

saag@IETF'98
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draft-goldbe-vrf-00

Verifiable Random Functions (VRF)

Sharon Goldberg (Boston University)

Dimitrios Papadopoulos (University of Maryland)

Jan Vcelak (ns1)

Contributors: Leonid Reyzin (Boston University), Shumon Huque (Salesforce),
David C. Lawrence (Akamai), Moni Naor & Asaf Ziv (Weizmann Institute)

hash function zoo

hash function:

SHA256

- no key
- hash = **H**(input)
- Verify: Check hash = **H**(input)

BLAKE

pseudorandom function:

HMAC

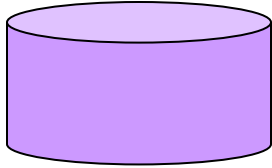
- symmetric key **k**
- hash = **H**(**k**, input)
- Verify: Cannot without **k**

verifiable random function (VRF):

- asymmetric key (**SK**, **PK**)
- hash = **VRF_hash**(**SK**, input)
- Verify: Use **PK**

VRF: verifiable random function

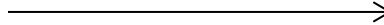
Verifier **PK**



Hasher **SK**



input



proof = **prove**(SK, input)

proof



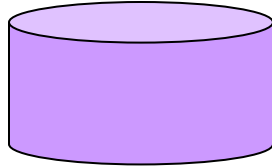
If **verify** (PK, input, proof)

hash = **proof2hash**(proof)

Else INVALID

VRF security: trusted uniqueness

Verifier **PK**



Hasher **SK**



input



proof = **prove**(**SK**, input)

proof



If **verify** (**PK**, input, proof)

hash = **proof2hash**(proof)

Else INVALID

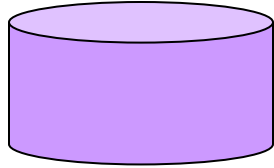
Like a regular hash function (eg SHA256)

Like a deterministic digital signature

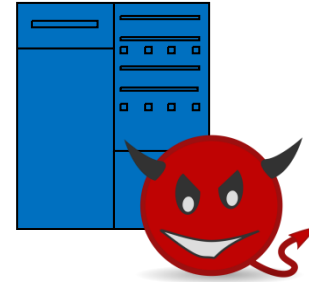
1-to-1 relationship between input and hash. (As with SHA-256!)

VRF security: trusted uniqueness

Verifier **PK**



Hasher **SK**



input



proof



proof = **prove**(SK, input)

If **verify**(PK, input, proof)

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Else INVALID

Like a regular hash function (eg SHA256)

Like a deterministic digital signature

1-to-1 relationship between input and hash. (As with SHA-256!)

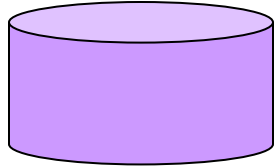
Trusted uniqueness:

Suppose the VRF keys (**PK,SK**) are generated in a trusted way.

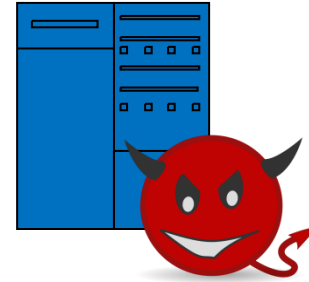
- If **PK** is fixed, then even an adversary that knows **SK** can't find
- ...two distinct VRF hash values that are valid for same input

VRF security: trusted collision resistance

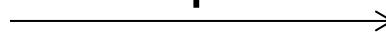
Verifier **PK**



Hasher **SK**



input



proof = **prove**(SK, input)

proof



If **verify** (PK, input, proof)

hash = **proof2hash**(proof)

Else INVALID

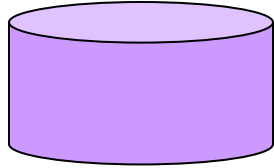
Like a regular hash function (eg SHA256)

Like a deterministic digital signature

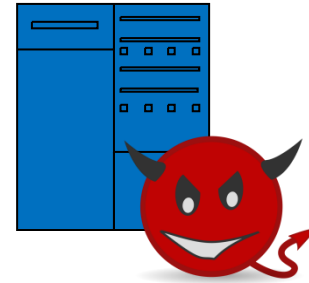
Collision resistance. (As with SHA-256!)

VRF security: trusted collision resistance

Verifier **PK**



Hasher **SK**



input



proof



proof = **prove**(SK, input)

Like a deterministic digital signature

If **verify** (PK, input, proof)

hash = **proof2hash**(proof)

Else INVALID

Like a regular hash function (eg SHA256)

Collision resistance. (As with SHA-256!)

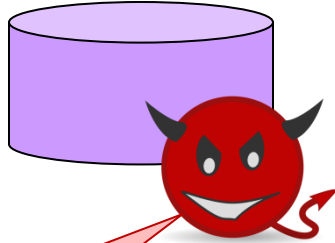
Trusted collision resistance:

Suppose the VRF keys (**PK,SK**) are generated in a trusted way.

- If **PK** is fixed, then even an adversary that knows **SK** can't find
- ...two distinct inputs that have the same valid VRF hash

VRF security: pseudorandomness

Verifier **PK**



I have no idea what input this hash corresponds to.

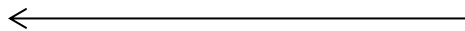
Hasher **SK**



proof = **prove** (**SK**, input)

hash = **proof2hash**(proof)

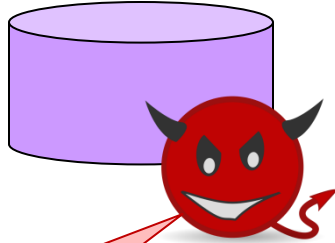
hash



Only the Hasher can compute the hash. (No dictionary attacks!)

VRF security: pseudorandomness

Verifier **PK**



I have no idea what input this hash corresponds to.

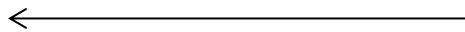
Hasher **SK**



proof = **prove** (**SK**, input)

hash = **proof2hash**(proof)

hash



Only the Hasher can compute the hash. (No dictionary attacks!)

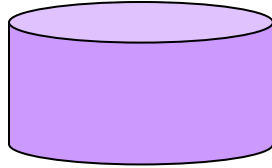
Pseudorandomness:

Suppose the VRF keys (**PK,SK**) are generated in a trusted way.

- Given an input, its VRF hash output looks pseudorandom
- ... to any adversary that does not know its proof or **SK**.

VRFs stop dictionary attacks on hash-based structures

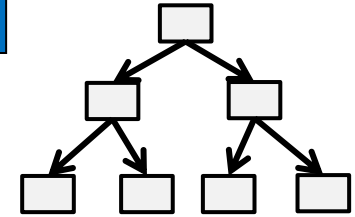
Verifier **PK**



Hasher **SK**



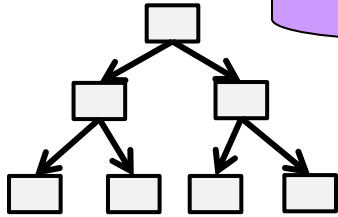
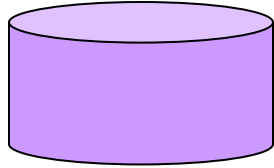
Build out of
VRF hashes



Hash-based data
structure

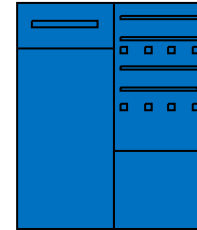
VRFs stop dictionary attacks on hash-based structures

Verifier **PK**

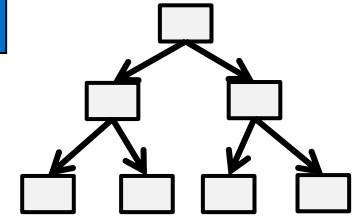


Hash-based data structure

Hasher **SK**



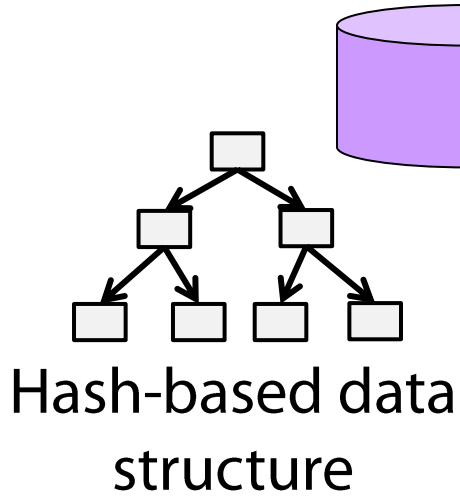
Build out of VRF hashes



Hash-based data structure

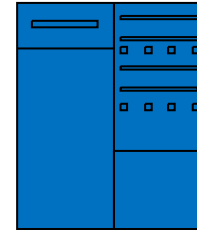
VRFs stop dictionary attacks on hash-based structures

Verifier **PK**

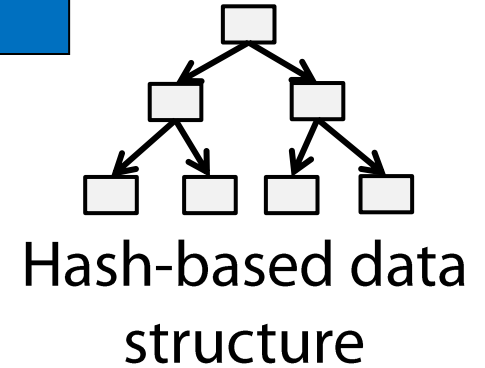


I can't use dictionary attacks to learn what is stored in this data structure.

Hasher **SK**

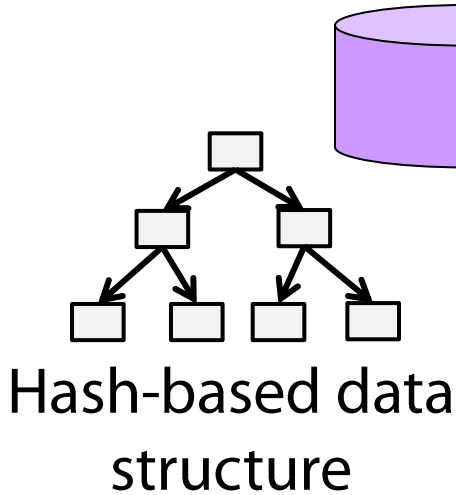


Build out of VRF hashes

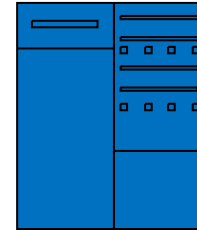


VRFs stop dictionary attacks on hash-based structures

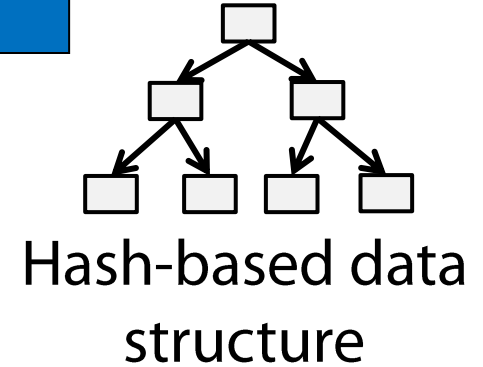
Verifier **PK**



Hasher **SK**

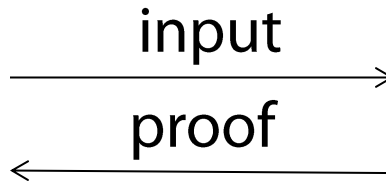


Build out of VRF hashes



I can't use dictionary attacks to learn what is stored in this data structure.

Is input in the data structure?



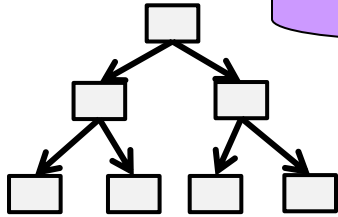
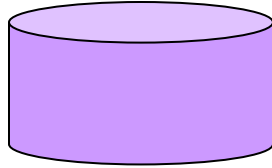
proof = **prove** (**SK**, input)

If **verify** (**PK**, input, proof)

hash = **proof2hash**(proof)

VRFs stop dictionary attacks on hash-based structures

Verifier **PK**

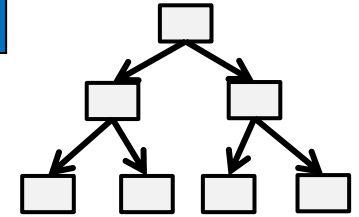


Hash-based data structure

Hasher **SK**



Build out of VRF hashes



Hash-based data structure

Is input in the data structure?

input

proof

proof = **prove** (**SK**, input)

If **verify** (**PK**, input, proof)

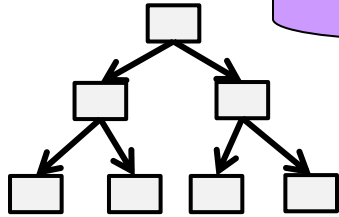
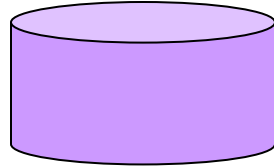
hash = **proof2hash**(proof)

Is hash in data structure?

Else INVALID

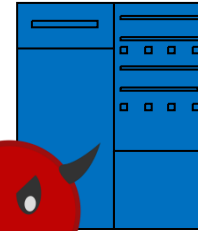
VRFs stop dictionary attacks on hash-based structures

Verifier **PK**

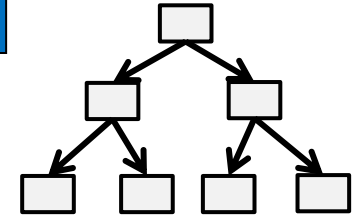


Hash-based data structure

Hasher **SK**



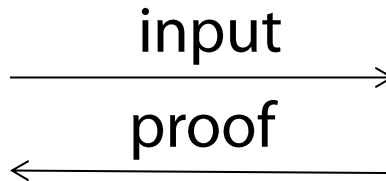
Build out of VRF hashes



Hash-based data structure

I can't lie about the hash output

Is input in the data structure?



proof = **prove** (**SK**, input)

If **verify** (**PK**, input, proof)

hash = **proof2hash**(proof)

Is hash in data structure?

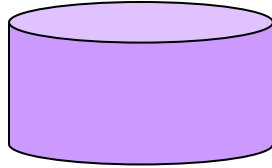
Else INVALID

-00 draft includes

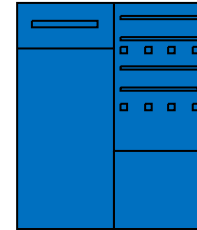
- VRF Security Definitions and Security Considerations
- Elliptic Curve VRF (**EC-VRF**)
 - Works with any cyclic group **G** of prime order **q** with generator **g**
 - Ciphersuites for NIST P-256 curve and Ed25519 curve
 - Algorithm is generic. Could support other curves
- RSA Full-Domain-Hash VRF (**RSA-FDH-VRF**)
- Also, we have:
 - Formal cryptographic security proofs: **<http://ia.cr/2017/099>**
 - Implementations: **<https://github.com/fcelda/nsec5-crypto>**

RSA-FDH-VRF (RSA full domain hash VRF)

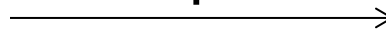
Verifier (N, e)



Hasher (N, d)



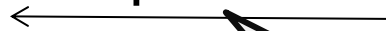
input



proof =

$$(\text{MGF1}(\text{input}))^d \bmod N$$

proof

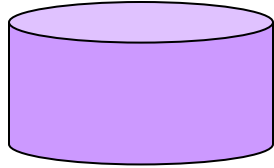


deterministic
RSA signature

[RFC8017]

RSA-FDH-VRF (RSA full domain hash VRF)

Verifier (N, e)



RSA signature verification

Hasher (N, d)



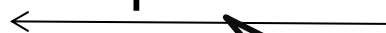
input



proof =

$$(\text{MGF1}(\text{input}))^d \bmod N$$

proof



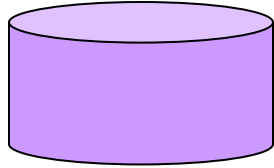
deterministic
RSA signature

[RFC8017]

$$\text{If } \text{MGF1}(\text{input}) = (\text{proof})^d \bmod N$$

RSA-FDH-VRF (RSA full domain hash VRF)

Verifier (N, e)



Hasher (N, d)



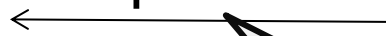
input



proof =

$$(\text{MGF1}(\text{input}))^d \bmod N$$

proof



RSA signature verification

If $\text{MGF1}(\text{input}) = (\text{proof})^d \bmod N$

hash = $\mathbf{H}(\text{proof})$

Else INVALID

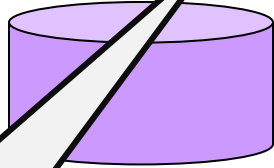
deterministic RSA signature

[RFC8017]

regular hash function
(eg SHA256)

EC-VRF (elliptic curve VRF)

Verifier g^x



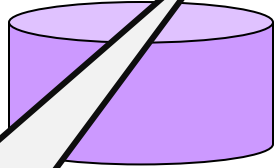
Hasher x



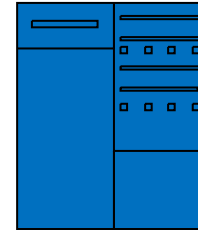
Cyclic group G of
prime order q
with generator g
(eg Ed25519)

EC-VRF (elliptic curve VRF)

Verifier g^x



Hasher x



input



Cyclic group G of
prime order q
with generator g
(eg Ed25519)

regular hash function
(eg SHA256)

$h = \text{hash_to_curve}(\text{input})$

$y = h^x$

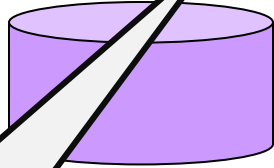
choose random nonce k

$c = H(g, g^x, h, h^x, g^k, h^k)$

$s = k - cx \text{ mod } q$

EC-VRF (elliptic curve VRF)

Verifier g^x



Hasher x



input

Cyclic group \mathbf{G} of
prime order \mathbf{q}
with generator \mathbf{g}
(eg Ed25519)

regular hash function
(eg SHA256)

$\mathbf{h} = \text{hash_to_curve}(\text{input})$

$\boldsymbol{\gamma} = \mathbf{h}^x$

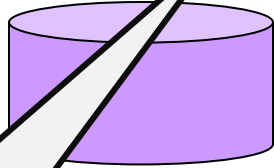
choose random nonce \mathbf{k}

$\mathbf{c} = \mathbf{H}(\mathbf{g}, \mathbf{g}^x, \mathbf{h}, \mathbf{h}^x, \mathbf{g}^{\mathbf{k}}, \mathbf{h}^{\mathbf{k}})$

proof: $(\boldsymbol{\gamma}, \mathbf{c}, \mathbf{s})$ $\mathbf{s} = \mathbf{k} - \mathbf{c}x \text{ mod } \mathbf{q}$

EC-VRF (elliptic curve VRF)

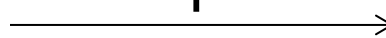
Verifier g^x



Hasher x



input



Cyclic group \mathbf{G} of
prime order \mathbf{q}
with generator \mathbf{g}
(eg Ed25519)

regular hash function
(eg SHA256)

$\mathbf{h} = \text{hash_to_curve}(\text{input})$

$\boldsymbol{\gamma} = \mathbf{h}^x$

choose random nonce \mathbf{k}

$\mathbf{c} = \mathbf{H}(\mathbf{g}, \mathbf{g}^x, \mathbf{h}, \boldsymbol{\gamma}, \mathbf{g}^{\mathbf{k}}, \mathbf{h}^{\mathbf{k}})$

proof: $(\boldsymbol{\gamma}, \mathbf{c}, \mathbf{s})$ $\mathbf{s} = \mathbf{k} - \mathbf{c}x \text{ mod } \mathbf{q}$

$\mathbf{u} = (\mathbf{g}^x)^{\mathbf{c}} \mathbf{g}^{\mathbf{s}}$

$\mathbf{h} = \text{hash_to_curve}(\text{input})$

$\mathbf{v} = \boldsymbol{\gamma}^{\mathbf{c}} \mathbf{h}^{\mathbf{s}}$

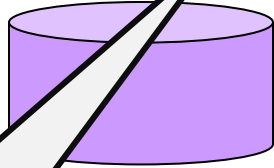
If $\mathbf{c} = \mathbf{H}(\mathbf{g}, \mathbf{g}^x, \mathbf{h}, \boldsymbol{\gamma}, \mathbf{u}, \mathbf{v})$

hash = x-coordinate of $\boldsymbol{\gamma}$

Else INVALID

EC-VRF (elliptic curve VRF)

Verifier g^x



Hasher x



input



Cyclic group G of
prime order q
with generator g
(eg Ed25519)

regular hash function
(eg SHA256)

$h = \text{hash_to_curve}(\text{input})$

$\gamma = h^x$

choose random nonce k

$c = H(g, g^x, h, h^x, g^k, h^k)$

proof: (γ, c, s) $s = k - cx \pmod q$

$u = (g^x)^c g^s$

$h = \text{hash_to_curve}(\text{input})$

$v = \gamma^c h^s$

If $c = H(g, g^x, h, \gamma, u, v)$

hash = x-coordinate of γ

Else INVALID

ciphersuites

- NIST P-256 curve with SHA256
- Ed25519 curve with SHA256
- Could add other curves (eg Ed448)