Diet-ESP: A Flexible and Wide Range Security protocol

draft-mglt-6lo-diet-esp-requirements draft-mglt-6lo-diet-esp

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Motivations

The current de-facto IoT security protocol is DTLS1.2 (DICE profile).

- Reasonable choice for:
  - Web based IoT applications,
  - End-to-end security
- but in our view DTLS1.2/DICE does not address all IoT segments
- There is a need for ESP in IoT
- Diet-ESP is the ESP implementation for IoT
ESP packet Description

<table>
<thead>
<tr>
<th>Security Parameters Index (SPI)</th>
<th>^Int.</th>
</tr>
</thead>
<tbody>
<tr>
<td>Sequence Number</td>
<td>^Cov-</td>
</tr>
<tr>
<td>Payload Data* (variable)</td>
<td>^Cov-*</td>
</tr>
<tr>
<td>Payload (0-255 bytes)</td>
<td>^Cov-*</td>
</tr>
<tr>
<td>Pad Length</td>
<td>^Cov-*</td>
</tr>
<tr>
<td>Integrity Check Value-ICV (variable)</td>
<td>^Cov-*</td>
</tr>
</tbody>
</table>
Diet-ESP

Diet-ESP aims at compressing:

- ESP Header
- ESP Trailer

Diet-ESP is not something new:

- ROHC/6LowPAN provides ways to compress the ESP Header
- ROHCoverIPsec provides ways to compress the clear text data
Diet-ESP

What is new with Diet-ESP is that:

- It enables to compress Padding, Pad length and Next Header fields
- Compression occurs before the encryption (similarly to ROHCoverIPsec)
- Compression occurs after the ESP fields are added (unlike ROHCoverIPsec)
  - There is no decompression of these fields. (one way compression)
- It takes advantage of IKEv2 to agree on the compression rules
  - Results in a light-compression framework
Diet-ESP/ROHC

Currently Diet-ESP is based on ROHC:

- ROHC compressor/decompressor
  - With initialized states
  - Without synchronization, initialization exchanges

- ROHC profiles

Open Discussion are:

- Does Diet-ESP defines ROHC compressor or a specific compressor?
  - Other alternatives exists like SCHC
  - Our profiles will be called rules

- Can we consider the profiles as an extension of ROHCoverIPsec?
  - This would make possible ROHC/Diet-ESP and Diet-ESP
Diet-ESP/ESP

Diet-ESP is based on ESP:

- Diet-ESP is based on standard ESP
- Diet-ESP is able to send ESP packet
  - Enable natural fall back to uncompressed ESP
  - Preserve ESP interoperability
- Diet-ESP DOES NOT modify cryptographic parameters, algorithms
  - Crypto is left untouched
Diet-ESP / AES-CCM

How Diet-ESP works with AES-CCM

- 1: Decompress ESP header.
- 2: Generate Diet-ESP ICV and check ICV send in the packet.
- 3: Check anti-replay
- 4: Remove compressed header.
- 5: Decrypt the Diet-ESP payload.

[...]
Questions

- Does ESP requires 32/64 bit alignment?
- How Padding is generated?
- Does an AES-CCM ESP packet has Padding?
- How the ICV is built?
Bit Alignment

Does ESP requires 32/64 bit alignment?

- RFC4303 section 2.4
  
  Padding also may be required, irrespective of encryption algorithm requirements, to ensure that the resulting ciphertext terminates on a 4-byte boundary. Specifically, the Pad Length and Next Header fields must be right aligned within a 4-byte word, as illustrated in the ESP packet format figures above, to ensure that the ICV field (if present) is aligned on a 4-byte boundary.

- draft-mglt-diet-esp-requirements

  IP extension headers MUST have 32 bit Byte-Alignment in IPv4 (section 3.1 of [RFC0791] - Padding description) and a 64 bit Byte-Alignment in IPv6 (section 4 of [RFC2460]). As ESP [RFC4303] is such an extension header, padding is mandatory to meet the alignment constraint.
Padding Generation

How Padding is generated?

- RFC4303 section 2.4

If Padding bytes are needed but the encryption algorithm does not specify the padding contents, then the following default processing MUST be used. The Padding bytes are initialized with a series of (unsigned, 1-byte) integer values. The first padding byte appended to the plaintext is numbered 1, with subsequent padding bytes making up a monotonically increasing sequence: 1, 2, 3, .... When this padding scheme is employed, the receiver SHOULD inspect the Padding field. (This scheme was selected because of its relative simplicity, ease of implementation in hardware, and because it offers limited protection against certain forms of "cut and paste" attacks in the absence of other integrity measures, if the receiver checks the padding values upon decryption.)
Padding in AES-CCM

Does an AES-CCM ESP packet has Padding?

- RFC4309

Padding:
The encrypted payload contains the ciphertext.

AES CCM mode does not require plaintext padding. However, ESP does require padding to 32-bit word-align the authentication data. The Padding, Pad Length, and Next Header fields MUST be concatenated with the plaintext before performing encryption, as described in [ESP]. When padding is required, it MUST be generated and checked in accordance with the conventions specified in [ESP].
Padding in AES-CCM

How ICV is built?

- RFC4309 section 3.3 Authentication Data

AES CCM provides an encrypted ICV. The ICV provided by CCM is carried in the Authentication Data fields without further encryption. Implementations MUST support ICV sizes of 8 octets and 16 octets. Implementations MAY also support ICV 12 octets.
Next

- ESP Payload Compression?
  * Implicit IV (presentation of Yoav)
  * Tunnel Header compression (similar to BEET-MODE/draft-mgl-6lo-diet-esp-payload-compression)
  * Transport Header Compression
Thank you for your attention