been received for
        which a response with an ERO object was sent. Such
        responses indicate that a path was successfully
        computed by the local PCEP entity.";
}

leaf num-req-rcvd-nopath-sent{
    type yang:counter32;
    description
        "The number of requests that have been received for
        which a response with a NO-PATH object was sent. Such
        responses indicate that the local PCEP entity could
        not find a path to satisfy the request.";
}

leaf num-req-rcvd-cancel-sent{
    type yang:counter32;
    description
        "The number of requests received that were cancelled
        by the local PCEP entity sending a PCNtf message.
        This might be different than num-pcntf-sent because
        not all PCNtf messages are used to cancel requests,
        and a single PCNtf message can cancel multiple
        requests.";
}

leaf num-req-rcvd-error-sent{
    type yang:counter32;
    description
        "The number of requests received that were cancelled
        by the local PCEP entity sending a PCErr message.
```

```
        This might be different than num-pcerr-sent because
        not all PCErr messages are used to cancel requests,
        and a single PCErr message can cancel multiple
        requests.";
    }

    leaf num-req-rcvd-cancel-rcvd{
        type yang:counter32;
        description
            "The number of requests that were received from the
            peer and explicitly cancelled by the peer sending
            a PCNTf.";
    }

    leaf num-req-rcvd-unknown{
        type yang:counter32;
        description
            "The number of responses to unknown requests
            received. A response to an unknown request is a
            response whose RP object does not contain the
            request ID of any request that is currently
            outstanding on the session.";
    }

    leaf num-req-rcvd-unknown{
        type yang:counter32;
        description
            "The number of unknown requests that have been
            received. An unknown request is a request
            whose RP object contains a request ID of
            zero.";
    }

    container svec{
        if-feature svec;
        description
            "If synchronized path computation is supported";
        leaf num-svec-sent{
            type yang:counter32;
            description
                "The number of SVEC objects sent in PCReq messages.
                An SVEC object represents a set of synchronized
                requests.";
        }

        leaf num-svec-req-sent{
            type yang:counter32;
            description
```

```
        "The number of requests sent that appeared in one
        or more SVEC objects.";
    }

    leaf num-svec-rcvd{
        type yang:counter32;
        description
            "The number of SVEC objects received in PCReq
            messages. An SVEC object represents a set of
            synchronized requests.";
    }

    leaf num-svec-req-rcvd{
        type yang:counter32;
        description
            "The number of requests received that appeared
            in one or more SVEC objects.";
    }
}
container stateful{
    if-feature stateful;
    description
        "Stateful PCE related statistics";
    leaf num-pcrpt-sent{
        type yang:counter32;
        description
            "The number of PCRpt messages sent.";
    }

    leaf num-pcrpt-rcvd{
        type yang:counter32;
        description
            "The number of PCRpt messages received.";
    }

    leaf num-pcupd-sent{
        type yang:counter32;
        description
            "The number of PCUpd messages sent.";
    }

    leaf num-pcupd-rcvd{
        type yang:counter32;
        description
            "The number of PCUpd messages received.";
    }

    leaf num-rpt-sent{
```

```
    type yang:counter32;
    description
      "The number of LSP Reports sent. A LSP report
       corresponds 1:1 with an LSP object in a PCRpt
       message. This might be greater than
       num-pcrpt-sent because multiple reports can
       be batched into a single PCRpt message.";
  }

  leaf num-rpt-rcvd{
    type yang:counter32;
    description
      "The number of LSP Reports received. A LSP report
       corresponds 1:1 with an LSP object in a PCRpt
       message.
       This might be greater than num-pcrpt-rcvd because
       multiple reports can be batched into a single
       PCRpt message.";
  }

  leaf num-rpt-rcvd-error-sent{
    type yang:counter32;
    description
      "The number of reports of LSPs received that were
       responded by the local PCEP entity by sending a
       PCErr message.";
  }

  leaf num-upd-sent{
    type yang:counter32;
    description
      "The number of LSP updates sent. A LSP update
       corresponds 1:1 with an LSP object in a PCUpd
       message. This might be greater than
       num-pcupd-sent because multiple updates can
       be batched into a single PCUpd message.";
  }

  leaf num-upd-rcvd{
    type yang:counter32;
    description
      "The number of LSP Updates received. A LSP update
       corresponds 1:1 with an LSP object in a PCUpd
       message.
       This might be greater than num-pcupd-rcvd because
       multiple updates can be batched into a single
       PCUpd message.";
  }
}
```

```
leaf num-upd-rcvd-unknown{
  type yang:counter32;
  description
    "The number of updates to unknown LSPs
    received. An update to an unknown LSP is a
    update whose LSP object does not contain the
    PLSP-ID of any LSP that is currently
    present.";
}

leaf num-upd-rcvd-undelegated{
  type yang:counter32;
  description
    "The number of updates to not delegated LSPs
    received. An update to an undelegated LSP is a
    update whose LSP object does not contain the
    PLSP-ID of any LSP that is currently
    delegated to current PCEP session.";
}

leaf num-upd-rcvd-error-sent{
  type yang:counter32;
  description
    "The number of updates to LSPs received that were
    responded by the local PCEP entity by sending a
    PCErr message.";
}

container initiation {
  if-feature pce-initiated;
  description
    "PCE-Initiated related statistics";
  leaf num-pcinitiate-sent{
    type yang:counter32;
    description
      "The number of PCInitiate messages sent.";
  }

  leaf num-pcinitiate-rcvd{
    type yang:counter32;
    description
      "The number of PCInitiate messages received.";
  }

  leaf num-initiate-sent{
    type yang:counter32;
    description
      "The number of LSP Initiation sent via PCE.
      A LSP initiation corresponds 1:1 with an LSP
```

```

        object in a PCInitiate message. This might be
        greater than num-pcinitiate-sent because
        multiple initiations can be batched into a
        single PCInitiate message.";
    }

    leaf num-initiate-rcvd{
        type yang:counter32;
        description
            "The number of LSP Initiation received from
            PCE. A LSP initiation corresponds 1:1 with
            an LSP object in a PCInitiate message. This
            might be greater than num-pcinitiate-rcvd
            because multiple initiations can be batched
            into a single PCInitiate message.";
    }

    leaf num-initiate-rcvd-error-sent{
        type yang:counter32;
        description
            "The number of initiations of LSPs received
            that were responded by the local PCEP entity
            by sending a PCErr message.";
    }
}
}
} //pcep-stats

grouping lsp-state{
    description
        "This grouping defines the attributes for LSP in LSP-DB.
        These are the attributes specifically from the PCEP
        perspective";
    leaf plsp-id{
        type uint32{
            range "1..1048575";
        }
        description
            "A PCEP-specific identifier for the LSP. A PCC
            creates a unique PLSP-ID for each LSP that is
            constant for the lifetime of a PCEP session.
            PLSP-ID is 20 bits with 0 and 0xFFFF are
            reserved";
    }
    leaf pcc-id{
        type inet:ip-address;
        description
            "The local internet address of the PCC, that

```

```
        generated the PLSP-ID.";
    }
container lsp-ref{
    description
        "reference to ietf-te lsp state";

    leaf source {
        type leafref {
            path "/te:te/te:lsps-state/te:lsp/te:source";
        }
        description
            "Tunnel sender address extracted from
            SENDER_TEMPLATE object";
        reference "RFC3209";
    }
    leaf destination {
        type leafref {
            path "/te:te/te:lsps-state/te:lsp/te:"
                + "destination";
        }
        description
            "Tunnel endpoint address extracted from
            SESSION object";
        reference "RFC3209";
    }
    leaf tunnel-id {
        type leafref {
            path "/te:te/te:lsps-state/te:lsp/te:tunnel-id";
        }
        description
            "Tunnel identifier used in the SESSION
            that remains constant over the life
            of the tunnel.";
        reference "RFC3209";
    }
    leaf lsp-id {
        type leafref {
            path "/te:te/te:lsps-state/te:lsp/te:lsp-id";
        }
        description
            "Identifier used in the SENDER_TEMPLATE
            and the FILTER_SPEC that can be changed
            to allow a sender to share resources with
            itself.";
        reference "RFC3209";
    }
    leaf extended-tunnel-id {
```



```

        type leafref {
            path "/te:te/te:lsps-state/te:lsp/te:"
                + "extended-tunnel-id";
        }
        description
            "Extended Tunnel ID of the LSP.";
        reference "RFC3209";
    }
    leaf type {
        type leafref {
            path "/te:te/te:lsps-state/te:lsp/te:type";
        }
        description "LSP type P2P or P2MP";
    }
}

leaf admin-state{
    type boolean;
    description
        "The desired operational state";
}
leaf operational-state{
    type operational-state;
    description
        "The operational status of the LSP";
}
container delegated{
    description
        "The delegation related parameters";
    leaf enabled{
        type boolean;
        description
            "LSP is delegated or not";
    }
    leaf pce{
        type leafref {
            path "/pcep-state/entity/peers/peer/addr";
        }
        must "((../enabled == true)" +
            " and " +
            "((../role == 'pcc')" +
            " or " +
            "(../role == 'pcc-and-pce')))"
        {
            error-message
                "The PCEP entity must be PCC
                and the LSP be delegated";
            description

```

```

        "When PCEP entity is PCC for
        delegated LSP";
    }
    description
        "The reference to the PCE peer to
        which LSP is delegated";
    }
    leaf srp-id{
        type uint32;
        description
            "The last SRP-ID-number associated with this
            LSP.";
    }
}
container initiation {
    if-feature pce-initiated;
    description
        "The PCE initiation related parameters";
    leaf enabled{
        type boolean;
        description
            "LSP is PCE-initiated or not";
    }
    leaf pce{
        type leafref {
            path "/pcep-state/entity/peers/peer/addr";
        }
        must "(../enabled == true)"
        {
            error-message
                "The LSP must be PCE-Initiated";
            description
                "When the LSP must be PCE-Initiated";
        }
        description
            "The reference to the PCE
            that initiated this LSP";
    }
}
leaf symbolic-path-name{
    type string;
    description
        "The symbolic path name associated with the LSP.";
}
leaf last-error{
    type lsp-error;
    description
        "The last error for the LSP.";
}

```

```
    }
        leaf pst{
            type pst;
            default "rsvp-te";
            description
                "The Path Setup Type";
        }
} //lsp-state

grouping notification-instance-hdr {
    description
        "This group describes common instance specific data
        for notifications.";

    leaf peer-addr {
        type leafref {
            path "/pcep-state/entity/peers/peer/addr";
        }
        description
            "Reference to peer address";
    }
} // notification-instance-hdr

grouping notification-session-hdr {
    description
        "This group describes common session instance specific
        data for notifications.";

    leaf session-initiator {
        type leafref {
            path "/pcep-state/entity/peers/peer/sessions/" +
                "session/initiator";
        }
        description
            "Reference to pcep session initiator leaf";
    }
} // notification-session-hdr

grouping stateful-pce-parameter {
    description
        "This group describes stateful PCE specific
        parameters.";
    leaf state-timeout{
        type uint32;
        units "seconds";
    }
}
```

```

        description
            "When a PCEP session is terminated, a PCC
            waits for this time period before flushing
            LSP state associated with that PCEP session
            and reverting to operator-defined default
            parameters or behaviours.";
    }
    leaf redelegation-timeout{
        type uint32;
        units "seconds";
        must "((../role == 'pcc')" +
            " or " +
            "(../role == 'pcc-and-pce'))"
        {
            error-message "The PCEP entity must be PCC";
            description
                "When PCEP entity is PCC";
        }
        description
            "When a PCEP session is terminated, a PCC
            waits for this time period before revoking
            LSP delegation to a PCE and attempting to
            redelegate LSPs associated with the
            terminated PCEP session to an alternate
            PCE.";
    }
    leaf rpt-non-pcep-lsp{
        type boolean;
        must "((../role == 'pcc')" +
            " or " +
            "(../role == 'pcc-and-pce'))"
        {
            error-message "The PCEP entity must be PCC";
            description
                "When PCEP entity is PCC";
        }
        description
            "If set, a PCC reports LSPs that are not
            controlled by any PCE (for example, LSPs
            that are statically configured at the
            PCC). ";
    }
}

grouping authentication {
    description "Authentication Information";
    choice auth-type-selection {

```

```
description
  "Options for expressing authentication setting.>";
case auth-key-chain {
  leaf key-chain {
    type key-chain:key-chain-ref;
    description
      "key-chain name.>";
  }
}
case auth-key {
  leaf key {
    type string;
    description
      "Key string in ASCII format.>";
  }
  container crypto-algorithm {
    uses key-chain:crypto-algorithm-types;
    description
      "Cryptographic algorithm associated
      with key.>";
  }
}
case auth-tls {
  if-feature tls;
  container tls {
    description
      "TLS related information - TBD";
  }
}
}
}

grouping association {
  description
    "Generic Association parameters";
  leaf type {
    type "assoc-type";
    description
      "The PCEP association type";
  }
  leaf id {
    type uint16;
    description
      "PCEP Association ID";
  }
  leaf source {
    type inet:ip-address;
    description

```

```
        "PCEP Association Source.";
    }
    leaf global-source {
        type uint32;
        description
            "PCEP Association Global
             Source.";
    }
    leaf extended-id{
        type string;
        description
            "Additional information to
             support unique identification.";
    }
}
grouping association-ref {
    description
        "Generic Association parameters";
    leaf id {
        type leafref {
            path "/pcep-state/entity/lsp-db/"
                + "association-list/id";
        }
        description
            "PCEP Association ID";
    }
    leaf source {
        type leafref {
            path "/pcep-state/entity/lsp-db/"
                + "association-list/source";
        }
        description
            "PCEP Association Source.";
    }
    leaf global-source {
        type leafref {
            path "/pcep-state/entity/lsp-db/"
                + "association-list/global-source";
        }
        description
            "PCEP Association Global
             Source.";
    }
    leaf extended-id{
        type leafref {
            path "/pcep-state/entity/lsp-db/"
                + "association-list/extended-id";
        }
    }
}
```

```
        description
            "Additional information to
            support unique identification.";
    }
}
/*
 * Configuration data nodes
 */
container pcep{

    presence
        "The PCEP is enabled";

    description
        "Parameters for list of configured PCEP entities
        on the device.";

    container entity {

        description
            "The configured PCEP entity on the device.";

        leaf addr {
            type inet:ip-address;
            mandatory true;
            description
                "The local Internet address of this PCEP
                entity.
                If operating as a PCE server, the PCEP
                entity listens on this address.
                If operating as a PCC, the PCEP entity
                binds outgoing TCP connections to this
                address.
                It is possible for the PCEP entity to
                operate both as a PCC and a PCE Server, in
                which case it uses this address both to
                listen for incoming TCP connections and to
                bind outgoing TCP connections.";
        }

        leaf enabled {
            type boolean;
            default true;
            description
                "The administrative status of this PCEP
                Entity.";
        }
    }
}
```

```

leaf role {
  type pcep-role;
  mandatory true;
  description
    "The role that this entity can play.
    Takes one of the following values.
    - unknown(0): this PCEP Entity role is not
      known.
    - pcc(1): this PCEP Entity is a PCC.
    - pce(2): this PCEP Entity is a PCE.
    - pcc-and-pce(3): this PCEP Entity is both
      a PCC and a PCE.";
}

leaf description {
  type string;
  description
    "Description of the PCEP entity configured
    by the user";
}

uses info {
  description
    "Local PCEP entity information";
}

container pce-info {
  must "((../role == 'pce')" +
    " or " +
    "(../role == 'pcc-and-pce'))"
  {
    error-message "The PCEP entity must be PCE";
    description
      "When PCEP entity is PCE";
  }
  uses pce-info {
    description
      "Local PCE information";
  }
  uses authentication {
    description
      "Local PCE authentication inform
ation";
  }
}

description
  "The Local PCE Entity PCE information";

```



```
    }

    uses pcep-entity-info {
        description
            "The configuration related to the PCEP
            entity.";
    }

    leaf pcep-notification-max-rate {
        type uint32;
        mandatory true;
        description
            "This variable indicates the maximum number of
            notifications issued per second. If events occur
            more rapidly, the implementation may simply fail
            to emit these notifications during that period,
            or may queue them until an appropriate time. A
            value of 0 means no notifications are emitted
            and all should be discarded (that is, not
            queued).";
    }

    container stateful-parameter{
        if-feature stateful;
        must "(../info/capability/stateful/active == true)"
        {
            error-message
                "The Active Stateful PCE must be enabled";
            description
                "When PCEP entity is active stateful
                enabled";
        }
        uses stateful-pce-parameter;

        description
            "The configured stateful parameters";
    }

    container peers{
        must "(../role == 'pcc') " +
            " or " +
            "(../role == 'pcc-and-pce') "
        {
            error-message
                "The PCEP entity must be PCC";
        }
    }
}
```

```
        description
            "When PCEP entity is PCC, as remote
            PCE peers are configured.";
    }
description
    "The list of configured peers for the
    entity (remote PCE)";
list peer{
    key "addr";

    description
        "The peer configured for the entity.
        (remote PCE)";

    leaf addr {
        type inet:ip-address;
        description
            "The local Internet address of this
            PCEP peer.";
    }

    leaf description {
        type string;
        description
            "Description of the PCEP peer
            configured by the user";
    }
    uses info {
        description
            "PCE Peer information";
    }
    uses pce-info {
        description
            "PCE Peer information";
    }
}

leaf delegation-pref{
    if-feature stateful;
    type uint8{
        range "0..7";
    }
    must "(../../info/capability/stateful/active"
        + " == true)"
    {
        error-message
            "The Active Stateful PCE must be
            enabled";
        description

```

```

        "When PCEP entity is active stateful
        enabled";
    }
    description
        "The PCE peer delegation preference.";
    }
    uses authentication {
        description
            "PCE Peer authentication";
    }
    } //peer
} //peers
} //entity
} //pcep

/*
 * Operational data nodes
 */

container pcep-state{
    config false;
    description
        "The list of operational PCEP entities on the
        device.";

    container entity{
        description
            "The operational PCEP entity on the device.";

        leaf addr {
            type inet:ip-address;
            description
                "The local Internet address of this PCEP
                entity.
                If operating as a PCE server, the PCEP
                entity listens on this address.
                If operating as a PCC, the PCEP entity
                binds outgoing TCP connections to this
                address.
                It is possible for the PCEP entity to
                operate both as a PCC and a PCE Server, in
                which case it uses this address both to
                listen for incoming TCP connections and to
                bind outgoing TCP connections.";
        }

        leaf index{
            type uint32;

```

```
        description
            "The index of the operational PECP
            entity";
    }

    leaf admin-status {
        type pcep-admin-status;
        description
            "The administrative status of this PCEP Entity.
            This is the desired operational status as
            currently set by an operator or by default in
            the implementation. The value of enabled
            represents the current status of an attempt
            to reach this desired status.";
    }

    leaf oper-status {
        type pcep-admin-status;
        description
            "The operational status of the PCEP entity.
            Takes one of the following values.
            - oper-status-up(1): the PCEP entity is
              active.
            - oper-status-down(2): the PCEP entity is
              inactive.
            - oper-status-going-up(3): the PCEP entity is
              activating.
            - oper-status-going-down(4): the PCEP entity is
              deactivating.
            - oper-status-failed(5): the PCEP entity has
              failed and will recover when possible.
            - oper-status-failed-perm(6): the PCEP entity
              has failed and will not recover without
              operator intervention.";
    }

    leaf role {
        type pcep-role;
        description
            "The role that this entity can play.
            Takes one of the following values.
            - unknown(0): this PCEP entity role is
              not known.
            - pcc(1): this PCEP entity is a PCC.
            - pce(2): this PCEP entity is a PCE.
            - pcc-and-pce(3): this PCEP entity is
              both a PCC and a PCE.";
```

```

    }
    uses info {
        description
            "Local PCEP entity information";
    }

    container pce-info {
        when "(../role == 'pce') " +
            " or " +
            "(../role == 'pcc-and-pce') "
        {
            description
                "When PCEP entity is PCE";
        }
        uses pce-info {
            description
                "Local PCE information";
        }
        uses authentication {
            description
                "Local PCE authentication inform
ation";
        }
        description
            "The Local PCE Entity PCE information";
    }

    uses pcep-entity-info{
        description
            "The operational information related to the
            PCEP entity.";
    }

    container stateful-parameter{
        if-feature stateful;
        must "(../info/capability/stateful/active == true)"
        {
            error-message
                "The Active Stateful PCE must be enabled";
            description
                "When PCEP entity is active stateful
                enabled";
        }
        uses stateful-pce-parameter;

        description
            "The operational stateful parameters";
    }

```

```

container lsp-db{
  if-feature stateful;
  description
    "The LSP-DB";
  list association-list {
    key "id source global-source extended-id";
    description
      "List of all PCEP associations";
    uses association {
      description
        "The Association attributes";
    }
    list lsp {
      key "plsp-id pcc-id";
      description
        "List of all LSP in this association";
      leaf plsp-id {
        type leafref {
          path "/pcep-state/entity/lsp-db/"
            + "lsp/plsp-id";
        }
        description
          "Reference to PLSP-ID in LSP-DB";
      }
      leaf pcc-id {
        type leafref {
          path "/pcep-state/entity/lsp-db/"
            + "lsp/pcc-id";
        }
        description
          "Reference to PCC-ID in LSP-DB";
      }
    }
  }
}
list lsp{
  key "plsp-id pcc-id";
  description
    "List of all LSPs in LSP-DB";
  uses lsp-state{
    description
      "The PCEP specific attributes for
        LSP-DB.";
  }
  list association-list {
    key "id source global-source extended-id";
    description
      "List of all PCEP associations";
    uses association-ref {

```

```

        description
            "Reference to the Association
            attributes";
    }
}
}
}
container peers{
    description
        "The list of peers for the entity";

    list peer{
        key "addr";

        description
            "The peer for the entity.";

        leaf addr {
            type inet:ip-address;
            description
                "The local Internet address of this PCEP
                peer.";
        }

        leaf role {
            type pcep-role;
            description
                "The role of the PCEP Peer.
                Takes one of the following values.
                - unknown(0): this PCEP peer role
                is not known.
                - pcc(1): this PCEP peer is a PCC.
                - pce(2): this PCEP peer is a PCE.
                - pcc-and-pce(3): this PCEP peer
                is both a PCC and a PCE.";
        }

        uses info {
            description
                "PCEP peer information";
        }

        container pce-info {
            when "(../role == 'pce') " +
            " or " +

```

```
    "(../role == 'pcc-and-pce'))"
  {
    description
      "When PCEP entity is PCE";
  }
  uses pce-info {
    description
      "PCE Peer information";
  }
  description
    "The PCE Peer information";
}

leaf delegation-pref{
  if-feature stateful;
  type uint8{
    range "0..7";
  }
  must "((../role == 'pcc')" +
    " or " +
    "(../role == 'pcc-and-pce'))"
  {
    error-message
      "The PCEP entity must be PCC";
    description
      "When PCEP entity is PCC";
  }
  must "(../info/capability/stateful/active"
    + " == true)"
  {
    error-message
      "The Active Stateful PCE must be
      enabled";
    description
      "When PCEP entity is active stateful
      enabled";
  }
  description
    "The PCE peer delegation preference.";
}

uses authentication {
  description
    "PCE Peer authentication";
}

leaf discontinuity-time {
  type yang:timestamp;
```



```
        description
            "The timestamp of the time when the
            information and statistics were
            last reset.";
    }

    leaf initiate-session {
        type boolean;
        description
            "Indicates whether the local PCEP
            entity initiates sessions to this peer,
            or waits for the peer to initiate a
            session.";
    }

    leaf session-exists{
        type boolean;
        description
            "Indicates whether a session with
            this peer currently exists.";
    }

    leaf num-sess-setup-ok{
        type yang:counter32;
        description
            "The number of PCEP sessions successfully
            successfully established with the peer,
            including any current session. This
            counter is incremented each time a
            session with this peer is successfully
            established.";
    }

    leaf num-sess-setup-fail{
        type yang:counter32;
        description
            "The number of PCEP sessions with the peer
            that have been attempted but failed
            before being fully established. This
            counter is incremented each time a
            session retry to this peer fails.";
    }

    leaf session-up-time{
        type yang:timestamp;
        must "(../num-sess-setup-ok != 0 or " +
            "(../num-sess-setup-ok = 0 and " +
            "session-up-time = 0))" {

```

```
        error-message
            "Invalid Session Up timestamp";
        description
            "If num-sess-setup-ok is zero,
             then this leaf contains zero.";
    }
    description
        "The timestamp value of the last time a
         session with this peer was successfully
         established.";
}

leaf session-fail-time{
    type yang:timestamp;
    must "(../num-sess-setup-fail != 0 or " +
        "(../num-sess-setup-fail = 0 and " +
        "session-fail-time = 0))" {
        error-message
            "Invalid Session Fail timestamp";
        description
            "If num-sess-setup-fail is zero,
             then this leaf contains zero.";
    }
    description
        "The timestamp value of the last time a
         session with this peer failed to be
         established.";
}

leaf session-fail-up-time{
    type yang:timestamp;
    must "(../num-sess-setup-ok != 0 or " +
        "(../num-sess-setup-ok = 0 and " +
        "session-fail-up-time = 0))" {
        error-message
            "Invalid Session Fail from
             Up timestamp";
        description
            "If num-sess-setup-ok is zero,
             then this leaf contains zero.";
    }
    description
        "The timestamp value of the last time a
         session with this peer failed from
         active.";
}

container pcep-stats {
```

```
description
  "The container for all statistics at peer
  level.";
uses pcep-stats{
  description
    "Since PCEP sessions can be
    ephemeral, the peer statistics tracks
    a peer even when no PCEP session
    currently exists to that peer. The
    statistics contained are an aggregate
    of the statistics for all successive
    sessions to that peer.";
}

leaf num-req-sent-closed{
  type yang:counter32;
  description
    "The number of requests that were
    sent to the peer and implicitly
    cancelled when the session they were
    sent over was closed.";
}

leaf num-req-rcvd-closed{
  type yang:counter32;
  description
    "The number of requests that were
    received from the peer and
    implicitly cancelled when the
    session they were received over
    was closed.";
}
} //pcep-stats
```

```
container sessions {
  description
    "This entry represents a single PCEP
    session in which the local PCEP entity
    participates.
    This entry exists only if the
    corresponding PCEP session has been
    initialized by some event, such as
    manual user configuration, auto-
    discovery of a peer, or an incoming
    TCP connection.";
```

```
list session {
  key "initiator";

  description
    "The list of sessions, note that
    for a time being two sessions
    may exist for a peer";

  leaf initiator {
    type pcep-initiator;
    description
      "The initiator of the session,
      that is, whether the TCP
      connection was initiated by
      the local PCEP entity or the
      peer.
      There is a window during
      session initialization where
      two sessions can exist between
      a pair of PCEP speakers, each
      initiated by one of the
      speakers. One of these
      sessions is always discarded
      before it leaves OpenWait state.
      However, before it is discarded,
      two sessions to the given peer
      appear transiently in this MIB
      module. The sessions are
      distinguished by who initiated
      them, and so this field is the
      key.";
  }

  leaf state-last-change {
    type yang:timestamp;
    description
      "The timestamp value at the
      time this session entered its
      current state as denoted by
      the state leaf.";
  }

  leaf state {
    type pcep-sess-state;
    description
      "The current state of the
      session.
      The set of possible states
```

```
        excludes the idle state since
        entries do not exist in the
        idle state.";
    }

    leaf session-creation {
        type yang:timestamp;
        description
            "The timestamp value at the
            time this session was
            created.";
    }

    leaf connect-retry {
        type yang:counter32;
        description
            "The number of times that the
            local PCEP entity has
            attempted to establish a TCP
            connection for this session
            without success. The PCEP
            entity gives up when this
            reaches connect-max-retry.";
    }

    leaf local-id {
        type uint32 {
            range "0..255";
        }
        description
            "The value of the PCEP session
            ID used by the local PCEP
            entity in the Open message
            for this session.
            If state is tcp-pending then
            this is the session ID that
            will be used in the Open
            message. Otherwise, this is
            the session ID that was sent
            in the Open message.";
    }

    leaf remote-id {
        type uint32 {
            range "0..255";
        }
        must "(../state != 'tcp-pending'" +
            "and " +
```

```

        "../state != 'open-wait' )" +
        "or " +
        "(../state = 'tcp-pending'" +
        " or " +
        "../state = 'open-wait' )" +
        "and remote-id = 0))" {
            error-message
                "Invalid remote-id";
            description
                "If state is tcp-pending
                or open-wait then this
                leaf is not used and
                MUST be set to zero.";
        }
    description
        "The value of the PCEP session
        ID used by the peer in its
        Open message for this
        session.";
}

leaf keepalive-timer {
    type uint32 {
        range "0..255";
    }
    units "seconds";
    must "(../state = 'session-up'" +
    "or " +
    "(../state != 'session-up'" +
    "and keepalive-timer = 0))" {
        error-message
            "Invalid keepalive
            timer";
        description
            "This field is used if
            and only if state is
            session-up. Otherwise,
            it is not used and
            MUST be set to
            zero.";
    }
}
description
    "The agreed maximum interval at
    which the local PCEP entity
    transmits PCEP messages on this
    PCEP session. Zero means that
    the local PCEP entity never
    sends Keepalives on this

```

```
        session.>";
    }

leaf peer-keepalive-timer {
    type uint32 {
        range "0..255";
    }
    units "seconds";
    must "(../state = 'session-up'" +
        "or " +
        "(../state != 'session-up'" +
        "and " +
        "peer-keepalive-timer = 0))" {
        error-message
            "Invalid Peer keepalive
            timer";
        description
            "This field is used if
            and only if state is
            session-up. Otherwise,
            it is not used and MUST
            be set to zero.";
    }
    description
        "The agreed maximum interval at
        which the peer transmits PCEP
        messages on this PCEP session.
        Zero means that the peer never
        sends Keepalives on this
        session.";
}

leaf dead-timer {
    type uint32 {
        range "0..255";
    }
    units "seconds";
    description
        "The dead timer interval for
        this PCEP session.";
}

leaf peer-dead-timer {
    type uint32 {
        range "0..255";
    }
    units "seconds";
    must "(../state != 'tcp-pending'" +
```

```

        "and " +
        "../state != 'open-wait' )" +
        "or " +
        "((../state = 'tcp-pending'" +
        " or " +
        "../state = 'open-wait' )" +
        "and " +
        "peer-dead-timer = 0))" {
            error-message
                "Invalid Peer Dead
                timer";
            description
                "If state is tcp-
                pending or open-wait
                then this leaf is not
                used and MUST be set to
                zero.";
        }
        description
            "The peer's dead-timer interval
            for this PCEP session.";
    }

    leaf ka-hold-time-rem {
        type uint32 {
            range "0..255";
        }
        units "seconds";
        must "((../state != 'tcp-pending'" +
        "and " +
        "../state != 'open-wait' )" +
        "or " +
        "((../state = 'tcp-pending'" +
        "or " +
        "../state = 'open-wait' )" +
        "and " +
        "ka-hold-time-rem = 0))" {
            error-message
                "Invalid Keepalive hold
                time remaining";
            description
                "If state is tcp-pending
                or open-wait then this
                field is not used and
                MUST be set to zero.";
        }
        description
            "The keep alive hold time

```



```
        remaining for this session.";
    }

    leaf overloaded {
        type boolean;
        description
            "If the local PCEP entity has
            informed the peer that it is
            currently overloaded, then this
            is set to true. Otherwise, it
            is set to false.";
    }

    leaf overload-time {
        type uint32;
        units "seconds";
        must "(../overloaded = true or" +
            "(../overloaded != true and" +
            " overload-time = 0))" {
            error-message
                "Invalid overload-time";
            description
                "This field is only used
                if overloaded is set to
                true. Otherwise, it is
                not used and MUST be set
                to zero.";
        }
        description
            "The interval of time that is
            remaining until the local PCEP
            entity will cease to be
            overloaded on this session.";
    }

    leaf peer-overloaded {
        type boolean;
        description
            "If the peer has informed the
            local PCEP entity that it is
            currently overloaded, then this
            is set to true. Otherwise, it
            is set to false.";
    }

    leaf peer-overload-time {
        type uint32;
        units "seconds";
    }
}
```

```
must "(../peer-overloaded = true" +
" or " +
"(../peer-overloaded != true" +
" and " +
"peer-overload-time = 0))" {
    error-message
        "Invalid peer overload
        time";
    description
        "This field is only used
        if peer-overloaded is
        set to true. Otherwise,
        it is not used and MUST
        be set to zero.";
}
description
    "The interval of time that is
    remaining until the peer will
    cease to be overloaded. If it
    is not known how long the peer
    will stay in overloaded state,
    this leaf is set to zero.";
}
leaf lspdb-sync {
    if-feature stateful;
    type sync-state;
    description
        "The LSP-DB state synchronization
        status.";
}
leaf discontinuity-time {
    type yang:timestamp;
    description
        "The timestamp value of the time
        when the statistics were last
        reset.";
}
}
container pcep-stats {
    description
        "The container for all statistics
        at session level.";
    uses pcep-stats{
        description
            "The statistics contained are
            for the current sessions to
            that peer. These are lost
            when the session goes down.
```

```

";
    }
  } // pcep-stats
} // session
} // sessions
} // peer
} // peers
} // entity
} // pcep-state

/*
 * Notifications
 */
notification pcep-session-up {
  description
    "This notification is sent when the value of
    '/pcep/pcep-state/peers/peer/sessions/session/state'
    enters the 'session-up' state.";

  uses notification-instance-hdr;

  uses notification-session-hdr;

  leaf state-last-change {
    type yang:timestamp;
    description
      "The timestamp value at the time this session entered
      its current state as denoted by the state leaf.";
  }

  leaf state {
    type pcep-sess-state;
    description
      "The current state of the session.
      The set of possible states excludes the idle state
      since entries do not exist in the idle state.";
  }
} // notification

notification pcep-session-down {
  description
    "This notification is sent when the value of
    '/pcep/pcep-state/peers/peer/sessions/session/state'
    leaves the 'session-up' state.";

  uses notification-instance-hdr;

```

```
leaf session-initiator {
    type pcep-initiator;
    description
        "The initiator of the session.";
}

leaf state-last-change {
    type yang:timestamp;
    description
        "The timestamp value at the time this session entered
        its current state as denoted by the state leaf.";
}

leaf state {
    type pcep-sess-state;
    description
        "The current state of the session.
        The set of possible states excludes the idle state
        since entries do not exist in the idle state.";
}
} //notification

notification pcep-session-local-overload {
    description
        "This notification is sent when the local PCEP entity
        enters overload state for a peer.";

    uses notification-instance-hdr;

    uses notification-session-hdr;

    leaf overloaded {
        type boolean;
        description
            "If the local PCEP entity has informed the peer that
            it is currently overloaded, then this is set to
            true. Otherwise, it is set to false.";
    }

    leaf overload-time {
        type uint32;
        units "seconds";
        must "(../overloaded = true or " +
            "(../overloaded != true and " +
            "overload-time = 0))" {
            error-message
                "Invalid overload-time";
            description

```

```
        "This field is only used if overloaded is
        set to true. Otherwise, it is not used
        and MUST be set to zero.";
    }
    description
        "The interval of time that is remaining until the
        local PCEP entity will cease to be overloaded on
        this session.";
}
} //notification

notification pcep-session-local-overload-clear {
    description
        "This notification is sent when the local PCEP entity
        leaves overload state for a peer.";

    uses notification-instance-hdr;

    leaf overloaded {
        type boolean;
        description
            "If the local PCEP entity has informed the peer
            that it is currently overloaded, then this is set
            to true. Otherwise, it is set to false.";
    }
} //notification

notification pcep-session-peer-overload {
    description
        "This notification is sent when a peer enters overload
        state.";

    uses notification-instance-hdr;

    uses notification-session-hdr;

    leaf peer-overloaded {
        type boolean;
        description
            "If the peer has informed the local PCEP entity that
            it is currently overloaded, then this is set to true.
            Otherwise, it is set to false.";
    }

    leaf peer-overload-time {
        type uint32;
        units "seconds";
        must "(../peer-overloaded = true or " +
```

```

        "(../peer-overloaded != true and " +
        "peer-overload-time = 0))" {
            error-message
                "Invalid peer-overload-time";
            description
                "This field is only used if
                peer-overloaded is set to true.
                Otherwise, it is not used and MUST
                be set to zero.";
        }
    description
        "The interval of time that is remaining until the
        peer will cease to be overloaded.  If it is not known
        how long the peer will stay in overloaded state, this
        leaf is set to zero.";
}
} //notification

notification pcep-session-peer-overload-clear {
    description
        "This notification is sent when a peer leaves overload
        state.";

    uses notification-instance-hdr;

    leaf peer-overloaded {
        type boolean;
        description
            "If the peer has informed the local PCEP entity that
            it is currently overloaded, then this is set to true.
            Otherwise, it is set to false.";
    }
} //notification
} //module

```

<CODE ENDS>

9. Security Considerations

The YANG module defined in this memo is designed to be accessed via the NETCONF protocol [RFC6241]. The lowest NETCONF layer is the secure transport layer and the mandatory-to-implement secure transport is SSH [RFC6242]. The NETCONF access control model [RFC6536] provides the means to restrict access for particular NETCONF users to a pre-configured subset of all available NETCONF protocol operations and content.

There are a number of data nodes defined in the YANG module which are writable/creatable/deletable (i.e., config true, which is the default). These data nodes may be considered sensitive or vulnerable in some network environments. Write operations (e.g., <edit-config>) to these data nodes without proper protection can have a negative effect on network operations.

TBD: List specific Subtrees and data nodes and their sensitivity/vulnerability.

10. Manageability Considerations

10.1. Control of Function and Policy

10.2. Information and Data Models

10.3. Liveness Detection and Monitoring

10.4. Verify Correct Operations

10.5. Requirements On Other Protocols

10.6. Impact On Network Operations

11. IANA Considerations

This document registers a URI in the "IETF XML Registry" [RFC3688]. Following the format in RFC 3688, the following registration has been made.

URI: urn:ietf:params:xml:ns:yang:ietf-pcep

Registrant Contact: The PCE WG of the IETF.

XML: N/A; the requested URI is an XML namespace.

This document registers a YANG module in the "YANG Module Names" registry [RFC6020].

Name:	ietf-pcep
Namespace:	urn:ietf:params:xml:ns:yang:ietf-pcep
Prefix:	pcep
Reference:	This I-D

12. Acknowledgements

The initial document is based on the PCEP MIB [RFC7420]. Further this document structure is based on Routing Yang Module [I-D.ietf-netmod-routing-cfg]. We would like to thank the authors of aforementioned documents.

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Appendix A. Contributor Addresses

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PCEP Extensions for SFC in SPRING Networks
draft-xu-pce-sr-sfc-02

Abstract

[I-D.xu-spring-pce-based-sfc-arch] describes a PCE-based SFC architecture in which the PCE is used to compute service function paths in SPRING networks. Based on the above architecture, this document describes extensions to the Path Computation Element Protocol (PCEP) that allow a PCE to compute and instantiate service function paths in SPRING networks. The extensions specified in this document are applicable to both the stateless PCE model and the stateful PCE model.

Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

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1. Introduction

Service Function Chaining [I-D.ietf-sfc-architecture] provides a flexible way to construct services. When applying a particular Service Function Chain (SFC) to the traffic classified by the Classifier, the traffic needs to be steered through an ordered set of Service Functions (SF) in the network. This ordered set of SFs in the network, referred to as a Service Function Path (SFP), is an instantiation of the SFC in the network. For example, as shown in Figure 1, an SFP corresponding to the SFC of {SF1, SF3} can be expressed as {SFF1, SF1, SFF2, SF3}.

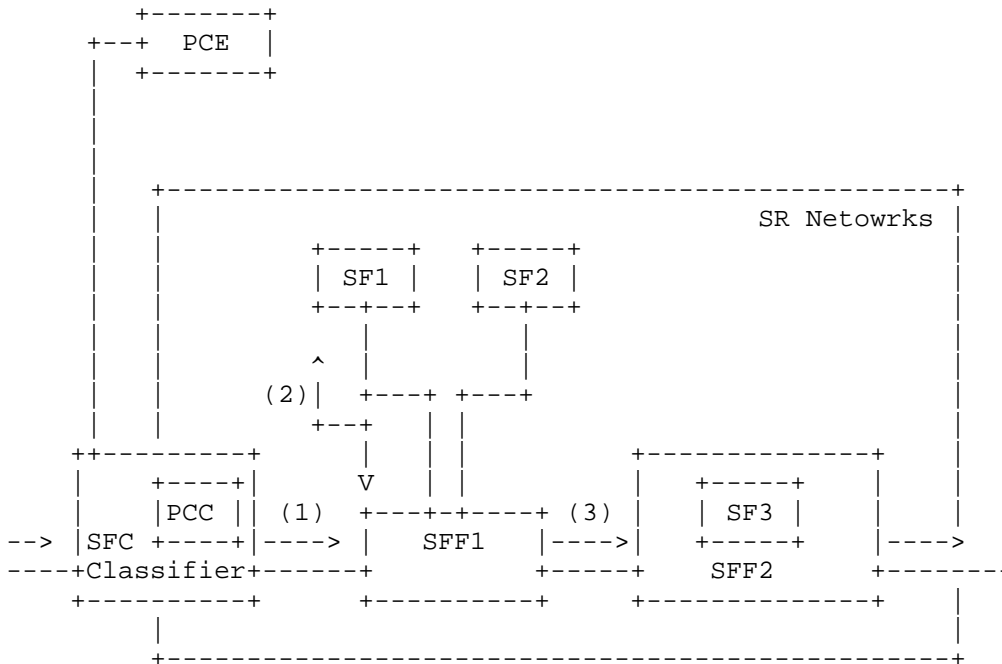


Figure 1: PCE-based Service Function Chaining in SR Network

[I-D.xu-spring-pce-based-sfc-arch] describes a PCE-based SFC architecture in which the PCE is used to compute an SFP (i.e., instantiate a service function chain) in SPRING networks (a.k.a., Segment Routing networks or SR networks in short). This document describes extensions to the PCEP on basis of that architecture. The extensions specified in this document are applicable to both the stateless PCE model defined in [RFC5440] and the stateful PCE model defined in [I-D.ietf-pce-stateful-pce].

2. Terminology

This section contains definitions for terms used frequently throughout this document. However, many additional definitions can be found in [RFC5440], [I-D.sivabalan-pce-segment-routing] and [I-D.xu-spring-pce-based-sfc-arch].

PCC: Path Computation Client

PCE: Path Computation Element

PCEP: Path Computation Element Protocol

ERO: Explicit Route Object

Service Function Chain (SFC): A service function chain defines a set of abstract service functions and ordering constraints that must be applied to packets and/or frames selected as a result of classification.

SF Identifier (SF ID): A unique identifier that represents a service function within an SFC-enabled domain.

Service Function Forwarder (SFF): A service function forwarder is responsible for delivering traffic received from the network to one or more connected service functions according to information carried in the SFC encapsulation, as well as handling traffic coming back from the SF.

Service Function Path (SFP): The SFP provides a level of indirection between the fully abstract notion of service chain as a sequence of abstract service functions to be delivered, and the fully specified notion of exactly which SFF/SFs the packet will visit when it actually traverses the network. Specifically, it is an ordered list of SFFs and SF IDs.

Compact SFP: An ordered list of SFFs.

SID: Segment Identifier

Service Function SID: A locally unique SID indicating a particular service function on an SFF.

SR: Segment Routing

SR-specific SFP: An ordered list of node SIDs (representing SFFs) and Service Function SIDs.

Compact SR-specific SFP: An ordered list of node SIDs (representing SFFs).

3. Overview of PCEP Extensions for SFC in SR Networks

As discussed in [I-D.xu-spring-pce-based-sfc-arch], the PCC provides an ordered list of SF IDs to the PCE and indicates to the PCE that what type SFs and paths are requested (e.g., an SFP, or a compact SFP, or an SR-specific SFP, or a compact SR-specific SFP) through the path computation request message, and then the PCE responds with a corresponding path through the path computation response message. This specification is applicable to both stateful and stateless PCEs.

4. PCEP Message Extensions for SR-based SFC

4.1. PCReq Message

This document does not specify any changes to the PCReq message format. This document requires the PATH-SETUP-TYPE TLV [I-D.sivabalan-pce-lsp-setup-type] to be carried in the RP Object in order for a PCC to request a particular type of path. Four new Path Setup Types need to be defined for SR-based SFC, or SR-SFC in short (Section 5.2). This document also requires the Include Route Object (IRO) to be carried in the PCReq message in order for a PCC to specify SFC. A new IRO sub-object type needs to be defined for SF (Section 5.3).

4.2. PCRep Message

This document defines the format of the PCRep message carrying an SFP. The message is sent by a PCE to a PCC in response to a previously received PCReq message, where the PCC requested an SFP. The format of the SFC-specific PCRep message is as follows:

```
<PCRep Message> ::= <Common Header>
                        <response-list>
```

Where:

```
<response-list> ::= <response> [<response-list>]
```

```
<response> ::= <RP>
                [<NO-PATH>]
                [<path-list>]
```

Where:

```
<path-list> ::= <SR-SFC-ERO> [<path-list>]
```

The RP and NO-PATH Objects are defined in [RFC5440]. The <SR-SFC-ERO> object contains the SFP and is defined in Section 5.4.

4.3. PCUpd Message

This document defines the format of the PCUpd message carrying an SFP update. The message is sent forwardly by a PCE to a PCC to update an previously computed SFP.

The format of the PCUpd message is as follows:

```
<PCUpd Message> ::= <Common Header>
                        <update-request-list>
```

Where:

```
<update-request-list> ::= <update-request> [<update-request-list>]
```

```
<update-request> ::= <SRP> <path-list>
```

Where:

```
<path-list> ::= <SR-SFC-ERO> [<path-list>]
```

4.4. PCRpt Message

PCRpt message sent from a PCC to PCE as a respond to a PCUpd message or in an unsolicited manner (e.g., during state synchronization).

The format of the PCUpd message is as follows:

```
<PCUpd Message> ::= <Common Header>
                        <state-report-list>
```

Where:

```
<state-report-list> ::= <state-report> [<state-report-list>]
<state-report> ::= [<SRP>] <path-list>
```

Where:

```
<path-list> ::= <SR-SFC-ERO> [<path-list>]
```

5. Object Formats

5.1. OPEN Object

This document defines a new optional TLV for use in the OPEN Object.

5.1.1. SR-SFC PCE Capability TLV

The SR-SFC-PCE-CAPABILITY TLV is an optional TLV for use in the OPEN Object to negotiate SR-SFC capability on the PCEP session. The format of the SR-SFC-PCE-CAPABILITY TLV is shown in the following Figure 2:

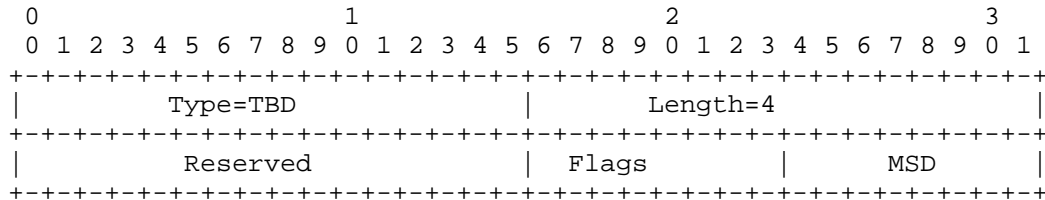


Figure 2: SR-SFC-PCE-CAPABILITY TLV format

The code point for the TLV type is to be defined by IANA. The TLV length is 4 octets. The 32-bit value is formatted as follows. The "Maximum SID Depth" (1 octet) field (MSD) specifies the maximum number of SIDs that a PCC is capable of imposing on a packet. The "Flags" (1 octet) and "Reserved" (2 octets) fields are currently unused, and MUST be set to zero and ignored on receipt.

5.1.1.1. Negotiating SR-SFC Capability

The SR-SFC capability TLV is contained in the OPEN object. By including the TLV in the OPEN message to a PCE, a PCC indicates its support for SFPs. By including the TLV in the OPEN message to a PCC, a PCE indicates that it is capable of computing SFPs.

5.2. RP/SRP Object

In order to setup an SFP, the RP or SRP object MUST carry a PATH-SETUP-TYPE TLV specified in [I-D.sivabalan-pce-lsp-setup-type]. This document defines four new Path Setup Types (PST) for SR-SFC as follows:

PST = 2: The path is an SFP.

PST = 3: The path is a compact SFP.

PST = 4: The path is an SR-specific SFP.

PST = 5: The path is a compact SR-specific SFP.

5.3. Include Route Object

The IRO (Include Route Object) MUST be carried within PCReq messages to indicate a particular SFC. Furthermore, the IRO MAY be carried in PCRep messages. When carried within a PCRep message with the NO-PATH object, the IRO indicates the set of service functions that cause the PCE to fail to find a path.

This document defines a new sub-object type for the SR-SFC as follows:

Type	Sub-object
5	Service Function ID

5.4. SR-SFC-ERO Object

Generally speaking, an SR-SFC-ERO object consists of one or more ERO subobjects described in the following sub-sections to represent a particular type of service function path. In the ERO subobject, each SID is associated with an identifier that represents either a service node or a service function. This identifier is referred to as the 'Node or Service Identifier' (NSI). As described later, an NSI can be represented in various formats (e.g., IPv4 address, IPv6 address, SF identifier, etc). Specifically, in the SFP case, the NSI of every ERO subobject contained in the SR-SFC-ERO object represents a service

node or a service function while the SID of each ERO subobject is set to null. In the compact SFP case, the NSI of every ERO subobject contained in the SR-SFC-ERO object only represents an SFF meanwhile the SID of every ERO subobject is set to null. In the SR-specific SFP, the NSI of every ERO subobject contained in the SR-SFC-ERO object represents an SFF or a service function while the SID of every ERO subject MUST NOT be null. In the compact SR-specific SFP, the NSI of every ERO subobject contained in the SR-SFC-ERO object represents an SFF meanwhile the SID of every ERO subobject MUST NOT be null.

5.4.1. SR-SFC-ERO Subobject

An SR-SFC-ERO subobject (as shown in Figure 3) consists of a 32-bit header followed by the SID and the NSI associated with the SID. The SID is a 32-bit or 128 bit number. The size of the NSI depends on its respective type, as described in the following sub-sections.

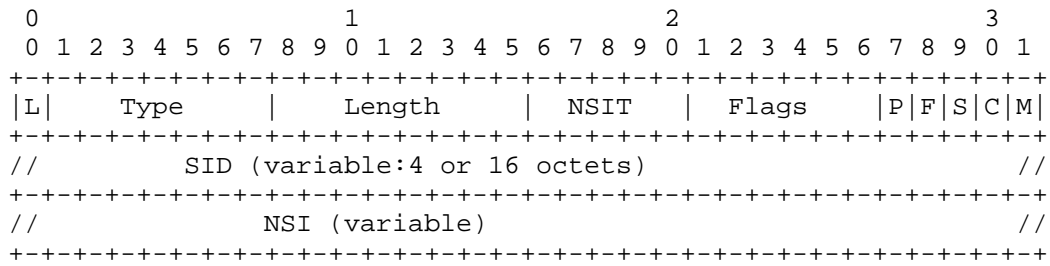


Figure 3: SR-SFC-ERO Subobject Format

The fields in the ERO Subobject are as follows:

'L' Flag: indicates whether the subobject represents a loose-hop in the explicit route [RFC3209]. If this flag is unset, a PCC MUST not overwrite the SID value present in the SR-SFC-ERO subobject. Otherwise, a PCC MAY expand or replace one or more SID value(s) in the received SR-SFC-ERO based on its local policy.

Type: is the type of the SR-SFC-ERO Subobject. This document defines the SR-SFC-ERO Subobject type. A new code point will be requested for the SR-SFC-ERO Subobject from IANA.

Length: contains the total length of the subobject in octets, including the L, Type and Length fields. Length MUST be at least 4, and MUST be a multiple of 4.

NSI Type (NSIT): indicates the type of NSI associated with the SID. The NSI-Type values are described later in this document.

Flags: is used to carry any additional information pertaining to SID. Currently, the following flag bits are defined:

M: When this bit is set, the SID value represents an MPLS label stack entry as specified in [RFC5462], where only the label value is specified by the PCE. Other fields (TC, S, and TTL) fields MUST be considered invalid, and PCC MUST set these fields according to its local policy and MPLS forwarding rules.

C: When this bit as well as the M bit are set, then the SID value represents an MPLS label stack entry as specified in [RFC5462], where all the entry's fields (Label, TC, S, and TTL) are specified by the PCE. However, a PCC MAY choose to override TC, S, and TTL values according its local policy and MPLS forwarding rules.

S: When this bit is set, the SID value in the subobject body is null. In this case, the PCC is responsible for choosing the SID value, e.g., by looking up its Traffic Engineering Database (TED) using node/service identifier in the subobject body.

F: When this bit is set, the NSI value in the subobject body is null.

P: When this bit is set, the SID value represents an IPv6 address.

SID: is the 4-octect or 16-octect Segment Identifier

NSI: contains the NSI associated with the SID. Depending on the value of NSIT, the NSI can have different format as described in the following sub-section.

5.4.2. NSI Associated with SID

This document defines the following NSIs:

'IPv4 Node ID': is specified as an IPv4 address. In this case, NSIT and Length are 1 and 12 respectively.

'IPv6 Node ID': is specified as an IPv6 address. In this case, NSIT and Length are 2 and 24 respectively.

'Service Function ID': is specified as an SF ID. In this case, NSIT and Length are TBD.

5.4.3. SR-SFC-ERO Processing

TBD.

6. IANA Considerations

6.1. PCEP Objects

IANA is requested to allocate an ERO subobject type (recommended value= 6) for the SR-SFC-ERO subobject.

6.2. PCEP-Error Object

TBD.

6.3. PCEP TLV Type Indicators

This document defines the following new PCEP TLVs:

Value	Meaning	Reference
27	SR-SFC-PCE-CAPABILITY	This document

6.4. New Path Setup Type

This document defines a new setup type for the PATH-SETUP-TYPE TLV as follows:

Value	Description	Reference
2	The path is an SFP.	This document
3	The path is a compact SFP.	This document
4	The path is an SR-specific SFP.	This document
5	The path is a compact SR-specific SFP.	This document

6.5. New IRO Sub-object Type

This document defines a new IRO sub-object type for the SFC as follows:

Type	Sub-object
5	Service Function ID

7. Security considerations

This document does not introduce any new security considerations.

8. Acknowledgement

TBD.

9. References

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PCEP Extensions for Unified Source Routing-based SFC
draft-xu-pce-sr-sfc-05

Abstract

MPLS-SPRING (a.k.a., MPLS Segment Routing) could be leveraged to realize a unified source routing mechanism across MPLS, IPv4 and IPv6 data planes by using a unified source routing instruction set while preserving backward compatibility with MPLS-SPRING. More specifically, the source routing instruction set information contained in a source routed packet could be uniformly encoded as an MPLS label stack no matter the underlay is IPv4, IPv6 or MPLS. The unified source routing mechanism could be leveraged to realize a transport-independent service function chaining by encoding the service function path information or service function chain information as an MPLS label stack. This document describes extensions to the Path Computation Element Protocol (PCEP) that allow a PCE to compute and instantiate service function paths in the unified source routing based service function chaining context. The extensions specified in this document are applicable to both the stateless PCE model and the stateful PCE model.

Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC 2119 [RFC2119].

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1. Introduction

Service Function Chaining [RFC7665] provides a flexible way to construct services. When applying a particular Service Function Chain (SFC) to the traffic classified by the Classifier, the traffic needs to be steered through an ordered set of Service Function Forwarders (SFF) and Service Functions (SF) in the network. This ordered set of SFFs and SFs in the network, referred to as a Service Function Path (SFP), is an instantiation of the SFC in the network. For example, as shown in Figure 1, an SFP corresponding to the SFC of {SF1, SF3} can be expressed as {SFF1, SF1, SFF2, SF3}.

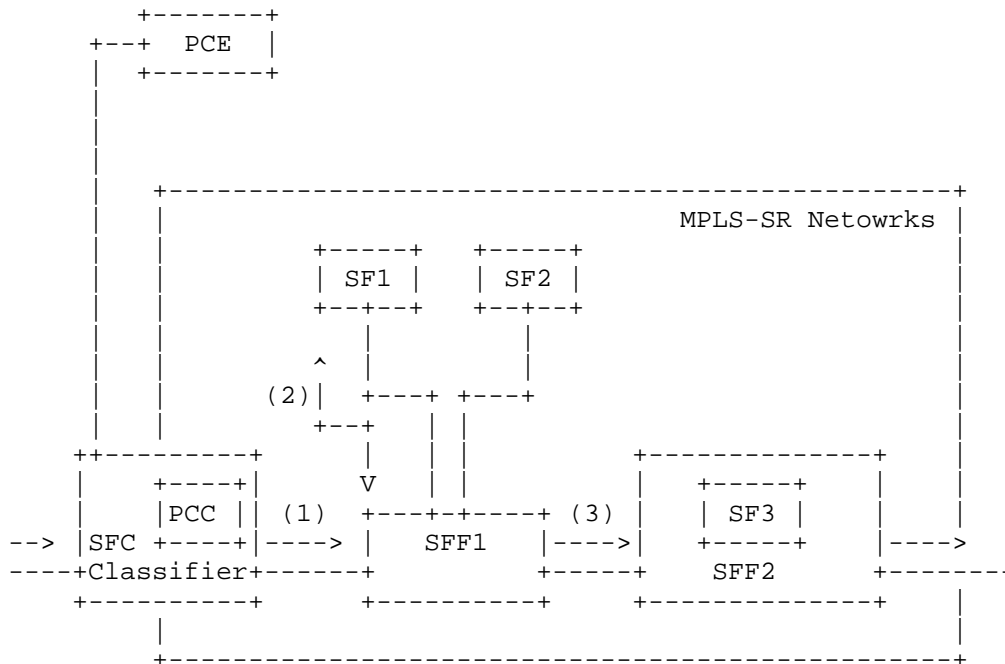


Figure 1: PCE-based Service Function Chaining in MPLS-SR Network

MPLS-SPRING (a.k.a., MPLS Segment Routing) could be leveraged to realize a unified source routing mechanism across MPLS, IPv4 and IPv6 data planes by using a unified source routing instruction set while preserving backward compatibility with MPLS-SPRING as described in [I-D.xu-mpls-unified-source-routing-instruction]. More specifically, the source routing instruction set information contained in a source routed packet could be uniformly encoded as an MPLS label stack no matter the underlay is IPv4, IPv6 or MPLS. The unified source routing mechanism in turn could be leveraged to realize a transport-independent service function chaining by encoding the service function path information or service function chain information as an MPLS label stack as described in [I-D.xu-mpls-service-chaining].

This document describes extensions to the Path Computation Element Protocol (PCEP) that allow a PCE to compute and instantiate service function paths in the MPLS source routing based service function chaining context. More specifically, the PCC provides an ordered list of SF IDs to the PCE and indicates to the PCE that what type SFs and paths are requested (e.g., an SFP, or a compact SFP, or an SR-specific SFP, or a compact SR-specific SFP) through the path computation request message, and then the PCE responds with a corresponding path through the path computation response message.

The extensions specified in this document are applicable to both the stateless PCE model [RFC5440] and the stateful PCE model [I-D.ietf-pce-stateful-pce].

2. Terminology

This memo makes use of the terms defined in [RFC5440], [I-D.ietf-pce-segment-routing] and [I-D.xu-mpls-service-chaining]. In addition, this memo defines the following two additional terms:

Compact SFP: An ordered list of SFPs.

SR-specific SFP: An ordered list of node SIDs (representing SFPs) and Service Function SIDs.

Compact SR-specific SFP: An ordered list of node SIDs (representing SFPs).

3. PCEP Message Extensions for MPLS Source Routing-based SFC

3.1. PCReq Message

This document does not specify any changes to the PCReq message format. This document requires the PATH-SETUP-TYPE TLV [I-D.ietf-pce-lsp-setup-type] to be carried in the RP Object in order for a PCC to request a particular type of path. Four new Path Setup Types need to be defined for MPLS source routing-based SFC (see Section 4.2). This document also requires the Include Route Object (IRO) to be carried in the PCReq message in order for a PCC to specify an SFC. A new IRO sub-object type needs to be defined for SF (see Section 4.3).

3.2. PCRep Message

This document defines the format of the PCRep message carrying an SFP. The message is sent by a PCE to a PCC in response to a previously received PCReq message, where the PCC requested an SFP. The format of the SFC-specific PCRep message is defined as follows:

```
<PCRep Message> ::= <Common Header>
                        <response-list>
Where:
  <response-list> ::= <response> [<response-list>]
  <response> ::= <RP>
                [<NO-PATH>]
                [<path-list>]
Where:
  <path-list> ::= <SR-SFC-ERO> [<path-list>]
```

The RP and NO-PATH Objects are defined in [RFC5440]. The <SR-SFC-ERO> object contains an SFP and is defined in Section 4.4.

3.3. PCUpd Message

This document defines the format of the PCUpd message carrying an SFP update. The message is sent forwardly by a PCE to a PCC to update an previously computed SFP.

The format of the PCUpd message is defined as follows:

```
<PCUpd Message> ::= <Common Header>
                        <udate-request-list>
Where:
  <udate-request-list> ::= <udate-request> [ <udate-request-list> ]
  <udate-request> ::= <SRP> <path-list>
Where:
  <path-list> ::= <SR-SFC-ERO> [ <path-list> ]
```

3.4. PCRpt Message

PCRpt message sent from a PCC to PCE as a respond to a PCUpd message or in an unsolicited manner (e.g., during state synchronization).

The format of the PCRpt message is defined as follows:

```
<PCRpt Message> ::= <Common Header>
                        <state-report-list>
Where:
  <state-report-list> ::= <state-report> [ <state-report-list> ]
  <state-report> ::= [ <SRP> ] <path-list>
Where:
  <path-list> ::= <SR-SFC-ERO> [ <path-list> ]
```

4. Object Formats

4.1. OPEN Object

This document defines a new optional TLV for use in the OPEN Object.

4.1.1. SR-SFC PCE Capability TLV

The SR-SFC-PCE-CAPABILITY TLV is an optional TLV for use in the OPEN Object to negotiate SR-SFC capability on the PCEP session. The format of the SR-SFC-PCE-CAPABILITY TLV is shown in the following Figure 2:

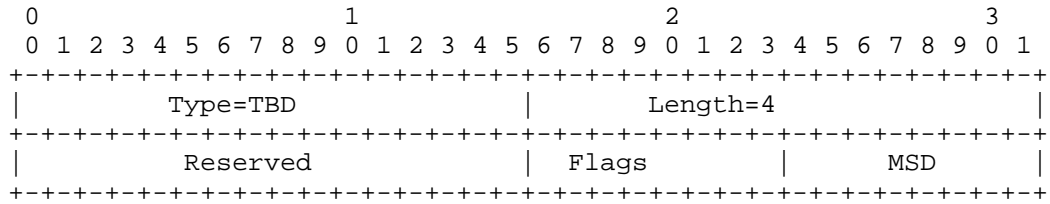


Figure 2: SR-SFC-PCE-CAPABILITY TLV Format

The code point for the TLV type is to be defined by IANA. The TLV length is 4 octets. The 32-bit value is formatted as follows. The "Maximum SID Depth" (1 octet) field (MSD) specifies the maximum number of SIDs that a PCC is capable of imposing on a packet. The "Flags" (1 octet) and "Reserved" (2 octets) fields are currently unused, and MUST be set to zero and ignored on receipt.

4.1.1.1. Negotiating SR-SFC Capability

The SR-SFC capability TLV is contained in the OPEN object. By including the TLV in the OPEN message to a PCE, a PCC indicates its support for SFPs. By including the TLV in the OPEN message to a PCC, a PCE indicates that it is capable of computing SFPs.

4.2. RP/SRP Object

In order to setup an SFP, the RP or SRP object MUST carry a PATH-SETUP-TYPE TLV specified in [I-D.ietf-pce-lsp-setup-type]. This document defines four new Path Setup Types (PST) for SR-SFC as follows:

- PST = 2: The path is an SFP.
- PST = 3: The path is a compact SFP.
- PST = 4: The path is an SR-specific SFP.
- PST = 5: The path is a compact SR-specific SFP.

4.3. Include Route Object

The IRO (Include Route Object) MUST be carried within PCReq messages to indicate a particular SFC. Furthermore, the IRO MAY be carried in PCRep messages. When carried within a PCRep message with the NO-PATH object, the IRO indicates the set of service functions that cause the PCE to fail to find a path. This document defines a new sub-object type for the SR-SFC as follows:

Type	Sub-object
5	Service Function ID

4.4. SR-SFC-ERO Object

Generally speaking, an SR-SFC-ERO object consists of one or more ERO subobjects described in the following sub-sections to represent a particular type of service function path. In the ERO subobject, each SID is associated with an identifier that represents either an SFF or an SF. This identifier is referred to as the 'Node or Service Identifier' (NSI). As described later, an NSI can be represented in various formats (e.g., IPv4 address, IPv6 address, SF identifier, etc). Specifically, in the SFP case, the NSI of every ERO subobject contained in the SR-SFC-ERO object represents an SFF or an SF while the SID of each ERO subobject is set to null. In the compact SFP case, the NSI of every ERO subobject contained in the SR-SFC-ERO object only represents an SFF meanwhile the SID of every ERO subobject is set to null. In the SR-specific SFP, the NSI of every ERO subobject contained in the SR-SFC-ERO object represents an SFF or an SF while the SID of every ERO subject MUST NOT be null. In the compact SR-specific SFP, the NSI of every ERO subobject contained in the SR-SFC-ERO object represents an SFF meanwhile the SID of every ERO subobject MUST NOT be null.

4.4.1. SR-SFC-ERO Subobject

An SR-SFC-ERO subobject (as shown in Figure 3) consists of a 32-bit header followed by the SID and the NSI associated with the SID. The SID is a 32-bit or 128 bit number. The size of the NSI depends on its respective type, as described in the following sub-sections.

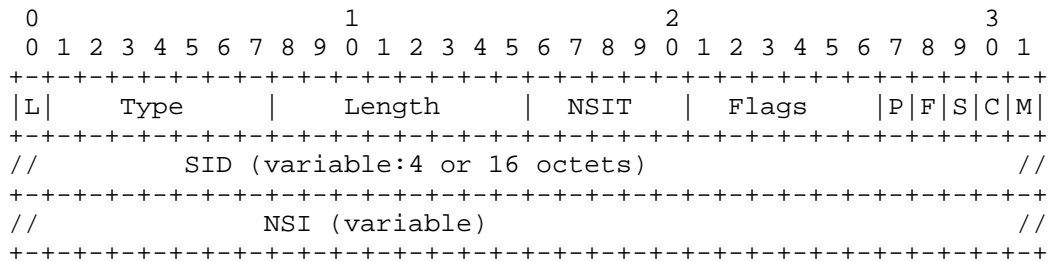


Figure 3: SR-SFC-ERO Subobject Format

'L' Flag: indicates whether the subobject represents a loose-hop in the explicit route [RFC3209]. If this flag is unset, a PCC MUST not overwrite the SID value present in the SR-SFC-ERO

subobject. Otherwise, a PCC MAY expand or replace one or more SID value(s) in the received SR-SFC-ERO based on its local policy.

Type: is the type of the SR-SFC-ERO Subobject. This document defines the SR-SFC-ERO Subobject type. A new code point will be requested for the SR-SFC-ERO Subobject from IANA.

Length: contains the total length of the subobject in octets, including the L, Type and Length fields. Length MUST be at least 4, and MUST be a multiple of 4.

NSI Type (NSIT): indicates the type of NSI associated with the SID. The NSI-Type values are described later in this document.

Flags: is used to carry any additional information pertaining to SID. Currently, the following flag bits are defined:

M: When this bit is set, the SID value represents an MPLS label stack entry as specified in [RFC5462], where only the label value is specified by the PCE. Other fields (TC, S, and TTL) fields MUST be considered invalid, and PCC MUST set these fields according to its local policy and MPLS forwarding rules.

C: When this bit as well as the M bit are set, then the SID value represents an MPLS label stack entry as specified in [RFC5462], where all the entry's fields (Label, TC, S, and TTL) are specified by the PCE. However, a PCC MAY choose to override TC, S, and TTL values according its local policy and MPLS forwarding rules.

S: When this bit is set, the SID value in the subobject body is null. In this case, the PCC is responsible for choosing the SID value, e.g., by looking up its Traffic Engineering Database (TED) using node/service identifier in the subobject body.

F: When this bit is set, the NSI value in the subobject body is null.// When will the NSI value is null?

P: When this bit is set, the SID value represents an IPv6 address.

SID: is the 4-octect or 16-octect Segment Identifier.

NSI: contains the NSI associated with the SID. Depending on the value of NSIT, the NSI can have different format as described in the following sub-section.

4.4.2. NSI Associated with SID

This document defines the following NSIs:

'IPv4 Node ID': is specified as an IPv4 address. In this case, NSIT and Length are 1 and 12 respectively.

'IPv6 Node ID': is specified as an IPv6 address. In this case, NSIT and Length are 2 and 24 respectively.

'Service Function ID': is specified as an SF ID. In this case, NSIT and Length are TBD.

4.4.3. SR-SFC-ERO Processing

TBD

5. Acknowledgements

TBD.

6. IANA Considerations

6.1. PCEP Objects

IANA is requested to allocate an ERO subobject type (recommended value= 6) for the SR-SFC-ERO subobject.

6.2. PCEP-Error Object

TBD

6.3. PCEP TLV Type Indicators

This document defines the following new PCEP TLV:

Value	Meaning	Reference
27	SR-SFC-PCE-CAPABILITY	This document

6.4. New Path Setup Type

This document defines the following four new setup types for the PATH-SETUP-TYPE TLV:

Value	Description	Reference
2	The path is an SFP.	This document
3	The path is a compact SFP.	This document
4	The path is an SR-specific SFP.	This document
5	The path is a compact SR-specific SFP.	This document

6.5. New IRO Sub-object Type

This document defines a new IRO sub-object type for SFC as follows:

Type	Sub-object
5	Service Function ID

7. Security Considerations

This document does not introduce any new security considerations.

8. References

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Extensions to Path Computation Element Protocol (PCEP) to Support
Resource Sharing-based Path Computation

draft-zhang-pce-resource-sharing-03.txt

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Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

Abstract

Resource sharing in a network means two or more Label Switched Paths (LSPs) use common piece(s) of resource along their paths. This can help save network resource and useful in scenarios such as LSP recovery or when two LSPs do not need to be active at the same time. A Path Computation Element (PCE) is a centralized entity, responsible for path computation. Given this feature and its access to the network resource information and possibly active LSPs information, it can be used to support resource-sharing-based path computation with better efficiency.

This document extends the Path Computation Element Protocol (PCEP) in order to support resource sharing-based path computation.

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1. Introduction and Motivation

A Path Computation Element (PCE) provides an alternative way for providing path computation function, and it is especially useful in the scenarios where complex constraints and/or a demanding amount of computation resource are required [RFC4655]. The development of PCE standardization has evolved from stateless to stateful. A stateful PCE has access to the LSP database information of the network(s) it serves as a computation engine [Stateful-PCE]. Unless specified, this document assumes a PCE mentioned is a stateful PCE (either passive or active).

Resource sharing denotes that two or more Label Switched Paths (LSPs) share common piece(s) of resource, (such as a common time slot of a link in an Optical Transport Network (OTN)). This is usually useful in the scenario where only one LSP is active and the benefit herein is to save network resources. A simple example of this is dynamically calculating a LSP for an existing LSP undergoing a link failure. Note that the resource sharing can be worked out using a stateless PCE, but the mechanism may be complex and is out the scope of this draft.

This document considers the following requirement: resource sharing with one or multiple existing LSPs.

In a single domain, this is a common requirement in the recovery cases especially in order to increase traffic resilience against failure while reducing the amount of network resource used for recovery purpose [RFC4428].

The current protocol supporting the communication between a PCE and a Path Computation Client (PCC), i.e. PCE Protocol (PCEP), allows for re-optimization of an existing LSP [RFC5440]. This is achieved by setting R bit in the Request Parameter (RP) object, together with some additional information if applicable, in the Path Computation Request (PCReq) message sent from a PCC to the PCE. To support this type of resource sharing, a PCC needs to ask a PCE to compute a new path with the constraints of sharing resource with one or multiple existing LSPs. It is worth noting the 'resource sharing' in this draft not only means one LSP re-using the same link(s) of another LSP, but also the same slice of bandwidth. This may occur when an LSP is required for re-routing, or online re-optimization. Current PCEP specifications do not provide such function.

As mentioned in [stateful-PCE], the PLSP-ID is unique during a PCEP session between PCC and PCE. Such identification is helpful in supporting the above resource sharing requirement for standardization of stateful PCEs. With a unique identifier, the configuration of PCCs is greatly simplified. Instead of determining all the resources to be shared, the PCC could request resource sharing directly from PCE.

The resource sharing can also be required in an inter-layer PCEP session. This is similar to the previous requirement. However, it is more complex and therefore deserves a more detailed explanation here.

In a multi-layer network, Label Switched Paths (LSPs) in a lower layer are used to carry higher-layer LSPs across the lower-layer network [RFC5623]. Therefore, the resource sharing constraints in the higher layer might actually relate to the resource sharing in the lower layer. Thus, it is useful to consider how this can be achieved and whether additional extensions are needed using the models defined in [RFC5623].

In the next sections, use cases are provided to show what information needs to be exchanged to fulfill these requirements. This memo then provides extensions to PCEP to enable this function.

2. Motivation

2.1. Single PCE Use Case

Figure 1 shows a single domain network with a stateful PCE. Assume a working LSP (N1-N2-N3) exists in the network, when there is failure on the link N2-N3, it is desired to set up a restoration path for this working LSP. Suppose N1 serves as the PCC and sends a request to the stateful PCE for such an LSP. Before sending the request, N1 may need to check what policy should be applied for the path computation. For example, it might value resource sharing and prefer to share as much resource with the working LSP as possible and specify this policy in the PCReq message. If resources are shared between the old and new LSPs, there will be some 'interruption' when the traffic is switched from the old LSP to the new LSP.

On the other hand, in some scenarios there are different policies, for example the LSP should be restored without any interruption with best effort. An example can be found in Fig. 1 without failure on N2-N3 link, instead, an online re-optimization is needed for the working LSP (N1-N2-N3) from the stateful PCE. In such cases, the best choice is to set up a backup LSP for the working LSP with

totally separate routing (for example N1-N5-N4-N3), and move the traffic to that backup LSP. After that the working LSP can be torn down, which will not result in any interruption during the optimization procedure. This can actually be implemented with existing PCEP mechanism. However, if there is no such separate path, existing PCEP will reply error. A secondary option for this case is to set up an LSP and complete such re-optimization with resource sharing, even if some interruption introduced. Given the resource from the LSP to be interrupted, there may be some solutions instead of Path Compute error due to the lack of resource.

A simple illustration is provided below:

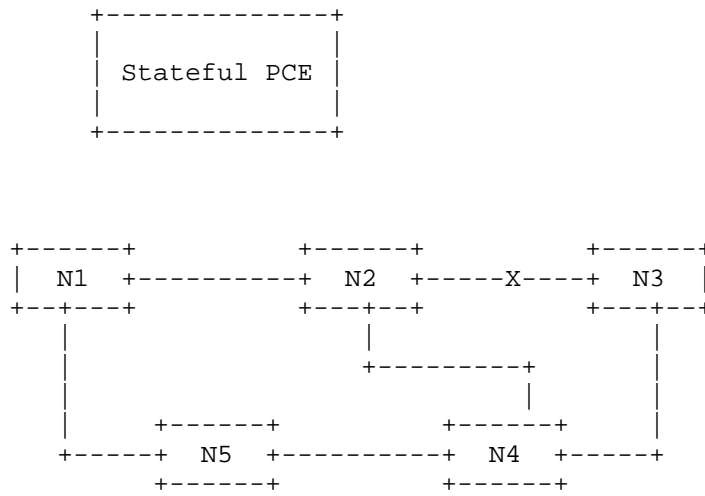


Figure 1: A Single Domain Example

Available recovery paths computed by the stateful PCE:

- LSP1: N1-N2-N4-N3
- LSP2: N1-N5-N4-N3

If resource sharing is preferred, the stateful PCE will reply with LSP1 information. Instead, if PCC prefer to have less interruption, PCE will reply with LSP2 information.

Another piece of information that needs to be conveyed to the PCE is the information about the working path LSP. Note this simple use case assumes end-to-end recovery. But in order to be applicable to use cases such as shared mesh protection purpose, where the head-end

or tail-end nodes may be different, this information is necessary in the message exchange between PCCs and PCEs, so that the stateful PCE knows which LSP the path computation request wants to share the resource.

Besides, parameter changes during the resource sharing computation also need to be considered. For example, the bandwidth of the request LSP may be different with the existing LSP, while resource sharing is still preferred by the PCC. PCE should consider the sharing request together with the policy and available resource(s) in the network. Details can be found in Section 3.3.

2.2. Multiple PCEs Use Case

Figure 2 shows a two-layer network example, with each layer managed by a PCE. As Discussed in Section 3 of [RFC5623], there are three models for inter-layer path computation. They are single PCE computation, multiple PCE with inter-PCE communication and multiple PCE without inter-PCE communication, respectively. For the single PCE computation, the process would be similar to that of the use case in Section 2.1. Thus, this model is not discussed further.

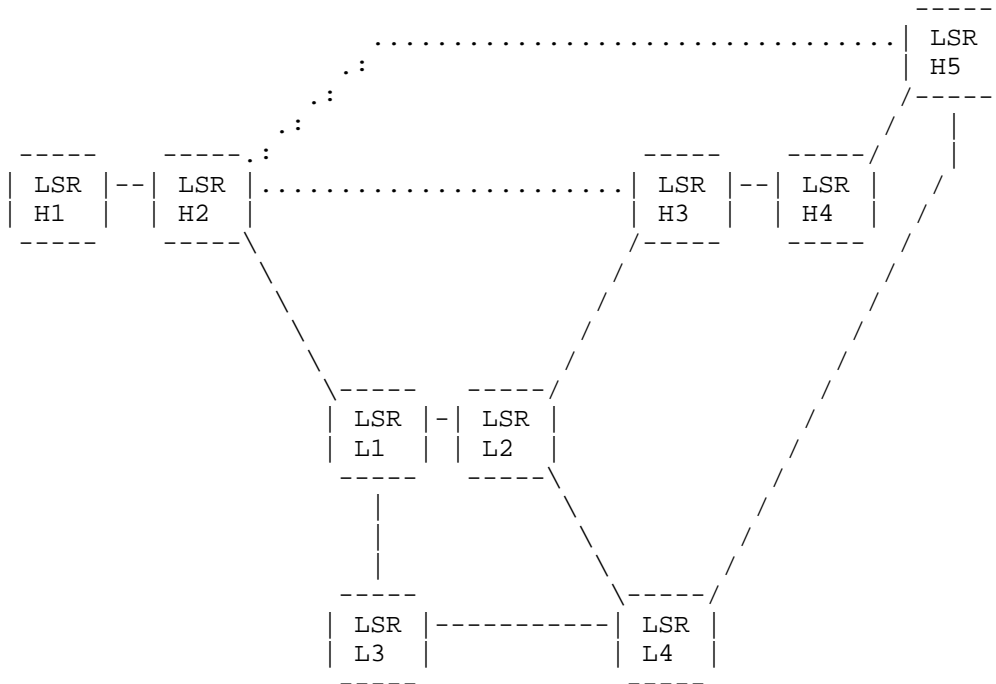


Figure 2: A Two-layer Network Example

An inter-layer path computation example is shown in Fig. 2, assume a LSP (LSP1: H2-H3) has been established already, visible as H2-H3 from view of higher-layer PCE and H2-L1-L2-H3 from the global view (or from the view of lower-layer PCE). A new request comes at H2 to establish a new LSP (LSP2: from H2 to H5), given the constraint it can share resource with LSP1. This requirement is possible if only one of the LSPs needs to be active and resource sharing is the target.

If multiple PCE with inter-PCE communication model is employed, the path computation request sent by H2 to higher-layer PCE will be forwarded to lower-layer PCE since there is no resource readily available in the higher layer. So it leaves the lower-layer PCE to compute a path in the lower layer in order to support the higher layer request. In this case, lower-layer PCE is required to compute a path between H2 and H5 under the constraint that it can share the resource with that of the LSP1. At this moment the lower-layer PCE has the knowledge on the explicit routing that LSP1 go through (H2-L1-L2-H3). So when lower-layer PCE computes the path for LSP2, it can consider the resource used by LSP1 as available with higher priority. For example, lower-layer PCE may choose H2-L1-L2-L4-H5 as the computation result. On the other hand, if the path computation policy is to have a separate path with LSP1, the lower-layer PCE may choose H2-L1-L3-L4-H5.

During this procedure higher-layer PCE can only use LSP1 information (such as its five-tuple LSP information) as the information, an issue to solve is how lower-layer PCE can resolve this information to the actual resource usage in its own layer, i.e. lower layer. This could be solved by edge LSR L1 reporting this higher-lower layer LSP correlation to the lower-layer PCE as part of the LSP information during the LSP state synchronization process. If needed, it can be later updated when there is a change in this information. Alternatively, the lower-layer PCE can get this information from other sources, such as network management system, where this information should be stored.

If multiple PCE without inter-PCE communication model is employed, the path computation request in the lower layer will be initiated the border LSR node, i.e., L1. The process would be similar to that of the previous scenario. A point worth noting is that the border LSR node may be able to resolve the higher layer LSP information itself, such as mapping it to the corresponding LSP in the lower layer, in this way lower-layer PCE does not need to perform this

function. Otherwise, the mapping method mentioned above can still be used.

3. Extensions to PCEP

This section provides PCEP extensions. Currently the text focuses only on passive stateful PCE and corresponding PCReq. But if active stateful PCE delegation is used, we would like to convey the same information via RSO in PCRpt. In the passive stateful PCE architecture, a PCC is allowed to specify resource sharing when sending a PCReq message. It also details the processing rule and error codes needed.

3.1. Resource Sharing Object

The PCEP Resource Sharing Object (RSO) is optional. It MAY be carried within a PCRep message from the network element (or other PCCs) so as to indicate the desired resource sharing requirements to be applied by the stateful PCE during path computation.

The RSO object format is compliant with the PCEP object format defined in [RFC5440].

The RSO Object-Class is TBA.

The RSO Object-type is 1.

The format of the RSO object body is:

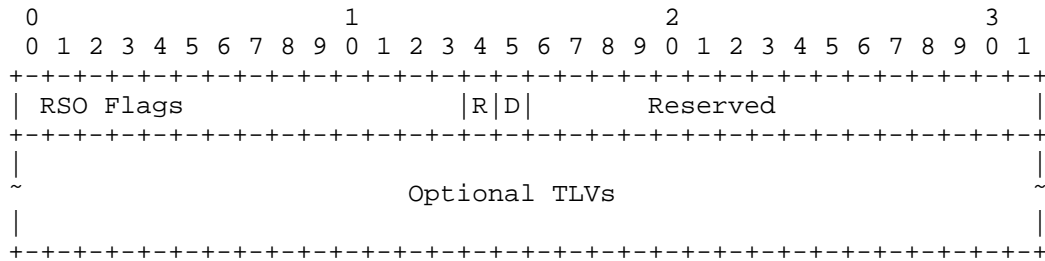


Figure 3: RSO Object Format

RSO flags (16 bits): the objective of the resource sharing. Currently, the following objectives are defined:

D (1 bit): sharing as little as possible.

R (1 bit): sharing as much as possible

It is possible that multiple computation results satisfy the request. Among these results, D set to 1 will select the most separate one, while R set to 1 will select the one sharing most resources. Both D and R set to 0 don't specify any constraint and will result in a random selection among these results. The combination of D=1 and R=1 is not allowed.

Reserved (2 bytes): This field MUST be set to zero on transmission and MUST be ignored on receipt.

Optional TLVs may be needed to indicate the LSP(s) with which the resource is shared. The LSP Info TLV, include the IPv4-LSP-IDENTIFIERS TLV and IPv6-LSP IDENTIFIERS TLV, are defined in the same way as in [stateful-pce].

3.2. Processing Rules

To request a path allowing sharing resource with one or multiple existing LSPs, a PCC includes a RSO object in the PCReq message.

On receipt of a PCReq message with a RSO object, a stateful PCE MUST proceed as follows:

- If the RSO object is unknown/unsupported, the PCE will follow procedures defined in [RFC5440]. That is, the PCE sends a PCErr message with error type 3 or 4 (Unknown / Not supported object) and error value 1 or 2 (unknown / unsupported object class / object type), and the related path computation request is discarded.
- If TLV(s) present in the RSO object are unknown/unsupported and the P bit is set, the PCE MUST send a PCErr message with error type 3 or 4 (Unknown / Not supported object) and error value 4 (Unrecognized/Unsupported parameter), and the related path computation request MUST be discarded as defined in [RFC5440].
- If the resource sharing information is extracted correctly, the PCE MUST apply the requested resource sharing requirement.

The procedure of setting R and/or D bit follows the rules defined in Section 3.1. The RSO flags may be locally configured on the requesting nodes via external entities, such as a network management system or the entity that impose the resource sharing requirement.

3.3. Carrying RSO in a PCEP Message

The RSO is applied to an individual path computation request and the format of the PCReq message is updated as follows:

```
<PCReq Message> ::= <Common Header>
                    [<svec-list>]
                    <request-list>
```

where:

```
<svec-list> ::= <SVEC>
                [<OF>]
                [<metric-list>]
                [<svec-list>]
```

```
<request-list> ::= <request> [<request-list>]
```

```
<request> ::= <RP>
              <END-POINTS>
              [<LSPA>]
              [<BANDWIDTH>]
              [<metric-list>]
              [<OF>]
              [<RRO>[<BANDWIDTH>]]
              [<IRO>]
              [<RSO>]
              [<LOAD-BALANCING>]
```

and where:

```
<metric-list> ::= <METRIC>[<metric-list>]
```

4. Security Considerations

Security of PCEP is discussed in [RFC5440] and [RFC6952]. The extensions in this document do not change the fundamentals of security for PCEP.

However, the introduction of the RSO provides a vector that may be used to probe for information from a network. For example, a PCC that wants to discover the path of an LSP with which it is not involved, can issue a PCReq with an RSO and may be able to get back quite a lot of information about the path of the LSP through issuing multiple such requests for different endpoints and analyzing the received results. To protect against this, a PCE should be configured with access and authorization controls such that only authorized PCCs (for example, those within the network) can make computation requests, only specifically authorized PCCs can make requests using the RSO, and resource sharing requests relating to specific LSPs are further limited to a select few PCCs. How such access controls and authorization is managed is outside the scope of this document, but it will at the least include Access Control Lists.

Furthermore, a PCC must be aware that setting up an LSP that shares resources with another LSP may be a way of attacking the other LSP, for example by depriving it of the resources it needs to operate correctly. Thus it is important that, both in PCEP and the associated signaling protocols, only authorized resource sharing is allowed.

5. IANA Considerations

5.1. New Object Type

IANA manages the PCEP Objects code point registry (see [RFC5440]). This is maintained as the "PCEP Objects" sub-registry of the "Path Computation Element Protocol (PCEP) Numbers" registry.

This document defines a new PCEP object, the RSO object, to be carried in PCReq messages. IANA is requested to make the following allocation in the "PCEP Objects" sub-registry:

Object Class	Name	Object Type	Name	Reference
--------------	------	-------------	------	-----------

 TBA RSO Resource Sharing [this document]

5.2 RSO flags field

IANA is requested to create and maintain a new sub-registry named "RSO flags". The following flags are defined in this document:

Bit	Flag name	Description	Reference
0	D	sharing as little as possible	[this document]
1	R	sharing as much as possible	[this document]

6. References

6.1. Normative References

- [RFC2119] Bradner, S., "Key words for use in RFCs to indicate requirements levels", RFC 2119, March 1997.
- [RFC4655] Farrel, A., Vasseur, J.-P., and Ash, J., "A Path Computation Element (PCE)-Based Architecture", RFC 4655, August 2006.
- [RFC5440] Vasseur, J.-P., and Le Roux, JL., "Path Computation Element (PCE) Communication Protocol (PCEP)", RFC 5440, March 2009.
- [Stateful-PCE] Crabbe, E., Medved, J., Minei, I., and R. Varga, "PCEP Extensions for Stateful PCE", draft-ietf-pce-stateful-pce-10 (work in progress), October 2014.

6.2. Informative References

- [RFC4428] Papadimitriou, D., Mannie., E., "Analysis of Generalized Multi-Protocol Label Switching (GMPLS)-based Recovery Mechanisms (including Protection and Restoration)", RFC4428, March 2006.

[RFC5623] Oki., E., Takeda, T., Le Roux, JL., Farrel, A., "Framework for PCE-Based Inter-Layer MPLS and GMPLS Traffic Engineering", RFC5623, September 2009.

[RFC6952] Jethanandani, M., Patel, K., Zheng, L., "Analysis of BGP, LDP, PCEP, and MSDP Issues According to the Keying and Authentication for Routing Protocols (KARP) Design Guide", RFC6952, May 2013.

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Extensions to the Path Computation Element Protocol (PCEP) to Support
Resource Sharing-based Path Computation

draft-zhang-pce-resource-sharing-12

Abstract

Resource sharing in a network means two or more Label Switched Paths (LSPs) use common pieces of resource along their paths. This can help save network resources and is useful in scenarios such as LSP recovery or when two LSPs do not need to be active at the same time. A Path Computation Element (PCE) is responsible for path computation with such requirement.

Existing extensions to the Path Computation Element Protocol (PCEP) allow one path computation request for an LSP to be associated with other (existing) LSPs through the use of the PCEP Association Object.

This document extends PCEP in order to support resource-sharing-based path computation as another use of the Association Object to enable better efficiency in the computation and in the resultant paths and network resource usage.

Status of this Memo

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1. Introduction and Motivation

A Path Computation Element (PCE) is a way to provide path computation function, and it is especially useful in the scenarios where complex constraints and/or a demanding amount of computation resource are required [RFC4655]. The development of PCE standardization has evolved from stateless to stateful. A stateful PCE has access to the LSP database information of the networks it serves as a computation engine [RFC8231]. Unless specified, this document assumes a PCE mentioned is a stateful PCE.

Resource sharing denotes that two or more Label Switched Paths (LSPs) share common pieces of resource, (such as a common time slot of a link in an Optical Transport Network (OTN)). This is usually useful in the scenario where only one of the LSPs is active and the benefit is to save network resources. A simple example of this is dynamically calculating a recovery LSP for an existing LSP undergoing a link failure. Note that resource sharing can be worked out using a stateless PCE, but the mechanism may be complex and is out the scope of this document.

This document considers the requirement that a new LSP may request for resource sharing with one or multiple existing LSPs. Furthermore, if there is resource sharing between a new LSP and existing an LSP, the two LSPs cannot be used to carry traffic simultaneously, the new LSP will take over the traffic from the existing LSP.

In a single domain, this is a common requirement in the recovery cases especially in order to increase traffic resilience against failure while reducing the amount of network resource used for recovery purposes [RFC4428].

The current protocol supporting the communication between a PCE and a Path Computation Client (PCC), i.e. PCE Protocol (PCEP), allows

for re-optimization of an existing LSP [RFC5440]. This is achieved by setting the R bit in the Request Parameter (RP) object, together with some additional information if applicable, in the Path Computation Request (PCReq) message sent from a PCC to the PCE. To support this type of resource sharing, a PCC needs to ask a PCE to compute a new path with the constraints of sharing resource with one or multiple existing LSPs. It is worth noting the "resource sharing" in this draft not only means one LSP re-using the same links of another LSP, but also the same slice of bandwidth in the network. This may occur when an LSP is required for re-routing, or online re-optimization. Current PCEP specifications do not provide such function. More specifically, this document describes the resource sharing issue during the procedure when a new LSP is required to replace an existing LSP for use together with Make-before-break (MBB) described in [RFC3209].

As mentioned in [RFC8231], the PLSP-ID provides a unique identifier for an LSP during a PCEP session between PCC and PCE. Such identification is helpful in supporting the resource sharing requirement for stateful PCEs because it greatly simplifies the operation of a PCC. Instead of the PCC determining all the resources to be shared, the PCC can request that the PCE share the resources of a specific LSP: the stateful PCE is able to determine those resource itself.

Resource sharing can also be required in an inter-layer PCEP session. This is similar to the previous requirement. However, it is more complex and therefore deserves a more detailed explanation here.

In a multi-layer network, LSPs in a lower layer are used to carry higher-layer LSPs across the lower-layer network [RFC5623]. Therefore, the resource sharing constraints in the higher layer might actually relate to resource sharing in the lower layer. Thus, it is useful to consider how this can be achieved and whether additional extensions are needed using the models defined in [RFC5623].

In the next sections, use cases are provided to show what information needs to be exchanged to fulfill these requirements. This memo then provides extensions to PCEP to enable this function.

1.1. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in

BCP 14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

2. Motivation

2.1. Single Domain Use Case

There are two potential cases that request resource to be shared: restoration and re-optimization. Figure 1 shows a single domain network with a stateful PCE, and is used as an example for the resource sharing application.

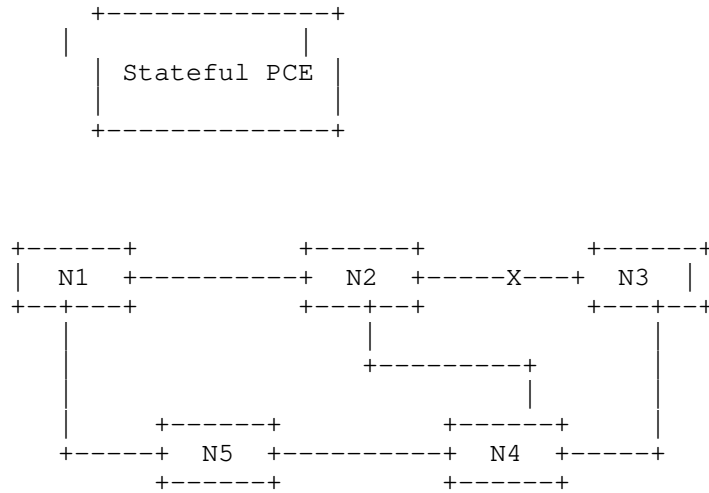


Figure 1: A Single Domain Example

- LSP0 (existing): N1-N2-N3
- LSP1 (restoration): N1-N2-N4-N3
- LSP2 (re-optimization): N1-N5-N4-N3

For the failure restoration, we can assume a working LSP (LSP0) exists in the network. When there is failure on the link N2-N3, it is desired to set up a restoration path for this working LSP. Suppose N1 serves as the PCC and sends a request to the stateful PCE for such an LSP. Before sending the request, N1 may need to check what policy should be applied for the restoration. For example, it might value resource sharing and prefer to share as much resource with the working LSP as possible and specify this policy in the PCReq message. Given such policy, a probable outcome from the path

computation would be LSP1, which shares the link 'N1-N2' with the existing LSP.

Re-optimization does not usually result from a specific failure in the network, but takes place on a stable network when more optimal paths may have become available. Thus switching from the existing LSP to the new LSP happens with live traffic. An example can be found in Figure 1 without failure on the link N2-N3. Instead, an online re-optimization is needed for the working LSP (LSP0) from the stateful PCE. In such cases, the best choice is to set up a backup LSP for the working LSP with totally separate routing (for example, LSP2), and move the traffic to that backup LSP. After that, the working LSP can be torn down, which will not result in any interruption during the optimization procedure. This can actually be implemented with existing PCEP mechanisms. However, if there is no such separate path, existing PCEP mechanisms will return an error. A secondary option for this case is to set up an LSP and complete re-optimization with resource sharing, even if some interruption is introduced.

In the example from Figure 1 it is assumed that the restored LSP or re-optimized LSP have the same source and destination nodes. But in some applications there is no restriction for this assumption, i.e., after an LSP is failed, it can be restored as a new LSP with different source/destination.

In the use cases above it is also assumed that the characteristics of the restored LSP or re-optimized LSP are unchanged. However, it is possible to have parameter changes during the resource sharing computation. For example, the bandwidth of the request LSP may be different from the existing LSP, while resource sharing is still preferred by the PCC. The PCE should consider the sharing request together with the policy and available resources in the network. Details can be found in Section 3.3.

Conversely to resource sharing, it may also be required to apply a disjoint constraint for the path computation. [ietf-pce-association-diversity] discusses the solution under such a scenario, which is a companion work to this document.

2.2. Multiple Layers/Domains Use Case

As Discussed in Section 3 of [RFC5623], there are three models for inter-layer path computation. They are single PCE computation, multiple PCE with inter-PCE communication, and multiple PCE without inter-PCE communication. For the single PCE computation, the process would be similar to that of the use case in Section 2.1.

An inter-layer path computation example is shown in Figure 2. Assume an LSP (LSP1: H2-H3) has been established already, visible as H2-H3 from the view of the higher-layer PCE, and as H2-L1-L2-H3 from the global view (or from the view of the lower-layer PCE). A new request is received by H2 to establish a new LSP (LSP2: from H2 to H5), given the constraint that it can share resources with LSP1. This requirement is possible if only one of the LSPs needs to be active and resource sharing is the target.

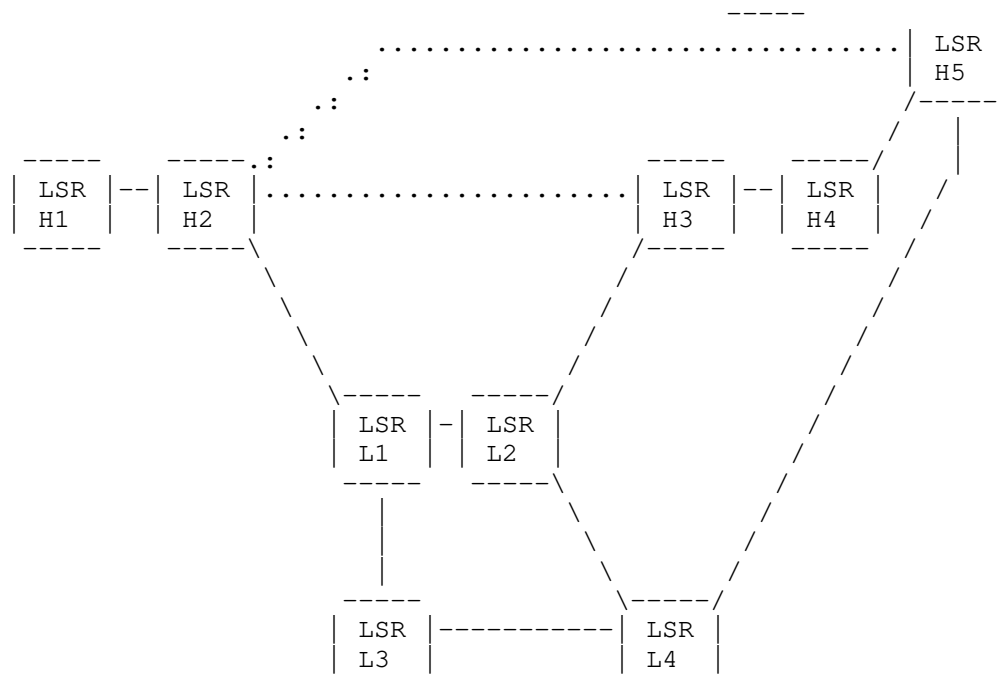


Figure 2: A Two-layer Network Example

If the model of multiple PCEs with inter-PCE communication is employed, the path computation request sent by H2 to higher-layer PCE will be forwarded to lower-layer PCE since there is no resource readily available in the higher layer. So it leaves the lower-layer PCE to compute a path in the lower layer in order to support the higher layer request. In this case, the lower-layer PCE is required to compute a path between H2 and H5 under the constraint that it can share the resource with that of LSP1. At this moment the lower-layer PCE has knowledge of the explicit route of LSP1 (H2-L1-L2-H3), and therefore can map the lower layer LSP with the higher-layer one. So when the lower-layer PCE computes the path for LSP2, it can consider

the resource used by LSP1 as available with higher priority. For example, the lower-layer PCE may choose H2-L1-L2-L4-H5 as the computation result. On the other hand, if the path computation policy is to have a separate path with LSP1, the lower-layer PCE may choose H2-L1-L3-L4-H5.

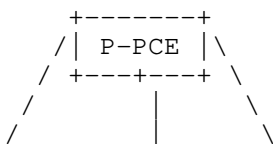
During this procedure the higher-layer PCE can only use information about LSP1 (such as its five-tuple LSP information). An issue to solve is how the lower-layer PCE can resolve this information to the actual resource usage in its own layer, i.e. the lower layer. This could be solved by the edge LSR (L1) reporting this higher-lower LSP correlation to the lower-layer PCE as part of the LSP information during the LSP state synchronization process. If needed, it can be updated later when there is a change in this information. Alternatively, the lower-layer PCE can get this information from other sources, such as a network management system, where this information should be stored.

If the model of multiple PCEs without inter-PCE communication is employed, the path computation request in the lower layer will be initiated by the border LSR node, i.e., L1. The process would be similar to that of the previous scenario. A point worth noting is that the border LSR node may be able to resolve the higher layer LSP information itself, such as by mapping it to the corresponding LSP in the lower layer, in this way the lower-layer PCE does not need to perform this function. Otherwise, the mapping method mentioned above can still be used.

2.3. Bulk Path Computation Use Case

There is a potential need for resource sharing during bulk path computation, especially the processing of the "sticky resources" in [RFC7399]. It would be useful to specify the resources that can be shared among different paths, i.e., the bandwidth information.

Considering the H-PCE architecture in [RFC8751], when the parent PCE asks for a single path across a few domains, such a request may become a bulk path computation to a certain child PCE. Figure 3 shows an example of 3 domains. The parent PCE will select one of these path for establishment.



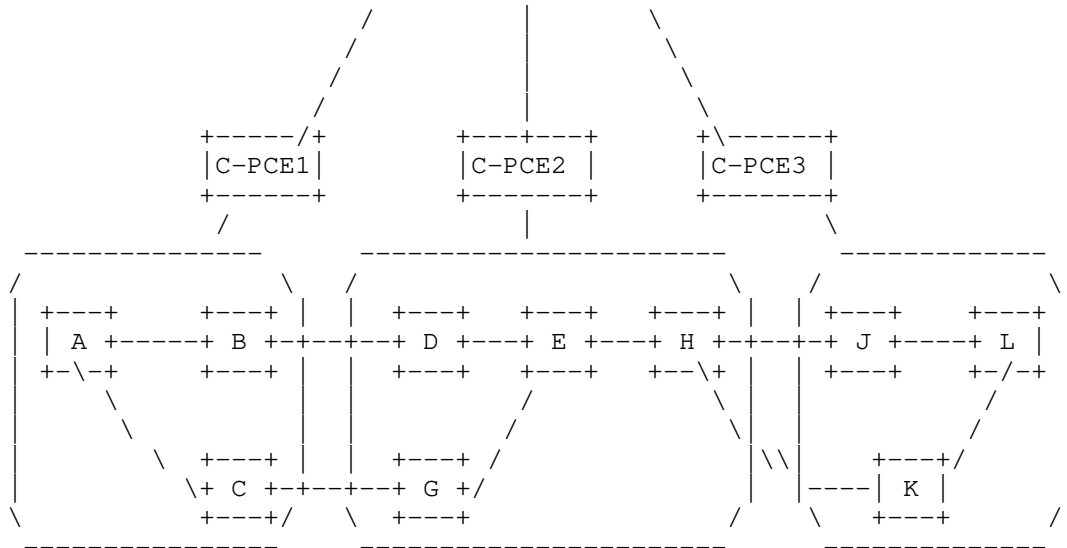


Figure 3: Bulk Request example with Hierarchical PCEs

A 3-domain example is shown in Figure 3, with the hierarchical PCE architecture. In this example nodes A/B/C belong to domain 1, nodes D/E/G/H belong to domain 2, and nodes J/K/L belong to domain 3. Inter-domain links are B-D/C-G between domains 1 and 2, and H-J/H-K between domains 2 and 3. Given a path computation request from A to L, a bulk request from P-PCE would be helpful to understand whether it is possible to have different combinations on the inter-domain links. However, the resources on some specific links become 'sticky' and have to be indicated as 'sharing allowed' to avoid unnecessary resource competition. For example, both the route A-B-D-E-H-J-L and A-C-G-E-H-K-L are qualified, but these routes are competing for the resource on the link E-H and cannot be established simultaneously, so there must be one route failed to be reported to P-PCE. Given the indication of allowing resource sharing on the link E-H, both of these routes can be reported for P-PCE's decision, and there will not be any competition as the P-PCE understands that only one path needs to be set up.

3. Extensions to PCEP

3.1. Association Group and Type

According to the definition in [RFC8697], the association group is used to associate multiple LSPs into one group for further path computation considerations, such as disjointness and resource sharing. An association ID will be used to identify the resource

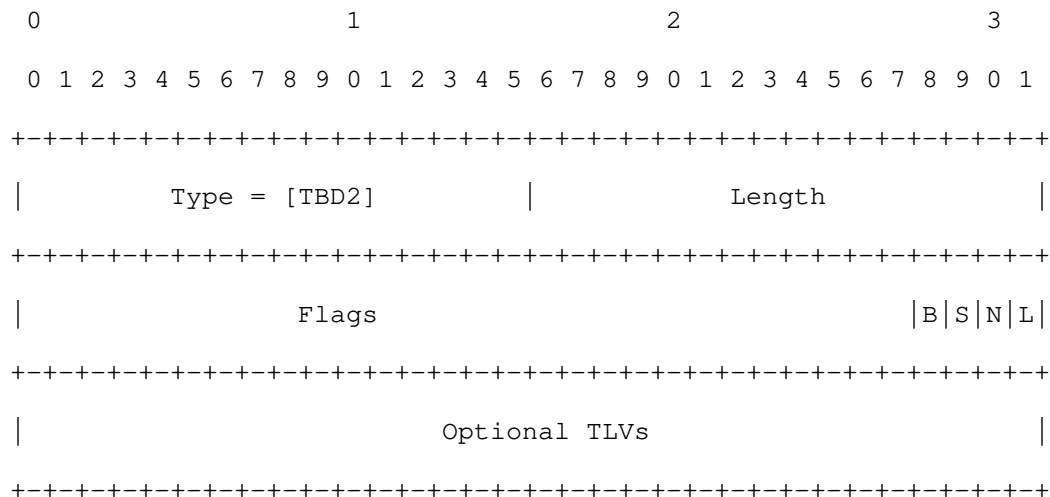
sharing group. An association type that described disjointness has been defined in [ietf-pce-association-diversity]. In this document, a new association type is defined as follows:

Association type = TBD1 ("Sharing Association Type").

A sharing group should have multiple LSPs. The number of LSPs and the criteria for how LSPs share among each other are dependent on local policy.

3.2. Resource Sharing TLV

The PCEP Resource Sharing group MAY carry the following TLV. It MAY be carried within a PCReq message from the network element (or other PCCs) so as to indicate the desired resource sharing requirements to be applied by the stateful PCE during path computation.



The following flags are defined:

- * L (Link share) bit: when set, this flag indicates that the PCE should prioritize the links shared by existing LSPs within the sharing group for path computation.
- * N (Node share) bit: when set, this flag indicates that the PCE should prioritize the nodes shared by existing LSPs within the sharing group for path computation.

* S (SRLG share) bit: when set, this flag indicates that the PCE should set the SRLG (Shared Risk Link Group) of the computed LSP to the same as existing LSPs within the sharing group for path computation.

* B (Bandwidth share) bit: when set, this flag indicates that the PCE should prioritize the bandwidth to be shared by LSPs within the sharing group for bulk path computation.

It is worth noting that there can be multiple flags set which may conflict with each. In this scenario, the result for path computation will be dependent on the policy of PCE.

Optional TLVs may be needed to indicate the LSPs with which the resource is shared. If multiple LSPs are required, the PCE may need to consider different sharing policies, which is implementation dependent and may result in a different computing result. The selection policy among multiple computation result is out of the scope of this document.

3.3. Processing Rules

To request a path allowing resource sharing with one or multiple existing LSPs, a PCC includes a Resource Sharing TLV in the Association Group Object in any kind of path computation request message, such as the PCReq, PCUpd, or PCInitiate messages specified in [RFC8231] and [RFC8281].

On receipt of a PCEP message with a Resource Sharing TLV, a stateful PCE MUST proceed as follows:

- If the Resource Sharing TLV is unknown/unsupported, the PCE will follow procedures defined in [RFC5440]. That is, the PCE sends a PCErr message with error type 26 (Association Error) and error value 6 (Association Information Mismatch), and the related path computation request is discarded.
- If the Resource Sharing TLV is extracted correctly, the PCE MUST apply the requested resource sharing requirement.

The procedure of setting flags follows the rules defined in Section 3.1. The flags in the Resource Sharing TLV may be locally configured on the requesting nodes via external entities, such as a network management system or the entity that imposes the resource sharing requirement.

It is worth noting that the Resource Sharing TLV can be used together with other path indication objects like the IRO/XRO, with different objectives. The first difference is, the use of the Resource Sharing TLV is to set up an alternative path, instead a new path. It is also dependent on the knowledge held by the PCC, e.g., if the PCC has full knowledge of the path information and has a strong preference on the route, it may send the request message with an IRO to specify the route. On the other hand, if the PCC does not know how the path should go but just wants to set up a new LSP to replace the old one, it may use the Resource Sharing TLV instead of an IRO. The second difference is that the Resource Sharing TLV is a loose requirement. For example, if the constraint specified in an IRO/XRO in an A-Z path computation request cannot be satisfied, the reply message from PCE to PCC would be unsuccessful. However it is still possible to have a path from the A-Z. If the target node/link/SRLG/Bandwidth is set in the Resource Sharing TLV rather than an IRO, the PCE may feedback a path from A-Z that does not share the target specified in the Resource Sharing TLV.

4. Implementation Status

[Note to the RFC Editor - remove this section before publication, as well as remove the reference to [RFC7942].

Currently the authors are not aware of any implementations.

5. Manageability Considerations

All manageability requirements and considerations listed in [RFC5440] and [RFC8231] apply to the PCEP protocol extensions defined in this document. In addition, requirements and considerations listed in this section apply.

5.1. Control of Function and Policy

A PCE or PCC implementation MUST allow operator-configured associations and SHOULD allow setting of the resource sharing TLV (Section 3.4) as described in this document.

5.2. Information and Data Models

An implementation SHOULD allow the operator to view the resource sharing configured or created dynamically. Further implementation SHOULD allow to view resource sharing associations reported by each peer, and the current set of LSPs in the association. The PCEP YANG module [ietf-pce-pcep-yang] includes association groups information.

5.3. Liveness Detection and Monitoring

Mechanisms defined in this document do not imply any new liveness detection and monitoring requirements in addition to those already listed in [RFC5440].

5.4. Verify Correct Operations

Mechanisms defined in this document do not imply any new operation verification requirements in addition to those already listed in [RFC5440] and [RFC8231].

5.5. Requirements on Other Protocols

Mechanisms defined in this document do not imply any new requirements on other protocols. The configuration on local policy may be accomplished by other protocols, such as Netconf.

5.6. Impact on Network Operations

Mechanisms defined in [RFC5440] and [RFC8231] also apply to PCEP extensions defined in this document.

6. Security Considerations

Security of PCEP is discussed in [RFC5440] and [RFC6952]. The extensions in this document do not change the fundamentals of security for PCEP.

However, the introduction of the Resource Sharing TLV in the Association Group Object provides a vector that may be used to probe for information from a network. For example, a PCC that wants to discover the path of an LSP with which it is not involved can issue a request message with a Resource Sharing TLV and may be able to get back quite a lot of information about the path of the LSP through issuing multiple such requests for different endpoints and analyzing the received results. To protect against this, a PCE SHOULD be configured with access and authorization controls such that only authorized PCCs (for example, those within the network) can make computation requests, only specifically authorized PCCs can make requests for resource sharing, and such requests relating to specific LSPs are further limited to a select few PCCs. How such access controls and authorization is managed is outside the scope of this document, but it will at the least include Access Control Lists.

Furthermore, a PCC must be aware that setting up an LSP that shares resources with another LSP may be a way of attacking the other LSP, for example by depriving it of the resources it needs to operate correctly. Thus it is important that, both in PCEP and the associated signaling protocols, only authorized resource sharing is allowed.

7. IANA Considerations

7.1. Association Object Type Indicators

IANA maintains a registry called the "Path Computation Element Protocol (PCEP) Numbers" registry with a subregistry called the "Association Type Field" subregistry. IANA is requested to make an assignment from that subregistry as follows:

Object Class	Name	Object Type	Reference
TBD1	Sharing-group	Association Type	[this document]

7.2. PCEP TLV Definitions

This document defines the following TLVs to support the resource sharing scenario:

Value	Name	Reference
TBD2	Resource-sharing TLV	[this document]

IANA is requested to allocate the following bit numbers in the flag spaces of Resource-sharing TLV:

Bit	Flag name	Reference
31	Link Share	[this document]
30	Node Share	[this document]
29	SRLG Share	[this document]
28	Bandwidth Share	[this document]

8. References

8.1. Normative References

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PCEP Procedures and Protocol Extensions for Using PCE as a Central
Controller (PCECC) of LSPs
draft-zhao-pce-pcep-extension-for-pce-controller-01

Abstract

In certain networks deployment scenarios, service providers would like to keep all the existing MPLS functionalities in both MPLS and GMPLS while removing the complexity of existing signaling protocols such as LDP and RSVP-TE. In [I-D.zhao-pce-central-controller-user-cases], we propose to use the PCE [RFC5440] as a central controller (PCECC) so that LSP can be calculated/ signaled/initiated and label forwarding entries are downloaded through a centralized PCE server to each network devices along the LSP path while leveraging the existing PCE technologies as much as possible.

This draft specify the procedures and PCEP protocol extensions for using the PCE as the central controller and user cases where LSPs are calculated/setup/initiated and label forwarding entries are downloaded through extending the existing PCE architectures and PCEP.

This document also discuss the role of PCECC in Segment Routing (SR).

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1. Introduction

In certain network deployment scenarios, service providers would like to have the ability to dynamically adapt to a wide range of customer's requests for the sake of flexible network service delivery, Software Defined Networks(SDN) has provides additional flexibility in how the network is operated compared to the traditional network.

The existing networking ecosystem has become awfully complex and highly demanding in terms of robustness, performance, scalability, flexibility, agility, etc. By migrating to the SDN enabled network from the existing network, service providers and network operators must have a solution which they can evolve easily from the existing network into the SDN enabled network while keeping the network services remain scalable, guarantee robustness and availability etc.

Taking the smooth transition between traditional network and the new SDN enabled network into account, especially from a cost impact assessment perspective, using the existing PCE components from the current network to function as the central controller of the SDN network is one choice, which not only achieves the goal of having a centralized controller, but also leverages the existing PCE network components.

The Path Computation Element communication Protocol (PCEP) provides mechanisms for Path Computation Elements (PCEs) to perform route computations in response to Path Computation Clients (PCCs) requests. PCEP Extensions for PCE-initiated LSP Setup in a Stateful PCE Model [I-D.ietf-pce-stateful-pce] describes a set of extensions to PCEP to enable active control of MPLS-TE and GMPLS tunnels.

[I-D.ietf-pce-pce-initiated-lsp] describes the setup and teardown of PCE-initiated LSPs under the active stateful PCE model, without the need for local configuration on the PCC, thus allowing for a dynamic MPLS network that is centrally controlled and deployed.

[I-D.ietf-pce-remote-initiated-gmpls-lsp] complements [I-D.ietf-pce-pce-initiated-lsp] by addressing the requirements for remote-initiated GMPLS LSPs.

Segment Routing (SR) technology leverages the source routing and tunneling paradigms. A source node can choose a path without relying on hop-by-hop signaling protocols such as LDP or RSVP-TE. Each path is specified as a set of "segments" advertised by link-state routing protocols (IS-IS or OSPF). [I-D.ietf-spring-segment-routing] provides an introduction to SR technology. The corresponding IS-IS and OSPF extensions are specified in [I-D.ietf-isis-segment-routing-extensions] and [I-D.ietf-ospf-segment-routing-extensions], respectively.

A Segment Routed path (SR path) can be derived from an IGP Shortest Path Tree (SPT). Segment Routed Traffic Engineering paths (SR-TE paths) may not follow IGP SPT. Such paths may be chosen by a suitable network planning tool and provisioned on the source node of the SR-TE path.

It is possible to use a stateful PCE for computing one or more SR-TE paths taking into account various constraints and objective functions. Once a path is chosen, the stateful PCE can instantiate an SR-TE path on a PCC using PCEP extensions specified in [I-D.ietf-pce-pce-initiated-lsp] using the SR specific PCEP extensions described in [I-D.ietf-pce-segment-routing].

PCECC may further use PCEP protocol for SR label distribution instead of IGP extensions with some benefits.

Current MPLS label has local meaning. That is, MPLS label is always allocated by the downstream node to the upstream node. Then the MPLS label is only identified by the neighboring upstream node and downstream node. The label allocation is done locally and signaled through the LDP/RSVP-TE/BGP protocol. To ease the label allocation and signaling mechanism, PCE can be conveniently used as a central

controller with Label download capability. Further PCE can also be used to manage the label range and SRGB etc.

The PCECC solution introduced in [I-D.zhao-pce-central-controller-user-cases] allow for a dynamic MPLS network that is eventually controlled and deployed without the deployment of RSVP-TE protocol or extended IGP protocol with node/adjacency segment identifiers signaling capability while providing all the key MPLS functionalities needed by the service providers.

This draft specify the procedures and PCEP protocol extensions for using the PCE as the central controller and user cases where LSPs are calculated/setup/initiated/downloaded through extending the existing PCE architectures and PCEP.

1.1. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

2. Terminology

The following terminology is used in this document.

IGP: Interior Gateway Protocol. Either of the two routing protocols, Open Shortest Path First (OSPF) or Intermediate System to Intermediate System (IS-IS).

PCC: Path Computation Client: any client application requesting a path computation to be performed by a Path Computation Element.

PCE: Path Computation Element. An entity (component, application, or network node) that is capable of computing a network path or route based on a network graph and applying computational constraints.

TE: Traffic Engineering.

3. PCECC Modes

The following PCECC modes are supported -

- o Basic PCECC.
- o PCECC SR.
 - * PCECC SR-BE (Best Effort).

- * PCECC SR-TE (Traffic Engineered).

In basic PCECC mode, the forwarding is similar to RSVP-TE signalled LSP without the RSVP-TE signaling. The PCECC allocates and download the label entries along the LSP. The rest of processing is similar to the existing stateful PCE mechanism.

In case of SR, there are two modes for SR-BE and SR-TE. For SR-BE, the forwarding is similar to LDP LSP without LDP signaling or IGP-SR extension. The SR Node label are allocated and distributed in the domain centrally by the PCE via PCEP. Each node (PCC) rely on local IGP for the nexthop calculation. For SR-TE, the forwarding uses label stack similar to IGP based SR-TE without IGP-SR extension. The SR node and adj labels are allocated and distributed in the domain centrally by the PCE via PCEP by PCECC. Rest of the processing is similar to existing stateful PCE with SR mechanism.

4. PCEP Requirements

Following key requirements associated PCECC should be considered when designing the PCECC based solution:

1. PCEP speaker supporting this draft MUST have the capability to advertise its PCECC capability to its peers.
2. Path Computation Client (PCC) supporting this draft MUST have a capability to communicate local label range or global label range or both to PCE.
3. Path Computation Element (PCE) supporting this draft SHOULD have the capability to negotiate a global label range for a group of clients and communicate the final global label range to PCC.
4. PCEP speaker not supporting this draft MUST be able to reject PCECC related message with a reason code that indicates no support for PCECC.
5. PCEP SHOULD provide a means to identify PCECC based LSP in the PCEP messages.
6. PCEP SHOULD provide a means to update (or cleanup) the label-download or label-map entry to the PCC.

5. Procedures for Using the PCE as the Central Controller (PCECC)

5.1. Stateful PCE Model

Active stateful PCE is described in [I-D.ietf-pce-stateful-pce]. PCE as a central controller (PCECC) reuses existing Active stateful PCE mechanism as much as possible to control the LSP.

5.2. New LSP Functions

This document defines the following new PCEP messages and extends the existing messages to support PCECC:

(PCLRRResv): a PCEP message sent by a PCC to a PCE to ask for the label range reservation or a PCE to a PCC to send the reserved label range. The PCLRRResv message described in Section 6.1.

(PCLLabelUpd): a PCEP message sent by a PCE to a PCC to download or cleanup the Label entry. The PCLLabelUpd message described in Section 6.2.

(PCRpt): a PCEP message described in [I-D.ietf-pce-stateful-pce]. PCRpt message MAYBE used to send PCECC LSP Reports.

(PCInitiate): a PCEP message described in [I-D.ietf-pce-pce-initiated-lsp]. PCInitiate message is used to setup PCE-Initiated LSP based on PCECC mechanism.

(PCUpd): a PCEP message described in [I-D.ietf-pce-stateful-pce]. PCUpd message is used to send PCECC LSP Update.

The new LSP functions defined in this document are mapped onto the messages as shown in the following table.

Function	Message
PCECC Capability advertisement	Open
Label Range Reservation	PCLRRResv
Label entry Update	PCLLabelUpd
Label entry Cleanup	PCLLabelUpd
PCECC Initiated LSP	PCInitiate
PCECC LSP Update	PCUpd
PCECC LSP State Report	PCRpt
PCECC LSP Delegation	PCRpt

5.3. PCECC Capability Advertisement

During PCEP Initialization Phase, PCEP Speakers (PCE or PCC) advertise their support of PCECC extensions. A PCEP Speaker includes the "PCECC Capability" TLV, described in Section 7.1.1 of this document, in the OPEN Object to advertise its support for PCECC extensions.

The presence of the PCECC Capability TLV in PCC's OPEN Object indicates that the PCC is willing to function as a PCECC client.

The presence of the PCECC Capability TLV in PCE's OPEN message indicates that the PCE is interested in function as a PCECC server.

The PCEP protocol extensions for PCECC MUST NOT be used if one or both PCEP Speakers have not included the PCECC Capability TLV in their respective OPEN message. If the PCEP Speakers support the extensions of this draft but did not advertise this capability then a PCErr message with Error-Type=19(Invalid Operation) and Error-Value=[TBD] (Attempted LSP setup/download/label-range reservation if PCECC capability was not advertised) will be generated and the PCEP session will be terminated.

L flag and G flag defined in PCECC Capability TLV specifies the local and global label range reservation capability.

A PCC or a PCE MUST include both PCECC-CAPABILITY TLV and STATEFUL-PCE-CAPABILITY TLV in OPEN Object to support the extensions defined in this document. If PCECC-CAPABILITY TLV is advertised and STATEFUL-PCE-CAPABILITY TLV is not advertised in OPEN Object, it SHOULD send a PCErr message with Error-Type=19 (Invalid Operation) and Error-value=[TBD](stateful pce capability was not advertised) and terminate the session.

5.4. Label Range Reservation

After PCEP initial state synchronization, the label range is reserved.

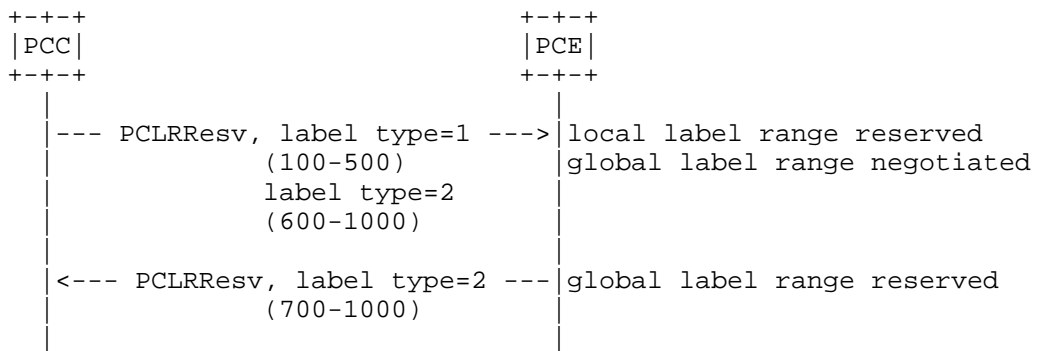
If L flag is advertised in OPEN Object by PCEP speakers, a PCC reserves a local label range and is communicated using PCLRRResv message to a PCE. The PCE maintains the local label range of each node and further during LSP setup, a label is assigned to each node from the corresponding local label range reserved.

If G flag is advertised in OPEN Object by PCEP speakers, a PCC reserves a global label range and is advertised in PCLRRResv message to a PCE. The PCE MAY negotiate and reserves the global label range

and also sends the negotiated global label range in PCLRRResv message to the PCC. Please refer [I-D.li-mpls-global-label-framework] for MPLS global label allocation.

A PCC MUST send PCLRRResv message immediately after the initial LSP synchronization completion. A PCE SHOULD not send PCLLabelUpd message to a PCC before PCLRRResv message received. If the PCC received PCLLabelUpd message and not initiated label range reservation, it SHOULD send a PCErr message with Error-type=[TBD] (label range not reserved) and Error-value=[TBD].

The label range reservation sequence is shown below.



[Editor’s Note: This section of the document would be updated with more details about Label Block Negotiation, Reservation, Adjustment etc in a future revision of the document.]

5.5. LSP Operations

The PCEP messages pertaining to PCECC MUST include PATH-SETUP-TYPE TLV [I-D.sivabalan-pce-lsp-setup-type] in the SRP object to clearly identify the PCECC LSP is intended.

5.5.1. Basic PCECC Mode

5.5.1.1. PCECC LSP Setup

Inorder to setup a LSP based on PCECC mechanism, a PCC MUST delegate the LSP by sending a PCRpt message with Path Setup Type set for basic PCECC (see Section 7.3) and D (Delegate) flag (see [I-D.ietf-pce-stateful-pce]) set in the LSP object.

The LSP-ID in LSP-IDENTIFIER TLV (which usually corresponds to the RSVP-TE LSP-ID) for PCECC LSP MUST always be generated by the PCE. In the first PCRpt message of PCECC LSP, LSP ID of LSP-IDENTIFIER TLV is set to zero.

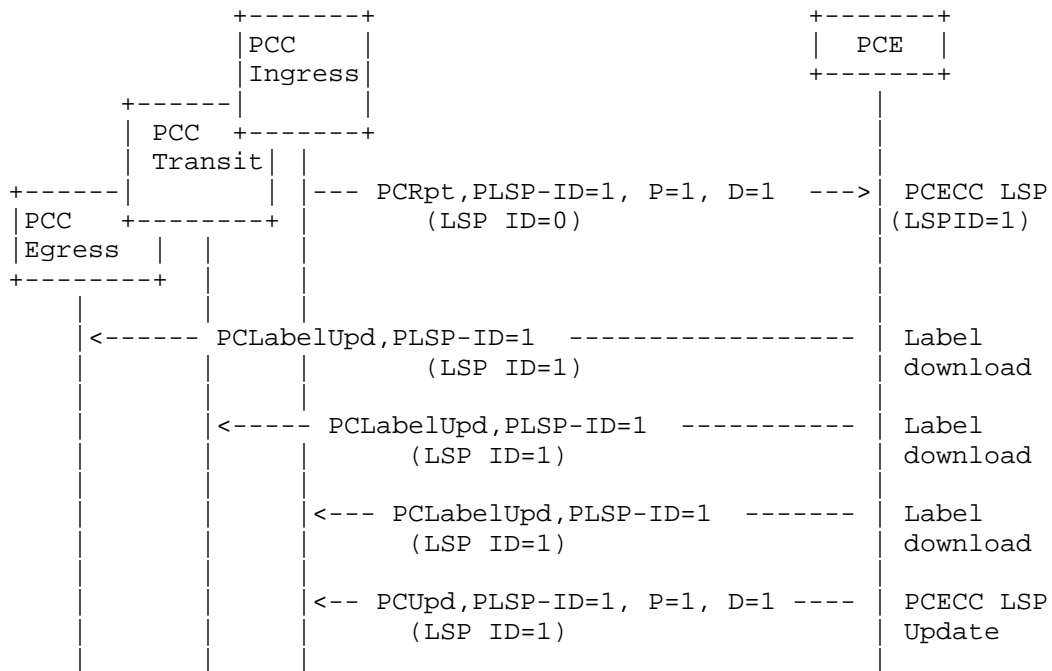
When a PCE received PCRpt message with P and D flags set, it generates LSP ID; calculates the path and assign labels along the path; and setup the path by sending PCLabelUpd message to each node along the path of the LSP.

The PCE SHOULD send the PCUpd message with the same PLSP-ID to the Ingress PCC in response to the delegate PCRpt message.

The PCECC LSPs MUST be delegated to a PCE at all times.

LSP deletion operation for PCECC LSP is same as defined in [I-D.ietf-pce-stateful-pce]. If the PCE received PCRpt message for LSP deletion then it does Label cleanup operation as described in Section 5.5.1.3 for the corresponding LSP.

The Basic PCECC LSP setup sequence is as shown below.



The PCECC LSP are considered to be 'up' by default. The Ingress MAY further choose to deploy a data plane check mechanism and report the status back to the PCE via PCRpt message.

5.5.1.2. Label Download

Inorder to setup an LSP based on PCECC, the PCE sends a PCLabelUpd message to each node of the LSP to download the Label entry as described in Section 5.5.1.1.

The LSP object in PCLabelUpd MUST include the LSP-IDENTIFIER TLV.

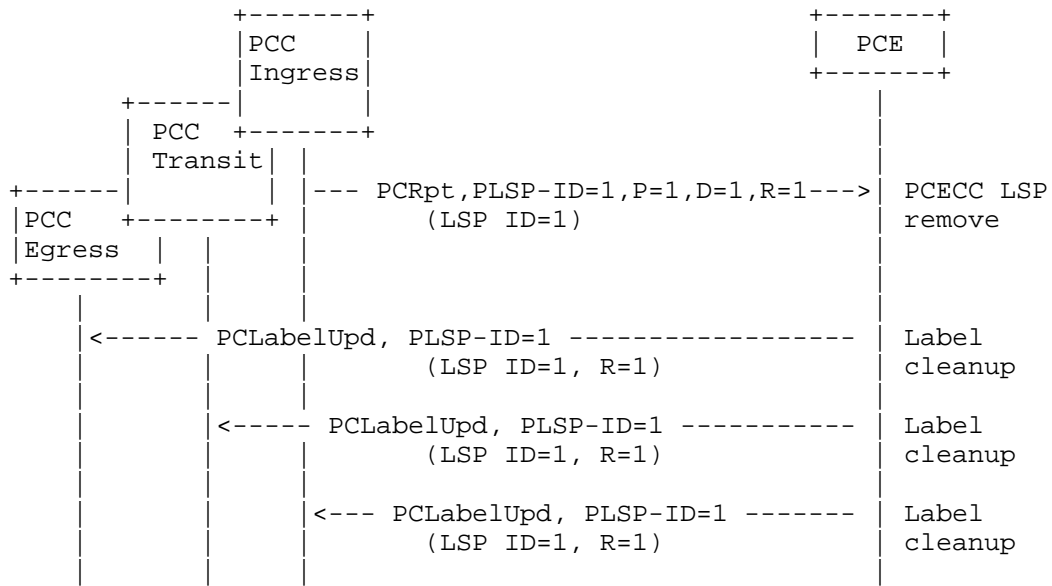
If a node (PCC) received a PCLabelUpd message but failed to download the Label entry, it MUST send a PCErr message with Error-type=[TBD] (label download failed) and Error-value=[TBD].

5.5.1.3. Label Cleanup

Inorder to delete an LSP based on PCECC, the PCE sends a PCLabelUpd message to each node along the path of the LSP to cleanup the Label entry.

If the PCC received a PCLabelUpd message but does not recognize the label, the PCC MUST generate a PCErr message with Error-Type 19(Invalid operation) and Error-Value=3, "Unknown Label".

The R flag in SRP object defined in [I-D.ietf-pce-pce-initiated-lsp] specifies the deletion of Label Entry in the PCLabelUpd message.



5.5.1.4. PCE Initiated PCECC LSP

The LSP Instantiation operation is same as defined in [I-D.ietf-pce-pce-initiated-lsp].

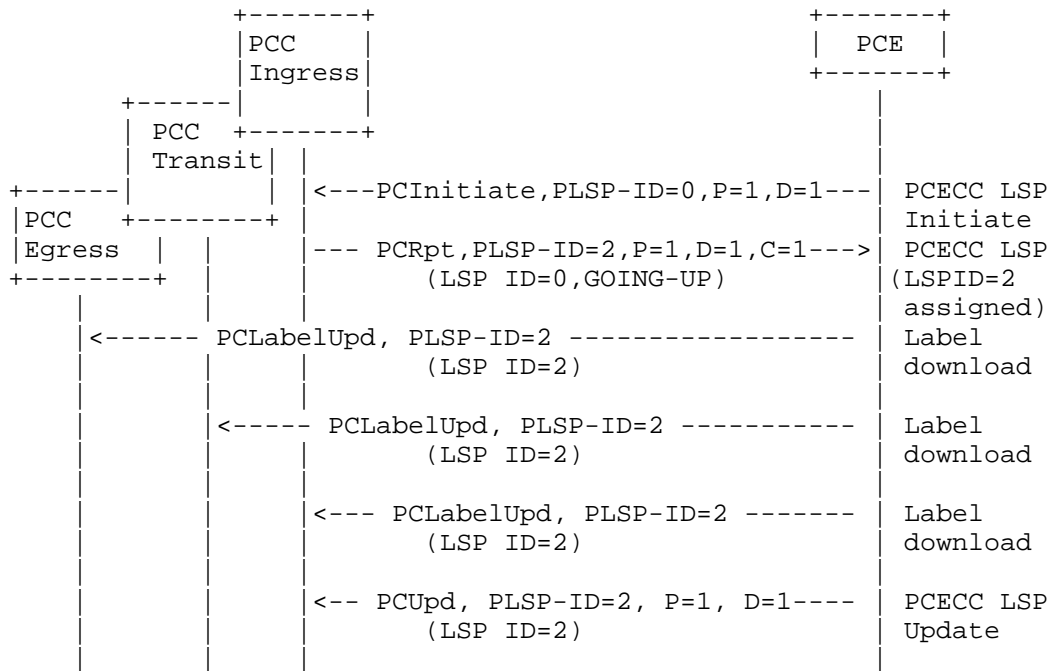
Inorder to setup a PCE Initiated LSP based on PCECC mechanism, a PCE sends PCInitiate message with Path Setup Type set for basic PCECC (see Section 7.3) to the Ingress PCC.

The Ingress PCC MUST also set D (Delegate) flag (see [I-D.ietf-pce-stateful-pce]) and C (Create) flag (see [I-D.ietf-pce-pce-initiated-lsp]) in LSP object of PCRpt message. The PCC responds with first PCRpt message with the status as "GOING-UP" and assigned PLSP-ID.

The rest of the PCECC LSP setup operations are same as those described in Section 5.5.1.1.

The LSP deletion operation for PCE Initiated PCECC LSP is same as defined in [I-D.ietf-pce-pce-initiated-lsp]. The PCE should further perform Label entry cleanup operation as described in Section 5.5.1.3 for the corresponding LSP.

The PCE Initiated PCECC LSP setup sequence is shown below.

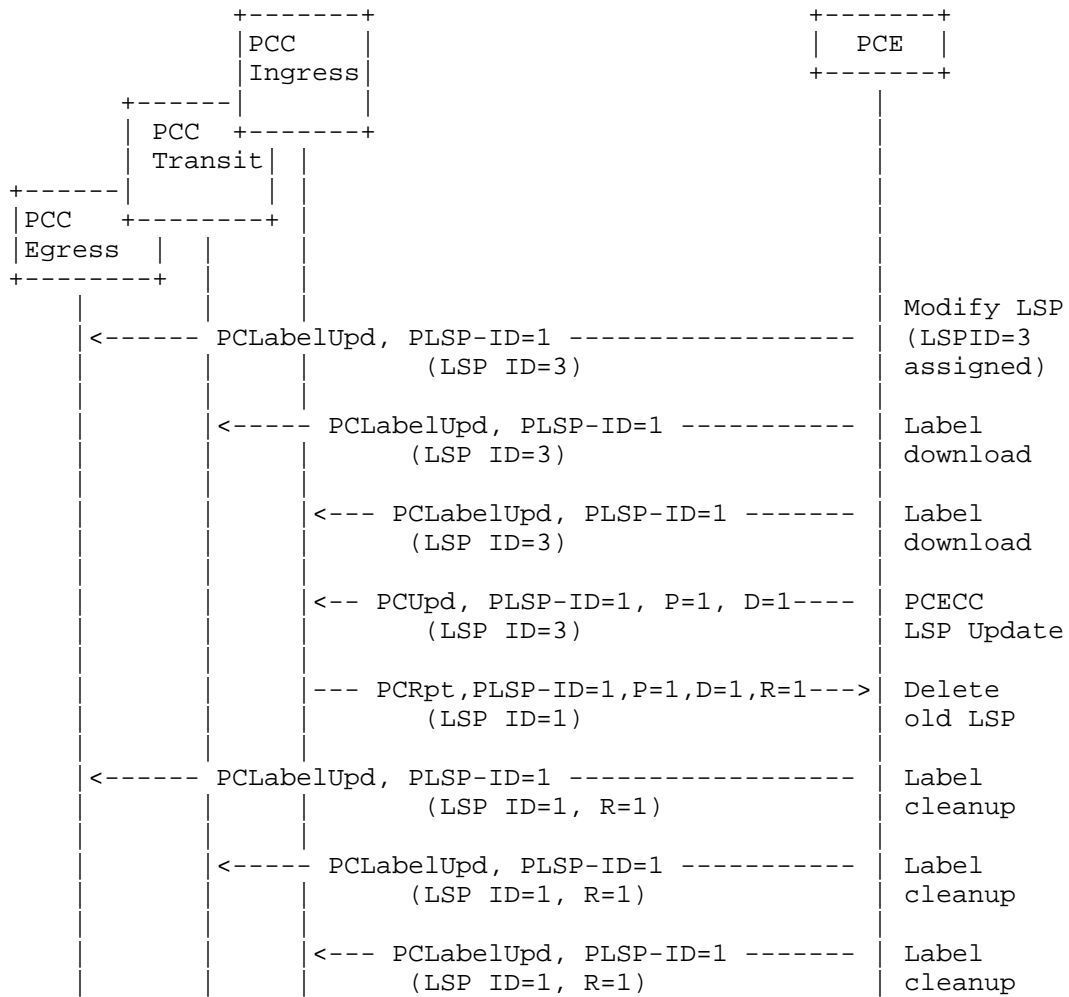


5.5.1.5. PCECC LSP Update

Incase of a modification of PCECC LSP with a new path, a PCE sends a PCUpd message to the Ingress PCC.

When a PCC received a PCUpd message for an existing LSP, a PCC MUST follow the make-before-break procedure. On successful traffic switch over to the new LSP, PCC sends a PCRpt message to the PCE for the deletion of old LSP. Further the PCE does cleanup operation for the old LSP described in Section 5.5.1.3.

The PCECC LSP Update and make-before-break sequence is shown below.



The modified PCECC LSP are considered to be 'up' by default. The Ingress MAY further choose to deploy a data plane check mechanism and report the status back to the PCE via PCRpt message.

5.5.1.6. PCECC LSP State Report

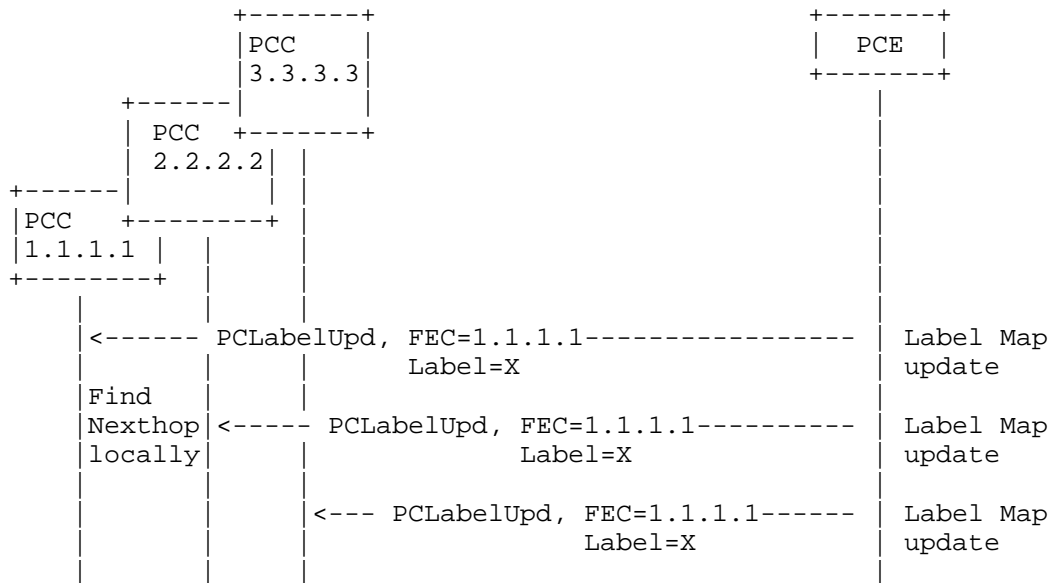
As mentioned before, an Ingress PCC MAY choose to apply any OAM mechanism to check the status of LSP in the Data plane and MAY further send its status in PCRpt message to the PCE.

5.5.2. PCECC Segment Routing (SR)

Segment Routing (SR) as described in [I-D.ietf-spring-segment-routing] depends on "segments" that are advertised by Interior Gateway Protocols (IGPs). The SR-node allocate and advertise the SID (node, adj etc) and flood via the IGP. This document proposes a new mechanism where PCE allocate the SID (label) centrally and uses PCEP to advertise the SID. In some deployments PCE (and PCEP) are better suited than IGP because of centralized nature of PCE and direct TCP based PCEP session to the node.

5.5.2.1. PCECC SR-BE

Each node (PCC) is allocated a node-SID (label) by the PCECC. The PCECC sends PCLabelUpd to update the label map of each node to all the nodes in the domain. Each node (PCC) uses the local information to determines the next-hop and download the label forwarding instructions accordingly. The PCLabelUpd message in this case MUST not have LSP object but uses new FEC object.



The forwarding behavior and the end result is similar to IGP based "Node-SID" in SR. Thus, from anywhere in the domain, it enforces the ECMP-aware shortest- path forwarding of the packet towards the related node.

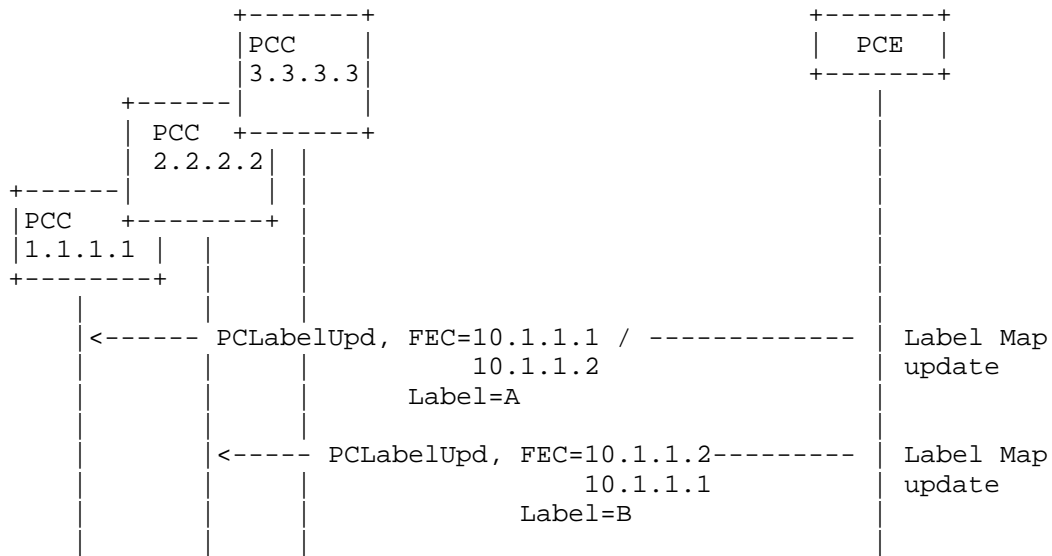
PCE rely on the Node label cleanup using the same PCLabelUpd message.

5.5.2.2. PCECC SR-TE

A Segment Routed Best Effort path (SR-BE path) can be derived from an IGP Shortest Path Tree (SPT) as explained above. On the other hand, SR-TE paths may not follow IGP SPT. Such paths may be chosen by a PCE and provisioned on the source node of the SR-TE path.

[I-D.ietf-pce-segment-routing] extends PCEP to allow a stateful PCE to compute and initiate SR-TE paths, as well as a PCC to request a path subject to certain constraint(s) and optimization criteria in SR networks.

For SR-TE, apart from node-SID, Adj-SID is used where each adjacency is allocated an Adj-SID (label) by the PCECC. The PCECC sends PCLabelUpd to update the label map of each Adj to the corresponding nodes in the domain. Each node (PCC) download the label forwarding instructions accordingly. Similar to SR-BE, the PCLabelUpd message in this case MUST not have LSP object but uses new FEC object.



The forwarding behavior and the end result is similar to IGP based "Adj-SID" in SR.

The Path Setup Type MUST be set for PCECC SR-TE (see Section 7.3). The rest of the PCEP procedures and mechanism are similar to [I-D.ietf-pce-segment-routing].

PCE rely on the Adj label cleanup using the same PCLLabelUpd message.

6. PCEP messages

As defined in [RFC5440], a PCEP message consists of a common header followed by a variable-length body made of a set of objects that can be either mandatory or optional. An object is said to be mandatory in a PCEP message when the object must be included for the message to be considered valid. For each PCEP message type, a set of rules is defined that specify the set of objects that the message can carry. An implementation MUST form the PCEP messages using the object ordering specified in this document.

6.1. The PCLRRResv message

A Label Range Reservation message (also referred to as PCLRRResv message) is a PCEP message sent by a PCC to a PCE for the reservation of label range or by PCE to PCC to send reserved label range for the network. The Message-Type field of the PCEP common header for the PCLRRResv message is set to [TBD].

The format of a PCLRRResv message is as follows:

```
PCLRRResv Message ::= <Common Header>
                               <label-range>
```

Where:

```
<label-range> ::= <SRP>
                  <labelrange-list>
```

Where

```
<labelrange-list> ::= <LABEL-RANGE> [<labelrange-list>]
```

There are two mandatory objects that MUST be included within each <label-range> in the PCLRRResv message: the SRP Object and LABEL-RANGE object.

SRP object is defined in [I-D.ietf-pce-stateful-pce] and this document extends the use of SRP object in PCLRRResv message. If the SRP object is missing, the receiving PCE MUST send a PCErr message with Error-type=6 (Mandatory Object missing) and Error-value=10 (SRP object missing).

PCC generates the value of SRP-ID-number in SRP object of PCLRResv message send to a PCE. The PCE MUST include the same SRP-ID-number in SRP object of PCLRResv message sent to the PCC in response to PCLRResv message.

LABEL-RANGE object is defined in Section 7.2. If the LABEL-RANGE object is missing, the receiving PCE MUST send a PCErr message with Error-type=6 (Mandatory Object missing) and Error-value=[TBD] (Label object missing).

[Editor's Note: This section of the document would be updated with more details about Label Block Negotiation, Reservation, Adjustment etc in a future revision of the document.]

6.2. The PCLabelUpd message

The Label Update Message (also referred to as PCLabelUpd) is a PCEP message sent by a PCE to a PCC to download label or update the label map. The same message is also used to cleanup the Label entry. The Message-Type field of the PCEP common header for the PCLabelUpd message is set to [TBD].

The format of the PCLabelUpd message is as follows:

```
<PCLabelUpd Message> ::= <Common Header>
                               <pce-label-update-list>
```

Where:

```
<pce-label-update-list> ::= <pce-label-update>
                               [<pce-label-update-list>]
```

```
<pce-label-update> ::= (<pce-label-download>|<pce-label-map>)
```

Where:

```
<pce-label-download> ::= <SRP>
                               <LSP>
                               <label-list>
```

```
<pce-label-map> ::= <SRP>
                               <LABEL>
                               <FEC>
```

```
<label-list > ::= <LABEL>
                               [<label-list>]
```

The PCLabelUpd message is used to download label along the path of the LSP for the basic PCECC mode, as well as to update the label map for the Node and Adjacency Label in case of SR.

The SRP object is defined in [I-D.ietf-pce-stateful-pce] and this document extends the use of SRP object in PCLabelUpd message. The SRP object is mandatory and MUST be included in PCLabelUpd message. If the SRP object is missing, the receiving PCC MUST send a PCerr message with Error-type=6 (Mandatory Object missing) and Error-value=10 (SRP object missing).

The LSP object is defined in [I-D.ietf-pce-stateful-pce] and this document extends the use of LSP object in PCLabelUpd message. The LSP is an optional object and used in the basic PCECC mode in PCLabelUpd message. LSP Identifiers TLV is defined in [I-D.ietf-pce-stateful-pce], it MUST be included in the LSP object in PCLabelUpd message. If the TLV is missing, the PCC will generate a PCerr message with Error-Type=6 (mandatory object missing) and Error-Value=11 (LSP-IDENTIFIERS TLV missing) and close the session.

The LABEL object is defined in Section 7.4. The LABEL is the mandatory object and MUST be included in PCLabelUpd message. If the LABEL object is missing, the receiving PCC MUST send a PCerr message with Error-type=6 (Mandatory Object missing) and Error-value=[TBD] (LABEL object missing). More than one LABEL object MAY be included in the PCLabelUpd message for the transit LSR.

The FEC object is defined in Section 7.5. The FEC is an optional object and used in PCECC SR mode in PCLabelUpd message. The FEC object encodes the Node and Adjacency information of the Label Map.

To cleanup the SRP object must set the R (remove) bit.

7. PCEP Objects

The PCEP objects defined in this document are compliant with the PCEP object format defined in [RFC5440]. The P flag and the I flag of the PCEP objects defined in this document MUST always be set to 0 on transmission and MUST be ignored on receipt since these flags are exclusively related to path computation requests.

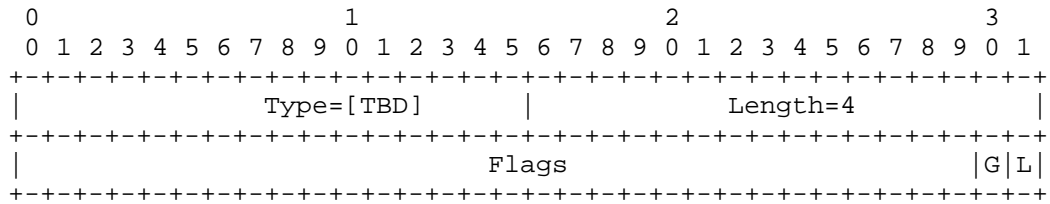
7.1. OPEN Object

This document defines a new optional TLV for use in the OPEN Object.

7.1.1. PCECC Capability TLV

The PCECC-CAPABILITY TLV is an optional TLV for use in the OPEN Object for PCECC capability advertisement. Advertisement of the PCECC capability implies support of LSPs that are setup through PCECC as per PCEP extensions defined in this document.

Its format is shown in the following figure:



The type of the TLV is [TBD] and it has a fixed length of 4 octets.

The value comprises a single field - Flags (32 bits):

L (LOCAL-LABEL-RANGE-CAPABILITY - 1 bit): If set to 1 by a PCEP speaker, it indicates that the PCEP speaker is capable for local label range reservation.

G (GLOBAL-LABEL-RANGE-CAPABILITY - 1 bit): If set to 1 by a PCEP speaker, it indicates that the PCEP speaker capable for global label range reservation.

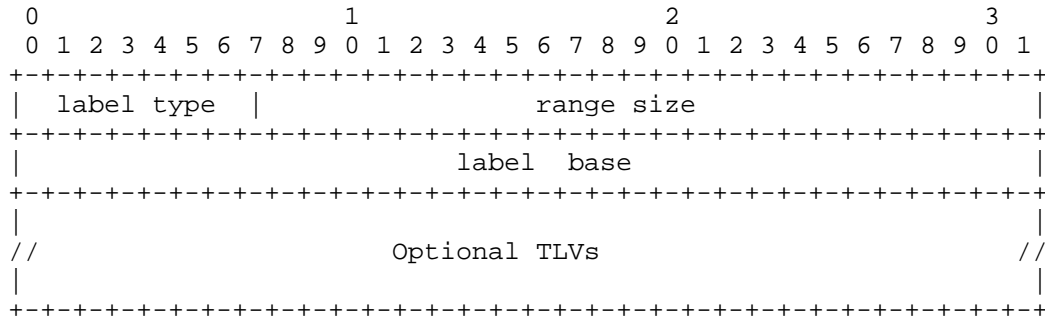
Unassigned bits are considered reserved. They MUST be set to 0 on transmission and MUST be ignored on receipt.

7.2. LABEL-RANGE Object

The LABEL-RANGE object MUST be carried within PCLRResv message. The LABEL-RANGE object is used to carry the label range information based on the label type.

LABEL-RANGE Object-Class is TBD.

LABEL-RANGE Object-Type is 1.



label type (8 bit): The values defined for label type are label type 1 specifies the local label. It means the label range is non negotiable. label type 2 specifies the global label. It means the label range is negotiable. Refer [I-D.li-mpls-global-label-framework] for global label.

Range size (24 bit): It specifies the size of label range.

Label base (32 bit): It specifies the minimum label of label range.

7.3. PATH-SETUP-TYPE TLV

The PATH-SETUP-TYPE TLV is defined in [I-D.sivabalan-pce-lsp-setup-type]; this document defines following new PST value:

- o PST = 2: Path is setup via Basic PCECC mode.
- o PST = 3: Path is setup via PCECC SR-TE mode.

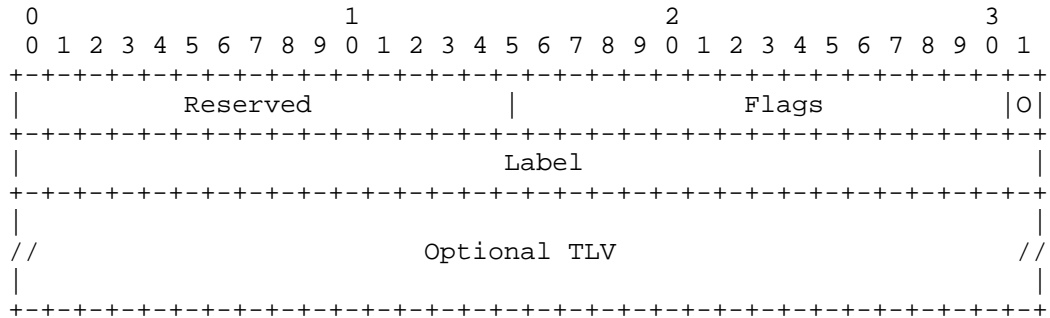
On a PCRpt or PCInitiate message, the PST=2 in PATH-SETUP-TYPE TLV in SRP object indicates that this LSP was setup via a basic PCECC based mechanism; the PST=3 indicates that this LSP was setup via a PCECC SR-TE based mechanism.

7.4. Label Object

The LABEL Object is used to specify the Label information and MUST be carried within PCLabelUpd message.

LABEL Object-Class is TBD.

LABEL Object-Type is 1.



The fields in the LABEL object are as follows:

Flags: is used to carry any additional information pertaining to the label. Currently, the following flag bit is defined:

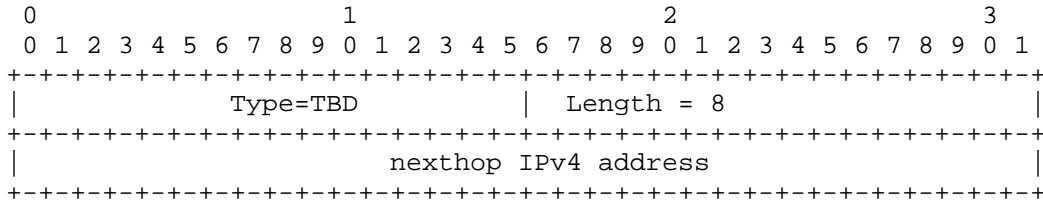
- * 0 bit(Out-label) : if the bit is set it specifies the label is the OUT label and it is mandatory to encode the nexthop information (via NEXTHOP-IPV4-ADDRESS TLV or NEXTHOP-IPV6-ADDRESS TLV or NEXTHOP-UNNUMBERED-IPV4-ID TLV in LABEL object).

Label (32-bit): The Label information encoded such that the 20 rightmost bits represent a label.

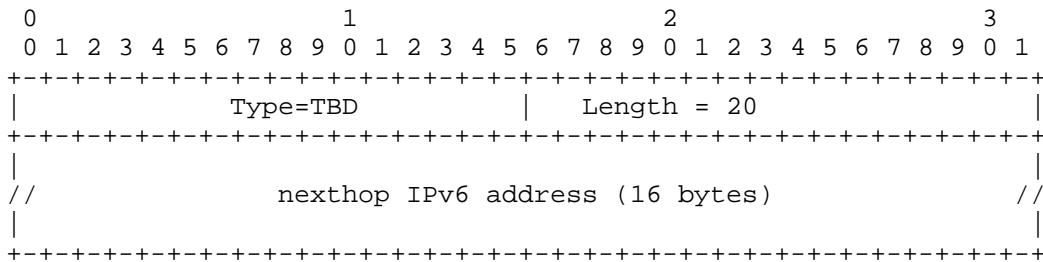
7.4.1. NextHop TLV

This document defines the following TLV for the LABEL object to associate the nexthop information incase of an outgoing label.

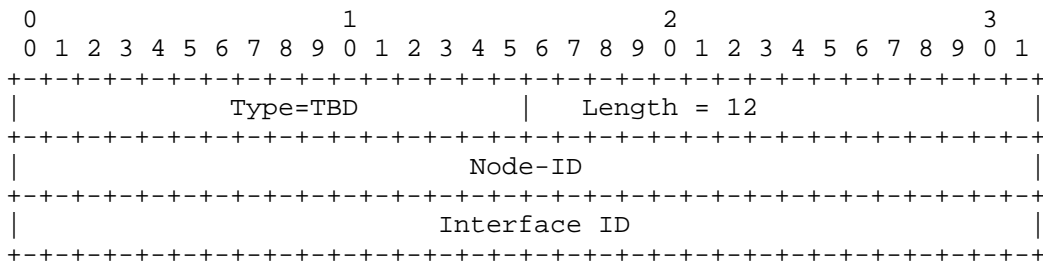
NEXTHOP-IPV4-ADDRESS TLV:



NEXTHOP-IPV6-ADDRESS TLV:



NEXTHOP-UNNUMBERED-IPV4-ID TLV:



The NextHop TLVs are as follows:

NEXTHOP-IPV4-ADDRESS TLV: where Nexthop IPv4 address is specified as an IPv4 address of the Nexthop.

NEXTHOP-IPV6-ADDRESS TLV: where Nexthop IPv6 address is specified as an IPv6 address of the Nexthop.

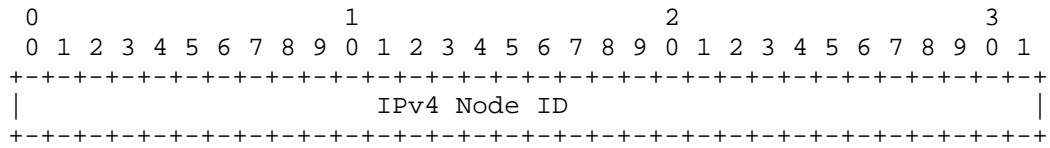
NEXTHOP-UNNUMBERED-IPV4-ID TLV: where a pair of Node ID / Interface ID tuples is used for the Nexthop.

7.5. FEC Object

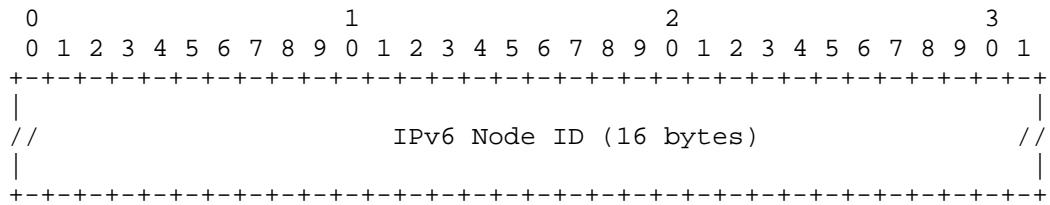
The FEC Object is used to specify the FEC information and MAY be carried within PCLabelUpd message.

FEC Object-Class is TBD.

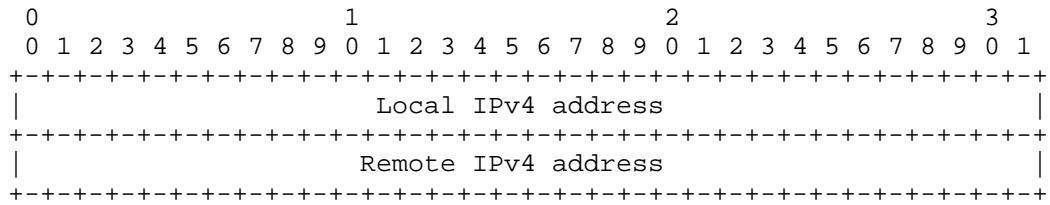
FEC Object-Type is 1 'IPv4 Node ID'.



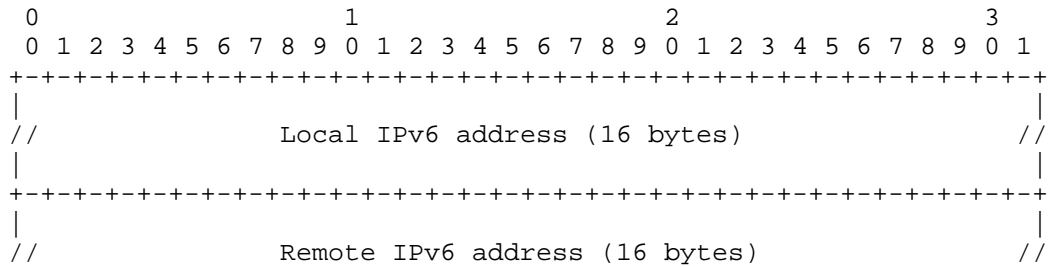
FEC Object-Type is 2 'IPv6 Node ID'.

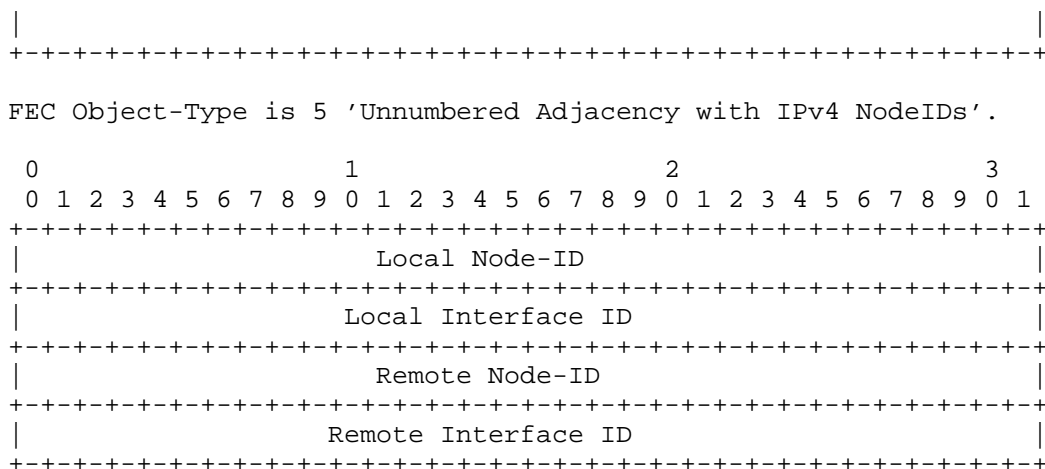


FEC Object-Type is 3 'IPv4 Adjacency'.



FEC Object-Type is 4 'IPv6 Adjacency'.





The FEC objects are as follows:

IPv4 Node ID: where IPv4 Node ID is specified as an IPv4 address of the Node. FEC Object-type is 1, and the Object-Length is 4 in this case.

IPv6 Node ID: where IPv6 Node ID is specified as an IPv6 address of the Node. FEC Object-type is 2, and the Object-Length is 16 in this case.

IPv4 Adjacency: where Local and Remote IPv4 address is specified as pair of IPv4 address of the adjacency. FEC Object-type is 3, and the Object-Length is 8 in this case.

IPv6 Adjacency: where Local and Remote IPv6 address is specified as pair of IPv6 address of the adjacency. FEC Object-type is 4, and the Object-Length is 32 in this case.

Unnumbered Adjacency with IPv4 NodeID: where a pair of Node ID / Interface ID tuples is used. FEC Object-type is 5, and the Object-Length is 16 in this case.

8. Security Considerations

TBD

9. Manageability Considerations

9.1. Control of Function and Policy

TBD.

9.2. Information and Data Models

TBD.

9.3. Liveness Detection and Monitoring

TBD.

9.4. Verify Correct Operations

TBD.

9.5. Requirements On Other Protocols

TBD.

9.6. Impact On Network Operations

TBD.

10. IANA Considerations

TBD

11. Acknowledgments

We would like to thank Robert Tao, Changjing Yan, Tieying Huang for their useful comments and suggestions.

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PCEP Procedures and Protocol Extensions for Using PCE as a Central
Controller (PCECC) of LSPs
draft-zhao-pce-pcep-extension-for-pce-controller-08

Abstract

The Path Computation Element (PCE) is a core component of Software-Defined Networking (SDN) systems. It can compute optimal paths for traffic across a network and can also update the paths to reflect changes in the network or traffic demands.

PCE was developed to derive paths for MPLS Label Switched Paths (LSPs), which are supplied to the head end of the LSP using the Path Computation Element Communication Protocol (PCEP). But SDN has a broader applicability than signaled (G)MPLS traffic-engineered (TE) networks, and the PCE may be used to determine paths in a range of use cases. PCEP has been proposed as a control protocol for use in these environments to allow the PCE to be fully enabled as a central controller.

A PCE-based central controller (PCECC) can simplify the processing of a distributed control plane by blending it with elements of SDN and without necessarily completely replacing it. Thus, the LSP can be calculated/setup/initiated and the label forwarding entries can also be downloaded through a centralized PCE server to each network devices along the path while leveraging the existing PCE technologies as much as possible.

This document specifies the procedures and PCEP protocol extensions for using the PCE as the central controller.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

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1. Introduction

The Path Computation Element (PCE) [RFC4655] was developed to offload path computation function from routers in an MPLS traffic-engineered network. Since then, the role and function of the PCE has grown to cover a number of other uses (such as GMPLS [RFC7025]) and to allow delegated control [RFC8231] and PCE-initiated use of network resources [RFC8281].

According to [RFC7399], Software-Defined Networking (SDN) refers to a separation between the control elements and the forwarding components so that software running in a centralized system, called a controller, can act to program the devices in the network to behave in specific ways. A required element in an SDN architecture is a component that plans how the network resources will be used and how the devices will be programmed. It is possible to view this

component as performing specific computations to place traffic flows within the network given knowledge of the availability of network resources, how other forwarding devices are programmed, and the way that other flows are routed. This is the function and purpose of a PCE, and the way that a PCE integrates into a wider network control system (including an SDN system) is presented in [RFC7491].

In early PCE implementations, where the PCE was used to derive paths for MPLS Label Switched Paths (LSPs), paths were requested by network elements (known as Path Computation Clients (PCCs)), and the results of the path computations were supplied to network elements using the Path Computation Element Communication Protocol (PCEP) [RFC5440]. This protocol was later extended to allow a PCE to send unsolicited requests to the network for LSP establishment [RFC8281].

[RFC8283] introduces the architecture for PCE as a central controller as an extension of the architecture described in [RFC4655] and assumes the continued use of PCEP as the protocol used between PCE and PCC. [RFC8283] further examines the motivations and applicability for PCEP as a Southbound Interface (SBI), and introduces the implications for the protocol. [I-D.ietf-teas-pcecc-use-cases] describes the use cases for the PCECC architecture.

A PCE-based central controller (PCECC) can simplify the processing of a distributed control plane by blending it with elements of SDN and without necessarily completely replacing it. Thus, the LSP can be calculated/setup/initiated and the label forwarding entries can also be downloaded through a centralized PCE server to each network devices along the path while leveraging the existing PCE technologies as much as possible.

This draft specify the procedures and PCEP protocol extensions for using the PCE as the central controller for static LSPs, where LSPs can be provisioned as explicit label instructions at each hop on the end-to-end path. Each router along the path must be told what label-forwarding instructions to program and what resources to reserve. The PCE-based controller keeps a view of the network and determines the paths of the end-to-end LSPs, and the controller uses PCEP to communicate with each router along the path of the end-to-end LSP.

The extension for PCECC in Segment Routing (SR) is specified in a separate draft [I-D.zhao-pce-pcep-extension-pce-controller-sr].

1.1. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

2. Terminology

Terminologies used in this document is same as described in the draft [RFC8283] and [I-D.ietf-teas-pcecc-use-cases].

3. Basic PCECC Mode

In this mode LSPs are provisioned as explicit label instructions at each hop on the end-to-end path. Each router along the path must be told what label forwarding instructions to program and what resources to reserve. The controller uses PCEP to communicate with each router along the path of the end-to-end LSP.

Note that the PCE-based controller will take responsibility for managing some part of the MPLS label space for each of the routers that it controls, and may take wider responsibility for partitioning the label space for each router and allocating different parts for different uses. This is also described in section 3.1.2. of [RFC8283]. For the purpose of this document, it is assumed that label range to be used by a PCE is known and set on both PCEP peers. A future extension could add this capability to advertise the range via possible PCEP extensions as well. The rest of processing is similar to the existing stateful PCE mechanism.

4. PCEP Requirements

Following key requirements associated PCECC should be considered when designing the PCECC based solution:

1. PCEP speaker supporting this draft MUST have the capability to advertise its PCECC capability to its peers.
2. PCEP speaker not supporting this draft MUST be able to reject PCECC related extensions with a error reason code that indicates that this feature is not supported.
3. PCEP speaker MUST provide a means to identify PCECC based LSP in the PCEP messages.

4. PCEP procedures SHOULD provide a means to update (or cleanup) the label- download entry to the PCC.
 5. PCEP procedures SHOULD provide a means to synchronize the labels between PCE to PCC in PCEP messages.
5. Procedures for Using the PCE as the Central Controller (PCECC)
- 5.1. Stateful PCE Model

Active stateful PCE is described in [RFC8231]. PCE as a central controller (PCECC) reuses existing Active stateful PCE mechanism as much as possible to control the LSP.

5.2. New LSP Functions

This document defines the following new PCEP messages and extends the existing messages to support PCECC:

(PCRpt): a PCEP message described in [RFC8231]. PCRpt message is used to send PCECC LSP Reports. It is also extended to report the set of Central Controller's Instructions (CCI) (label forwarding instructions in the context of this document) received from the PCE. See Section 5.4.6 for more details.

(PCInitiate): a PCEP message described in [RFC8281]. PCInitiate message is used to setup PCE-Initiated LSP based on PCECC mechanism. It is also extended for Central Controller's Instructions (CCI) (download or cleanup the Label forwarding instructions in the context of this document) on all nodes along the path.

(PCUpd): a PCEP message described in [RFC8231]. PCUpd message is used to send PCECC LSP Update.

The new LSP functions defined in this document are mapped onto the messages as shown in the following table.

Function	Message
PCECC Capability advertisement	Open
Label entry Add	PCInitiate
Label entry Cleanup	PCInitiate
PCECC Initiated LSP	PCInitiate
PCECC LSP Update	PCUpd
PCECC LSP State Report	PCRpt
PCECC LSP Delegation	PCRpt
PCECC Label Report	PCRpt

This document specifies a new object CCI (see Section 7.3) for the encoding of central controller's instructions. In the scope of this document this is limited to Label forwarding instructions. The CC-ID is the unique identifier for the central controller's instructions in PCEP. The PCEP messages are extended in this document to handle the PCECC operations.

5.3. PCECC Capability Advertisement

During PCEP Initialization Phase, PCEP Speakers (PCE or PCC) advertise their support of PCECC extensions.

This document defines a new Path Setup Type (PST) [I-D.ietf-pce-lsp-setup-type] for PCECC, as follows:

- o PST = TBD: Path is setup via PCECC mode.

A PCEP speaker MUST indicate its support of the function described in this document by sending a PATH-SETUP-TYPE-CAPABILITY TLV in the OPEN object with this new PST included in the PST list.

This document also defines the PCECC Capability sub-TLV Section 7.1.1. PCEP speakers use this sub-TLV to exchange information about their PCECC capability. If a PCEP speaker includes PST=TBD in the PST List of the PATH-SETUP-TYPE-CAPABILITY TLV then it MUST also include the PCECC Capability sub-TLV inside the PATH-SETUP-TYPE-CAPABILITY TLV.

The presence of the PST and PCECC Capability sub-TLV in PCC's OPEN Object indicates that the PCC is willing to function as a PCECC client.

The presence of the PST and PCECC Capability sub-TLV in PCE's OPEN message indicates that the PCE is interested in function as a PCECC server.

The PCEP protocol extensions for PCECC MUST NOT be used if one or both PCEP Speakers have not included the PST or the PCECC Capability sub-TLV in their respective OPEN message. If the PCEP Speakers support the extensions of this draft but did not advertise this capability then a PCerr message with Error-Type=19(Invalid Operation) and Error-Value=TBD (Attempted PCECC operations when PCECC capability was not advertised) will be generated and the PCEP session will be terminated.

A PCC or a PCE MUST include both PCECC-CAPABILITY sub-TLV and STATEFUL-PCE-CAPABILITY TLV ([RFC8231]) (with I flag set [RFC8281]) in OPEN Object to support the extensions defined in this document. If PCECC-CAPABILITY sub-TLV is advertised and STATEFUL-PCE-CAPABILITY TLV is not advertised in OPEN Object, it SHOULD send a PCerr message with Error-Type=19 (Invalid Operation) and Error-value=TBD (stateful PCE capability was not advertised) and terminate the session.

5.4. LSP Operations

The PCEP messages pertaining to PCECC MUST include PATH-SETUP-TYPE TLV [I-D.ietf-pce-lsp-setup-type] in the SRP object to clearly identify the PCECC LSP is intended.

5.4.1. Basic PCECC LSP Setup

In order to setup a LSP based on PCECC mechanism, a PCC MUST delegate the LSP by sending a PCRpt message with PST set for PCECC (see Section 7.2) and D (Delegate) flag (see [RFC8231]) set in the LSP object.

LSP-IDENTIFIER TLV MUST be included for PCECC LSP, the tuple uniquely identifies the LSP in the network. The LSP object is included in central controller's instructions (label download) to identify the PCECC LSP for this instruction. The PLSP-ID is the original identifier used by the ingress PCC, so the transit LSR could have multiple central controller instructions that have the same PLSP-ID. The PLSP-ID in combination with the source (in LSP-IDENTIFIER TLV) MUST be unique. The PLSP-ID is included for maintainability reasons. As per [RFC8281], the LSP object could include SPEAKER-ENTITY-ID TLV to identify the PCE that initiated these instructions. Also the CC-ID is unique on the PCEP session as described in Section 7.3.

When a PCE receives PCRpt message with D flags and PST Type set, it calculates the path and assigns labels along the path; and set up the

path by sending PCInitiate message to each node along the path of the LSP. The PCC generates a Path Computation State Report (PCRpt) and include the central controller's instruction (CCI) and the identified LSP. The CC-ID is uniquely identify the central controller's instruction within PCEP. The PCC further responds with the PCRpt messages including the CCI and LSP objects.

Once the central controller's instructions (label operations) are completed, the PCE SHOULD send the PCUpd message to the Ingress PCC. The PCUpd message is as per [RFC8231] SHOULD include the path information as calculated by the PCE.

Note that the PCECC LSPs MUST be delegated to a PCE at all times.

LSP deletion operation for PCECC LSP is same as defined in [RFC8231]. If the PCE receives PCRpt message for LSP deletion then it does Label cleanup operation as described in Section 5.4.2.2 for the corresponding LSP.

The Basic PCECC LSP setup sequence is as shown below.

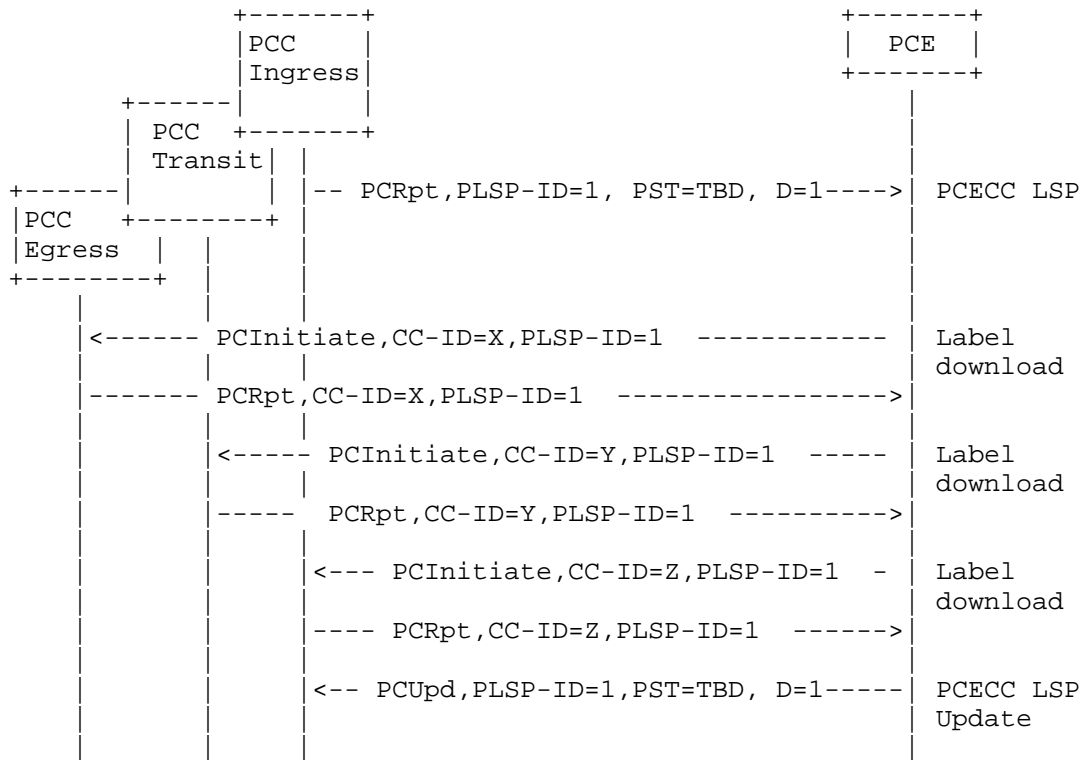


Figure 2: Basic PCECC LSP setup

The PCECC LSP are considered to be 'up' by default (on receipt of PCUpd message from PCE). The Ingress MAY further choose to deploy a data plane check mechanism and report the status back to the PCE via PCRpt message.

5.4.2. Central Control Instructions

The new central controller's instructions (CCI) for the label operations in PCEP is done via the PCInitiate message, by defining a new PCEP Objects for CCI operations. Local label range of each PCC is assumed to be known at both the PCC and the PCE.

5.4.2.1. Label Download

In order to setup an LSP based on PCECC, the PCE sends a PCInitiate message to each node along the path to download the Label instruction as described in Section 5.4.1.

The CCI object MUST be included, along with the LSP object in the PCInitiate message. The LSP-IDENTIFIER TLV MUST be included in LSP object. The SPEAKER-ENTITY-ID TLV SHOULD be included in LSP object.

If a node (PCC) receives a PCInitiate message which includes a Label to download as part of CCI, that is out of the range set aside for the PCE, it MUST send a PCErr message with Error-type=TBD (PCECC failure) and Error-value=TBD (Label out of range) and MUST include the SRP object to specify the error is for the corresponding label update via PCInitiate message. If a PCC receives a PCInitiate message but failed to download the Label entry, it MUST send a PCErr message with Error-type=TBD (PCECC failure) and Error-value=TBD (instruction failed) and MUST include the SRP object to specify the error is for the corresponding label update via PCInitiate message.

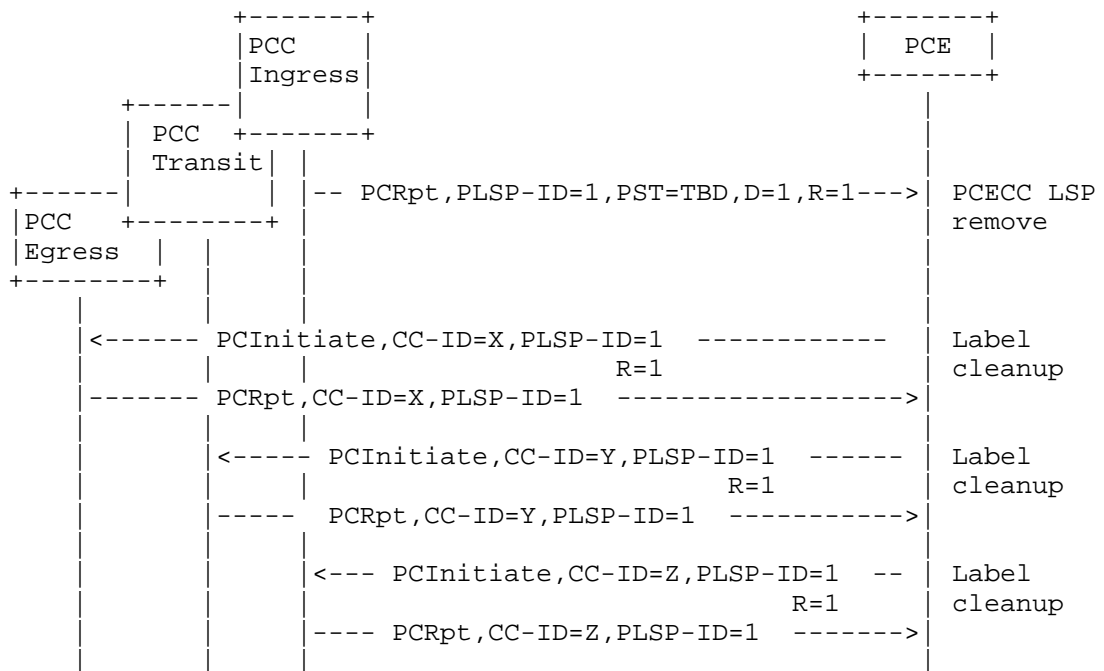
New PCEP object for central control instructions (CCI) is defined in Section 7.3.

5.4.2.2. Label Cleanup

In order to delete an LSP based on PCECC, the PCE sends a central controller instructions via a PCInitiate message to each node along the path of the LSP to cleanup the Label forwarding instruction.

If the PCC receives a PCInitiate message but does not recognize the label in the CCI, the PCC MUST generate a PCErr message with Error-Type 19(Invalid operation) and Error-Value=TBD, "Unknown Label" and MUST include the SRP object to specify the error is for the corresponding label cleanup (via PCInitiate message).

The R flag in the SRP object defined in [RFC8281] specifies the deletion of Label Entry in the PCInitiate message.



As per [RFC8281], following the removal of the Label forwarding instruction, the PCC MUST send a PCRpt message. The SRP object in the PCRpt MUST include the SRP-ID-number from the PCInitiate message that triggered the removal. The R flag in the SRP object MUST be set.

5.4.3. PCE Initiated PCECC LSP

The LSP Instantiation operation is same as defined in [RFC8281].

In order to setup a PCE Initiated LSP based on the PCECC mechanism, a PCE sends PCInitiate message with Path Setup Type set for PCECC (see Section 7.2) to the Ingress PCC.

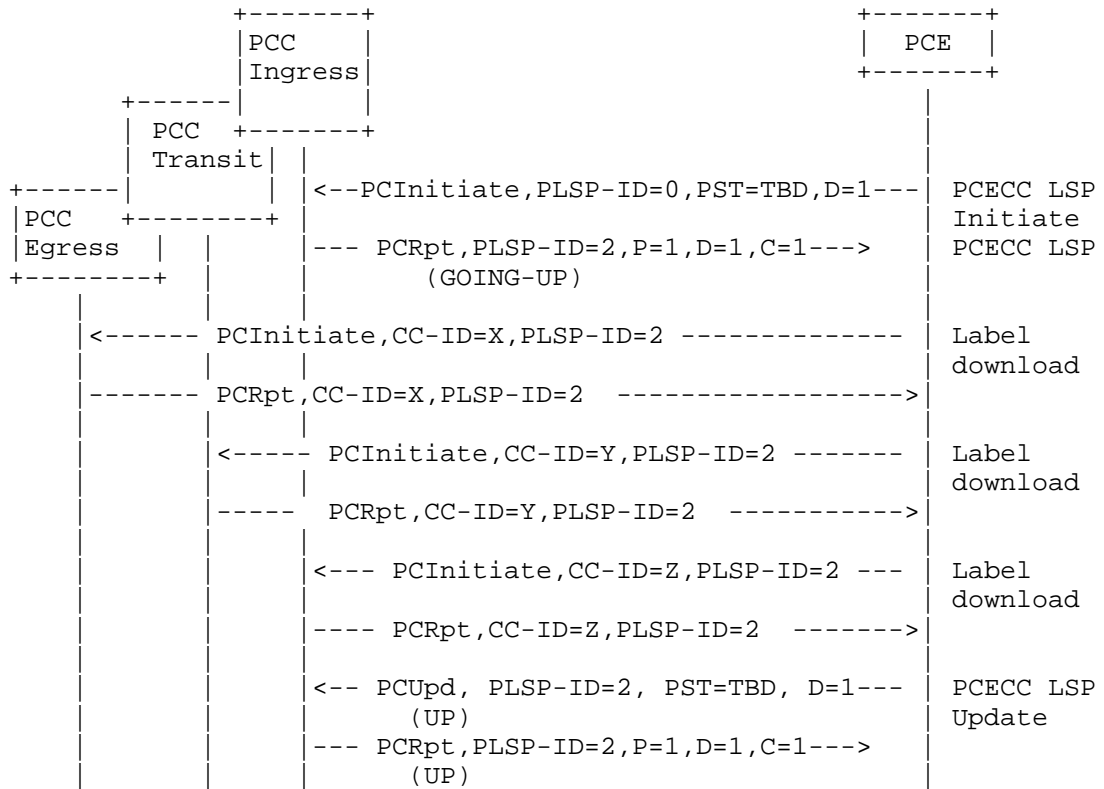
The Ingress PCC MUST also set D (Delegate) flag (see [RFC8231]) and C (Create) flag (see [RFC8281]) in LSP object of PCRpt message. The PCC responds with first PCRpt message with the status as "GOING-UP" and assigned PLSP-ID.

Note that the label forwarding instructions from PCECC are send after the initial PCInitiate and PCRpt exchange. This is done so that the PLSP-ID and other LSP identifiers can be obtained from the ingress and can be included in the label forwarding instruction in the next

PCInitiate message. The rest of the PCECC LSP setup operations are same as those described in Section 5.4.1.

The LSP deletion operation for PCE Initiated PCECC LSP is same as defined in [RFC8281]. The PCE should further perform Label entry cleanup operation as described in Section 5.4.2.2 for the corresponding LSP.

The PCE Initiated PCECC LSP setup sequence is shown below -

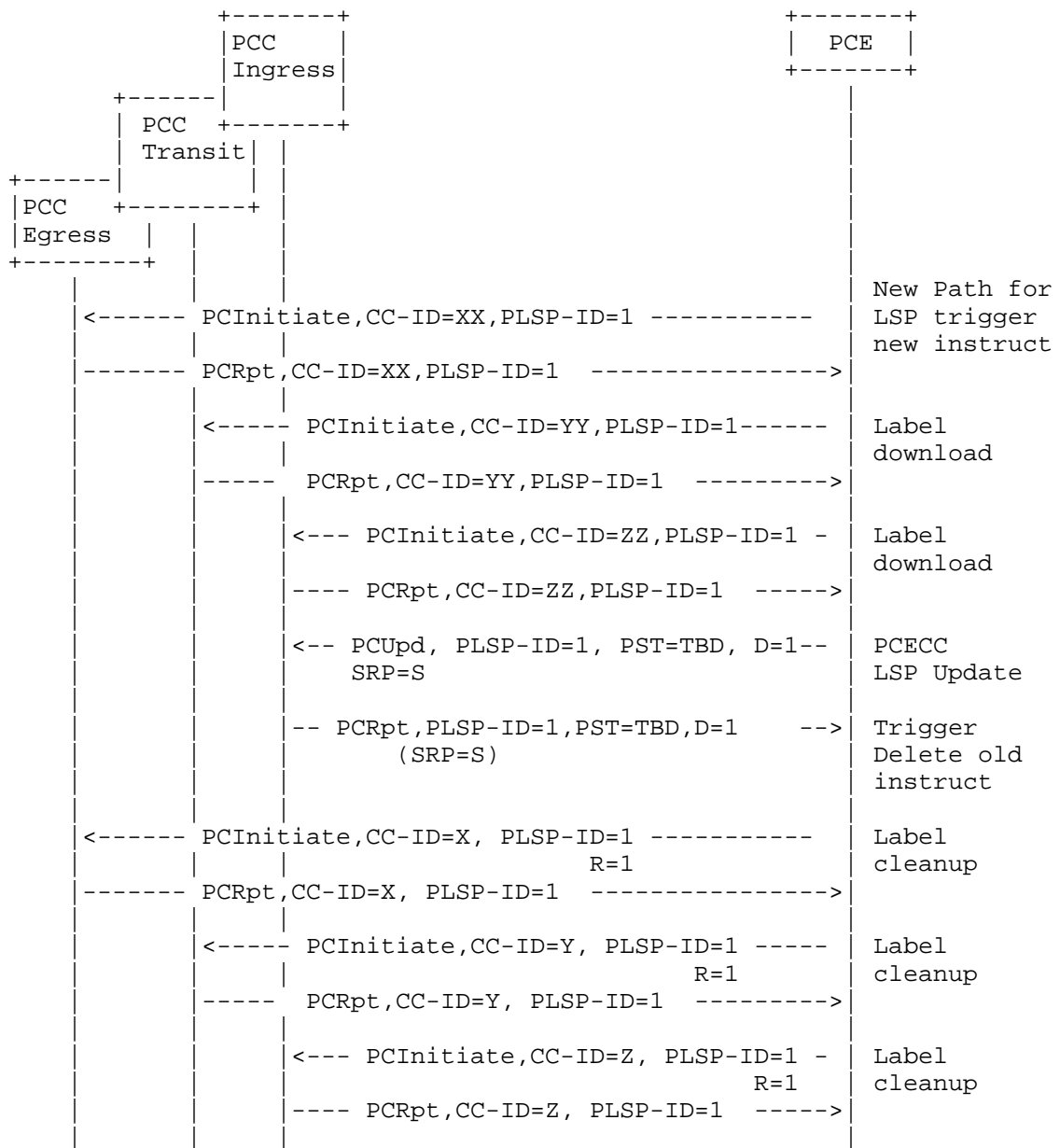


Once the label operations are completed, the PCE SHOULD send the PCUpd message to the Ingress PCC. The PCUpd message is as per [RFC8231].

5.4.4. PCECC LSP Update

In case of a modification of PCECC LSP with a new path, a PCE sends a PCUpd message to the Ingress PCC. But to follow the make-before-break procedures, the PCECC first update new instructions based on the updated LSP and then update to ingress to switch traffic, before cleaning up the old instructions. A new CC-ID is used to identify the updated instruction, the existing identifiers in the LSP object identify the existing LSP. Once new instructions are downloaded, the PCE further updates the new path at the ingress which triggers the traffic switch on the updated path. The Ingress PCC acknowledges with a PCRpt message, on receipt of PCRpt message, the PCE does cleanup operation for the old LSP as described in Section 5.4.2.2.

The PCECC LSP Update sequence is shown below -



The modified PCECC LSP are considered to be 'up' by default. The Ingress MAY further choose to deploy a data plane check mechanism and report the status back to the PCE via PCRpt message.

5.4.5. Re Delegation and Cleanup

As described in [RFC8281], a new PCE can gain control over the orphaned LSP. In case of PCECC LSP, the new PCE MUST also gain control over the central controllers instructions in the same way by sending a PCInitiate message that includes the SRP, LSP and CCI objects and carries the CC-ID and PLSP-ID identifying the instruction, it wants to take control of.

Further, as described in [RFC8281], the State Timeout Interval timer ensures that a PCE crash does not result in automatic and immediate disruption for the services using PCE-initiated LSPs. Similarly the central controller instructions are not removed immediately upon PCE failure. Instead, they are cleaned up on the expiration of this timer. This allows for network cleanup without manual intervention. The PCC MUST support removal of CCI as one of the behaviors applied on expiration of the State Timeout Interval timer.

5.4.6. Synchronization of Central Controllers Instructions

The purpose of Central Controllers Instructions synchronization (labels in the context of this document) is to make sure that the PCE's view of CCI (Labels) matches with the PCC's Label allocation. This synchronization is performed as part of the LSP state synchronization as described in [RFC8231] and [RFC8233].

As per LSP State Synchronization [RFC8231], a PCC reports the state of its LSPs to the PCE using PCRpt messages and as per [RFC8281], PCE would initiate any missing LSPs and/or remove any LSPs that are not wanted. The same PCEP messages and procedure is also used for the Central Controllers Instructions synchronization. The PCRpt message includes the CCI and the LSP object to report the label forwarding instructions. The PCE would further remove any unwanted instructions or initiate any missing instructions.

5.4.7. PCECC LSP State Report

As mentioned before, an Ingress PCC MAY choose to apply any OAM mechanism to check the status of LSP in the Data plane and MAY further send its status in PCRpt message to the PCE.

6. PCEP messages

As defined in [RFC5440], a PCEP message consists of a common header followed by a variable-length body made of a set of objects that can be either mandatory or optional. An object is said to be mandatory in a PCEP message when the object must be included for the message to be considered valid. For each PCEP message type, a set of rules is

defined that specify the set of objects that the message can carry. An implementation MUST form the PCEP messages using the object ordering specified in this document.

LSP-IDENTIFIERS TLV MUST be included in the LSP object for PCECC LSP.

6.1. The PCInitiate message

The PCInitiate message [RFC8281] can be used to download or remove the labels, the message has been extended as shown below -

```
<PCInitiate Message> ::= <Common Header>
                          <PCE-initiated-lsp-list>
```

Where:

```
<Common Header> is defined in [RFC5440]
```

```
<PCE-initiated-lsp-list> ::= <PCE-initiated-lsp-request>
                              [<PCE-initiated-lsp-list>]
```

```
<PCE-initiated-lsp-request> ::=
    (<PCE-initiated-lsp-instantiation>|
     <PCE-initiated-lsp-deletion>|
     <PCE-initiated-lsp-central-control>)
```

```
<PCE-initiated-lsp-central-control> ::= <SRP>
                                          <LSP>
                                          <cci-list>
```

```
<cci-list> ::= <CCI>
                [<cci-list>]
```

Where:

```
<PCE-initiated-lsp-instantiation> and
<PCE-initiated-lsp-deletion> are as per
[RFC8281].
```

The LSP and SRP object is defined in [RFC8231].

When PCInitiate message is used for central controller's instructions (labels), the SRP, LSP and CCI objects MUST be present. The SRP object is defined in [RFC8231] and if the SRP object is missing, the receiving PCC MUST send a PCErr message with Error-type=6 (Mandatory Object missing) and Error-value=10 (SRP object missing). The LSP object is defined in [RFC8231] and if the LSP object is missing, the receiving PCC MUST send a PCErr message with Error-type=6 (Mandatory Object missing) and Error-value=8 (LSP object missing). The CCI

object is defined in Section 7.3 and if the CCI object is missing, the receiving PCC MUST send a PCErr message with Error-type=6 (Mandatory Object missing) and Error-value=TBD (CCI object missing). More than one CCI object MAY be included in the PCInitiate message for the transit LSR.

To cleanup the SRP object must set the R (remove) bit.

At max two instances of CCI object would be included in case of transit LSR to encode both in-coming and out-going label forwarding instructions. Other instances MUST be ignored.

6.2. The PCRpt message

The PCRpt message can be used to report the labels that were allocated by the PCE, to be used during the state synchronization phase.

```
<PCRpt Message> ::= <Common Header>
                    <state-report-list>
```

Where:

```
<state-report-list> ::= <state-report>[<state-report-list>]
```

```
<state-report> ::= (<lsp-state-report>|
                    <central-control-report>)
```

```
<lsp-state-report> ::= [<SRP>]
                    <LSP>
                    <path>
```

```
<central-control-report> ::= [<SRP>]
                    <LSP>
                    <cci-list>
```

```
<cci-list> ::= <CCI>
               [<cci-list>]
```

Where:

<path> is as per [RFC8231] and the LSP and SRP object are also defined in [RFC8231].

When PCRpt message is used to report the central controller's instructions (labels), the LSP and CCI objects MUST be present. The LSP object is defined in [RFC8231] and if the LSP object is missing, the receiving PCE MUST send a PCErr message with Error-type=6 (Mandatory Object missing) and Error-value=8 (LSP object missing).

The CCI object is defined in Section 7.3 and if the CCI object is missing, the receiving PCC MUST send a PCErr message with Error-type=6 (Mandatory Object missing) and Error-value=TBD (CCI object missing). Two CCI object can be included in the PCRpt message for the transit LSR.

7. PCEP Objects

The PCEP objects defined in this document are compliant with the PCEP object format defined in [RFC5440].

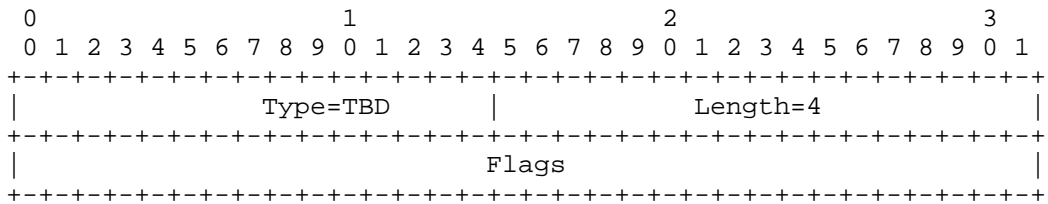
7.1. OPEN Object

This document defines a new optional TLVs for use in the OPEN Object.

7.1.1. PCECC Capability sub-TLV

The PCECC-CAPABILITY sub-TLV is an optional TLV for use in the OPEN Object for PCECC capability advertisement in PATH-SETUP-TYPE-CAPABILITY TLV. Advertisement of the PCECC capability implies support of LSPs that are setup through PCECC as per PCEP extensions defined in this document.

Its format is shown in the following figure:



The type of the TLV is TBD and it has a fixed length of 4 octets.

The value comprises a single field - Flags (32 bits).

No flags are assigned right now.

Unassigned bits are considered reserved. They MUST be set to 0 on transmission and MUST be ignored on receipt.

7.2. PATH-SETUP-TYPE TLV

The PATH-SETUP-TYPE TLV is defined in [I-D.ietf-pce-lsp-setup-type]; this document defines a new PST value:

- o PST = TBD: Path is setup via PCECC mode.

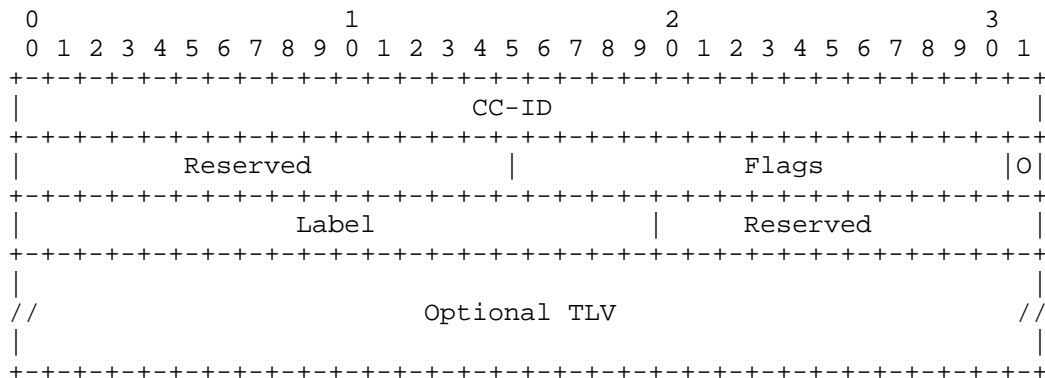
On a PCRpt/PCUpd/PCInitiate message, the PST=TBD in PATH-SETUP-TYPE TLV in SRP object indicates that this LSP was setup via a PCECC based mechanism.

7.3. CCI Object

The Central Control Instructions (CCI) Object is used by the PCE to specify the forwarding instructions (Label information in the context of this document) to the PCC, and MAY be carried within PCInitiate or PCRpt message for label download.

CCI Object-Class is TBD.

CCI Object-Type is 1 for the MPLS Label.



The fields in the CCI object are as follows:

CC-ID: A PCEP-specific identifier for the CCI information. A PCE creates an CC-ID for each instruction, the value is unique within the scope of the PCE and is constant for the lifetime of a PCEP session. The values 0 and 0xFFFFFFFF are reserved and MUST NOT be used.

Flags: is used to carry any additional information pertaining to the CCI. Currently, the following flag bit is defined:

- * O bit(Out-label) : If the bit is set, it specifies the label is the OUT label and it is mandatory to encode the next-hop information (via IPV4-ADDRESS TLV or IPV6-ADDRESS TLV or UNNUMBERED-IPV4-ID-ADDRESS TLV in the CCI object). If the bit is not set, it specifies the label is the IN label and it is optional to encode the local interface information (via IPV4-ADDRESS TLV or IPV6-ADDRESS TLV or UNNUMBERED-IPV4-ID-ADDRESS TLV in the CCI object).

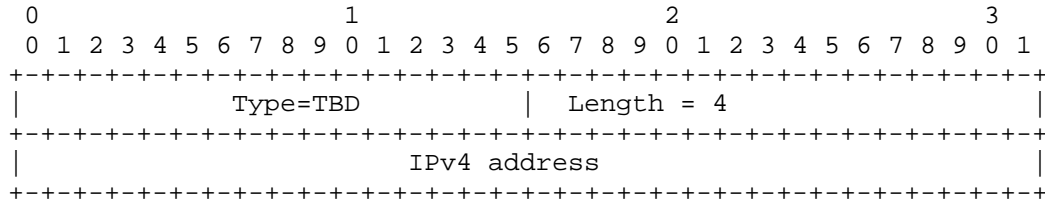
Label (20-bit): The Label information.

Reserved (12 bit): Set to zero while sending, ignored on receive.

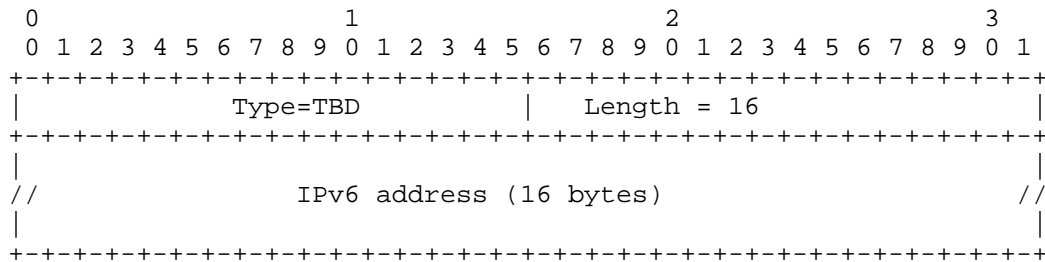
7.3.1. Address TLVs

This document defines the following TLVs for the CCI object to associate the next-hop information in case of an outgoing label and local interface information in case of an incoming label.

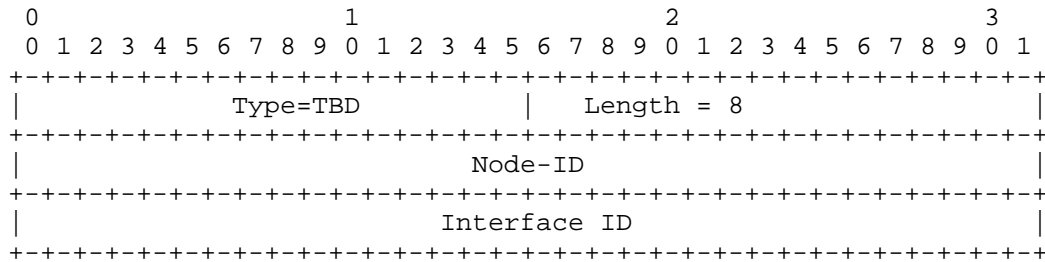
IPV4-ADDRESS TLV:



IPV6-ADDRESS TLV:



UNNUMBERED-IPV4-ID-ADDRESS TLV:



The address TLVs are as follows:

IPV4-ADDRESS TLV: an IPv4 address.

IPV6-ADDRESS TLV: an IPv6 address.

UNNUMBERED-IPV4-ID-ADDRESS TLV: a pair of Node ID / Interface ID tuples.

8. Security Considerations

The security considerations described in [RFC8231] and [RFC8281] apply to the extensions described in this document. Additional considerations related to a malicious PCE are introduced.

8.1. Malicious PCE

PCE has complete control over PCC to update the labels and can cause the LSP's to behave inappropriate and cause major impact to the network. As a general precaution, it is RECOMMENDED that these PCEP extensions only be activated on authenticated and encrypted sessions across PCEs and PCCs belonging to the same administrative authority, using Transport Layer Security (TLS) [RFC8253], as per the recommendations and best current practices in [RFC7525].

9. Manageability Considerations

9.1. Control of Function and Policy

A PCE or PCC implementation SHOULD allow to configure to enable/disable PCECC capability as a global configuration.

9.2. Information and Data Models

[RFC7420] describes the PCEP MIB, this MIB can be extended to get the PCECC capability status.

The PCEP YANG module [I-D.ietf-pce-pcep-yang] could be extended to enable/disable PCECC capability.

9.3. Liveness Detection and Monitoring

Mechanisms defined in this document do not imply any new liveness detection and monitoring requirements in addition to those already listed in [RFC5440].

9.4. Verify Correct Operations

Mechanisms defined in this document do not imply any new operation verification requirements in addition to those already listed in [RFC5440] and [RFC8231].

9.5. Requirements On Other Protocols

PCEP extensions defined in this document do not put new requirements on other protocols.

9.6. Impact On Network Operations

PCEP extensions defined in this document do not put new requirements on network operations.

10. IANA Considerations

10.1. PCEP TLV Type Indicators

IANA is requested to confirm the early allocation of the following TLV Type Indicator values within the "PCEP TLV Type Indicators" sub-registry of the PCEP Numbers registry, and to update the reference in the registry to point to this document, when it is an RFC:

Value	Meaning	Reference
TBD	PCECC-CAPABILITY	This document
TBD	IPV4-ADDRESS TLV	This document
TBD	IPV6-ADDRESS TLV	This document
TBD	UNNUMBERED-IPV4-ID-ADDRESS TLV	This document

10.2. New Path Setup Type Registry

IANA is requested to allocate new PST Field in PATH- SETUP-TYPE TLV. The allocation policy for this new registry should be by IETF Consensus. The new registry should contain the following value:

Value	Description	Reference
TBD	Traffic engineering path is setup using PCECC mode	This document

10.3. PCEP Object

IANA is requested to allocate new registry for CCI PCEP object.

Object-Class	Value	Name	Reference
TBD		CCI Object-Type	This document
		1	MPLS Label

10.4. CCI Object Flag Field

IANA is requested to create a registry to manage the Flag field of the CCI object.

One bit to be defined for the CCI Object flag field in this document:

Codespace of the Flag field (CCI Object)

Bit	Description	Reference
7	Specifies label is out label	This document

10.5. PCEP-Error Object

IANA is requested to allocate new error types and error values within the "PCEP-ERROR Object Error Types and Values" sub-registry of the PCEP Numbers registry for the following errors:

Error-Type	Meaning	
-----	-----	
19	Invalid operation.	
	Error-value = TBD :	Attempted PCECC operations when PCECC capability was not advertised
	Error-value = TBD :	Stateful PCE capability was not advertised
6	Mandatory Object missing.	Unknown Label
	Error-value = TBD :	CCI object missing
TBD	PCECC failure.	
	Error-value = TBD :	Label out of range.
	Error-value = TBD :	Instruction failed.

11. Acknowledgments

We would like to thank Robert Tao, Changjing Yan, Tieying Huang and Avantika for their useful comments and suggestions.

12. References

12.1. Normative References

- [RFC2119] Bradner, S., "Key words for use in RFCs to Indicate Requirement Levels", BCP 14, RFC 2119, DOI 10.17487/RFC2119, March 1997, <<https://www.rfc-editor.org/info/rfc2119>>.

- [RFC5440] Vasseur, JP., Ed. and JL. Le Roux, Ed., "Path Computation Element (PCE) Communication Protocol (PCEP)", RFC 5440, DOI 10.17487/RFC5440, March 2009, <<https://www.rfc-editor.org/info/rfc5440>>.
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