

# An Improvement of ECN to Enhance TCP Fairness Performance

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# Background

- ❑ RFC 3168 allows E2E notification of network congestion:
  - CE can be marked by any of the network devices in the path
  - TCP receiver echoes CE to sender by ACK
  - Sender adjusts SWND according to CWND & ECN labeled packets ratio
  
- ❑ CANNOT accurately reflect network congestion status:
  - If multiple nodes exceed threshold value  $\implies$   
Sender received congestion status  $\geq$  Actual link congestion(the worst node status)

# Other ECN Limitations – 1/2

- ❑ Lack of good solution for light load in network:
  - Current ECN only handles with network congestion
  - If light loaded, lack of rapid notification to TCP sender
  - Unable to rapidly adjust SWND for better utilization of idled bandwidth

# Other ECN Limitations - 2/2

## ❑ TCP fairness problem:

- In Current ECN, different flows on a same path may attain discrepant link congestion status:
  - Packet transmission order depends on flows' own sending rate
  - Amounts of packets sent in 1 RTT vary → Proportions of ECN labelled packets vary
- Unfairness between flows' SWND adjustment ratio:
  - Especially for flows in same business scenarios/protocols

Example:

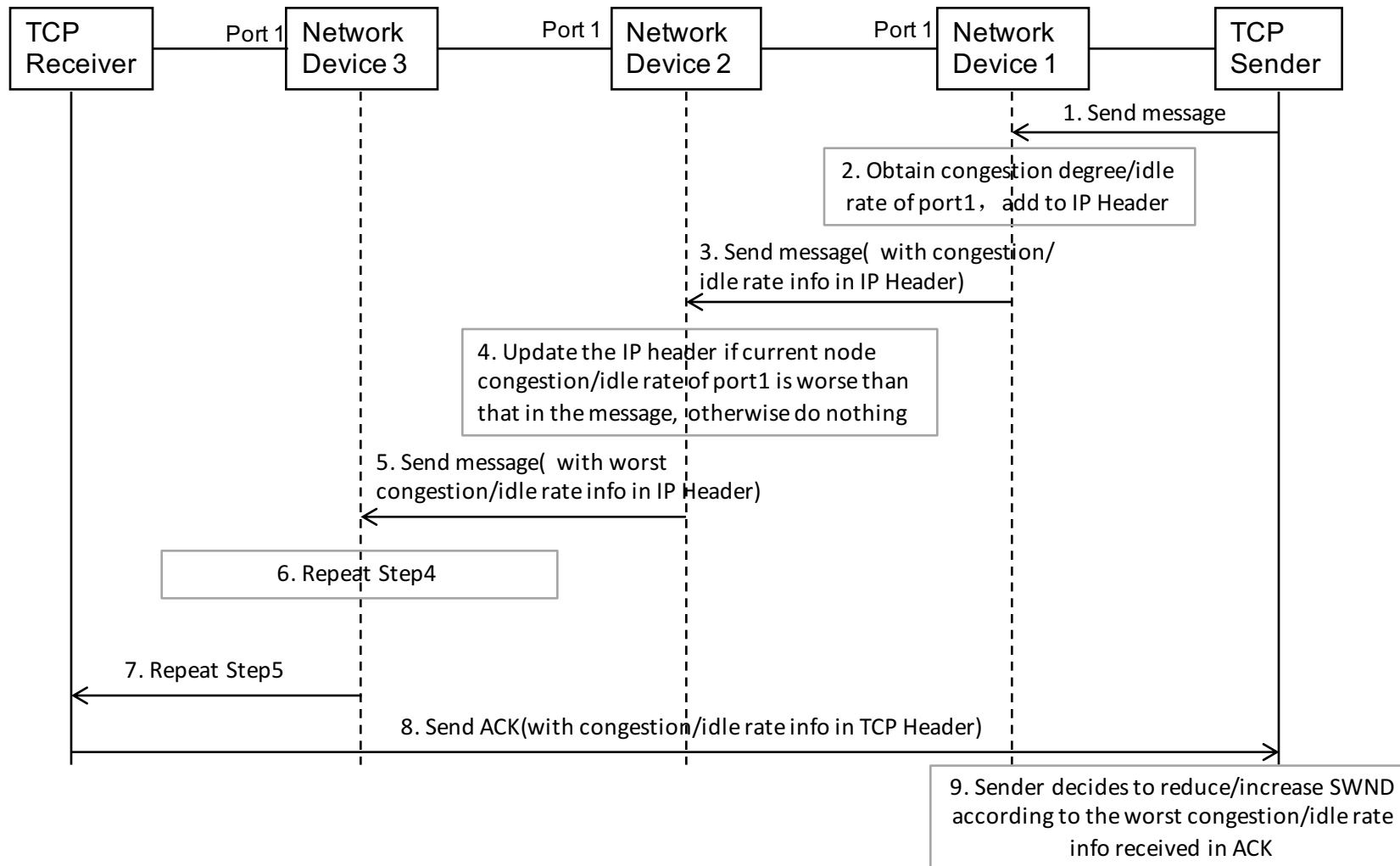
Assuming that 3 receivers are requesting the same sports live show stream in 4K resolution, If the 3 flows attain different link congestion status, flow rates will vary although they are in the same contents/protocols.

# Main Goal

- ❑ Optimize the ECN scheme, which can:
  - Reflect the worst congested node status for more accurate congestion control
  - Fully utilize the link idle rate in light load network situation
  - Achieve more fairness for different streams

# Improvement of ECN

The worst case (congestion or idle rate) of network devices is notified to TCP sender



- ❑ Congestion degree:
  - Time it takes to complete message transmission in the link cache
  - Example:
    - Cache size: 20MB, Link Bandwidth: 1Gbps
    - $20 \times 8 \text{ Mb} / 1\text{Gbps} = 160 / 1024 = 0.16\text{s}$
- ❑ Link idle rate:
  - Opposite to link usage
  - Example:
    - 1Gbps link with 600Mbps traffic
    - Link usage 60%, Link idle rate 40%

# Send Window Adjusting Method

## ❑ Option 1: Using the Worst Congestion Degree to Adjust SWND

- TCP sender adjusts the decrease rate of the window according to the worst congestion degree in the received TCP ACK and the current SWND

## ❑ Option 2: Using the Worst Idle Rate to Adjust SWND

- TCP sender adjusts the increase in window size based on the worst idle rate in the received TCP ACK and the current SWND

Example: Assuming that,

The worst idle rate in TCP ACK: 40%

The current window of Flow1: 1000

The current window of Flow2: 200

The current link utilization (total flow rate):  $1 - 40\% = 60\%$

The **flow rate can be increased by**  $40\% / 60\% = 66.67\%$

For Flow 1: the window should be increased by  $1000 \times 66.67\% = 667$

For Flow 2: the window should be increased by  $200 \times 66.67\% = 133$

# TCP/IP Option Extension

## ❑ Congestion Degree Extend

- Congestion degree carried in IP packets can be achieved by extending the IP option

```
+-----+-----+-----+
|           Type           |Length|Value|
+-----+-----+-----+
|node congestion degree|4 bytes| 0.1 |
+-----+-----+-----+
```

- The worst congestion degree carried by TCP ACK can be extended by TCP option

```
+-----+-----+-----+
|           Type           |Length|Value|
+-----+-----+-----+
|the worst congestion degree|4 bytes| 0.1 |
+-----+-----+-----+
```

## ❑ Idle Rate IP Extend

- Idle rate carried in IP packets can be achieved by extending the IP option

```
+-----+-----+-----+
|           Type           |Length|Value|
+-----+-----+-----+
|   node idle rate   |4 bytes| 0.45 |
+-----+-----+-----+
```

- The worst idle rate carried by TCP ACK can be extended by TCP option

```
+-----+-----+-----+
|           Type           |Length|Value|
+-----+-----+-----+
|the worst idle rate|4 bytes| 0.45 |
+-----+-----+-----+
```



# TCP Fairness Enhancement

- ❑ Each message is carrying the SAME worst node/port status (congestion/idle rate) in the same path
- ❑ TCP sender is no longer calculating the ratio of ECN labelled messages to modify the SWND, but adjusting window size according to the accurate worst node/port status
- ❑ Since each TCP flow attains the same worst node/port status, TCP sender give them same proportion on SWND adjustment, which indicates that their TCP fairness is guaranteed.

# Next Step

- ❑ Read and comment, please – thx for useful comments so far
  
- ❑ More experiments & verification:
  - Give some more verification supports in next version
  - Methods of more rapid notification to sender
  - Alternative method to present the worst link node status
  
- ❑ Etc.