

TICTOC Working Group
Internet-Draft
Intended status: Experimental
Expires: April 18, 2016

S. Davari
A. Oren
Broadcom Corp.
M. Bhatia
P. Roberts
Alcatel-Lucent
L. Montini
Cisco Systems
October 16, 2015

Transporting Timing messages over MPLS Networks
draft-ietf-tictoc-1588overmpls-07

Abstract

This document defines a method for transporting timing messages, such as Precision Time Protocol (PTP) or Network Time Protocol (NTP), over a Multiprotocol Label Switched (MPLS) network. The method facilitates efficient recognition of timing packets to enable their port level processing in both Label Edge Routers (LERs) and Label Switched Routers (LSRs).

The basic mechanism is to transport timing messages inside "Timing LSPs", which are dedicated MPLS Label Switched Paths (LSPs) that carry only timing, and possibly related Operations, Administration and Maintenance (OAM) or management packets, but do not carry customer traffic.

Two encapsulations methods are defined. The first transports UDP/IP encapsulated timing messages directly over the dedicated LSP. The second transports Ethernet encapsulated timing messages inside an Ethernet pseudowire.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF). Note that other groups may also distribute working documents as Internet-Drafts. The list of current Internet-Drafts is at <http://datatracker.ietf.org/drafts/current/>.

Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as "work in progress."

This Internet-Draft will expire on April 18, 2016.

Copyright Notice

Copyright (c) 2015 IETF Trust and the persons identified as the document authors. All rights reserved.

This document is subject to BCP 78 and the IETF Trust's Legal Provisions Relating to IETF Documents (<http://trustee.ietf.org/license-info>) in effect on the date of publication of this document. Please review these documents carefully, as they describe your rights and restrictions with respect to this document. Code Components extracted from this document must include Simplified BSD License text as described in Section 4.e of the Trust Legal Provisions and are provided without warranty as described in the Simplified BSD License.

Table of Contents

1. Introduction	3
2. Terminology	4
3. Problem Statement	5
4. Timing over MPLS Architecture	5
5. Dedicated LSPs for Timing messages	7
6. Timing over LSP Encapsulation	8
6.1. Timing over UDP/IP over MPLS Encapsulation	8
6.2. Timing over PW Encapsulation	8
7. Timing message Processing	9
8. Protection and Redundancy	10
9. ECMP and Entropy	10
10. PHP	11
11. OAM, Control and Management	11
12. QoS Considerations	11
13. FCS and Checksum Recalculation	11
14. Behavior of LER/LSRs	12
14.1. Behavior of Timing-capable/aware LERs/LSRs	12
14.2. Behavior of non-Timing-capable/aware LSR	12
15. Other considerations	13
16. Security Considerations	13
17. Applicability Statement	14
18. Acknowledgements	14
19. IANA Considerations	14
20. References	15
20.1. Normative References	15
20.2. Informative References	16
Appendix A. Appendix	17
A.1. Routing extensions for Timing-aware Routers	17
A.2. Signaling Extensions for Creating Timing LSPs	17

Authors' Addresses 18

1. Introduction

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC2119 [RFC2119].

When used in lower case, these words convey their typical use in common language, and are not to be interpreted as described in RFC2119 [RFC2119].

The objective of timing distribution protocols, such as Precision Time Protocol (PTP) and Network Timing Protocol (NTP), is to synchronize clocks running on nodes of a distributed system.

Timing distribution protocols are presently transported over IP or Ethernet. The present document presents a mechanism for transport over Multiprotocol Label Switched (MPLS) networks. Our solution involves transporting timing messages over dedicated "Timing Label Switched Paths (LSPs)". These are ordinary LSPs that carry timing messages and MAY carry Operations, Administration and Maintenance (OAM) or management messages, but do not carry any other traffic.

Timing LSPs may be established statically or via signaling. When using signaling, extensions to routing protocols (e.g., OSPF, ISIS) are required to enable routers to distribute their timing processing capabilities, and extensions to path set up protocols (e.g., RSVP-TE) are required for establishing the LSPs. All such extensions are beyond the scope of this document.

High accuracy timing distribution requires on-path support, e.g., Transparent Clocks (TCs) or Boundary Clocks (BCs), at intermediate nodes. These intermediate nodes need to recognize and appropriately process timing distribution packets. To facilitate efficient recognition of timing messages transported over MPLS, this document restricts the specific encapsulations to be used.

[IEEE-1588] defines PTP messages for frequency, phase and time synchronization. PTP messages may be transported over UDP/IP (Annex D and E of [IEEE-1588]) or over Ethernet (Annex F of [IEEE-1588]). This document defines two methods to transport PTP messages over MPLS networks.

PTP defines several clock types, including ordinary clocks, boundary clocks, end-to-end transparent clocks, and peer-to-peer transparent clocks. Transparent clocks are situated at intermediate nodes and

update the Correction Field inside PTP messages in order to reflect the time required to transit the node.

[RFC5905] defines NTP messages for clock and time synchronization. NTP messages are transported over UDP/IP. This document defines a method to transport NTP messages over MPLS networks.

It can be expected that only a subset of LSR ports will be capable of processing timing messages. Timing LSPs MUST be set up (either by manual provisioning or via signaling) to traverse these ports. While Timing LSPs are designed to optimize timing distribution, the performance of slave clocks is beyond the scope of this document.

Presently on-path support is only defined for PTP, and therefore much of our discussion will focus on PTP. NTP timing distribution may benefit from transport in a Timing LSP due to prioritization or selection of ports or nodes with minimal delay or delay asymmetry.

2. Terminology

1588: The timing distribution protocol defined in IEEE 1588.

Boundary Clock: A device with one timing port to receive timing messages and at least one port to re-distribute timing messages.

CF: Correction Field, a field inside certain PTP messages that holds the accumulated transit time.

Master Clock: The source of 1588 timing messages to a set of slave clocks.

NTP: The timing distribution protocol defined in RFC 5905.

Ordinary Clock: A master or slave clock. Note that ordinary clocks have only a single PTP port.

PTP: Precision Time Protocol. See 1588.

Slave Clock: A receiver of 1588 timing messages from a master clock.

Timing LSP: An MPLS LSP dedicated to carry timing messages.

Timing messages: Timing distribution protocol messages that are exchanged between clocks.

Timing port: A port on a (master, slave, transparent, or boundary) clock.

Timing PW: A PW within a Timing LSP that is dedicated to carry timing messages.

Transparent Clock: An intermediate node that forwards timing messages while updating their CF.

3. Problem Statement

[IEEE-1588] defines methods for transporting PTP messages over Ethernet and IP networks. [RFC5905] defines a method of transporting NTP messages over IP networks. There is a need to transport timing messages over MPLS networks while supporting the Transparent Clock (TC), Boundary Clock (BC) and Ordinary Clock (OC) functionalities in LER and LSRs of the MPLS network.

There are potentially many ways of transporting timing packets over MPLS. However, it is advisable to limit the number of possible encapsulation options to simplify recognition and processing of timing packets.

The solution herein described transports timing messages over dedicated "Timing Label Switched Paths (LSPs)". Were timing packets to share LSPs with other traffic, intermediate LSRs would be required to perform some deeper inspection to differentiate between timing packets and other packets. The method herein proposed avoids this complexity, and can readily detect all PTP messages (one-step or two-step), and supports ordinary, boundary and transparent clocks.

4. Timing over MPLS Architecture

Timing messages are exchanged between timing ports on ordinary and boundary clocks. Boundary clocks terminate the timing messages and act as master clock for other boundary clocks or slave clocks. End-to-End transparent clocks do not terminate the timing messages but do modify the contents of the timing messages in transit.

OC, BC and TC functionality may be implemented in either LERs or LSRs.

An example is shown in Figure 1, where the LERs act as OCs and are the initiating/terminating points for timing messages. The ingress LER encapsulates timing messages in a Timing LSP and the egress LER terminates this Timing LSP. Intermediate LSRs (only one is shown here) act as TCs, updating the CF of transiting timing messages, as well as performing label switching operations.

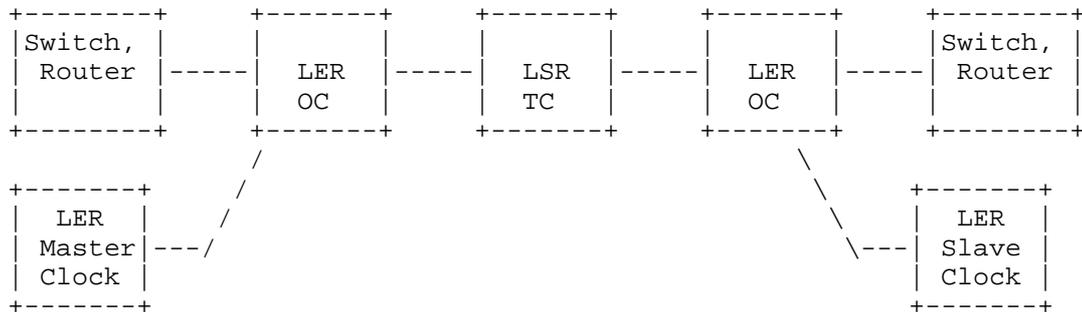


Figure (1) - Deployment example 1 of timing over MPLS network

Another example is shown in Figure 2, where LERs act as BCs, and switches/routers outside of the MPLS network, act as OCs or BCs. The ingress LER BC recovers timing and initiates timing messages encapsulated in the Timing LSP toward the MPLS network, an intermediate LSR acts as a TC, and the egress LER acts as a BC sending timing messages to equipment outside the MPLS network.

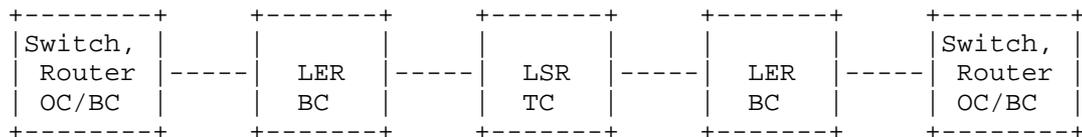


Figure (2) - Deployment example 2 of timing over MPLS network

Yet another example is shown in Figure 3, where both LERs and LSRs act as TCs. The ingress LER updates the CF and encapsulates the timing message in an MPLS packet, intermediate LSRs update the CF and perform label switching, and the egress LER updates the CF and sends the timing messages to equipment outside the MPLS network.

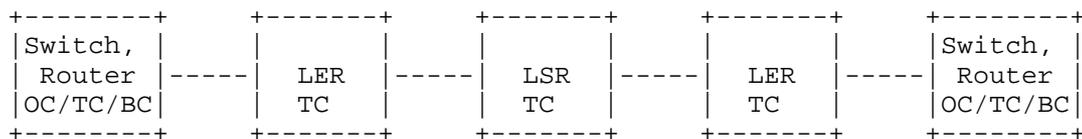


Figure (3) - Deployment example 3 of timing over MPLS network

A final example is shown in Figure 4, where all nodes act as BCs. Single-hop LSPs are created between every two adjacent LSRs. Of course, PTP transport over Ethernet MAY be used between two network elements.

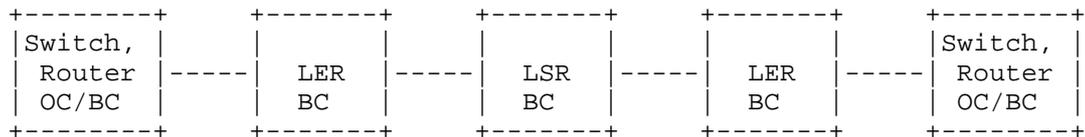


Figure (4) - Deployment example 3 of timing over MPLS network

An MPLS domain MAY serve multiple customers, each having its own Timing domain. In these cases the MPLS domain (maintained by a service provider) MUST provide dedicated timing services to each customer.

The timing over MPLS architecture assumes a full mesh of Timing LSPs between all LERs supporting this specification. It supports point-to-point (VPWS) and Multipoint (VPLS) services. This means that a customer may purchase a point-to-point timing service between two customer sites or a multipoint timing service between more than two customer sites.

The Timing over MPLS architecture supports P2P or P2MP Timing LSPs. This means that the Timing Multicast messages such as PTP Multicast event messages MAY be transported over P2MP Timing LSPs or MAY be replicated and transported over multiple P2P Timing LSPs.

Timing LSPs, as defined by this specification, MAY be used for timing messages that do not require time-stamping or CF updating.

PTP Announce messages that determine the Timing LSP terminating point behavior such as BC/OC/TC SHOULD be transported over the Timing LSP to simplify hardware and software.

5. Dedicated LSPs for Timing messages

The method defined in this document is used by LER and LSRs to identify timing messages by observing the top label of the MPLS label stack. Compliant implementations MUST use dedicated LSPs to carry timing messages over MPLS. Such LSPs are herein referred to as "Timing LSPs" and the labels associated with these LSPs as "Timing LSP labels".

Timing distribution requires symmetrical bidirectional communications. Co-routing of the two directions is required to limit delay asymmetry. Thus timing messages **MUST** be transported either over two co-routed unidirectional Timing LSPs, or a single bidirectional co-routed Timing LSP.

Timing LSPs **MAY** be configured using RSVP-TE. Extensions to RSVP-TE are required for this purpose, but are beyond the scope of this document.

6. Timing over LSP Encapsulation

We define two methods for carrying timing messages over MPLS. The first method transports UDP/IP-encapsulated timing messages over Timing LSPs, and the second method transports Ethernet encapsulated timing messages over Ethernet PWs placed in Timing LSPs.

6.1. Timing over UDP/IP over MPLS Encapsulation

The first method directly encapsulates UDP/IP timing messages in a Timing LSP. The UDP/IP encapsulation of PTP messages **MUST** comply to Annex D and E of [IEEE-1588], and the UDP/IP encapsulation of NTP messages **MUST** comply to [RFC5905]. This format is shown in Figure 4.

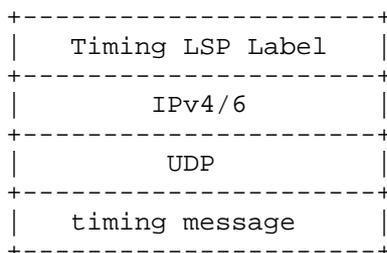


Figure (4) - Timing over UDP/IP over MPLS Encapsulation

In order for an LER/LSR to process timing messages, the Timing LSP Label must be the top label of the label stack. The LER/LSR **MUST** know that this label is a Timing LSP Label. It can learn this by static configuration or via RSVP-TE signaling.

6.2. Timing over PW Encapsulation

Another method of transporting timing over MPLS networks is to use Ethernet encapsulated timing messages, and to transport these in an Ethernet PW which in turn is transported over a Timing LSP. In the

case of PTP, the Ethernet encapsulation MUST comply to Annex F of [IEEE-1588] and the Ethernet PW encapsulation to [RFC4448], resulting in the format shown in Figure 5(A).

Either the Raw mode or Tagged mode defined in [RFC-4448] MAY be used and the payload MAY have 0, 1, or 2 VLAN tags. The Timing over PW encapsulation MUST use the Control Word (CW) as specified in [RFC4448]. The use of Sequence Number in the CW is optional.

NTP MAY be transported using an IP PW (as defined in [RFC4447]) as shown in Fig 5(B).

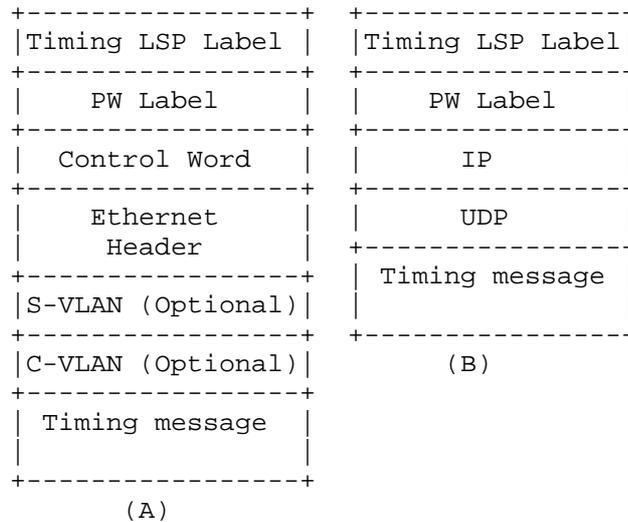


Figure (5) - Timing over PW Encapsulations

7. Timing message Processing

Each Timing protocol such as PTP and NTP, defines a set of timing messages. PTP defines SYNC, DELAY_REQ, DELAY_RESP, FOLLOW_UP, etc.

Some timing messages require per-packet processing, such as time-stamping or CF updating. A compliant LER/LSR parses each timing message to determine the required processing.

For example, the following PTP messages (event messages) require time-stamping or CF updating:

- o SYNC

- o DELAY_REQ (Delay Request)
- o PDELAY_REQ (Peer Delay Request)
- o PDELAY_RESP (Peer Delay Response)

SYNC and DELAY_REQ are exchanged between a Master Clock and a Slave Clock and MUST be transported over Timing LSPs. PDELAY_REQ and PDELAY_RESP are exchanged between adjacent PTP clocks (master, slave, boundary, or transparent) and SHOULD be transported over single hop Timing LSPs. If two-Step PTP clocks are present, then the FOLLOW_UP, and PDELAY_RESP_FOLLOW_UP messages MUST also be transported over Timing LSPs.

For a given instance of the 1588 protocol, SYNC and DELAY_REQ MUST be transported in opposite directions. As aforementioned, two co-routed unidirectional LSPs or a single bidirectional co-routed LSP MAY be used.

Except as indicated above for two-step PTP clocks, PTP messages that are not "event messages" need not be processed by intermediate routers. These message types MAY be carried in PTP Tunnel LSPs.

8. Protection and Redundancy

In order to ensure continuous uninterrupted operation of timing distribution, slave clocks often track redundant master clocks. Prolonged outages of Timing LSPs trigger switching to a redundant master clock. It is the responsibility of the network operator to ensure that physically disjoint Timing LSPs are established between a slave clock and redundant master clocks.

LSP or PW layer protection, such as linear protection Switching, ring protection switching or MPLS Fast Reroute (FRR), will lead to changes in propagation delay between master and slave clocks. Such a change, if undetected by the slave clock, would negatively impact timing performance. While it is expected that slave clocks will often be able to detect such delay changes, this specification RECOMMENDS that automatic protection switching NOT be used for Timing LSPs, unless the operator can ensure that it will not negatively impact timing performance.

9. ECMP and Entropy

To ensure the correct operation of slave clocks and avoid error introduced by forward and reverse path delay asymmetry, the physical path taken by timing messages MUST be the same for all timing

messages. In particular, the PTP event messages listed in section 7 MUST be routed in the same way.

Therefore the Timing LSPs MUST not be subject to ECMP (Equal Cost Multipath). Entropy labels MUST NOT be used for the Timing LSP [RFC6790] and MUST NOT be used for PWs inside the Timing LSP [RFC6391].

10. PHP

To ensure that the label on the top of the label stack is the Timing LSP Label, PHP MUST not be employed.

11. OAM, Control and Management

In order to monitor Timing LSPs or PWs, it is necessary to enable them to carry OAM messages. OAM packets MUST be differentiated from timing messages by already defined IETF methods.

For example BFD [RFC5880], [RFC5884] and LSP-Ping [RFC4389] MAY run over Timing LSPs via UDP/IP encapsulation or via GAL/G-ACh. These protocols can easily be identified by the UDP Destination port number or by GAL/G-ACh respectively.

Also BFD, LSP-Ping and other messages MAY run over Timing PWs via VCCV [RFC5085]. In this case these messages are recognized according to the VCCV type.

12. QoS Considerations

There may be deployments where timing messages traverse LSR/LEs that are not capable of the required processing. In order to minimize the negative impact on the timing performance of the slave clock timing messages MUST be treated with the highest priority. This can be achieved by proper setup of Timing LSPs.

It is recommended that Timing LSPs be configured to indicate EF-PHB [RFC3246] for the CoS and "green" [RFC2697] for drop eligibility.

13. FCS and Checksum Recalculation

Since Boundary and Transparent Clocks modify packets, when the MPLS packets are transported over Ethernet the processing MUST include recalculation of the Ethernet FCS. FCS retention as described in [RFC4720] MUST NOT be used.

For the UDP/IP encapsulation mode, calculation of the UDP checksum will generally be required. After updating the CF a Transparent

Clock MUST either incrementally update the UDP checksum or completely recalculate the checksum before transmission to downstream node.

14. Behavior of LER/LSRs

Timing-aware LERs or LSRs are MPLS routers that are able to recognize timing packets. Timing-capable LERs and LSRs further have one or more interfaces that can perform timing processing (OC/BC/TC) on timing packets. Timing-capable/aware LERs and LSRs MAY advertise the timing capabilities of their interfaces via control plane protocols such as OSPF or IS-IS, and timing-aware LERs can then be set up Timing LSPs via RSVP-TE signaling. Alternatively the timing capabilities of LERs and LSRs may be known by a centralized controller or management system, and Timing LSPs may be manually configured, or set up by a management platform or a Software Defined Networking (SDN) controller.

14.1. Behavior of Timing-capable/aware LERs/LSRs

When a timing-capable ingress LER acting as a TC receives a timing message packet from a timing-capable non-MPLS interface, the LER updates the CF, encapsulates and forwards the packet over a previously established Timing LSP. When a timing-capable egress LER acting as a TC receives a timing message packet on timing-capable MPLS interface, the LER updates the CF, decapsulates the MPLS encapsulation, and forwards the packet via a non-MPLS interface. When a timing-capable LSR acting as a TC receives a timing message from a timing-capable MPLS interface, the LSR updates the CF and forwards the timing message over another MPLS interface.

When a timing-capable LER acting as a BC receives a timing message packet from a timing-capable interface, the LER time-stamps the packet and sends it to the BC processing module.

When a timing-capable LER acting as an OC receives a timing message from a timing-capable MPLS interface, the LER time-stamps the packet and sends it to the OC processing module.

14.2. Behavior of non-Timing-capable/aware LSR

It is most beneficial when all LSRs in the path of a Timing LSP be timing-Capable/aware LSRs. This would ensure the highest quality time and clock synchronization by slave clocks. However, this specification does not mandate that all LSRs in path of a Timing LSP be timing-capable/aware.

Non-timing-capable/aware LSRs just perform label switching on the packets encapsulated in Timing LSPs and don't perform any timing

related processing. However, as explained in QoS section, timing packets MUST be still be treated with the highest priority based on their Traffic Class marking.

15. Other considerations

[IEEE-1588] defines an optional peer-to-peer transparent clocking (P2P TC) mode that compensates both for residence time in the network node and for propagation time on the link between nodes. To support P2P TC, delay measurement must be performed between two adjacent timing-capable/aware LSRs. Thus, in addition to the TC functionality detailed above on transit PTP timing messages, adjacent peer to peer TCs MUST engage in single-hop peer delay measurement.

For single hop peer delay measurement a single-hop LSP SHOULD be created between the two adjacent LSRs. Other methods MAY be used; for example, if the link between the two adjacent routers is Ethernet, PTP transport over Ethernet MAY be used.

To support P2P TC, a timing-capable/aware LSR MUST maintain a list of all neighbors to which it needs to send a PDelay_Req, and maintain a single-hop timing LSP to each.

The use of Explicit Null Label (label 0 or 2) is acceptable as long as either the Explicit Null label is the bottom of stack label (for the UDP/IP encapsulation) or the label below the Explicit Null label (for the PW case).

16. Security Considerations

Security considerations for MPLS and pseudowires are discussed in [RFC3985] and [RFC4447]. Security considerations for timing are discussed in [RFC7384]. Everything discussed in those documents applies to the Timing LSP of this document.

An experimental security protocol is defined in [IEEE-1588]. The PTP security extension and protocol provides group source authentication, message integrity, and replay attack protection for PTP messages.

When the MPLS network (provider network) serves multiple customers, it is important to distinguish between timing messages belonging to different customers. For example if an LER BC is synchronized to a grandmaster belonging to customer A, then the LER MUST only use that BC for slaves of customer A, to ensure that customer A cannot adversely affect the timing distribution of other customers.

Timing messages MAY be encrypted or authenticated, provided that the timing-capable LERs/LSRs can authenticate/ decrypt the timing messages.

17. Applicability Statement

The Timing over MPLS transport methods described in this document apply to the following network Elements:

- o An ingress LER that receives IP or Ethernet encapsulated timing messages from a non-MPLS interface and forwards them as MPLS encapsulated timing messages over Timing LSP, optionally performing TC functionality.
- o An egress LER that receives MPLS encapsulated timing messages from a Timing LSP and forwards them to non-MPLS interface as IP or Ethernet encapsulated timing messages, optionally performing TC functionality.
- o An ingress LER that receives MPLS encapsulated timing messages from a non-MPLS interface, performs BC functionality, and sends timing messages over a Timing LSP.
- o An egress LER that receives MPLS encapsulated timing messages from a Timing LSP, performs BC functionality, and sends timing messages over a non-MPLS interface.
- o An LSR on a Timing LSP that receives MPLS encapsulated timing messages from one MPLS interface and forwards them to another MPLS interface, optionally performing TC functionality.

This document also supports the case where not all LSRs are timing-capable/aware, or not all LER/LSR interfaces are timing-capable/aware.

18. Acknowledgements

The authors would like to thank Yaakov Stein, Luca Martini, Ron Cohen, Tal Mizrahi, Stefano Ruffini, Peter Meyer and other IETF participants for reviewing and providing feedback on this draft.

19. IANA Considerations

There are no IANA requirements in this specification.

20. References

20.1. Normative References

- [IEEE-1588] IEEE 1588-2008, "IEEE Standard for a Precision Clock Synchronization Protocol for Networked Measurement and Control Systems", July 2008.
- [RFC2119] Bradner, S., "Key words for use in RFCs to Indicate Requirement Levels", BCP 14, RFC 2119, DOI 10.17487/RFC2119, March 1997, <<http://www.rfc-editor.org/info/rfc2119>>.
- [RFC3985] Bryant, S., Ed. and P. Pate, Ed., "Pseudo Wire Emulation Edge-to-Edge (PWE3) Architecture", RFC 3985, DOI 10.17487/RFC3985, March 2005, <<http://www.rfc-editor.org/info/rfc3985>>.
- [RFC4389] Thaler, D., Talwar, M., and C. Patel, "Neighbor Discovery Proxies (ND Proxy)", RFC 4389, DOI 10.17487/RFC4389, April 2006, <<http://www.rfc-editor.org/info/rfc4389>>.
- [RFC4447] Martini, L., Ed., Rosen, E., El-Aawar, N., Smith, T., and G. Heron, "Pseudowire Setup and Maintenance Using the Label Distribution Protocol (LDP)", RFC 4447, DOI 10.17487/RFC4447, April 2006, <<http://www.rfc-editor.org/info/rfc4447>>.
- [RFC4448] Martini, L., Ed., Rosen, E., El-Aawar, N., and G. Heron, "Encapsulation Methods for Transport of Ethernet over MPLS Networks", RFC 4448, DOI 10.17487/RFC4448, April 2006, <<http://www.rfc-editor.org/info/rfc4448>>.
- [RFC4720] Malis, A., Allan, D., and N. Del Regno, "Pseudowire Emulation Edge-to-Edge (PWE3) Frame Check Sequence Retention", RFC 4720, DOI 10.17487/RFC4720, November 2006, <<http://www.rfc-editor.org/info/rfc4720>>.
- [RFC5085] Nadeau, T., Ed. and C. Pignataro, Ed., "Pseudowire Virtual Circuit Connectivity Verification (VCCV): A Control Channel for Pseudowires", RFC 5085, DOI 10.17487/RFC5085, December 2007, <<http://www.rfc-editor.org/info/rfc5085>>.
- [RFC5880] Katz, D. and D. Ward, "Bidirectional Forwarding Detection (BFD)", RFC 5880, DOI 10.17487/RFC5880, June 2010, <<http://www.rfc-editor.org/info/rfc5880>>.

- [RFC5884] Aggarwal, R., Kompella, K., Nadeau, T., and G. Swallow, "Bidirectional Forwarding Detection (BFD) for MPLS Label Switched Paths (LSPs)", RFC 5884, DOI 10.17487/RFC5884, June 2010, <<http://www.rfc-editor.org/info/rfc5884>>.

20.2. Informative References

- [ISO] ISO/IEC 10589:1992, "Intermediate system to Intermediate system routing information exchange protocol for use in conjunction with the Protocol for providing the Connectionless-mode Network Service (ISO 8473)", April 1992.
- [RFC1195] Callon, R., "Use of OSI IS-IS for routing in TCP/IP and dual environments", RFC 1195, DOI 10.17487/RFC1195, December 1990, <<http://www.rfc-editor.org/info/rfc1195>>.
- [RFC2328] Moy, J., "OSPF Version 2", STD 54, RFC 2328, DOI 10.17487/RFC2328, April 1998, <<http://www.rfc-editor.org/info/rfc2328>>.
- [RFC2697] Heinanen, J. and R. Guerin, "A Single Rate Three Color Marker", RFC 2697, DOI 10.17487/RFC2697, September 1999, <<http://www.rfc-editor.org/info/rfc2697>>.
- [RFC3246] Davie, B., Charny, A., Bennet, J., Benson, K., Le Boudec, J., Courtney, W., Davari, S., Firoiu, V., and D. Stiliadis, "An Expedited Forwarding PHB (Per-Hop Behavior)", RFC 3246, DOI 10.17487/RFC3246, March 2002, <<http://www.rfc-editor.org/info/rfc3246>>.
- [RFC5340] Coltun, R., Ferguson, D., Moy, J., and A. Lindem, "OSPF for IPv6", RFC 5340, DOI 10.17487/RFC5340, July 2008, <<http://www.rfc-editor.org/info/rfc5340>>.
- [RFC5905] Mills, D., Martin, J., Ed., Burbank, J., and W. Kasch, "Network Time Protocol Version 4: Protocol and Algorithms Specification", RFC 5905, DOI 10.17487/RFC5905, June 2010, <<http://www.rfc-editor.org/info/rfc5905>>.
- [RFC6391] Bryant, S., Ed., Filsfils, C., Drafz, U., Kompella, V., Regan, J., and S. Amante, "Flow-Aware Transport of Pseudowires over an MPLS Packet Switched Network", RFC 6391, DOI 10.17487/RFC6391, November 2011, <<http://www.rfc-editor.org/info/rfc6391>>.

[RFC6790] Kompella, K., Drake, J., Amante, S., Henderickx, W., and L. Yong, "The Use of Entropy Labels in MPLS Forwarding", RFC 6790, DOI 10.17487/RFC6790, November 2012, <<http://www.rfc-editor.org/info/rfc6790>>.

[RFC7384] Mizrahi, T., "Security Requirements of Time Protocols in Packet Switched Networks", RFC 7384, DOI 10.17487/RFC7384, October 2014, <<http://www.rfc-editor.org/info/rfc7384>>.

Appendix A. Appendix

A.1. Routing extensions for Timing-aware Routers

MPLS-TE routing relies on extensions to OSPF [RFC2328] [RFC5340] and IS-IS [ISO] [RFC1195] in order to advertise Traffic Engineering (TE) link information used for constraint-based routing.

Timing related capabilities, such as the capability for a router to perform time-stamping, and OC, TC or BC processing, need to be advertised in order for them to be taken into account during path computation. A management system or SDN controller cognizant of timing related capabilities, can prefer or even require a Timing LSP to traverse links or nodes or interfaces with the required capabilities. The optimal path will optimize the performance of the slave clock.

Extensions are required to OSPF and IS-IS in order to advertise timing related capabilities of a link. Such extensions are outside the scope of this document; however such extensions SHOULD be able to signal the following information per Router Link:

- o Capable of processing PTP, NTP or other timing flows
- o Capable of performing TC operation
- o Capable of performing BC operation

A.2. Signaling Extensions for Creating Timing LSPs

RSVP-TE signaling MAY be used to set up Timing LSPs. Extensions are required to RSVP-TE for this purpose. Such extensions are outside the scope of this document; however, the following information MAY be included in such extensions:

- o Offset from Bottom of Stack (BoS) to the start of the Time-stamp field
- o Number of VLANs in case of PW encapsulation

- o Time-stamp field Type
 - * Correction Field, time-stamp
- o Time-stamp Field format
 - * 64-bit PTPv1, 80-bit PTPv2, 32-bit NTP, 64-bit NTP, 128-bit NTP, etc.

Note that when the above optional information is signaled with RSVP-TE for a Timing LSP, all the timing packets carried in that LSP must have the same signaled characteristics. For example if time-stamp format is signaled as 64-bit PTPv1, then all timing packets must use 64-bit PTPv1 time-stamp.

Authors' Addresses

Shahram Davari
Broadcom Corp.
San Jose, CA 95134
USA

Email: davari@broadcom.com

Amit Oren
Broadcom Corp.
San Jose, CA 95134
USA

Email: amito@broadcom.com

Manav Bhatia
Alcatel-Lucent
Bangalore
India

Email: manav.bhatia@alcatel-lucent.com

Peter Roberts
Alcatel-Lucent
Kanata
Canada

Email: peter.roberts@alcatel-lucent.com

Laurent Montini
Cisco Systems
San Jose CA
USA

Email: lmontini@cisco.com

Network Working Group
Internet Draft
Intended status: Experimental
Expires: April 2016

Technion - Israel Institute of Technology
A. Shpiner
R. Tse
C. Schelp
PMC-Sierra
T. Mizrahi
Marvell
October 18, 2015

Multi-Path Time Synchronization
draft-ietf-tictoc-multi-path-synchronization-03.txt

Abstract

Clock synchronization protocols are very widely used in IP-based networks. The Network Time Protocol (NTP) has been commonly deployed for many years, and the last few years have seen an increasingly rapid deployment of the Precision Time Protocol (PTP). As time-sensitive applications evolve, clock accuracy requirements are becoming increasingly stringent, requiring the time synchronization protocols to provide high accuracy. Slave Diversity is a recently introduced approach, where the master and slave clocks (also known as server and client) are connected through multiple network paths, and the slave combines the information received through all paths to obtain a higher clock accuracy compared to the conventional one-path approach. This document describes a multi-path approach to PTP and NTP over IP networks, allowing the protocols to run concurrently over multiple communication paths between the master and slave clocks. The multi-path approach can significantly contribute to clock accuracy, security and fault tolerance. The Multi-Path Precision Time Protocol (MPPTP) and Multi-Path Network Time Protocol (MPNTP) define an additional layer that extends the existing PTP and NTP without the need to modify these protocols. MPPTP and MPNTP also allow backward compatibility with nodes that do not support the multi-path extension.

Status of this Memo

This Internet-Draft is submitted to IETF in full conformance with the provisions of BCP 78 and BCP 79.

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF), its areas, and its working groups. Note that other groups may also distribute working documents as Internet-Drafts.

Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as "work in progress."

The list of current Internet-Drafts can be accessed at <http://www.ietf.org/ietf/lid-abstracts.txt>.

The list of Internet-Draft Shadow Directories can be accessed at <http://www.ietf.org/shadow.html>.

This Internet-Draft will expire on April 18, 2016.

Copyright Notice

Copyright (c) 2015 IETF Trust and the persons identified as the document authors. All rights reserved.

This document is subject to BCP 78 and the IETF Trust's Legal Provisions Relating to IETF Documents (<http://trustee.ietf.org/license-info>) in effect on the date of publication of this document. Please review these documents carefully, as they describe your rights and restrictions with respect to this document. Code Components extracted from this document must include Simplified BSD License text as described in Section 4.e of the Trust Legal Provisions and are provided without warranty as described in the Simplified BSD License.

Table of Contents

1. Introduction.....	3
2. Conventions Used in this Document.....	5
2.1. Abbreviations.....	5
2.2. Terminology.....	5
3. Multiple Paths in IP Networks.....	5
3.1. Load Balancing.....	5
3.2. Using Multiple Paths Concurrently.....	6
3.3. Two-Way Paths.....	6
4. Solution Overview.....	6
4.1. Path Configuration and Identification.....	6
4.2. Combining.....	7
5. Multi-Path Time Synchronization Protocols over IP Networks.....	7
5.1. Single-Ended Multi-Path Synchronization.....	8
5.1.1. Single-Ended MPPTP Synchronization Message Exchange..	8
5.1.2. Single-Ended MPNTP Synchronization Message Exchange..	9
5.2. Dual-Ended Multi-Path Synchronization.....	10
5.2.1. Dual-Ended MPPTP Synchronization Message Exchange...	10

5.2.2. Dual-Ended MPNTP Synchronization Message Exchange...	11
5.3. Using Traceroute for Path Discovery.....	12
5.4. Using Unicast Discovery for MPPTP.....	12
6. Combining Algorithm.....	13
7. Security Considerations.....	13
8. IANA Considerations.....	13
9. Acknowledgments.....	13
10. References.....	13
10.1. Normative References.....	13
10.2. Informative References.....	14

1. Introduction

The two most common time synchronization protocols in IP networks are the Network Time Protocol [NTP], and the Precision Time Protocol (PTP), defined in the IEEE 1588 standard [IEEE1588].

The accuracy of the time synchronization protocols directly depends on the stability and the symmetry of propagation delays on both directions between the master and slave clocks. Depending on the nature of the underlying network, time synchronization protocol packets can be subject to variable network latency or path asymmetry (e.g. [ASSYMETRY], [ASSYMETRY2]). As time sensitive applications evolve, accuracy requirements are becoming increasingly stringent.

Using a single network path in a clock synchronization protocol closely ties the slave clock accuracy to the behavior of the specific path, which may suffer from temporal congestion, faults or malicious attacks. Relying on multiple clock servers as in NTP solves these problems, but requires active maintenance of multiple accurate sources in the network, which is not always possible. The usage of Transparent Clocks (TC) in PTP solves the congestion problem by eliminating the queueing time from the delay calculations, but does not address security or fault-tolerance aspects.

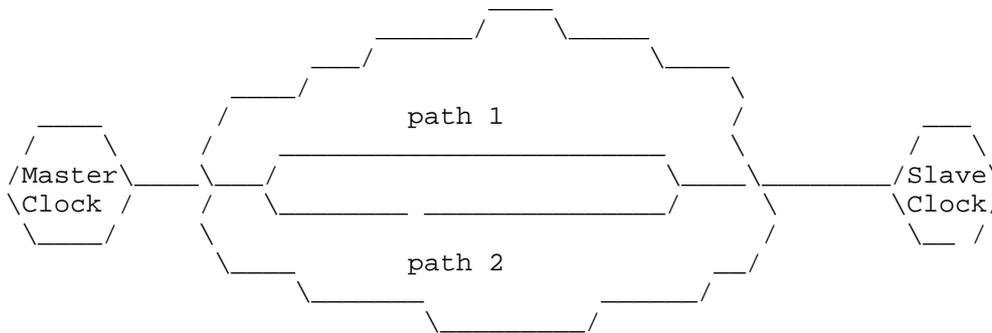


Figure 1 Multi-Path Connection

Since master and slave clocks are often connected through more than one path in the network, as shown in Figure 1, [SLAVEDIV] suggested that a time synchronization protocol can be run over multiple paths, providing several advantages. First, it can significantly increase the clock accuracy as shown in [SLAVEDIV]. Second, this approach provides additional security, allowing to mitigate man-in-the-middle attacks against the time synchronization protocol [DELAY-ATT]. Third, using multiple paths concurrently provides an inherent failure protection mechanism.

This document introduces Multi-Path PTP (MPPTP) and Multi-Path NTP (MPNTP), respectively. These extensions are defined at the network layer and do not require any changes in the PTP or in the NTP protocols.

MPPTP and MPNTP are defined over IP networks. As IP networks typically combine ECMP routing, this property is leveraged for the multiple paths used in MPPTP and MPNTP. The key property of the multi-path extension is that clocks in the network can use more than one IP address. Each {master IP, slave IP} address pair defines a path. Depending on the network topology and configuration, the IP combination pairs can form multiple diverse paths used by the multi-path synchronization protocols. It has been shown [MULTI] that using multiple IP addresses over the wide Internet indeed allows two endpoints to attain multiple diverse paths.

This document introduces two variants for each of the two multi-path protocols; a variant that requires both master and slave nodes to support the multi-path protocol, referred to as the dual-ended

variant, and a backward compatible variant that allows a multi-path clock to connect to a conventional single-path clock, referred to as the single-ended variant.

2. Conventions Used in this Document

2.1. Abbreviations

BMC	Best Master Clock [IEEE1588]
ECMP	Equal Cost Multiple Path
LAN	Local Area Network
MPNTP	Multi-Path Network Time Protocol
MPPTP	Multi-Path Precision Time Protocol
NTP	Network Time Protocol [NTP]
PTP	Precision Time Protocol [IEEE1588]

2.2. Terminology

In the NTP terminology, a time synchronization protocol is run between a client and a server, while PTP uses the terms master and slave. Throughout this document, the sections that refer to both PTP and NTP generically use the terms master and slave.

3. Multiple Paths in IP Networks

3.1. Load Balancing

Traffic sent across IP networks is often load balanced across multiple paths. The load balancing decisions are typically based on packet header fields: source and destination addresses, Layer 4 ports, the Flow Label field in IPv6, etc.

Three common load balancing criteria are per-destination, per-flow and per-packet. The per-destination load balancers take a load balancing decision based on the destination IP address. Per-flow load balancers use various fields in the packet header, e.g., IP addresses and Layer 4 ports, for the load balancing decision. Per-packet load balancers use flow-blind techniques such as round-robin without basing the choice on the packet content.

3.2. Using Multiple Paths Concurrently

To utilize the diverse paths that traverse per-destination load-balancers or per-flow load-balancers, the packet transmitter can vary the IP addresses in the packet header. The analysis in [PARIS2] shows that a significant majority of the flows on the internet traverse per-destination or per-flow load-balancing. It presents statistics that 72% of the flows traverse per-destination load balancing and 39% of the flows traverse per-flow load-balancing, while only a negligible part of the flows traverse per-packet load balancing. These statistics show that the vast majority of the traffic on the internet is load balanced based on packet header fields.

The approaches in this draft are based on varying the source and destination IP addresses in the packet header. Possible extensions have been considered that also vary the UDP ports. However some of the existing implementations of PTP and NTP use fixed UDP port values in both the source and destination UDP port fields, and thus do not allow this approach.

3.3. Two-Way Paths

A key property of IP networks is that packets forwarded from A to B do not necessarily traverse the same path as packets from B to A. Thus, we define a two-way path for a master-slave connection as a pair of one-way paths: the first from master to slave and the second from slave to master.

If possible, a traffic engineering approach can be used to verify that time synchronization traffic is always forwarded through bidirectional two-way paths, i.e., that each two-way path uses the same route on the forward and reverse directions, thus allowing propagation time symmetry. However, in the general case two-way paths do not necessarily use the same path for the forward and reverse directions.

4. Solution Overview

The multi-path time synchronization protocols we present are comprised of two building blocks; one is the path configuration and identification, and the other is the algorithm used by the slave to combine the information received from the various paths.

4.1. Path Configuration and Identification

The master and slave clocks must be able to determine the path of transmitted protocol packets, and to identify the path of incoming

protocol packets. A path is determined by a {master IP, slave IP} address pair. The synchronization protocol message exchange is run independently through each path.

Each IP address pair defines a two-way path, and thus allows the clocks to bind a transmitted packet to a specific path, or to identify the path of an incoming packet.

If possible, the routing tables across the network should be configured with multiple traffic engineered paths between the pair of clocks. By carefully configuring the routers in such networks it is possible to create diverse paths for each of the IP address pairs between two clocks in the network. However, in public and provider networks the load balancing behavior is hidden from the end users. In this case the actual number of paths may be less than the number of IP address pairs, since some of the address pairs may share common paths.

4.2. Combining

Various methods can be used for combining the time information received from the different paths. The output of the combining algorithm is the accurate time offset. Combining methods are further discussed in Section 6.

5. Multi-Path Time Synchronization Protocols over IP Networks

This section presents two variants of MPPTP and MPNTP; single-ended multi-path time synchronization and dual-ended multi-path time synchronization. In the first variant, the multi-path protocol is run only by the slave and the master is not aware of its usage. In the second variant, all clocks must support the multi-path protocol.

The dual-ended protocol provides higher path diversity by using multiple IP addresses at both ends, the master and slave, while the single-ended protocol only uses multiple addresses at the slave. On the other hand, the dual-ended protocol can only be deployed when both the master and the slave support this protocol. Dual-ended and single-ended protocols can co-exist in the same network. Each slave selects the connection(s) it wants to make with the available masters. A dual-ended slave could switch to single-ended mode if it does not see any dual-ended masters available. A single-ended slave could connect to a single IP address of a dual-ended master.

Multi-path time synchronization, in both variants, requires clocks to use multiple IP addresses. If possible, the set of IP addresses for each clock should be chosen in a way that enables the establishment

of paths that are the most different. It is applicable if the load balancing rules in the network are known. Using multiple IP addresses introduces a tradeoff. A large number of IP addresses allows a large number of diverse paths, providing the advantages of slave diversity discussed in Section 1. On the other hand, a large number of IP addresses is more costly, requires the network topology to be more redundant, and exacts extra management overhead.

The descriptions in this section refer to the end-to-end scheme of PTP, but are similarly applicable to the peer-to-peer scheme. The MPNTP protocol described in this document refers to the NTP client-server mode, although the concepts described here can be extended to include the symmetric variant as well.

Multi-path synchronization protocols by nature require protocol messages to be sent as unicast. Specifically in PTP, the following messages must be sent as unicast in MPPTP: Sync, Delay_Req, Delay_Resp, PDelay_Req, PDelay_Resp, Follow_Up, and PDelay_Resp_Follow_Up. Note that [IEEE1588] allows these messages to be sent either as multicast or as unicast.

5.1. Single-Ended Multi-Path Synchronization

In the single-ended approach, only the slave is aware of the fact that multiple paths are used, while the master is agnostic to the usage of multiple paths. This approach allows a hybrid network, where some of the clocks are multi-path clocks, and others are conventional one-path clocks. A single-ended multi-path clock presents itself to the network as N independent clocks, using N IP addresses, as well as N clock identity values (in PTP). Thus, the usage of multiple slave identities by a slave clock is transparent from the master's point of view, such that it treats each of the identities as a separate slave clock.

5.1.1. Single-Ended MPPTP Synchronization Message Exchange

The single-ended MPPTP message exchange procedure is as follows.

- o Each single-ended MPPTP clock has a fixed set of N IP addresses and N corresponding clockIdentities. Each clock arbitrarily defines one of its IP addresses and clockIdentity values as the clock primary identity.
- o A single-ended MPPTP port sends Announce messages only from its primary identity, according to the BMC algorithm.

- o The BMC algorithm at each clock determines the master, based on the received Announce messages.
- o A single-ended MPPTP port that is in the 'slave' state uses unicast negotiation to request the master to transmit unicast messages to each of the N slave clock identities. The slave port periodically sends N Signaling messages to the master, using each of its N identities. The Signaling message includes the REQUEST_UNICAST_TRANSMISSION_TLV.
- o The master periodically sends unicast Sync messages from its primary identity, identified by the sourcePortIdentity and IP address, to each of the slave identities.
- o The slave, upon receiving a Sync message, identifies its path according to the destination IP address. The slave sends a Delay_Req unicast message to the primary identity of the master. The Delay_Req is sent using the slave identity corresponding to the path the Sync was received through. Note that the rate of Delay_Req messages may be lower than the Sync message rate, and thus a Sync message is not necessarily followed by a Delay_Req.
- o The master, in response to a Delay_Req message from the slave, responds with a Delay_Resp message using the IP address and sourcePortIdentity from the Delay_Req message.
- o Upon receiving the Delay_Resp message, the slave identifies the path using the destination IP address and the requestingPortIdentity. The slave can then compute the corresponding path delay and the offset from the master.
- o The slave combines the information from all negotiated paths.

5.1.2. Single-Ended MPNTP Synchronization Message Exchange

The single-ended MPNTP message exchange procedure is as follows.

- o A single-ended MPNTP client has N separate identities, i.e., N IP addresses. The assumption is that the server information, including its IP address is known to the NTP clients.
- o A single-ended MPNTP client initiates the NTP protocol with an NTP server N times, using each of its N identities.
- o The NTP protocol is maintained between the server and each of the N client identities.

- o The client sends NTP messages to the master using each of its N identities.
- o The server responds to the client's NTP messages using the IP address from the received NTP packet.
- o The client, upon receiving an NTP packet, uses the IP destination address to identify the path it came through, and uses the time information accordingly.
- o The client combines the information from all paths.

5.2. Dual-Ended Multi-Path Synchronization

In dual-ended multi-path synchronization each clock has N IP addresses. Time synchronization messages are exchanged between some of the combinations of {master IP, slave IP} addresses, allowing multiple paths between the master and slave. Note that the actual number of paths between the master and slave may be less than the number of chosen {master, slave} IP address pairs.

Once the multiple two-way connections are established, a separate synchronization protocol exchange instance is run through each of them.

5.2.1. Dual-Ended MPPTP Synchronization Message Exchange

The dual-ended MPPTP message exchange procedure is as follows.

- o Every clock has N IP addresses, but uses a single clockIdentity.
- o The BMC algorithm at each clock determines the master. The master is identified by its clockIdentity, allowing other clocks to know the multiple IP addresses it uses.
- o When a clock sends an Announce message, it sends it from each of its IP addresses with its clockIdentity.
- o A dual-ended MPPTP port that is in the 'slave' state uses unicast negotiation to request the master to transmit unicast messages to some or all of its N_s IP addresses. This negotiation is done individually between a slave IP address and the corresponding master IP address that the slave desires a connection with. The slave port periodically sends Signaling messages to the master, using some or all of its N_s IP addresses as source, to the corresponding master's N_m IP addresses. The Signaling message includes the REQUEST_UNICAST_TRANSMISSION_TLV.

- o The master periodically sends unicast Sync messages from each of its IP addresses to the corresponding slave IP addresses for which a unicast connection was negotiated.
- o The slave, upon receiving a Sync message, identifies its path according to the {source, destination} IP addresses. The slave sends a Delay_Req unicast message, swapping the source and destination IP addresses from the Sync message. Note that the rate of Delay_Req messages may be lower than the Sync message rate, and thus a Sync message is not necessarily followed by a Delay_Req.
- o The master, in response to a Delay_Req message from the slave, responds with a Delay_Resp message using the sourcePortIdentity from the Delay_Req message, and swapping the IP addresses from the Delay_Req.
- o Upon receiving the Delay_Resp message, the slave identifies the path using the {source, destination} IP address pair. The slave can then compute the corresponding path delay and the offset from the master.
- o The slave combines the information from all negotiated paths.

5.2.2. Dual-Ended MPNTP Synchronization Message Exchange

The MPNTP message exchange procedure is as follows.

- o Each NTP clock has a set of N IP addresses. The assumption is that the server information, including its multiple IP addresses is known to the NTP clients.
- o The MPNTP client chooses N_{svr} of the N server IP addresses and N_c of the N client IP addresses and initiates the $N_{svr} * N_c$ instances of the protocol, one for each {server IP, client IP} pair, allowing the client to combine the information from the $N_s * N_c$ paths.
(N_{svr} and N_c indicate the number of IP addresses of the server and client, respectively, which a client chooses to connect with)
- o The client sends NTP messages to the master using each of the source-destination address combinations.
- o The server responds to the client's NTP messages using the IP address combination from the received NTP packet.

- o Using the {source, destination} IP address pair in the received packets, the client identifies the path, and performs its computations for each of the paths accordingly.
- o The client combines the information from all paths.

5.3. Using Traceroute for Path Discovery

The protocols presented above use multiple IP addresses in a single clock to create multiple paths. However, although each two-way path is defined by a different {master, slave} address pair, some of the IP address pairs may traverse exactly the same network path, making them redundant. Traceroute-based path discovery can be used for filtering only the IP addresses that obtain diverse paths. 'Paris Traceroute' [PARIS] and 'TraceFlow' [TRACEFLOW] are examples of tools that discover the paths between two points in the network.

The Traceroute-based filtering can be implemented by both master and slave nodes, or it can be restricted to run only on slave nodes to reduce the overhead on the master. For networks that guarantee the path of the timing packets in the forward and reverse direction are the same, path discovery should only be performed at the slave.

5.4. Using Unicast Discovery for MPPTP

As presented above, MPPTP uses Announce messages and the BMC algorithm to discover the master. The unicast discovery option of PTP can be used as an alternative.

When using unicast discovery the MPPTP slave ports maintain a list of the IP addresses of the master. The slave port uses unicast negotiation to request unicast service from the master, as follows:

- o In single-ended MPPTP, the slave uses unicast negotiation from each of its identities to the master's (only) identity.
- o In dual-ended MPPTP, the slave uses unicast negotiation from its IP addresses, each to a corresponding master IP address to request unicast synchronization messages.

Afterwards, the message exchange continues as described in sections 5.1.1. and 5.2.1.

The unicast discovery option can be used in networks that do not support multicast or in networks in which the master clocks are known in advance. In particular, unicast discovery avoids multicasting Announce messages.

6. Combining Algorithm

Previous sections discussed the methods of creating the multiple paths and obtaining the time information required by the slave algorithm. Once the time information is received through each of the paths, the slave should use a combining algorithm, which consolidates the information from the different paths into a single clock. Various methods have been suggested for combining information from different paths or from different clocks, e.g., [NTP], [SLAVEDIV], [HIGH-AVAI], [KALMAN]. The choice of the combining algorithm is local to the slave, and does not affect the interoperability of the protocol. Hence, this document does not define a specific method to be used by the slave.

7. Security Considerations

The security aspects of time synchronization protocols are discussed in detail in [TICTOCSEC]. The methods describe in this document propose to run a time synchronization protocol through redundant paths, and thus allow to detect and mitigate man-in-the-middle attacks, as described in [DELAY-ATT].

8. IANA Considerations

There are no IANA actions required by this document.

RFC Editor: please delete this section before publication.

9. Acknowledgments

The authors gratefully acknowledge the useful comments provided by Peter Meyer and Doug Arnold, as well as other comments received from the TICTOC working group participants.

This document was prepared using 2-Word-v2.0.template.dot.

10. References

10.1. Normative References

- [IEEE1588] IEEE Instrumentation and Measurement Society, "IEEE Standard for a Precision Clock Synchronization Protocol for Networked Measurement and Control Systems", IEEE Std 1588, 2008.

- [NTP] D. Mills, J. Martin, J. Burbank, W. Kasch, "Network Time Protocol Version 4: Protocol and Algorithms Specification", IETF, RFC 5905, 2010.

10.2. Informative References

- [ASSYMETRY] Yihua He and Michalis Faloutsos and Srikanth Krishnamurthy and Bradley Huffaker, "On routing asymmetry in the internet", IEEE Globecom, 2005.
- [ASSYMETRY2] Abhinav Pathak, Himabindu Pucha, Ying Zhang, Y. Charlie Hu, and Z. Morley Mao, "A measurement study of internet delay asymmetry", PAM'08, 2008.
- [DELAY-ATT] T. Mizrahi, "A Game Theoretic Analysis of Delay Attacks against Time Synchronization Protocols", ISPCS, 2012.
- [HIGH-AVAI] P. Ferrari, A. Flammini, S. Rinaldi, G. Prytz, "High availability IEEE 1588 nodes over IEEE 802.1 aq Shortest Path Bridging networks" ISPCS, 2013.
- [KALMAN] G. Giorgi, C. Narduzzi, "Kalman filtering for multi-path network synchronization" ISPCS, 2014.
- [MULTI] A. Shpiner, Y. Revah, T. Mizrahi, "Multi-path Time Protocols" ISPCS, 2013.
- [PARIS] Brice Augustin, Timur Friedman and Renata Teixeira, "Measuring Load-balanced Paths in the Internet", IMC, 2007.
- [PARIS2] B. Augustin, T. Friedman, and R. Teixeira, "Measuring Multipath Routing in the Internet", IEEE/ACM Transactions on Networking, 19(3), p. 830 - 840, June 2011.
- [SLAVEDIV] T. Mizrahi, "Slave Diversity: Using Multiple Paths to Improve the Accuracy of Clock Synchronization Protocols", ISPCS, 2012.
- [TICTOCSEC] T. Mizrahi, "Security Requirements of Time Protocols in Packet Switched Networks", IETF, RFC 7384, 2014.
- [TRACEFLOW] J. Narasimhan, B. V. Venkataswami, R. Groves and P. Hoose, "Traceflow", IETF, draft-janapath-intarea-traceflow, work in progress, 2012.

Authors' Addresses

Alex Shpiner
Department of Electrical Engineering
Technion - Israel Institute of Technology
Haifa, 32000 Israel

Email: shalex@tx.technion.ac.il

Richard Tse
PMC-Sierra
8555 Baxter Place
Burnaby, BC
Canada
V5A 4V7

Email: Richard.Tse@pmcs.com

Craig Schelp
PMC-Sierra
8555 Baxter Place
Burnaby, BC
Canada
V5A 4V7

Email: craig.schelp@pmcs.com

Tal Mizrahi
Marvell
6 Hamada St.
Yokneam, 20692 Israel

Email: talmi@marvell.com

Internet-Draft

Enterprise Profile for PTP

Feb 2015

TICTOC Working Group

Internet Draft

Intended status: Standards Track

Expires: August 4, 2015

Doug Arnold

Meinberg-USA

Heiko Gerstung

Meinberg

Feb 4, 2015

Enterprise Profile for the Precision Time Protocol
With Mixed Multicast and Unicast Messages

draft-ietf-tictoc-ntp-enterprise-profile-05.txt

Status of this Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79. This document may not be modified, and derivative works of it may not be created, except to publish it as an RFC and to translate it into languages other than English.

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF), its areas, and its working groups. Note that other groups may also distribute working documents as Internet-Drafts.

Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as "work in progress."

The list of current Internet-Drafts can be accessed at <http://www.ietf.org/ietf/lid-abstracts.txt>

The list of Internet-Draft Shadow Directories can be accessed at <http://www.ietf.org/shadow.html>

This Internet-Draft will expire on August 4, 2015.

Copyright Notice

Copyright (c) 2015 IETF Trust and the persons identified as the document authors. All rights reserved.

This document is subject to BCP 78 and the IETF Trust's Legal Provisions Relating to IETF Documents (<http://trustee.ietf.org/license-info>) in effect on the date of publication of this document. Please review these documents carefully, as they describe your rights and restrictions with respect to this document. Code Components extracted from this document must include Simplified BSD License text as described in Section 4.e of the Trust Legal Provisions and are provided without warranty as described in the Simplified BSD License.

Abstract

This document describes a profile for the use of the Precision Time Protocol in an IPV4 or IPV6 Enterprise information system environment. The profile uses the End to End Delay Measurement Mechanism, allows both multicast and unicast Delay Request and Delay Response Messages.

Table of Contents

1.	Introduction	2
2.	Conventions used in this document	3
3.	Technical Terms	3
4.	Problem Statement	5
5.	Network Technology	6
6.	Time Transfer and Delay Measurement	7
7.	Default Message Rates	8
8.	Requirements for Master Clocks	8
9.	Requirements for Slave Clocks	9
10.	Requirements for Transparent Clocks	9
11.	Requirements for Boundary Clocks	10
12.	Management and Signaling Messages	10
13.	Forbidden PTP Options	10
14.	Interoperation with Other PTP Profiles	10
15.	Security Considerations	11
16.	IANA Considerations	11
17.	References	11
	17.1. Normative References	11
	17.2. Informative References	12
18.	Acknowledgments	12
19.	Authors addresses	12

1. Introduction

The Precision Time Protocol ("PTP"), standardized in IEEE 1588, has been designed in its first version (IEEE 1588-2002) with the goal to minimize configuration on the participating nodes. Network communication was based solely on multicast messages, which unlike NTP did not require that a receiving node ("slave clock") in [IEEE1588] needs to know the identity of the time sources in the network (the Master Clocks).

The so-called "Best Master Clock Algorithm" ([IEEE1588] Clause 9.3), a mechanism that all participating PTP nodes must follow, set up strict rules for all members of a PTP domain to determine which node shall be the active sending time source (Master Clock). Although the multicast communication model has advantages in smaller networks, it complicated the application of PTP in larger networks, for example in environments like IP based telecommunication networks or financial data centers. It is considered inefficient that, even if the content of a message applies only to one receiver, it is forwarded by the underlying

network (IP) to all nodes, requiring them to spend network bandwidth and other resources like CPU cycles to drop the message.

The second revision of the standard (IEEE 1588-2008) is the current version (also known as PTPv2) and introduced the possibility to use unicast communication between the PTP nodes in order to overcome the limitation of using multicast messages for the bi-directional information exchange between PTP nodes. The unicast approach avoided that, in PTP domains with a lot of nodes, devices had to throw away up to 99% of the received multicast messages because they carried information for some other node. PTPv2 also introduced so-called "PTP profiles" ([IEEE1588] Clause 19.3). This construct allows organizations to specify selections of attribute values and optional features, simplifying the configuration of PTP nodes for a specific application. Instead of having to go through all possible parameters and configuration options and individually set them up, selecting a profile on a PTP node will set all the parameters that are specified in the profile to a defined value. If a PTP profile definition allows multiple values for a parameter, selection of the profile will set the profile-specific default value for this parameter. Parameters not allowing multiple values are set to the value defined in the PTP profile. A number of PTP features and functions are optional and a profile should also define which optional features of PTP are required, permitted or prohibited. It is possible to extend the PTP standard with a PTP profile by using the TLV mechanism of PTP (see [IEEE1588] Clause 13.4), defining an optional Best Master Clock Algorithm and a few other ways. PTP has its own management protocol (defined in [IEEE1588] Clause 15.2) but allows a PTP profile specify an alternative management mechanism, for example SNMP.

2. Conventions used in this document

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC-2119 [RFC2119].

In this document, these words will appear with that interpretation only when in ALL CAPS. Lower case uses of these words are not to be interpreted as carrying RFC-2119 significance.

3. Technical Terms

Acceptable Master Table: A PTP Slave Clock may maintain a list of masters which it is willing to synchronize to.

Alternate Master: A PTP Master Clock, which is not the Best Master, may act as a master with the Alternate Master flag set on the messages it sends.

Announce message: Contains the master clock properties of a Master clock. Used to determine the Best Master.

Best Master: A clock with a port in the master state, operating consistently with the Best Master Clock Algorithm.

Best Master Clock Algorithm: A method for determining which state a port of a PTP clock should be in. The algorithm works by identifying which of several PTP Master capable clocks is the best master. Clocks have priority to become the acting Grandmaster, based on the properties each Master Clock sends in its Announce Message.

Boundary Clock: A device with more than one PTP port. Generally boundary clocks will have one port in slave state to receive timing and then other ports in master state to re-distribute the timing.

Clock Identity: In IEEE 1588-2008 this is a 64-bit number assigned to each PTP clock which must be unique. Often the Ethernet MAC address is used since there is already an international infrastructure for assigning unique numbers to each device manufactured.

Domain: Every PTP message contains a domain number. Domains are treated as separate PTP systems in the network. Slaves, however, can combine the timing information derived from multiple domains.

End to End Delay Measurement Mechanism: A network delay measurement mechanism in PTP facilitated by an exchange of messages between a Master Clock and Slave Clock.

Grandmaster: the primary master clock within a domain of a PTP system

IEEE 1588: The timing and synchronization standard which defines PTP, and describes The node, system, and communication properties necessary to support PTP.

Master clock: a clock with at least one port in the master state.

NTP: Network Time Protocol, defined by RFC 5905, see [NTP].

Ordinary Clock: A clock that has a single Precision Time Protocol (PTP) port in a domain and maintains the timescale used in the domain. It may serve as a master clock, or be a slave clock.

Peer to Peer Delay Measurement Mechanism: A network delay measurement mechanism in PTP facilitated by an exchange of messages between adjacent devices in a network.

Preferred Master: A device intended to act primarily as the Grandmaster of a PTP system, or as a back up to a Grandmaster.

PTP: The Precision Time Protocol, the timing and synchronization protocol define by IEEE 1588.

PTP port: An interface of a PTP clock with the network. Note that there may be multiple PTP ports running on one physical interface, for example a unicast slave which talks to several Grandmaster clocks in parallel.

PTPv2: Refers specifically to the second version of PTP defined by IEEE 1588-2008.

Rogue Master: A clock with a port in the master state, even though it should not be in the master state according to the Best Master Clock Algorithm, and does not set the alternate master flag.

Slave clock: a clock with at least one port in the slave state, and no ports in the master state.

Slave Only Clock: An Ordinary clock which cannot become a Master clock.

TLV: Type Length Value, a mechanism for extending messages in networked communications.

Transparent Clock. A device that measures the time taken for a PTP event message to transit the device and then updates the message with a correction for this transit time.

Unicast Discovery: A mechanism for PTP slaves to establish a unicast communication with PTP masters using a configures table of master IP addresses and Unicast Message Negotiation.

Unicast Negotiation: A mechanism in PTP for Slave Clocks to negotiate unicast Sync, announce and Delay Request Message Rates from a Master Clock.

4. Problem Statement

This document describes a version of PTP intended to work in large enterprise networks. Such networks are deployed, for example, in financial corporations. It is becoming increasingly common in such networks to perform distributed time tagged measurements, such as one-way packet latencies and cumulative delays on software systems spread across multiple computers. Furthermore there is often a desire to check the age of information time tagged by a different machine. To perform these measurements it is necessary to deliver a common precise time to multiple devices on a network. Accuracy currently required in the Financial Industry range from 100 microseconds to 500 nanoseconds to the Grandmaster. This profile does not specify timing performance requirements, but such requirements explain why the needs cannot always be met by NTP, as commonly implemented. Such accuracy cannot usually be achieved with a traditional time transfer such as NTP, without adding

non-standard customizations such as hardware time stamping, and on path support. These features are currently part of PTP, or are allowed by it. Because PTP has a complex range of features and options it is necessary to create a profile for enterprise networks to achieve interoperability between equipment manufactured by different vendors.

Although enterprise networks can be large, it is becoming increasingly common to deploy multicast protocols, even across multiple subnets. For this reason it is desired to make use of multicast whenever the information going to many destinations is the same. It is also advantageous to send information which is unique to one device as a unicast message. The latter can be essential as the number of PTP slaves becomes hundreds or thousands.

PTP devices operating in these networks need to be robust. This includes the ability to ignore PTP messages which can be identified as improper, and to have redundant sources of time.

5. Network Technology

This PTP profile SHALL operate only in networks characterized by UDP [RFC768] over either IPv4 [RFC791] or IPv6 [RFC2460], as described by Annexes D and E in [IEEE1588] respectively. If a network contains both IPv4 and IPv6, then they SHALL be treated as separate communication paths. Clocks which communicate using IPv4 can interact with clocks using IPv6 if there is an intermediary device which simultaneously communicates with both IP versions. A boundary clock might perform this function, for example. A PTP domain SHALL use either IPv4 or IPv6 over a communication path, but not both. The PTP system MAY include switches and routers. These devices MAY be transparent clocks, boundary clocks, or neither, in any combination. PTP Clocks MAY be Preferred Masters, Ordinary Clocks, or Boundary Clocks. The ordinary clocks may be Slave Only Clocks, or be master capable.

Note that clocks SHOULD always be identified by their clock ID and not the IP or Layer 2 address. This is important in IPv6 networks since Transparent clocks are required to change the source address of any packet which they alter. In IPv4 networks some clocks might be hidden behind a NAT, which hides their IP addresses from the rest of the network. Note also that the use of NATs may place limitations on the topology of PTP networks, depending on the port forwarding scheme employed. Details of implementing PTP with NATs are out of scope of this document.

Similar to NTP, PTP makes the assumption that the one way network delay for Sync Messages and Delay Response Messages are the same. When this is not true it can cause errors in the transfer of time from the Master to the Slave. It is up to the system integrator to design the network so that such effects do not prevent the PTP system from meeting the timing requirements. The details of

network asymmetry are outside the scope of this document. See for example, [G8271].

6. Time Transfer and Delay Measurement

Master clocks, Transparent clocks and Boundary clocks MAY be either one-step clocks or two-step clocks. Slave clocks MUST support both behaviors. The End to End Delay Measurement Method MUST be used.

Note that, in IP networks, Sync messages and Delay Request messages exchanged between a master and slave do not necessarily traverse the same physical path. Thus, wherever possible, the network SHOULD be traffic engineered so that the forward and reverse routes traverse the same physical path. Traffic engineering techniques for path consistency are out of scope of this document.

Sync messages MUST be sent as PTP event multicast messages (UDP port 319) to the PTP primary IP address. Two step clocks SHALL send Follow-up messages as PTP general messages (UDP port 320). Announce messages MUST be sent as multicast messages (UDP port 320) to the PTP primary address. The PTP primary IP address is 224.0.1.129 for IPv4 and FF0X:0:0:0:0:0:0:181 for Ipv6, where X can be a value between 0x0 and 0xF, see [IEEE1588] Annex E, Section E.3.

Delay Request Messages MAY be sent as either multicast or unicast PTP event messages. Master clocks SHALL respond to multicast Delay Request messages with multicast Delay Response PTP general messages. Master clocks SHALL respond to unicast Delay Request PTP event messages with unicast Delay Response PTP general messages. This allow for the use of Ordinary clocks which do not support the Enterprise Profile, as long as they are slave Only Clocks.

Clocks SHOULD include support for multiple domains. The purpose is to support multiple simultaneous masters for redundancy. Leaf devices (non-forwarding devices) can use timing information from multiple masters by combining information from multiple instantiations of a PTP stack, each operating in a different domain. Redundant sources of timing can be ensembled, and/or compared to check for faulty master clocks. The use of multiple simultaneous masters will help mitigate faulty masters reporting as healthy, network delay asymmetry, and security problems. Security problems include man-in-the-middle attacks such as delay attacks, packet interception / manipulation attacks. Assuming the path to each master is different, failures malicious or otherwise would have to happen at more than one path simultaneously. Whenever feasible, the underlying network transport technology SHOULD be configured so that timing messages in different domains traverse different network paths.

7. Default Message Rates

The Sync, Announce and Delay Request default message rates SHALL each be once per second. The Sync and Delay Request message rates MAY be set to other values, but not less than once every 128 seconds, and not more than 128 messages per second. The Announce message rate SHALL NOT be changed from the default value. The Announce Receipt Timeout Interval SHALL be three Announce Intervals for Preferred Masters, and four Announce Intervals for all other masters. Unicast Discovery and Unicast Message Negotiation options MUST NOT be utilized.

8. Requirements for Master Clocks

Master clocks SHALL obey the standard Best Master Clock Algorithm from [IEEE1588]. PTP systems using this profile MAY support multiple simultaneous Grandmasters as long as each active Grandmaster is operating in a different PTP domain.

A port of a clock SHALL NOT be in the master state unless the clock has a current value for the number of UTC leap seconds. A clock with a port in the master state SHOULD indicate the maximum adjustment to its internal clock within one sync interval. The maximum phase adjustment is indicated in the Enterprise Profile announce TLV field for Maximum Phase Adjustment.

The Announce Messages SHALL include a TLV which indicates that the clock is operating in the Enterprise Profile. The TLV shall have the following structure:

TLV Type (Enumeration16): ORGANIZATION_EXTENSION value = 0003 hex

Length Field (UInteger16): value = 10. The number of TLV octets

Organization Unique Identifier (3 Octets): The Organization ID value for IETF assigned by IEEE = 00005Ehex

IETF Profile number (UInteger8): value = 1

Revision number (UInteger8): value = 1

Port Number (UInteger16): The Port Number of the port transmitting the TLV. The all-ones Port Number, with value FFFFhex, is used to indicate that the identified profile is applicable to all ports on the clock.

Maximum Absolute Phase Adjustment Value within one sync interval (UInteger16): value

Maximum Phase Adjustment Units (Enumeration8):

- Value 0 = unknown
- Value 1 = seconds
- Value 3 = milliseconds
- Value 6 = microseconds
- Value 9 = nanoseconds
- Value 12 = picoseconds
- Value 15 = femtoseconds
- All other values reserved for future use

Slaves can use the Maximum Phase Adjustment to determine if the clock is slewing to rapidly for the applications which are of interest. For example if the clock set by slave is used to measure time intervals then it might be desired that that the amount which the time changes during the intervals is limited.

9. Requirements for Slave Clocks

Slave clocks MUST be able to operate properly in a network which contains multiple Masters in multiple domains. Slaves SHOULD make use of information from the all Masters in their clock control subsystems. Slave Clocks MUST be able to operate properly in the presence of a Rogue Master. Slaves SHOULD NOT Synchronize to a Master which is not the Best Master in its domain. Slaves will continue to recognize a Best Master for the duration of the Announce Time Out Interval. Slaves MAY use an Acceptable Master Table. If a Master is not an Acceptable Master, then the Slave MUST NOT synchronize to it. Note that IEEE 1588-2008 requires slave clocks to support both two-step or one-step Master clocks. See [IEEE1588], section 11.2.

Since Announce messages are sent as multicast messages slaves can obtain the IP addresses of master from the Announce messages. Note that the IP source addresses of Sync and Follow-up messages may have been replaced by the source addresses of a transparent clock, so slaves MUST send Delay Request messages to the IP address in the Announce message. Sync and Follow-up messages can be correlated with the Announce message using the clock ID, which is never altered by Transparent clocks in this profile.

10. Requirements for Transparent Clocks

Transparent clocks SHALL NOT change the transmission mode of an Enterprise Profile PTP message. For example a Transparent clock SHALL NOT change a unicast message to a multicast message. Transparent clocks SHALL NOT alter the Enterprise Profile TLV of an Announce message, or any other part of an Announce message. Transparent Clocks SHOULD support multiple domains. Transparent Clocks which syntonize to the master clock will need to maintain separate clock rate offsets for each of the supported domains.

11. Requirements for Boundary Clocks

Boundary Clocks SHOULD support multiple simultaneous PTP domains. This will require them to maintain servo loops for each of the domains supported, at least in software. Boundary clocks MUST NOT combine timing information from different domains.

12. Management and Signaling Messages

PTP Management messages MAY be used. Any PTP management message which is sent with the `targetPortIdentity.clockIdentity` set to all 1s (all clocks) MUST be sent as a multicast message. Management messages with any other value of for the Clock Identity is intended for a specific clock and MUST be sent as a unicast message. Similarly, if any signaling messages are used they MUST also be sent as unicast messages whenever the message is intended for a specific clock.

13. Forbidden PTP Options

Clocks operating in the Enterprise Profile SHALL NOT use peer to peer timing for delay measurement. Clocks operating in the Enterprise Profile SHALL NOT use Unicast Message Negotiation to determine message rates. Slave clocks operating in the Enterprise Profile SHALL NOT use Unicast Discovery to establish connection to Master clocks. Grandmaster Clusters are NOT ALLOWED. The Alternate Master option is also forbidden. Clocks operating in the Enterprise Profile SHALL NOT use Alternate Timescales.

14. Interoperation with Other PTP Profiles

Clocks operating in the Enterprise Profile will not interoperate with clocks operating in the Power Profile [C37.238], because the Enterprise Profile requires the End to End Delay Measurement Mechanism and the Power Profile requires the Peer to Peer Delay Measurement Mechanism.

Clocks operating in the Enterprise Profile will not interoperate with clocks operating in the Telecom Profile for Frequency Synchronization[G8265.1], because the Enterprise Profile forbids Unicast Message Negotiation. This feature is part of the default manner of operation for the Telecom Profile for Frequency Synchronization.

Clocks operating in the Enterprise Profile will interoperate with clocks operating in the Default Profile described in [IEEE1588] Annex J.3. This variant of the Default Profile uses the End to End Delay Measurement Mechanism. In addition the Default Profile would have to operate over IPv4 or IPv6 networks, and use management messages in unicast when those messages are directed at a specific clock. If either of these requirements are not met than Enterprise

Profile clocks will not interoperate with Annex J.3 Default Profile Clocks. The Enterprise Profile Profile will will not interoperate with the Annex J.4 variant of the Default Profile which requires use of the Peer to Peer Delay Measurement Mechanism.

Enterprise Profile Clocks will interoperate with clocks operating in other profiles if the clocks in the other profiles obey the rules of the Enterprise Profile. These rules MUST NOT be changed to achieve interoperability with other profiles.

15. Security Considerations

Protocols used to transfer time, such as PTP and NTP can be important to security mechanisms which use time windows for keys and authorization. Passing time through the networks poses a security risk since time can potentially be manipulated. The use of multiple simultaneous masters, using multiple PTP domains can mitigate problems from rogue masters and man-in-the-middle attacks. See sections 9 and 10. Additional security mechanisms are outside the scope of this document.

16. IANA Considerations

There are no IANA requirements in this specification.

17. References

17.1. Normative References

- [IEEE1588] IEEE std. 1588-2008, "IEEE Standard for a Precision Clock Synchronization for Networked Measurement and Control Systems." July, 2008.
- [RFC768] Postel, J., "User Datagram Protocol," RFC 768, August, 1980.
- [RFC791] "Internet Protocol DARPA Internet Program Protocol Specification," RFC 791, September, 1981.
- [RFC2119] Bradner, S., "Key words for use in RFCs to Indicate Requirement Levels", BCP 14, RFC 2119, March 1997.
- [RFC2460] Deering, S., Hinden, R., "Internet Protocol, Version 6 (IPv6) Specification," RFC 2460, December, 1998.

17.2. Informative References

- [C37.238] IEEE std. C37.238-2011, "IEEE Standard Profile for Use of IEEE 1588 Precision Time Protocol in Power System Applications," July 2011.
- [G8265.1] ITU-T G.8265.1/Y.1365.1, "Precision time protocol telecom profile for frequency synchronization," October 2011.
- [G8271] ITU-T G.8271/Y.1366, "Time and Phase Synchronization Aspects of Packet Networks" February, 2012.
- [NTP] Mills, D., Martin, J., Burbank, J., Kasch, W., "Network Time Protocol Version 4: Protocol and Algorithms Specification," RFC 5905, June 2010.

18. Acknowledgments

The authors would like to thank members of IETF for reviewing and providing feedback on this draft.

This document was initially prepared using
2-Word-v2.0.template.dot.

19. Authors' Addresses

Doug Arnold
Meinberg USA
228 Windsor River Rd
Windsor, CA 95492
USA

Email: doug.arnold@meinberg-usa.com

Heiko Gerstung
Meinberg Funkuhren GmbH & Co. KG
Lange Wand 9
D-31812 Bad Pyrmont
Germany

Email: Heiko.gerstung@meinberg.de

TICTOC Working Group
INTERNET DRAFT
Intended status: Standards Track

Vinay Shankarkumar
Laurent Montini
Cisco Systems

Tim Frost
Calnex Solutions Ltd.

Greg Dowd
Microsemi

Expires: September 24, 2015

March 25, 2015

Precision Time Protocol Version 2 (PTPv2)
Management Information Base
draft-ietf-tictoc-ntp-mib-07.txt

Status of this Memo

This Internet-Draft is submitted to IETF in full conformance with the provisions of BCP 78 and BCP 79.

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF), its areas, and its working groups. Note that other groups may also distribute working documents as Internet-Drafts.

Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as "work in progress."

The list of current Internet-Drafts can be accessed at
<http://www.ietf.org/ietf/lid-abstracts.txt>

The list of Internet-Draft Shadow Directories can be accessed at
<http://www.ietf.org/shadow.html>

This Internet-Draft will expire on September 24, 2015.

Copyright Notice

Copyright (c) 2015 IETF Trust and the persons identified as the document authors. All rights reserved.

This document is subject to BCP 78 and the IETF Trust's Legal Provisions Relating to IETF Documents

(<http://trustee.ietf.org/license-info>) in effect on the date of publication of this document. Please review these documents carefully, as they describe your rights and restrictions with respect to this document. Code Components extracted from this document must include Simplified BSD License text as described in Section 4.e of the Trust Legal Provisions and are provided without warranty as described in the Simplified BSD License.

Abstract

This memo defines a portion of the Management Information Base (MIB) for use with network management protocols in TCP/IP-based internets. In particular, it defines objects for managing networks using Precision Time Protocol, specified in IEEE Std. 1588(TM)-2008.

This memo specifies a MIB module in a manner that is both compliant to the SNMPv2 SMI, and semantically identical to the peer SNMPv1 definitions.

Table of Contents

1. Introduction.....	2
1.1. Relationship to other Profiles and MIBs.....	3
1.2. Change Log.....	3
2. The SNMP Management Framework.....	4
3. Overview.....	5
4. IETF PTP MIB Definition.....	6
5. Security Considerations.....	75
6. IANA Considerations.....	75
7. References.....	76
7.1. Normative References.....	76
7.2. Informative References.....	76
8. Acknowledgements.....	78
9. Author's Addresses.....	78

1. Introduction

This memo defines a portion of the Management Information Base (MIB) for use with network management protocols in the Internet Community. In particular, it describes managed objects used for managing PTP devices including the ordinary clock, transparent clock, boundary clocks.

This MIB is restricted to reading standard PTP data elements, as described in [IEEE 1588-2008]. This enables it to monitor the operation of PTP clocks within the network. It is envisioned this MIB will complement other managed objects to be defined that will provide more detailed information on the performance of PTP clocks

supporting the Telecom Profile defined in [G.8265.1], and any future profiles that may be defined. Those objects are considered out of scope for the current draft.

Similarly, this MIB is read-only and not intended to provide the ability to configure PTP clocks. Since PTP clocks are often embedded in other network elements such as routers, switches and gateways, this ability is generally provided via the configuration interface for the network element.

1.1. Relationship to other Profiles and MIBs

This MIB is intended to be used with the default PTP profile described in [IEEE 1588-2008] when running over the IP network layer. As stated above, it is envisioned this MIB will complement other managed objects to be defined to monitor and measure the performance of PTP clocks supporting specific PTP profiles, e.g. the Telecom Profile defined in [G.8265.1].

Some other PTP profiles have their own MIBs defined as part of the profile, and this MIB is not intended to replace those MIBs.

1.2. Change Log

This section tracks changes made to the revisions of the Internet Drafts of this document. It will be **deleted** when the document is published as an RFC.

draft-vinay-tictoc-ntp-mib

-00 Mar 11 Initial version; showed structure of MIB

draft-ietf-tictoc-ntp-mib

-00 Jul 11 First full, syntactically correct and compileable MIB

-01 Jan 12 Revised following comments from Bert Wijnen:
- revised introduction to clarify the scope, and the relationship to other MIBs and profiles
- changed name to "ntpbases"
- corrected some data types
- corrected references and typos

-02 Jul 12 Revised following comment at IETF83:
- changed "ntpbasesClockPortRunningIPversion" to the more generic "ntpbasesClockPortRunningTransport", covering all transport types defined in [IEEE 1588-2008] (i.e. IPv4, IPv6, Ethernet, DeviceNet and ControlNet).
- changed addresses associated with transports from

"InetAddress" (for the IP transport) to a string, to allow for the different transport types.

- 03 Jul 12 Minor changes following comments from Andy Bierman:
 - corrected some compilation errors
 - moved OBJECT-GROUP and MODULE-COMPLIANCE macros to the end
- 04 Jan 13 Changes:
 - Use of 'AutonomousType' import
 - Display hint being specified for ClockIdentity, ClockInterval, ClockPortTransportTypeAddress Textual Conventions
 - Removal of the Textual convention ClockPortTransportType, replaced with the wellKnownTransportTypes
 - Modified ptpbaseClockPortCurrentPeerAddressType, ptpbaseClockPortRunningTransport, ptpbaseClockPortAssociateAddressType, to use AutonomousType.
 - various textual changes to descriptive text in response to comments
- 05 Feb 13 Several changes in response to comments from Alun Luchuk and Kevin Gross:
 - Modified the use of wellKnownTransportTypes and wellKnownEncapsulationTypes
 - changed ptpbaseClockPortSyncOneStep to ptpbaseClockPortSyncTwoStep to match IEEE1588 semantics
 - Re-ordered textual conventions to be alphabetic
 - Changed some types from Integer32 to use defined textual conventions
 - various minor descriptive text changes
- 06 Mar 14 Updated author information, and fixed typos
- 07 Mar 15 Updated author information, and fixed typo/enum

2. The SNMP Management Framework

The SNMP Management Framework presently consists of five major components:

- o An overall architecture, described in STD62, [RFC 3411].
- o Mechanisms for describing and naming objects and events for the purpose of management. The first version of this Structure of

Management Information (SMI) is called SMIV1 and described in STD 16: [RFC 1155], [RFC 1212] and [RFC 1215].

The second version, called SMIV2, is described in STD 58: [RFC 2578], [RFC 2579] and [RFC 2580].

- o Message protocols for transferring management information. The first version of the SNMP message protocol is called SNMPv1 and described in STD 15 [RFC 1157]. A second version of the SNMP message protocol, which is not an Internet standards track protocol, is called SNMPv2c and described in [RFC 1901] and [RFC 1906]. The third version of the message protocol is called SNMPv3 and described in STD62: [RFC 3417], [RFC 3412] and [RFC 3414].
- o Protocol operations for accessing management information. The first set of protocol operations and associated PDU formats is described in STD 15 [RFC 1157]. A second set of protocol operations and associated PDU formats is described in STD 62 [RFC 3416].
- o A set of fundamental applications described in STD 62 [RFC 3413] and the view-based access control mechanism described in STD 62 [RFC 3415].

Managed objects are accessed via a virtual information store, termed the Management Information Base or MIB. Objects in the MIB are defined using the mechanisms defined in the SMI.

This memo specifies a MIB module that is compliant to the SMIV2. A MIB conforming to the SMIV1 can be produced through the appropriate translations. The resulting translated MIB must be semantically equivalent, except where objects or events are omitted because no translation is possible (e.g., use of Counter64). Some machine readable information in SMIV2 will be converted into textual descriptions in SMIV1 during the translation process. However, this loss of machine readable information is not considered to change the semantics of the MIB.

3. Overview

The objects defined in this MIB are to be used when describing the Precision Time Protocol (PTPv2).

4. IETF PTP MIB Definition

```
PTPBASE-MIB DEFINITIONS ::= BEGIN
```

```
IMPORTS
```

```
    MODULE-IDENTITY,
    OBJECT-TYPE,
    OBJECT-IDENTITY,
    Gauge32,
    Unsigned32,
    Counter32,
    Counter64,
    mib-2,
    Integer32
    FROM SNMPv2-SMI
    OBJECT-GROUP,
    MODULE-COMPLIANCE
    FROM SNMPv2-CONF
    TEXTUAL-CONVENTION,
    TruthValue,
    DisplayString,
    AutonomousType
    FROM SNMPv2-TC
    InterfaceIndexOrZero
    FROM IF-MIB;
```

```
ptpbasesMIB MODULE-IDENTITY
```

```
    LAST-UPDATED      "201207230000Z"
    ORGANIZATION      "TICTOC Working Group"
    CONTACT-INFO
        "WG Email: tictoc@ietf.org
```

```

        Vinay Shankarkumar
        Cisco Systems,
        Email: vinays@cisco.com
```

```

        Laurent Montini,
        Cisco Systems,
        Email: lmontini@cisco.com
```

```

        Tim Frost,
        Calnex Solutions Ltd.,
```

Email: tim.frost@calnexsol.com

Greg Dowd,
Microsemi Inc.,
Email: greg.dowd@microsemi.com"

DESCRIPTION

"The MIB module for PTP version 2 (IEEE Std. 1588(TM)-2008)

Overview of PTP version 2 (IEEE Std. 1588(TM)-2008)

[IEEE 1588-2008] defines a protocol enabling precise synchronization of clocks in measurement and control systems implemented with packet-based networks, the Precision Time Protocol Version 2 (PTPv2). This MIB does not address the earlier version IEEE Std. 1588(TM)-2002 (PTPv1). The protocol is applicable to network elements communicating using IP. The protocol enables heterogeneous systems that include clocks of various inherent precision, resolution, and stability to synchronize to a grandmaster clock.

The protocol supports system-wide synchronization accuracy in the sub-microsecond range with minimal network and local clock computing resources. [IEEE 1588-2008] uses UDP/IP or Ethernet and can be adapted to other mappings. It includes formal mechanisms for message extensions, higher sampling rates, correction for asymmetry, a clock type to reduce error accumulation in large topologies, and specifications on how to incorporate the resulting additional data into the synchronization protocol. The [IEEE 1588-2008] defines conformance and management capability also.

MIB description

This MIB is to support the Precision Time Protocol version 2 (PTPv2, hereafter designated as PTP) features of network element system devices, when using the default PTP profile described in [IEEE 1588-2008] when running over the IP network layer.

It is envisioned this MIB will complement other managed objects to be defined to monitor and measure the performance of the PTP devices and telecom clocks supporting specific PTP profiles.

Some other PTP profiles have their own MIBs defined as part of the profile, and this MIB is not intended to replace those MIBs.

Acronyms:

ARB	Arbitrary Timescale
E2E	End-to-End
EUI	Extended Unique Identifier.
GPS	Global Positioning System
IANA	Internet Assigned Numbers Authority
IP	Internet Protocol
MAC	Media Access Control according to [IEEE 802.3-2008]
NIST	National Institute of Standards and Technology
NTP	Network Time Protocol (see IETF [RFC 5905])
OUI	Organizational Unique Identifier (allocated by the IEEE)
P2P	Peer-to-Peer
PTP	Precision Time Protocol
TAI	International Atomic Time
TC	Transparent Clock
UDP	User Datagram Protocol
UTC	Coordinated Universal Time

References:

[IEEE 1588-2008] IEEE Standard for A Precision Clock Synchronization Protocol for Networked Measurement and Control Systems, IEEE Std. 1588(TM)-2008, 24 July 2008.

As defined in [IEEE 1588-2008]:

Accuracy:

The mean of the time or frequency error between the clock under test and a perfect reference clock, over an ensemble of measurements. Stability is a measure of how the mean varies with respect to variables such as time, temperature, and so on, while the precision is a measure of the deviation of the error from the mean.

Atomic process:

A process is atomic if the values of all inputs to the process are not permitted to change until all of the results of the

process are instantiated, and the outputs of the process are not visible to other processes until the processing of each output is complete.

Boundary clock:

A clock that has multiple Precision Time Protocol (PTP) ports in a domain and maintains the timescale used in the domain. It may serve as the source of time, i.e., be a master clock, and may synchronize to another clock, i.e., be a slave clock.

Boundary node clock:

A clock that has multiple Precision Time Protocol (PTP) ports in a domain and maintains the timescale used in the domain. It differs from a boundary clock in that the clock roles can change.

Clock:

A node participating in the Precision Time Protocol (PTP) that is capable of providing a measurement of the passage of time since a defined epoch.

Domain:

A logical grouping of clocks that synchronize to each other using the protocol, but that are not necessarily synchronized to clocks in another domain.

End-to-end transparent clock:

A transparent clock that supports the use of the end-to-end delay measurement mechanism between slave clocks and the master clock. Each node must measure the residence time of PTP event messages and accumulate it in Correction Field.

Epoch:

The origin of a timescale.

Event:

An abstraction of the mechanism by which signals or conditions are generated and represented.

Foreign master:

An ordinary or boundary clock sending Announce messages to another clock that is not the current master recognized by the other clock.

Grandmaster clock:

Within a domain, a clock that is the ultimate source of time for clock synchronization using the protocol.

Holdover:

A clock previously synchronized/syntonized to another clock (normally a primary reference or a master clock) but now free-running based on its own internal oscillator, whose frequency is being adjusted using data acquired while it had been synchronized/syntonized to the other clock. It is said to be in holdover or in the holdover mode, as long as it is within its accuracy requirements.

Link:

A network segment between two Precision Time Protocol ports supporting the peer delay mechanism of this standard. The peer delay mechanism is designed to measure the propagation time over such a link.

Management node:

A device that configures and monitors clocks.

Master clock:

In the context of a single Precision Time Protocol communication path, a clock that is the source of time to which all other clocks on that path synchronize.

Message timestamp point:

A point within a Precision Time Protocol event message serving as a reference point in the message. A timestamp is defined by the instant a message timestamp point passes the reference plane of a clock.

Multicast communication:

A communication model in which each Precision Time Protocol message sent from any PTP port is capable of being received and processed by all PTP ports on the same PTP communication path.

Node:

A device that can issue or receive Precision Time Protocol communications on a network.

One-step clock:

A clock that provides time information using a single event message.

On-pass support:

Indicates that each node in the synchronization chain from master to slave can support IEEE-1588.

Ordinary clock:

A clock that has a single Precision Time Protocol port in a domain and maintains the timescale used in the domain. It may serve as a source of time, i.e., be a master clock, or may synchronize to another clock, i.e., be a slave clock.

Parent clock:

The master clock to which a clock is synchronized.

Peer-to-peer transparent clock:

A transparent clock that, in addition to providing Precision Time Protocol event transit time information, also provides corrections for the propagation delay of the link connected to the port receiving the PTP event message. In the presence of peer-to-peer transparent clocks, delay measurements between slave clocks and the master clock are performed using the peer-to-peer delay measurement mechanism.

Phase change rate:

The observed rate of change in the measured time with respect to the reference time. The phase change rate is equal to the fractional frequency offset between the measured frequency and the reference frequency.

PortNumber:

An index identifying a specific Precision Time Protocol port on a PTP node.

Primary reference:

A source of time and or frequency that is traceable to international standards.

Profile:

The set of allowed Precision Time Protocol features applicable to a device.

Precision Time Protocol communication:
Information used in the operation of the protocol, transmitted in a PTP message over a PTP communication path.

Precision Time Protocol communication path:
The signaling path portion of a particular network enabling direct communication among ordinary and boundary clocks.

Precision Time Protocol node:
PTP ordinary, boundary, or transparent clock or a device that generates or parses PTP messages.

Precision Time Protocol port:
A logical access point of a clock for PTP communications to the communications network.

Recognized standard time source:
A recognized standard time source is a source external to Precision Time Protocol that provides time and/or frequency as appropriate that is traceable to the international standards laboratories maintaining clocks that form the basis for the International Atomic Time and Universal Coordinated Time timescales. Examples of these are GPS, NTP, and NIST timeservers.

Requestor:
The port implementing the peer-to-peer delay mechanism that initiates the mechanism by sending a Pdelay_Req message.

Responder:
The port responding to the receipt of a Pdelay_Req message as part of the operation of the peer-to-peer delay mechanism.

Synchronized clocks:
Two clocks are synchronized to a specified uncertainty if they have the same epoch and their measurements of the time of a single event at an arbitrary time differ by no more than that uncertainty.

Syntonized clocks:
Two clocks are syntonized if the duration of the second is the same on both, which means the time as measured by each advances

at the same rate. They may or may not share the same epoch.

Timeout:

A mechanism for terminating requested activity that, at least from the requester's perspective, does not complete within the specified time.

Timescale:

A linear measure of time from an epoch.

Traceability:

A property of the result of a measurement or the value of a standard whereby it can be related to stated references, usually national or international standards, through an unbroken chain of comparisons all having stated uncertainties.

Translation device:

A boundary clock or, in some cases, a transparent clock that translates the protocol messages between regions implementing different transport and messaging protocols, between different versions of [IEEE 1588-2008], or different PTP profiles.

Transparent clock:

A device that measures the time taken for a Precision Time Protocol event message to transit the device and provides this information to clocks receiving this PTP event message.

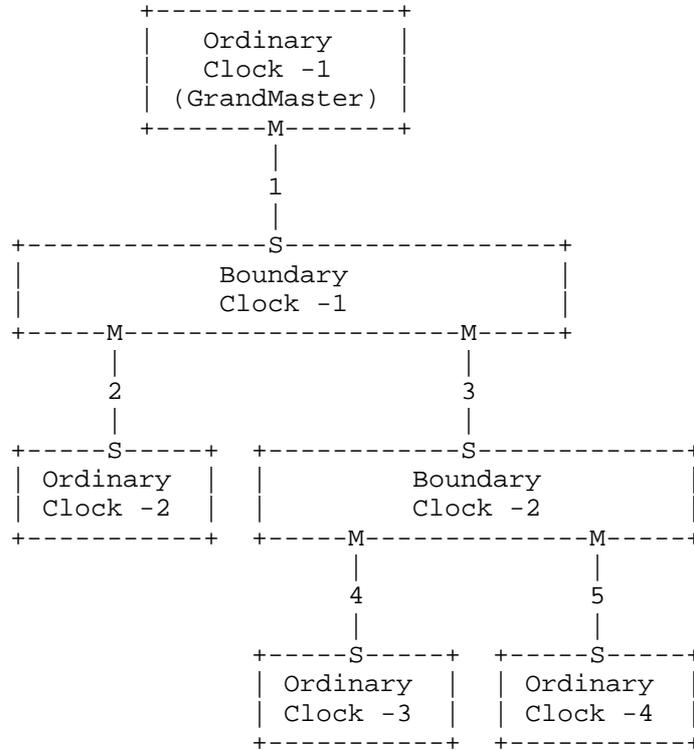
Two-step clock:

A clock that provides time information using the combination of an event message and a subsequent general message.

The below table specifies the object formats of the various textual conventions used.

Data type mapping	Textual Convention	SYNTAX
5.3.2 TimeInterval	ClockTimeInterval	OCTET STRING(SIZE(1..255))
5.3.3 Timestamp	ClockTimestamp	OCTET STRING(SIZE(6))
5.3.4 ClockIdentity	ClockIdentity	OCTET STRING(SIZE(1..255))
5.3.5 PortIdentity	ClockPortNumber	INTEGER(1..65535)
5.3.7 ClockQuality	ClockQualityClassType	

Simple master-slave hierarchy, section 6.6.2.4 [IEEE 1588-2008]:



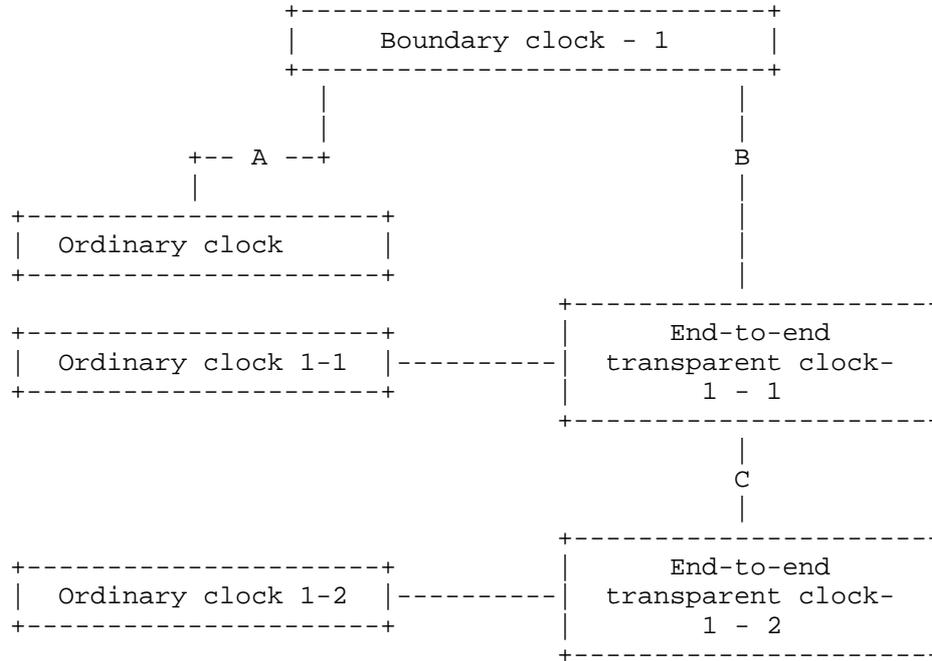
Grandmaster

Boundary Clock(0-N) Ordinary Clocks(0-N)
 Ordinary Clocks(0-N)

Relationship cardinality:

- PTP system 1 : N PTP Clocks
- PTP Clock 1 : 1 Domain
- PTP Clock 1 : N PTP Ports
- PTP Ports N : M Physical Ports (interface in IF-MIB)

Transparent clock diagram, section 6.7.1.3 of [IEEE 1588-2008]:



The MIB refers to the sections of [IEEE 1588-2008]."

-- revision log

::= { mib-2 XXX } -- XXX to be assigned by IANA

-- Textual Conventions

ClockDomainType ::= TEXTUAL-CONVENTION

DISPLAY-HINT "d"

STATUS current

DESCRIPTION

"The Domain is identified by an integer, the domainNumber, in the range of 0 to 255. An integer value that is used to assign each PTP device to a particular domain. The following values

define the valid domains.

Value	Definition
-----	-----
0	Default domain
1	Alternate domain 1
2	Alternate domain 2
3	Alternate domain 3
4 - 127	User-defined domains
128 - 255	Reserved"

REFERENCE "Section 7.1 Domains, Table 2 of [IEEE 1588-2008]"
 SYNTAX Unsigned32 (0..255)

ClockIdentity ::= TEXTUAL-CONVENTION

DISPLAY-HINT "255a"
 STATUS current
 DESCRIPTION

"The clock Identity is an 8-octet array and will be presented in the form of a character array. Network byte order is assumed.

The value of the ClockIdentity should be taken from the IEEE EUI-64 individual assigned numbers as indicated in Section 7.5.2.2.2 of [IEEE 1588-2008].

The EUI-64 address is divided into the following fields:

OUI bytes (0-2)
 Extension identifier bytes (3-7)

The clock identifier can be constructed from existing EUI-48 assignments and here is an abbreviated example extracted from section 7.5.2.2.2 [IEEE 1588-2008].

Company EUI-48 = 0xACDE4823456716
 EUI-64 = ACDE48FFFE23456716

It is important to note the IEEE Registration Authority has deprecated the use of MAC-48 in any new design."

REFERENCE "Section 7.5.2.2.1 of [IEEE 1588-2008]"
 SYNTAX OCTET STRING (SIZE (1..255))

ClockInstanceType ::= TEXTUAL-CONVENTION
DISPLAY-HINT "d"
STATUS current
DESCRIPTION
"The instance of the Clock of a given clock type in a given domain."
SYNTAX Unsigned32 (0..255)

ClockIntervalBase2 ::= TEXTUAL-CONVENTION
DISPLAY-HINT "d"
STATUS current
DESCRIPTION
"The interval included in message types Announce, Sync, Delay_Req, and Pdelay_Req as indicated in section 7.7.2.1 of [IEEE 1588-2008].

The mean time interval between successive messages shall be represented as the logarithm to the base 2 of this time interval measured in seconds on the local clock of the device sending the message. The values of these logarithmic attributes shall be selected from integers in the range -128 to 127 subject to further limits established in an applicable PTP profile."

REFERENCE "Section 7.7.2.1 General interval specification of [IEEE 1588-2008]"
SYNTAX Integer32 (-128..127)

ClockMechanismType ::= TEXTUAL-CONVENTION
STATUS current
DESCRIPTION
"The clock type based on whether End to End or peer to peer mechanisms are used. The mechanism used to calculate the Mean Path Delay as indicated in Table 9 of [IEEE 1588-2008].

Table with 3 columns: Delay mechanism, Value(hex), Specification. Rows include E2E, P2P, and DISABLED with their respective hex values and descriptions.

REFERENCE "Sections 8.2.5.4.4, 6.6.4, 7.4.2 of [IEEE 1588-2008]."
 SYNTAX INTEGER {
 e2e(1),
 p2p(2),
 disabled(254)
 }

ClockPortNumber ::= TEXTUAL-CONVENTION

DISPLAY-HINT "d"
 STATUS current
 DESCRIPTION
 "An index identifying a specific Precision Time Protocol (PTP)
 port on a PTP node."

REFERENCE "Sections 7.5.2.3 and 5.3.5 of [IEEE 1588-2008]"
 SYNTAX Unsigned32 (0..65535)

ClockPortState ::= TEXTUAL-CONVENTION

STATUS current
 DESCRIPTION
 "This is the value of the current state of the protocol engine
 associated with this port."

Port state	Value	Description
initializing	1	In this state a port initializes its data sets, hardware, and communication facilities.
faulty	2	The fault state of the protocol.
disabled	3	The port shall not place any messages on its communication path.
listening	4	The port is waiting for the announceReceiptTimeout to expire or to receive an Announce message from a master.
preMaster	5	The port shall behave in all respects as though it were in the MASTER state except that it shall not place any messages on its communication path except for Pdelay_Req, Pdelay_Resp, Pdelay_Resp_Follow_Up, signaling, or management messages.
master	6	The port is behaving as a master port.
passive	7	The port shall not place any messages

on its communication path except for Pdelay_Req, Pdelay_Resp, Pdelay_Resp_Follow_Up, or signaling messages, or management messages that are a required response to another management message

uncalibrated	8	The local port is preparing to synchronize to the master port.
slave	9	The port is synchronizing to the selected master port."

REFERENCE "Section 8.2.5.3.1 portState and 9.2.5 of [IEEE 1588-2008]"

SYNTAX INTEGER {
initializing(1),
faulty(2),
disabled(3),
listening(4),
preMaster(5),
master(6),
passive(7),
uncalibrated(8),
slave(9)
}

ClockPortTransportTypeAddress ::= TEXTUAL-CONVENTION

DISPLAY-HINT "255a"
STATUS current
DESCRIPTION

"The Clock port transport protocol address used for this communication between the clock nodes. This is a string corresponding to the address type as specified by the Transport type used. The transport types can be defined elsewhere, in addition to the ones defined in this document. This can be address of type IP version 4, IP version 6, Ethernet, DeviceNET, ControlNET and IEC61158."

REFERENCE "Annex D (IPv4), Annex E (IPv6), Annex F (Ethernet), Annex G (DeviceNET), Annex H (ControlNET) and Annex I (IEC61158) of [IEEE 1588-2008]"

SYNTAX OCTET STRING (SIZE (1..255))

ClockProfileType ::= TEXTUAL-CONVENTION

STATUS current
DESCRIPTION
"Clock Profile used. A profile is the set of allowed Precision Time Protocol (PTP) features applicable to a device."

REFERENCE "Section 3.1.30 and 19.3 PTP profiles of [IEEE 1588-2008]"
SYNTAX INTEGER {
 default(1),
 telecom(2),
 vendorspecific(3)
}

ClockQualityAccuracyType ::= TEXTUAL-CONVENTION

STATUS current
DESCRIPTION
"The ClockQuality as specified in section 5.3.7, 7.6.2.5 and Table 6 of [IEEE 1588-2008].

The following values are not represented in the enumerated values.

0x01-0x1F Reserved
0x32-0x7F Reserved

It is important to note that section 7.1.1 RFC2578 allows for gaps and enumerate values to start with zero when indicated by the protocol."

REFERENCE "Section 5.3.7, 7.6.2.5 and Table 6 of [IEEE 1588-2008]"
SYNTAX INTEGER {
 -- reserved00(0:31), -- 0x00 to 0x1F
 nanoSecond25(32), -- 0x20
 nanoSecond100(33), -- 0x21
 nanoSecond250(34), -- 0x22
 microSec1(35), -- 0x23
 microSec2dot5(36), -- 0x24
 microSec10(37), -- 0x25
 microSec25(38), -- 0x26
 microSec100(39), -- 0x27
 microSec250(40), -- 0x28
 milliSec1(41), -- 0x29

```

        milliSec2dot5(42),    -- 0x2A
        milliSec10(43),      -- 0x2B
        milliSec25(44),      -- 0x2C
        milliSec100(45),     -- 0x2D
        milliSec250(46),     -- 0x2E
        second1(47),         -- 0x2F
        second10(48),        -- 0x30
        secondGreater10(49), -- 0x31
        unknown(254),        -- 0xFE
        -- reserved255(255)  -- 0xFF
    }

```

ClockQualityClassType ::= TEXTUAL-CONVENTION

DISPLAY-HINT "d"

STATUS current

DESCRIPTION

"The ClockQuality as specified in section 5.3.7, 7.6.2.4 and Table 5 of [IEEE 1588-2008]."

Value	Description
0	Reserved to enable compatibility with future versions.
1-5	Reserved
6	Shall designate a clock that is synchronized to a primary reference time source. The timescale distributed shall be PTP. A clockClass 6 clock shall not be a slave to another clock in the domain.
7	Shall designate a clock that has previously been designated as clockClass 6 but that has lost the ability to synchronize to a primary reference time source and is in holdover mode and within holdover specifications. The timescale distributed shall be PTP. A clockClass 7 clock shall not be a slave to another clock in the domain.
8	Reserved.
9-10	Reserved to enable compatibility with future versions.
11-12	Reserved.
13	Shall designate a clock that is synchronized to an application-specific source of time.

- The timescale distributed shall be ARB. A clockClass 13 clock shall not be a slave to another clock in the domain.
- 14 Shall designate a clock that has previously been designated as clockClass 13 but that has lost the ability to synchronize to an application-specific source of time and is in holdover mode and within holdover specifications. The timescale distributed shall be ARB. A clockClass 14 clock shall not be a slave to another clock in the domain.
 - 15-51 Reserved.
 - 52 Degradation alternative A for a clock of clockClass 7 that is not within holdover specification. A clock of clockClass 52 shall not be a slave to another clock in the domain.
 - 53-57 Reserved.
 - 58 Degradation alternative A for a clock of clockClass 14 that is not within holdover specification. A clock of clockClass 58 shall not be a slave to another clock in the domain.
 - 59-67 Reserved.
 - 68-122 For use by alternate PTP profiles.
 - 123-127 Reserved.
 - 128-132 Reserved.
 - 133-170 For use by alternate PTP profiles.
 - 171-186 Reserved.
 - 187 Degradation alternative B for a clock of clockClass 7 that is not within holdover specification. A clock of clockClass 187 may be a slave to another clock in the domain.
 - 188-192 Reserved.
 - 193 Degradation alternative B for a clock of clockClass 14 that is not within holdover specification. A clock of clockClass 193 may be a slave to another clock in the domain.
 - 194-215 Reserved.
 - 216-232 For use by alternate PTP profiles.
 - 233-247 Reserved.
 - 248 Default. This clockClass shall be used if none of the other clockClass definitions apply.

249-250 Reserved.
 251 Reserved for version 1 compatibility; see Clause 18.
 252-254 Reserved.
 255 Shall be the clockClass of a slave-only clock; see 9.2.2."

REFERENCE "Section 5.3.7, 7.6.2.4 and Table 5 of [IEEE 1588-2008]."

SYNTAX Unsigned32 (0..255)

ClockRoleType ::= TEXTUAL-CONVENTION

STATUS current

DESCRIPTION

"The Clock Role. The protocol generates a Master Slave relationship among the clocks in the system.

Clock Role	Value	Description
Master clock	1	A clock that is the source of time to which all other clocks on that path synchronize.
Slave clock	2	A clock which synchronizes to another clock (master)."

SYNTAX INTEGER {
 master(1),
 slave(2)
 }

ClockStateType ::= TEXTUAL-CONVENTION

STATUS current

DESCRIPTION

"The clock state returned by PTP engine.

Clock State	Value	Description
Freerun state	1	Applies to a slave device that is not locked to a master. This is the initial state a slave starts out with when it is not getting any PTP packets from the master or because of some other input error (erroneous packets, etc).

- Holdover state 2 In this state the slave device is locked to a master but communication with the master has been lost or the timestamps in the ntp packets are incorrect. Since the slave was locked to the master, it can run in this state, with similar accuracy for some time. If communication with the master is not restored for an extended period (dependent on the clock implementation), the device should move to the FREERUN state.
- Acquiring state 3 The slave device is receiving packets from a master and is trying to acquire a lock.
- Freq_locked state 4 Slave device is locked to the Master with respect to frequency, but not phase aligned
- Phase_aligned state 5 Locked to the master with respect to frequency and phase."

```
SYNTAX          INTEGER {
                    freerun(1),
                    holdover(2),
                    acquiring(3),
                    frequencyLocked(4),
                    phaseAligned(5)
                  }
```

ClockTimeInterval ::= TEXTUAL-CONVENTION

```
DISPLAY-HINT    "255a"
STATUS          current
DESCRIPTION
```

"This textual convention corresponds to the TimeInterval structure indicated in section 5.3.2 of [IEEE 1588-2008]. It will be presented in the form of a character array. Network byte order is assumed.

The TimeInterval type represents time intervals.

```
struct TimeInterval
```

```

    {
      Integer64 scaledNanoseconds;
    };

```

The scaledNanoseconds member is the time interval expressed in units of nanoseconds and multiplied by 2**16.

Positive or negative time intervals outside the maximum range of this data type shall be encoded as the largest positive and negative values of the data type, respectively.

For example, 2.5 ns is expressed as 0000 0000 0002 8000 in Base16."

REFERENCE

"Section 5.3.2 and section 7.7.2.1 Timer interval specification of [IEEE 1588-2008]"

SYNTAX OCTET STRING (SIZE (1..255))

ClockTimeSourceType ::= TEXTUAL-CONVENTION

STATUS current

DESCRIPTION

"The ClockQuality as specified in section 5.3.7, 7.6.2.6 and Table 7 of [IEEE 1588-2008].

The following values are not represented in the enumerated values.

0xF0-0xFE For use by alternate PTP profiles

0xFF Reserved

It is important to note that section 7.1.1 RFC2578 allows for gaps and enumerate values to start with zero when indicated by the protocol."

REFERENCE "Section 5.3.7, 7.6.2.6 and Table 7 of [IEEE 1588-2008]."

SYNTAX INTEGER {
 atomicClock(16), -- 0x10
 gps(32), -- 0x20
 terrestrialRadio(48), -- 0x22
 ntp(64), -- 0x40
 ntp(80), -- 0x50

```
        handSet(96), -- 0x60
        other(144), -- 0x90
        internalOscillator(160) -- 0xA0
    }

ClockTxModeType ::= TEXTUAL-CONVENTION
    STATUS          current
    DESCRIPTION
        "Transmission mode.

        Unicast:      Using unicast communication channel.
        Multicast:    Using Multicast communication channel.
        multicast-mix: Using multicast-unicast communication channel"
    SYNTAX          INTEGER {
        unicast(1),
        multicast(2),
        multicastmix(3)
    }

ClockType ::= TEXTUAL-CONVENTION
    STATUS          current
    DESCRIPTION
        "The clock types as defined in the MIB module description."

    REFERENCE      "Section 6.5.1 of [IEEE 1588-2008]."
    SYNTAX          INTEGER {
        ordinaryClock(1),
        boundaryClock(2),
        transparentClock(3),
        boundaryNode(4)
    }

ptpbaseMIBNotifs OBJECT IDENTIFIER
    ::= { ptpbaseMIB 0 }

ptpbaseMIBObjects OBJECT IDENTIFIER
    ::= { ptpbaseMIB 1 }

ptpbaseMIBConformance OBJECT IDENTIFIER
    ::= { ptpbaseMIB 2 }

ptpbaseMIBSystemInfo OBJECT IDENTIFIER
    ::= { ptpbaseMIBObjects 1 }
```

```
ptpbasesystemClockInfo OBJECT IDENTIFIER
 ::= { ptpbaseMIBObjects 2 }

ptpbasesystemTable OBJECT-TYPE
    SYNTAX          SEQUENCE OF PtpbaseSystemEntry
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "Table of count information about the PTP system for all
        domains."
    ::= { ptpbaseMIBSystemInfo 1 }

ptpbasesystemEntry OBJECT-TYPE
    SYNTAX          PtpbaseSystemEntry
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "An entry in the table, containing count information about a
        single domain. New row entries are added when the PTP clock for
        this domain is configured, while the unconfiguration of the PTP
        clock removes it."
    INDEX           {
                    ptpDomainIndex,
                    ptpInstanceIndex
                  }
    ::= { ptpbaseSystemTable 1 }

PtpbaseSystemEntry ::= SEQUENCE {
    ptpDomainIndex      ClockDomainType,
    ptpInstanceIndex    ClockInstanceType,
    ptpDomainClockPortsTotal Gauge32
}

ptpDomainIndex OBJECT-TYPE
    SYNTAX          ClockDomainType
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the domain number used to create logical
        group of PTP devices. The Clock Domain is a logical group of
        clocks and devices that synchronize with each other using the
```

PTP protocol.

0	Default domain
1	Alternate domain 1
2	Alternate domain 2
3	Alternate domain 3
4 - 127	User-defined domains
128 - 255	Reserved"

::= { ptpbaseSystemEntry 1 }

ptpInstanceIndex OBJECT-TYPE

SYNTAX	ClockInstanceType
MAX-ACCESS	not-accessible
STATUS	current

DESCRIPTION

"This object specifies the instance of the Clock for this domain."

::= { ptpbaseSystemEntry 2 }

ptpDomainClockPortsTotal OBJECT-TYPE

SYNTAX	Gauge32
UNITS	"ntp ports"
MAX-ACCESS	read-only
STATUS	current

DESCRIPTION

"This object specifies the total number of clock ports configured within a domain in the system."

::= { ptpbaseSystemEntry 3 }

ptpbaseSystemDomainTable OBJECT-TYPE

SYNTAX	SEQUENCE OF PtpbaseSystemDomainEntry
MAX-ACCESS	not-accessible
STATUS	current

DESCRIPTION

"Table of information about the PTP system for all clock modes -- ordinary, boundary or transparent."

::= { ptpbaseMIBSystemInfo 2 }

ptpbaseSystemDomainEntry OBJECT-TYPE

SYNTAX	PtpbaseSystemDomainEntry
MAX-ACCESS	not-accessible

```
STATUS          current
DESCRIPTION
  "An entry in the table, containing information about a single
  clock mode for the PTP system. A row entry gets added when PTP
  clocks are configured on the router."
INDEX           { ptpbaseSystemDomainClockTypeIndex }
 ::= { ptpbaseSystemDomainTable 1 }

PtpbaseSystemDomainEntry ::= SEQUENCE {
  ptpbaseSystemDomainClockTypeIndex ClockType,
  ptpbaseSystemDomainTotals         Unsigned32
}

ptpbaseSystemDomainClockTypeIndex OBJECT-TYPE
SYNTAX          ClockType
MAX-ACCESS      not-accessible
STATUS          current
DESCRIPTION
  "This object specifies the clock type as defined in the
  Textual convention description."
 ::= { ptpbaseSystemDomainEntry 1 }

ptpbaseSystemDomainTotals OBJECT-TYPE
SYNTAX          Unsigned32
UNITS           "domains"
MAX-ACCESS      read-only
STATUS          current
DESCRIPTION
  "This object specifies the total number of PTP domains for this
  particular clock type configured in this node."
 ::= { ptpbaseSystemDomainEntry 2 }

ptpbaseSystemProfile OBJECT-TYPE
SYNTAX          ClockProfileType
MAX-ACCESS      read-only
STATUS          current
DESCRIPTION
  "This object specifies the PTP Profile implemented on the
  system."
REFERENCE       "Section 19.3 PTP profiles of [IEEE 1588-2008]"
 ::= { ptpbaseMIBSystemInfo 3 }
```

```

ptpbasedClockCurrentDSTable OBJECT-TYPE
    SYNTAX          SEQUENCE OF PtpbasedClockCurrentDSEntry
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "Table of information about the PTP clock Current Datasets for
        all domains."
    ::= { ptpbaseMIBClockInfo 1 }

ptpbasedClockCurrentDSEntry OBJECT-TYPE
    SYNTAX          PtpbasedClockCurrentDSEntry
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "An entry in the table, containing information about a single
        PTP clock Current Datasets for a domain."
    REFERENCE
        "1588 Version 2.0 Section 8.2.2 currentDS data set member
        specifications of [IEEE 1588-2008]"
    INDEX
        {
            ptpbasedClockCurrentDSDomainIndex,
            ptpbasedClockCurrentDSClockTypeIndex,
            ptpbasedClockCurrentDSInstanceIndex
        }
    ::= { ptpbasedClockCurrentDSTable 1 }

PtpbasedClockCurrentDSEntry ::= SEQUENCE {
    ptpbasedClockCurrentDSDomainIndex      ClockDomainType,
    ptpbasedClockCurrentDSClockTypeIndex   ClockType,
    ptpbasedClockCurrentDSInstanceIndex    ClockInstanceType,
    ptpbasedClockCurrentDSStepsRemoved     Unsigned32,
    ptpbasedClockCurrentDSOffsetFromMaster ClockTimeInterval,
    ptpbasedClockCurrentDSMeanPathDelay    ClockTimeInterval
}

ptpbasedClockCurrentDSDomainIndex OBJECT-TYPE
    SYNTAX          ClockDomainType
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the domain number used to create logical
        group of PTP devices."

```

```
::= { ptpbaseClockCurrentDSEntry 1 }
```

```
ptpbaseClockCurrentDSClockTypeIndex OBJECT-TYPE
```

```
SYNTAX          ClockType
```

```
MAX-ACCESS      not-accessible
```

```
STATUS          current
```

```
DESCRIPTION
```

```
"This object specifies the clock type as defined in the  
Textual convention description."
```

```
::= { ptpbaseClockCurrentDSEntry 2 }
```

```
ptpbaseClockCurrentDSInstanceIndex OBJECT-TYPE
```

```
SYNTAX          ClockInstanceType
```

```
MAX-ACCESS      not-accessible
```

```
STATUS          current
```

```
DESCRIPTION
```

```
"This object specifies the instance of the clock for this clock  
type in the given domain."
```

```
::= { ptpbaseClockCurrentDSEntry 3 }
```

```
ptpbaseClockCurrentDSStepsRemoved OBJECT-TYPE
```

```
SYNTAX          Unsigned32
```

```
UNITS           "Steps"
```

```
MAX-ACCESS      read-only
```

```
STATUS          current
```

```
DESCRIPTION
```

```
"The current clock dataset StepsRemoved value.
```

```
This object specifies the distance measured by the number of  
Boundary clocks between the local clock and the Foreign master  
as indicated in the stepsRemoved field of Announce messages."
```

```
REFERENCE       "1588 Version 2.0 Section 8.2.2.2 stepsRemoved"
```

```
::= { ptpbaseClockCurrentDSEntry 4 }
```

```
ptpbaseClockCurrentDSOffsetFromMaster OBJECT-TYPE
```

```
SYNTAX          ClockTimeInterval
```

```
UNITS           "Time Interval"
```

```
MAX-ACCESS      read-only
```

```
STATUS          current
```

```
DESCRIPTION
```

```
"This object specifies the current clock dataset ClockOffset  
value. The value of the computation of the offset in time  
between a slave and a master clock."
```

REFERENCE "1588 Version 2.0 Section 8.2.2.3 of
[IEEE 1588-2008]"
 ::= { ptpbaseClockCurrentDSEntry 5 }

ptpbaseClockCurrentDSMeanPathDelay OBJECT-TYPE

SYNTAX ClockTimeInterval

UNITS "Time Interval"

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies the current clock dataset
MeanPathDelay value.

The mean path delay between a pair of ports as measure by the
delay request-response mechanism."

REFERENCE "1588 Version 2.0 Section 8.2.2.4 mean path delay"
 ::= { ptpbaseClockCurrentDSEntry 6 }

ptpbaseClockParentDSTable OBJECT-TYPE

SYNTAX SEQUENCE OF PtpbaseClockParentDSEntry

MAX-ACCESS not-accessible

STATUS current

DESCRIPTION

"Table of information about the PTP clock Parent Datasets for
all domains."

::= { ptpbaseMIBClockInfo 2 }

ptpbaseClockParentDSEntry OBJECT-TYPE

SYNTAX PtpbaseClockParentDSEntry

MAX-ACCESS not-accessible

STATUS current

DESCRIPTION

"An entry in the table, containing information about a single
PTP clock Parent Datasets for a domain."

REFERENCE

"Section 8.2.3 parentDS data set member specifications of
[IEEE 1588-2008]"

INDEX

{
 ptpbaseClockParentDSDomainIndex,
 ptpbaseClockParentDSClockTypeIndex,
 ptpbaseClockParentDSInstanceIndex

```

    }
    ::= { ptpbaseClockParentDSTable 1 }

PtpbaseClockParentDSEntry ::= SEQUENCE {
    ptpbaseClockParentDSDomainIndex      ClockDomainType,
    ptpbaseClockParentDSClockTypeIndex   ClockType,
    ptpbaseClockParentDSInstanceIndex    ClockInstanceType,
    ptpbaseClockParentDSParentPortIdentity OCTET STRING,
    ptpbaseClockParentDSParentStats      TruthValue,
    ptpbaseClockParentDSOffset           ClockIntervalBase2,
    ptpbaseClockParentDSClockPhChRate    Integer32,
    ptpbaseClockParentDSGMClockIdentity  ClockIdentity,
    ptpbaseClockParentDSGMClockPriority1  Unsigned32,
    ptpbaseClockParentDSGMClockPriority2  Unsigned32,
    ptpbaseClockParentDSGMClockQualityClass ClockQualityClassType,
    ptpbaseClockParentDSGMClockQualityAccuracy ClockQualityAccuracyType,
    ptpbaseClockParentDSGMClockQualityOffset Unsigned32
}

ptpbaseClockParentDSDomainIndex OBJECT-TYPE
    SYNTAX      ClockDomainType
    MAX-ACCESS  not-accessible
    STATUS      current
    DESCRIPTION
        "This object specifies the domain number used to create logical
        group of PTP devices."
    ::= { ptpbaseClockParentDSEntry 1 }

ptpbaseClockParentDSClockTypeIndex OBJECT-TYPE
    SYNTAX      ClockType
    MAX-ACCESS  not-accessible
    STATUS      current
    DESCRIPTION
        "This object specifies the clock type as defined in the
        Textual convention description."
    ::= { ptpbaseClockParentDSEntry 2 }

ptpbaseClockParentDSInstanceIndex OBJECT-TYPE
    SYNTAX      ClockInstanceType
    MAX-ACCESS  not-accessible
    STATUS      current
    DESCRIPTION
        "This object specifies the instance of the clock for this clock

```

type in the given domain."
 ::= { ptpbaseClockParentDSEntry 3 }

ptpbaseClockParentDSParentPortIdentity OBJECT-TYPE

SYNTAX OCTET STRING(SIZE(1..256))

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies the value of portIdentity of the port on the master that issues the Sync messages used in synchronizing this clock."

REFERENCE

"Section 8.2.3.2 parentDS.parentPortIdentity of [IEEE 1588-2008]"

::= { ptpbaseClockParentDSEntry 4 }

ptpbaseClockParentDSParentStats OBJECT-TYPE

SYNTAX TruthValue

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies the Parent Dataset ParentStats value.

This value indicates whether the values of ParentDSOffset and ParentDSClockPhChRate have been measured and are valid. A TRUE value shall indicate valid data."

REFERENCE "Section 8.2.3.3 parentDS.parentStats of [IEEE 1588-2008]"

::= { ptpbaseClockParentDSEntry 5 }

ptpbaseClockParentDSOffset OBJECT-TYPE

SYNTAX ClockIntervalBase2 (-128..127)

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies the Parent Dataset ParentOffsetScaledLogVariance value.

This value is the variance of the parent clocks phase as measured by the local clock."

REFERENCE

"Section 8.2.3.4 parentDS.observedParentOffsetScaledLogVariance

```
[IEEE 1588-2008]"  
 ::= { ptpbaseClockParentDSEntry 6 }
```

ptpbaseClockParentDSClockPhChRate OBJECT-TYPE

SYNTAX Integer32

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies the clock's parent dataset
ParentClockPhaseChangeRate value.

This value is an estimate of the parent clocks phase change
rate as measured by the slave clock."

REFERENCE

"Section 8.2.3.5

parentDS.observedParentClockPhaseChangeRate of

[IEEE 1588-2008]"

```
 ::= { ptpbaseClockParentDSEntry 7 }
```

ptpbaseClockParentDSGMClockIdentity OBJECT-TYPE

SYNTAX ClockIdentity

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies the parent dataset Grandmaster clock
identity."

REFERENCE

"Section 8.2.3.6 parentDS.grandmasterIdentity of

[IEEE 1588-2008]"

```
 ::= { ptpbaseClockParentDSEntry 8 }
```

ptpbaseClockParentDSGMClockPriority1 OBJECT-TYPE

SYNTAX Unsigned32

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies the parent dataset Grandmaster clock
priority1."

REFERENCE

"Section 8.2.3.8 parentDS.grandmasterPriority1 of

[IEEE 1588-2008]"

```
 ::= { ptpbaseClockParentDSEntry 9 }
```

```
ptpbasedClockParentDSGMClockPriority2 OBJECT-TYPE
    SYNTAX          Unsigned32
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the parent dataset grandmaster clock
        priority2."
    REFERENCE
        "Section 8.2.3.9 parentDS.grandmasterPriority2 of
        [IEEE 1588-2008]"
    ::= { ptpbasedClockParentDSEntry 10 }

ptpbasedClockParentDSGMClockQualityClass OBJECT-TYPE
    SYNTAX          ClockQualityClassType (0..255)
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the parent dataset grandmaster clock
        quality class."
    REFERENCE
        "Section 8.2.3.7 parentDS.grandmasterClockQuality of
        [IEEE 1588-2008]"
    ::= { ptpbasedClockParentDSEntry 11 }

ptpbasedClockParentDSGMClockQualityAccuracy OBJECT-TYPE
    SYNTAX          ClockQualityAccuracyType
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the parent dataset grandmaster clock
        quality accuracy."
    REFERENCE
        "Section 8.2.3.7 parentDS.grandmasterClockQuality of
        [IEEE 1588-2008]"
    ::= { ptpbasedClockParentDSEntry 12 }

ptpbasedClockParentDSGMClockQualityOffset OBJECT-TYPE
    SYNTAX          Unsigned32
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the parent dataset grandmaster clock
        quality offset."
```

REFERENCE

"Section 8.2.3.7 parentDS.grandmasterClockQuality of
[IEEE 1588-2008]"

::= { ptpbaseClockParentDSEntry 13 }

ptpbaseClockDefaultDSTable OBJECT-TYPE

SYNTAX SEQUENCE OF PtpbaseClockDefaultDSEntry

MAX-ACCESS not-accessible

STATUS current

DESCRIPTION

"Table of information about the PTP clock Default Datasets for
all domains."

::= { ptpbaseMIBClockInfo 3 }

ptpbaseClockDefaultDSEntry OBJECT-TYPE

SYNTAX PtpbaseClockDefaultDSEntry

MAX-ACCESS not-accessible

STATUS current

DESCRIPTION

"An entry in the table, containing information about a single
PTP clock Default Datasets for a domain."

INDEX {
 ptpbaseClockDefaultDSDomainIndex,
 ptpbaseClockDefaultDSClockTypeIndex,
 ptpbaseClockDefaultDSInstanceIndex
}

::= { ptpbaseClockDefaultDSTable 1 }

PtpbaseClockDefaultDSEntry ::= SEQUENCE {
 ptpbaseClockDefaultDSDomainIndex ClockDomainType,
 ptpbaseClockDefaultDSClockTypeIndex ClockType,
 ptpbaseClockDefaultDSInstanceIndex ClockInstanceType,
 ptpbaseClockDefaultDSTwoStepFlag TruthValue,
 ptpbaseClockDefaultDSClockIdentity ClockIdentity,
 ptpbaseClockDefaultDSPriority1 Unsigned32,
 ptpbaseClockDefaultDSPriority2 Unsigned32,
 ptpbaseClockDefaultDSSlaveOnly TruthValue,
 ptpbaseClockDefaultDSQualityClass ClockQualityClassType,
 ptpbaseClockDefaultDSQualityAccuracy ClockQualityAccuracyType,
 ptpbaseClockDefaultDSQualityOffset Integer32
}

```
ptpbasedefaultDSDomainIndex OBJECT-TYPE
    SYNTAX          ClockDomainType
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the domain number used to create logical
        group of PTP devices."
    ::= { ptpbasedefaultDSEntry 1 }

ptpbasedefaultDSClockTypeIndex OBJECT-TYPE
    SYNTAX          ClockType
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the clock type as defined in the
        Textual convention description."
    ::= { ptpbasedefaultDSEntry 2 }

ptpbasedefaultDSInstanceIndex OBJECT-TYPE
    SYNTAX          ClockInstanceType
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the instance of the clock for this clock
        type in the given domain."
    ::= { ptpbasedefaultDSEntry 3 }

ptpbasedefaultDSTwoStepFlag OBJECT-TYPE
    SYNTAX          TruthValue
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies whether the Two Step process is used."
    ::= { ptpbasedefaultDSEntry 4 }

ptpbasedefaultDSClockIdentity OBJECT-TYPE
    SYNTAX          ClockIdentity
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the default Datasets clock identity."
    ::= { ptpbasedefaultDSEntry 5 }
```

```
ptpbasedefaultDSPriority1 OBJECT-TYPE
    SYNTAX      Unsigned32
    MAX-ACCESS  read-only
    STATUS      current
    DESCRIPTION
        "This object specifies the default Datasets clock Priority1."
    ::= { ptpbasedefaultDSEntry 6 }
```

```
ptpbasedefaultDSPriority2 OBJECT-TYPE
    SYNTAX      Unsigned32
    MAX-ACCESS  read-only
    STATUS      current
    DESCRIPTION
        "This object specifies the default Datasets clock Priority2."
    ::= { ptpbasedefaultDSEntry 7 }
```

```
ptpbasedefaultDSSlaveOnly OBJECT-TYPE
    SYNTAX      TruthValue
    MAX-ACCESS  read-only
    STATUS      current
    DESCRIPTION
        "Whether the SlaveOnly flag is set."
    ::= { ptpbasedefaultDSEntry 8 }
```

```
ptpbasedefaultDSQualityClass OBJECT-TYPE
    SYNTAX      ClockQualityClassType (0..255)
    MAX-ACCESS  read-only
    STATUS      current
    DESCRIPTION
        "This object specifies the default dataset Quality Class."
    ::= { ptpbasedefaultDSEntry 9 }
```

```
ptpbasedefaultDSQualityAccuracy OBJECT-TYPE
    SYNTAX      ClockQualityAccuracyType
    MAX-ACCESS  read-only
    STATUS      current
    DESCRIPTION
        "This object specifies the default dataset Quality Accuracy."
    ::= { ptpbasedefaultDSEntry 10 }
```

```
ptpbasedefaultDSQualityOffset OBJECT-TYPE
    SYNTAX      Integer32
```

```

MAX-ACCESS      read-only
STATUS          current
DESCRIPTION
    "This object specifies the default dataset Quality offset."
 ::= { ptpbaseClockDefaultDSEntry 11 }

```

```

ptpbaseClockRunningTable OBJECT-TYPE
SYNTAX          SEQUENCE OF PtpbaseClockRunningEntry
MAX-ACCESS      not-accessible
STATUS          current
DESCRIPTION
    "Table of information about the PTP clock Running Datasets for
    all domains."
 ::= { ptpbaseMIBClockInfo 4 }

```

```

ptpbaseClockRunningEntry OBJECT-TYPE
SYNTAX          PtpbaseClockRunningEntry
MAX-ACCESS      not-accessible
STATUS          current
DESCRIPTION
    "An entry in the table, containing information about a single
    PTP clock running Datasets for a domain."
INDEX          {
                ptpbaseClockRunningDomainIndex,
                ptpbaseClockRunningClockTypeIndex,
                ptpbaseClockRunningInstanceIndex
            }
 ::= { ptpbaseClockRunningTable 1 }

```

```

PtpbaseClockRunningEntry ::= SEQUENCE {
    ptpbaseClockRunningDomainIndex      ClockDomainType,
    ptpbaseClockRunningClockTypeIndex  ClockType,
    ptpbaseClockRunningInstanceIndex   ClockInstanceType,
    ptpbaseClockRunningState           ClockStateType,
    ptpbaseClockRunningPacketsSent     Counter64,
    ptpbaseClockRunningPacketsReceived Counter64
}

```

```

ptpbaseClockRunningDomainIndex OBJECT-TYPE
SYNTAX          ClockDomainType
MAX-ACCESS      not-accessible

```

STATUS current
DESCRIPTION
"This object specifies the domain number used to create logical
group of PTP devices."
 ::= { ptpbaseClockRunningEntry 1 }

ptpbaseClockRunningClockTypeIndex OBJECT-TYPE

SYNTAX ClockType
MAX-ACCESS not-accessible
STATUS current
DESCRIPTION
"This object specifies the clock type as defined in the
Textual convention description."
 ::= { ptpbaseClockRunningEntry 2 }

ptpbaseClockRunningInstanceIndex OBJECT-TYPE

SYNTAX ClockInstanceType
MAX-ACCESS not-accessible
STATUS current
DESCRIPTION
"This object specifies the instance of the clock for this clock
type in the given domain."
 ::= { ptpbaseClockRunningEntry 3 }

ptpbaseClockRunningState OBJECT-TYPE

SYNTAX ClockStateType
MAX-ACCESS read-only
STATUS current
DESCRIPTION
"This object specifies the Clock state returned by PTP engine
which was described earlier.

Freerun state. Applies to a slave device that is not locked to
a master. This is the initial state a slave starts out with
when it is not getting any PTP packets from the master, or
because of some other input error (erroneous packets, etc).

Holdover state. In this state the slave device is locked to a
master but communication with the master has been lost or the
timestamps in the ptp packets are incorrect. Since the
slave was previously locked to the master, it can run in this
state, with similar accuracy for some time. If communication
with the master is not restored for an extended period

(dependent on the clock implementation), the device should move to the FREERUN state.

Acquiring state. The slave device is receiving packets from a master and is trying to acquire a lock.

Freq_locked state. Slave device is locked to the Master with respect to frequency, but not phase aligned

Phase_aligned state. Locked to the master with respect to frequency and phase."

```
::= { ptpbaseClockRunningEntry 4 }
```

ptpbaseClockRunningPacketsSent OBJECT-TYPE

SYNTAX Counter64

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies the total number of all packet Unicast and multicast that have been sent out for this clock in this domain for this type."

```
::= { ptpbaseClockRunningEntry 5 }
```

ptpbaseClockRunningPacketsReceived OBJECT-TYPE

SYNTAX Counter64

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies the total number of all packet Unicast and multicast that have been received for this clock in this domain for this type."

```
::= { ptpbaseClockRunningEntry 6 }
```

ptpbaseClockTimePropertiesDSTable OBJECT-TYPE

SYNTAX SEQUENCE OF PtpbaseClockTimePropertiesDSEntry

MAX-ACCESS not-accessible

STATUS current

DESCRIPTION

"Table of information about the PTP clock time properties datasets for all domains."

```
::= { ptpbaseMIBClockInfo 5 }
```

```

ptpbasedClockTimePropertiesDSEntry OBJECT-TYPE
    SYNTAX          PtpbasedClockTimePropertiesDSEntry
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "An entry in the table, containing information about a single
        PTP clock timeproperties Datasets for a domain."
    REFERENCE      "Section 8.2.4 of [IEEE 1588-2008]"
    INDEX          {
        ptpbasedClockTimePropertiesDSDomainIndex,
        ptpbasedClockTimePropertiesDSClockTypeIndex,
        ptpbasedClockTimePropertiesDSInstanceIndex
    }
    ::= { ptpbasedClockTimePropertiesDSTable 1 }

PtpbasedClockTimePropertiesDSEntry ::= SEQUENCE {
    ptpbasedClockTimePropertiesDSDomainIndex      ClockDomainType,
    ptpbasedClockTimePropertiesDSClockTypeIndex   ClockType,
    ptpbasedClockTimePropertiesDSInstanceIndex    ClockInstanceType,
    ptpbasedClockTimePropertiesDSCurrentUTCOffsetValid TruthValue,
    ptpbasedClockTimePropertiesDSCurrentUTCOffset Integer32,
    ptpbasedClockTimePropertiesDSLeap59          TruthValue,
    ptpbasedClockTimePropertiesDSLeap61         TruthValue,
    ptpbasedClockTimePropertiesDSTimeTraceable   TruthValue,
    ptpbasedClockTimePropertiesDSFreqTraceable   TruthValue,
    ptpbasedClockTimePropertiesDSPTPTimescale    TruthValue,
    ptpbasedClockTimePropertiesDSSource          ClockTimeSourceType
}

ptpbasedClockTimePropertiesDSDomainIndex OBJECT-TYPE
    SYNTAX          ClockDomainType
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the domain number used to create logical
        group of PTP devices."
    ::= { ptpbasedClockTimePropertiesDSEntry 1 }

ptpbasedClockTimePropertiesDSClockTypeIndex OBJECT-TYPE
    SYNTAX          ClockType
    MAX-ACCESS      not-accessible
    STATUS          current

```

DESCRIPTION

"This object specifies the clock type as defined in the Textual convention description."

::= { ptpbaseClockTimePropertiesDSEntry 2 }

ptpbaseClockTimePropertiesDSInstanceIndex OBJECT-TYPE

SYNTAX ClockInstanceType

MAX-ACCESS not-accessible

STATUS current

DESCRIPTION

"This object specifies the instance of the clock for this clock type in the given domain."

::= { ptpbaseClockTimePropertiesDSEntry 3 }

ptpbaseClockTimePropertiesDSCurrentUTCOffsetValid OBJECT-TYPE

SYNTAX TruthValue

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies the timeproperties dataset value of whether current UTC offset is valid."

REFERENCE "Section 8.2.4.2 of [IEEE 1588-2008]"

::= { ptpbaseClockTimePropertiesDSEntry 4 }

ptpbaseClockTimePropertiesDSCurrentUTCOffset OBJECT-TYPE

SYNTAX Integer32

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies the timeproperties dataset value of current UTC offset."

In PTP systems whose epoch is the PTP epoch, the value of timePropertiesDS.currentUtcOffset is the offset between TAI and UTC; otherwise the value has no meaning. The value shall be in units of seconds.

The initialization value shall be selected as follows:

- a) If the timePropertiesDS.ptpTimescale (see 8.2.4.8) is TRUE, the value is the value obtained from a primary reference if the value is known at the time of initialization, else,
- b) The value shall be the current number of leap seconds (7.2.3) when the node is designed."

REFERENCE "Section 8.2.4.3 of [IEEE 1588-2008]"
 ::= { ptpbaseClockTimePropertiesDSEntry 5 }

ptpbaseClockTimePropertiesDSLeap59 OBJECT-TYPE
SYNTAX TruthValue
MAX-ACCESS read-only
STATUS current
DESCRIPTION
 "This object specifies the Leap59 value in the clock Current
 Dataset."
REFERENCE "Section 8.2.4.4 of [IEEE 1588-2008]"
 ::= { ptpbaseClockTimePropertiesDSEntry 6 }

ptpbaseClockTimePropertiesDSLeap61 OBJECT-TYPE
SYNTAX TruthValue
MAX-ACCESS read-only
STATUS current
DESCRIPTION
 "This object specifies the Leap61 value in the clock Current
 Dataset."
REFERENCE "Section 8.2.4.5 of [IEEE 1588-2008]"
 ::= { ptpbaseClockTimePropertiesDSEntry 7 }

ptpbaseClockTimePropertiesDSTimeTraceable OBJECT-TYPE
SYNTAX TruthValue
MAX-ACCESS read-only
STATUS current
DESCRIPTION
 "This object specifies the Timetraceable value in the clock
 Current Dataset."
REFERENCE "Section 8.2.4.6 of [IEEE 1588-2008]"
 ::= { ptpbaseClockTimePropertiesDSEntry 8 }

ptpbaseClockTimePropertiesDSFreqTraceable OBJECT-TYPE
SYNTAX TruthValue
MAX-ACCESS read-only
STATUS current
DESCRIPTION
 "This object specifies the Frequency Traceable value in the
 clock Current Dataset."
REFERENCE "Section 8.2.4.7 of [IEEE 1588-2008]"
 ::= { ptpbaseClockTimePropertiesDSEntry 9 }

ptpbasedClockTimePropertiesDSPTPTimescale OBJECT-TYPE

SYNTAX TruthValue
MAX-ACCESS read-only
STATUS current
DESCRIPTION
"This object specifies the PTP Timescale value in the clock
Current Dataset."
REFERENCE "Section 8.2.4.8 of [IEEE 1588-2008]"
 ::= { ptpbasedClockTimePropertiesDSEntry 10 }

ptpbasedClockTimePropertiesDSSource OBJECT-TYPE

SYNTAX ClockTimeSourceType
MAX-ACCESS read-only
STATUS current
DESCRIPTION
"This object specifies the Timesource value in the clock Current
Dataset."
REFERENCE "Section 8.2.4.9 of [IEEE 1588-2008]"
 ::= { ptpbasedClockTimePropertiesDSEntry 11 }

ptpbasedClockTransDefaultDSTable OBJECT-TYPE

SYNTAX SEQUENCE OF PtpbasedClockTransDefaultDSEntry
MAX-ACCESS not-accessible
STATUS current
DESCRIPTION
"Table of information about the PTP Transparent clock Default
Datasets for all domains."
 ::= { ptpbasedMIBClockInfo 6 }

ptpbasedClockTransDefaultDSEntry OBJECT-TYPE

SYNTAX PtpbasedClockTransDefaultDSEntry
MAX-ACCESS not-accessible
STATUS current
DESCRIPTION
"An entry in the table, containing information about a single
PTP Transparent clock Default Datasets for a domain."
REFERENCE "Section 8.3.2 of [IEEE 1588-2008]"
INDEX {
 ptpbasedClockTransDefaultDSDomainIndex,
 ptpbasedClockTransDefaultDSInstanceIndex
 }

```
 ::= { ptpbaseClockTransDefaultDSTable 1 }

PtpbaseClockTransDefaultDSEntry ::= SEQUENCE {
    ptpbaseClockTransDefaultDSDomainIndex  ClockDomainType,
    ptpbaseClockTransDefaultDSInstanceIndex ClockInstanceType,
    ptpbaseClockTransDefaultDSClockIdentity ClockIdentity,
    ptpbaseClockTransDefaultDSNumOfPorts   Counter32,
    ptpbaseClockTransDefaultDSDelay        ClockMechanismType,
    ptpbaseClockTransDefaultDSPrimaryDomain ClockDomainType
}

ptpbaseClockTransDefaultDSDomainIndex OBJECT-TYPE
    SYNTAX      ClockDomainType
    MAX-ACCESS  not-accessible
    STATUS      current
    DESCRIPTION
        "This object specifies the domain number used to create logical
        group of PTP devices."
    ::= { ptpbaseClockTransDefaultDSEntry 1 }

ptpbaseClockTransDefaultDSInstanceIndex OBJECT-TYPE
    SYNTAX      ClockInstanceType
    MAX-ACCESS  not-accessible
    STATUS      current
    DESCRIPTION
        "This object specifies the instance of the clock for this clock
        type in the given domain."
    ::= { ptpbaseClockTransDefaultDSEntry 2 }

ptpbaseClockTransDefaultDSClockIdentity OBJECT-TYPE
    SYNTAX      ClockIdentity
    MAX-ACCESS  read-only
    STATUS      current
    DESCRIPTION
        "This object specifies the value of the clockIdentity attribute
        of the local clock."
    REFERENCE   "Section 8.3.2.2.1 of [IEEE 1588-2008]"
    ::= { ptpbaseClockTransDefaultDSEntry 3 }

ptpbaseClockTransDefaultDSNumOfPorts OBJECT-TYPE
    SYNTAX      Counter32
    MAX-ACCESS  read-only
    STATUS      current
```

DESCRIPTION

"This object specifies the number of PTP ports of the device."

REFERENCE "Section 8.3.2.2.2 of [IEEE 1588-2008]"

::= { ptpbaseClockTransDefaultDSEntry 4 }

ptpbaseClockTransDefaultDSDelay OBJECT-TYPE

SYNTAX ClockMechanismType

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object, if the transparent clock is an end-to-end transparent clock, has the value shall be E2E; If the transparent clock is a peer-to-peer transparent clock, the value shall be P2P."

REFERENCE "Section 8.3.2.3.1 of [IEEE 1588-2008]"

::= { ptpbaseClockTransDefaultDSEntry 5 }

ptpbaseClockTransDefaultDSPrimaryDomain OBJECT-TYPE

SYNTAX ClockDomainType

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies the value of the primary syntonization domain. The initialization value shall be 0."

REFERENCE "Section 8.3.2.3.2 of [IEEE 1588-2008]"

::= { ptpbaseClockTransDefaultDSEntry 6 }

ptpbaseClockPortTable OBJECT-TYPE

SYNTAX SEQUENCE OF PtpbaseClockPortEntry

MAX-ACCESS not-accessible

STATUS current

DESCRIPTION

"Table of information about the clock ports for a particular domain."

::= { ptpbaseMIBClockInfo 7 }

ptpbaseClockPortEntry OBJECT-TYPE

SYNTAX PtpbaseClockPortEntry

MAX-ACCESS not-accessible

STATUS current

```

DESCRIPTION
    "An entry in the table, containing information about a single
    clock port."
INDEX
    {
        ptpbaseClockPortDomainIndex,
        ptpbaseClockPortClockTypeIndex,
        ptpbaseClockPortClockInstanceIndex,
        ptpbaseClockPortTablePortNumberIndex
    }
 ::= { ptpbaseClockPortTable 1 }

PtpbaseClockPortEntry ::= SEQUENCE {
    ptpbaseClockPortDomainIndex          ClockDomainType,
    ptpbaseClockPortClockTypeIndex       ClockType,
    ptpbaseClockPortClockInstanceIndex   ClockInstanceType,
    ptpbaseClockPortTablePortNumberIndex ClockPortNumber,
    ptpbaseClockPortName                  DisplayString,
    ptpbaseClockPortRole                   ClockRoleType,
    ptpbaseClockPortSyncTwoStep           TruthValue,
    ptpbaseClockPortCurrentPeerAddressType AutonomousType,
    ptpbaseClockPortCurrentPeerAddress
ClockPortTransportTypeAddress,
    ptpbaseClockPortNumOfAssociatedPorts  Gauge32
}

ptpbaseClockPortDomainIndex OBJECT-TYPE
    SYNTAX          ClockDomainType
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the domain number used to create logical
        group of PTP devices."
    ::= { ptpbaseClockPortEntry 1 }

ptpbaseClockPortClockTypeIndex OBJECT-TYPE
    SYNTAX          ClockType
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the clock type as defined in the
        Textual convention description."
    ::= { ptpbaseClockPortEntry 2 }

```

ptpbasedClockPortClockInstanceIndex OBJECT-TYPE
SYNTAX ClockInstanceType
MAX-ACCESS not-accessible
STATUS current
DESCRIPTION
"This object specifies the instance of the clock for this clock
type in the given domain."
 ::= { ptpbasedClockPortEntry 3 }

ptpbasedClockPortTablePortNumberIndex OBJECT-TYPE
SYNTAX ClockPortNumber
MAX-ACCESS not-accessible
STATUS current
DESCRIPTION
"This object specifies the PTP Portnumber for this port."
 ::= { ptpbasedClockPortEntry 4 }

ptpbasedClockPortName OBJECT-TYPE
SYNTAX DisplayString (SIZE (1..64))
MAX-ACCESS read-only
STATUS current
DESCRIPTION
"This object specifies the PTP clock port name configured on the
router."
 ::= { ptpbasedClockPortEntry 5 }

ptpbasedClockPortRole OBJECT-TYPE
SYNTAX ClockRoleType
MAX-ACCESS read-only
STATUS current
DESCRIPTION
"This object describes the current role (slave/master) of the
port."
 ::= { ptpbasedClockPortEntry 6 }

ptpbasedClockPortSyncTwoStep OBJECT-TYPE
SYNTAX TruthValue
MAX-ACCESS read-only
STATUS current
DESCRIPTION
"This object specifies that two-step clock operation between
the PTP master and slave device is enabled."
 ::= { ptpbasedClockPortEntry 7 }

ptpbaseClockPortCurrentPeerAddressType OBJECT-TYPE

SYNTAX AutonomousType

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies the current peer's network address type used for PTP communication."

::= { ptpbaseClockPortEntry 8 }

ptpbaseClockPortCurrentPeerAddress OBJECT-TYPE

SYNTAX ClockPortTransportTypeAddress

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies the current peer's network address used for PTP communication."

::= { ptpbaseClockPortEntry 9 }

ptpbaseClockPortNumOfAssociatedPorts OBJECT-TYPE

SYNTAX Gauge32

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies -

For a master port - the number of PTP slave sessions (peers) associated with this PTP port.

For a slave port - the number of masters available to this slave port (might or might not be peered)."

::= { ptpbaseClockPortEntry 10 }

ptpbaseClockPortDSTable OBJECT-TYPE

SYNTAX SEQUENCE OF PtpbaseClockPortDSEntry

MAX-ACCESS not-accessible

STATUS current

DESCRIPTION

"Table of information about the clock ports dataset for a particular domain."

::= { ptpbaseMIBClockInfo 8 }

ptpbaseClockPortDSEntry OBJECT-TYPE

```

SYNTAX          PtpbaseClockPortDSEntry
MAX-ACCESS      not-accessible
STATUS          current
DESCRIPTION     "An entry in the table, containing port dataset information for
                a single clock port."
INDEX           {
                ptpbaseClockPortDSDomainIndex,
                ptpbaseClockPortDSClockTypeIndex,
                ptpbaseClockPortDSClockInstanceIndex,
                ptpbaseClockPortDSPortNumberIndex
                }
 ::= { ptpbaseClockPortDSTable 1 }

```

```

PtpbaseClockPortDSEntry ::= SEQUENCE {
    ptpbaseClockPortDSDomainIndex      ClockDomainType,
    ptpbaseClockPortDSClockTypeIndex   ClockType,
    ptpbaseClockPortDSClockInstanceIndex ClockInstanceType,
    ptpbaseClockPortDSPortNumberIndex  ClockPortNumber,
    ptpbaseClockPortDSName             DisplayString,
    ptpbaseClockPortDSPortIdentity     OCTET STRING,
    ptpbaseClockPortDSlogAnnouncementInterval ClockIntervalBase2,
    ptpbaseClockPortDSAnnounceRctTimeout Integer32,
    ptpbaseClockPortDSlogSyncInterval  ClockIntervalBase2,
    ptpbaseClockPortDSMinDelayReqInterval Integer32,
    ptpbaseClockPortDSPeerDelayReqInterval Integer32,
    ptpbaseClockPortDSDelayMech        ClockMechanismType,
    ptpbaseClockPortDSPeerMeanPathDelay ClockTimeInterval,
    ptpbaseClockPortDSGrantDuration    Unsigned32,
    ptpbaseClockPortDSPTPVersion       Unsigned32
}

```

ptpbaseClockPortDSDomainIndex OBJECT-TYPE

```

SYNTAX          ClockDomainType
MAX-ACCESS      not-accessible
STATUS          current
DESCRIPTION     "This object specifies the domain number used to create logical
                group of PTP devices."
 ::= { ptpbaseClockPortDSEntry 1 }

```

ptpbaseClockPortDSClockTypeIndex OBJECT-TYPE

```

SYNTAX          ClockType

```

MAX-ACCESS not-accessible
STATUS current
DESCRIPTION
"This object specifies the clock type as defined in the
Textual convention description."
 ::= { ptpbaseClockPortDSEntry 2 }

ptpbaseClockPortDSClockInstanceIndex OBJECT-TYPE
SYNTAX ClockInstanceType
MAX-ACCESS not-accessible
STATUS current
DESCRIPTION
"This object specifies the instance of the clock for this clock
type in the given domain."
 ::= { ptpbaseClockPortDSEntry 3 }

ptpbaseClockPortDSPortNumberIndex OBJECT-TYPE
SYNTAX ClockPortNumber
MAX-ACCESS not-accessible
STATUS current
DESCRIPTION
"This object specifies the PTP portnumber associated with this
PTP port."
 ::= { ptpbaseClockPortDSEntry 4 }

ptpbaseClockPortDSName OBJECT-TYPE
SYNTAX DisplayString (SIZE (1..64))
MAX-ACCESS read-only
STATUS current
DESCRIPTION
"This object specifies the PTP clock port dataset name."
 ::= { ptpbaseClockPortDSEntry 5 }

ptpbaseClockPortDSPortIdentity OBJECT-TYPE
SYNTAX OCTET STRING(SIZE(1..256))
MAX-ACCESS read-only
STATUS current
DESCRIPTION
"This object specifies the PTP clock port Identity."
 ::= { ptpbaseClockPortDSEntry 6 }

ptpbaseClockPortDSlogAnnouncementInterval OBJECT-TYPE
SYNTAX ClockIntervalBase2

UNITS "Time Interval"
MAX-ACCESS read-only
STATUS current
DESCRIPTION
"This object specifies the Announce message transmission
interval associated with this clock port."
 ::= { ptpbaseClockPortDSEntry 7 }

ptpbaseClockPortDSAnnounceRctTimeout OBJECT-TYPE

SYNTAX Integer32
MAX-ACCESS read-only
STATUS current
DESCRIPTION
"This object specifies the Announce receipt timeout associated
with this clock port."
 ::= { ptpbaseClockPortDSEntry 8 }

ptpbaseClockPortDSlogSyncInterval OBJECT-TYPE

SYNTAX ClockIntervalBase2
UNITS "Time Interval"
MAX-ACCESS read-only
STATUS current
DESCRIPTION
"This object specifies the Sync message transmission interval."
 ::= { ptpbaseClockPortDSEntry 9 }

ptpbaseClockPortDSMinDelayReqInterval OBJECT-TYPE

SYNTAX Integer32
MAX-ACCESS read-only
STATUS current
DESCRIPTION
"This object specifies the Delay_Req message transmission
interval."
 ::= { ptpbaseClockPortDSEntry 10 }

ptpbaseClockPortDSPeerDelayReqInterval OBJECT-TYPE

SYNTAX Integer32
MAX-ACCESS read-only
STATUS current
DESCRIPTION
"This object specifies the Pdelay_Req message transmission
interval."
 ::= { ptpbaseClockPortDSEntry 11 }

```
ptpbasedelayMech OBJECT-TYPE
    SYNTAX          ClockMechanismType
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the delay mechanism used. If the clock
        is an end-to-end clock, the value of the is e2e, else if the
        clock is a peer to-peer clock, the value shall be p2p."
    ::= { ptpbaseClockPortDSEntry 12 }

ptpbasedelayPeerMeanPathDelay OBJECT-TYPE
    SYNTAX          ClockTimeInterval
    UNITS           "Time Interval"
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the peer meanPathDelay."
    ::= { ptpbaseClockPortDSEntry 13 }

ptpbasedelayGrantDuration OBJECT-TYPE
    SYNTAX          Unsigned32
    UNITS           "seconds"
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the grant duration allocated by the
        master."
    ::= { ptpbaseClockPortDSEntry 14 }

ptpbasedelayPTPVersion OBJECT-TYPE
    SYNTAX          Unsigned32
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the PTP version being used."
    ::= { ptpbaseClockPortDSEntry 15 }

ptpbasedelayRunningTable OBJECT-TYPE
    SYNTAX          SEQUENCE OF PtpbaseClockPortRunningEntry
    MAX-ACCESS      not-accessible
```

```

STATUS          current
DESCRIPTION
    "Table of information about the clock ports running dataset for
    a particular domain."
 ::= { ptpbaseMIBClockInfo 9 }

```

```

ptpbaseClockPortRunningEntry OBJECT-TYPE
SYNTAX          PtpbaseClockPortRunningEntry
MAX-ACCESS     not-accessible
STATUS         current
DESCRIPTION
    "An entry in the table, containing running dataset information
    about a single clock port."
INDEX          {
                ptpbaseClockPortRunningDomainIndex,
                ptpbaseClockPortRunningClockTypeIndex,
                ptpbaseClockPortRunningClockInstanceIndex,
                ptpbaseClockPortRunningPortNumberIndex
            }
 ::= { ptpbaseClockPortRunningTable 1 }

```

```

PtpbaseClockPortRunningEntry ::= SEQUENCE {
    ptpbaseClockPortRunningDomainIndex      ClockDomainType,
    ptpbaseClockPortRunningClockTypeIndex   ClockType,
    ptpbaseClockPortRunningClockInstanceIndex ClockInstanceType,
    ptpbaseClockPortRunningPortNumberIndex  ClockPortNumber,
    ptpbaseClockPortRunningName             DisplayString,
    ptpbaseClockPortRunningState            ClockPortState,
    ptpbaseClockPortRunningRole             ClockRoleType,
    ptpbaseClockPortRunningInterfaceIndex   InterfaceIndexOrZero,
    ptpbaseClockPortRunningTransport        AutonomousType,
    ptpbaseClockPortRunningEncapsulationType AutonomousType,
    ptpbaseClockPortRunningTxMode          ClockTxModeType,
    ptpbaseClockPortRunningRxMode          ClockTxModeType,
    ptpbaseClockPortRunningPacketsReceived Counter64,
    ptpbaseClockPortRunningPacketsSent     Counter64
}

```

```

ptpbaseClockPortRunningDomainIndex OBJECT-TYPE
SYNTAX          ClockDomainType
MAX-ACCESS     not-accessible
STATUS         current
DESCRIPTION

```

```

    "This object specifies the domain number used to create logical
    group of PTP devices."
 ::= { ptpbaseClockPortRunningEntry 1 }

ptpbaseClockPortRunningClockTypeIndex OBJECT-TYPE
    SYNTAX          ClockType
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the clock type as defined in the
        Textual convention description."
 ::= { ptpbaseClockPortRunningEntry 2 }

ptpbaseClockPortRunningClockInstanceIndex OBJECT-TYPE
    SYNTAX          ClockInstanceType
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the instance of the clock for this clock
        type in the given domain."
 ::= { ptpbaseClockPortRunningEntry 3 }

ptpbaseClockPortRunningPortNumberIndex OBJECT-TYPE
    SYNTAX          ClockPortNumber
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the PTP portnumber associated with this
        clock port."
 ::= { ptpbaseClockPortRunningEntry 4 }

ptpbaseClockPortRunningName OBJECT-TYPE
    SYNTAX          DisplayString (SIZE (1..64))
    MAX-ACCESS      read-only
    STATUS          current
    DESCRIPTION
        "This object specifies the PTP clock port name."
 ::= { ptpbaseClockPortRunningEntry 5 }

ptpbaseClockPortRunningState OBJECT-TYPE
    SYNTAX          ClockPortState
    MAX-ACCESS      read-only
    STATUS          current
```

DESCRIPTION

"This object specifies the port state returned by PTP engine.

- initializing - In this state a port initializes its data sets, hardware, and communication facilities.
- faulty - The fault state of the protocol.
- disabled - The port shall not place any messages on its communication path.
- listening - The port is waiting for the announceReceiptTimeout to expire or to receive an Announce message from a master.
- preMaster - The port shall behave in all respects as though it were in the MASTER state except that it shall not place any messages on its communication path except for Pdelay_Req, Pdelay_Resp, Pdelay_Resp_Follow_Up, signaling, or management messages.
- master - The port is behaving as a master port.
- passive - The port shall not place any messages on its communication path except for Pdelay_Req, Pdelay_Resp, Pdelay_Resp_Follow_Up, or signaling messages, or management messages that are a required response to another management message
- uncalibrated - The local port is preparing to synchronize to the master port.
- slave - The port is synchronizing to the selected master port."

::= { ptpbaseClockPortRunningEntry 6 }

ptpbaseClockPortRunningRole OBJECT-TYPE

SYNTAX ClockRoleType

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies the Clock Role."

::= { ptpbaseClockPortRunningEntry 7 }

ptpbaseClockPortRunningInterfaceIndex OBJECT-TYPE

SYNTAX InterfaceIndexOrZero
MAX-ACCESS read-only
STATUS current
DESCRIPTION
 "This object specifies the interface on the router being used by
 the PTP Clock for PTP communication."
 ::= { ptpbaseClockPortRunningEntry 8 }

ptpbaseClockPortRunningTransport OBJECT-TYPE

SYNTAX AutonomousType
MAX-ACCESS read-only
STATUS current
DESCRIPTION
 "This object specifies the transport protocol being used for PTP
 communication (the mapping used)."
 ::= { ptpbaseClockPortRunningEntry 9 }

ptpbaseClockPortRunningEncapsulationType OBJECT-TYPE

SYNTAX AutonomousType
MAX-ACCESS read-only
STATUS current
DESCRIPTION
 "This object specifies the type of encapsulation if the
 interface is adding extra layers (eg. VLAN, Pseudowire
 encapsulation...) for the PTP messages."
 ::= { ptpbaseClockPortRunningEntry 10 }

ptpbaseClockPortRunningTxMode OBJECT-TYPE

SYNTAX ClockTxModeType
MAX-ACCESS read-only
STATUS current
DESCRIPTION
 "This object specifies the clock transmission mode as

 unicast: Using unicast communication channel.
 multicast: Using Multicast communication channel.
 multicast-mix: Using multicast-unicast communication channel"
 ::= { ptpbaseClockPortRunningEntry 11 }

ptpbaseClockPortRunningRxMode OBJECT-TYPE

SYNTAX ClockTxModeType
MAX-ACCESS read-only
STATUS current

DESCRIPTION

"This object specifies the clock receive mode as

unicast: Using unicast communication channel.
multicast: Using Multicast communication channel.
multicast-mix: Using multicast-unicast communication channel"

::= { ptpbaseClockPortRunningEntry 12 }

ptpbaseClockPortRunningPacketsReceived OBJECT-TYPE

SYNTAX Counter64

UNITS "packets"

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies the packets received on the clock port (cumulative)."

::= { ptpbaseClockPortRunningEntry 13 }

ptpbaseClockPortRunningPacketsSent OBJECT-TYPE

SYNTAX Counter64

UNITS "packets"

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies the packets sent on the clock port (cumulative)."

::= { ptpbaseClockPortRunningEntry 14 }

ptpbaseClockPortTransDSTable OBJECT-TYPE

SYNTAX SEQUENCE OF PtpbaseClockPortTransDSEntry

MAX-ACCESS not-accessible

STATUS current

DESCRIPTION

"Table of information about the Transparent clock ports running dataset for a particular domain."

::= { ptpbaseMIBClockInfo 10 }

ptpbaseClockPortTransDSEntry OBJECT-TYPE

SYNTAX PtpbaseClockPortTransDSEntry

MAX-ACCESS not-accessible

STATUS current

DESCRIPTION

"An entry in the table, containing clock port Transparent dataset information about a single clock port"

INDEX

```

{
    ptpbaseClockPortTransDSDomainIndex,
    ptpbaseClockPortTransDSInstanceIndex,
    ptpbaseClockPortTransDSPortNumberIndex
}
 ::= { ptpbaseClockPortTransDSTable 1 }

```

```

PtpbaseClockPortTransDSEntry ::= SEQUENCE {
    ptpbaseClockPortTransDSDomainIndex      ClockDomainType,
    ptpbaseClockPortTransDSInstanceIndex    ClockInstanceType,
    ptpbaseClockPortTransDSPortNumberIndex  ClockPortNumber,
    ptpbaseClockPortTransDSPortIdentity     ClockIdentity,
    ptpbaseClockPortTransDSlogMinPdelayReqInt ClockIntervalBase2,
    ptpbaseClockPortTransDSFaultyFlag      TruthValue,
    ptpbaseClockPortTransDSPeerMeanPathDelay ClockTimeInterval
}

```

ptpbaseClockPortTransDSDomainIndex OBJECT-TYPE

SYNTAX ClockDomainType

MAX-ACCESS not-accessible

STATUS current

DESCRIPTION

"This object specifies the domain number used to create logical group of PTP devices."

```
 ::= { ptpbaseClockPortTransDSEntry 1 }
```

ptpbaseClockPortTransDSInstanceIndex OBJECT-TYPE

SYNTAX ClockInstanceType

MAX-ACCESS not-accessible

STATUS current

DESCRIPTION

"This object specifies the instance of the clock for this clock type in the given domain."

```
 ::= { ptpbaseClockPortTransDSEntry 2 }
```

ptpbaseClockPortTransDSPortNumberIndex OBJECT-TYPE

SYNTAX ClockPortNumber

MAX-ACCESS not-accessible

STATUS current

DESCRIPTION

"This object specifies the PTP port number associated with this port."

REFERENCE "Section 7.5.2 Port Identity [IEEE 1588-2008]"
 ::= { ptpbaseClockPortTransDSEntry 3 }

ptpbaseClockPortTransDSPortIdentity OBJECT-TYPE

SYNTAX ClockIdentity

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies the value of the PortIdentity attribute of the local port."

REFERENCE "Section 8.3.3.2.1 of [IEEE 1588-2008]"
 ::= { ptpbaseClockPortTransDSEntry 4 }

ptpbaseClockPortTransDSlogMinPdelayReqInt OBJECT-TYPE

SYNTAX ClockIntervalBase2

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies the value of the logarithm to the base 2 of the minPdelayReqInterval."

REFERENCE "Section 8.3.3.3.1 of [IEEE 1588-2008]"
 ::= { ptpbaseClockPortTransDSEntry 5 }

ptpbaseClockPortTransDSFaultyFlag OBJECT-TYPE

SYNTAX TruthValue

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies the value TRUE if the port is faulty and FALSE if the port is operating normally."

REFERENCE "Section 8.3.3.3.2 of [IEEE 1588-2008]"
 ::= { ptpbaseClockPortTransDSEntry 6 }

ptpbaseClockPortTransDSPeerMeanPathDelay OBJECT-TYPE

SYNTAX ClockTimeInterval

UNITS "Time Interval"

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies, (if the delayMechanism used is P2P) the value is the estimate of the current one-way propagation delay,

i.e., <meanPathDelay> on the link attached to this port, computed using the peer delay mechanism. If the value of the delayMechanism used is E2E, then the value will be zero."

REFERENCE "Section 8.3.3.3.3 of [IEEE 1588-2008]"

```
::= { ptpbaseClockPortTransDSEntry 7 }
```

ptpbaseClockPortAssociateTable OBJECT-TYPE

SYNTAX SEQUENCE OF PtpbaseClockPortAssociateEntry

MAX-ACCESS not-accessible

STATUS current

DESCRIPTION

"Table of information about a given port's associated ports.

For a master port: multiple slave ports which have established sessions with the current master port.

For a slave port: the list of masters available for a given slave port.

Session information (packets, errors) to be displayed based on availability and scenario."

```
::= { ptpbaseMIBClockInfo 11 }
```

--

-- Well Known transport types for PTP communication.

--

```
ptpbaseWellKnownTransportTypes OBJECT IDENTIFIER ::= {  
ptpbaseMIBClockInfo 12 }
```

ptpbaseTransportTypeIPversion4 OBJECT-IDENTITY

STATUS current

DESCRIPTION

"IP version 4"

```
::= { ptpbaseWellKnownTransportTypes 1 }
```

ptpbaseTransportTypeIPversion6 OBJECT-IDENTITY

STATUS current

DESCRIPTION

"IP version 6"

```
::= { ptpbaseWellKnownTransportTypes 2 }
```

```
ptpbaseTransportTypeEthernet OBJECT-IDENTITY
  STATUS current
  DESCRIPTION
    "Ethernet"
    ::= { ptpbaseWellKnownTransportTypes 3 }

ptpbaseTransportTypeDeviceNET OBJECT-IDENTITY
  STATUS current
  DESCRIPTION
    "Device NET"
    ::= { ptpbaseWellKnownTransportTypes 4 }

ptpbaseTransportTypeControlNET OBJECT-IDENTITY
  STATUS current
  DESCRIPTION
    "Control NET"
    ::= { ptpbaseWellKnownTransportTypes 5 }

ptpbaseTransportTypeIEC61158 OBJECT-IDENTITY
  STATUS current
  DESCRIPTION
    "IEC61158"
    ::= { ptpbaseWellKnownTransportTypes 6 }

--
-- Well Known encapsulation types for PTP communication.
--
ptpbaseWellKnownEncapsulationTypes OBJECT IDENTIFIER ::= {
  ptpbaseMIBClockInfo 13 }

ptpbaseEncapsulationTypeEthernet OBJECT-IDENTITY
  STATUS current
  DESCRIPTION
    "Ethernet Encapsulation type."
    ::= { ptpbaseWellKnownEncapsulationTypes 1 }

ptpbaseEncapsulationTypeVLAN OBJECT-IDENTITY
  STATUS current
  DESCRIPTION
    "VLAN Encapsulation type."
```

```

 ::= { ptpbaseWellKnownEncapsulationTypes 2 }

ptpbaseEncapsulationTypeUDPIPLSP OBJECT-IDENTITY
  STATUS current
  DESCRIPTION
    "UDP/IP over MPLS Encapsulation type."
    ::= { ptpbaseWellKnownEncapsulationTypes 3 }

ptpbaseEncapsulationTypePWUDPIPLSP OBJECT-IDENTITY
  STATUS current
  DESCRIPTION
    "UDP/IP Pseudowire over MPLS Encapsulation type."
    ::= { ptpbaseWellKnownEncapsulationTypes 4 }

ptpbaseEncapsulationTypePWEthernetLSP OBJECT-IDENTITY
  STATUS current
  DESCRIPTION
    "Ethernet Pseudowire over MPLS Encapsulation type."
    ::= { ptpbaseWellKnownEncapsulationTypes 5 }

ptpbaseClockPortAssociateEntry OBJECT-TYPE
  SYNTAX          PtpbaseClockPortAssociateEntry
  MAX-ACCESS      not-accessible
  STATUS          current
  DESCRIPTION
    "An entry in the table, containing information about a single
    associated port for the given clockport."
  INDEX           {
                  ptpClockPortCurrentDomainIndex,
                  ptpClockPortCurrentClockTypeIndex,
                  ptpClockPortCurrentClockInstanceIndex,
                  ptpClockPortCurrentPortNumberIndex,
                  ptpbaseClockPortAssociatePortIndex
                }
  ::= { ptpbaseClockPortAssociateTable 1 }

PtpbaseClockPortAssociateEntry ::= SEQUENCE {
  ptpClockPortCurrentDomainIndex      ClockDomainType,
  ptpClockPortCurrentClockTypeIndex   ClockType,
  ptpClockPortCurrentClockInstanceIndex ClockInstanceType,
  ptpClockPortCurrentPortNumberIndex  ClockPortNumber,

```

```
        ptpbaseClockPortAssociatePortIndex      Unsigned32,
        ptpbaseClockPortAssociateAddressType    AutonomousType,
        ptpbaseClockPortAssociateAddress
ClockPortTransportTypeAddress,
        ptpbaseClockPortAssociatePacketsSent    Counter64,
        ptpbaseClockPortAssociatePacketsReceived Counter64,
        ptpbaseClockPortAssociateInErrors       Counter64,
        ptpbaseClockPortAssociateOutErrors      Counter64
    }

ptpClockPortCurrentDomainIndex OBJECT-TYPE
    SYNTAX          ClockDomainType
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the given port's domain number."
    ::= { ptpbaseClockPortAssociateEntry 1 }

ptpClockPortCurrentClockTypeIndex OBJECT-TYPE
    SYNTAX          ClockType
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the given port's clock type."
    ::= { ptpbaseClockPortAssociateEntry 2 }

ptpClockPortCurrentClockInstanceIndex OBJECT-TYPE
    SYNTAX          ClockInstanceType
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the instance of the clock for this clock
        type in the given domain."
    ::= { ptpbaseClockPortAssociateEntry 3 }

ptpClockPortCurrentPortNumberIndex OBJECT-TYPE
    SYNTAX          ClockPortNumber
    MAX-ACCESS      not-accessible
    STATUS          current
    DESCRIPTION
        "This object specifies the PTP Port Number for the given port."
    ::= { ptpbaseClockPortAssociateEntry 4 }
```

ptpbaseClockPortAssociatePortIndex OBJECT-TYPE

SYNTAX Unsigned32 (1..65535)

MAX-ACCESS not-accessible

STATUS current

DESCRIPTION

"This object specifies the associated port's serial number in the current port's context."

::= { ptpbaseClockPortAssociateEntry 5 }

ptpbaseClockPortAssociateAddressType OBJECT-TYPE

SYNTAX AutonomousType

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies the peer port's network address type used for PTP communication."

::= { ptpbaseClockPortAssociateEntry 6 }

ptpbaseClockPortAssociateAddress OBJECT-TYPE

SYNTAX ClockPortTransportTypeAddress

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies the peer port's network address used for PTP communication."

::= { ptpbaseClockPortAssociateEntry 7 }

ptpbaseClockPortAssociatePacketsSent OBJECT-TYPE

SYNTAX Counter64

UNITS "packets"

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"The number of packets sent to this peer port from the current port."

::= { ptpbaseClockPortAssociateEntry 8 }

ptpbaseClockPortAssociatePacketsReceived OBJECT-TYPE

SYNTAX Counter64

UNITS "packets"

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"The number of packets received from this peer port by the current port."
 ::= { ptpbaseClockPortAssociateEntry 9 }

ptpbaseClockPortAssociateInErrors OBJECT-TYPE

SYNTAX Counter64

UNITS "packets"

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies the input errors associated with the peer port."

::= { ptpbaseClockPortAssociateEntry 10 }

ptpbaseClockPortAssociateOutErrors OBJECT-TYPE

SYNTAX Counter64

UNITS "packets"

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"This object specifies the output errors associated with the peer port."

::= { ptpbaseClockPortAssociateEntry 11 }

-- Conformance Information Definition

ptpbaseMIBCompliances OBJECT IDENTIFIER

::= { ptpbaseMIBConformance 1 }

ptpbaseMIBGroups OBJECT IDENTIFIER

::= { ptpbaseMIBConformance 2 }

ptpbaseMIBCompliancesSystemInfo MODULE-COMPLIANCE

STATUS current

DESCRIPTION

"Compliance statement for agents that provide read-only support for PTPBASE-MIB to provide system level information of clock devices.

Such devices can only be monitored using this MIB module.

The Module is implemented with support for read-only. In other

```
words, only monitoring is available by implementing this
MODULE-COMPLIANCE."
MODULE          -- this module
MANDATORY-GROUPS { ntpbaseMIBSystemInfoGroup }
 ::= { ntpbaseMIBCompliances 1 }

ntpbaseMIBCompliancesClockInfo MODULE-COMPLIANCE
STATUS          current
DESCRIPTION
  "Compliance statement for agents that provide read-only support
  for NTPBASE-MIB to provide clock related information.
  Such devices can only be monitored using this MIB module.

  The Module is implemented with support for read-only. In other
  words, only monitoring is available by implementing this
  MODULE-COMPLIANCE."
MODULE          -- this module
MANDATORY-GROUPS {
    ntpbaseMIBClockCurrentDSGroup,
    ntpbaseMIBClockParentDSGroup,
    ntpbaseMIBClockDefaultDSGroup,
    ntpbaseMIBClockRunningGroup,
    ntpbaseMIBClockTimepropertiesGroup
  }
 ::= { ntpbaseMIBCompliances 2 }

ntpbaseMIBCompliancesClockPortInfo MODULE-COMPLIANCE
STATUS          current
DESCRIPTION
  "Compliance statement for agents that provide read-only support
  for NTPBASE-MIB to provide clock port related information.
  Such devices can only be monitored using this MIB module.

  The Module is implemented with support for read-only. In other
  words, only monitoring is available by implementing this
  MODULE-COMPLIANCE."
MODULE          -- this module
MANDATORY-GROUPS {
    ntpbaseMIBClockPortGroup,
    ntpbaseMIBClockPortDSGroup,
    ntpbaseMIBClockPortRunningGroup,
    ntpbaseMIBClockPortAssociateGroup
  }
```

```
 ::= { ptpbaseMIBCompliances 3 }

ptpbaseMIBCompliancesTransparentClockInfo MODULE-COMPLIANCE
  STATUS          current
  DESCRIPTION
    "Compliance statement for agents that provide read-only support
    for PTPBASE-MIB to provide Transparent clock related
    information.
    Such devices can only be monitored using this MIB module.

    The Module is implemented with support for read-only. In other
    words, only monitoring is available by implementing this
    MODULE-COMPLIANCE."
  MODULE          -- this module
  MANDATORY-GROUPS {
    ptpbaseMIBClockTransparentDSGroup,
    ptpbaseMIBClockPortTransDSGroup
  }
  ::= { ptpbaseMIBCompliances 4 }

ptpbaseMIBSystemInfoGroup OBJECT-GROUP
  OBJECTS          {
    ptpbaseSystemDomainTotals,
    ptpDomainClockPortsTotal,
    ptpbaseSystemProfile
  }
  STATUS          current
  DESCRIPTION
    "Group which aggregates objects describing system-wide
    information"
  ::= { ptpbaseMIBGroups 1 }

ptpbaseMIBClockCurrentDSGroup OBJECT-GROUP
  OBJECTS          {
    ptpbaseClockCurrentDSStepsRemoved,
    ptpbaseClockCurrentDSOffsetFromMaster,
    ptpbaseClockCurrentDSMeanPathDelay
  }
  STATUS          current
  DESCRIPTION
    "Group which aggregates objects describing PTP Current Dataset
    information"
  ::= { ptpbaseMIBGroups 2 }
```

```
ptpbasesMIBClockParentDSGroup OBJECT-GROUP
  OBJECTS          {
                    ptpbaseClockParentDSParentPortIdentity,
                    ptpbaseClockParentDSParentStats,
                    ptpbaseClockParentDSOffset,
                    ptpbaseClockParentDSClockPhChRate,
                    ptpbaseClockParentDSGMClockIdentity,
                    ptpbaseClockParentDSGMClockPriority1,
                    ptpbaseClockParentDSGMClockPriority2,
                    ptpbaseClockParentDSGMClockQualityClass,
                    ptpbaseClockParentDSGMClockQualityAccuracy,
                    ptpbaseClockParentDSGMClockQualityOffset
                    }
  STATUS           current
  DESCRIPTION     "Group which aggregates objects describing PTP Parent Dataset
                  information"
  ::= { ptpbaseMIBGroups 3 }

ptpbasesMIBClockDefaultDSGroup OBJECT-GROUP
  OBJECTS          {
                    ptpbaseClockDefaultDSTwoStepFlag,
                    ptpbaseClockDefaultDSClockIdentity,
                    ptpbaseClockDefaultDSPriority1,
                    ptpbaseClockDefaultDSPriority2,
                    ptpbaseClockDefaultDSSlaveOnly,
                    ptpbaseClockDefaultDSQualityClass,
                    ptpbaseClockDefaultDSQualityAccuracy,
                    ptpbaseClockDefaultDSQualityOffset
                    }
  STATUS           current
  DESCRIPTION     "Group which aggregates objects describing PTP Default Dataset
                  information"
  ::= { ptpbaseMIBGroups 4 }

ptpbasesMIBClockRunningGroup OBJECT-GROUP
  OBJECTS          {
                    ptpbaseClockRunningState,
                    ptpbaseClockRunningPacketsSent,
                    ptpbaseClockRunningPacketsReceived
                    }
}
```

```
STATUS          current
DESCRIPTION
  "Group which aggregates objects describing PTP running state
  information"
 ::= { ptpbaseMIBGroups 5 }

ptpbaseMIBClockTimepropertiesGroup OBJECT-GROUP
OBJECTS {
  ptpbaseClockTimePropertiesDSCurrentUTCOffsetValid,
  ptpbaseClockTimePropertiesDSCurrentUTCOffset,
  ptpbaseClockTimePropertiesDSLeap59,
  ptpbaseClockTimePropertiesDSLeap61,
  ptpbaseClockTimePropertiesDSTimeTraceable,
  ptpbaseClockTimePropertiesDSFreqTraceable,
  ptpbaseClockTimePropertiesDSPTPTimescale,
  ptpbaseClockTimePropertiesDSSource
}
STATUS          current
DESCRIPTION
  "Group which aggregates objects describing PTP Time Properties
  information"
 ::= { ptpbaseMIBGroups 6 }

ptpbaseMIBClockTranparentDSGroup OBJECT-GROUP
OBJECTS {
  ptpbaseClockTransDefaultDSClockIdentity,
  ptpbaseClockTransDefaultDSNumOfPorts,
  ptpbaseClockTransDefaultDSDelay,
  ptpbaseClockTransDefaultDSPrimaryDomain
}
STATUS          current
DESCRIPTION
  "Group which aggregates objects describing PTP Transparent
  Dataset
  information"
 ::= { ptpbaseMIBGroups 7 }

ptpbaseMIBClockPortGroup OBJECT-GROUP
OBJECTS {
  ptpbaseClockPortName,
  ptpbaseClockPortSyncTwoStep,
  ptpbaseClockPortCurrentPeerAddress,
  ptpbaseClockPortNumOfAssociatedPorts,
```

```

        ptpbaseClockPortCurrentPeerAddressType,
        ptpbaseClockPortRole
    }
STATUS          current
DESCRIPTION
    "Group which aggregates objects describing information for a
    given PTP Port."
 ::= { ptpbaseMIBGroups 8 }

ptpbaseMIBClockPortDSGroup OBJECT-GROUP
OBJECTS        {
    ptpbaseClockPortDSName,
    ptpbaseClockPortDSPortIdentity,
    ptpbaseClockPortDSlogAnnouncementInterval,
    ptpbaseClockPortDSAnnounceRctTimeout,
    ptpbaseClockPortDSlogSyncInterval,
    ptpbaseClockPortDSMinDelayReqInterval,
    ptpbaseClockPortDSPeerDelayReqInterval,
    ptpbaseClockPortDSDelayMech,
    ptpbaseClockPortDSPeerMeanPathDelay,
    ptpbaseClockPortDSGrantDuration,
    ptpbaseClockPortDSPTPVersion
}
STATUS          current
DESCRIPTION
    "Group which aggregates objects describing PTP Port Dataset
    information"
 ::= { ptpbaseMIBGroups 9 }

ptpbaseMIBClockPortRunningGroup OBJECT-GROUP
OBJECTS        {
    ptpbaseClockPortRunningName,
    ptpbaseClockPortRunningState,
    ptpbaseClockPortRunningRole,
    ptpbaseClockPortRunningInterfaceIndex,
    ptpbaseClockPortRunningTransport,
    ptpbaseClockPortRunningEncapsulationType,
    ptpbaseClockPortRunningTxMode,
    ptpbaseClockPortRunningRxMode,
    ptpbaseClockPortRunningPacketsReceived,
    ptpbaseClockPortRunningPacketsSent
}
STATUS          current
```

DESCRIPTION

"Group which aggregates objects describing PTP running interface information"

::= { ptpbaseMIBGroups 10 }

ptpbaseMIBClockPortTransDSGroup OBJECT-GROUP

```
OBJECTS      {
                ptpbaseClockPortTransDSPortIdentity,
                ptpbaseClockPortTransDSlogMinPdelayReqInt,
                ptpbaseClockPortTransDSFaultyFlag,
                ptpbaseClockPortTransDSPeerMeanPathDelay
            }
STATUS      current
```

DESCRIPTION

"Group which aggregates objects describing PTP TransparentDS Dataset information"

::= { ptpbaseMIBGroups 11 }

ptpbaseMIBClockPortAssociateGroup OBJECT-GROUP

```
OBJECTS      {
                ptpbaseClockPortAssociatePacketsSent,
                ptpbaseClockPortAssociatePacketsReceived,
                ptpbaseClockPortAssociateAddress,
                ptpbaseClockPortAssociateAddressType,
                ptpbaseClockPortAssociateInErrors,
                ptpbaseClockPortAssociateOutErrors
            }
STATUS      current
```

DESCRIPTION

"Group which aggregates objects describing information on peer PTP ports for a given PTP clock-port."

::= { ptpbaseMIBGroups 12 }

END

5. Security Considerations

This MIB contains readable objects whose values provide information related to PTP objects. It does not contain writable objects.

Some of the readable objects in this MIB module (i.e., objects with a MAX-ACCESS other than not-accessible) may be considered sensitive or vulnerable in some network environments. It is thus important to control even GET and/or NOTIFY access to these objects and possibly to even encrypt the values of these objects when sending them over the network via SNMP.

SNMP versions prior to SNMPv3 did not include adequate security. Even if the network itself is secure (for example by using IPsec), even then, there is no control as to who on the secure network is allowed to access and GET (read) the objects in this MIB module.

It is recommended that the implementers consider the security features as provided by the SNMPv3 framework (see [RFC 3410], section 8). Specifically, the use of the User-based Security Model [RFC 3414] and the View-based Access Control Model [RFC 3415] is recommended.

Further, deployment of SNMP versions prior to SNMPv3 is NOT recommended. Instead, it is recommended to deploy SNMPv3 and to enable cryptographic security. It is then a customer/operator responsibility to ensure that the SNMP entity giving access to an instance of this MIB is properly configured to give access to those objects only to those principals (users) that have legitimate rights to access them.

6. IANA Considerations

The MIB module defined in this document uses the following IANA-assigned OBJECT IDENTIFIER value recorded in the SMI Numbers registry:

Descriptor	OBJECT IDENTIFIER value
-----	-----
ptpbasesMIB	{ mib-2 xxx }

[NOTE for IANA: Please allocate an object identifier at <http://www.iana.org/assignments/smi-numbers> for object ptpbasesMIB.]

7. References

7.1. Normative References

[IEEE 1588-2008] "IEEE Standard for A Precision Clock Synchronization Protocol for Networked Measurement and Control Systems", IEEE Std. 1588(TM)-2008, 24 July 2008

7.2. Informative References

[RFC 1155] Rose, M., and K. McCloghrie, "Structure and Identification of Management Information for TCP/IP-based Internets", STD 16, RFC 1155, Performance Systems International, Hughes LAN Systems, May 1990

[RFC 1157] Case, J., Fedor, M., Schoffstall, M., and J. Davin, "Simple Network Management Protocol", STD 15, RFC 1157, SNMP Research, Performance Systems International, Performance Systems International, MIT Laboratory for Computer Science, May 1990.

[RFC 1212] Rose, M., and K. McCloghrie, "Concise MIB Definitions", STD 16, RFC 1212, Performance Systems International, Hughes LAN Systems, March 1991

[RFC 1215] M. Rose, "A Convention for Defining Traps for use with the SNMP", RFC 1215, Performance Systems International, March 1991

[RFC 1901] SNMPv2 Working Group, Case, J., McCloghrie, K., Rose, M., and S. Waldbusser, "Introduction to Community-based SNMPv2", RFC 1901, SNMP Research, Inc., Cisco Systems, Inc., Dover Beach Consulting, Inc., International Network Services, January 1996.

[RFC 1906] SNMPv2 Working Group, Case, J., McCloghrie, K., Rose, M., and S. Waldbusser, "Transport Mappings for Version 2 of the Simple Network Management Protocol (SNMPv2)", RFC 1906, SNMP Research, Inc., Cisco Systems, Inc., Dover Beach Consulting, Inc., International Network Services, January 1996.

[RFC 2578] McCloghrie, K., Perkins, D., and J. Schoenwaelder, "Structure of Management Information Version 2 (SMIV2)", STD 58, RFC 2578, April 1999.

[RFC 2579] McCloghrie, K., Perkins, D., and J. Schoenwaelder, "Textual Conventions for SMIV2", STD 58, RFC 2579, April 1999.

[RFC 2580] McCloghrie, K., Perkins, D., and J. Schoenwaelder, "Conformance Statements for SMIV2", STD 58, RFC 2580, April 1999.

[RFC 3410] Case, J., Mundy, R., Partain, D., and B. Stewart, "Introduction and Applicability Statements for Internet Standard

Management Framework", RFC 3410 SNMP Research, Inc., Network Associates Laboratories, Ericsson, December 2002.

[RFC 3411] Harrington, D., Presuhn, R., and B. Wijnen, "An Architecture for Describing Simple Network Management Protocol (SNMP) Management Frameworks", STD 62, RFC 3411, Enterasys Networks, BMC Software, Inc., Lucent Technologies, December 2002

[RFC 3412] Case, J., Harrington D., Presuhn R., and B. Wijnen, "Message Processing and Dispatching for the Simple Network Management Protocol (SNMP)", STD 62, RFC 3412, SNMP Research, Inc., Enterasys Networks, BMC Software, Inc., Lucent Technologies, December 2002.

[RFC 3413] Levi, D., Meyer, P., and B. Stewart, "Simple Network Management Protocol (SNMP) Applications", STD 62, RFC 3413, Nortel Networks, Secure Computing Corporation, December 2002.

[RFC 3414] Blumenthal, U., and B. Wijnen, "User-based Security Model (USM) for version 3 of the Simple Network Management Protocol (SNMPv3)", STD 62, RFC 3414, Lucent Technologies, December 2002.

[RFC 3415] Wijnen, B., Presuhn, R., and K. McCloghrie, "View-based Access Control Model (VACM) for the Simple Network Management Protocol (SNMP)", STD 62, RFC 3415, Lucent Technologies, BMC Software, Inc., Cisco Systems, Inc., December 2002.

[RFC 3416] Presuhn, R. (Ed.), "Version 2 of the Protocol Operations for the Simple Network Management Protocol (SNMP)", STD 62, RFC 3416, BMC Software, Inc., December 2002.

[RFC 3417] Presuhn, R. (Ed.), "Transport Mappings for the Simple Network Management Protocol (SNMP)", STD 62, RFC 3417, BMC Software, Inc., December 2002.

[RFC 5905] David L. Mills, " Network Time Protocol Version 4: Protocol and Algorithms Specification", RFC 5905, University of Delaware, June 2010.

[IEEE 802.3-2008] "IEEE Standard for Information technology - Telecommunications and information exchange between systems - Local and Metropolitan area networks - Specific requirements Part 3: Carrier sense multiple access with Collision Detection (CSMA/CD) Access Method and Physical Layer Specifications", IEEE Std. 802.3 - 2008, 26 December 2008

[G.8265.1] "Precision time protocol telecom profile for frequency synchronization", ITU-T Recommendation G.8265.1, October 2010.

8. Acknowledgements

Thanks to John Linton and Danny Lee for valuable comments, and to Bert Wijnen, Kevin Gross, Alan Luchuk and Chris Elliot for their reviews of this MIB.

9. Author's Addresses

Vinay Shankarkumar
Cisco Systems,
7025-4 Kit Creek Road,
Research Triangle Park,
NC 27560,
USA.

Email: vinays@cisco.com

Laurent Montini,
Cisco Systems,
11, rue Camille Desmoulins,
92782 Issy-les-Moulineaux,
France.

Email: lmontini@cisco.com

Tim Frost,
Calnex Solutions Ltd.,
Oracle Campus,
Linlithgow,
EH49 7LR,
UK.

Email: tim.frost@calnexsol.com

Greg Dowd,
Microsemi Inc.,
3870 N. First Str.,
San Jose,
CA 95134,
USA.

Email: greg.dowd@microsemi.com

