

# Fibbing: Central Control over Distributed Routing

[www.fibbing.net](http://www.fibbing.net)



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IRTF Open Meeting  
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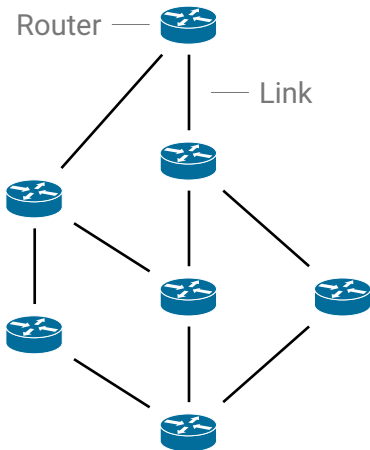
Joint work with  
S. Vissicchio (UCL), L. Vanbever (ETH Zürich) and J. Rexford (Princeton)

Fibbing

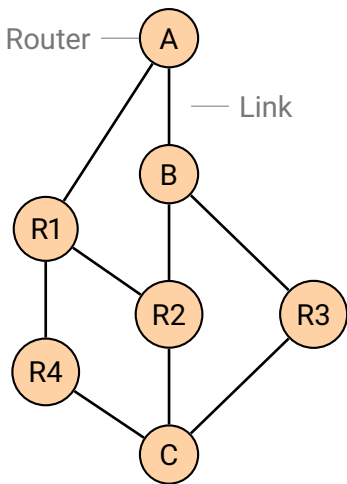
# Fibbing

Control routers' **FIB, lying** to routers

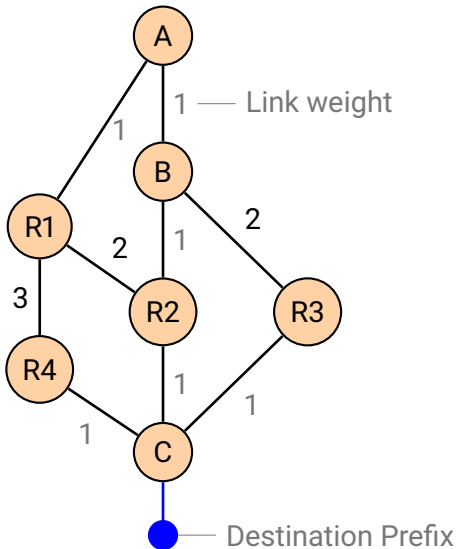
Consider this example network.



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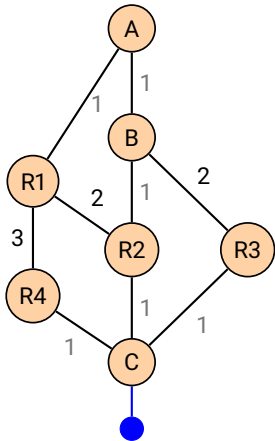


**Link-state Interior Gateway Protocols (IGPs)** exchange reachability information to infer the topology of the network.

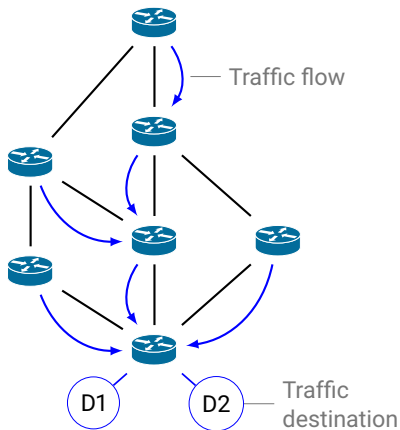


# The intra-domain traffic flows along the shortest path on the shared topology.

## Control-Plane



## Data-Plane



**IGPs cause operator to follow a *descriptive* management process.**

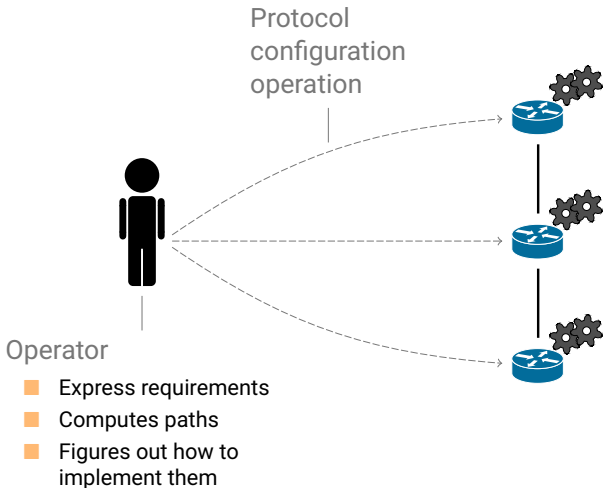


Operator

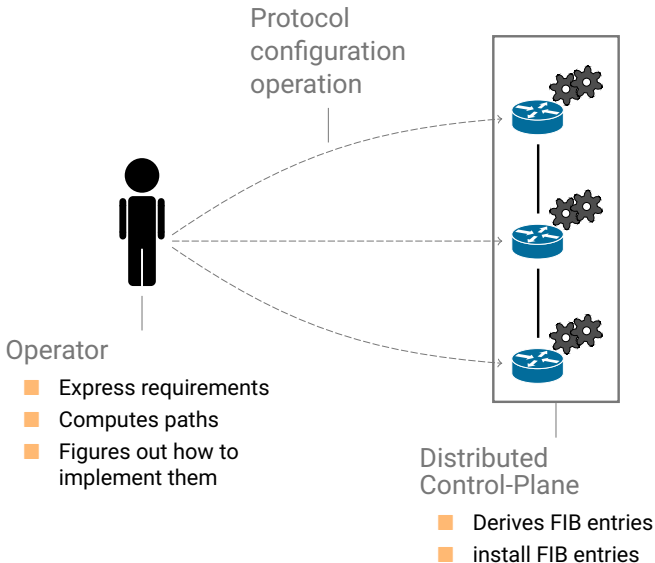
■ Express requirements



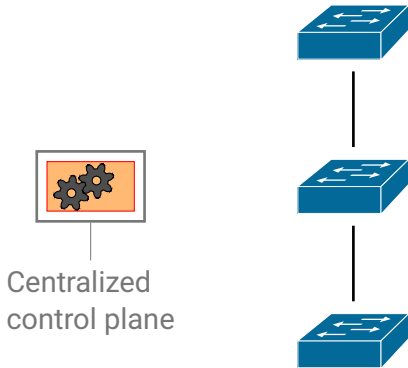
# IGPs cause operator to follow a *descriptive* management process.



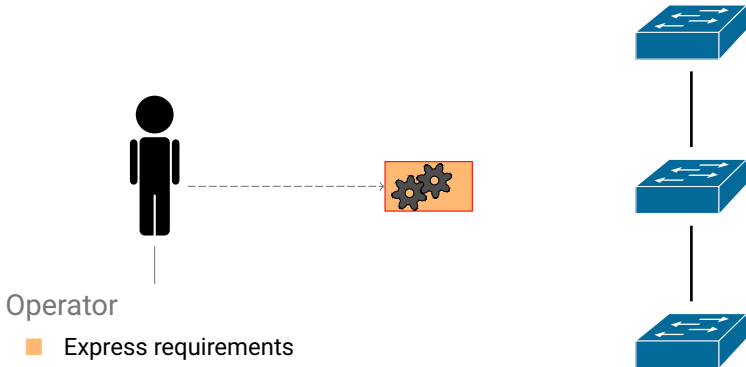
# IGPs cause operator to follow a *descriptive* management process.



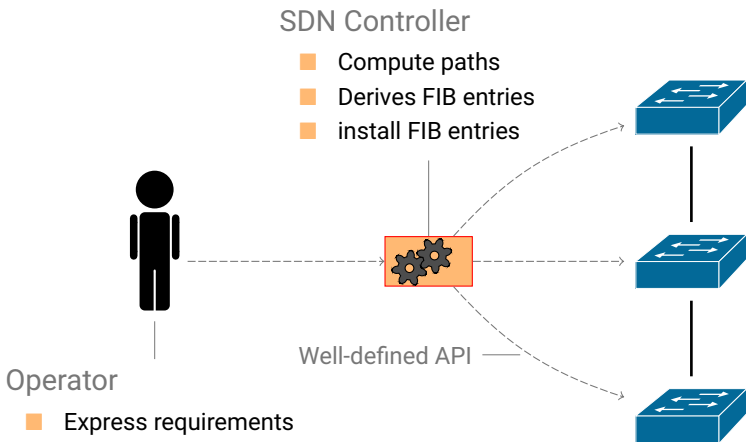
# Software-Defined Networking (SDN) enables *declarative* management.



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## SDN *sacrifices* the robustness and scalability of distributed protocols.

	Traditional	SDN
Manageability	low	<b>high</b>
Flexibility	low	<b>highest</b>
Scalability	<b>by design</b>	ad hoc
Robustness	<b>high</b>	low

## The networking world has two paradigm, based on opposed principles.

	Traditional		SDN
Manageability	low		<b>high</b>
Flexibility	low		<b>highest</b>
Scalability	<b>by design</b>		ad hoc
Robustness	<b>high</b>		low

We propose a middleground approach,  
named *Fibbing*.

	Traditional	Fibbing	SDN
Manageability	low	high	high
Flexibility	low	high	highest
Scalability	by design	by design	ad hoc
Robustness	high	high	low



# Fibbing: Central Control over Distributed Routing

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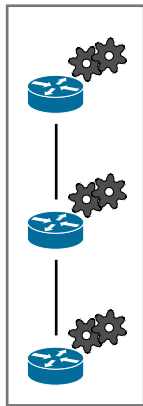


1. Controlling distributed protocols
2. Case study: surviving flash crowds
3. Fibbing today's networks
4. Food for thoughts

# Fibbing uses an hybrid control plane.

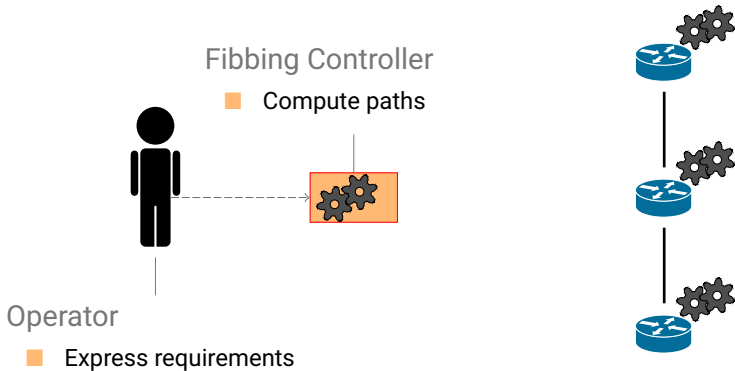


Centralized  
control plane

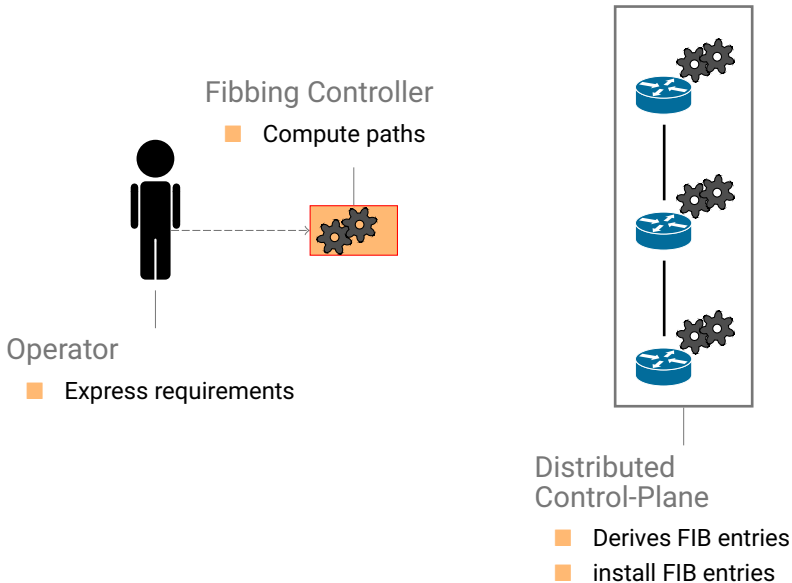


Distributed  
control plane

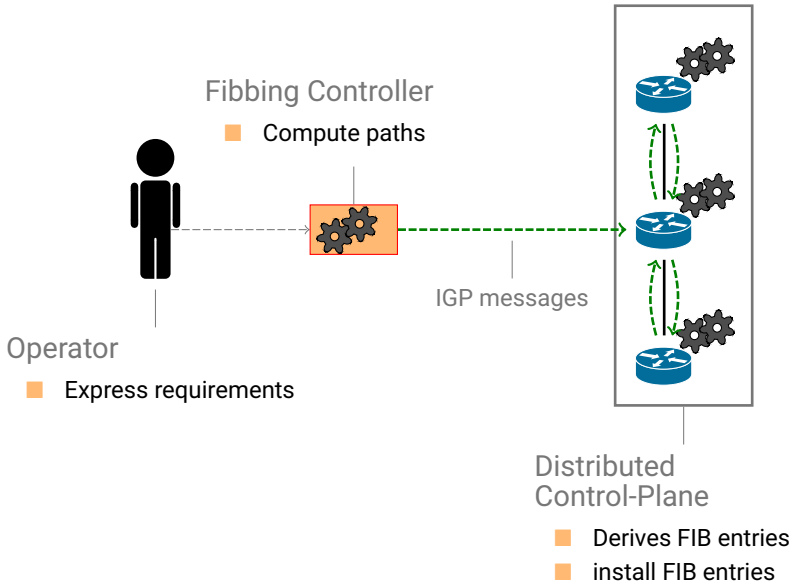
# Fibbing centralizes high level routing decisions.



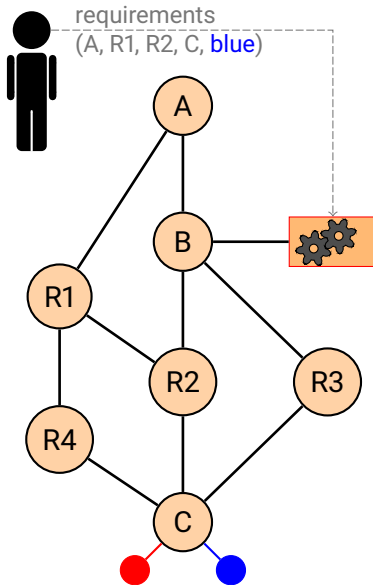
# Fibbing keeps the route installation distributed.



# We study which IGP messages to inject.

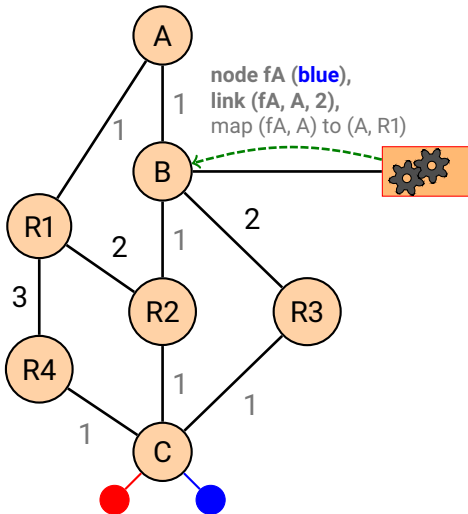


# Operators specify paths that must be enforced.



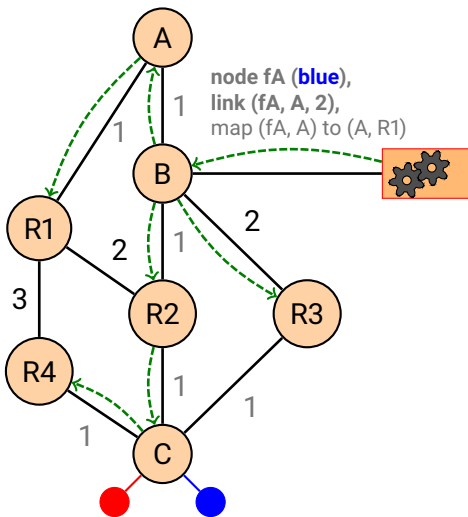
# The controller injects one IGP message adding a fake node and links.

requirements  
(A, R1, R2, C, blue)



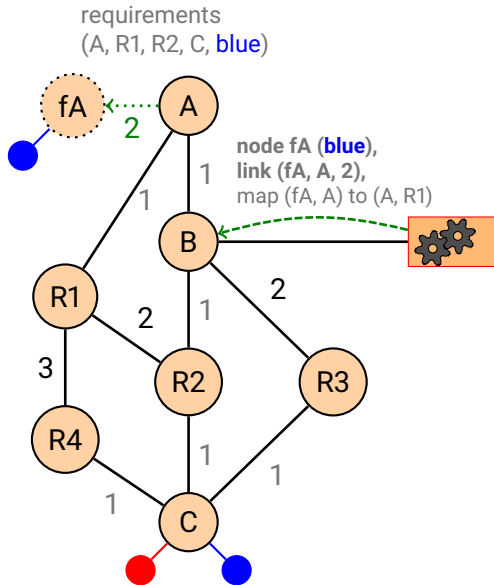
# IGP flooding propagates the information.

requirements  
(A, R1, R2, C, blue)

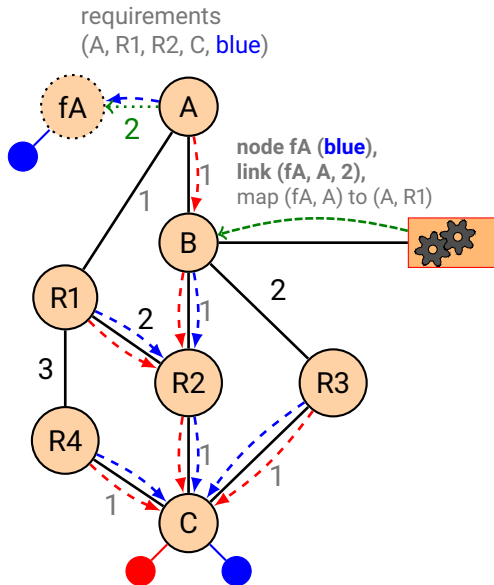




# The Fibbing message *augments* the topology.

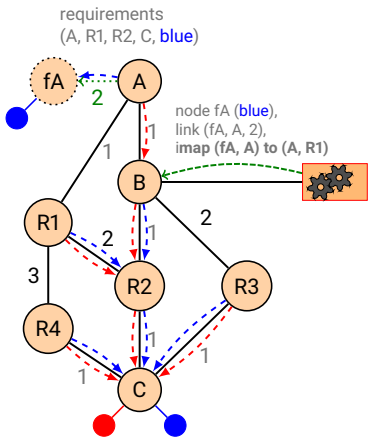


# Augmented topologies translate into new control-plane paths.

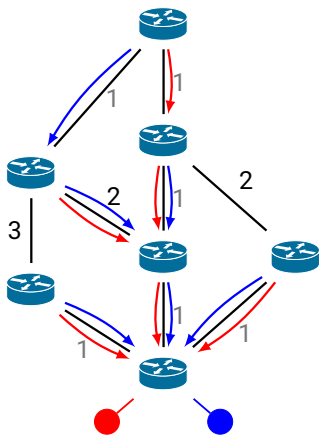


# Augmented topologies translate into new data-plane paths.

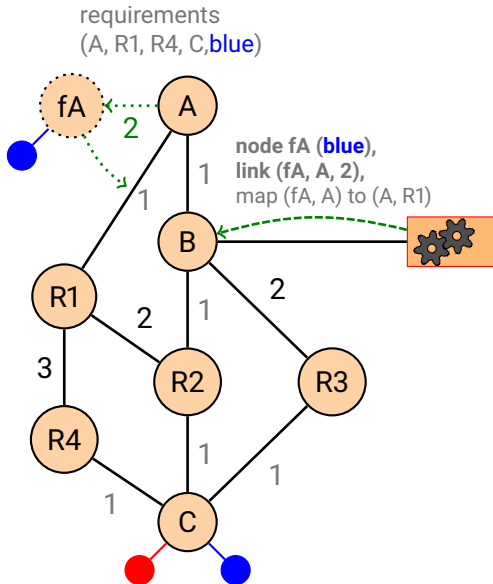
## Control-Plane



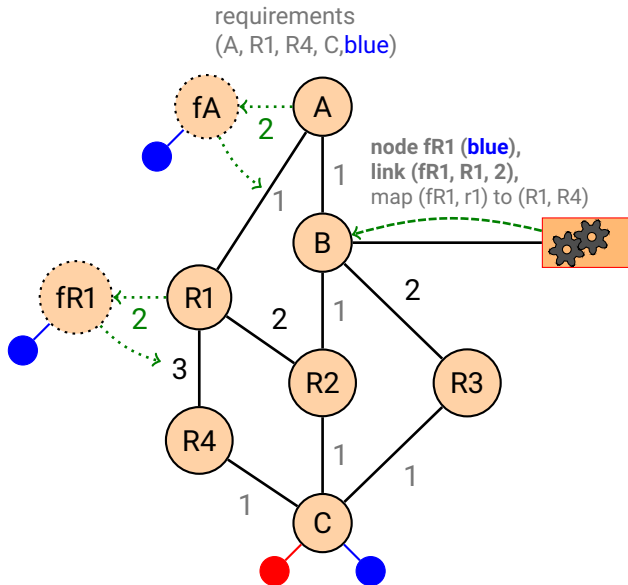
## Data-Plane



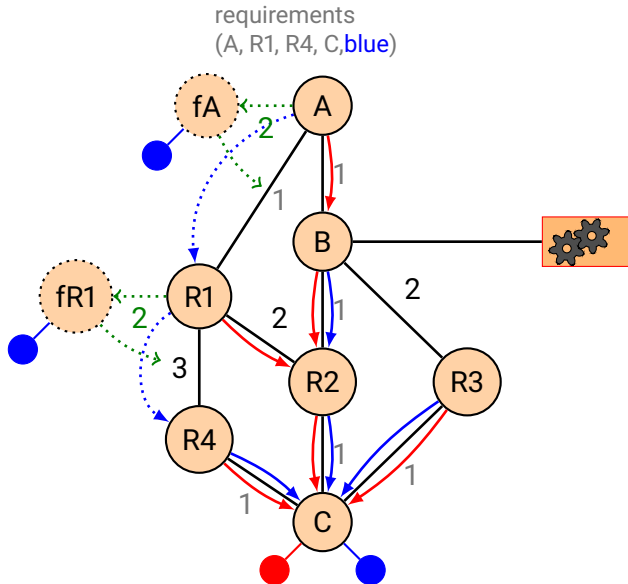
# Chaining multiple fake nodes enables to program complex paths.



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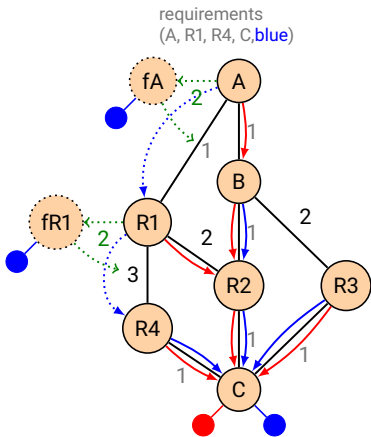


# Chaining multiple fake nodes enables to program complex paths.

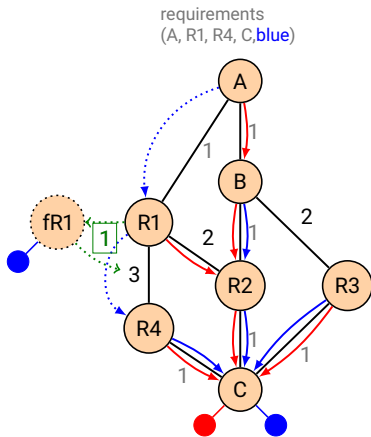


# Augmented topologies can be reduced to optimize the number of fake nodes.

## Naive augmentation



## Reduced augmentation



## Fibbing preserves the scalability of IGP.

- We can compute augmented topologies in  $\mathcal{O}(ms)$   
Ensures quick reaction to changes
  
- We can *reduce* augmented topologies in  $\mathcal{O}(s)$   
Ensures limited control-plane overhead



## Fibbing leverages the robustness of IGPs.

- Fast failure detection and recovery
- Survive controller failure
- Support fail-close and fail-open semantics

Fibbing can enforce any set of loop-free paths, on a per destination basis.

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## Flash crowds cause service disruption.

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- Video delivery services require good network performance
- Protecting the services against flash crowds is challenging:
  1. Traditional traffic engineering techniques perform poorly;
  2. Over-provisioning is expensive.

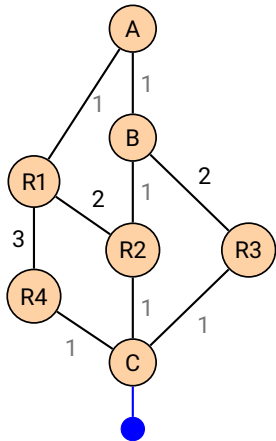
# Fibbing reduces the need for over-provisioning by enabling real-time traffic engineering.

## Experiment setup

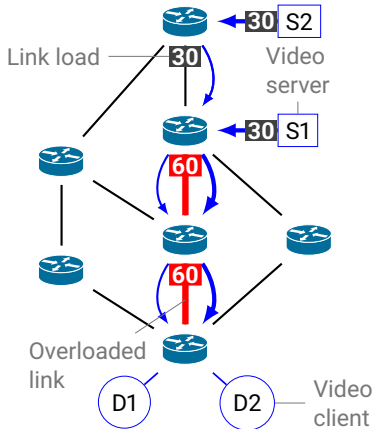
- Network with 2 video streaming servers
- Multiple clients are competing for bandwidth
- The network controller is able to detect flash crowds

# The initial IGP configuration has a bottleneck towards router C.

## Control-Plane

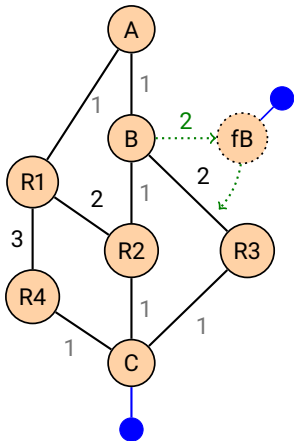


## Data-Plane

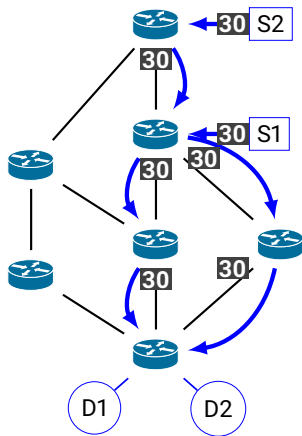


# Fibbing can program on-demand ECMP to spread the load

## Control-Plane



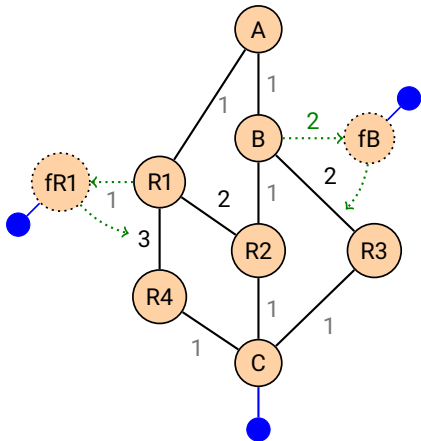
## Data-Plane



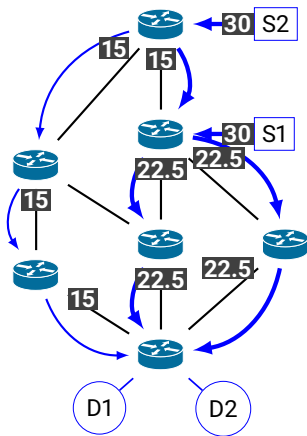


# Fibbing can program on-demand ECMP to spread the load

## Control-Plane

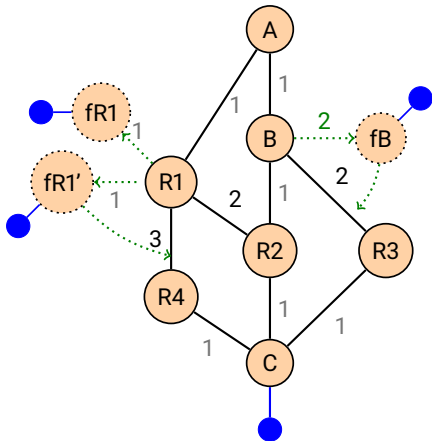


## Data-Plane

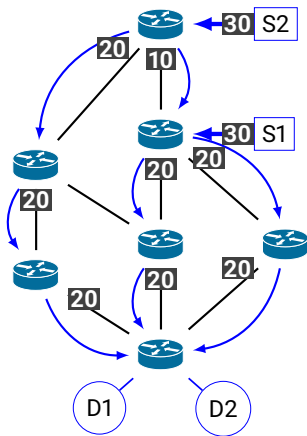


# Fibbing controls the splitting ratios across equal-cost paths.

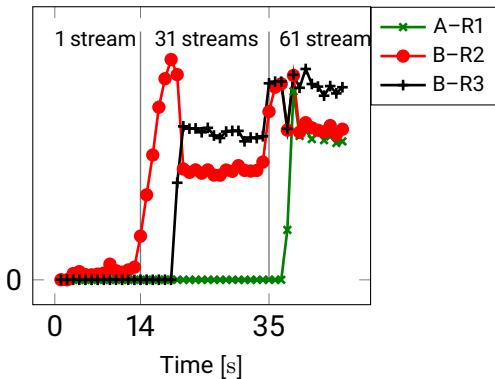
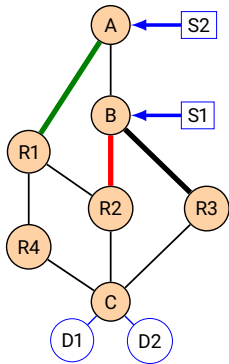
## Control-Plane



## Data-Plane



As the demand increases, the Fibbing controller adds more paths to spread the load.



- We initially have 1 video stream from S1 to D1.
  - At time  $t = 14s$ , we start 30 new streams from S1 to D1.
  - At time  $t = 35s$ , we start 30 streams from S2 to D2.

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## We have a working Fibbing controller prototype

- The controller maintains an OSPF adjacency to one router
- Topology discovery using the adjacency
- Tested against IOS, NX-OS, JunOS

## Fake nodes can be injected using LSA types 5/7

- Leverages the forwarding address field  
Advertise reachability towards prefix, with cost, using specified IP next hop
  
- The controller multiplexes multiple virtual routers  
 $\mathcal{N}$  successive fake nodes towards the same prefix require  
 $\mathcal{N}$  different router-ids

## Using T5/7 LSAs has (almost) no overhead on routers and is fast.

- No measurable impact on SPF duration
- 10 000 LSAs eat 14.5 MB of DRAM
- 900 $\mu$ s to push one fibbed route to the FIB

## Using T5/7 LSAs comes at a price

- Different expressivity model
- Can only affect prefixes from other T5/T7 LSAs
- Does not work with IS-IS!



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# Centrally modifying the shared topology is powerful

- Gives *some* control over BGP/MPLS-LDP
- Simplify configurations through exception-based routing
- Enables optimal, real-time TE

What would be the right abstraction?

# Fibbing: Central Control over Distributed Routing

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Tell me lies, tell me sweet little lies  
— Fleetwood Mac

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