

# Proof of Transit & In-Band OAM

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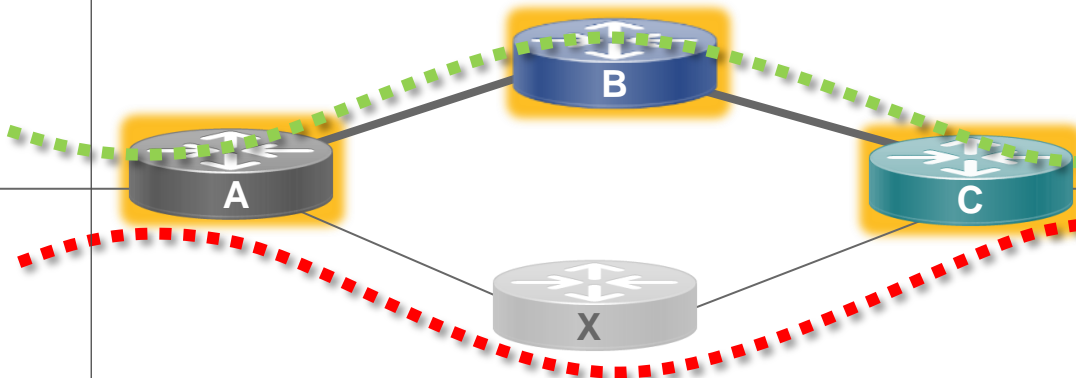
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[draft-brockners-proof-of-transit-01.txt](#)  
[draft-brockners-inband-oam-requirements-01.txt](#)  
[draft-brockners-inband-oam-data-01.txt](#)  
[draft-brockners-inband-oam-transport-01.txt](#)

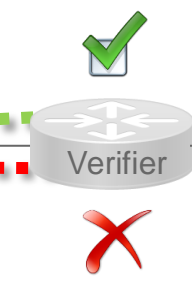
Consider a segment route / service chain:

“How do you *prove* that traffic follows the suggested path?”

## Trust Domain A

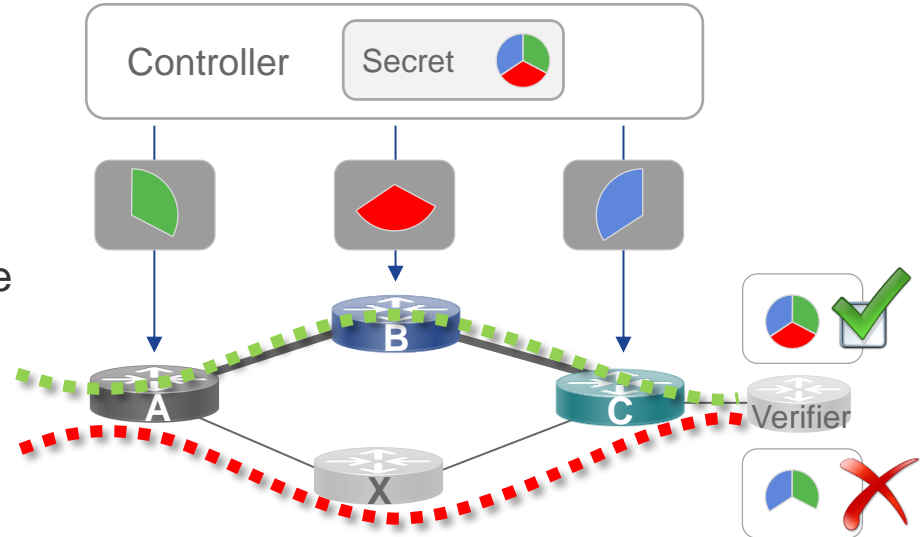


## Trust Domain B

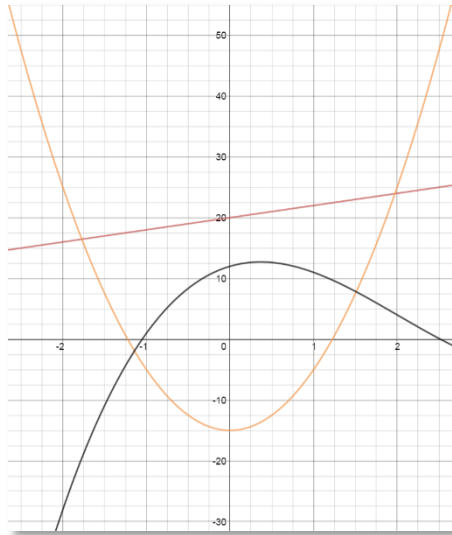


# Ensuring Path and/or Service Chain Integrity Approach

- Meta-data added to all user traffic
  - Based on “Share of a secret”
  - Provisioned by controller over secure channel to segment hops where “proof of transit” is required
  - Updated at every segment hop where proof of transit is required
- Verifier checks whether collected meta-data allows retrieval of secret
  - ➡ “Proof of Transit”: Path verified



# Solution Approach: Leveraging Shamir's Secret Sharing Polynomials 101



$$f2(x) = 10x^2 - 15$$

- Parabola: Min 3 points

$$f1(x) = 2x + 20$$

- Line: Min 2 points

$$f3(x) = x^3 - 6x^2 + 4x - 12$$

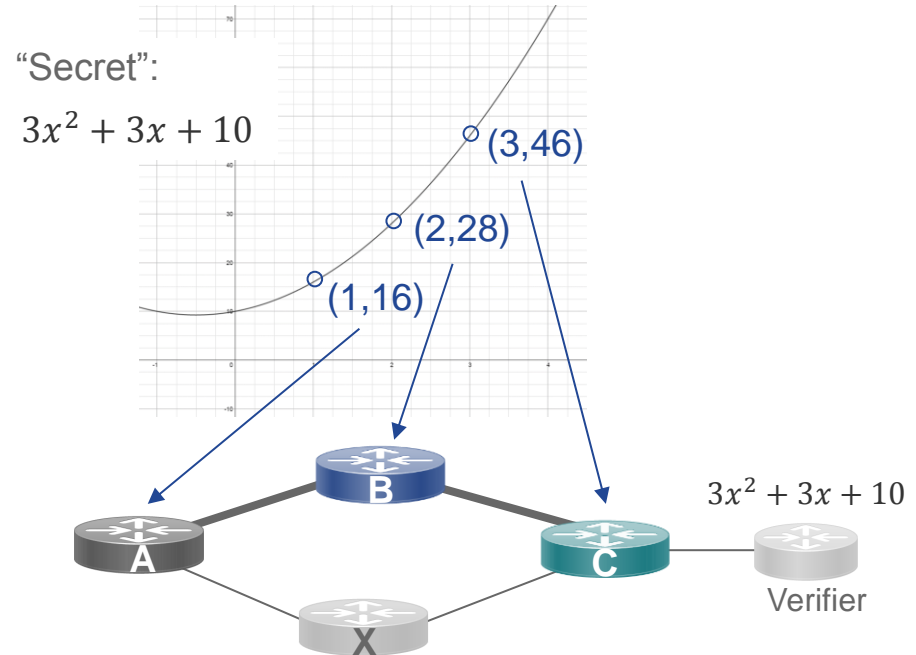
- Cubic function: Min 4 points

General: It takes  $k+1$  points to defines a polynomial of degree  $k$ .

# Solution Approach: Leverage Shamir's Secret Sharing

## “A polynomial as secret”

- Each service is given a point on the curve
- When the packet travels through each service it collects these points
- A verifier can reconstruct the curve using the collected points
- Operations done over a finite field (mod prime) to protect against differential analysis



# Operationalizing the Solution

- Leverage two polynomials:
  - POLY-1 secret, constant: Each hop gets a point on POLY-1  
Only the verifier knows POLY-1
  - POLY-2 public, random and per packet.  
Each hop generates a point on POLY-2 each time a packet crosses it.
- Each service function calculates (Point on POLY-1 + Point on POLY-2) to get (Point on POLY-3) and passes it to verifier by adding it to each packet.
- The verifier constructs POLY-3 from the points given by all the services and cross checks whether  $POLY-3 = POLY-1 + POLY-2$
- Computationally efficient: 3 additions, 1 multiplication, mod prime per hop

POLY-1  
Secret – Constant

+

POLY-2  
Public – Per Packet

=

POLY-3  
Secret – Per Packet

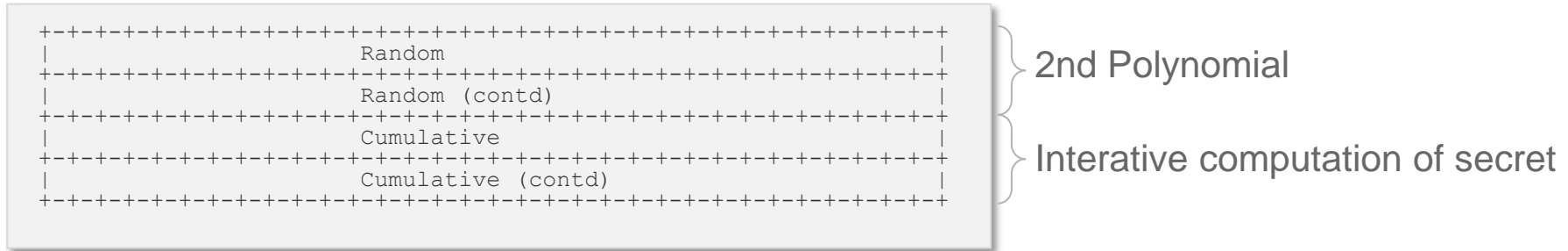
# Meta Data for Service/Path Verification

- Verification secret is the independent coefficient of POLY-1
  - Computation/retrieval through a cumulative computation at every hop (“cumulative”)
- For POLY-2 the independent coefficient is carried within the packet (typically a combination of timestamp and random number)
  - n bits can service a maximum of  $2^n$  packets
- Verification secret and POLY-2 coefficient (“random”) are of the same size
  - Secret size is bound by prime number

Transfer Rate	RND/ Secret Size	Max # of packets (assuming 64 byte packets)	Time that “random” lasts at maximum
1 Gbps	<b>64</b>	$2^{64} \approx 2 * 10^{19}$	approx. $10^{13}$ seconds, approx. 310000 years
10 Gbps	<b>64</b>	$2^{64} \approx 2 * 10^{19}$	approx. $10^{12}$ seconds, approx. 31000 years
100 Gbps	<b>64</b>	$2^{64} \approx 2 * 10^{19}$	approx. $10^{11}$ seconds, approx. 3100 years
10 Gbps	56	$2^{56} \approx 7 * 10^{16}$	approx. $10^9$ seconds, approx. 120 years
10 Gbps	48	$2^{48} \approx 2 * 10^{14}$	approx. $10^7$ seconds, approx. 5.5 months
10 Gbps	40	$2^{40} \approx 1 * 10^{12}$	approx. $5 * 10^4$ seconds, approx. 15 hours
1 Gbps	32	$2^{32} \approx 4 * 10^9$	2200 seconds, 36 minutes
10 Gbps	32	$2^{32} \approx 4 * 10^9$	220 seconds, 3.5 minutes
100 GBps	32	$2^{32} \approx 4 * 10^9$	22 seconds



# Proof of Transit: Meta-Data Transport Options

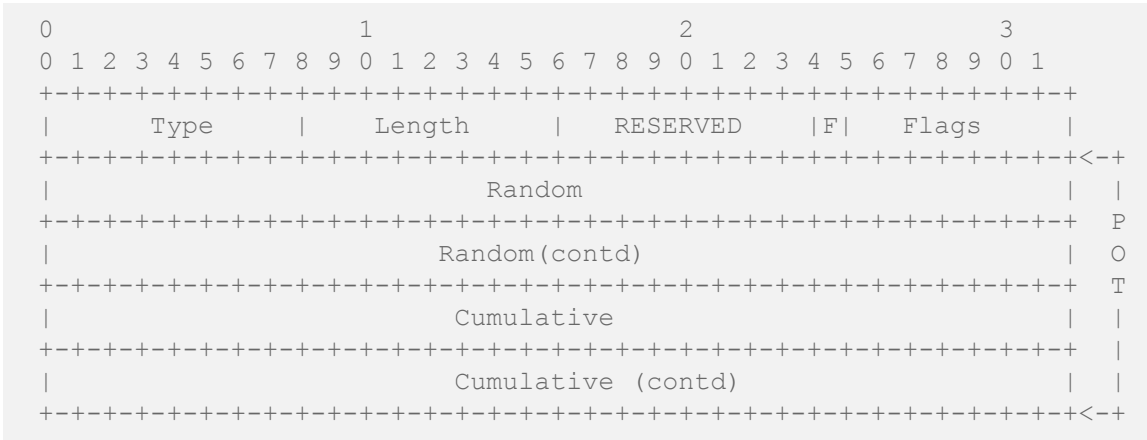


- 16\* Bytes of Meta-Data for SCV
  - Random – Unique random number (e.g. Timestamp or combination of Timestamp and Sequence number)
  - Cumulative (algorithm dependent)

- Transport options for different protocols
  - Segment Routing: New TLV in SRH header
  - Network Service Header: Type-2 Meta-Data
  - In-band OAM for IPv6: Proof-of-transit extension header
  - VXLAN-GPE Proof-of-transit embedded-telemetry header
  - ... more to be added (incl. IPv4)

\*Note: Smaller numbers are feasible, but require a more frequent renewal of the polynomials/secrets.

# SRH: Proposed new POT TLV



**Type:** To be assigned by IANA.

**Length:** 18.

**RESERVED:** 8 bits. SHOULD be unset on transmission and MUST be ignored on receipt.

**F:** 1 bit. Indicates which POT-profile is active. 0 means the even POT-profile is active, 1 means the odd POT-profile is active.

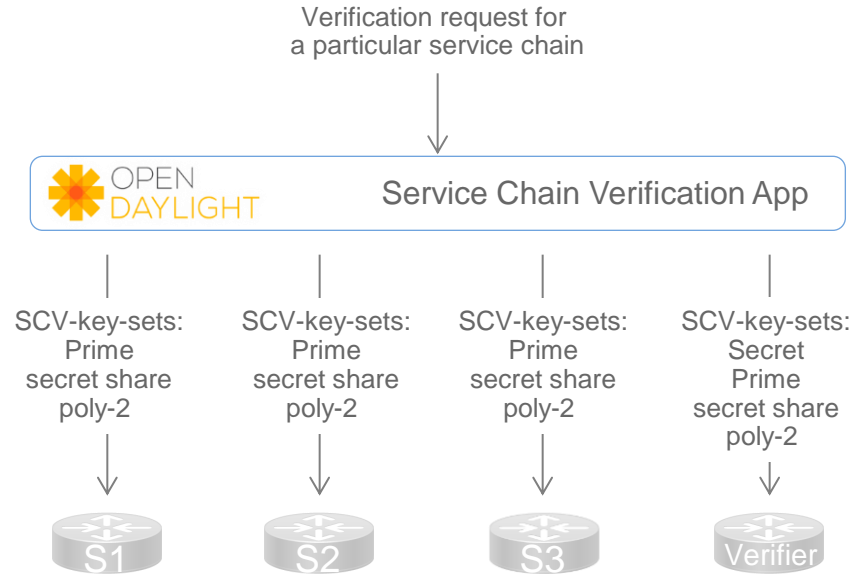
**Flags:** 7 bits. No flags are defined in this document.

**Random:** 64-bit per packet random number.

**Cumulative:** 64-bit cumulative value that is updated at specific nodes that form the service path to be verified.

# Meta-Data Provisioning

- Meta-Data for Segment list provisioned through a controller (e.g. OpenDaylight App)
- Netconf/YANG based protocol
- Provisioned information from Controller to Service Function / Verifier
  - Path/Segment/Service-Chain Identifier
  - Service count (number of services in the chain)
  - 2 x SCV-key-set
    - Secret (in case of communication to the verifier)
    - Share of a secret, service index
    - 2nd polynomial coefficients
    - Prime number



Enter...

# In-Band OAM

# What if you could collect operational meta-data within your traffic?

## Example use-cases...

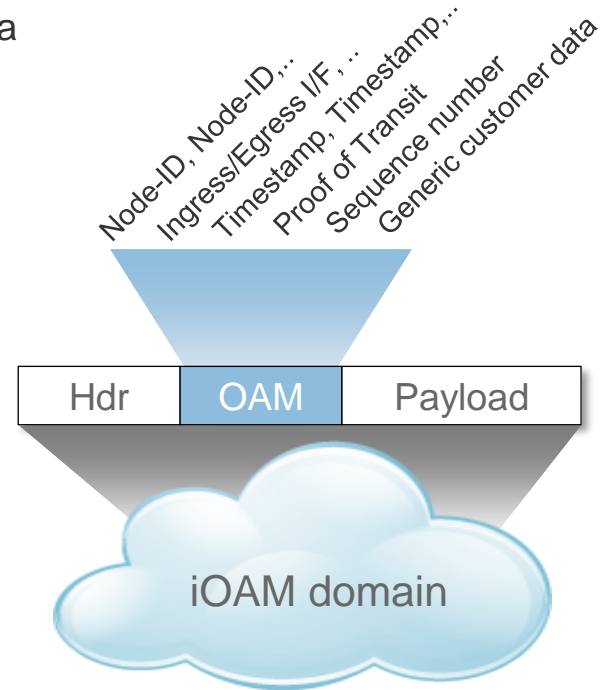
- Path Tracing for ECMP networks
- Service/Path Verification
- Derive Traffic Matrix
- SLA proof: Delay, Jitter, Loss
- Custom data: Geo-Location,..

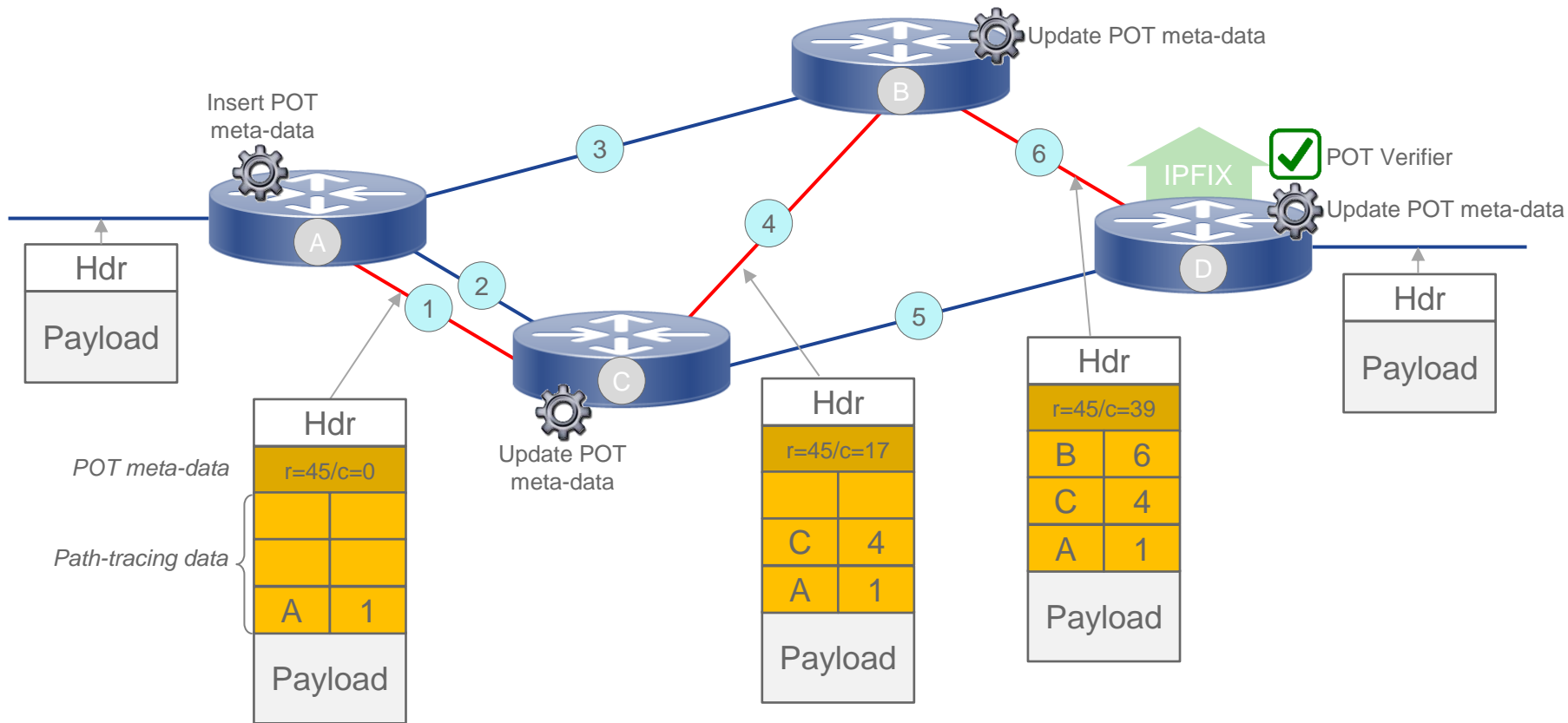
## Meta-data required...

- Node-ID, ingress i/f, egress i/f
- Proof of Transit (random, cumulative)
- Node-ID
- Sequence numbers, Timestamps
- Custom meta-data

# In-Band OAM

- Gather telemetry and OAM information along the path **within** the data packet, as part of an existing/additional header
  - **No** extra probe-traffic (as with ping, trace, ipsla)
- Transport options
  - IPv6: Native v6 HbyH extension header or double-encap
  - VXLAN-GPE: Embedded telemetry protocol header
  - SRv6: Policy-Element (proof-of-transit only)
  - NSH: Type-2 Meta-Data (proof-of-transit only)
  - ... additional encapsulations being considered (incl. IPv4, MPLS)
- Deployment
  - Domain-ingress, domain-egress, and select devices within a domain insert/remove/update the extension header
  - Information export via IPFIX/Flexible-Netflow/publish into Kafka
  - Fast-path implementation





# Next Steps

- The authors appreciate thoughts, feedback, and text on the content of the documents from the SPRING WG
- Consider dedicated draft to specify POT TLV for SRH.