CCNx Key Exchange

IETF 94 - Yokohama - ICNRG Christopher A. Wood November 5, 2015

Motivation and Goals

Motivation

- We need a way to establish session keys between consumers and producers that makes use of CCN properties
 - (D)TLS, QUIC, etc. are a good start

Requirements

- Session keys must be forward secure
 - Compromising long-term secrets does not put session keys at risk
- At most 2 RTTs to establish a session key, with the possibility for session resumption in 0 RTT
- Allow extensions for client authentication in addition to server authentication

TLS and QUIC Overview

- Support 0-, 1-, and 2-RTT forward secure key derivation
 - Long-term public key shares enable faster handshakes
- Different keys are used to encrypt (and MAC) different parts of the protocol
 - A short-term ephemeral key is used for exchanging random key shares to derive a master key
- Server is authenticated to the client
- Prevents address spoofing (via SYN cookies) and replay attacks (via QUIC Source Address Tokens and TLS nonce)

CCNx Key Exchange (CCNx-KE)

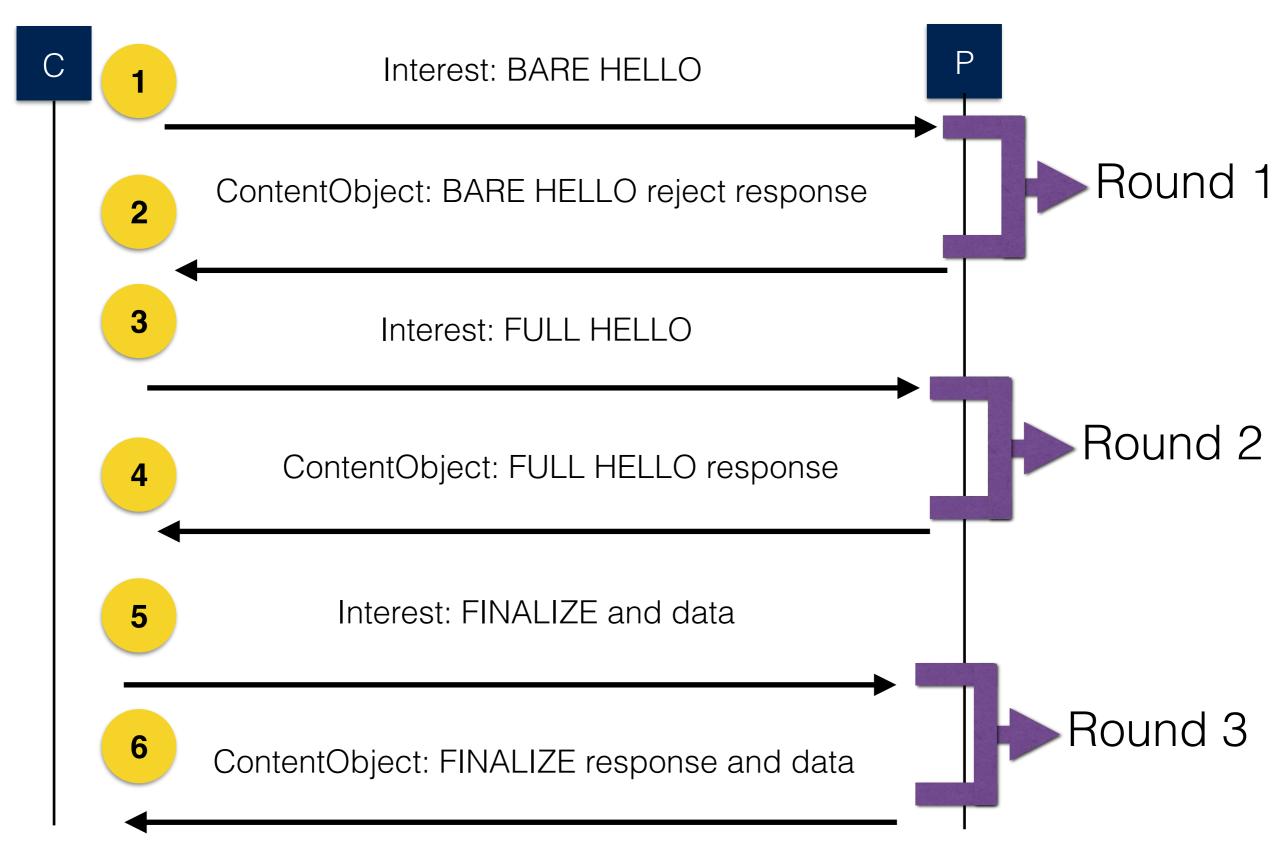
Assumptions

- Consumers know the prefix of the target producer, e.g., /prefix/
- Consumers possess the appropriate trust anchors to authenticate the server
- ... that's it

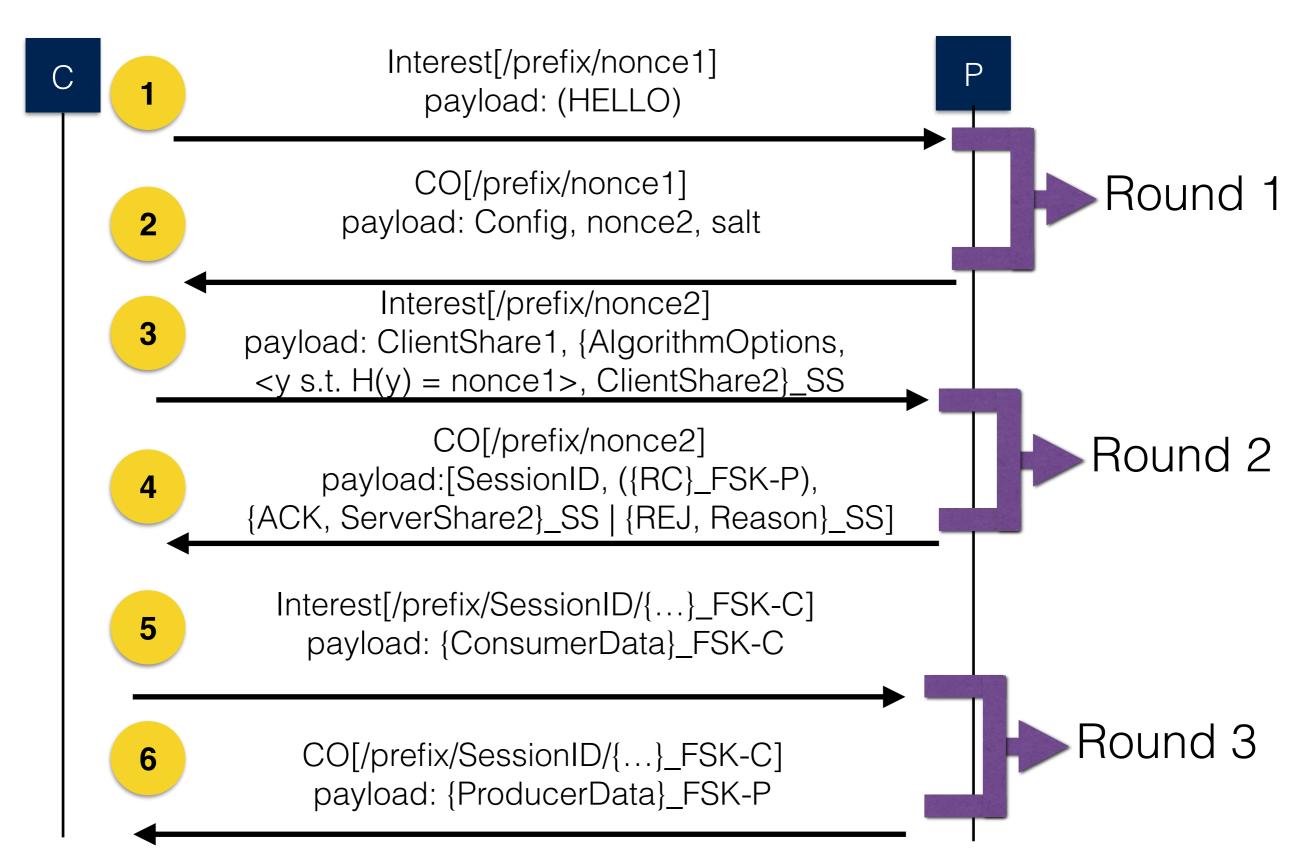
Protocol Overview

- Round 1: Obtain the server config (if not available or it has expired)
- Round 2: FULL HELLO handshake and establish ephemeral keys
- Round 3: Final exchange to derive forward-secure secrets for all subsequent communication

Sketch of the Full Protocol



Sketch of the Full Protocol



SS Derivation

```
SS = HKDF(Salt, IKM) Salt = CSALT1||PSALT1|| ``ss generation" \\ IKM = 32 - byte key-exchange output
```

FSK-C/P Derivation

Second key exchange uses the ServerShare2 and ClientShare2 inputs

$$FSK = HKDF(Salt, IKM)$$

Salt = CSALT2||PSALT2||"fsk generation"

IKM = Second 32-byte key-exchange output

FSP-C/P and IVs are pumped from FSK in the following order:

- 1. FSK-C
- 2. FSK-P
- 3. FSK-CIV (client IV)
- 4. FSK-PIV (producer IV)

SessionID and RC Properties

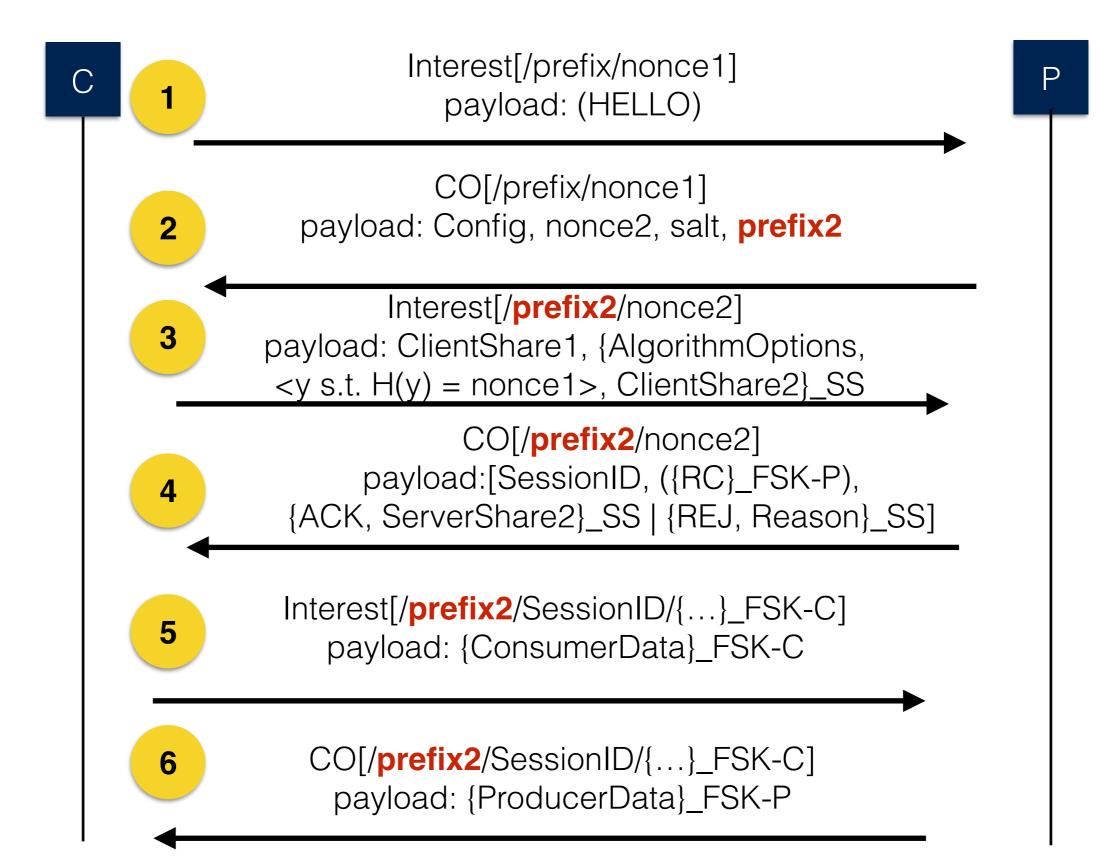
SessionID

- Used to uniquely identifies a single session
- ... a random string/number suffices

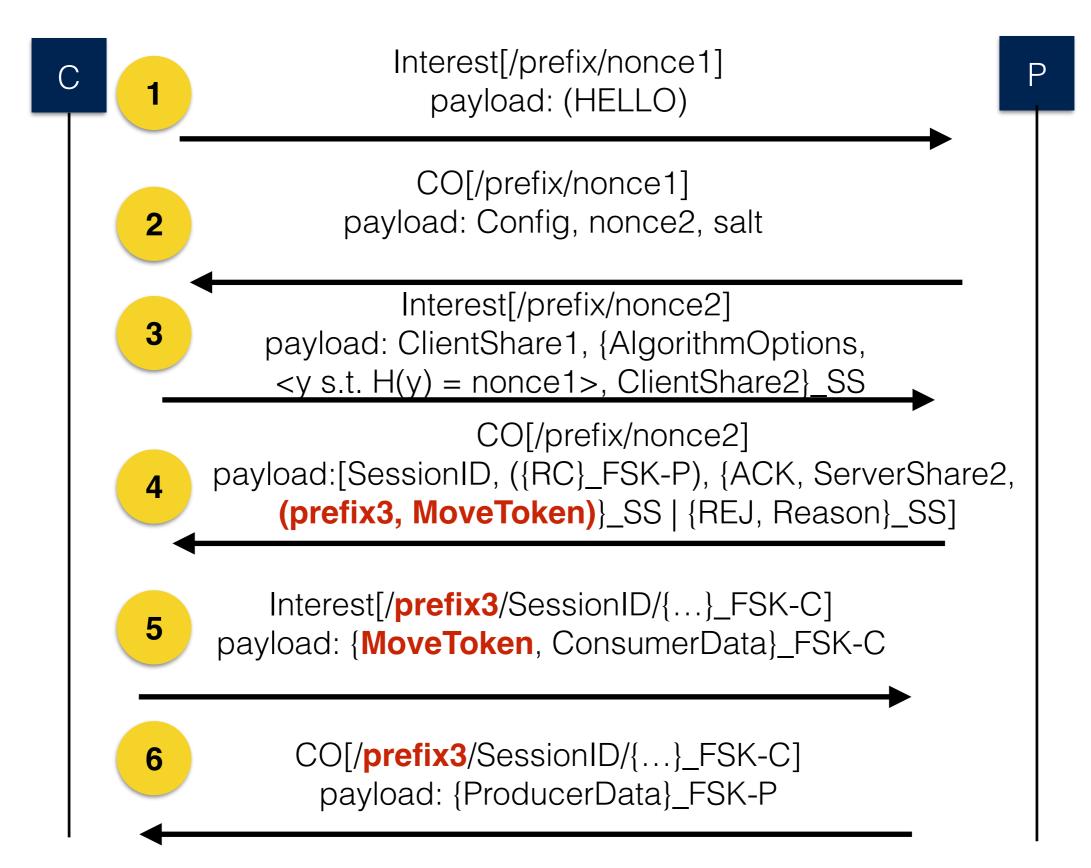
RC

Used to recover SS and FSK for a given session

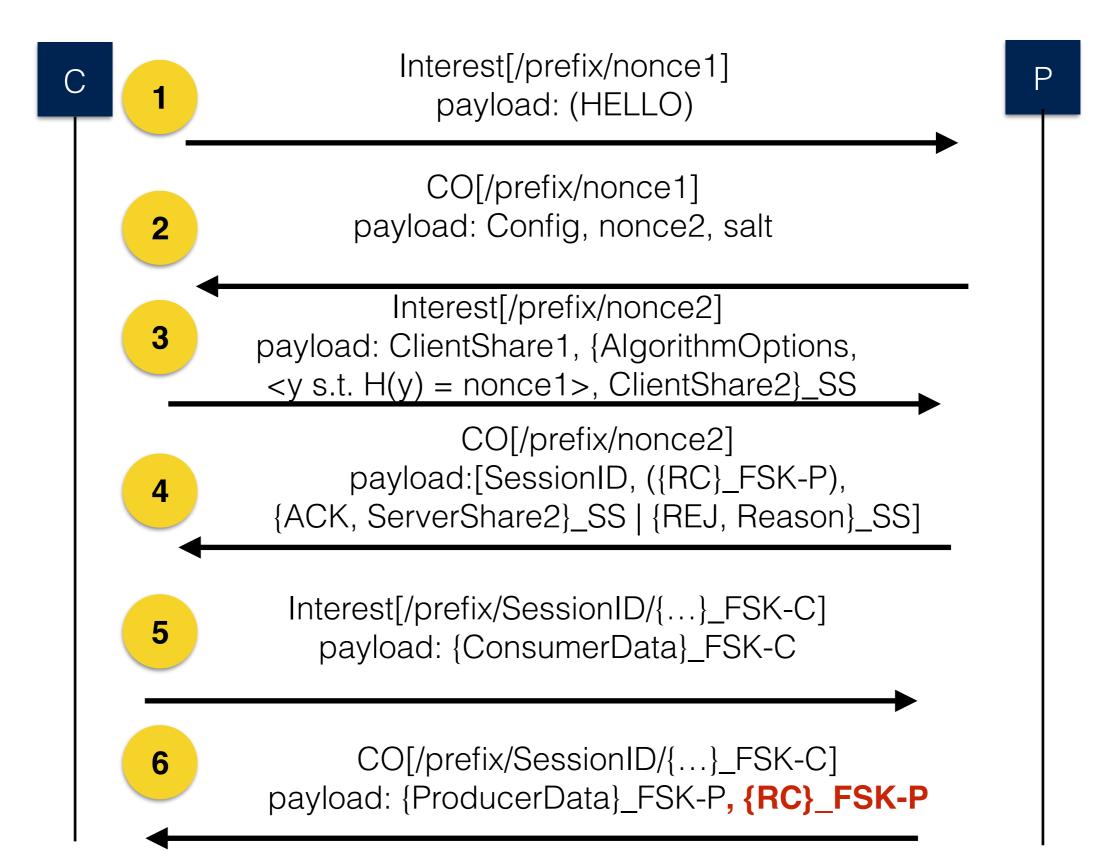
Option #1: HELLO prefix redirection



Option #2: Final prefix redirection



Option #3: Resumption Cookie Echo



CCNx-KE Properties

- Minimal deviation from TLS and QUIC.
- Forward-secure session keys derived similar to TLS and QUIC.
- Server-to-client authentication.
 - Client-to-server authentication is future work.
- Clients are securely bound to the protocol execution (via the hash-based tokens).
- Session state can be securely migrated from the producer to a trusted party.

Backup

SessionID*

Structure: Generated as encryption of the hash digest of a sever secret, FSK, and optional prefix (e.g., Prefix3). Encryption happens with a long-term, private key held by the server.

SessionID = $Enc(k, H(secret||FSK||(Prefix3| \perp)))$

Usage: Append to service prefix (in the name) to indicate what key is used for encrypting payload data

*** This is only one way to create the SessionID

Resumption Cookie (RC)*

Structure: Encryption of H(server secret), SS, FSK, and the (Prefix3, MoveToken) tuple (if provided), with a producer secret key that is also known to the service operating under Prefix3 (if provided)

$$RC = Enc(k, SS||FSK||((Prefix3||MoveToken)| \perp))$$

Usage: The SessionID and RC are needed to resume a session (i.e., recompute SessionID and check for equality):

$$(SS||FSK||((Prefix3||MoveToken)|\perp)) = Dec(RC)$$

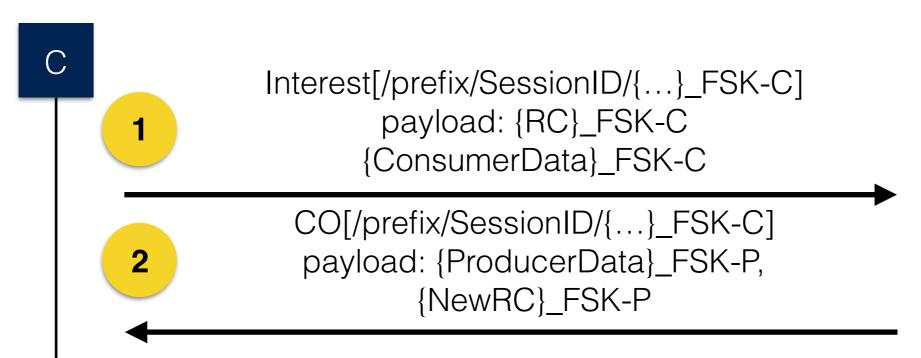
SessionID =
$$?$$
Enc $(k, H(secret||FSK||(Prefix3| $\perp)))$$

*** This is only one way to create the RC

Session Resumption

- Approach 0: If client has nothing, start with HELLO
 [2 RTT delay]
- Approach 1: If the client already has the config, start at the second step [1 RTT delay]
- Approach 2: If the client already has the SessionID and the ResumptionCookie, provide both to resume sessions after long periods of inactivity (requires producer state) [0 RTT delay]

Session Resumption (cont'd)



Note: SessionID is used to verify ownership of the SessionID, the MSK-C encryption key, and the real RC.