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IPsec profile of ACE
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Abstract

This document defines a profile of the ACE framework for authentication and authorization. It uses the IPsec protocol suite and the IKEv2 protocol to ensure secure communication, server authentication and proof-of-possession for a key bound to an OAuth 2.0 access token. The profile also defines an alternative key management approach to establish IPsec Security Associations using Ephemeral Diffie-Hellman Over COSE (EDHOC).

Status of This Memo

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1. Introduction

The IPsec protocol suite [RFC4301] allows communications based on the Constrained Application Protocol (CoAP) [RFC7252] to fulfill a number of security goals at the network layer, i.e. integrity and IP spoofing protection, confidentiality of traffic flows, and message replay protection. In several resource-constrained platforms, this can leverage security operations directly provided by hardware

crypto-modules, including mandatory-to-implement cipher suites defined in [RFC4835].

This document defines a profile of the ACE framework for authentication and authorization [I-D.ietf-ace-oauth-authz], where a client (C) and a resource server (RS) communicate using CoAP [RFC7252] over IPsec [RFC4301]. In particular, C uses an Access Token released by an Authorization Server (AS) and bound to a key (proof-of-possession key) to authorize its access to RS and its protected resources.

The establishment of an IPsec channel between C and RS provides secure communication, proof-of-possession as well as RS and C mutual authentication. Furthermore, this profile preserves the flexibility of IPsec as to the selection of specific security protocols, i.e. Encapsulating Security Payload (ESP) [RFC4303] and IP Authentication Header (AH) [RFC4302], key management, and modes of operations, i.e. tunnel or transport. Those parameters are specified in the IPsec Security Association (SA) pair established between C and RS.

This specification supports different key management methods for setting up SA pairs, namely direct provisioning of SA pairs and establishment of SA pairs based on symmetric or asymmetric key authentication. The latter approach can use the Internet Key Exchange Protocol version 2 (IKEv2) [RFC7296] or the Ephemeral Diffie-Hellman Over COSE (EDHOC) key exchange [I-D.selander-ace-cose-ecdhe].

1.1. Terminology

In this document, the key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" are to be interpreted as described in BCP 14, RFC 2119 [RFC2119] and indicate requirement levels for compliant CoAP-IPsec profile implementations.

Readers are expected to be familiar with terminology such as client (C), resource server (RS), authentication server (AS), and endpoint which are defined in [RFC6749] and [I-D.ietf-ace-actors].

The concept of IPsec Security Association (Section 4.1. of [RFC4301]) plays a key role, and this profile uses it extensively. An SA indicates how to secure a one-way communication between two parties. Hence, two SAs are required to be created and coordinated, in order to secure a two-way communication channel. This document refers to a SA pair as the two IPsec SAs used to protect the two-way communication channel between two IPsec peers.

The SA parameters described in section 4.4.2.1 of [RFC4301] are divided into the following two sets.

- Network Parameters: the parameters defining the network properties of the IPsec channel, e.g. DSCP filtering;
- Security Parameters: the parameters defining the security properties of the IPsec channel.

This document refers to SA-C as the SA for securing communication from C to RS, and to SA-RS as the SA for securing communication from RS to C. Thus, a SA pair consists of an SA-C and an SA-RS.

2. Methods for Setting Up SA Pairs

The following key management methods are supported for setting up a SA pair between C and RS.

1. Direct Provisoning (DP). The SA pair is pre-defined by the AS. Then, SA-RS and SA-C are specified in the Access Token Response and in the Access Token issued by the AS.
2. Establishment with symmetric key authentication. A symmetric Pre-Shared Key (PSK) is used to authenticate both parties during the SA pair establishment and is bound to the Access Token as proof-of-possession key. If C is interacting for the first time with the RS, then the AS MUST include a PSK and a unique key identifier in the Access Token Response. Otherwise, C MUST include the unique key identifier pointing at the previously established PSK in the Access Token Request.
3. Establishment with asymmetric key authentication. An asymmetric Raw Public Key (RPK) or Certificate-based Public Key (CPK) is used to authenticate both parties during the SA pair establishment and is bound to the Access Token as proof-of-possession key. If the AS does not know C’s asymmetric authentication information, then C MUST include its RPK or CPK in the Access Token Request. Otherwise, C MUST include a key identifier linked to its own RPK or CPK available at the AS.

Every SA MUST include the following Security Parameters.

- A Security Parameter Index (SPI);
- IPsec protocol mode: tunnel or transport;
- Security protocol: AH or ESP;

- "AH-authentication", "ESP-encryption", "ESP-integrity" or "ESP-combined" algorithm;
- Source and destination, if tunnel mode is selected;
- Cryptographic keys;
- SA lifetime.

As assumed in Section 5.5.2 of [I-D.ietf-ace-oauth-authz], the AS has knowledge of C's and RS's capabilities, and of RS's preferred communications settings. Therefore, the AS MUST set the values of Security Parameters and Network Parameters in the SA pair.

2.1. The "ipsec" Structure

This document defines the "ipsec" structure as a field of the "cnf" parameter of the Access Token and Access Token Response. This structure encodes the Network and Security Parameters of the SA pair as defined in Figure 1. The Network Parameters are not discussed in this specification.

```
ipsec{
        <Security Parameters>,
        <Network Parameters>
}
```

Figure 1: "ipsec" structure overview.

The AS builds the "ipsec" structure as follows:

- The Security Parameters MUST always include the set of parameters sec_A shown in Figure 2.
- The Security Parameters MUST include the set of parameters sec_B shown in Figure 3 if the AS uses the Direct Provisioning method.

```
sec_A{
        mode,
        protocol,
        life,
        IP_C,  (if mode == tunnel)
        IP_RS  (if mode == tunnel)
}
```

Figure 2: Set sec_A of Security Parameters

```
sec_B{
        SPI_SA_C,
        SPI_SA_RS,
        alg,
        seed
}
```

Figure 3: Set sec_B of Security Parameters

In sec_A, the IP_C field is the IP address of C, while IP_RS is the IP address of RS. In tunnel mode, the RS MUST use IP_C as the destination address and IP_RS as source address of outgoing IPsec messages. Similarly, C MUST use IP_RS as destination address and IP_C as source address of incoming IPsec messages.

In sec_B, the field "SPI_SA_C" is the SPI of SA-C. Similarly, "SPI_SA_RS" is the SPI of SA-RS. The field "alg" indicates the algorithm used for securing communications over IPsec. The "seed" field MUST reflect the SKEYSEED secret defined in Section 2.14 of [RFC7296]. Thus, C and RS MUST use the same key derivation techniques to generate the necessary SA keys from "seed".

Note that if the Direct Provisioning method is used, the AS cannot guarantee the uniqueness of the "SPI_SA_C" value at the RS and of the "SPI_SA_RS" value at C. In such a case, the AS MUST randomly generate the "SPI_SA_C" value and the "SPI_SA_RS" value, so that the probability of a collision to occur is negligible.

If RS receives an "SPI_SA_C" value which results in a collision, then RS MUST reply to C with en error response, and both C and RS MUST abort the set up of the IPsec channel. In order to overcome this issue, the AS can manage a pool of "SPI_SA_C" reserved values, intended only for use with the Direct Provisioning method. Then, in case of SA termination, the RS asks the AS to set back the identifier of that SA-C as available.

If C receives an "SPI_SA_RS" value which results in a collision, then C sends a second Token Request to the AS, asking for a Token Update. This Token Request includes also an "ipsec" structure, which contains only the field "SPI_SA_RS" specifying an available value to use. Then, the AS replies with an Access Token and an Access Token Response both updated as to the "SPI_SA_RS" value only.

3. Protocol Description

This profile considers a client C that intends to access a protected resource hosted by a resource server RS. The resource access is authorized through an Access Token issued by the AS as specified in

[I-D.ietf-ace-oauth-authz] and indicating that IPsec is used to secure communications between C and RS. In particular, this profile defines how C and RS set up a SA pair, using the key management methods introduced in Section 2.

The protocol is composed of three parts, as shown in Figure 4.

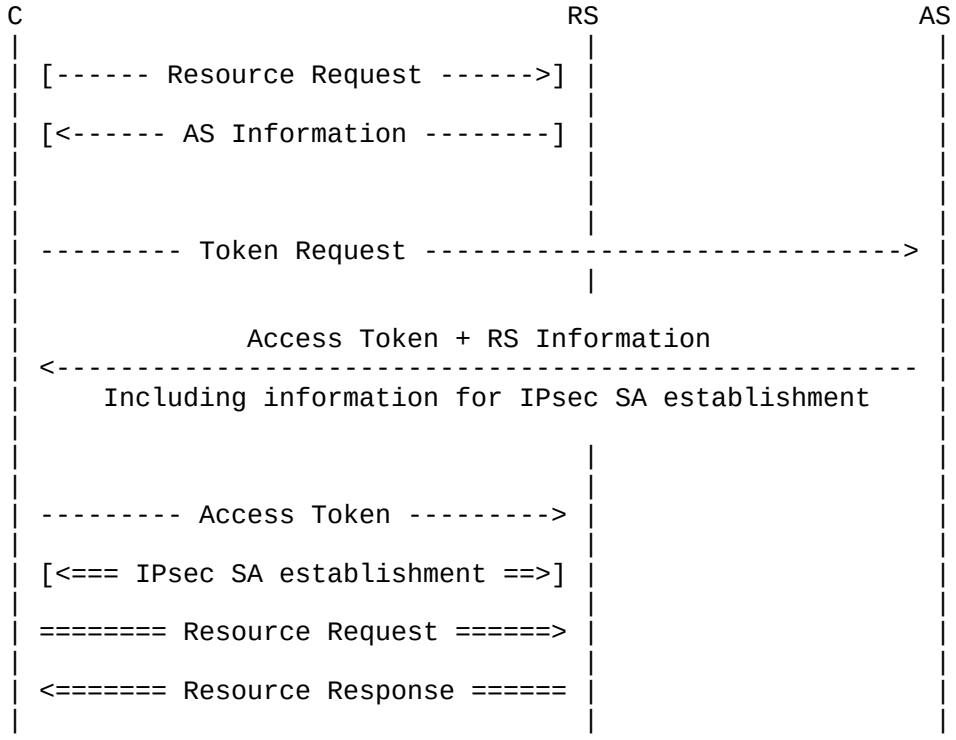

Figure 4: Protocol Overview

3.1. Unauthorized Client to RS

Phase (1) in Figure 4 is OPTIONAL and aims at providing C with the necessary information to contact the AS, in case C does not know AS's address. Through an unauthorized request to RS, C determines which AS is responsible for granting authorization to that particular RS. When doing so, C learns to which address the Access Token Request has to be addressed. The unauthorized request is denied by RS, which sends back to C a response containing the information to contact the AS.

3.2. Client to AS

Phase (2) in Figure 4 starts with C sending the Access Token Request to the /token endpoint at the AS, as specified in Section 5.5.1 of [I-D.ietf-ace-oauth-authz]. Figure 2 and Figure 3 of [I-D.ietf-ace-oauth-authz] provide examples of such request.

If the AS successfully verifies the Access Token Request and C is authorized to access the resource specified in the Token Request, then the AS issues the corresponding Access Token and includes it in a CoAP response with code 2.01 (Created) as specified in Section 5.5.2 of [I-D.ietf-ace-oauth-authz]. The AS can signal that IPsec is REQUIRED to secure communications between C and RS by including the "profile" parameter with the value "coap_ipsec" in the Access Token Response. Together with authorization information, the Access Token also includes the same information for the set up of the IPsec channel included in the Access Token Response. The error responses are handled as described in Section 5.5.3 of [I-D.ietf-ace-oauth-authz].

The information exchanged between C and the AS depends on the specific method used to set up the SA pair (see Section 3.2.1, Section 3.2.2 and Section 3.2.3). Note that, unless Direct Provisioning of SAs is used, C and RS are required to finalize the SA pair set up by running a Key Management Protocol such as IKEv2 (see Section 3.3.2). The AS indicates to use IKEv2 for establishing a SA pair by setting the "kmp" field to "ikev2" in the "cnf" parameter in the Access Token Response. An alternative key management method based on Ephemeral Diffie-Hellman Over COSE (EDHOC) [I-D.selander-ace-cose-ecdhe] is described in Appendix B.

The communications between the AS and C (/token endpoint) rely on the CoAP protocol and MUST be secured. In particular, they can be secured by means of Object Security of OSCOAP [I-D.ietf-core-object-security]} as described in section 5.5 of [I-D.ietf-ace-oauth-authz], or through other means such as IPSec [RFC4301] or DTLS [RFC6347].

3.2.1. Direct Provisioning of SA pairs

If the AS selects this key management method, it encodes the SA pair in the Access Token and in the Access Token Response as an "ipsec" structure in the "cnf" parameter.

Figure 5 shows an example of an Access Token Response, signaling C to set up an IPsec channel with RS based on the ESP protocol in transport mode.

```
Header: Created (Code=2.01)
Content-Type: "application/cose+cbor"
Payload : {
    "access_token" : b64'YiksuH&=1GFfg ...
    (remainder of Access Token omitted for brevity)',
    "profile" : "coap_ipsec",
    "expires_in" : "3600",
    "cnf" : {
       "ipsec" : {

             "mode"      : "transport",
             "protocol"  : "ESP",
             "life"      : "3600",
             "SPI_SA_C"  : "87615",
             "SPI_SA_RS" : "87616",
             "seed"      : b64'+a+Dg2jjU+eIiOFCa9lObw',
             "alg"       : "AES-CCM-16-64-128",
             ... (the Network Parameters are omitted for brevity),
       }
    }
}
```

Figure 5: Example of Access Token Response with DP of SA pair

3.2.2. SA Establishment Based of Symmetric Keys

If the AS selects this key management method, it specifies the following pieces of information in the Access Token Response and in the Access Token:

- a symmetric key to be used as proof-of-possession key;
- a key identifier associated to the symmetric key;
- SA pair's Network Parameters and Security Parameters, as an "ipsec" structure in the "cnf" parameter (see Section 2.1).

If C has previously received a PSK from the AS, then C MUST provide a key identifier of that PSK either directly in the "kid" field of "cnf" parameter or in the "kid" field of the "COSE_Key" object of the Access Token Request. In this case, the AS omits the PSK and its identifier in the Access Token Response.

The AS indicates the use of symmetric cryptography for the key management message exchange in the "kty" field of the "COSE_Key" object, including also the PSK in the "k" field as well as its key identifier in the "kid" field, as shown in Figure 6.

```
Header: Created (Code=2.01)
Content-Type: "application/cose+cbor"
Payload:
  {
    "access_token" : b64'YiksuH&=1GFfg ...
    (remainder of Access Token omitted for brevity)',
    "profile" : "coap_ipsec",
    "expires_in" : "3600",
    "cnf" : {
      "COSE_Key" : {
      "kty" : "Symmetric",
      "kid" : b64'6kwi42ec',
      "k"   : b64'+pAd48jU+eIiOF23gd=',
      }
      "kmp": "ikev2",
      "ipsec" : {
           "mode"     : "tunnel",
           "protocol" : "ESP",
           "life"     : "1800",
           "IP_C"     : "a.b.c.d2",
           "IP_RS"    : "a.b.c.d1",
           ... (the Network Parameters are omitted for brevity),
      }
    }
  }
```

Figure 6: Example of Access Token Response with a symmetric key as proof-of-possession key.

3.2.3. SA Establishment Based on Asymmetric Keys

C MUST include its own public key in the Access Token Request, as shown in Figure 7. As an alternative, C MUST provide the key identifier of its own public key, previously shared with the AS.

The AS specifies in the Access Token and in the Access Token Response the SA pair's Network Parameters and Security Parameters, as an "ipsec" structure in the "cnf" parameter (see Section 2.1).

In addition, the AS specifies the RS's public key in the Access Token Response, and the C's public key to be used as proof-of-possession key in the Access Token.

The AS indicates the use of asymmetric cryptography for the key management message exchange in the "kty" field of the "COSE_Key" object, which includes also the RS's public key in the Access Token Response and the C's public key in the Access Token.

```
Header: POST (Code=0.02)
Uri-Host: "server.example.com"
Uri-Path: "token"
Content-Type: "application/cose+cbor"
Payload:
{
   "grant_type" : "client_credentials",
   "cnf" : {
     "COSE_Key" : {
             "kty" : "EC",
             "crv" : "P-256",
             "x"   : b64'CaFadPPavdtjRH3YqaTqm0FrFtNV0',
             "y"   : b64'ehekJBwciJdeT6cKieycnk6kg4pHC'
     }
   }
 }
```

Figure 7: Example of Access Token Request with an asymmetric key as proof-of-possession key.

3.3. Client to RS

Phase (3) in Figure 4 starts with C posting the Access Token by means of a POST CoAP message to the /authz-info endpoint at RS, as specified in Section 8 of [I-D.ietf-ace-oauth-authz]. The processing details of this request, as well as the handling of invalid Access Tokens at RS, are defined in Section 5.7.1 of [I-D.ietf-ace-oauth-authz] and in the rest of this section. The Access Token and Access Token Response specify one of the SA setup methods defined in Section 2. In particular, C and RS determine the specific SA setup method as follows:

- In case of Direct Provisioning, the "ipsec" structure is present, while the "COSE_Key" object is not present.
- If the SA pair set up based on Symmetric Keys through IKEv2 is used, then:
  - the "COSE_Key" object is present with the "kty" field set to "Symmetric"; and
  - the "kmp" parameter is set to "ikev2".
- If the SA pair set up based on Asymmetric Keys through IKEv2 is used, then:

* the "COSE_Key" object is present with the "kty" field set to a value that indicates the use of an asymmetric key, e.g. "EC"; and

* the "kmp" parameter is set to "ikev2".

If the Direct Provisioning method is used, then C and RS do not perform the SA establishment shown in Figure 4. Otherwise, C and RS perform the key management protocol indicated by the "kmp" parameter (such as IKEv2), in the authentication mode indicated by the "kty" field of the "COSE_key" object.

Regardless the chosen SA setup method and the successful establishment of the IPsec channel, if C holds a valid Access Token but this does not grant access to the requested protected resource, RS MUST send a 4.03 (Forbidden) response. Similarly, if the Access Token does not cover the intended action, RS MUST send a 4.05 (Method Not Allowed) response.

3.3.1. SA Direct Provisioning

Once received a positive Access Token Response from the AS, C derives the necessary IPsec key material from the "seed" field of the "ipsec" structure in the Access Token Response, as discussed in Section 2.1. Similarly, RS performs the same key derivation process upon receiving and successfully verifying the Access Token. After that, RS replies to C with a 2.01 (Created) response, using the IPsec channel specified by the SA pair. Thereafter, Resource Requests and Responses are also sent using the IPsec channel.

3.3.2. Authenticated SA Establishment

If an Authenticated Key Management method is used (see Section 3.2.2 and Section 3.2.3), C and RS MUST run a Key Management Protocol to finalize the establishment of the SA pair and the IPsec channel, i.e. the required keys and algorithms. As shown in Figure 8, the first message IKE_SA_INIT of the IKEv2 protocol is used to acknowledge the Access Token submission. Depending on the used authentication method, i.e. symmetric or asymmetric, the proof-of-possession key MUST be used accordingly to authenticate the IKEv2 message exchange as defined in Section 2.15 of [RFC7296]. The rest of the IKEv2 protocol MUST be executed between C and RS as described in Section 2 of [RFC7296], with no further modifications. If IKEv2 is successfully completed, C and RS agree on keys and algorithms to use, and thus the IPsec channel between C and RS is ready to be used.

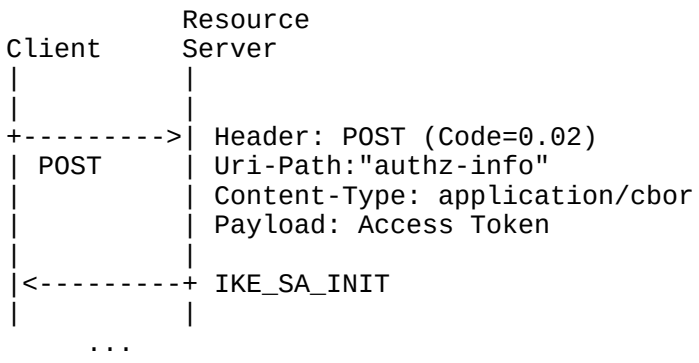

Figure 8: IKEv2 used as Key Management Protocol.

3.4. RS to AS

The communication between RS and the AS occurs through the /introspect endpoint, as defined in Section 5.6 of [I-D.ietf-ace-oauth-authz]. This channel SHOULD be confidential, authenticated and integrity protected. Therefore, RS and the AS are expected to have pre-established credentials and to use security protocols such as OSCOAP, DTLS or IPsec to protect their communications.

4. Security Considerations

This document inherits the security considerations of [RFC4301], [RFC4302] and [RFC4303]. Furthermore, if IKEv2 is used as key establishment method (see Section 3.3.2), the same considerations discussed in [RFC7296] hold.

4.1. Privacy Considerations

The message exchange in Phase (1) of Figure 4 is unprotected and MAY disclose the relation between the AS, RS and C, as well as network related information, such as IP addresses. Thus RS SHOULD only include the necessary information to contact the AS in the unprotected response.

5. IANA Considerations

This document requires the following IANA considerations:

name	label	CBOR type	value	description
kmp	TBD	bstr	ikev2	Indicates the key management protocol to be used to establish a SA pair
ipsec	TBD	struct		Contains Security and Network Parameters of an SA pair

5.1. CoAP-IPsec Profile registration

- Profile name: CoAP-IPsec
- Profile description: ACE Framework profile
- Profile ID: coap_ipsec
- Change Controller: IESG
- Specification Document: This document

5.2. Confirmation Methods registration

5.2.1. IPsec field

- Confirmation Method Name: "ipsec"
- Confirmation Method Value: TBD
- Confirmation Method Description: A structure containing the corresponding information of an IPsec Security Association Pair.
- Change Controller: IESG
- Specification Document: This document

5.2.2. Key Management Protocol field

- Confirmation Method Name: "kmp"
- Confirmation Method Value: TBD

- Confirmation Method Description: Key management protocol.

- Change Controller: IESG

- Specification Document: This document

5.3. Key Management Protocol Methods Registry

This specification establishes the IANA "Key Management Protocol Methods" registry for the "kmp" member values. The registry records the confirmation method member and a reference to the spec that defines it.

5.3.1. Registration Template

Key Management Protocol Method Name:

The name requested (e.g. "ikev2"). This name is intended to be human readable and be used for debugging purposes. It is case sensitive. Names may not match other registered names in a case-insensitive manner unless the Designated Experts state that there is a compelling reason to allow an exception.

Key Management Protocol Method Value:

Integer representation for the confirmation method value. Intended for use to uniquely identify the confirmation method. The value MUST be an integer in the range of 1 to 65536.

Key Management Protocol Method Description:

Brief description of the confirmation method (e.g. "Key Identifier").

Change Controller:

For Standards Track RFCs, list the "IESG". For others, give the name of the responsible party. Other details (e.g. postal address, email address, home page URI) may also be included.

Specification Document(s):

Reference to the document or documents that specify the parameter, preferably including URIs that can be used to retrieve copies of the documents. An indication of the relevant sections may also be included but is not required.

5.3.2. Initial Registry Contents

- Key Management Protocol Method Name: "ikev2"
- Key Management Protocol Method Value: TBD
- Key Management Protocol Method Description: Defines IKEv2 as key management protocol.
- Change Controller: IESG
- Specification Document: this document

6. Acknowledgments

The authors sincerely thank Max Maass for his comments and feedback.

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7.1. Normative References

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7.2. Informative References

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Appendix A. Coexistence of OSCOAP and IPsec

Object Security of CoAP (OSCOAP) [I-D.ietf-core-object-security] is a data object based security protocol that protects CoAP messages end-to-end while allowing proxy operations. It encloses unprotected CoAP messages, and selected CoAP options and headers fields into a CBOR Object Signing and Encryption (COSE) object [I-D.ietf-cose-msg]. This section describes a scenario where communications between C and RS are secured by means of OSCOAP and IPsec. Figure 9 depicts a scenario where a Client needs to access a Resource Server which is behind an untrusted CoAP-Proxy. This scenario requires that:

1. the Proxy has access to the selected CoAP options to perform management and support operations;

2. the integrity of messages and their IP headers can be verified by the Resource Server;

3. the confidentiality of the Resource Server address and CoAP request has to be guaranteed between the Client and the Proxy.

The first requirement is addressed by means of an OSCOAP channel between the Client and the Resource Server established as described in [I-D.seitz-ace-oscoap-profile]), by marking as Class E the sensitive fields of the CoAP payload as defined in [I-D.ietf-core-object-security].

To address the second requirement, a SA pair between the Client and the Resource Server is established, as specified in Section 3, by using the IPsec AH protocol in transport mode. Finally, the third requirement is fulfilled by means of a SA pair between the Client and the CoAP-Proxy, as specified in Section 3, by using the IPsec ESP protocol in tunnel mode.

This profile can be used to establish the necessary SA pairs. After that, C can request a token update to the AS, in order to establish an OSCOAP security context with RS, as specified in Section 2.2 of [I-D.seitz-ace-oscoap-profile].

Figure 9 overviews the involved secure communication channels. Logical links such as the SA pair shared between the Client and the Proxy are represented by dotted lines. IPsec traffic is depicted with double-dashed lines, and an example of the packets going through these links is represented with numbers, e.g. (1). The destination address included in the IP headers is also specified, e.g. "IP:P" indicates the Proxy’s address as destination address. The source address of the IP header is omitted, since all the IP packets have the Client’s address as source address.

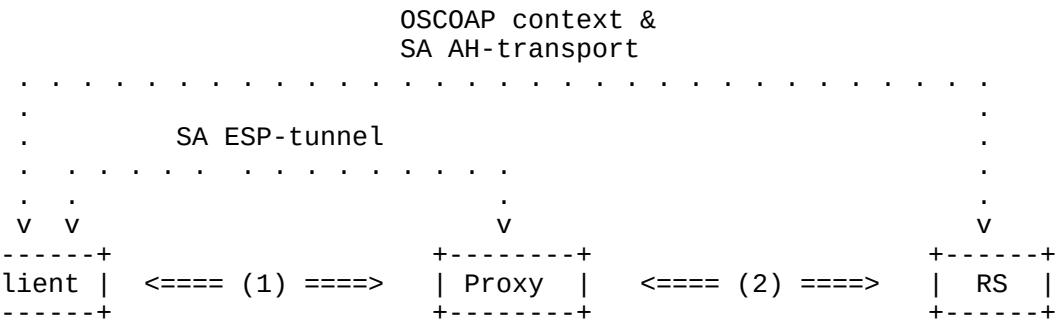

Figure 9: OSCOAP and IPsec - Scenario overview

Appendix B. SA Establishment with EDHOC

As discussed in Appendix A, securing communications between C and RS with both OSCOAP and IPsec makes it possible to fulfill a number of additional security requirements. An OSCOAP security context between C and RS can be established using Ephemeral Diffie-Hellman Over COSE (EDHOC) as defined in Appendix C.2 of [I-D.selander-ace-cose-ecdhe] and according to [I-D.seitz-ace-oscoap-profile]. This section proposes a method to establish also IPsec SA pairs by means of EDHOC. This makes it possible for constrained devices running the scenario described in Appendix A to rely solely on EDHOC for establishing both OSCOAP contexts and IPsec SA pairs, thus avoiding to include the implementation of IKEv2 as further key management protocol.

In particular, C and RS can refer to the SA Authenticated Establishment methods described in this specification, and then use EDHOC to finalize the SA pair, i.e. by deriving the encryption and authentication keys for the security protocols specified in the SA pair. This is possible thanks to IPsec's independence from specific key management protocols. In addition, the same security consideration discussed in [I-D.selander-ace-cose-ecdhe] hold.

The AS, C and RS refer to the same protocol shown in Figure 4, with the following changes.

B.1. Client to AS

The AS specifies the fields "alg", "SPI_SA_C" and "SPI_SA_RS" of the "ipsec" structure in the Access Token and in the Access Token Response, in addition to the pieces of information defined in

Section 3.2.2 or Section 3.2.3, in case the proof-of-possession key is symmetric or asymmetric, respectively.

The AS signals that EDHOC MUST be used, by setting the "kmp" field to "edhoc" in the Access Token and the Access Token Response. Then, C and RS MUST perform EDHOC as described in Section 4 or Section 5 of [I-D.selander-ace-cose-ecdhe], in case the proof-of-possession key is asymmetric or symmetric, respectively.

B.2. Client to RS

Figure 10 shows how EDHOC message_1 is sent through a POST Access Token Request to the /authz-info at the RS. The RS SHALL process the Access Token according to [I-D.ietf-ace-oauth-authz], and, if valid, continue with the EDHOC protocol as defined in Appendix C.1 of [I-D.selander-ace-cose-ecdhe]. Otherwise, RS aborts EDHOC and responds with an error code as specified in [I-D.ietf-ace-oauth-authz]. At the end of the EDHOC protocol, C and RS MUST derive an IPsec seed from the EDHOC shared secret. The seed is derived as specified in Section 3.2 of [I-D.selander-ace-cose-ecdhe], with other=exchange_hash, AlgorithmID="EDHOC IKE seed" and keyDataLength equal to the key length of the SKEYSEED secret defined in Section 2.14 of [RFC7296]. After that, the derived seed is written in the "seed" field of the "ipsec" structure, and accordingly used to derive IPsec key material as described in Section 2.1.

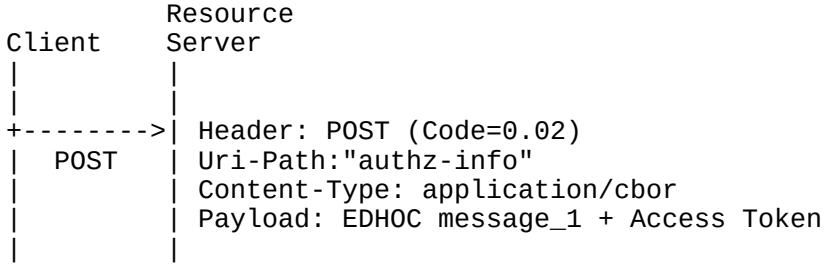

Figure 10: EDHOC used as Key Management Protocol

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Raw-Public-Key and Pre-Shared-Key as OAuth client credentials
draft-erdtman-ace-rpcc-01

Abstract

This document describes Transport Layer Security (TLS) authentication using Raw-Public-Key and Pre-Shared-Key as new mechanisms for OAuth client authentication. Although defined for TLS the mechanisms are equally applicable for DTLS.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF). Note that other groups may also distribute working documents as Internet-Drafts. The list of current Internet-Drafts is at http://datatracker.ietf.org/drafts/current/.

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This Internet-Draft will expire on February 17, 2018.

Table of Contents

1. Introduction

This document describes Transport Layer Security (TLS) authentication using Raw-Public-Key and Pre-Shared-Key as the mechanism for OAuth client authentication. Examples of endpoint requiering client authentication are token and introspection.

The OAuth 2.0 Authorization Framework [RFC6749] defines a shared secret method of client authentication but also allows for the definition and use of additional client authentication mechanisms when interacting with the authorization server's token endpoint. This document describes two additional mechanisms of client authentication utilizing Raw-Public-Key [RFC7250] and Pre-Shared-Key TLS [RFC4279], which provide better security characteristics than shared secrets.

1.1. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC 2119 [RFC2119].

2. Pre-Shared-Key for Client Authentication

The following section defines, as an extension of OAuth 2.0, Section 2.3 [RFC6749], using Pre-Shared-Key with TLS [RFC4279] to authenticate the client. This method is registered as 'tls_client_psk' in "OAuth Token Endpoint Authentication Methods" registry.

The (D)TLS handshake MUST be done according to [RFC4279], with the client indicating support for one or more Pre-Shared-Key cipher suites and authorization server selecting a Pre-Shared-Key cipher suite. In order to enable authorization server to select the correct pre-shared-key the client MUST send its client identifier in the psk-identity field of the ClientKeyExchange message. How the authorization server maps a client identifier to the pre-shared-key is out of scope for this specification.

Note that the client identity MUST be 2^16 bytes or shorter, in order to fit into the psk-identity field.

3. Raw-Public-Key for Client Authentication

The following section defines, as an extension of OAuth 2.0, Section 2.3 [RFC6749], the use of Raw-Public-Key with (D)TLS [RFC7250] to authenticate the client. This method is registered as'tls_client_rpk' in "OAuth Token Endpoint Authentication Methods" registry.

The (D)TLS handshake MUST be done according to [RFC7250], with the client indicating support for Raw-Public-Key certificates and the authorization server asking client send its Raw Public Key certificate. Since the client cannot send an explicit client identifier in the handshake, the authorization server MUST the derive a client identifier from RPK that the client uses.

Authorization servers MAY use the following method to map a Raw Public Key to a client identifier: The client identifier is generated from the Raw Public Key using the procedure specified in section 3 of [RFC6920]. The digest is calculated on the Raw Public Key only (not on the SubjectPublicKeyInfo used in the handshake). An example is shown in Figure 1.

```
Raw Public Key (Base64 encoded):
MFkwEwYHKoZIzj0CAQYIKoZIzj0DAQcDQgAEEtboxNKPgxEKV9JTNzy
tUvAbxEfkCTVB9kOzheF5wRAoOz2NKP+ln+XLVAQSp1D6jfo09tppvN
poQA1nnBNH6A==";

Encoding:
ni:///sha-256;xzLa24yOBeCkos3VFzD2gd83Urohr9TsXqY9nhdDN0
```

Figure 1: Example encoding of a raw public key in the Named Information URI Format

4. Acknowledgements

This document is highly inspired by [I-D.ietf-oauth-mtls] written by B. Campbell, J. Bradley, N. Sakimura and T. Lodderstedt.

5. IANA Considerations

5.1. Token Endpoint Authentication Method Registration

This specification requests registration of the following value in the IANA "OAuth Token Endpoint Authentication Methods" registry [IANA.OAuth.Parameters] established by [RFC7591].

5.1.1. Registry Contents

- Token Endpoint Authentication Method Name: "tls_client_rpk"
- Change Controller: IESG
- Specification Document(s): [[this specification]]

5.1.2. Registry Contents

- Token Endpoint Authentication Method Name: "tls_client_psk"
- Change Controller: IESG
- Specification Document(s): [[this specification]]

6. Security Considerations

TBD

7. References

7.1. Normative References

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[RFC6749] Hardt, D., Ed., "The OAuth 2.0 Authorization Framework", RFC 6749, DOI 10.17487/RFC6749, October 2012, <http://www.rfc-editor.org/info/rfc6749>.

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[RFC7250] Wouters, P., Ed., Tschofenig, H., Ed., Gilmore, J., Weiler, S., and T. Kivinen, "Using Raw Public Keys in Transport Layer Security (TLS) and Datagram Transport Layer Security (DTLS)", RFC 7250, DOI 10.17487/RFC7250, June 2014, <https://www.rfc-editor.org/info/rfc7250>.

[RFC7591] Richer, J., Ed., Jones, M., Bradley, J., Machulak, M., and P. Hunt, "OAuth 2.0 Dynamic Client Registration Protocol", RFC 7591, DOI 10.17487/RFC7591, July 2015, <https://www.rfc-editor.org/info/rfc7591>.

7.2. Informative References

[I-D.ietf-oauth-mtls] Campbell, B., Bradley, J., Sakimura, N., and T. Lodderstedt, "Mutual TLS Profile for OAuth 2.0", draft-ietf-oauth-mtls-03 (work in progress), July 2017.

[IANA.OAuth.Parameters] IANA, "OAuth Parameters", <http://www.iana.org/assignments/oauth-parameters>.

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CBOR Web Token (CWT)
draft-ietf-ace-cbor-web-token-08

Abstract

CBOR Web Token (CWT) is a compact means of representing claims to be transferred between two parties. The claims in a CWT are encoded in the Concise Binary Object Representation (CBOR) and CBOR Object Signing and Encryption (COSE) is used for added application layer security protection. A claim is a piece of information asserted about a subject and is represented as a name/value pair consisting of a claim name and a claim value. CWT is derived from JSON Web Token (JWT), but uses CBOR rather than JSON.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF). Note that other groups may also distribute working documents as Internet-Drafts. The list of current Internet-Drafts is at http://datatracker.ietf.org/drafts/current/.

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This Internet-Draft will expire on February 17, 2018.

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Table of Contents

1. Introduction

The JSON Web Token (JWT) [RFC7519] is a standardized security token format that has found use in OAuth 2.0 and OpenID Connect deployments, among other applications. JWT uses JSON Web Signature (JWS) [RFC7515] and JSON Web Encryption (JWE) [RFC7516] to secure the contents of the JWT, which is a set of claims represented in JSON. The use of JSON for encoding information is popular for Web and native applications, but it is considered inefficient for some Internet of Things (IoT) systems that use low power radio technologies.

An alternative encoding of claims is defined in this document. Instead of using JSON, as provided by JWTs, this specification uses CBOR [RFC7049] and calls this new structure "CBOR Web Token (CWT)", which is a compact means of representing secured claims to be transferred between two parties. CWT is closely related to JWT. It references the JWT claims and both its name and pronunciation are derived from JWT. To protect the claims contained in CWTs, the CBOR Object Signing and Encryption (COSE) [RFC8152] specification is used.

The suggested pronunciation of CWT is the same as the English word "cot".

1.1. CBOR Related Terminology

In JSON, maps are called objects and only have one kind of map key: a string. CBOR uses strings, negative integers, and unsigned integers as map keys. The integers are used for compactness of encoding and easy comparison. The inclusion of strings allows for an additional range of short encoded values to be used.

2. Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and

"OPTIONAL" in this document are to be interpreted as described in "Key words for use in RFCs to Indicate Requirement Levels" [RFC2119].

This document reuses terminology from JWT [RFC7519] and COSE [RFC8152].

StringOrURI
: The "StringOrURI" term has the same meaning, syntax, and processing rules as the "StringOrUri" term defined in Section 2 of JWT [RFC7519], except that it uses a CBOR text string instead of a JSON string value.

NumericDate
: The "NumericDate" term has the same meaning, syntax, and processing rules as the "NumericDate" term defined in Section 2 of JWT [RFC7519], except that the CBOR numeric date representation (from Section 2.4.1 of [RFC7049]) is used. The encoding is modified so that the leading tag 1 (epoch-based date/time) MUST be omitted.

Claim Name
: The human-readable name used to identify a claim.

Claim Key
: The CBOR map key used to identify a claim.

Claim Value
: The CBOR map value representing the value of the claim.

CWT Claims Set
: The CBOR map that contains the claims conveyed by the CWT.

3. Claims

The set of claims that a CWT must contain to be considered valid is context dependent and is outside the scope of this specification. Specific applications of CWTs will require implementations to understand and process some claims in particular ways. However, in the absence of such requirements, all claims that are not understood by implementations MUST be ignored.

To keep CWTs as small as possible, the Claim Keys are represented using integers or text strings. Section 4 summarizes all keys used to identify the claims defined in this document.

3.1. Registered Claims

None of the claims defined below are intended to be mandatory to use or implement. They rather provide a starting point for a set of useful, interoperable claims. Applications using CWTs should define which specific claims they use and when they are required or optional.

3.1.1. iss (Issuer) Claim

The "iss" (issuer) claim has the same meaning, syntax, and processing rules as the "iss" claim defined in Section 4.1.1 of JWT [RFC7519], except that the value is of type StringOrURI. The Claim Key 1 is used to identify this claim.

3.1.2. sub (Subject) Claim

The "sub" (subject) claim has the same meaning, syntax, and processing rules as the "sub" claim defined in Section 4.1.2 of JWT [RFC7519], except that the value is of type StringOrURI. The Claim Key 2 is used to identify this claim.

3.1.3. aud (Audience) Claim

The "aud" (audience) claim has the same meaning, syntax, and processing rules as the "aud" claim defined in Section 4.1.3 of JWT [RFC7519], except that the value is of type StringOrURI. The Claim Key 3 is used to identify this claim.

3.1.4. exp (Expiration Time) Claim

The "exp" (expiration time) claim has the same meaning, syntax, and processing rules as the "exp" claim defined in Section 4.1.4 of JWT [RFC7519], except that the value is of type NumericDate. The Claim Key 4 is used to identify this claim.

3.1.5. nbf (Not Before) Claim

The "nbf" (not before) claim has the same meaning, syntax, and processing rules as the "nbf" claim defined in Section 4.1.5 of JWT [RFC7519], except that the value is of type NumericDate. The Claim Key 5 is used to identify this claim.

3.1.6. iat (Issued At) Claim

The "iat" (issued at) claim has the same meaning, syntax, and processing rules as the "iat" claim defined in Section 4.1.6 of JWT

[RFC7519], except that the value is of type NumericDate. The Claim Key 6 is used to identify this claim.

3.1.7. cti (CWT ID) Claim

The "cti" (CWT ID) claim has the same meaning, syntax, and processing rules as the "jti" claim defined in Section 4.1.7 of JWT [RFC7519], except that the value is of type binary string. The Claim Key 7 is used to identify this claim.

4. Summary of the claim names, keys, and value types

Name	Key	Value type
iss	1	text string
sub	2	text string
aud	3	text string
exp	4	integer or floating-point number
nbf	5	integer or floating-point number
iat	6	integer or floating-point number
cti	7	binary string

Figure 1: Summary of the claim names, keys, and value types

5. CBOR Tags and Claim Values

The claim values defined in this specification MUST NOT be prefixed with any CBOR tag. For instance, while CBOR tag 1 (epoch-based date/time) could logically be prefixed to values of the "exp", "nbf", and "iat" claims, this is unnecessary, since the representation of the claim values is already specified by the claim definitions. Tagging claim values would only take up extra space without adding information. However, this does not prohibit future claim definitions from requiring the use of CBOR tags for those specific claims.

6. CWT CBOR Tag

How to determine that a CBOR data structure is a CWT is application-dependent. In some cases, this information is known from the application context, such as from the position of the CWT in a data structure at which the value must be a CWT. One method of indicating that a CBOR object is a CWT is the use of the "application/cwt" content type by a transport protocol.

This section defines the CWT CBOR tag as another means for applications to declare that a CBOR data structure is a CWT. Its use is optional, and is intended for use in cases in which this information would not otherwise be known.

If present, the CWT tag MUST prefix a tagged object using one of the COSE CBOR tags. In this example, the COSE_Mac0 tag is used. The actual COSE_Mac0 object has been excluded from this example.

```
/ CWT CBOR tag / 61(
  / COSE_Mac0 CBOR tag / 17(
    / COSE_Mac0 object /
  )
)
```

Figure 2: Example of a CWT tag usage

7. Creating and Validating CWTs

7.1. Creating a CWT

To create a CWT, the following steps are performed. The order of the steps is not significant in cases where there are no dependencies between the inputs and outputs of the steps.

1. Create a CWT Claims Set containing the desired claims.

2. Let the Message be the binary representation of the CWT Claims Set.

3. Create a COSE Header containing the desired set of Header Parameters. The COSE Header MUST be valid per the [RFC8152] specification.

4. Depending upon whether the CWT is signed, MACed, or encrypted, there are three cases:

   * If the CWT is signed, create a COSE_Sign/COSE_Sign1 object using the Message as the COSE_Sign/COSE_Sign1 Payload; all steps specified in [RFC8152] for creating a COSE_Sign/ COSE_Sign1 object MUST be followed.

   * Else, if the CWT is MACed, create a COSE_Mac/COSE_Mac0 object using the Message as the COSE_Mac/COSE_Mac0 Payload; all steps specified in [RFC8152] for creating a COSE_Mac/COSE_Mac0 object MUST be followed.

   * Else, if the CWT is a COSE_Encrypt/COSE_Encrypt0 object, create a COSE_Encrypt/COSE_Encrypt0 using the Message as the plaintext for the COSE_Encrypt/COSE_Encrypt0 object; all steps specified in [RFC8152] for creating a COSE_Encrypt/ COSE_Encrypt0 object MUST be followed.

5. If a nested signing, MACing, or encryption operation will be performed, let the Message be the COSE_Sign/COSE_Sign1, COSE_Mac/ COSE_Mac0, or COSE_Encrypt/COSE_Encrypt0, add the matching COSE CBOR tag, and return to Step 3.

6. If needed by the application, add the appropriate COSE CBOR tag to the COSE object to indicate the type of the COSE object. If needed by the application, add the CWT CBOR tag to indicate that the COSE object is a CWT.

7.2. Validating a CWT

When validating a CWT, the following steps are performed. The order of the steps is not significant in cases where there are no dependencies between the inputs and outputs of the steps. If any of the listed steps fail, then the CWT MUST be rejected -- that is, treated by the application as invalid input.

1. Verify that the CWT is a valid CBOR object.

2. If the object begins with the CWT CBOR tag, remove it and verify that one of the COSE CBOR tags follows it.

3. If the object is tagged with one of the COSE CBOR tags, remove it and use it to determine the type of the CWT, COSE_Sign/ COSE_Sign1, COSE_Mac/COSE_Mac0, or COSE_Encrypt/COSE_Encrypt0. If the object does not have a COSE CBOR tag, the COSE message type is determined from the application context.

4. Verify that the resulting COSE Header includes only parameters and values whose syntax and semantics are both understood and supported or that are specified as being ignored when not understood.

5. Depending upon whether the CWT is a signed, MACed, or encrypted, there are three cases:

   * If the CWT is a COSE_Sign/COSE_Sign1, follow the steps specified in [RFC8152] Section 4 (Signing Objects) for validating a COSE_Sign/COSE_Sign1 object. Let the Message be the COSE_Sign/COSE_Sign1 payload.

   * Else, if the CWT is a COSE_Mac/COSE_Mac0, follow the steps specified in [RFC8152] Section 6 (MAC Objects) for validating a COSE_Mac/COSE_Mac0 object. Let the Message be the COSE_Mac/COSE_Mac0 payload.

   * Else, if the CWT is a COSE_Encrypt/COSE_Encrypt0 object, follow the steps specified in [RFC8152] Section 5 (Encryption Objects) for validating a COSE_Encrypt/COSE_Encrypt0 object. Let the Message be the resulting plaintext.

6. If the Message begins with a COSE CBOR tag, then the Message is a CWT that was the subject of nested signing, MACing, or encryption operations. In this case, return to Step 1, using the Message as the CWT.

7. Verify that the Message is a valid CBOR map; let the CWT Claims Set be this CBOR map.

8. Security Considerations

The security of the CWT relies upon on the protections offered by COSE. Unless the claims in a CWT are protected, an adversary can modify, add, or remove claims.

Since the claims conveyed in a CWT may be used to make authorization decisions, it is not only important to protect the CWT in transit but also to ensure that the recipient can authenticate the party that assembled the claims and created the CWT. Without trust of the recipient in the party that created the CWT, no sensible authorization decision can be made. Furthermore, the creator of the CWT needs to carefully evaluate each claim value prior to including it in the CWT so that the recipient can be assured of the validity of the information provided.

While syntactically, the signing and encryption operations for Nested CWTs may be applied in any order, if both signing and encryption are necessary, normally producers should sign the message and then encrypt the result (thus encrypting the signature). This prevents attacks in which the signature is stripped, leaving just an encrypted message, as well as providing privacy for the signer. Furthermore, signatures over encrypted text are not considered valid in many jurisdictions.

9. IANA Considerations

9.1. CBOR Web Token (CWT) Claims Registry

This section establishes the IANA "CBOR Web Token (CWT) Claims" registry.

Values are registered on a Specification Required [RFC5226] basis after a three-week review period on the cwt-reg-review@ietf.org mailing list, on the advice of one or more Designated Experts. However, to allow for the allocation of values prior to publication, the Designated Experts may approve registration once they are satisfied that such a specification will be published. [[Note to the RFC Editor: The name of the mailing list should be determined in consultation with the IESG and IANA. Suggested name: cwt-reg-review@ietf.org.]]

Registration requests sent to the mailing list for review should use an appropriate subject (e.g., "Request to register claim: example"). Registration requests that are undetermined for a period longer than 21 days can be brought to the IESG's attention (using the iesg@ietf.org mailing list) for resolution.

Criteria that should be applied by the Designated Experts includes determining whether the proposed registration duplicates existing functionality, whether it is likely to be of general applicability or whether it is useful only for a single application, and whether the registration description is clear.

It is suggested that multiple Designated Experts be appointed who are able to represent the perspectives of different applications using this specification in order to enable broadly informed review of registration decisions. In cases where a registration decision could be perceived as creating a conflict of interest for a particular Expert, that Expert should defer to the judgment of the other Experts.

9.1.1. Registration Template

Claim Name:
: The human-readable name requested (e.g., "iss").

Claim Description:
: Brief description of the claim (e.g., "Issuer").

JWT Claim Name:
: Claim Name of the equivalent JWT claim, as registered in [IANA.JWT.Claims]. CWT claims should normally have a corresponding JWT claim. If a corresponding JWT claim would not

make sense, the Designated Experts can choose to accept registrations for which the JWT Claim Name is listed as "N/A".

Claim Key:
CBOR map key for the claim. Integer values between -256 and 255 and strings of length 1 are designated as Standards Track Document required. Integer values from -65536 to 65535 and strings of length 2 are designated as Specification Required. Integer values of greater than 65535 and strings of length greater than 2 are designated as expert review. Integer values less than -65536 are marked as private use.

Claim Value Type(s):
CBOR types that can be used for the claim value.

Change Controller:
For Standards Track RFCs, list the "IESG". For others, give the name of the responsible party. Other details (e.g., postal address, email address, home page URI) may also be included.

Specification Document(s):
Reference to the document or documents that specify the parameter, preferably including URIs that can be used to retrieve copies of the documents. An indication of the relevant sections may also be included but is not required.

9.1.2. Initial Registry Contents

- Claim Name: (RESERVED)
- Claim Description: This registration reserves the key value 0.
- JWT Claim Name: N/A
- Claim Key: 0
- Claim Value Type(s): N/A
- Change Controller: IESG
- Specification Document(s): [[this specification]]

- Claim Name: "iss"
- Claim Description: Issuer
- JWT Claim Name: "iss"
- Claim Key: 1
- Claim Value Type(s): text string
- Change Controller: IESG
- Specification Document(s): Section 3.1.1 of [[this specification]]

- Claim Name: "sub"
- Claim Description: Subject
- JWT Claim Name: "sub"

- Claim Key: 2
- Claim Value Type(s): text string
- Change Controller: IESG
- Specification Document(s): Section 3.1.2 of [[this specification]]

- Claim Name: "aud"
- Claim Description: Audience
- JWT Claim Name: "aud"
- Claim Key: 3
- Claim Value Type(s): text string
- Change Controller: IESG
- Specification Document(s): Section 3.1.3 of [[this specification]]

- Claim Name: "exp"
- Claim Description: Expiration Time
- JWT Claim Name: "exp"
- Claim Key: 4
- Claim Value Type(s): integer or floating-point number
- Change Controller: IESG
- Specification Document(s): Section 3.1.4 of [[this specification]]

- Claim Name: "nbf"
- Claim Description: Not Before
- JWT Claim Name: "nbf"
- Claim Key: 5
- Claim Value Type(s): integer or floating-point number
- Change Controller: IESG
- Specification Document(s): Section 3.1.5 of [[this specification]]

- Claim Name: "iat"
- Claim Description: Issued At
- JWT Claim Name: "iat"
- Claim Key: 6
- Claim Value Type(s): integer or floating-point number
- Change Controller: IESG
- Specification Document(s): Section 3.1.6 of [[this specification]]

- Claim Name: "cti"
- Claim Description: CWT ID
- JWT Claim Name: "jti"
- Claim Key: 7
- Claim Value Type(s): binary string
- Change Controller: IESG

- Specification Document(s): Section 3.1.7 of [[this specification]]

9.2. Media Type Registration

This section registers the "application/cwt" media type in the "Media Types" registry [IANA.MediaTypes] in the manner described in RFC 6838 [RFC6838], which can be used to indicate that the content is a CWT.

9.2.1. Registry Contents

- Type name: application
- Subtype name: cwt
- Required parameters: N/A
- Optional parameters: N/A
- Encoding considerations: binary
- Security considerations: See the Security Considerations section of [[this specification]]
- Interoperability considerations: N/A
- Published specification: [[this specification]]
- Applications that use this media type: IoT applications sending security tokens over HTTP(S) and other transports.
- Fragment identifier considerations: N/A
- Additional information:

  Magic number(s): N/A
  File extension(s): N/A
  Macintosh file type code(s): N/A

- Person & email address to contact for further information: IESG, iesg@ietf.org
- Intended usage: COMMON
- Restrictions on usage: none
- Author: Michael B. Jones, mbj@microsoft.com
- Change controller: IESG
- Provisional registration? No

9.3. CoAP Content-Formats Registration

This section registers the CoAP Content-Format ID for the "application/cwt" media type in the "CoAP Content-Formats" registry [IANA.CoAP.Content-Formats].

9.3.1. Registry Contents

- Media Type: application/cwt
- Encoding: -
- Id: TBD (maybe 61)

- Reference: [[this specification]]

9.4. CBOR Tag registration

This section registers the CWT CBOR tag in the "CBOR Tags" registry [IANA.CBOR.Tags].

9.4.1. Registry Contents

- CBOR Tag: TBD (maybe 61 to use the same value as the Content-Format)
- Data Item: CBOR Web Token (CWT)
- Semantics: CBOR Web Token (CWT), as defined in [[this specification]]
- Reference: [[this specification]]
- Point of Contact: Michael B. Jones, mbj@microsoft.com

10. References

10.1. Normative References

[IANA.CBOR.Tags]
IANA, "Concise Binary Object Representation (CBOR) Tags", <http://www.iana.org/assignments/cbor-tags/cbor-tags.xhtml>.

[IANA.CoAP.Content-Formats]
IANA, "CoAP Content-Formats", <http://www.iana.org/assignments/core-parameters/core-parameters.xhtml#content-formats>.

[IANA.MediaTypes]
IANA, "Media Types", <http://www.iana.org/assignments/media-types>.

[RFC2119] Bradner, S., "Key words for use in RFCs to Indicate Requirement Levels", BCP 14, RFC 2119, DOI 10.17487/RFC2119, March 1997, <http://www.rfc-editor.org/info/rfc2119>.

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[RFC7519] Jones, M., Bradley, J., and N. Sakimura, "JSON Web Token (JWT)", RFC 7519, DOI 10.17487/RFC7519, May 2015, <http://www.rfc-editor.org/info/rfc7519>.

[RFC8152] Schaad, J., "CBOR Object Signing and Encryption (COSE)", RFC 8152, DOI 10.17487/RFC8152, July 2017, <http://www.rfc-editor.org/info/rfc8152>.

10.2. Informative References

[IANA.JWT.Claims] IANA, "JSON Web Token Claims", <http://www.iana.org/assignments/jwt>.

[RFC5226] Narten, T. and H. Alvestrand, "Guidelines for Writing an IANA Considerations Section in RFCs", RFC 5226, DOI 10.17487/RFC5226, May 2008, <http://www.rfc-editor.org/info/rfc5226>.

[RFC6838] Freed, N., Klensin, J., and T. Hansen, "Media Type Specifications and Registration Procedures", BCP 13, RFC 6838, DOI 10.17487/RFC6838, January 2013, <http://www.rfc-editor.org/info/rfc6838>.

[RFC7515] Jones, M., Bradley, J., and N. Sakimura, "JSON Web Signature (JWS)", RFC 7515, DOI 10.17487/RFC7515, May 2015, <http://www.rfc-editor.org/info/rfc7515>.

[RFC7516] Jones, M. and J. Hildebrand, "JSON Web Encryption (JWE)", RFC 7516, DOI 10.17487/RFC7516, May 2015, <http://www.rfc-editor.org/info/rfc7516>.

Appendix A. Examples

This appendix includes a set of CWT examples that show how the CWT Claims Set can be protected. There are examples that are signed, MACed, encrypted, and that use nested signing and encryption. To make the examples easier to read, they are presented both as hex strings and in the extended CBOR diagnostic notation described in Section 6 of [RFC7049].

Where a byte string is to carry an embedded CBOR-encoded item, the diagnostic notation for this CBOR data item can be enclosed in '<<' and '>>' to notate the byte string resulting from encoding the data item, e.g., h'63666F6F' translates to <<"foo">>.

A.1. Example CWT Claims Set

The CWT Claims Set used for the different examples displays usage of all the defined claims. For signed and MACed examples, the CWT Claims Set is the CBOR encoding as a binary string.

```
a70175636f61703a2f2f61732e6578616d706c652e636f6d02656572696b7703
7818636f61703a2f2f6c696768742e6578616d706c652e636f6d041a5612aeb0
051a5610d9f0061a5610d9f007420b71
```

Figure 3: Example CWT Claims Set as hex string

```
{
  / iss / 1: "coap://as.example.com",
  / sub / 2: "erikw",
  / aud / 3: "coap://light.example.com",
  / exp / 4: 1444064944,
  / nbf / 5: 1443944944,
  / iat / 6: 1443944944,
  / cti / 7: h'0b71'
}
```

Figure 4: Example CWT Claims Set in CBOR diagnostic notation

A.2. Example keys

This section contains the keys used to sign, MAC, and encrypt the messages in this appendix. Line breaks are for display purposes only.

A.2.1. 128-bit Symmetric Key as Hex Encoded String

```
231f4c4d4d3051fdc2ec0a3851d5b383
```

A.2.2. 256-bit Symmetric Key as Hex Encoded String

```
403697de87af64611c1d32a05dab0fe1fcb715a86ab435f1ec99192d79569388
```

A.2.3. ECDSA P-256 256-bit COSE Key

```
a622582060f7f1a780d8a783bfb7a2dd6b2796e8128dbbcef9d3d168db952997
1a36e7b92358206c1382765aec5358f117733d281c1c7bdc39884d04a45a1e6c
67c858bc206c19032601022158201432 9cce7868e416927599cf65a34f3ce2f
fda55a7eca69ed8919a394d42f0f2001
```

Figure 5: ECDSA 256-bit COSE Key as hex string

```
{
  / d /   -4: h'6c1382765aec5358f117733d281c1c7bdc39884d04a45a1e
              6c67c858bc206c19',
  / y /   -3: h'60f7f1a780d8a783bfb7a2dd6b2796e8128dbbcef9d3d168
              db9529971a36e7b9',
  / x /   -2: h'143329cce7868e416927599cf65a34f3ce2ffda55a7eca69
              ed8919a394d42f0f',
  / crv / -1: 1 / P-256 /,
  / kty /  1: 2 / EC2 /,
  / alg /  3: -7 / ECDSA 256 /
}
```

Figure 6: ECDSA 256-bit COSE Key in CBOR diagnostic notation

A.3. Example Signed CWT

This section shows a signed CWT with a single recipient and a full CWT Claims Set.

The signature is generated using the private key listed in Appendix A.2.3 and it can be validated using the public key from Appendix A.2.3. Line breaks are for display purposes only.

```
d28443a10126a05850a70175636f61703a2f2f61732e6578616d706c652e636f6
d02656572696b77037818636f61703a2f2f6c696768742e6578616d706c652e63
6f6d041a5612aeb0051a5610d9f0061a5610d9f007420b7158405427c1ff28d23
fbad1f29c4c7c6a555e601d6fa29f9179bc3d7438bacaca5acd08c8d4d4f96131
680c429a01f85951ecee743a52b9b63632c57209120e1c9e30
```

Figure 7: Signed CWT as hex string

```
18(
  [
    / protected / << {
      / alg / 1: -7 / ECDSA 256 /
    } >>,
    / unprotected / {},
    / payload / << {
      / iss / 1: "coap://as.example.com",
      / sub / 2: "erikw",
      / aud / 3: "coap://light.example.com",
      / exp / 4: 1444064944,
      / nbf / 5: 1443944944,
      / iat / 6: 1443944944,
      / cti / 7: h'0b71'
    } >>,
    / signature / h'5427c1ff28d23fbad1f29c4c7c6a555e601d6fa29f
                    9179bc3d7438bacaca5acd08c8d4d4f96131680c42
                    9a01f85951ecee743a52b9b63632c57209120e1c9e
                    30'
  ]
)
```

Figure 8: Signed CWT in CBOR diagnostic notation

A.4. Example MACed CWT

This section shows a MACed CWT with a single recipient, a full CWT Claims Set, and a CWT tag.

The MAC is generated using the 256-bit symmetric key from Appendix A.2.2 with a 64-bit truncation. Line breaks are for display purposes only.

```
d83dd18443a10104a05850a70175636f61703a2f2f61732e6578616d706c652e
636f6d02656572696b77037818636f61703a2f2f6c696768742e6578616d706c
652e636f6d041a5612aeb0051a5610d9f0061a5610d9f007420b7148093101ef
6d789200
```

Figure 9: MACed CWT with CWT tag as hex string

```
61(
  17(
    [
      / protected / << {
        / alg / 1: 4 / HMAC-256-64 /
      } >>,
      / unprotected / {},
      / payload / << {
        / iss / 1: "coap://as.example.com",
        / sub / 2: "erikw",
        / aud / 3: "coap://light.example.com",
        / exp / 4: 1444064944,
        / nbf / 5: 1443944944,
        / iat / 6: 1443944944,
        / cti / 7: h'0b71'
      } >>,
      / tag / h'093101ef6d789200'
    ]
  )
)
```

Figure 10: MACed CWT with CWT tag in CBOR diagnostic notation

A.5. Example Encrypted CWT

This section shows an encrypted CWT with a single recipient and a full CWT Claims Set.

The encryption is done with AES-CCM mode using the 128-bit symmetric key from Appendix A.2.1 with a 64-bit tag and 13-byte nonce, i.e., COSE AES-CCM-16-64-128. Line breaks are for display purposes only.

```
d08343a1010aa1054d99a0d7846e762c49ffe8a63e0b5858b918a11fd81e438b
7f973d9e2e119bcb22424ba0f38a80f27562f400ee1d0d6c0fdb559c02421fd3
84fc2ebe22d7071378b0ea7428fff157444d45f7e6afcda1aae5f6495830c586
27087fc5b4974f319a8707a635dd643b
```

Figure 11: Encrypted CWT as hex string

```
16(
  [
    / protected / << {
      / alg / 1: 10 / AES-CCM-16-64-128 /
    } >>,
    / unprotected / {
      / iv / 5: h'99a0d7846e762c49ffe8a63e0b'
    },
    / ciphertext / h'b918a11fd81e438b7f973d9e2e119bcb22424ba0f38
                     a80f27562f400ee1d0d6c0fdb559c02421fd384fc2e
                     be22d7071378b0ea7428fff157444d45f7e6afcda1a
                     ae5f6495830c58627087fc5b4974f319a8707a635dd
                     643b'
  ]
)
```

Figure 12: Encrypted CWT in CBOR diagnostic notation

A.6. Example Nested CWT

This section shows a Nested CWT, signed and then encrypted, with a single recipient and a full CWT Claims Set.

The signature is generated using the private ECDSA key from Appendix A.2.3 and it can be validated using the public ECDSA parts from Appendix A.2.3. The encryption is done with AES-CCM mode using the 128-bit symmetric key from Appendix A.2.1 with a 64-bit tag and 13-byte nonce, i.e., COSE AES-CCM-16-64-128. The content type is set to CWT to indicate that there are multiple layers of COSE protection before finding the CWT Claims Set. The decrypted ciphertext will be a COSE_sign1 structure. In this example, it is the same one as in Appendix A.3, i.e., a Signed CWT Claims Set. Note that there is no limitation to the number of layers; this is an example with two layers. Line breaks are for display purposes only.

```
d08343a1010aa1054d86bbd41cc32604396324b7f38058a372439fbff538aa7b
601ebfb29454050a3c99fd13b27216d084556496c7355c4bb462510f8e0e8479
dbe08722d620e96bcb7764d75140d96220f062679b46b897e7abe0c325dc2c96
d8bb2c8334e3b92a42c0078983e753c054e647ad5387ed149f802f52b5a95ebf
5f153c4fd64854ab7531e082b7f22721f939d257c94f8bc248e1d9cf04f9dd4e
5de7ab62df37842fabec230a657d4abf7162bc786345ebb8eb3af0
```

Figure 13: Signed and Encrypted CWT as hex string

```
16(
  [
    / protected / << {
      / alg / 1: 10 / AES-CCM-16-64-128 /
    } >>,
    / unprotected / {
      / iv / 5: h'86bbd41cc32604396324b7f380'
    },
    / ciphertext / h'72439fbff538aa7b601ebfb29454050a3c99fd13b27
                     216d084556496c7355c4bb462510f8e0e8479dbe087
                     22d620e96bcb7764d75140d96220f062679b46b897e
                     7abe0c325dc2c96d8bb2c8334e3b92a42c0078983e7
                     53c054e647ad5387ed149f802f52b5a95ebf5f153c4
                     fd64854ab7531e082b7f22721f939d257c94f8bc248
                     e1d9cf04f9dd4e5de7ab62df37842fabec230a657d4
                     abf7162bc786345ebb8eb3af0'
  ]
)
```

Figure 14: Signed and Encrypted CWT in CBOR diagnostic notation

A.7. Example MACed CWT with a floating-point value

This section shows a MACed CWT with a single recipient and a simple CWT Claims Set. The CWT Claims Set with a floating-point 'iat' value.

The MAC is generated using the 256-bit symmetric key from Appendix A.2.2 with a 64-bit truncation. Line breaks are for display purposes only.

```
d18443a10104a04ba106fb41d584367c20000048b8816f34c0542892
```

Figure 15: MACed CWT with a floating-point value as hex string

```
17(
  [
    / protected / << {
      / alg / 1: 4 / HMAC-256-64 /
    } >>,
    / unprotected / {},
    / payload / << {
      / iat / 6: 1443944944.5
    } >>,
    / tag / h'b8816f34c0542892'
  ]
)
```

Figure 16: MACed CWT with a floating-point value in CBOR diagnostic notation

Appendix B. Acknowledgements

This specification is based on JSON Web Token (JWT) [RFC7519], the authors of which also include Nat Sakimura and John Bradley. It also incorporates suggestions made by many people, notably Carsten Bormann, Jim Schaad, Ludwig Seitz, and Goeran Selander.

Appendix C. Document History

[[to be removed by the RFC Editor before publication as an RFC]]

-08

- Updated the diagnostic notation for embedded objects in the examples, addressing feedback by Carsten Bormann.

-07

- Updated examples for signing and encryption. Signatures are now deterministic as recommended by COSE specification.

-06

- Addressed review comments by Carsten Bormann and Jim Schaad. All changes were editorial in nature.

-05

- Addressed working group last call comments with the following changes:

- Say that CWT is derived from JWT, rather than CWT is a profile of JWT.
- Used CBOR type names in descriptions, rather than major/minor type numbers.
- Clarified the NumericDate and StringOrURI descriptions.
- Changed to allow CWT claim names to use values of any legal CBOR map key type.
- Changed to use the CWT tag to identify nested CWTs instead of the CWT content type.
- Added an example using a floating-point date value.
- Acknowledged reviewers.

-04

- Specified that the use of CBOR tags to prefix any of the claim values defined in this specification is NOT RECOMMENDED.

-03

- Reworked the examples to include signed, MACed, encrypted, and nested CWTs.
- Defined the CWT CBOR tag and explained its usage.

-02

- Added IANA registration for the application/cwt media type.
- Clarified the nested CWT language.
- Corrected nits identified by Ludwig Seitz.

-01

- Added IANA registration for CWT Claims.
- Added IANA registration for the application/cwt CoAP content-format type.
- Added Samuel Erdtman as an editor.
- Changed Erik’s e-mail address.

-00

- Created the initial working group version based on draft-wahlstroem-ace-cbor-web-token-00.

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Datagram Transport Layer Security (DTLS) Profile for Authentication and Authorization for Constrained Environments (ACE) draft-ietf-ace-dtls-authorize-01

Abstract

This specification defines a profile for delegating client authentication and authorization in a constrained environment by establishing a Datagram Transport Layer Security (DTLS) channel between resource-constrained nodes. The protocol relies on DTLS for communication security between entities in a constrained network. A resource-constrained node can use this protocol to delegate management of authorization information to a trusted host with less severe limitations regarding processing power and memory.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF). Note that other groups may also distribute working documents as Internet-Drafts. The list of current Internet-Drafts is at http://datatracker.ietf.org/drafts/current/.

Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as "work in progress."

This Internet-Draft will expire on January 4, 2018.

Table of Contents

1. Introduction

This specification defines a profile of the ACE framework [I-D.ietf-ace-oauth-authz]. In this profile, a client and a resource server use CoAP [RFC7252] over DTLS [RFC6347] to communicate. The client uses an access token, bound to a key (the proof-of-possession key) to authorize its access to the resource server. DTLS provides communication security, proof of possession, and server authentication. Optionally the client and the resource server may also use CoAP over DTLS to communicate with the authorization server. This specification supports the DTLS handshake with Raw Public Keys (RPK) [RFC7250] and the DTLS PSK handshake [RFC4279].

The DTLS RPK handshake [RFC7250] requires client authentication to provide proof-of-possession for the key tied to the access token. Here the access token needs to be transferred to the resource server before the handshake is initiated, as described in section 8.1 of draft-ietf-ace-oauth-authz. [1]

The DTLS PSK handshake [RFC4279] provides the proof-of-possession for the key tied to the access token. Furthermore the psk_identity parameter in the DTLS PSK handshake is used to transfer the access token from the client to the resource server.

Note: While the scope of this draft is on client and resource server

communicating using CoAP over DTLS, it is expected that it applies also to CoAP over TLS, possibly with minor modifications. However, that is out of scope for this version of the draft.

1.1. Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC 2119 [RFC2119].

Readers are expected to be familiar with the terms and concepts described in [I-D.ietf-ace-oauth-authz].

2. Protocol Overview

The CoAP-DTLS profile for ACE specifies the transfer of authentication and, if necessary, authorization information between C and RS during setup of a DTLS session for CoAP messaging. It also specifies how a Client can use CoAP over DTLS to retrieve an Access Token from AS for a protected resource hosted on RS.

This profile requires a Client (C) to retrieve an Access Token for the resource(s) it wants to access on a Resource Server (RS) as specified in [I-D.ietf-ace-oauth-authz]. Figure 1 shows the typical message flow in this scenario (messages in square brackets are optional):

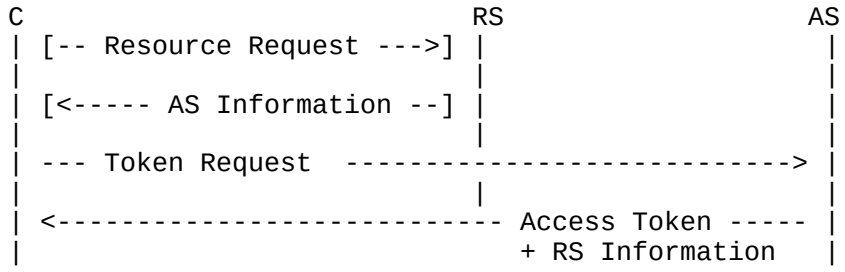

Figure 1: Retrieving an Access Token

To determine the AS in charge of a resource hosted at the RS, C MAY send an initial Unauthorized Resource Request message to RS. RS then denies the request and sends the address of its AS back to C.

Instead of the initial Unauthorized Resource Request message, C MAY look up the desired resource in a resource directory (cf. [I-D.ietf-core-resource-directory]).

Once C knows AS's address, it can send an Access Token request to the /token endpoint at the AS as specified in [I-D.ietf-ace-oauth-authz]. If C wants to use the CoAP RawPublicKey mode as described in Section 9 of RFC 7252 [2] it MUST provide a key or key identifier within a "cnf" object in the token request. If AS decides that the request is to be authorized it generates an access token response for C containing a "profile" parameter with the value "coap_dtls" to indicate that this profile MUST be used for communication between C and RS. Is also adds a "cnf" parameter with additional data for the establishment of a secure DTLS channel between C and RS. The semantics of the 'cnf' parameter depend on the type of key used between C and RS, see Section 3 and Section 4.

The Access Token returned by AS then can be used by C to establish a new DTLS session with RS. When C intends to use asymmetric cryptography in the DTLS handshake with RS, C MUST upload the Access Token to the "/authz-info" resource on RS before starting the DTLS handshake, as described in section 8.1 of draft-ietf-ace-oauth-authz [3]. If only symmetric cryptography is used between C and RS, the Access Token MAY instead be transferred in the DTLS ClientKeyExchange message (see Section 4.1).

Figure 2 depicts the common protocol flow for the DTLS profile after C has retrieved the Access Token from AS.

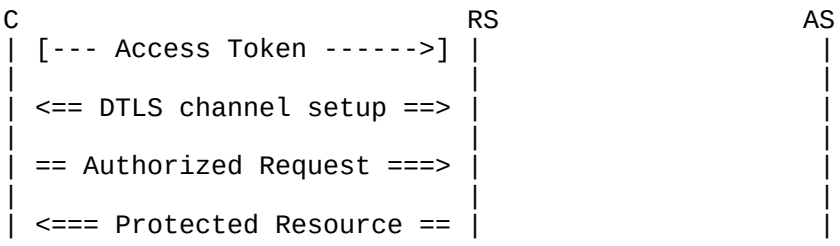

Figure 2: Protocol overview

The following sections specify how CoAP is used to interchange access-related data between RS and AS so that AS can provide C and RS with sufficient information to establish a secure channel, and convey authorization information specific for this communication relationship to RS.

Depending on the desired CoAP security mode, the Client-to-AS request, AS-to-Client response and DTLS session establishment carry slightly different information. Section 3 addresses the use of raw public keys while Section 4 defines how pre-shared keys are used in this profile.

2.1. Unauthorized Resource Request Message

The optional Unauthorized Resource Request message is a request for a resource hosted by RS for which no proper authorization is granted. RS MUST treat any CoAP request for a resource other than "/authz-info" as Unauthorized Resource Request message when any of the following holds:

- The request has been received on an unprotected channel.
- RS has no valid access token for the sender of the request regarding the requested action on that resource.
- RS has a valid access token for the sender of the request, but this does not allow the requested action on the requested resource.

Note: These conditions ensure that RS can handle requests autonomously once access was granted and a secure channel has been established between C and RS. The resource "/authz-info" is publicly accessible to be able to upload new access tokens to RS (cf. [I-D.ietf-ace-oauth-authz]).

Unauthorized Resource Request messages MUST be denied with a client error response. In this response, the Resource Server SHOULD provide proper AS Information to enable the Client to request an access token from RS's Authorization Server as described in Section 2.2.

The response code MUST be 4.01 (Unauthorized) in case the sender of the Unauthorized Resource Request message is not authenticated, or if RS has no valid access token for C. If RS has an access token for C but not for the resource that C has requested, RS MUST reject the request with a 4.03 (Forbidden). If RS has an access token for C but it does not cover the action C requested on the resource, RS MUST reject the request with a 4.05 (Method Not Allowed).

Note: The use of the response codes 4.03 and 4.05 is intended to prevent infinite loops where a dumb Client optimistically tries to access a requested resource with any access token received from AS. As malicious clients could pretend to be C to determine C's privileges, these detailed response codes must be used only when a certain level of security is already available which can be achieved only when the Client is authenticated.

2.2. AS Information

The AS Information is sent by RS as a response to an Unauthorized Resource Request message (see Section 2.1) to point the sender of the Unauthorized Resource Request message to RS's AS. The AS information is a set of attributes containing an absolute URI (see Section 4.3 of [RFC3986]) that specifies the AS in charge of RS.

TBD: We might not want to add more parameters in the AS information because
this would not only reveal too much information about RS's capabilities to unauthorized peers but also be of little value as C cannot really trust that information anyway.

The message MAY also contain a nonce generated by RS to ensure freshness in case that the RS and AS do not have synchronized clocks.

Figure 3 shows an example for an AS Information message payload using CBOR [RFC7049] diagnostic notation.

```
4.01 Unauthorized
Content-Format: application/ace+cbor
{AS: "coaps://as.example.com/token",
 nonce: h'e0a156bb3f'}
```

Figure 3: AS Information payload example

In this example, the attribute AS points the receiver of this message to the URI "coaps://as.example.com/token" to request access permissions. The originator of the AS Information payload (i.e., RS) uses a local clock that is loosely synchronized with a time scale common between RS and AS (e.g., wall clock time). Therefore, it has included a parameter "nonce" for replay attack prevention (c.f. Section 5.2).

Note: There is an ongoing discussion how freshness of access tokens can be achieved in constrained environments. This specification for now assumes that RS and AS do not have a common understanding of time that allows RS to achieve its security objectives without explicitly adding a nonce.

The examples in this document are written in CBOR diagnostic notation to improve readability. Figure 4 illustrates the binary encoding of the message payload shown in Figure 3.

```
a2                                    # map(2)
    00                                # unsigned(0) (=AS)
    78 1c                             # text(28)
       636f6170733a2f2f61732e657861
       6d706c652e636f6d2f746f6b656e   # "coaps://as.example.com/token"
    05                                # unsigned(5) (=nonce)
    45                                # bytes(5)
       e0a156bb3f
```

Figure 4: AS Information example encoded in CBOR

2.3. Resource Access

Once a DTLS channel has been established as described in Section 3 and Section 4, respectively, C is authorized to access resources covered by the Access Token it has uploaded to the "/authz-info" resource hosted by RS.

On the server side (i.e., RS), successful establishment of the DTLS channel binds C to the access token, functioning as a proof-of-possession associated key. Any request that RS receives on this channel MUST be checked against these authorization rules that are associated with the identity of C. Incoming CoAP requests that are not authorized with respect to any Access Token that is associated with C MUST be rejected by RS with 4.01 response as described in Section 2.1.

Note: The identity of C is determined by the authentication process

during the DTLS handshake. In the asymmetric case, the public key will define C's identity, while in the PSK case, C's identity is defined by the session key generated by AS for this communication.

RS SHOULD treat an incoming CoAP request as authorized if the following holds:

1. The message was received on a secure channel that has been established using the procedure defined in this document.

2. The authorization information tied to the sending peer is valid.

3. The request is destined for RS.

4. The resource URI specified in the request is covered by the authorization information.

5. The request method is an authorized action on the resource with respect to the authorization information.

Incoming CoAP requests received on a secure DTLS channel MUST be rejected

1. with response code 4.03 (Forbidden) when the resource URI specified in the request is not covered by the authorization information, and

2. with response code 4.05 (Method Not Allowed) when the resource URI specified in the request covered by the authorization information but not the requested action.

C cannot always know a priori if a Authorized Resource Request will succeed. If C repeatedly gets AS Information messages (cf. Section 2.2) as response to its requests, it SHOULD request a new Access Token from AS in order to continue communication with RS.

2.4. Dynamic Update of Authorization Information

The Client can update the authorization information stored at RS at any time. To do so, the Client requests from AS a new Access Token for the intended action on the respective resource and uploads this Access Token to the "/authz-info" resource on RS.

Figure 5 depicts the message flow where C requests a new Access Token after a security association between C and RS has been established using this protocol.

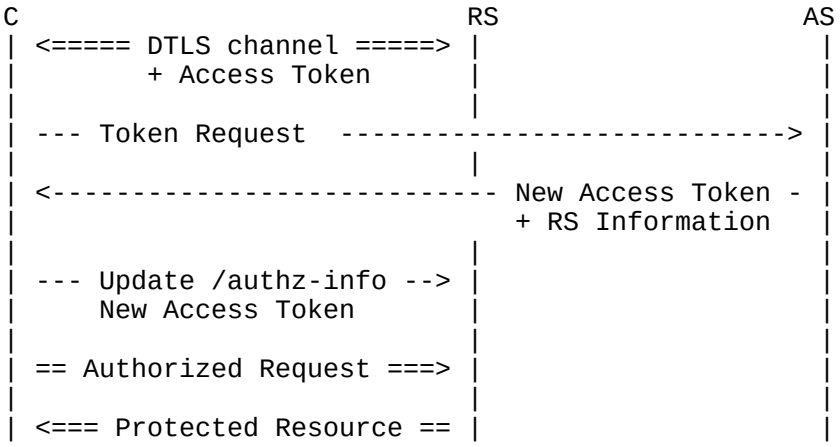

Figure 5: Overview of Dynamic Update Operation

3. RawPublicKey Mode

To retrieve an access token for the resource that C wants to access, C requests an Access Token from AS. C MUST add a "cnf" object carrying either its raw public key or a unique identifier for a public key that it has previously made known to AS.

An example Access Token request from C to RS is depicted in Figure 6.

```
POST coaps://as.example.com/token
Content-Format: application/cbor
{
  grant_type:    client_credentials,
  aud:           "tempSensor4711",
  cnf: {
    COSE_Key: {
      kty: EC2,
      crv: P-256,
      x:   h'TODOX',
      y:   h'TODOY'
    }
  }
}
```

Figure 6: Access Token Request Example for RPK Mode

The example shows an Access Token request for the resource identified by the audience string "tempSensor4711" on the AS using a raw public key.

When AS authorizes a request, it will return an Access Token and a "cnf" object in the AS-to-Client response. Before C initiates the DTLS handshake with RS, it MUST send a "POST" request containing the new Access Token to the "/authz-info" resource hosted by RS. If this operation yields a positive response, C SHOULD proceed to establish a new DTLS channel with RS. To use raw public key mode, C MUST pass the same public key that was used for constructing the Access Token with the SubjectPublicKeyInfo structure in the DTLS handshake as specified in [RFC7250].

Note: According to [RFC7252], CoAP implementations MUST support the ciphersuite TLS_ECDHE_ECDSA_WITH_AES_128_CCM_8 [RFC7251] and the NIST P-256 curve. C is therefore expected to offer at least this ciphersuite to RS.

The Access Token is constructed by AS such that RS can associate the Access Token with the Client’s public key. If CBOR web tokens [I-D.ietf-ace-cbor-web-token] are used as recommended in [I-D.ietf-ace-oauth-authz], the AS MUST include a "COSE_Key" object in the "cnf" claim of the Access Token. This "COSE_Key" object MAY contain a reference to a key for C that is already known by RS (e.g., from previous communication). If the AS has no certain knowledge that the Client’s key is already known to RS, the Client’s public key MUST be included in the Access Token’s "cnf" parameter.

4. PreSharedKey Mode

To retrieve an access token for the resource that C wants to access, C MAY include a "cnf" object carrying an identifier for a symmetric key in its Access Token request to AS. This identifier can be used by AS to determine the session key to construct the proof-of-possession token and therefore MUST specify a symmetric key that was previously generated by AS as a session key for the communication between C and RS.

Depending on the requested token type and algorithm in the Access Token request, AS adds RS Information to the response that provides C with sufficient information to setup a DTLS channel with RS. For symmetric proof-of-possession keys (c.f. [I-D.ietf-ace-oauth-authz]), C must ensure that the Access Token request is sent over a secure channel that guarantees authentication, message integrity and confidentiality.

When AS authorizes C it returns an AS-to-Client response with the profile parameter set to "coap_dtls" and a "cnf" parameter carrying a "COSE_Key" object that contains the symmetric session key to be used between C and RS as illustrated in Figure 7.

```
2.01 Created
Content-Format: application/cbor
Location-Path: /token/asdjbaskd
Max-Age: 86400
{
   access_token: b64'SlAV32hkKG ...
   (remainder of CWT omitted for brevity;
   token_type:   pop,
   alg:          HS256,
   expires_in:   86400,
   profile:      coap_dtls,
   cnf: {
     COSE_Key: {
       kty: symmetric,
       k: h'73657373696f6e6b6579'
     }
   }
}
```

Figure 7: Example Access Token response

In this example, AS returns a 2.01 response containing a new Access Token. The information is transferred as a CBOR data structure as specified in [I-D.ietf-ace-oauth-authz]. The Max-Age option tells the receiving Client how long this token will be valid.

A response that declines any operation on the requested resource is constructed according to Section 5.2 of RFC 6749 [4], (cf. Section 5.5.3 of [I-D.ietf-ace-oauth-authz]).

```
4.00 Bad Request
Content-Format: application/cbor
{
  error: invalid_request
}
```

Figure 8: Example Access Token response with reject

4.1. DTLS Channel Setup Between C and RS

When C receives an Access Token from AS, it checks if the payload contains an "access_token" parameter and a "cnf" parameter. With this information C can initiate establishment of a new DTLS channel with RS. To use DTLS with pre-shared keys, C follows the PSK key exchange algorithm specified in Section 2 of [RFC4279] using the key conveyed in the "cnf" parameter of the AS response as PSK when constructing the premaster secret.

In PreSharedKey mode, the knowledge of the session key by C and RS is used for mutual authentication between both peers. Therefore, RS must be able to determine the session key from the Access Token. Following the general ACE authorization framework, C can upload the Access Token to RS's "/authz-info" resource before starting the DTLS handshake. Alternatively, C MAY provide the most recent base64-encoded Access Token in the "psk_identity" field of the ClientKeyExchange message.

If RS receives a ClientKeyExchange message that contains a "psk_identity" with a length greater zero, it MUST base64-decode its contents and check if the "psk_identity" field contains a key identifier or Access Token according to the following CDDL specification:

```
psk_identity = {
  kid => bstr // access_token => bstr
}
```

The identifiers for the map keys "kid" and "access_token" are used with the same meaning as in COSE [I-D.ietf-cose-msg] and the ACE framework [I-D.ietf-ace-oauth-authz] respectively. The identifier "kid" thus has the value 4 (see [I-D.ietf-cose-msg]), and the identifier "access_token" has the value 19, respectively (see [I-D.ietf-ace-oauth-authz]).

If the "psk_identity" field contains a key identifier, the receiver MUST check if it has one or more Access Tokens that are associated with the specified key. If no valid Access Token is available for this key, the DTLS session setup is terminated with an "illegal_parameter" DTLS alert message.

If instead the "psk_identity" field contains an Access Token, it must processed in the same way as an Access Token that has been uploaded to its "/authz-info" resource. In this case, RS continues processing the ClientKeyExchange message if the contents of the "psk_identity" contained a valid Access Token. Otherwise, the DTLS session setup is terminated with an "illegal_parameter" DTLS alert message.

Note1: As RS cannot provide C with a meaningful PSK identity hint in response to C's ClientHello message, RS SHOULD NOT send a ServerKeyExchange message.

Note2: According to [RFC7252], CoAP implementations MUST support the ciphersuite TLS_PSK_WITH_AES_128_CCM_8 [RFC6655]. C is therefore expected to offer at least this ciphersuite to RS.

This specification assumes that the Access Token is a PoP token as described in [I-D.ietf-ace-oauth-authz] unless specifically stated otherwise. Therefore, the Access Token is bound to a symmetric PoP key that is used as session key between C and RS.

While C can retrieve the session key from the contents of the "cnf" parameter in the AS-to-Client response, RS uses the information contained in the "cnf" claim of the Access Token to determine the actual session key when no explicit "kid" was provided in the "psk_identity" field. Usually, this is done by including a "COSE_Key" object carrying either a key that has been encrypted with a shared secret between AS and RS, or a key identifier that can be used by RS to lookup the session key.

Instead of the "COSE_Key" object, AS MAY include a "COSE_Encrypt" structure to enable RS to calculate the session key from the Access Token. The "COSE_Encrypt" structure MUST use the _Direct Key with KDF_ method as described in Section 12.1.2 of draft-ietf-cose-msg [5]. The AS MUST include a Context information structure carrying a PartyU "nonce" parameter carrying the nonce that has been used by AS to construct the session key.

This specification mandates that at least the key derivation algorithm "HKDF SHA-256" as defined in [I-D.ietf-cose-msg] MUST be supported. This key derivation function is the default when no "alg" field is included in the "COSE_Encrypt" structure for RS.

4.2. Updating Authorization Information

Usually, the authorization information that RS keeps for C is updated by uploading a new Access Token as described in Section 2.4.

If the security association with RS still exists and RS has indicated support for session renegotiation according to [RFC5746], the new Access Token MAY be used to renegotiate the existing DTLS session. In this case, the Access Token is used as "psk_identity" as defined in Section 4.1. The Client MAY also perform a new DTLS handshake according to Section 4.1 that replaces the existing DTLS session.

After successful completion of the DTLS handshake RS updates the existing authorization information for C according to the new Access Token.

5. Security Considerations

TODO

5.1. Unprotected AS Information

Initially, no secure channel exists to protect the communication between C and RS. Thus, C cannot determine if the AS information contained in an unprotected response from RS to an unauthorized request (c.f. Section 2.2) is authentic. It is therefore advisable to provide C with a (possibly hard-coded) list of trustworthy authorization servers. AS information responses referring to a URI not listed there would be ignored.

5.2. Use of Nonces for Replay Protection

RS may add a nonce to the AS Information message sent as a response to an unauthorized request to ensure freshness of an Access Token subsequently presented to RS. While a timestamp of some granularity would be sufficient to protect against replay attacks, using randomized nonce is preferred to prevent disclosure of information about RS's internal clock characteristics.

5.3. Privacy

An unprotected response to an unauthorized request (c.f. Section 2.2) may disclose information about RS and/or its existing relationship with C. It is advisable to include as little information as possible in an unencrypted response. When a DTLS session between C and RS already exists, more detailed information may be included with an error response to provide C with sufficient information to react on that particular error.

6. IANA Considerations

This document has no actions for IANA.

7. References

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[3] https://tools.ietf.org/html/draft-ietf-ace-oauth-authz-03#section-8.1

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Authentication and Authorization for Constrained Environments (ACE) draft-ietf-ace-oauth-authz-07

Abstract

This specification defines a framework for authentication and authorization in Internet of Things (IoT) environments. The framework is based on a set of building blocks including OAuth 2.0 and CoAP, thus making a well-known and widely used authorization solution suitable for IoT devices. Existing specifications are used where possible, but where the constraints of IoT devices require it, extensions are added and profiles are defined.

Table of Contents

1. Introduction

Authorization is the process for granting approval to an entity to access a resource [RFC4949]. The authorization task itself can best be described as granting access to a requesting client, for a resource hosted on a device, the resource server (RS). This exchange is mediated by one or multiple authorization servers (AS). Managing authorization for a large number of devices and users is a complex task.

While prior work on authorization solutions for the Web and for the mobile environment also applies to the IoT environment many IoT devices are constrained, for example in terms of processing capabilities, available memory, etc. For web applications on constrained nodes this specification makes use of CoAP [RFC7252].

A detailed treatment of constraints can be found in [RFC7228], and the different IoT deployments present a continuous range of device and network capabilities. Taking energy consumption as an example: At one end there are energy-harvesting or battery powered devices which have a tight power budget, on the other end there are mains-powered devices, and all levels in between.

Hence, IoT devices may be very different in terms of available processing and message exchange capabilities and there is a need to support many different authorization use cases [RFC7744].

This specification describes a framework for authentication and authorization in constrained environments (ACE) built on re-use of OAuth 2.0 [RFC6749], thereby extending authorization to Internet of Things devices. This specification contains the necessary building blocks for adjusting OAuth 2.0 to IoT environments.

More detailed, interoperable specifications can be found in profiles. Implementations may claim conformance with a specific profile, whereby implementations utilizing the same profile interoperate while implementations of different profiles are not expected to be interoperable. Some devices, such as mobile phones and tablets, may implement multiple profiles and will therefore be able to interact with a wider range of low end devices. Requirements on profiles are described at contextually appropriate places througout this memo, and also summarized in Appendix C.

2. Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

Certain security-related terms such as "authentication", "authorization", "confidentiality", "(data) integrity", "message authentication code", and "verify" are taken from [RFC4949].

Since we describe exchanges as RESTful protocol interactions HTTP [RFC7231] offers useful terminology.

Terminology for entities in the architecture is defined in OAuth 2.0 [RFC6749] and [I-D.ietf-ace-actors], such as client (C), resource server (RS), and authorization server (AS).

Note that the term "endpoint" is used here following its OAuth definition, which is to denote resources such as /token and /introspect at the AS and /authz-info at the RS. The CoAP [RFC7252] definition, which is "An entity participating in the CoAP protocol" is not used in this memo.

Since this specification focuses on the problem of access control to resources, we simplify the actors by assuming that the client authorization server (CAS) functionality is not stand-alone but subsumed by either the authorization server or the client (see section 2.2 in [I-D.ietf-ace-actors]).

We call the specifications of this memo the "framework" or "ACE framework". When referring to "profiles of this framework" we mean additional memo's that define the use of this specification with concrete transport, and communication security protocols (e.g. CoAP over DTLS).

3. Overview

This specification defines the ACE framework for authorization in the Internet of Things environment. It consists of a set of building blocks.

The basic block is the OAuth 2.0 [RFC6749] framework, which enjoys widespread deployment. Many IoT devices can support OAuth 2.0 without any additional extensions, but for certain constrained settings additional profiling is needed.

Another building block is the lightweight web transfer protocol CoAP [RFC7252] for those communication environments where HTTP is not appropriate. CoAP typically runs on top of UDP which further reduces overhead and message exchanges. While this specification defines extensions for the use of OAuth over CoAP, we do envision further underlying protocols to be supported in the future, such as HTTP/2, MQTT and QUIC.

A third building block is CBOR [RFC7049] for encodings where JSON [RFC7159] is not sufficiently compact. CBOR is a binary encoding designed for small code and message size, which may be used for encoding of self contained tokens, and also for encoding CoAP POST parameters and CoAP responses.

A fourth building block is the compact CBOR-based secure message format COSE [RFC8152], which enables application layer security as an alternative or complement to transport layer security (DTLS [RFC6347] or TLS [RFC5246]). COSE is used to secure self contained tokens such as proof-of-possession (PoP) tokens, which is an extension to the OAuth access tokens, and "client tokens" which are defined in this framework (see Section 5.6.4). The default access token format is defined in CBOR web token (CWT) [I-D.ietf-ace-cbor-web-token]. Application layer security for CoAP using COSE can be provided with OSCOAP [I-D.ietf-core-object-security].

With the building blocks listed above, solutions satisfying various IoT device and network constraints are possible. A list of constraints is described in detail in RFC 7228 [RFC7228] and a description of how the building blocks mentioned above relate to the various constraints can be found in Appendix A.

Luckily, not every IoT device suffers from all constraints. The ACE framework nevertheless takes all these aspects into account and allows several different deployment variants to co-exist rather than mandating a one-size-fits-all solution. We believe this is important to cover the wide range of possible interworking use cases and the different requirements from a security point of view. Once IoT deployments mature, popular deployment variants will be documented in form of ACE profiles.

In the subsections below we provide further details about the different building blocks.

3.1. OAuth 2.0

The OAuth 2.0 authorization framework enables a client to obtain limited access to a resource with the permission of a resource owner. Authorization information, or references to it, is passed between the nodes using access tokens. These access tokens are issued to clients by an authorization server with the approval of the resource owner. The client uses the access token to access the protected resources hosted by the resource server.

A number of OAuth 2.0 terms are used within this specification:

The token and introspect Endpoints:

> The AS hosts the /token endpoint that allows a client to request access tokens. The client makes a POST request to the /token endpoint on the AS and receives the access token in the response (if the request was successful).

The token introspection endpoint, /introspect, is used by the RS when requesting additional information regarding a received access token. The RS makes a POST request to /introspect on the AS and receives information about the access token in the response. (See "Introspection" below.)

Access Tokens:

Access tokens are credentials needed to access protected resources. An access token is a data structure representing authorization permissions issued by the AS to the client. Access tokens are generated by the authorization server and consumed by the resource server. The access token content is opaque to the client.

Access tokens can have different formats, and various methods of utilization (e.g., cryptographic properties) based on the security requirements of the given deployment.

Proof of Possession Tokens:

An access token may be bound to a cryptographic key, which is then used by an RS to authenticate requests from a client. Such tokens are called proof-of-possession tokens (or PoP tokens).

The proof-of-possession (PoP) security concept assumes that the AS acts as a trusted third party that binds keys to access tokens. These so called PoP keys are then used by the client to demonstrate the possession of the secret to the RS when accessing the resource. The RS, when receiving an access token, needs to verify that the key used by the client matches the one bound to the access token. When this specification uses the term "access token" it is assumed to be a PoP token unless specifically stated otherwise.

The key bound to the access token (aka PoP key) may be based on symmetric as well as on asymmetric cryptography. The appropriate choice of security depends on the constraints of the IoT devices as well as on the security requirements of the use case.

Symmetric PoP key: The AS generates a random symmetric PoP key. The key is either stored to be returned on introspection calls or encrypted and included in the access token. The PoP key is also encrypted for the client and sent together with the access token to the client.

Asymmetric PoP key: An asymmetric key pair is generated on the client and the public key is sent to the AS (if it does not already have knowledge of the client's public key). Information about the public key, which is the PoP key in this case, is either stored to be returned on introspection calls or included inside the access token and sent back to the requesting client. The RS can identify the client's public key from the information in the token, which allows the client to use the corresponding private key for the proof of possession.

The access token is either a simple reference, or a structured information object (e.g. CWT [I-D.ietf-ace-cbor-web-token]), protected by a cryptographic wrapper (e.g. COSE [RFC8152]). The choice of PoP key does not necessarily imply a specific credential type for the integrity protection of the token.

Scopes and Permissions:

In OAuth 2.0, the client specifies the type of permissions it is seeking to obtain (via the scope parameter) in the access token request. In turn, the AS may use the scope response parameter to inform the client of the scope of the access token issued. As the client could be a constrained device as well, this specification uses CBOR encoded messages for CoAP, defined in Section 5, to request scopes and to be informed what scopes the access token actually authorizes.

The values of the scope parameter are expressed as a list of space- delimited, case-sensitive strings, with a semantic that is well-known to the AS and the RS. More details about the concept of scopes is found under Section 3.3 in [RFC6749].

Claims:

Information carried in the access token or returned from introspection, called claims, is in the form of type-value pairs. An access token may, for example, include a claim identifying the AS that issued the token (via the "iss" claim) and what audience the access token is intended for (via the "aud" claim). The audience of an access token can be a specific resource or one or many resource servers. The resource owner policies influence what claims are put into the access token by the authorization server.

While the structure and encoding of the access token varies throughout deployments, a standardized format has been defined with the JSON Web Token (JWT) [RFC7519] where claims are encoded as a JSON object. In [I-D.ietf-ace-cbor-web-token] an equivalent format using CBOR encoding (CWT) has been defined.

Introspection:

Introspection is a method for a resource server to query the authorization server for the active state and content of a received access token. This is particularly useful in those cases where the authorization decisions are very dynamic and/or where the received access token itself is a reference rather than a self-contained token. More information about introspection in OAuth 2.0 can be found in [RFC7662].

3.2. CoAP

CoAP is an application layer protocol similar to HTTP, but specifically designed for constrained environments. CoAP typically uses datagram-oriented transport, such as UDP, where reordering and loss of packets can occur. A security solution need to take the latter aspects into account.

While HTTP uses headers and query-strings to convey additional information about a request, CoAP encodes such information in so-called 'options'.

CoAP supports application-layer fragmentation of the CoAP payloads through blockwise transfers [RFC7959]. However, block-wise transfer does not increase the size limits of CoAP options, therefore data encoded in options has to be kept small.

Transport layer security for CoAP can be provided by DTLS 1.2 [RFC6347] or TLS 1.2 [RFC5246]. CoAP defines a number of proxy operations which requires transport layer security to be terminated at the proxy. One approach for protecting CoAP communication end-to-end through proxies, and also to support security for CoAP over a different transport in a uniform way, is to provide security on application layer using an object-based security mechanism such as COSE [RFC8152].

One application of COSE is OSCOAP [I-D.ietf-core-object-security], which provides end-to-end confidentiality, integrity and replay protection, and a secure binding between CoAP request and response messages. In OSCOAP, the CoAP messages are wrapped in COSE objects and sent using CoAP.

4. Protocol Interactions

The ACE framework is based on the OAuth 2.0 protocol interactions using the /token and /introspect endpoints. A client obtains an access token from an AS using the /token endpoint and subsequently presents the access token to a RS to gain access to a protected

resource. The RS, after receiving an access token, may present it to the AS via the /introspect endpoint to get information about the access token. In other deployments the RS may process the access token locally without the need to contact an AS. These interactions are shown in Figure 1. An overview of various OAuth concepts is provided in Section 3.1.

The OAuth 2.0 framework defines a number of "protocol flows" via grant types, which have been extended further with extensions to OAuth 2.0 (such as RFC 7521 [RFC7521] and [I-D.ietf-oauth-device-flow]). What grant types works best depends on the usage scenario and RFC 7744 [RFC7744] describes many different IoT use cases but there are two preferred grant types, namely the Authorization Code Grant (described in Section 4.1 of [RFC7521]) and the Client Credentials Grant (described in Section 4.4 of [RFC7521]). The Authorization Code Grant is a good fit for use with apps running on smart phones and tablets that request access to IoT devices, a common scenario in the smart home environment, where users need to go through an authentication and authorization phase (at least during the initial setup phase). The native apps guidelines described in [I-D.ietf-oauth-native-apps] are applicable to this use case. The Client Credential Grant is a good fit for use with IoT devices where the OAuth client itself is constrained. In such a case the resource owner or another person on his or her behalf have arranged with the authorization server out-of-band, which is often accomplished using a commissioning tool.

The consent of the resource owner, for giving a client access to a protected resource, can be provided dynamically as in the traditional OAuth flows, or it could be pre-configured by the resource owner as authorization policies at the AS, which the AS evaluates when a token request arrives. The resource owner and the requesting party (i.e. client owner) are not shown in Figure 1.

This framework supports a wide variety of communication security mechanisms between the ACE entities, such as client, AS, and RS. We assume that the client has been registered (also called enrolled or onboarded) to an AS using a mechanism defined outside the scope of this document. In practice, various techniques for onboarding have been used, such as factory-based provisioning or the use of commissioning tools. Regardless of the onboarding technique, this registration procedure implies that the client and the AS share credentials, and configuration parameters. These credentials are used to mutually authenticate each other and to protect messages exchanged between the client and the AS.

It is also assumed that the RS has been registered with the AS, potentially in a similar way as the client has been registered with

the AS. Established keying material between the AS and the RS allows the AS to apply cryptographic protection to the access token to ensure that its content cannot be modified, and if needed, that the content is confidentiality protected.

The keying material necessary for establishing communication security between C and RS is dynamically established as part of the protocol described in this document.

At the start of the protocol there is an optional discovery step where the client discovers the resource server and the resources this server hosts. In this step the client might also determine what permissions are needed to access the protected resource. The detailed procedures for this discovery process may be defined in an ACE profile and depend on the protocols being used and the specific deployment environment.

In Bluetooth Low Energy, for example, advertisements are broadcasted by a peripheral, including information about the primary services. In CoAP, as a second example, a client can make a request to "/.well-known/core" to obtain information about available resources, which are returned in a standardized format as described in [RFC6690].

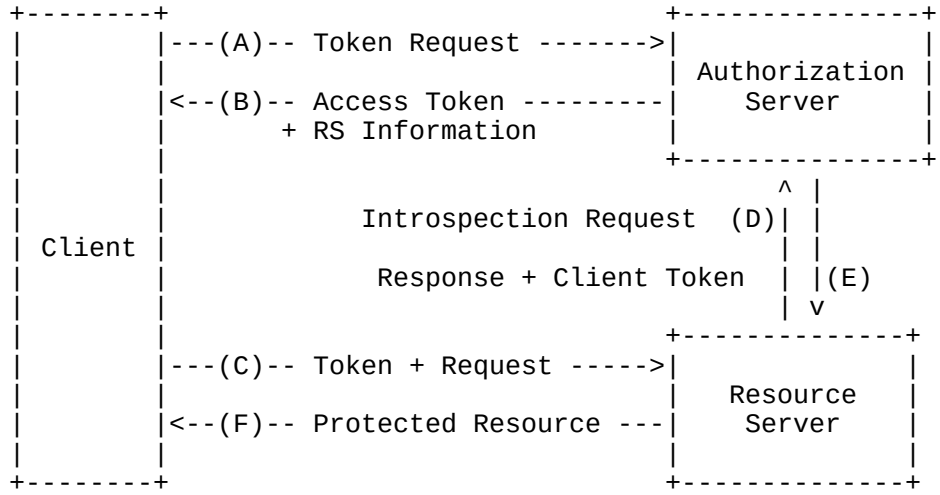

Figure 1: Basic Protocol Flow.

Requesting an Access Token (A):

The client makes an access token request to the /token endpoint at the AS. This framework assumes the use of PoP tokens (see Section 3.1 for a short description) wherein the AS binds a key to

an access token. The client may include permissions it seeks to obtain, and information about the credentials it wants to use (e.g., symmetric/asymmetric cryptography or a reference to a specific credential).

Access Token Response (B):

If the AS successfully processes the request from the client, it returns an access token. It also returns various parameters, referred as "RS Information". In addition to the response parameters defined by OAuth 2.0 and the PoP token extension, further response parameters, such as information on which profile the client should use with the resource server(s). More information about these parameters can be found in Section 5.5.4.

Resource Request (C):

The client interacts with the RS to request access to the protected resource and provides the access token. The protocol to use between the client and the RS is not restricted to CoAP. HTTP, HTTP/2, QUIC, MQTT, Bluetooth Low Energy, etc., are also viable candidates.

Depending on the device limitations and the selected protocol this exchange may be split up into two parts:

(1) the client sends the access token containing, or referencing, the authorization information to the RS, that may be used for subsequent resource requests by the client, and (2) the client makes the resource access request, using the communication security protocol and other RS Information obtained from the AS.

The Client and the RS mutually authenticate using the security protocol specified in the profile (see step B) and the keys obtained in the access token or the RS Information or the client token. The RS verifies that the token is integrity protected by the AS and compares the claims contained in the access token with the resource request. If the RS is online, validation can be handed over to the AS using token introspection (see messages D and E) over HTTP or CoAP, in which case the different parts of step C may be interleaved with introspection.

Token Introspection Request (D):

A resource server may be configured to introspect the access token by including it in a request to the /introspect endpoint at that

AS. Token introspection over CoAP is defined in Section 5.6 and for HTTP in [RFC7662].

Note that token introspection is an optional step and can be omitted if the token is self-contained and the resource server is prepared to perform the token validation on its own.

Token Introspection Response (E):

The AS validates the token and returns the most recent parameters, such as scope, audience, validity etc. associated with it back to the RS. The RS then uses the received parameters to process the request to either accept or to deny it. The AS can additionally return information that the RS needs to pass on to the client in the form of a client token. The latter is used to establish keys for mutual authentication between client and RS, when the client has no direct connectivity to the AS, see Section 5.6.4 for details.

Protected Resource (F):

If the request from the client is authorized, the RS fulfills the request and returns a response with the appropriate response code. The RS uses the dynamically established keys to protect the response, according to used communication security protocol.

5. Framework

The following sections detail the profiling and extensions of OAuth 2.0 for constrained environments which constitutes the ACE framework.

Credential Provisioning

For IoT we cannot generally assume that the client and RS are part of a common key infrastructure, so the AS provisions credentials or associated information to allow mutual authentication. These credentials need to be provided to the parties before or during the authentication protocol is executed, and may be re-used for subsequent token requests.

Proof-of-Possession

The ACE framework by default implements proof-of-possession for access tokens, i.e. that the token holder can prove being a holder of the key bound to the token. The binding is provided by the "cnf" claim indicating what key is used for proof-of-possession. If clients need to update a token, e.g. to get additional rights,

they can request that the AS binds the new access token to the same key as the previous token.

ACE Profiles

The client or RS may be limited in the encodings or protocols it supports. To support a variety of different deployment settings, specific interactions between client and RS are defined in an ACE profile. In ACE framework the AS is expected to manage the matching of compatible profile choices between a client and an RS. The AS informs the client of the selected profile using the "profile" parameter in the token response.

OAuth 2.0 requires the use of TLS both to protect the communication between AS and client when requesting an access token; between client and RS when accessing a resource and between AS and RS for introspection. In constrained settings TLS is not always feasible, or desirable. Nevertheless it is REQUIRED that the data exchanged with the AS is encrypted and integrity protected. It is furthermore REQUIRED that the AS and the endpoint communicating with it (client or RS) perform mutual authentication.

Profiles MUST specify how mutual authentication is done, depending e.g. on the communication protocol and the credentials used by the client or the RS.

In OAuth 2.0 the communication with the Token and the Introspection endpoints at the AS is assumed to be via HTTP and may use Uri-query parameters. This framework RECOMMENDS to use CoAP instead and RECOMMENDS the use of the following alternative instead of Uri-query parameters: The sender (client or RS) encodes the parameters of its request as a CBOR map and submits that map as the payload of the POST request. The Content-format depends on the security applied to the content and MUST be specified by the profile that is used.

The OAuth 2.0 AS uses a JSON structure in the payload of its responses both to client and RS. This framework RECOMMENDS the use of CBOR [RFC7049] instead. The requesting device can explicitly request this encoding by setting the CoAP Accept option in the request to "application/cbor". Depending on the profile, the content MAY arrive in a different format wrapping a CBOR payload.

5.1. Authorization Grants

To request an access token, the client obtains authorization from the resource owner or uses its client credentials as grant. The authorization is expressed in the form of an authorization grant.

The OAuth framework defines four grant types. The grant types can be split up into two groups, those granted on behalf of the resource owner (password, authorization code, implicit) and those for the client (client credentials).

The grant type is selected depending on the use case. In cases where the client acts on behalf of the resource owner, authorization code grant is recommended. If the client acts on behalf of the resource owner, but does not have any display or very limited interaction possibilities it is recommended to use the device code grant defined in [I-D.ietf-oauth-device-flow]. In cases where the client does not act on behalf of the resource owner, client credentials grant is recommended.

For details on the different grant types see the OAuth 2.0 framework. The OAuth 2.0 framework provides an extension mechanism for defining additional grant types so profiles of this framework MAY define additional grant types if needed.

5.2. Client Credentials

Authentication of the client is mandatory independent of the grant type when requesting the access token from the token endpoint. In the case of client credentials grant type the authentication and grant coincides.

Client registration and provisioning of client credentials to the client is out of scope for this specification.

The OAuth framework, [RFC6749], defines one client credential type, client id and client secret. Profiles of this framework MAY extend with additional client credentials such as DTLS pre-shared keys or client certificates.

5.3. AS Authentication

Client credential does not by default authenticate the AS that the client connects to. In classic OAuth the AS is authenticated with a TLS server certificate.

Profiles of this framework SHOULD specify how clients authenticate the AS and how communication security is implemented, otherwise server side TLS certificates as defined by OAuth 2.0 is required.

5.4. The 'Authorize' Endpoint

The authorization endpoint is used to interact with the resource owner and obtain an authorization grant in certain grant flows. Since it requires the use of a user agent (i.e. browser), it is not expected that these types of grant flow will be used by constrained clients. This endpoint is therefore out of scope for this specification. Implementations should use the definition and recommendations of [RFC6749] and [RFC6819].

If clients involved cannot support HTTP and TLS profiles MAY define mappings for the authorization endpoint.

5.5. The 'Token' Endpoint

In plain OAuth 2.0 the AS provides the /token endpoint for submitting access token requests. This framework extends the functionality of the /token endpoint, giving the AS the possibility to help client and RS to establish shared keys or to exchange their public keys. Furthermore this framework defines encodings using CoAP and CBOR, in addition to HTTP and JSON.

For the AS to be able to issue a token the client MUST be authenticated and present a valid grant for the scopes requested.

The figures of this section uses CBOR diagnostic notation without the integer abbreviations for the parameters or their values for better readability.

5.5.1. Client-to-AS Request

The client sends a CoAP POST request to the token endpoint at the AS, the profile MUST specify the Content-Type and wrapping of the payload. The content of the request consists of the parameters specified in section 4 of the OAuth 2.0 specification [RFC6749] encoded as a CBOR map.

In addition to these parameters, this framework defines the following parameters for requesting an access token from a /token endpoint:

aud
: OPTIONAL. Specifies the audience for which the client is requesting an access token. If this parameter is missing it is assumed that the client and the AS have a pre-established understanding of the audience that an access token should address. If a client submits a request for an access token without specifying an "aud" parameter, and the AS does not have a default

"aud" value for this client, then the AS MUST respond with an error message with the CoAP response code 4.00 (Bad Request).

cnf
OPTIONAL. This field contains information about the key the client would like to bind to the access token for proof-of-possession. It is RECOMMENDED that an AS reject a request containing a symmetric key value in the 'cnf' field. See Section 5.5.4.5 for more details on the formatting of the 'cnf' parameter.

The following examples illustrate different types of requests for proof-of-possession tokens.

Figure 2 shows a request for a token with a symmetric proof-of-possession key. Note that in this example we assume a DTLS-based communication security profile, therefore the Content-Type is "application/cbor". The content is displayed in CBOR diagnostic notation, without abbreviations for better readability.

```
Header: POST (Code=0.02)
Uri-Host: "server.example.com"
Uri-Path: "token"
Content-Type: "application/cbor"
Payload:
{
  "grant_type" : "client_credentials",
  "aud" : "tempSensor4711",
 }
```

Figure 2: Example request for an access token bound to a symmetric key.

Figure 3 shows a request for a token with an asymmetric proof-of-possession key. Note that in this example we assume an object security-based profile, therefore the Content-Type is "application/cose".

```
Header: POST (Code=0.02)
Uri-Host: "server.example.com"
Uri-Path: "token"
Content-Type: "application/cose"
Payload:
"COSE_Encrypted" : {
   16(
      [ h'a1010a', # protected header: {"alg" : "AES-CCM-16-64-128"}
      {5 : b64'ifUvZaHFgJM7UmGnjA'},  # unprotected header, IV
      b64'WXThuZo6TMCaZZqi6ef/8WHTjOdGk8kNzaIhIQ' # ciphertext
      ]
   )
}

Decrypted payload:
{
  "grant_type" : "client_credentials",
  "cnf" : {
    "COSE_Key" : {
      "kty" : "EC",
      "kid" : h'11',
      "crv" : "P-256",
      "x" : b64'usWxHK2PmfnHKwXPS54m0kTcGJ90UiglWiGahtagnv8',
      "y" : b64'IBOL+C3BttVivg+lSreASjpkttcsz+1rb7btKLv8EX4'
    }
  }
}
```

Figure 3: Example token request bound to an asymmetric key.

Figure 4 shows a request for a token where a previously communicated proof-of-possession key is only referenced. Note that we assume a DTLS-based communication security profile for this example, therefore the Content-Type is "application/cbor". Also note that the client performs a password based authentication in this example by submitting its client_secret (see section 2.3.1. of [RFC6749]).

```
Header: POST (Code=0.02)
Uri-Host: "server.example.com"
Uri-Path: "token"
Content-Type: "application/cbor"
Payload:
{
  "grant_type" : "client_credentials",
  "client_id" : "myclient",
  "client_secret" : "mysecret234",
  "aud" : "valve424",
  "scope" : "read",
  "cnf" : {
    "kid" : b64'6kg0dXJM13U'
  }
}
```

Figure 4: Example request for an access token bound to a key reference.

5.5.2. AS-to-Client Response

If the access token request has been successfully verified by the AS and the client is authorized to obtain an access token corresponding to its access token request, the AS sends a response with the CoAP response code 2.01 (Created). If client request was invalid, or not authorized, the AS returns an error response as described in Section 5.5.3.

Note that the AS decides which token type and profile to use when issuing a successful response. It is assumed that the AS has prior knowledge of the capabilities of the client, and the RS (see Appendix D. This prior knowledge may, for example, be set by the use of a dynamic client registration protocol exchange [RFC7591].

The content of the successful reply is the RS Information. It MUST be encoded as CBOR map, containing parameters as specified in section 5.1 of [RFC6749]. In addition to these parameters, the following parameters are also part of a successful response:

profile
: REQUIRED. This indicates the profile that the client MUST use towards the RS. See Section 5.5.4.4 for the formatting of this parameter.

cnf
: REQUIRED if the token type is 'pop'. OPTIONAL otherwise. If a symmetric proof-of-possession algorithms was selected, this field contains the proof-of-possession key. If an asymmetric algorithm

was selected, this field contains information about the public key used by the RS to authenticate. See Section 5.5.4.5 for the formatting of this parameter.

token_type
OPTIONAL. By default implementations of this framework SHOULD assume that the token_type is 'pop'. If a specific use case requires another token_type (e.g. 'Bearer') to be used then this parameter is REQUIRED.

Note that if CBOR Web Tokens [I-D.ietf-ace-cbor-web-token] are used, the access token can also contain a 'cnf' claim. This claim is however consumed by a different party. The access token is created by the AS and processed by the RS (and opaque to the client) whereas the RS Information is created by the AS and processed by the client; it is never forwarded to the resource server.

Figure Figure 5 summarizes the parameters that may be part of the RS Information.

Parameter name	Specified in
access_token	RFC 6749
token_type	RFC 6749
expires_in	RFC 6749
refresh_token	RFC 6749
scope	RFC 6749
state	RFC 6749
profile	[[this specification]]
cnf	[[this specification]]

Figure 5: RS Information parameters

Figure 6 shows a response containing a token and a 'cnf' parameter with a symmetric proof-of-possession key. Note that we assume a DTLS-based communication security profile for this example, therefore the Content-Type is "application/cbor".

```
Header: Created (Code=2.01)
Content-Type: "application/cbor"
Payload:
{
  "access_token" : b64'SlAV32hkKG ...
   (remainder of CWT omitted for brevity;
   CWT contains COSE_Key in the 'cnf' claim)',
  "profile" : "coap_dtls",
  "expires_in" : "3600",
  "cnf" : {
    "COSE_Key" : {
      "kty" : "Symmetric",
      "kid" : b64'39Gqlw',
      "k" : b64'hJtXhkV8FJG+Onbc6mxCcQh'
    }
  }
}
```

Figure 6: Example AS response with an access token bound to a symmetric key.

5.5.3. Error Response

The error responses for CoAP-based interactions with the AS are equivalent to the ones for HTTP-based interactions as defined in section 5.2 of [RFC6749], with the following differences:

- The Content-Type MUST be specified by the communication security profile used between client and AS. The raw payload before being processed by the communication security protocol MUST be encoded as a CBOR map.
- The CoAP response code 4.00 (Bad Request) MUST be used for all error responses, except for invalid_client where the CoAP response code 4.01 (Unauthorized) MAY be used under the same conditions as specified in section 5.2 of [RFC6749].
- The parameters "error", "error_description" and "error_uri" MAY be abbreviated using the codes specified in table Figure 13.
- The error codes MAY be abbreviated using the codes specified in table Figure 7.

error code	CBOR Key	Major Type
invalid_request	0	0 (uint)
invalid_client	1	0
invalid_grant	2	0
unauthorized_client	3	0
unsupported_grant_type	4	0
invalid_scope	5	0
unsupported_pop_key	6	0

Figure 7: CBOR abbreviations for common error codes

In addition to the error responses defined in OAuth 2.0, the follwoing behaviour MUST be implemented by the AS: If the client submits an asymmetric key in the token request that the RS cannot process, the AS MUST reject that request with the CoAP response code 4.00 (Bad Request) with the error code "unsupported_pop_key" defined in figure Figure 7.

5.5.4. Request and Response Parameters

This section provides more detail about the new parameters that can be used in access token requests and responses, as well as abbreviations for more compact encoding of existing parameters and common parameter values.

5.5.4.1. Audience

This parameter specifies for which audience the client is requesting a token. It should be encoded as CBOR text string (major type 3). The formatting and semantics of these strings are application specific.

5.5.4.2. Grant Type

The abbreviations in Figure 8 MAY be used in CBOR encodings instead of the string values defined in [RFC6749].

grant_type	CBOR Key	Major Type
password	0	0 (uint)
authorization_code	1	0
client_credentials	2	0
refresh_token	3	0

Figure 8: CBOR abbreviations for common grant types

5.5.4.3. Token Type

The token_type parameter is defined in [RFC6749], allowing the AS to indicate to the client which type of access token it is receiving (e.g. a bearer token).

This document registers the new value "pop" for the OAuth Access Token Types registry, specifying a Proof-of-Possession token. How the proof-of-possession is performed MUST be specified by the profiles.

The values in the 'token_type' parameter MUST be CBOR text strings (major type 3).

In this framework token type 'pop' MUST be assumed by default if the AS does not provide a different value.

5.5.4.4. Profile

Profiles of this framework MUST define the communication protocol and the communication security protocol between the client and the RS. Furthermore profiles MUST define proof-of-possession methods, if they support proof-of-possession tokens.

A profile MUST specify an identifier that is used to uniquely identify itself in the 'profile' parameter.

Profiles MAY define additional parameters for both the token request and the RS Information in the access token response in order to support negotiation or signalling of profile specific parameters.

5.5.4.5. Confirmation

The "cnf" parameter identifies or provides the key used for proof-of-possession or for authenticating the RS depending on the proof-of-possession algorithm and the context cnf is used in. This framework extends the definition of 'cnf' from [RFC7800] by adding CBOR/COSE

encodings and the use of 'cnf' for transporting keys in the RS Information.

The "cnf" parameter is used in the following contexts with the following meaning:

- In the access token, to indicate the proof-of-possession key bound to this token.
- In the token request C -> AS, to indicate the client's raw public key, or the key-identifier of a previously established key between C and RS.
- In the token response AS -> C, to indicate either the symmetric key generated by the AS for proof-of-possession or the raw public key used by the RS to authenticate.
- In the introspection response AS -> RS, to indicate the proof-of-possession key bound to the introspected token.
- In the client token AS -> RS -> C, to indicate the proof-of-possession key bound to the access token.

A CBOR encoded payload MAY contain the 'cnf' parameter with the following contents:

COSE_Key In this case the 'cnf' parameter contains the proof-of-possession key to be used by the client. An example is shown in Figure 9.

```
"cnf" : {
  "COSE_Key" : {
    "kty" : "EC",
    "kid" : h'11',
    "crv" : "P-256",
    "x" : b64'usWxHK2PmfnHKwXPS54m0kTcGJ90UiglWiGahtagnv8',
    "y" : b64'IBOL+C3BttVivg+lSreASjpkttcsz+1rb7btKLv8EX4'
  }
}
```

Figure 9: Confirmation parameter containing a public key

Note that the COSE_Key structure may contain an "alg" or "key_ops" parameter. If such parameters are present, a client MUST NOT use a key that is not compatible with the profile or proof-of-possession algorithm according to those parameters.

COSE_Encrypted In this case the 'cnf' parameter contains an encrypted symmetric key destined for the client. The client is assumed to be able to decrypt the ciphertext of this parameter. The parameter is encoded as COSE_Encrypted object wrapping a COSE_Key object. Figure 10 shows an example of this type of encoding.

```
"cnf" : {
  "COSE_Encrypted" : {
    993(
      [ h'a1010a' # protected header : {"alg" : "AES-CCM-16-64-128"}
        "iv" : b64'ifUvZaHFgJM7UmGnjA',  # unprotected header
       b64'WXThuZo6TMCaZZqi6ef/8WHTjOdGk8kNzaIhIQ' # ciphertext
      ]
    )
  }
}
```

Figure 10: Confirmation parameter containing an encrypted symmetric key

The ciphertext here could e.g. contain a symmetric key as in Figure 11.

```
{
  "kty" : "Symmetric",
  "kid" : b64'39Gqlw',
  "k" : b64'hJtXhkV8FJG+Onbc6mxCcQh'
}
```

Figure 11: Example plaintext of an encrypted cnf parameter

Key Identifier In this case the 'cnf' parameter references a key that is assumed to be previously known by the recipient. This allows clients that perform repeated requests for an access token for the same audience but e.g. with different scopes to omit key transport in the access token, token request and token response. Figure 12 shows such an example.

```
"cnf" : {
  "kid" : b64'39Gqlw'
}
```

Figure 12: A Confirmation parameter with just a key identifier

This specification establishes the IANA "CWT Confirmation Methods" registry for these types of confirmation methods in Section 8.10 and registers the methods defined by this specification. Other specifications can register other methods used for confirmation. The registry is meant to be analogous to the "JWT Confirmation Methods" registry defined by [RFC7800].

5.5.5. Mapping parameters to CBOR

All OAuth parameters in access token requests and responses are mapped to CBOR types as follows and are given an integer key value to save space.

Parameter name	CBOR Key	Major Type
aud	3	3
client_id	8	3 (text string)
client_secret	9	2 (byte string)
response_type	10	3
redirect_uri	11	3
scope	12	3
state	13	3
code	14	2
error	15	3
error_description	16	3
error_uri	17	3
grant_type	18	0
access_token	19	3
token_type	20	0
expires_in	21	0
username	22	3
password	23	3
refresh_token	24	3
cnf	25	5 (map)
profile	26	3

Figure 13: CBOR mappings used in token requests

5.6. The 'Introspect' Endpoint

Token introspection [RFC7662] is used by the RS and potentially the client to query the AS for metadata about a given token e.g. validity or scope. Analogous to the protocol defined in RFC 7662 [RFC7662] for HTTP and JSON, this section defines adaptations to more constrained environments using CoAP and CBOR.

Communication between the RS and the introspection endpoint at the AS MUST be integrity protected and encrypted. Furthermore AS and RS MUST perform mutual authentication. Finally the AS SHOULD verify that the RS has the right to access introspection information about the provided token. Profiles of this framework that support introspection MUST specify how authentication and communication security between RS and AS is implemented.

The figures of this section uses CBOR diagnostic notation without the integer abbreviations for the parameters or their values for better readability.

5.6.1. RS-to-AS Request

The RS sends a CoAP POST request to the introspection endpoint at the AS, the profile MUST specify the Content-Type and wrapping of the payload. The payload MUST be encoded as a CBOR map with a 'token' parameter containing the access token along with optional parameters representing additional context that is known by the RS to aid the AS in its response.

The same parameters are required and optional as in section 2.1 of RFC 7662 [RFC7662].

For example, Figure 14 shows a RS calling the token introspection endpoint at the AS to query about an OAuth 2.0 proof-of-possession token. Note that we assume a object security-based communication security profile for this example, therefore the Content-Type is "application/cose+cbor".

```
Header: POST (Code=0.02)
Uri-Host: "server.example.com"
Uri-Path: "introspect"
Content-Type: "application/cose+cbor"
Payload:
{
  "token" : b64'7gj0dXJQ43U',
  "token_type_hint" : "pop"
}
```

Figure 14: Example introspection request.

5.6.2. AS-to-RS Response

If the introspection request is authorized and successfully processed, the AS sends a response with the CoAP response code 2.01 (Created). If the introspection request was invalid, not authorized or couldn't be processed the AS returns an error response as described in Section 5.6.3.

In a successful response, the AS encodes the response parameters in a CBOR map including with the same required and optional parameters as in section 2.2. of RFC 7662 [RFC7662] with the following additions:

cnf

OPTIONAL. This field contains information about the proof-of-possession key that binds the client to the access token. See Section 5.5.4.5 for more details on the formatting of the 'cnf' parameter.

profile
OPTIONAL. This indicates the profile that the RS MUST use with the client. See Section 5.5.4.4 for more details on the formatting of this parameter.

client_token
OPTIONAL. This parameter contains information that the RS MUST pass on to the client. See Section 5.6.4 for more details.

For example, Figure 15 shows an AS response to the introspection request in Figure 14. Note that we assume a DTLS-based communication security profile for this example, therefore the Content-Type is "application/cbor".

```
Header: Created Code=2.01)
Content-Type: "application/cbor"
Payload:
{
  "active" : true,
  "scope" : "read",
  "profile" : "coap_dtls",
  "client_token" : b64'2QPhg0OhAQo ...
  (remainder of client token omitted for brevity)',
  "cnf" : {
    "COSE_Key" : {
      "kty" : "Symmetric",
      "kid" : b64'39Gqlw',
      "k" : b64'hJtXhkV8FJG+Onbc6mxCcQh'
    }
  }
}
```

Figure 15: Example introspection response.

5.6.3. Error Response

The error responses for CoAP-based interactions with the AS are equivalent to the ones for HTTP-based interactions as defined in section 2.3 of [RFC7662], with the following differences:

- If content is sent, the Content-Type MUST be set according to the specification of the communication security profile, and the content payload MUST be encoded as a CBOR map.

- o If the credentials used by the RS are invalid the AS MUST respond with the CoAP response code 4.01 (Unauthorized) and use the required and optional parameters from section 5.2 in RFC 6749 [RFC6749].
- o If the RS does not have the right to perform this introspection request, the AS MUST respond with the CoAP response code 4.03 (Forbidden). In this case no payload is returned.
- o The parameters "error", "error_description" and "error_uri" MAY be abbreviated using the codes specified in table Figure 13.
- o The error codes MAY be abbreviated using the codes specified in table Figure 7.

Note that a properly formed and authorized query for an inactive or otherwise invalid token does not warrant an error response by this specification. In these cases, the authorization server MUST instead respond with an introspection response with the "active" field set to "false".

5.6.4. Client Token

EDITORIAL NOTE: We have tentatively introduced this concept and would specifically like feedback whether this is viewed as a useful addition to the framework.

In cases where the client has limited connectivity and needs to get access to a previously unknown resource servers, this framework suggests the following approach: The client is pre-configured with a generic, long-term access token when it is commissioned. When the client then tries to access a RS it transmits this access token. The RS then performs token introspection to learn what access this token grants. In the introspection response, the AS also relays information for the client, such as the proof-of-possession key, through the RS. The RS passes on this Client Token to the client in response to the submission of the token.

The client_token parameter is designed to carry such information, and is intended to be used as described in Figure 16.

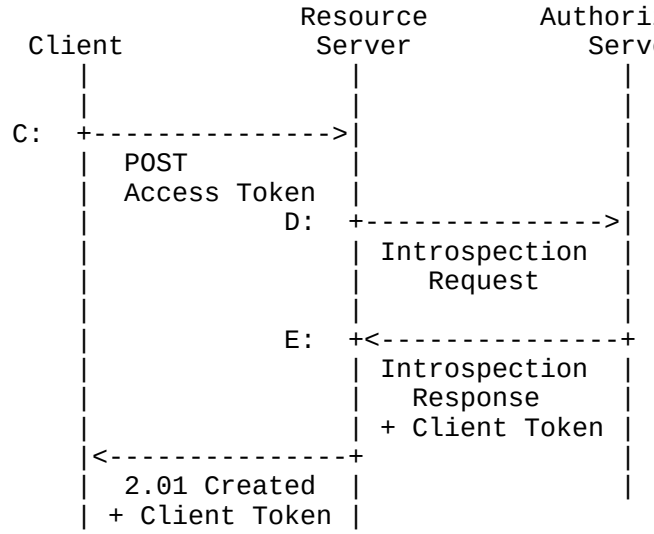

Figure 16: Use of the client_token parameter.

The client token is a COSE_Encrypted object, containing as payload a CBOR map with the following claims:

cnf
: REQUIRED if the token type is 'pop', OPTIONAL otherwise. Contains information about the proof-of-possession key the client is to use with its access token. See Section 5.5.4.5.

token_type
: OPTIONAL. See Section 5.5.4.3.

profile
: REQUIRED. See Section 5.5.4.4.

rs_cnf
: OPTIONAL. Contains information about the key that the RS uses to authenticate towards the client. If the key is symmetric then this claim MUST NOT be part of the Client Token, since this is the same key as the one specified through the 'cnf' claim. This claim uses the same encoding as the 'cnf' parameter. See Section 5.5.4.4.

The AS encrypts this token using a key shared between the AS and the client, so that only the client can decrypt it and access its payload. How this key is established is out of scope of this framework, however it can be established at the same time at which the client's long term token is created.

5.6.5. Mapping Introspection parameters to CBOR

The introspection request and response parameters are mapped to CBOR types as follows and are given an integer key value to save space.

Parameter name	CBOR Key	Major Type
iss	1	3 (text string)
sub	2	3
aud	3	3
exp	4	6 tag value 1
nbf	5	6 tag value 1
iat	6	6 tag value 1
cti	7	2 (byte string)
client_id	8	3
scope	12	3
token_type	20	3
username	22	3
cnf	25	5 (map)
profile	26	0 (uint)
token	27	3
token_type_hint	28	3
active	29	0
client_token	30	3
rs_cnf	31	5

Figure 17: CBOR Mappings to Token Introspection Parameters.

5.7. The Access Token

This framework RECOMMENDS the use of CBOR web token (CWT) as specified in [I-D.ietf-ace-cbor-web-token].

In order to facilitate offline processing of access tokens, this draft specifies the "cnf" and "scope" claims for CBOR web tokens.

The "scope" claim explicitly encodes the scope of a given access token. This claim follows the same encoding rules as defined in section 3.3 of [RFC6749]. The meaning of a specific scope value is application specific and expected to be known to the RS running that application.

The "cnf" claim follows the same rules as specified for JOSE web token in RFC7800 [RFC7800], except that it is encoded in COSE in the same way as specified for the "cnf" parameter in Section 5.5.4.5.

5.7.1. The 'Authorization Information' Endpoint

The access token, containing authorization information and information about the key used by the client, needs to be transported to the RS so that the RS can authenticate and authorize the client request.

This section defines a method for transporting the access token to the RS using CoAP. Profiles of this framework MAY define other methods for token transport.

The method consists of an /authz-info endpoint, implemented by the RS. A client using this method MUST make a POST request to /authz-info at the RS with the access token in the payload. The RS receiving the token MUST verify the validity of the token. If the token is valid, the RS MUST respond to the POST request with 2.01 (Created). This response MAY contain the identifier of the token (e.g. the cti for a CWT) as a payload.

If the token is not valid, the RS MUST respond with the CoAP response code 4.01 (Unauthorized). If the token is valid but the audience of the token does not match the RS, the RS MUST respond with the CoAP response code 4.03 (Forbidden). If the token is valid but is associated to claims that the RS cannot process (e.g. an unknown scope) the RS MUST respond with the CoAP response code 4.00 (Bad Request). In the latter case the RS MAY provide additional information in the error response, in order to clarify what went wrong.

The RS MAY make an introspection request to validate the token before responding to the POST /authz-info request. If the introspection response contains a client token (Section 5.6.4) then this token SHALL be included in the payload of the 2.01 (Created) response.

Profiles MUST specify how the /authz-info endpoint is protected. Note that since the token contains information that allow the client and the RS to establish a security context in the first place, mutual authentication may not be possible at this point.

The RS MUST be prepared to store more than one token for each client, and MUST apply the combined permissions granted by all applicable, valid tokens to client requests.

5.7.2. Token Expiration

Depending on the capabilities of the RS, there are various ways in which it can verify the validity of a received access token. We list

the possibilities here including what functionality they require of the RS.

- The token is a CWT/JWT and includes a 'exp' claim and possibly the 'nbf' claim. The RS verifies these by comparing them to values from its internal clock as defined in [RFC7519]. In this case the RS's internal clock must reflect the current date and time, or at least be synchronized with the AS's clock. How this clock synchronization would be performed is out of scope for this memo.
- The RS verifies the validity of the token by performing an introspection request as specified in Section 5.6. This requires the RS to have a reliable network connection to the AS and to be able to handle two secure sessions in parallel (C to RS and AS to RS).
- The RS and the AS both store a sequence number linked to their common security association. The AS increments this number for each access token it issues and includes it in the access token, which is a CWT. The RS keeps track of the most recently received sequence number, and only accepts tokens as valid, that are in a certain range around this number. This method does only require the RS to keep track of the sequence number. The method does not provide timely expiration, but it makes sure that older tokens cease to be valid after a certain number of newer ones got issued. For a constrained RS with no network connectivity and no means of reliably measuring time, this is the best that can be achieved.

If a token, that authorizes a long running request such as e.g. a CoAP Observe [RFC7641], expires, the RS MUST send an error response with the response code 4.01 Unauthorized to the client and then terminate processing the long running request.

6. Security Considerations

Security considerations applicable to authentication and authorization in RESTful environments provided in OAuth 2.0 [RFC6749] apply to this work, as well as the security considerations from [I-D.ietf-ace-actors]. Furthermore [RFC6819] provides additional security considerations for OAuth which apply to IoT deployments as well.

A large range of threats can be mitigated by protecting the contents of the access token by using a digital signature or a keyed message digest (MAC) or an AEAD algorithm. Consequently, the token integrity protection MUST be applied to prevent the token from being modified, particularly since it contains a reference to the symmetric key or the asymmetric key. If the access token contains the symmetric key, this symmetric key MUST be encrypted by the authorization server so that only the resource server can decrypt it. Note that using an

AEAD algorithm is preferrable over using a MAC unless the message needs to be publicly readable.

It is important for the authorization server to include the identity of the intended recipient (the audience), typically a single resource server (or a list of resource servers), in the token. Using a single shared secret with multiple resource servers to simplify key management is NOT RECOMMENDED since the benefit from using the proof-of-possession concept is significantly reduced.

The authorization server MUST offer confidentiality protection for any interactions with the client. This step is extremely important since the client may obtion the proof-of-possession key from the authorization server for use with a specific access token. Not using confidentiality protection exposes this secret (and the access token) to an eavesdropper thereby completely negating proof-of-possession security. Profiles MUST specify how confidentiality protection is provided, and additional protection can be applied by encrypting the token, for example encryption of CWTs is specified in section 5.1 of [I-D.ietf-ace-cbor-web-token].

Developers MUST ensure that the ephemeral credentials (i.e., the private key or the session key) are not leaked to third parties. An adversary in possession of the ephemeral credentials bound to the access token will be able to impersonate the client. Be aware that this is a real risk with many constrained environments, since adversaries can often easily get physical access to the devices.

Clients can at any time request a new proof-of-possession capable access token. If clients have that capability, the AS can keep the lifetime of the access token and the associated proof-of-possesion key short and therefore use shorter proof-of-possession key sizes, which translate to a performance benefit for the client and for the resource server. Shorter keys also lead to shorter messages (particularly with asymmetric keying material).

When authorization servers bind symmetric keys to access tokens, they SHOULD scope these access tokens to a specific permissions. Furthermore access tokens using symmetric keys for proof-of-possession SHOULD NOT be targeted at an audience that contains more than one RS, since otherwise any RS in the audience that receives that access token can impersonate the client towards the other members of the audience.

7. Privacy Considerations

Implementers and users should be aware of the privacy implications of the different possible deployments of this framework.

The AS is in a very central position can potentially learn sensitive information about the clients requesting access tokens. If the client credentials grant is used, the AS can track what kind of access the client intends to perform. With other grants this can be prevented by the Resource Owner. To do so the resource owner needs to bind the grants it issues to anonymous, ephemeral credentials, that do not allow the AS to link different grants and thus different access token requests by the same client.

If access tokens are only integrity protected and not encrypted, they may reveal information to attackers listening on the wire, or able to acquire the access tokens in some other way. In the case of CWTs or JWTs the token may e.g. reveal the audience, the scope and the confirmation method used by the client. The latter may reveal the client's identity.

Clients using asymmetric keys for proof-of-possession should be aware of the consequences of using the same key pair for proof-of-possession towards different RS. A set of colluding RS or an attacker able to obtain the access tokens will be able to link the requests, or even to determine the client's identity.

8. IANA Considerations

This specification registers new parameters for OAuth and establishes registries for mappings to CBOR.

8.1. OAuth Introspection Response Parameter Registration

This specification registers the following parameters in the OAuth introspection response parameters

- Name: "cnf"
- Description: Key to prove the right to use an access token, as defined in [RFC7800].
- Change Controller: IESG
- Specification Document(s): this document

- Name: "aud"
- Description: Reference to intended receiving RS, as defined in PoP token specification.
- Change Controller: IESG
- Specification Document(s): this document

- Name: "profile"
- Description: The communication and communication security profile used between client and RS, as defined in ACE profiles.
- Change Controller: IESG
- Specification Document(s): this document

- Name: "client_token"
- Description: Information that the RS MUST pass to the client e.g. about the proof-of-possession keys.
- Change Controller: IESG
- Specification Document(s): this document

- Name: "rs_cnf"
- Description: Describes the public key the RS uses to authenticate.
- Change Controller: IESG
- Specification Document(s): this document

8.2. OAuth Parameter Registration

This specification registers the following parameters in the OAuth Parameters Registry

- Parameter name: "profile"
- Parameter usage location: token request, and token response
- Change Controller: IESG
- Specification Document(s): this document

- Name: "cnf"
- Description: Key to prove the right to use an access token, as defined in [RFC7800].
- Change Controller: IESG
- Specification Document(s): this document

8.3. OAuth Access Token Types

This specification registers the following new token type in the OAuth Access Token Types Registry

- Name: "PoP"
- Description: A proof-of-possession token.
- Change Controller: IESG
- Specification Document(s): this document

8.4. Token Type Mappings

A new registry will be requested from IANA, entitled "Token Type Mappings". The registry is to be created as Expert Review Required.

8.4.1. Registration Template

Token Type:
: Name of token type as registered in the OAuth token type registry e.g. "Bearer".

Mapped value:
: Integer representation for the token type value. The key value MUST be an integer in the range of 1 to 65536.

Change Controller:
: For Standards Track RFCs, list the "IESG". For others, give the name of the responsible party. Other details (e.g., postal address, email address, home page URI) may also be included.

Specification Document(s):
: Reference to the document or documents that specify the parameter,preferably including URIs that can be used to retrieve copies of the documents. An indication of the relevant sections may also be included but is not required.

8.4.2. Initial Registry Contents

- Parameter name: "Bearer"
- Mapped value: 1
- Change Controller: IESG
- Specification Document(s): this document

- Parameter name: "pop"
- Mapped value: 2
- Change Controller: IESG
- Specification Document(s): this document

8.5. CBOR Web Token Claims

This specification registers the following new claims in the CBOR Web Token (CWT) registry:

- Claim Name: "scope"
- Claim Description: The scope of an access token as defined in [RFC6749].
- Change Controller: IESG
- Specification Document(s): this document

- Claim Name: "cnf"
- Claim Description: The proof-of-possession key of an access token as defined in [RFC7800].
- Change Controller: IESG
- Specification Document(s): this document

8.6. ACE Profile Registry

A new registry will be requested from IANA, entitled "ACE Profile Registry". The registry is to be created as Expert Review Required.

8.6.1. Registration Template

Profile name:
: Name of the profile to be included in the profile attribute.

Profile description:
: Text giving an overview of the profile and the context it is developed for.

Profile ID:
: Integer value to identify the profile. The value MUST be an integer in the range of 1 to 65536.

Change Controller:
: For Standards Track RFCs, list the "IESG". For others, give the name of the responsible party. Other details (e.g., postal address, email address, home page URI) may also be included.

Specification Document(s):
: Reference to the document or documents that specify the parameter,preferably including URIs that can be used to retrieve copies of the documents. An indication of the relevant sections may also be included but is not required.

8.7. OAuth Parameter Mappings Registry

A new registry will be requested from IANA, entitled "Token Endpoint CBOR Mappings Registry". The registry is to be created as Expert Review Required.

8.7.1. Registration Template

Parameter name:
: OAuth Parameter name, refers to the name in the OAuth parameter registry e.g. "client_id".

CBOR key value:
: Key value for the claim. The key value MUST be an integer in the range of 1 to 65536.

Change Controller:
: For Standards Track RFCs, list the "IESG". For others, give the name of the responsible party. Other details (e.g., postal address, email address, home page URI) may also be included.

Specification Document(s):
: Reference to the document or documents that specify the parameter,preferably including URIs that can be used to retrieve copies of the documents. An indication of the relevant sections may also be included but is not required.

8.7.2. Initial Registry Contents

- Parameter name: "aud"
- CBOR key value: 3
- Change Controller: IESG
- Specification Document(s): this document

- Parameter name: "client_id"
- CBOR key value: 8
- Change Controller: IESG
- Specification Document(s): this document

- Parameter name: "client_secret"
- CBOR key value: 9
- Change Controller: IESG
- Specification Document(s): this document

- Parameter name: "response_type"
- CBOR key value: 10
- Change Controller: IESG
- Specification Document(s): this document

- Parameter name: "redirect_uri"
- CBOR key value: 11
- Change Controller: IESG
- Specification Document(s): this document

- Parameter name: "scope"
- CBOR key value: 12
- Change Controller: IESG
- Specification Document(s): this document

- Parameter name: "state"
- CBOR key value: 13
- Change Controller: IESG
- Specification Document(s): this document

- Parameter name: "code"
- CBOR key value: 14
- Change Controller: IESG
- Specification Document(s): this document

- Parameter name: "error"
- CBOR key value: 15
- Change Controller: IESG
- Specification Document(s): this document

- Parameter name: "error_description"

- CBOR key value: 16
- Change Controller: IESG
- Specification Document(s): this document

- Parameter name: "error_uri"
- CBOR key value: 17
- Change Controller: IESG
- Specification Document(s): this document

- Parameter name: "grant_type"
- CBOR key value: 18
- Change Controller: IESG
- Specification Document(s): this document

- Parameter name: "access_token"
- CBOR key value: 19
- Change Controller: IESG
- Specification Document(s): this document

- Parameter name: "token_type"
- CBOR key value: 20
- Change Controller: IESG
- Specification Document(s): this document

- Parameter name: "expires_in"
- CBOR key value: 21
- Change Controller: IESG
- Specification Document(s): this document

- Parameter name: "username"
- CBOR key value: 22
- Change Controller: IESG
- Specification Document(s): this document

- Parameter name: "password"
- CBOR key value: 23
- Change Controller: IESG
- Specification Document(s): this document

- Parameter name: "refresh_token"
- CBOR key value: 24
- Change Controller: IESG
- Specification Document(s): this document

- Parameter name: "cnf"
- CBOR key value: 25
- Change Controller: IESG
- Specification Document(s): this document

- Parameter name: "profile"
- CBOR key value: 26
- Change Controller: IESG
- Specification Document(s): this document

8.8. Introspection Endpoint CBOR Mappings Registry

A new registry will be requested from IANA, entitled "Introspection Endpoint CBOR Mappings Registry". The registry is to be created as Expert Review Required.

8.8.1. Registration Template

Response parameter name:
: Name of the response parameter as defined in the "OAuth Token Introspection Response" registry e.g. "active".

CBOR key value:
: Key value for the claim. The key value MUST be an integer in the range of 1 to 65536.

Change Controller:
: For Standards Track RFCs, list the "IESG". For others, give the name of the responsible party. Other details (e.g., postal address, email address, home page URI) may also be included.

Specification Document(s):
: Reference to the document or documents that specify the parameter,preferably including URIs that can be used to retrieve copies of the documents. An indication of the relevant sections may also be included but is not required.

8.8.2. Initial Registry Contents

- Response parameter name: "iss"
- CBOR key value: 1
- Change Controller: IESG
- Specification Document(s): this document

- Response parameter name: "sub"
- CBOR key value: 2
- Change Controller: IESG
- Specification Document(s): this document

- Response parameter name: "aud"
- CBOR key value: 3
- Change Controller: IESG
- Specification Document(s): this document

- Response parameter name: "exp"
- CBOR key value: 4

- Change Controller: IESG
- Specification Document(s): this document

- Response parameter name: "nbf"
- CBOR key value: 5
- Change Controller: IESG
- Specification Document(s): this document

- Response parameter name: "iat"
- CBOR key value: 6
- Change Controller: IESG
- Specification Document(s): this document

- Response parameter name: "cti"
- CBOR key value: 7
- Change Controller: IESG
- Specification Document(s): this document

- Response parameter name: "client_id"
- CBOR key value: 8
- Change Controller: IESG
- Specification Document(s): this document

- Response parameter name: "scope"
- CBOR key value: 12
- Change Controller: IESG
- Specification Document(s): this document

- Response parameter name: "token_type"
- CBOR key value: 20
- Change Controller: IESG
- Specification Document(s): this document

- Response parameter name: "username"
- CBOR key value: 22
- Change Controller: IESG
- Specification Document(s): this document

- Parameter name: "cnf"
- CBOR key value: 25
- Change Controller: IESG
- Specification Document(s): this document

- Parameter name: "profile"
- CBOR key value: 26
- Change Controller: IESG
- Specification Document(s): this document

- Response parameter name: "token"
- CBOR key value: 27
- Change Controller: IESG
- Specification Document(s): this document

- Response parameter name: "token_type_hint"
- CBOR key value: 28
- Change Controller: IESG
- Specification Document(s): this document

- Response parameter name: "active"
- CBOR key value: 29
- Change Controller: IESG
- Specification Document(s): this document

- Response parameter name: "client_token"
- CBOR key value: 30
- Change Controller: IESG
- Specification Document(s): this document

- Response parameter name: "rs_cnf"
- CBOR key value: 31
- Change Controller: IESG
- Specification Document(s): this document

8.9. CoAP Option Number Registration

This section registers the "Access-Token" CoAP Option Number in the "CoRE Parameters" sub-registry "CoAP Option Numbers" in the manner described in [RFC7252].

Name

Access-Token

Number

TBD

Reference

[This document].

Meaning in Request

Contains an Access Token according to [This document] containing access permissions of the client.

Meaning in Response

Not used in response

Safe-to-Forward

Yes

Format

Based on the observer the format is perceived differently. Opaque data to the client and CWT or reference token to the RS.

Length

Less then 255 bytes

8.10. CWT Confirmation Methods Registry

This specification establishes the IANA "CWT Confirmation Methods" registry for CWT "cnf" member values. The registry records the confirmation method member and a reference to the specification that defines it.

8.10.1. Registration Template

Confirmation Method Name:
The name requested (e.g., "kid"). This name is intended to be human readable and be used for debugging purposes. It is case sensitive. Names may not match other registered names in a case-insensitive manner unless the Designated Experts state that there is a compelling reason to allow an exception.

Confirmation Method Value:
Integer representation for the confirmation method value. Intended for use to uniquely identify the confirmation method. The value MUST be an integer in the range of 1 to 65536.

Confirmation Method Description:
Brief description of the confirmation method (e.g. "Key Identifier").

Change Controller:
For Standards Track RFCs, list the "IESG". For others, give the name of the responsible party. Other details (e.g., postal address, email address, home page URI) may also be included.

Specification Document(s):
Reference to the document or documents that specify the parameter, preferably including URIs that can be used to retrieve copies of the documents. An indication of the relevant sections may also be included but is not required.

8.10.2. Initial Registry Contents

- Confirmation Method Name: "COSE_Key"
- Confirmation Method Value: 1
- Confirmation Method Description: A COSE_Key that is either a public key or a symmetric key.
- Change Controller: IESG
- Specification Document(s): this document

- Confirmation Method Name: "COSE_Encrypted"
- Confirmation Method Value: 2
- Confirmation Method Description: A COSE_Encrypted structure that wraps a COSE_Key containing a symmetric key.
- Change Controller: IESG
- Specification Document(s): this document

- Confirmation Method Name: "Key Identifier"
- Confirmation Method Value: 3
- Confirmation Method Description: A key identifier.
- Change Controller: IESG
- Specification Document(s): this document

9. Acknowledgments

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Appendix A. Design Justification

This section provides further insight into the design decisions of the solution documented in this document. Section 3 lists several building blocks and briefly summarizes their importance. The justification for offering some of those building blocks, as opposed to using OAuth 2.0 as is, is given below.

Common IoT constraints are:

Low Power Radio:

Many IoT devices are equipped with a small battery which needs to last for a long time. For many constrained wireless devices the highest energy cost is associated to transmitting or receiving messages. It is therefore important to keep the total communication overhead low, including minimizing the number and size of messages sent and received, which has an impact of choice on the message format and protocol. By using CoAP over UDP, and CBOR encoded messages some of these aspects are addressed. Security protocols contribute to the communication overhead and can in some cases be optimized. For example authentication and

key establishment may in certain cases where security requirements so allows be replaced by provisioning of security context by a trusted third party, using transport or application layer security.

Low CPU Speed:

Some IoT devices are equipped with processors that are significantly slower than those found in most current devices on the Internet. This typically has implications on what timely cryptographic operations a device is capable to perform, which in turn impacts e.g. protocol latency. Symmetric key cryptography may be used instead of the computationally more expensive public key cryptography where the security requirements so allows, but this may also require support for trusted third party assisted secret key establishment using transport or application layer security.

Small Amount of Memory:

Microcontrollers embedded in IoT devices are often equipped with small amount of RAM and flash memory, which places limitations what kind of processing can be performed and how much code can be put on those devices. To reduce code size fewer and smaller protocol implementations can be put on the firmware of such a device. In this case, CoAP may be used instead of HTTP, symmetric key cryptography instead of public key cryptography, and CBOR instead of JSON. Authentication and key establishment protocol, e.g. the DTLS handshake, in comparison with assisted key establishment also has an impact on memory and code.

User Interface Limitations:

Protecting access to resources is both an important security as well as privacy feature. End users and enterprise customers do not want to give access to the data collected by their IoT device or to functions it may offer to third parties. Since the classical approach of requesting permissions from end users via a rich user interface does not work in many IoT deployment scenarios these functions need to be delegated to user controlled devices that are better suitable for such tasks, such as smart phones and tablets.

Communication Constraints:

In certain constrained settings an IoT device may not be able to communicate with a given device at all times. Devices may be sleeping, or just disconnected from the Internet because of general lack of connectivity in the area, for cost reasons, or for

security reasons, e.g. to avoid an entry point for Denial-of-Service attacks.

The communication interactions this framework builds upon (as shown graphically in Figure 1) may be accomplished using a variety of different protocols, and not all parts of the message flow are used in all applications due to the communication constraints. While we envision deployments to make use of CoAP we explicitly want to support HTTP, HTTP/2 or specific protocols, such as Bluetooth Smart communication, which does not necessarily use IP. The latter raises the need for application layer security over the various interfaces.

Appendix B. Roles and Responsibilities

Resource Owner

* Make sure that the RS is registered at the AS. This includes making known to the AS which profiles, token_types, scopes, and key types (symmetric/asymmetric) the RS supports. Also making it known to the AS which audience(s) the RS identifies itself with.
* Make sure that clients can discover the AS which is in charge of the RS.
* If the client-credentials grant is used, make sure that the AS has the necessary, up-to-date, access control policies for the RS.

Requesting Party

* Make sure that the client is provisioned the necessary credentials to authenticate to the AS.
* Make sure that the client is configured to follow the security requirements of the Requesting Party, when issuing requests (e.g. minimum communication security requirements, trust anchors).
* Register the client at the AS. This includes making known to the AS which profiles, token_types, and key types (symmetric/asymmetric) the client.

Authorization Server

* Register RS and manage corresponding security contexts.
* Register clients and including authentication credentials.
* Allow Resource Owners to configure and update access control policies related to their registered RS’
* Expose the /token endpoint to allow clients to request tokens.
* Authenticate clients that wish to request a token.

* Process a token request against the authorization policies configured for the RS.
* Optionally: Expose the /introspection endpoint that allows RS's to submit token introspection requests.
* If providing an introspection endpoint: Authenticate RS's that wish to get an introspection response.
* If providing an introspection endpoint: Process token introspection requests.
* Optionally: Handle token revocation.

Client

* Discover the AS in charge of the RS that is to be targeted with a request.
* Submit the token request (A).
  + Authenticate towards the AS.
  + Optionally (if not pre-configured): Specify which RS, which resource(s), and which action(s) the request(s) will target.
  + If raw public key (rpk) or certificate is used, make sure the AS has the right rpk or certificate for this client.
* Process the access token and RS Information (B)
  + Check that the RS Information provides the necessary security parameters (e.g. PoP key, information on communication security protocols supported by the RS).
* Send the token and request to the RS (C)
  + Authenticate towards the RS (this could coincide with the proof of possession process).
  + Transmit the token as specified by the AS (default is to the /authz-info endpoint, alternative options are specified by profiles).
  + Perform the proof-of-possession procedure as specified by the profile in use (this may already have been taken care of through the authentication procedure).
* Process the RS response (F) requirements of the Requesting Party, when issuing requests (e.g. minimum communication security requirements, trust anchors).
* Register the client at the AS.

Resource Server

* Expose a way to submit access tokens. By default this is the /authz-info endpoint.
* Process an access token.
  + Verify the token is from the right AS.

    + Verify that the token applies to this RS.
    + Check that the token has not expired (if the token provides expiration information).
    + Check the token's integrity.
    + Store the token so that it can be retrieved in the context of a matching request.
  * Process a request.

    + Set up communication security with the client.
    + Authenticate the client.
    + Match the client against existing tokens.
    + Check that tokens belonging to the client actually authorize the requested action.
    + Optionally: Check that the matching tokens are still valid, using introspection (if this is possible.)
  * Send a response following the agreed upon communication security.

Appendix C. Requirements on Profiles

This section lists the requirements on profiles of this framework, for the convenience of a profile designer.

- Optionally Specify the discovery process of how the client finds the right AS for an RS it wants to send a request to. Section 4
- Specify the communication protocol the client and RS the must use (e.g. CoAP). Section 5 and Section 5.5.4.4
- Specify the security protocol the client and RS must use to protect their communication (e.g. OSCOAP or DTLS over CoAP). This must provide encryption and integrity protection. Section 5.5.4.4
- Specify how the client and the RS mutually authenticate. Section 4
- Specify the Content-format of the protocol messages (e.g. "application/cbor" or "application/cose+cbor"). Section 4
- Specify the proof-of-possession protocol(s) and how to select one, if several are available. Also specify which key types (e.g. symmetric/asymmetric) are supported by a specific proof-of-possession protocol. Section 5.5.4.3
- Specify a unique profile identifier. Section 5.5.4.4
- Optionally specify how the RS talks to the AS for introspection.Section 5.6
- Optionally specify how the client talks to the AS for requesting a token. Section 5.5
- Specify how/if the /authz-info endpoint is protected. Section 5.7.1
- Optionally define other methods of token transport than the /authz-info endpoint. Section 5.7.1

Appendix D. Assumptions on AS knowledge about C and RS

This section lists the assumptions on what an AS should know about a client and a RS in order to be able to respond to requests to the /token and /introspect endpoints. How this information is established is out of scope for this document.

- The identifier of the client or RS.
- The profiles that the client or RS supports.
- The scopes that the RS supports.
- The audiences that the RS identifies with.
- The key types (e.g. pre-shared symmetric key, raw public key, key length, other key parameters) that the client or RS supports.
- The types of access tokens the RS supports (e.g. CWT).
- If the RS supports CWTs, the COSE parameters for the crypto wrapper (e.g. algorithm, key-wrap algorithm, key-length).
- The expiration time for access tokens issued to this RS (unless the RS accepts a default time chosen by the AS).
- The symmetric key shared between client or RS and AS (if any).
- The raw public key of the client or RS (if any).

Appendix E. Deployment Examples

There is a large variety of IoT deployments, as is indicated in Appendix A, and this section highlights a few common variants. This section is not normative but illustrates how the framework can be applied.

For each of the deployment variants there are a number of possible security setups between clients, resource servers and authorization servers. The main focus in the following subsections is on how authorization of a client request for a resource hosted by a RS is performed. This requires the the security of the requests and responses between the clients and the RS to consider.

Note: CBOR diagnostic notation is used for examples of requests and responses.

E.1. Local Token Validation

In this scenario we consider the case where the resource server is offline, i.e. it is not connected to the AS at the time of the access request. This access procedure involves steps A, B, C, and F of Figure 1.

Since the resource server must be able to verify the access token locally, self-contained access tokens must be used.

This example shows the interactions between a client, the authorization server and a temperature sensor acting as a resource server. Message exchanges A and B are shown in Figure 18.

A: The client first generates a public-private key pair used for communication security with the RS.
The client sends the POST request to /token at the AS. The security of this request can be transport or application layer, it is up the the communication security profile to define. In the example transport layer identification of the AS is done and the client identifies with client_id and client_secret as in classic OAuth. The request contains the public key of the client and the Audience parameter set to "tempSensorInLivingRoom", a value that the temperature sensor identifies itself with. The AS evaluates the request and authorizes the client to access the resource.
B: The AS responds with a PoP token and RS Information. The PoP token contains the public key of the client, and the RS Information contains the public key of the RS. For communication security this example uses DTLS RawPublicKey between the client and the RS. The issued token will have a short validity time, i.e. 'exp' close to 'iat', to protect the RS from replay attacks. The token includes the claim such as "scope" with the authorized access that an owner of the temperature device can enjoy. In this example, the 'scope' claim, issued by the AS, informs the RS that the owner of the token, that can prove the possession of a key is authorized to make a GET request against the /temperature resource and a POST request on the /firmware resource. Note that the syntax and semantics of the scope claim are application specific.
Note: In this example we assume that the client knows what resource it wants to access, and is therefore able to request specific audience and scope claims for the access token.

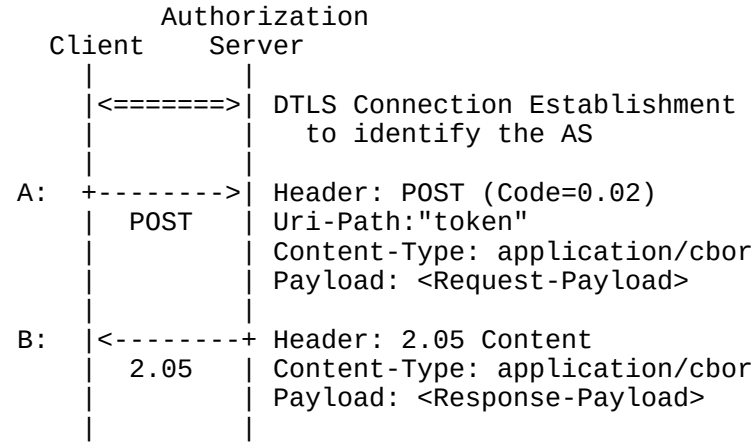

Figure 18: Token Request and Response Using Client Credentials.

The information contained in the Request-Payload and the Response-Payload is shown in Figure 19. Note that we assume a DTLS-based communication security profile for this example, therefore the Content-Type is "application/cbor".

```
Request-Payload :
{
  "grant_type" : "client_credentials",
  "aud" : "tempSensorInLivingRoom",
  "client_id" : "myclient",
  "client_secret" : "qwerty"
}

Response-Payload :
{
  "access_token" : b64'SlAV32hkKG ...',
  "token_type" : "pop",
  "csp" : "DTLS",
  "cnf" : {
    "COSE_Key" : {
      "kid" : b64'c29tZSBwdWJsaWMga2V5IGlk',
      "kty" : "EC",
      "crv" : "P-256",
      "x"   : b64'MKBCTNIcKUSDii11ySs3526iDZ8AiTo7Tu6KPAqv7D4',
      "y"   : b64'4Etl6SRW2YiLUrN5vfvVHuhp7x8PxltmWWlbbM4IFyM'
    }
  }
}
```

Figure 19: Request and Response Payload Details.

The content of the access token is shown in Figure 20.

```
{
  "aud" : "tempSensorInLivingRoom",
  "iat" : "1360189224",
  "exp" : "1360289224",
  "scope" :  "temperature_g firmware_p",
  "cnf" : {
    "jwk" : {
      "kid" : b64'1Bg8vub9tLe1gHMzV76e8',
      "kty" : "EC",
      "crv" : "P-256",
      "x" : b64'f83OJ3D2xF1Bg8vub9tLe1gHMzV76e8Tus9uPHvRVEU',
      "y" : b64'x_FEzRu9m36HLN_tue659LNpXW6pCyStikYjKIWI5a0'
    }
  }
}
```

Figure 20: Access Token including Public Key of the Client.

Messages C and F are shown in Figure 21 - Figure 22.

C: The client then sends the PoP token to the /authz-info endpoint at the RS. This is a plain CoAP request, i.e. no transport or application layer security between client and RS, since the token is integrity protected between AS and RS. The RS verifies that the PoP token was created by a known and trusted AS, is valid, and responds to the client. The RS caches the security context together with authorization information about this client contained in the PoP token.

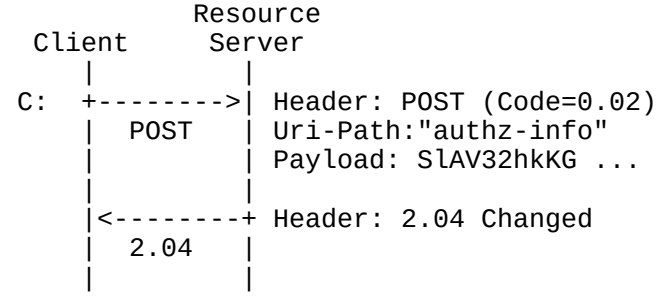

Figure 21: Access Token provisioning to RS

The client and the RS runs the DTLS handshake using the raw public keys established in step B and C.

The client sends the CoAP request GET to /temperature on RS over DTLS. The RS verifies that the request is authorized, based on previously established security context.
F: The RS responds with a resource representation over DTLS.

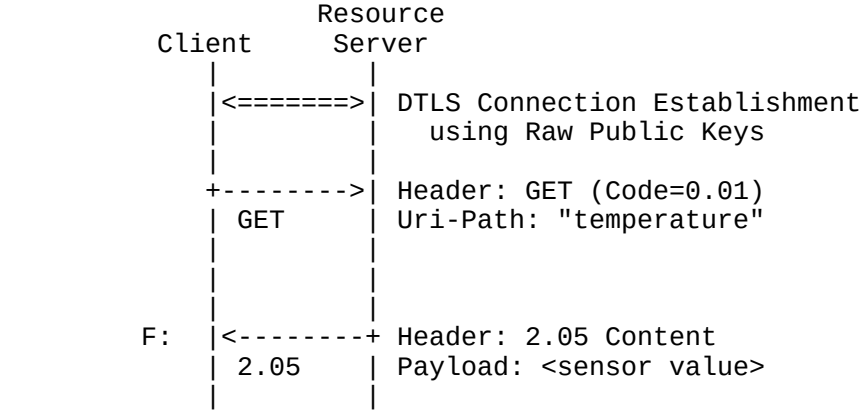

Figure 22: Resource Request and Response protected by DTLS.

E.2. Introspection Aided Token Validation

In this deployment scenario we assume that a client is not able to access the AS at the time of the access request. Since the RS is, however, connected to the back-end infrastructure it can make use of token introspection. This access procedure involves steps A-F of Figure 1, but assumes steps A and B have been carried out during a phase when the client had connectivity to AS.

Since the client is assumed to be offline, at least for a certain period of time, a pre-provisioned access token has to be long-lived. The resource server may use its online connectivity to validate the access token with the authorization server, which is shown in the example below.

In the example interactions between an offline client (key fob), a RS (online lock), and an AS is shown. We assume that there is a provisioning step where the client has access to the AS. This corresponds to message exchanges A and B which are shown in Figure 23.

Authorization consent from the resource owner can be pre-configured, but it can also be provided via an interactive flow with the resource owner. An example of this for the key fob case could be that the resource owner has a connected car, he buys a generic key that he wants to use with the car. To authorize the key fob he connects it

to his computer that then provides the UI for the device. After that OAuth 2.0 implicit flow can used to authorize the key for his car at the the car manufacturers AS.

Note: In this example the client does not know the exact door it will be used to access since the token request is not send at the time of access. So the scope and audience parameters is set quite wide to start with and new values different form the original once can be returned from introspection later on.

A: The client sends the request using POST to /token at AS. The request contains the Audience parameter set to "PACS1337" (PACS, Physical Access System), a value the that the online door in question identifies itself with. The AS generates an access token as on opaque string, which it can match to the specific client, a targeted audience and a symmetric key. The security is provided by identifying the AS on transport layer using a pre shared security context (psk, rpk or certificate) and then the client is identified using client_id and client_secret as in classic OAuth
B: The AS responds with the an access token and RS Information, the latter containing a symmetric key. Communication security between C and RS will be DTLS and PreSharedKey. The PoP key being used as the PreSharedKey.

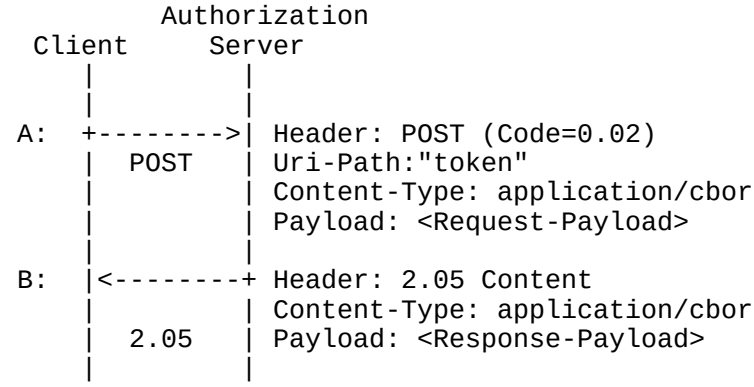

Figure 23: Token Request and Response using Client Credentials.

The information contained in the Request-Payload and the Response-Payload is shown in Figure 24.

```
Request-Payload:
{
  "grant_type" : "client_credentials",
  "aud" : "lockOfDoor4711",
  "client_id" : "keyfob",
  "client_secret" : "qwerty"
}

Response-Payload:
{
  "access_token" : b64'SlAV32hkKG ...'
  "token_type" : "pop",
  "csp" : "DTLS",
  "cnf" : {
    "COSE_Key" : {
      "kid" : b64'c29tZSBwdWJsaWMga2V5IGlk',
      "kty" : "oct",
      "alg" : "HS256",
      "k": b64'ZoRSOrFzN_FzUA5XKMYoVHyzff5oRJxl-IXRtztJ6uE'
    }
  }
}
```

Figure 24: Request and Response Payload for C offline

The access token in this case is just an opaque string referencing the authorization information at the AS.

C: Next, the client POSTs the access token to the /authz-info endpoint in the RS. This is a plain CoAP request, i.e. no DTLS between client and RS. Since the token is an opaque string, the RS cannot verify it on its own, and thus defers to respond the client with a status code until after step E.
D: The RS forwards the token to the /introspect endpoint on the AS. Introspection assumes a secure connection between the AS and the RS, e.g. using transport of application layer security. In the example AS is identified using pre shared security context (psk, rpk or certificate) while RS is acting as client and is identified with client_id and client_secret.
E: The AS provides the introspection response containing parameters about the token. This includes the confirmation key (cnf) parameter that allows the RS to verify the client's proof of possession in step F.
After receiving message E, the RS responds to the client's POST in step C with the CoAP response code 2.01 (Created).

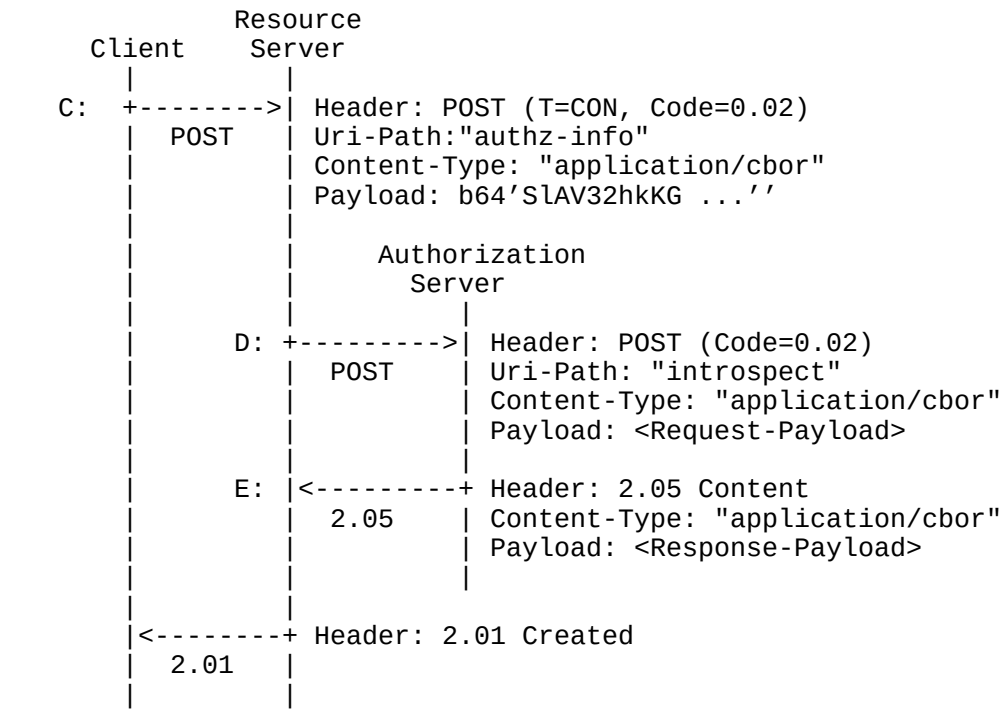

Figure 25: Token Introspection for C offline

The information contained in the Request-Payload and the Response-Payload is shown in Figure 26.

```
Request-Payload:
{
  "token" : b64'SlAV32hkKG...',
  "client_id" : "FrontDoor",
  "client_secret" : "ytrewq"
}

Response-Payload:
{
  "active" : true,
  "aud" : "lockOfDoor4711",
  "scope" : "open, close",
  "iat" : 1311280970,
  "cnf" : {
    "kid" : b64'JDLUhTMjU2IiwiY3R5Ijoi ...'
  }
}
```

Figure 26: Request and Response Payload for Introspection

The client uses the symmetric PoP key to establish a DTLS PreSharedKey secure connection to the RS. The CoAP request PUT is sent to the uri-path /state on RS changing state of the door to locked.
F: The RS responds with a appropriate over the secure DTLS channel.

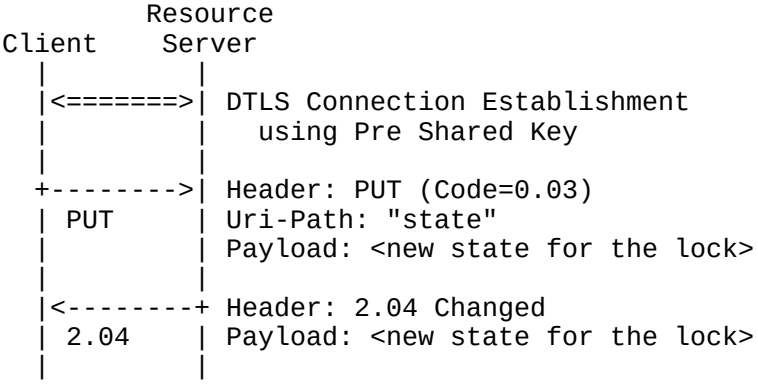

Figure 27: Resource request and response protected by OSCOAP

Appendix F. Document Updates

F.1. Version -06 to -07

- Various clarifications added.
- Fixed erroneous author email.

F.2. Version -05 to -06

- Moved sections that define the ACE framework into a subsection of the framework Section 5.
- Split section on client credentials and grant into two separate sections, Section 5.1, and Section 5.2.
- Added Section 5.3 on AS authentication.
- Added Section 5.4 on the Authorize endpoint.

F.3. Version -04 to -05

- Added RFC 2119 language to the specification of the required behavior of profile specifications.
- Added Section 5.2 on the relation to the OAuth2 grant types.
- Added CBOR abbreviations for error and the error codes defined in OAuth2.
- Added clarification about token expiration and long-running requests in Section 5.7.2

- Added security considerations about tokens with symmetric pop keys valid for more than one RS.
- Added privacy considerations section.
- Added IANA registry mapping the confirmation types from RFC 7800 to equivalent COSE types.
- Added appendix D, describing assumptions about what the AS knows about the client and the RS.

F.4. Version -03 to -04

- Added a description of the terms "framework" and "profiles" as used in this document.
- Clarified protection of access tokens in section 3.1.
- Clarified uses of the 'cnf' parameter in section 6.4.5.
- Clarified intended use of Client Token in section 7.4.

F.5. Version -02 to -03

- Removed references to draft-ietf-oauth-pop-key-distribution since the status of this draft is unclear.
- Copied and adapted security considerations from draft-ietf-oauth-pop-key-distribution.
- Renamed "client information" to "RS information" since it is information about the RS.
- Clarified the requirements on profiles of this framework.
- Clarified the token endpoint protocol and removed negotiation of 'profile' and 'alg' (section 6).
- Renumbered the abbreviations for claims and parameters to get a consistent numbering across different endpoints.
- Clarified the introspection endpoint.
- Renamed token, introspection and authz-info to 'endpoint' instead of 'resource' to mirror the OAuth 2.0 terminology.
- Updated the examples in the appendices.

F.6. Version -01 to -02

- Restructured to remove communication security parts. These shall now be defined in profiles.
- Restructured section 5 to create new sections on the OAuth endpoints /token, /introspect and /authz-info.
- Pulled in material from draft-ietf-oauth-pop-key-distribution in order to define proof-of-possession key distribution.
- Introduced the 'cnf' parameter as defined in RFC7800 to reference or transport keys used for proof of possession.
- Introduced the 'client-token' to transport client information from the AS to the client via the RS in conjunction with introspection.
- Expanded the IANA section to define parameters for token request, introspection and CWT claims.

- Moved deployment scenarios to the appendix as examples.

F.7. Version -00 to -01

- Changed 5.1. from "Communication Security Protocol" to "Client Information".
- Major rewrite of 5.1 to clarify the information exchanged between C and AS in the PoP token request profile for IoT.
  - Allow the client to indicate preferences for the communication security protocol.
  - Defined the term "Client Information" for the additional information returned to the client in addition to the access token.
  - Require that the messages between AS and client are secured, either with (D)TLS or with COSE_Encrypted wrappers.
  - Removed dependency on OSCOAP and added generic text about object security instead.
  - Defined the "rpk" parameter in the client information to transmit the raw public key of the RS from AS to client.
  - (D)TLS MUST use the PoP key in the handshake (either as PSK or as client RPK with client authentication).
  - Defined the use of x5c, x5t and x5tS256 parameters when a client certificate is used for proof of possession.
  - Defined "tktn" parameter for signaling for how to transfer the access token.
- Added 5.2. the CoAP Access-Token option for transferring access tokens in messages that do not have payload.
- 5.3.2. Defined success and error responses from the RS when receiving an access token.
- 5.6.:Added section giving guidance on how to handle token expiration in the absence of reliable time.
- Appendix B Added list of roles and responsibilities for C, AS and RS.

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June 30, 2017

Proof-of-Possession Key Semantics for CBOR Web Tokens (CWTs)
draft-jones-ace-cwt-proof-of-possession-01

Abstract

This specification describes how to declare in a CBOR Web Token (CWT) that the presenter of the CWT possesses a particular proof-of-possession key. Being able to prove possession of a key is also sometimes described as the presenter being a holder-of-key. This specification provides equivalent functionality to "Proof-of-Possession Key Semantics for JSON Web Tokens (JWTs)" (RFC 7800), but using CBOR and CWTs rather than JSON and JWTs.

Table of Contents

1. Introduction

This specification describes how a CBOR Web Token [CWT] can declare that the presenter of the CWT possesses a particular proof-of-possession (PoP) key. Proof of possession of a key is also sometimes

described as the presenter being a holder-of-key. This specification provides equivalent functionality to "Proof-of-Possession Key Semantics for JSON Web Tokens (JWTs)" [RFC7800], but using CBOR [RFC7049] and CWTs [CWT] rather than JSON [RFC7159] and JWTs [JWT].

1.1. Notational Conventions

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

Unless otherwise noted, all the protocol parameter names and values are case sensitive.

2. Terminology

This specification uses terms defined in the CBOR Web Token [CWT], [I-D.ietf-cose-msg], and Concise Binary Object Representation (CBOR) [RFC7049] specifications.

These terms are defined by this specification:

Issuer
: Party that creates the CWT and binds the proof-of-possession key to it.

Presenter
: Party that proves possession of a private key (for asymmetric key cryptography) or secret key (for symmetric key cryptography) to a recipient.

Recipient
: Party that receives the CWT containing the proof-of-possession key information from the presenter.

3. Representations for Proof-of-Possession Keys

By including a "cnf" (confirmation) claim in a CWT, the issuer of the CWT declares that the presenter possesses a particular key and that the recipient can cryptographically confirm that the presenter has possession of that key. The value of the "cnf" claim is a CBOR map and the members of that map identify the proof-of-possession key.

The presenter can be identified in one of several ways by the CWT depending upon the application requirements. If the CWT contains a "sub" (subject) claim [CWT], the presenter is normally the subject identified by the CWT. (In some applications, the subject identifier

will be relative to the issuer identified by the "iss" (issuer) claim [CWT].) If the CWT contains no "sub" claim, the presenter is normally the issuer identified by the CWT using the "iss" claim. The case in which the presenter is the subject of the CWT is analogous to Security Assertion Markup Language (SAML) 2.0 [OASIS.saml-core-2.0-os] SubjectConfirmation usage. At least one of the "sub" and "iss" claims is typically present in the CWT and some use cases may require that both be present.

3.1. Confirmation Claim

The "cnf" claim is used in the CWT to contain members used to identify the proof-of-possession key. Other members of the "cnf" map may be defined because a proof-of-possession key may not be the only means of confirming the authenticity of the token. This is analogous to the SAML 2.0 [OASIS.saml-core-2.0-os] SubjectConfirmation element in which a number of different subject confirmation methods can be included (including proof-of-possession key information).

The set of confirmation members that a CWT must contain to be considered valid is context dependent and is outside the scope of this specification. Specific applications of CWTs will require implementations to understand and process some confirmation members in particular ways. However, in the absence of such requirements, all confirmation members that are not understood by implementations MUST be ignored.

This specification establishes the IANA "CWT Confirmation Methods" registry for these members in Section 6.2 and registers the members defined by this specification. Other specifications can register other members used for confirmation, including other members for conveying proof-of-possession keys using different key representations.

The "cnf" claim value MUST represent only a single proof-of-possession key; thus, at most one of the "COSE_Key" and "Encrypted_COSE_Key" confirmation values defined below may be present. Note that if an application needs to represent multiple proof-of-possession keys in the same CWT, one way for it to achieve this is to use other claim names, in addition to "cnf", to hold the additional proof-of-possession key information. These claims could use the same syntax and semantics as the "cnf" claim. Those claims would be defined by applications or other specifications and could be registered in the IANA "CBOR Web Token Claims" registry [IANA.CWT.Claims].

3.2. Representation of an Asymmetric Proof-of-Possession Key

When the key held by the presenter is an asymmetric private key, the "COSE_Key" member is a COSE_Key [I-D.ietf-cose-msg] representing the corresponding asymmetric public key. The following example (using JSON notation) demonstrates such a declaration in the CWT Claims Set of a CWT:

```
{
 "iss": "https://server.example.com",
 "aud": "https://client.example.org",
 "exp": 1361398824,
 "cnf":{
   "COSE_Key":{
     "kty": "EC",
     "crv": "P-256",
     "x": "18wHLeIgW9wVN6VD1Txgpqy2LszYkMf6J8njVAibvhM",
     "y": "-V4dS4UaLMgP_4fY4j8ir7cl1TXlFdAgcx55o7TkcSA"
    }
  }
}
```

The COSE_Key MUST contain the required key members for a COSE_Key of that key type and MAY contain other COSE_Key members, including the "kid" (Key ID) member.

The "COSE_Key" member MAY also be used for a COSE_Key representing a symmetric key, provided that the CWT is encrypted so that the key is not revealed to unintended parties. The means of encrypting a CWT is explained in [CWT]. If the CWT is not encrypted, the symmetric key MUST be encrypted as described below.

3.3. Representation of an Encrypted Symmetric Proof-of-Possession Key

When the key held by the presenter is a symmetric key, the "Encrypted_COSE_Key" member is an encrypted COSE_Key [I-D.ietf-cose-msg] representing the symmetric key encrypted to a key known to the recipient using COSE_Encrypt or COSE_Encrypt0.

The following example (using JSON notation) illustrates a symmetric key that could subsequently be encrypted for use in the "Encrypted_COSE_Key" member:

```
{
 "kty": "oct",
 "alg": "HS256",
 "k": "ZoRSOrFzN_FzUA5XKMYoVHyzff5oRJxl-IXRtztJ6uE"
}
```

The COSE_Key representation is used as the plaintext when encrypting the key. The COSE_Key could, for instance, be encrypted using a COSE_Encrypt0 representation using the AES-CCM-16-64-128 algorithm.

The following example CWT Claims Set of a CWT (using JSON notation) illustrates the use of an encrypted symmetric key as the "Encrypted_COSE_Key" member value:

```
{
 "iss": "https://server.example.com",
 "sub": "24400320",
 "aud": "s6BhdRkqt3",
 "exp": 1311281970,
 "iat": 1311280970,
 "cnf":{
   "Encrypted_COSE_Key":
     "(TBD)"
   }
}
```

3.4. Representation of a Key ID for a Proof-of-Possession Key

The proof-of-possession key can also be identified by the use of a Key ID instead of communicating the actual key, provided the recipient is able to obtain the identified key using the Key ID. In this case, the issuer of a CWT declares that the presenter possesses a particular key and that the recipient can cryptographically confirm proof of possession of the key by the presenter by including a "cnf" claim in the CWT whose value is a CBOR map with the CBOR map containing a "kid" member identifying the key.

The following example (using JSON notation) demonstrates such a declaration in the CWT Claims Set of a CWT:

```
{
 "iss": "https://server.example.com",
 "aud": "https://client.example.org",
 "exp": 1361398824,
 "cnf":{
   "kid": "dfd1aa97-6d8d-4575-a0fe-34b96de2bfad"
  }
}
```

The content of the "kid" value is application specific. For instance, some applications may choose to use a cryptographic hash of the public key value as the "kid" value.

3.5. Specifics Intentionally Not Specified

Proof of possession is typically demonstrated by having the presenter sign a value determined by the recipient using the key possessed by the presenter. This value is sometimes called a "nonce" or a "challenge".

The means of communicating the nonce and the nature of its contents are intentionally not described in this specification, as different protocols will communicate this information in different ways. Likewise, the means of communicating the signed nonce is also not specified, as this is also protocol specific.

Note that another means of proving possession of the key when it is a symmetric key is to encrypt the key to the recipient. The means of obtaining a key for the recipient is likewise protocol specific.

4. Security Considerations

All of the security considerations that are discussed in [CWT] also apply here. In addition, proof of possession introduces its own unique security issues. Possessing a key is only valuable if it is kept secret. Appropriate means must be used to ensure that unintended parties do not learn private key or symmetric key values.

Applications utilizing proof of possession should also utilize audience restriction, as described in Section 4.1.3 of [JWT], as it provides different protections. Proof of possession can be used by recipients to reject messages from unauthorized senders. Audience restriction can be used by recipients to reject messages intended for different recipients.

A recipient might not understand the "cnf" claim. Applications that require the proof-of-possession keys communicated with it to be understood and processed must ensure that the parts of this specification that they use are implemented.

Proof of possession via encrypted symmetric secrets is subject to replay attacks. This attack can, for example, be avoided when a signed nonce or challenge is used since the recipient can use a distinct nonce or challenge for each interaction. Replay can also be avoided if a sub-key is derived from a shared secret that is specific to the instance of the PoP demonstration.

As is the case with other information included in a CWT, it is necessary to apply data origin authentication and integrity protection (via a keyed message digest or a digital signature). Data origin authentication ensures that the recipient of the CWT learns

about the entity that created the CWT since this will be important for any policy decisions. Integrity protection prevents an adversary from changing any elements conveyed within the CWT payload. Special care has to be applied when carrying symmetric keys inside the CWT since those not only require integrity protection but also confidentiality protection.

5. Privacy Considerations

A proof-of-possession key can be used as a correlation handle if the same key is used with multiple parties. Thus, for privacy reasons, it is recommended that different proof-of-possession keys be used when interacting with different parties.

6. IANA Considerations

The following registration procedure is used for all the registries established by this specification.

Values are registered on a Specification Required [RFC5226] basis after a three-week review period on the cwt-reg-review@ietf.org mailing list, on the advice of one or more Designated Experts. However, to allow for the allocation of values prior to publication, the Designated Experts may approve registration once they are satisfied that such a specification will be published. [[Note to the RFC Editor: The name of the mailing list should be determined in consultation with the IESG and IANA. Suggested name: cwt-reg-review@ietf.org.]]

Registration requests sent to the mailing list for review should use an appropriate subject (e.g., "Request to Register CWT Confirmation Method: example"). Registration requests that are undetermined for a period longer than 21 days can be brought to the IESG's attention (using the iesg@ietf.org mailing list) for resolution.

Criteria that should be applied by the Designated Experts include determining whether the proposed registration duplicates existing functionality, determining whether it is likely to be of general applicability or whether it is useful only for a single application, and evaluating the security properties of the item being registered and whether the registration makes sense.

It is suggested that multiple Designated Experts be appointed who are able to represent the perspectives of different applications using this specification in order to enable broadly informed review of registration decisions. In cases where a registration decision could be perceived as creating a conflict of interest for a particular

Expert, that Expert should defer to the judgment of the other Experts.

6.1. CBOR Web Token Claims Registration

This specification registers the "cnf" claim in the IANA "CBOR Web Token Claims" registry [IANA.CWT.Claims] established by [CWT].

6.1.1. Registry Contents

- Claim Name: "cnf"
- Claim Description: Confirmation
- JWT Claim Name: "cnf"
- Claim Key: TBD (maybe 8)
- Claim Value Type(s): map
- Change Controller: IESG
- Specification Document(s): Section 3.1 of [[this document]]

6.2. CWT Confirmation Methods Registry

This specification establishes the IANA "CWT Confirmation Methods" registry for CWT "cnf" member values. The registry records the confirmation method member and a reference to the specification that defines it.

6.2.1. Registration Template

Confirmation Method Name:
: The human-readable name requested (e.g., "kid").

Confirmation Method Description:
: Brief description of the confirmation method (e.g., "Key Identifier").

JWT Confirmation Method Name:
: Claim Name of the equivalent JWT confirmation method value, as registered in [IANA.JWT.Claims]. CWT claims should normally have a corresponding JWT claim. If a corresponding JWT claim would not make sense, the Designated Experts can choose to accept registrations for which the JWT Claim Name is listed as "N/A".

Confirmation Key:
: CBOR map key value for the confirmation method.

Confirmation Value Type(s):
: CBOR types that can be used for the confirmation method value.

Change Controller:

For Standards Track RFCs, list the "IESG". For others, give the name of the responsible party. Other details (e.g., postal address, email address, home page URI) may also be included.

Specification Document(s):
Reference to the document or documents that specify the parameter, preferably including URIs that can be used to retrieve copies of the documents. An indication of the relevant sections may also be included but is not required.

6.2.2. Initial Registry Contents

- Confirmation Method Name: "COSE_Key"
- Confirmation Method Description: COSE_Key Representing Public Key
- JWT Confirmation Method Name: "jwk"
- Confirmation Key: 1
- Confirmation Value Type(s): map
- Change Controller: IESG
- Specification Document(s): Section 3.2 of [[this document]]

- Confirmation Method Name: "Encrypted_COSE_Key"
- Confirmation Method Description: Encrypted COSE_Key
- JWT Confirmation Method Name: "jwe"
- Confirmation Key: 2
- Confirmation Value Type(s): array (with an optional COSE_Encrypt or COSE_Encrypt0 tag)
- Change Controller: IESG
- Specification Document(s): Section 3.3 of [[this document]]

- Confirmation Method Name: "kid"
- Confirmation Method Description: Key Identifier
- JWT Confirmation Method Name: "kid"
- Confirmation Key: 3
- Confirmation Value Type(s): binary string
- Change Controller: IESG
- Specification Document(s): Section 3.4 of [[this document]]

7. References

7.1. Normative References

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7.2. Informative References

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Acknowledgements

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Open Issues

- Convert the examples from JSON/JWT to CBOR/CWT.

Document History

[[to be removed by the RFC Editor before publication as an RFC]]

-01

- Tracked CBOR Web Token (CWT) Claims Registry updates.
- Addressed review comments by Michael Richardson and Jim Schaad.
- Added co-authors.

-00

- Created the initial draft from RFC 7800.

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OSCOAP profile of the Authentication and Authorization for Constrained Environments Framework
draft-seitz-ace-oscoap-profile-04

Abstract

This memo specifies a profile for the Authentication and Authorization for Constrained Environments (ACE) framework. It utilizes Object Security of CoAP (OSCOAP) and Ephemeral Diffie-Hellman over COSE (EDHOC) to provide communication security, server authentication, and proof-of-possession for a key owned by the client and bound to an OAuth 2.0 access token.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF). Note that other groups may also distribute working documents as Internet-Drafts. The list of current Internet-Drafts is at http://datatracker.ietf.org/drafts/current/.

Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as "work in progress."

This Internet-Draft will expire on January 24, 2018.

Copyright Notice

Table of Contents

1. Introduction

This memo specifies a profile of the ACE framework [I-D.ietf-ace-oauth-authz]. In this profile, a client and a resource server use CoAP [RFC7252] to communicate. The client uses an access token, bound to a key (the proof-of-possession key) to authorize its access to the resource server. In order to provide communication security, proof of possession, and server authentication they use Object Security of CoAP (OSCOAP) [I-D.ietf-core-object-security] and Ephemeral Diffie-Hellman Over COSE (EDHOC) [I-D.selander-ace-cose-ecdhe]. Optionally the client and the resource server may also use CoAP and OSCOAP to communicate with the authorization server. The use of EDHOC in this profile in addition to OSCOAP, provides perfect forward secrecy (PFS) and the initial proof-of-possession, which ties the proof-of-possession key to an OSCOAP security context.

OSCOAP specifies how to use CBOR Object Signing and Encryption (COSE) [RFC8152] to secure CoAP messages. In order to provide replay and reordering protection OSCOAP also introduces sequence numbers that are used together with COSE. EDHOC specifies an authenticated Diffie-Hellman protocol that allows two parties to use CBOR [RFC7049]

and COSE in order to establish a shared secret key with perfect forward secrecy.

1.1. Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119]. These words may also appear in this document in lowercase, absent their normative meanings.

Certain security-related terms such as "authentication", "authorization", "confidentiality", "(data) integrity", "message authentication code", and "verify" are taken from [RFC4949].

Since we describe exchanges as RESTful protocol interactions HTTP [RFC7231] offers useful terminology.

Terminology for entities in the architecture is defined in OAuth 2.0 [RFC6749] and [I-D.ietf-ace-actors], such as client (C), resource server (RS), and authorization server (AS). It is assumed in this document that a given resource on a specific RS is associated to a unique AS.

Note that the term "endpoint" is used here following its OAuth definition, which is to denote resources such as /token and /introspect at the AS and /authz-info at the RS. The CoAP [RFC7252] definition, which is "An entity participating in the CoAP protocol" is not used in this memo.

2. Client to Resource Server

The use of OSCOAP for arbitrary CoAP messages is specified in [I-D.ietf-core-object-security]. This section defines the specific uses and their purpose for securing the communication between a client and a resource server, and the parameters needed to negotiate the use of this profile with the token endpoint at the authorization server as specified in section 5.5 of the ACE framework [I-D.ietf-ace-oauth-authz].

2.1. Signaling the use of OSCOAP

A client requests a token at an AS via the /token endpoint. This follows the message formats specified in section 5.5.1 of the ACE framework [I-D.ietf-ace-oauth-authz].

The AS responding to a successful access token request as defined in section 5.5.2 of the ACE framework can signal that the use of OSCOAP

is REQUIRED for a specific access token by including the "profile" parameter with the value "coap_oscoap" or "coap_oscoap_edhoc" in the access token response. This means that the client MUST use OSCOAP towards all resource servers for which this access token is valid, and follow Section 2.2.1 to derive the security context to run OSCOAP, if the profile has value "coap_oscoap", or follow Section 2.2.2 to derive the security context to run OSCOAP, if the profile has value "coap_oscoap_edhoc".

The error response procedures defined in section 5.5.3 of the ACE framework are unchanged by this profile.

Note the the client and the authorization server MAY OPTIONALLY use OSCOAP to protect the interaction via the /token endpoint. See section 3 for details.

2.2. Key establishment for OSCOAP

Section 3.2 of OSCOAP [I-D.ietf-core-object-security] defines how to derive a security context based on a shared master secret and a few other parameters, established between client and server. The proof-of-possession key (pop-key) provisioned from the AS MAY, in case of pre-shared keys, be used directly as master secret in OSCOAP. Alternatively the pop-key (symmetric or asymmetric) MAY be used to authenticate the messages in the key exchange protocol EDHOC [I-D.selander-ace-cose-ecdhe], from which a master secret is derived.

2.2.1. Using the pop-key with OSCOAP directly (OSCOAP)

If OSCOAP is used directly with the symmetric pop-key as master secret, then the AS MUST provision the following data, in response to the access token request:

- a master secret
- the sender identifier
- the recipient identifier

Additionally, the AS MAY provision the following data, in the same response. In case these parameters are omitted, the default values are used as described in section 3.2. of [I-D.ietf-core-object-security].

- an AEAD algorithm
- a KDF algorithm

o a salt

The master secret MUST be communicated as COSE_Key in the 'cnf' parameter of the access token response as defined in section 5.5.4.5 of [I-D.ietf-ace-oauth-authz]. The AEAD algorithm MAY be included as the 'alg' parameter in the COSE_Key; the KDF algorithm MAY be included as the 'kdf' parameter of the COSE_Key and the salt MAY be included as the 'slt' parameter of the COSE_Key as defined in table 1. The same parameters MUST be included as metadata of the access token, if the token is a CWT [I-D.ietf-ace-cbor-web-token], the same COSE_Key structure MUST be placed in the 'cnf' claim of this token. The AS MUST also assign identifiers to both client and RS, which are then used as Sender ID and Recipient ID in the OSCOAP context as described in section 3.1. of [I-D.ietf-core-object-security]. These identifiers MUST be unique in the set of all clients and RS identifiers for a certain AS. Moreover, these MUST be included in the COSE_Key as header parameters, as defined in table 1.

We assume in this document that a resource is associated to one single AS, which makes it possible to assume unique identifiers for each client requesting a particular resource to a RS. If this is not the case, collisions of identifiers may appear in the RS, in which case the RS needs to have a mechanism in place to disambiguate identifiers or mitigate their effect.

Note that C should set the Sender ID of its security context to the clientId value received and the Recipient ID to the serverId value, and RS should do the opposite.

name	label	CBOR type	registry	description
clientId	TBD	bstr		Identifies the client in an OSCOAP context using this key
serverId	TBD	bstr		Identifies the server in an OSCOAP context using this key
kdf	TBD	bstr		Identifies the KDF algorithm in an OSCOAP context using this key
slt	TBD	bstr		Identifies the master salt in an OSCOAP context using this key

Table 1: Additional common header parameters for COSE_Key

Figure 1 shows an example of such an AS response, in CBOR diagnostic notation without the tag and value abbreviations.

```
Header: Created (Code=2.01)
Content-Type: "application/cose+cbor"
Payload:
{
  "access_token" : b64'SlAV32hkKG ...
   (remainder of access token omitted for brevity)',
  "profile" : "coap_oscoap",
  "expires_in" : "3600",
  "cnf" : {
    "COSE_Key" : {
      "kty" : "Symmetric",
      "alg" : "AES-CCM-16-64-128",
      "clientId" : b64'qA',
      "serverId" : b64'Qg',
      "k" : b64'+a+Dg2jjU+eIiOFCa9lObw'
    }
  }
}
```

Figure 1: Example AS response with OSCOAP parameters.

Figure 2 shows an example CWT, containing the necessary OSCOAP parameters in the 'cnf' claim, in CBOR diagnostic notation without tag and value abbreviations.

```
{
  "aud" : "tempSensorInLivingRoom",
  "iat" : "1360189224",
  "exp" : "1360289224",
  "scope" :  "temperature_g firmware_p",
  "cnf" : {
    "COSE_Key" : {
      "kty" : "Symmetric",
      "alg" : "AES-CCM-16-64-128",
      "clientId" : b64'Qg',
      "serverId" : b64'qA',
      "k" : b64'+a+Dg2jjU+eIiOFCa9lObw'
  }
}
```

Figure 2: Example CWT with OSCOAP parameters.

2.2.2. Using the pop-key with EDHOC (EDHOC+OSCOAP)

If EDHOC is used together with OSCOAP, and the pop-key (symmetric or asymmetric) is used to authenticate the messages in EDHOC, then the AS MUST provision the following data, in response to the access token request:

- a symmetric or asymmetric key (associated to the RS)

- a key identifier;

How these parameters are communicated depends on the type of key (asymmetric or symmetric).

Note that in the case described in this section, the 'expires_in' parameter, defined in section 4.2.2. of [RFC6749] defines the lifetime in seconds of both the access token and the shared secret. After expiration, C MUST acquire a new access token from the AS, and run EDHOC again, as specified in this section

2.2.2.1. Using Asymmetric Keys

In case of an asymmetric key, C MUST communicate its own asymmetric key to the AS in the 'cnf' parameter of the access token request, as specified in section 5.5.1 of [I-D.ietf-ace-oauth-authz]; the AS MUST communicate the RS's public key to C in the response, in the 'rs_cnf' parameter, as specified in section 5.5.1 of [I-D.ietf-ace-oauth-authz]. Note that the RS's public key MUST include a 'kid' parameter, and that the value of the 'kid' MUST be included in the access token, to let the RS know which of its public keys C used. If the access token is a CWT [I-D.ietf-ace-cbor-web-token], the key identifier MUST be placed directly in the 'cnf' structure (if the key is only referenced).

Figure 3 shows an example of such a request in CBOR diagnostic notation without tag and value abbreviations.

```
Header: POST (Code=0.02)
Uri-Host: "server.example.com"
Uri-Path: "token"
Content-Type: "application/cose+cbor"
Payload:
{
  "grant_type" : "client_credentials",
  "cnf" : {
    "COSE_Key" : {
      "kid" : "client_key"
      "kty" : "EC",
      "crv" : "P-256",
      "x" : b64'usWxHK2PmfnHKwXPS54m0kTcGJ90UiglWiGahtagnv8',
      "y" : b64'IBOL+C3BttVivg+lSreASjpkttcsz+1rb7btKLv8EX4'
    }
  }
}
```

Figure 3: Example access token request (OSCOAP+EDHOC, asymmetric).

Figure 4 shows an example of a corresponding response in CBOR diagnostic notation without tag and value abbreviations.

```
Header: Created (Code=2.01)
Content-Type: "application/cose+cbor"
Payload:
{
  "access_token" : b64'SlAV32hkKG ...
   (contains "kid" : "server_key")',
  "profile" : "coap_oscoap_edhoc",
  "expires_in" : "3600",
  "rs_cnf" : {
    "COSE_Key" : {
      "kid" : "server_key"
      "kty" : "EC",
      "crv" : "P-256",
      "x" : b64'cGJ90UiglWiGahtagnv8usWxHK2PmfnHKwXPS54m0kT',
      "y" : b64'reASjpkttcsz+1rb7btKLv8EX4IBOL+C3BttVivg+lS'
    }
  }
}
```

Figure 4: Example AS response (EDHOC+OSCOAP, asymmetric).

2.2.2.2. Using Symmetric Keys

In the case of a symmetric key, the AS MUST communicate the key to the client in the 'cnf' parameter of the access token response, as specified in section 5.5.2. of [I-D.ietf-ace-oauth-authz]. AS MUST also select a key identifier, that MUST be included as the 'kid' parameter either directly in the 'cnf' structure, as in figure 4 of [I-D.ietf-ace-oauth-authz], or as the 'kid' parameter of the COSE_key, as in figure 6 of [I-D.ietf-ace-oauth-authz].

Figure 5 shows an example of the necessary parameters in the AS response to the access token request when EDHOC is used. The example uses CBOR diagnostic notation without tag and value abbreviations.

```
Header: Created (Code=2.01)
Content-Type: "application/cose+cbor"
Payload:
{
  "access_token" : b64'SlAV32hkKG ...
   (remainder of access token omitted for brevity)',
  "profile" : "coap_oscoap_edhoc",
  "expires_in" : "3600",
  "cnf" : {
    "COSE_Key" : {
      "kty" : "Symmetric",
      "kid" : b64'5tOS+h42dkw',
      "k" : b64'+a+Dg2jjU+eIiOFCa9lObw'
    }
  }
}
```

Figure 5: Example AS response (EDHOC+OSCOAP, symmetric).

In both cases, the AS MUST also include the same key identifier as 'kid' parameter in the access token metadata. If the access token is a CWT [I-D.ietf-ace-cbor-web-token], the key identifier MUST be placed inside the 'cnf' claim as 'kid' parameter of the COSE_Key or directly in the 'cnf' structure (if the key is only referenced).

Figure 6 shows an example CWT containing the necessary EDHOC+OSCOAP parameters in the 'cnf' claim, in CBOR diagnostic notation without tag and value abbreviations.

```
{
  "aud" : "tempSensorInLivingRoom",
  "iat" : "1360189224",
  "exp" : "1360289224",
  "scope" :  "temperature_g firmware_p",
  "cnf" : {
    "COSE_Key" : {
      "kty" : "Symmetric",
      "kid" : b64'5tOS+h42dkw',
      "k" : b64'+a+Dg2jjU+eIiOFCa9lObw'
  }
}
```

Figure 6: Example CWT with EDHOC+OSCOAP, symmetric case.

All other parameters defining OSCOAP security context are derived from EDHOC message exchange, including the master secret (see Appendix C.2 of [I-D.selander-ace-cose-ecdhe]).

2.2.2.3. Processing

To provide forward secrecy and mutual authentication in the case of pre-shared keys, pre-established raw public keys or with X.509 certificates it is RECOMMENDED to use EDHOC [I-D.selander-ace-cose-ecdhe] to generate the keying material. EDHOC MUST be used as defined in Appendix C, with the following additions and modifications.

The first EDHOC message is sent after the access token is posted to the /authz-info endpoint of the RS as specified in section 5.7.1 of, as defined in [I-D.ietf-ace-oauth-authz]. Then the EDHOC message_1 is sent and the EDHOC protocol is initiated [I-D.selander-ace-cose-ecdhe]).

Before the RS continues with the EDHOC protocol and responds to this token submission request, additional verifications on the access token are done: the RS SHALL process the access token according to [I-D.ietf-ace-oauth-authz]. If the token is valid then the RS continues processing EDHOC following Appendix C of [I-D.selander-ace-cose-ecdhe], otherwise it discontinues EDHOC and responds with the error code as specified in [I-D.ietf-ace-oauth-authz].

- In case the EDHOC verification fails, the RS MUST return an error response to the client with code 4.01 (Unauthorized).

- If RS has an access token for C but not for the resource that C has requested, RS MUST reject the request with a 4.03 (Forbidden).

- If RS has an access token for C but it does not cover the action C requested on the resource, RS MUST reject the request with a 4.05 (Method Not Allowed).

- If all verifications above succeeds, further communication between client and RS is protected with OSCOAP, including the RS response to the OSCOAP request.

In the case of EDHOC being used with symmetric keys, the protocol in section 5 of [I-D.selander-ace-cose-ecdhe] MUST be used. If the key is asymmetric, the RS MUST also use an asymmetric key for authentication. This key is known to the client through the access token response (see section 5.5.2 of the ACE framework). In this case the protocol in section 4 of [I-D.selander-ace-cose-ecdhe] MUST be used.

Figure 7 illustrates the message exchanges for using OSCOAP+EDHOC (step C in figure 1 of [I-D.ietf-ace-oauth-authz]).

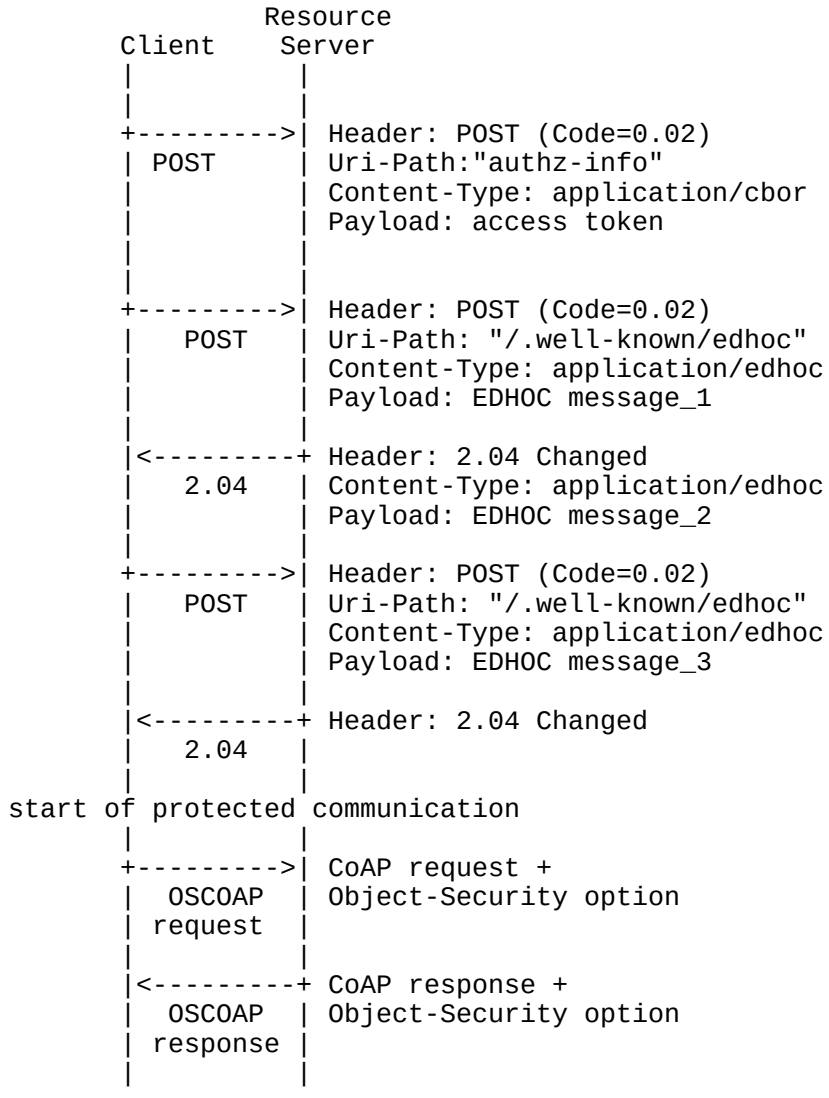

Figure 7: Access token and key establishment with EDHOC

2.3. Client to Authorization Server

As specified in the ACE framework section 5.5 [I-D.ietf-ace-oauth-authz], the Client and AS can also use CoAP instead of HTTP to communicate via the token endpoint. This section specifies how to use OSCOAP between Client and AS together with CoAP. The use of OSCOAP for this communication is OPTIONAL in this profile, other security protocols (such as DTLS) MAY be used instead.

The client and the AS are expected to have pre-established credentials (e.g. raw public keys). How these credentials are established is out of scope for this profile. Furthermore the client and the AS communicate using CoAP through the token endpoint as specified in section 5.5 of [I-D.ietf-ace-oauth-authz]. At first point of contact, prior to making the token request and response, the client and the AS MAY perform an EDHOC exchange with the pre-established credentials to create forward secret keying material for use with OSCOAP. Subsequent requests and the responses MUST be protected with OSCOAP.

3. Resource Server to Authorization Server

As specified in the ACE framework section 5.6 [I-D.ietf-ace-oauth-authz], the RS and AS can also use CoAP instead of HTTP to communicate via the introspection endpoint. This section specifies how to use OSCOAP between RS and AS together with CoAP. The use of OSCOAP for this communication is OPTIONAL in this profile, other security protocols (such as DTLS) MAY be used instead.

The RS and the AS are expected to have pre-established credentials (e.g. symmetric keys). How these credentials are established is out of scope for this profile. Furthermore the RS and the AS communicate using CoAP through the introspection endpoint as specified in section 5.6 of [I-D.ietf-ace-oauth-authz]. At first point of contact, prior to making the introspection request and response, the RS and the AS MAY perform an EDHOC exchange with the pre-established credentials to create forward secret keying material for use with OSCOAP. Subsequent requests and the responses MUST be protected with OSCOAP

4. Security Considerations

TBD.

5. Privacy Considerations

TBD.

6. IANA Considerations

TBD. ’coap_oscoap’ as profile id. Header parameters ’sid’, ’rid’, ’kdf’ and ’slt’ for COSE_Key.

7. Acknowledgments

The author wishes to thank Jim Schaad, Goeran Selander and Marco Tiloca for the input on this memo. The error responses specified in

section 2.2. were originally specified by Gerdes et al. in [I-D.gerdes-ace-dcaf-authorize].

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July 03, 2017

Ephemeral Diffie-Hellman Over COSE (EDHOC)
draft-selander-ace-cose-ecdhe-07

Abstract

This document specifies Ephemeral Diffie-Hellman Over COSE (EDHOC), a compact, and lightweight authenticated Diffie-Hellman key exchange with ephemeral keys that can be used over any layer. EDHOC messages are encoded with CBOR and COSE, allowing reuse of existing libraries.

Table of Contents

1. Introduction

Security at the application layer provides an attractive option for protecting Internet of Things (IoT) deployments, for example where transport layer security is not sufficient [I-D.hartke-core-e2e-security-reqs] or where the protocol needs to work on a variety of underlying protocols. IoT devices may be constrained in various ways, including memory, storage, processing capacity, and energy [RFC7228]. A method for protecting individual messages at the application layer suitable for constrained devices, is provided by CBOR Object Signing and Encryption (COSE) [I-D.ietf-cose-msg]), which builds on the Concise Binary Object Representation (CBOR) [RFC7049].

In order for a communication session to provide forward secrecy, the communicating parties can run an Elliptic Curve Diffie-Hellman (ECDH) key exchange protocol with ephemeral keys, from which shared key material can be derived. This document specifies Ephemeral Diffie-Hellman Over COSE (EDHOC), an authenticated ECDH protocol using CBOR and COSE objects. Authentication is based on credentials established out of band, e.g. from a trusted third party, such as an Authorization Server as specified by [I-D.ietf-ace-oauth-authz]. EDHOC supports authentication using pre-shared keys (PSK), raw public keys (RPK), and certificates (Cert). Note that this document focuses on authentication and key establishment: for integration with authorization of resource access, refer to [I-D.seitz-ace-oscoap-profile]. This document also specifies the derivation of shared key material.

The ECDH exchange and the key derivation follow [SIGMA], NIST SP-800-56a [SP-800-56a], and HKDF [RFC5869]. CBOR [RFC7049] and COSE [I-D.ietf-cose-msg] are used to implement these standards.

1.1. Terminology

This document use the same informational CBOR Data Definition Language (CDDL) [I-D.greevenbosch-appsawg-cbor-cddl] grammar as COSE (see Section 1.3 of [I-D.ietf-cose-msg]). A vertical bar | denotes byte string concatenation.

1.2. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119]. These words may also appear in this document in lowercase, absent their normative meanings.

2. Protocol Overview

SIGMA (SIGn-and-MAc) is a family of theoretical protocols with a large number of variants [SIGMA]. Like IKEv2 and TLS 1.3, EDHOC is built on a variant of the SIGMA protocol which provide identity protection, and like TLS 1.3, EDHOC implements the SIGMA-I variant as Sign-then-MAC. The SIGMA-I protocol using an AEAD algorithm is shown in Figure 1.

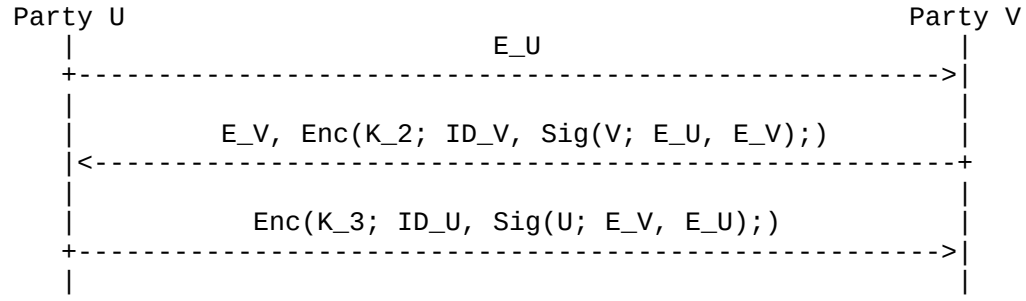

Figure 1: AEAD variant of the SIGMA-I protocol

The parties exchanging messages are called "U" and "V". They exchange identities and ephemeral public keys, compute the shared secret, and derive the keying material. The messages are signed, MACed, and encrypted.

- E_U and E_V are the ECDH ephemeral public keys of U and V, respectively.
- ID_U and ID_V are identifiers for the public keys of U and V, respectively.
- Sig(U; .) and S(V; .) denote signatures made with the private key of U and V, respectively.
- Enc(K; P; A) denotes AEAD encryption of plaintext P and additional authenticated data A using the key K derived from the shared secret. The AEAD MUST NOT be replaced by plain encryption, see Section 8.

As described in Appendix B of [SIGMA], in order to create a "full-fledged" protocol some additional protocol elements are needed. EDHOC adds:

- Explicit session identifiers S_U, S_V different from other concurrent session identifiers (EDHOC or other used protocol identifier) chosen by U and V, respectively.
- Explicit nonces N_U, N_V chosen freshly and anew with each session by U and V, respectively.
- Computationally independent keys derived from the ECDH shared secret and used for encryption of different messages.

EDHOC also makes the following additions:

- Negotiation of key derivation, encryption, and signature algorithms:
  - U proposes one or more algorithms of the following kinds:
    - HKDF
    - AEAD
    - Signature verification
    - Signature generation
  - V selects one algorithm of each kind
- Verification of common preferred ECDH curve:
  - U lists supported ECDH curves in order of preference
  - V verifies that the ECDH curve of the ephemeral key is the most preferred common curve
- Transport of opaque application defined data.

EDHOC is designed to encrypt and integrity protect as much information as possible, and all symmetric keys are derived using as much previous information as possible. EDHOC is furthermore designed to be as compact and lightweight as possible, in terms of message sizes, processing, and the ability to reuse already existing CBOR and COSE libraries. EDHOC does not put any requirement on the lower layers and can therefore be also be used e.g. in environments without IP.

This paper is organized as follows: Section 3 specifies general properties of EDHOC, including formatting of the ephemeral public keys and key derivation, Section 4 specifies EDHOC with asymmetric key authentication, Section 5 specifies EDHOC with symmetric key authentication, and Appendix A provides a wealth of test vectors to ease implementation and ensure interoperability.

3. EDHOC Overview

EDHOC consists of three messages (message_1, message_2, message_3) that maps directly to the three messages in SIGMA-I, plus an EDHOC error message. All EDHOC messages consists of a CBOR array where the first element is an int specifying the message type (MSG_TYPE). After creating EDHOC message_3, Party U can derive the traffic key (master secret) and protected application data can therefore be sent

in parallel with EDHOC message_3. The application data may be protected using the negotiated AEAD algorithm and the explicit session identifiers S_U and S_V. EDHOC may be used with the media type application/edhoc defined in Section 7.

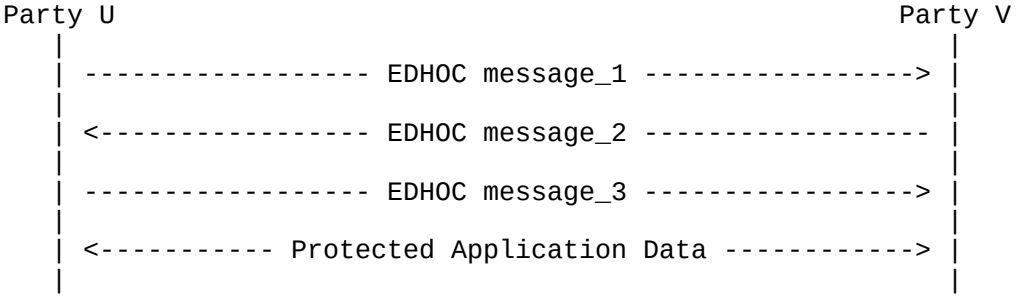

Figure 2: EDHOC message flow

The EDHOC message exchange may be authenticated using pre-shared keys (PSK), raw public keys (RPK), or certificates (Cert). EDHOC assumes the existence of mechanisms (certification authority, manual distribution, etc.) for binding identities with authentication keys (public or pre-shared). EDHOC with symmetric key authentication is very similar to EDHOC with asymmetric key authentication, the difference being that information is only MACed, not signed.

EDHOC also allows opaque application data (APP_1, APP_2, APP_3) to be sent in the respective messages. APP_1 is unprotected, APP_2 is protected (encrypted and integrity protected), and APP_3 is protected and mutually authenticated. When EDHOC is used with asymmetric key authentication APP_2 is sent to an unauthenticated party, but with symmetric key authentication APP_2 is mutually authenticated.

3.1. Formatting of the Ephemeral Public Keys

The ECDH ephemeral public key SHALL be formatted as a COSE_Key of type EC2 or OKP according to section 13.1 and 13.2 of [I-D.ietf-cose-msg]. The curve X25519 is mandatory to implement. For Elliptic Curve Keys of type EC2, compact representation and compact output as per [RFC6090] SHALL be used, i.e. the ’y’ parameter SHALL NOT be present in the The COSE_Key object. COSE [I-D.ietf-cose-msg] always use compact output for Elliptic Curve Keys of type EC2.

3.2. Key Derivation

Key and IV derivation SHALL be done as specified in Section 11.1 of [I-D.ietf-cose-msg] with the following input:

- The PRF SHALL be the HKDF [RFC5869] in the ECDH-SS w/ HKDF negotiated during the message exchange (HKDF_V).
- The secret SHALL be the ECDH shared secret as defined in Section 12.4.1 of [I-D.ietf-cose-msg].
- The salt SHALL be the PSK when EDHOC is authenticated with symmetric keys and nil when EDHOC is authenticated with asymmetric keys.
- The fields in the context information COSE_KDF_Context SHALL have the following values:
  - AlgorithmID is a tstr as defined below
  - PartyUInfo = PartyVInfo = (nil, nil, nil)
  - keyDataLength is a uint as defined below
  - protected SHALL be a zero length bstr
  - other is a bstr SHALL be aad_2, aad_3, or exchange_hash

where exchange_hash, in non-CDDL notation, is:

exchange_hash = H(H(message_1 | message_2) | message_3)

where H() is the hash function in HKDF_V.

For message_i the key, called K_i, SHALL be derived using other = aad_i, where i = 2 or 3. The key SHALL be derived using AlgorithmID set to the name of the negotiated AEAD (AEAD_V), and keyDataLength equal to the key length of AEAD_V.

If the AEAD algorithm requires an IV, then IV_i for message_i SHALL be derived using other = aad_i, where i = 2 or 3. The IV SHALL be derived using AlgorithmID = "IV-GENERATION" as specified in section 12.1.2. of [I-D.ietf-cose-msg], and keyDataLength equal to the IV length of AEAD_V.

Application specific traffic keys and other data SHALL be derived using other = exchange_hash. AlgorithmID SHALL be a tstr defined by the application and SHALL be different for different data being derived (an example is given in Appendix C.2). keyDataLength is set to the length of the data being derived.

4. EDHOC Authenticated with Asymmetric Keys

4.1. Overview

EDHOC supports authentication with raw public keys (RPK) and certificates (Cert) with the requirements that:

- Party U SHALL be able to identify Party V's public key using ID_V.
- Party V SHALL be able to identify Party U's public key using ID_U.

ID_U and ID_V SHALL either contain the credential used for authentication (e.g. x5bag or x5chain) or uniquely identify the credential used for authentication (e.g. x5t), see [I-D.schaad-cose-x509]. Party U and V MAY retrieve the other party's credential out of band. Optionally, ID_U and ID_V are complemented with the additional parameters HINT_ID_U and HINT_ID_V containing information about how to retrieve the credential of Party U and Party V, respectively (e.g. x5u), see [I-D.schaad-cose-x509].

Party U and Party V MAY use different type of credentials, e.g. one uses RPK and the other uses Cert. Party U and Party V MAY use different signature algorithms.

EDHOC with asymmetric key authentication is illustrated in Figure 3.

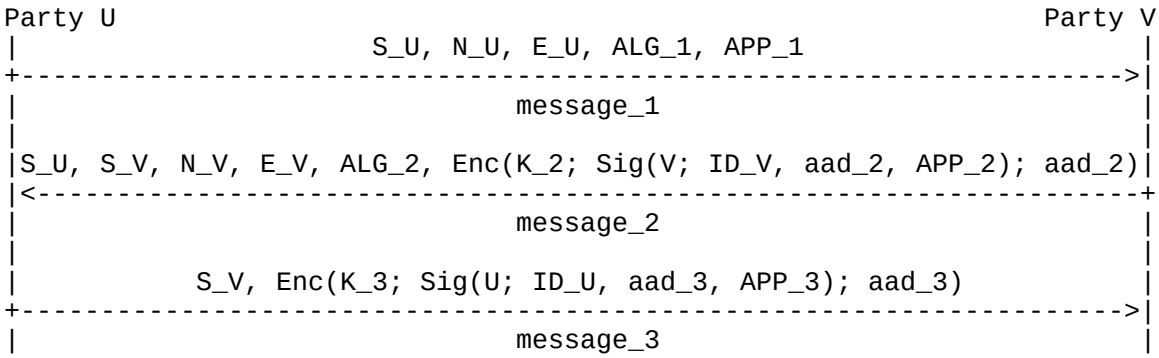

Figure 3: EDHOC with asymmetric key authentication.

4.1.1. Mandatory to Implement Algorithms

For EDHOC authenticated with asymmetric keys, the COSE algorithms ECDH-SS + HKDF-256, AES-CCM-64-64-128, and EdDSA are mandatory to implement.

4.2. EDHOC Message 1

4.2.1. Formatting of Message 1

message_1 SHALL be a CBOR array as defined below

```
message_1 = [
  MSG_TYPE : int,
  S_U : bstr,
  N_U : bstr,
  E_U : serialized_COSE_Key,
  ECDH-Curves_U : alg_array,
  HKDFs_U : alg_array,
  AEADs_U : alg_array,
  SIGs_V : alg_array,
  SIGs_U : alg_array,
  ? APP_1 : bstr
]

serialized_COSE_Key = bstr .cbor COSE_Key

alg_array = [ + alg : int / tstr ]
```

where:

- MSG_TYPE = 1
- S_U - variable length session identifier
- N_U - 64-bit random nonce
- E_U - the ephemeral public key of Party U
- ECDH-Curves_U - EC curves for ECDH which Party U supports, in the order of decreasing preference
- HKDFs_U - supported ECDH-SS w/ HKDF algorithms
- AEADs_U - supported AEAD algorithms
- SIGs_V - signature algorithms, with which Party U supports verification

- SIGs_U - signature algorithms, with which Party U supports signing

- APP_1 - bstr containing opaque application data

4.2.2. Party U Processing of Message 1

Party U SHALL compose message_1 as follows:

- Determine which ECDH curve to use with Party V. If U previously received from Party V an error message to message_1 with diagnostic payload identifying an ECDH curve in ECDH-Curves_U, then U SHALL retrieve an ephemeral from that curve. Otherwise the first curve in ECDH-Curves_U MUST be used.

- Retrieve an ephemeral ECDH key pair generated as specified in Section 5 of [SP-800-56a] and format the ephemeral public key E_U as a COSE_key as specified in Section 3.1.

- Generate the pseudo-random nonce N_U.

- Choose a session identifier S_U which is not in use and store it for the length of the protocol. The session identifier SHOULD be different from other concurrent session identifiers used by Party U. The session identifier MAY be used with the protocol for which EDHOC establishes traffic keys/master secret, in which case S_U SHALL be different from the concurrently used session identifiers of that protocol.

- Format message_1 as specified in Section 4.2.1.

4.2.3. Party V Processing of Message 1

Party V SHALL process message_1 as follows:

- Verify (OPTIONAL) that N_U has not been received before.

- Verify that at least one of each kind of the proposed algorithms are supported.

- Verify that the ECDH curve used in E_U is supported, and that no prior curve in ECDH-Curves_U is supported.

- For EC2 curves, validate that E_U is a valid point by verifying that there is a solution to the curve definition for the given parameter 'x'.

If any verification step fails, Party V MUST send an EDHOC error message back, formatted as defined in Section 6.1, and the protocol

MUST be discontinued. If V does not support the ECDH curve used in E_U, but supports another ECDH curves in ECDH-Curves_U, then the error message MUST include the following diagnostic payload describing the first supported ECDH curve in ECDH-Curves_U:

ERR_MSG = "Curve not supported; X"

where X is the first curve in ECDH-Curves_U that V supports, encoded as in Table 22 of {{I-D.ietf-cose-msg}}.

o Pass APP_1 to the application.

4.3. EDHOC Message 2

4.3.1. Formatting of Message 2

message_2 SHALL be a CBOR array as defined below

```
message_2 = [
  data_2,
  COSE_ENC_2 : COSE_Encrypt0
]

data_2 = (
  MSG_TYPE : int,
  S_U : bstr,
  S_V : bstr,
  N_V : bstr,
  E_V : serialized_COSE_Key,
  HKDF_V : int / tstr,
  AEAD_V : int / tstr,
  SIG_V : int / tstr,
  SIG_U : int / tstr
)

aad_2 : bstr
```

where aad_2, in non-CDDL notation, is:

aad_2 = H(message_1 | [data_2])

where:

o MSG_TYPE = 2

o S_V - variable length session identifier

o N_V - 64-bit random nonce

- E_V - the ephemeral public key of Party V
- HKDF_V - a single chosen algorithm from HKDFs_U
- AEAD_V - a single chosen algorithm from AEADs_U
- SIG_V - a single chosen algorithm from SIGs_V with which Party V signs
- SIG_U - a single chosen algorithm from SIGs_U with which Party U signs
- COSE_ENC_2 has the following fields and values:
  - external_aad = aad_2
  - plaintext = [COSE_SIG_V, ? APP_2]
- COSE_SIG_V is a COSE_Sign1 object with the following fields and values:
  - protected = { abc : ID_V, ? xyz : HINT_ID_V }
  - detached payload = aad_2, ? APP_2
- abc - any COSE map label that can identify a public key, see Section 4.1
- ID_V - identifier for the public key of Party V
- xyz - any COSE map label for information about how to retrieve the credential of Party V, see Section 4.1
- HINT_ID_V - information about how to retrieve the credential of Party V
- APP_2 - bstr containing opaque application data
- H() - the hash function in HKDF_V

4.3.2. Party V Processing of Message 2

Party V SHALL compose message_2 as follows:

- Retrieve an ephemeral ECDH key pair generated as specified in Section 5 of [SP-800-56a] using same curve as used in E_U. Format the ephemeral public key E_V as a COSE_key as specified in Section 3.1.

- Generate the pseudo-random nonce N_V.

- Choose a session identifier S_V which is not in use and store it for the length of the protocol. The session identifier SHOULD be different from other relevant concurrent session identifiers used by Party V. The session identifier MAY be used with the protocol for which EDHOC establishes traffic keys/master secret, in which case S_V SHALL be different from the concurrently used session identifiers of that protocol.

- Select HKDF_V, AEAD_V, SIG_V, and SIG_U from the algorithms proposed in HKDFs_U, AEADs_U, SIGs_V, and SIGs_U.

- Format message_2 as specified in Section 4.3.1:

  * COSE_Sign1 is computed as defined in section 4.4 of [I-D.ietf-cose-msg], using algorithm SIG_V and the private key of Party V.

  * COSE_Encrypt0 is computed as defined in section 5.3 of [I-D.ietf-cose-msg], with AEAD_V, K_2, and IV_2. The AEAD algorithm MUST NOT be replaced by plain encryption, see Section 8.

4.3.3. Party U Processing of Message 2

Party U SHALL process message_2 as follows:

- Use the session identifier S_U to retrieve the protocol state.

- Verify that HKDF_V, AEAD_V, SIG_V, and SIG_U were proposed in HKDFs_U, AEADs_U, SIGs_V, and SIGs_U.

- Verify (OPTIONAL) that N_V has not been received before.

- For EC2 curves, validate that E_V is a valid point by verifying that there is a solution to the curve definition for the given parameter 'x'.

- Verify message_2 as specified in Section 4.3.1:

  * COSE_Encrypt0 is decrypted defined in section 5.3 of [I-D.ietf-cose-msg], with AEAD_V, K_2, and IV_2.

  * COSE_Sign1 is verified as defined in section 4.4 of [I-D.ietf-cose-msg], using algorithm SIG_V and the public key of Party V.

If any verification step fails, Party U MUST send an EDHOC error message back, formatted as defined in Section 6.1, and the protocol MUST be discontinued.

- Pass APP_2 to the application.

4.4. EDHOC Message 3

4.4.1. Formatting of Message 3

message_3 SHALL be a CBOR array as defined below

```
message_3 = [
  data_3,
  COSE_ENC_3 : COSE_Encrypt0
]

data_3 = (
  MSG_TYPE : int,
  S_V : bstr
)

aad_3 : bstr
```

where aad_3, in non-CDDL notation, is:

```
aad_3 = H( H( message_1 | message_2 ) | [ data_3 ] )
```

where:

- MSG_TYPE = 3
- COSE_ENC_3 has the following fields and values:
  - external_aad = aad_3
  - plaintext = [COSE_SIG_U, ? APP_3]
- COSE_SIG_U is a COSE_Sign1 object with the following fields and values:
  - protected = { abc : ID_U, ? xyz : HINT_ID_U }
  - detached payload = aad_3, ? APP_3
- abc - any COSE map label that can identify a public key, see Section 4.1

- ID_U - identifier for the public key of Party U
- xyz - any COSE map label for information about how to retrieve the credential of Party U, see Section 4.1
- HINT_ID_V - information about how to retrieve the credential of Party U
- APP_3 - bstr containing opaque application data

4.4.2. Party U Processing of Message 3

Party U SHALL compose message_3 as follows:

- Format message_3 as specified in Section 4.4.1:
  - COSE_Sign1 is computed as defined in section 4.4 of [I-D.ietf-cose-msg], using algorithm SIG_U and the private key of Party U.
  - COSE_Encrypt0 is computed as defined in section 5.3 of [I-D.ietf-cose-msg], with AEAD_V, K_3, and IV_3. The AEAD algorithm MUST NOT be replaced by plain encryption, see Section 8.

4.4.3. Party V Processing of Message 3

Party V SHALL process message_3 as follows:

- Use the session identifier S_V to retrieve the protocol state.
- Verify message_3 as specified in Section 4.4.1:
  - COSE_Encrypt0 is decrypted as defined in section 5.3 of [I-D.ietf-cose-msg], with AEAD_V, K_3, and IV_3.
  - COSE_Sign1 is verified as defined in section 4.4 of [I-D.ietf-cose-msg], using algorithm SIG_U and the public key of Party U.

If any verification step fails, Party V MUST send an EDHOC error message back, formatted as defined in Section 6.1, and the protocol MUST be discontinued.

- Pass APP_3 to the application.

5. EDHOC Authenticated with Symmetric Keys

5.1. Overview

EDHOC supports authentication with pre-shared keys. Party U and V are assumed to have a pre-shared uniformly random key (PSK) with the requirement that:

- Party V SHALL be able to identify the PSK using KID.

KID may optionally contain information about how to retrieve the PSK.

EDHOC with symmetric key authentication is illustrated in Figure 4.

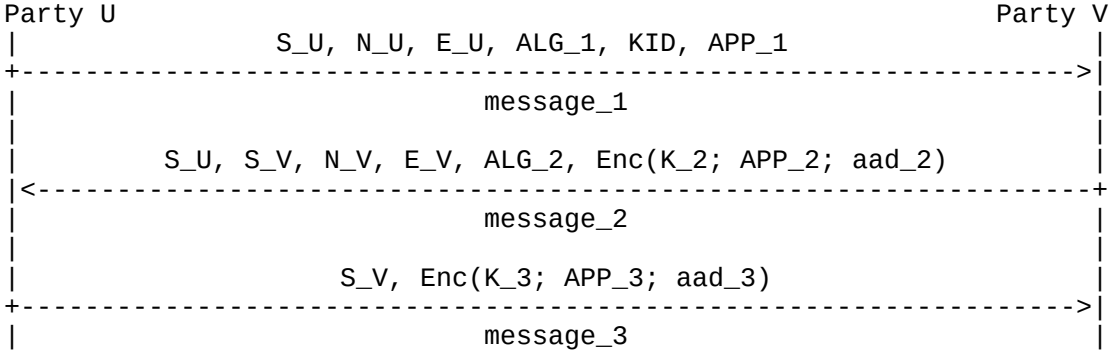

Figure 4: EDHOC with symmetric key authentication.

5.1.1. Mandatory to Implement Algorithms

For EDHOC authenticated with symmetric keys, the COSE algorithms ECDH-SS + HKDF-256 and AES-CCM-64-64-128 are mandatory to implement.

5.2. EDHOC Message 1

5.2.1. Formatting of Message 1

message_1 SHALL be a CBOR array as defined below

```
message_1 = [
  data_1
]

data_1 = (
  MSG_TYPE : int,
  S_U : bstr,
  N_U : bstr,
  E_U : serialized_COSE_Key,
  ECDH-Curves_U : alg_array,
  HKDFs_U : alg_array,
  AEADs_U : alg_array,
  KID : bstr,
  ? APP_1 : bstr
)

serialized_COSE_Key = bstr .cbor COSE_Key

alg_array = [ + alg : int / tstr ]
```

where:

- MSG_TYPE = 4
- S_U - variable length session identifier
- N_U - 64-bit random nonce
- E_U - the ephemeral public key of Party U
- ECDH-Curves_U - EC curves for ECDH which Party U supports, in the order of decreasing preference
- HKDFs_U - supported ECDH-SS w/ HKDF algorithms
- AEADs_U - supported AEAD algorithms
- KID - identifier of the pre-shared key
- APP_1 - bstr containing opaque application data

5.2.2. Party U Processing of Message 1

Party U SHALL compose message_1 as follows:

- Determine which ECDH curve to use with Party V. If U previously received from Party V an error message to message_1 with diagnostic payload identifying an ECDH curve in ECDH-Curves_U,

then U SHALL retrieve an ephemeral from that curve. Otherwise the first curve in ECDH-Curves_U MUST be used.

- Retrieve an ephemeral ECDH key pair generated as specified in Section 5 of [SP-800-56a] and format the ephemeral public key E_U as a COSE_key as specified in Section 3.1.

- Generate the pseudo-random nonce N_U.

- Choose a session identifier S_U which is not in use and store it for the length of the protocol. The session identifier SHOULD be different from other relevant concurrent session identifiers used by Party U. The session identifier MAY be used with the protocol for which EDHOC establishes traffic keys/master secret, in which case S_U SHALL be different from the concurrently used session identifiers of that protocol.

- Format message_1 as specified in Section 5.2.1.

5.2.3. Party V Processing of Message 1

Party V SHALL process message_1 as follows:

- Verify (OPTIONAL) that N_U has not been received before.

- Verify that at least one of each kind of the proposed algorithms are supported.

- Verify that the ECDH curve used in E_U is supported, and that no prior curve in ECDH-Curves_U is supported.

- For EC2 curves, validate that E_U is a valid point by verifying that there is a solution to the curve definition for the given parameter 'x'.

If any verification step fails, Party V MUST send an EDHOC error message back, formatted as defined in Section 6.1, and the protocol MUST be discontinued. If V does not support the ECDH curve used in E_U, but supports another ECDH curves in ECDH-Curves_U, then the error message SHOULD include a diagnostic payload describing the first supported ECDH curve in ECDH-Curves_U.

- Pass APP_1 to the application.

5.3. EDHOC Message 2

5.3.1. Formatting of Message 2

message_2 SHALL be a CBOR array as defined below

```
message_2 = [
  data_2,
  COSE_ENC_2 : COSE_Encrypt0
]

data_2 = (
  MSG_TYPE : int,
  S_U : bstr,
  S_V : bstr,
  N_V : bstr,
  E_V : serialized_COSE_Key,
  HKDF_V : int / tstr,
  AEAD_V : int / tstr
)

aad_2 : bstr
```

where aad_2, in non-CDDL notation, is:

aad_2 = H(message_1 | [data_2])

where:

- MSG_TYPE = 5
- S_V - variable length session identifier
- N_V - 64-bit random nonce
- E_V - the ephemeral public key of Party V
- HKDF_V - an single chosen algorithm from HKDFs_U
- AEAD_V - an single chosen algorithm from AEADs_U
- COSE_ENC_2 has the following fields and values:
  - external_aad = aad_2
  - plaintext = ? APP_2
- APP_2 - bstr containing opaque application data

- H() - the hash function in HKDF_V

5.3.2. Party V Processing of Message 2

Party V SHALL compose message_2 as follows:

- Retrieve an ephemeral ECDH key pair generated as specified in Section 5 of [SP-800-56a] using same curve as used in E_U. Format the ephemeral public key E_V as a COSE_key as specified in Section 3.1.

- Generate the pseudo-random nonce N_V.

- Choose a session identifier S_V which is not in use and store it for the length of the protocol. The session identifier SHOULD be different from other relevant concurrent session identifiers used by Party V. The session identifier MAY be used with the protocol for which EDHOC establishes traffic keys/master secret, in which case S_V SHALL be different from the concurrently used session identifiers of that protocol.

- Select HKDF_V and AEAD_V from the algorithms proposed in HKDFs_U and AEADs_U.

- Format message_2 as specified in Section 5.3.1 where COSE_Encrypt0 is computed as defined in section 5.3 of [I-D.ietf-cose-msg], with AEAD_V, K_2, and IV_2.

5.3.3. Party U Processing of Message 2

Party U SHALL process message_2 as follows:

- Use the session identifier S_U to retrieve the protocol state.

- For EC2 curves, validate that E_V is a valid point by verifying that there is a solution to the curve definition for the given parameter 'x'.

- Verify message_2 as specified in Section 5.3.1 where COSE_Encrypt0 is decrypted defined in section 5.3 of [I-D.ietf-cose-msg], with AEAD_V, K_2, and IV_2.

If any verification step fails, Party U MUST send an EDHOC error message back, formatted as defined in Section 6.1, and the protocol MUST be discontinued.

- Pass APP_2 to the application.

5.4. EDHOC Message 3

5.4.1. Formatting of Message 3

message_3 SHALL be a CBOR array as defined below

```
message_3 = [
  data_3,
  COSE_ENC_3 : COSE_Encrypt0
]

data_3 = (
  MSG_TYPE : int,
  S_V : bstr
)

aad_3 : bstr
```

where aad_3, in non-CDDL notation, is:

```
aad_3 = H( H( message_1 | message_2 ) | [ data_3 ] )
```

where:

- MSG_TYPE = 6
- COSE_ENC_3 has the following fields and values:
  - external_aad = aad_3
  - plaintext = ? APP_3
- APP_3 - bstr containing opaque application data

5.4.2. Party U Processing of Message 3

Party U SHALL compose message_3 as follows:

- Format message_3 as specified in Section 5.4.1 where COSE_Encrypt0 is computed as defined in section 5.3 of [I-D.ietf-cose-msg], with AEAD_V, K_3, and IV_3.

5.4.3. Party V Processing of Message 3

Party V SHALL process message_3 as follows:

- Use the session identifier S_V to retrieve the protocol state.

- Verify message_3 as specified in Section 5.4.1 where COSE_Encrypt0 is decrypted and verified as defined in section 5.3 of [I-D.ietf-cose-msg], with AEAD_V, K_3, and IV_3.

If any verification step fails, Party V MUST send an EDHOC error message back, formatted as defined in Section 6.1, and the protocol MUST be discontinued.

- Pass APP_3 to the application.

6. Error Handling

6.1. Error Message Format

This section defines a message format for an EDHOC error message, used during the protocol. This is an error on EDHOC level and is independent of the lower layers used. An advantage of using such a construction is to avoid issues created by usage of cross protocol proxies (e.g. UDP to TCP).

error SHALL be a CBOR array as defined below

```
error = [
  MSG_TYPE : int,
  ? ERR_MSG : tstr
]
```

where:

- MSG_TYPE = 0

- ERR_MSG is an optional text string containing the diagnostic payload, defined in the same way as in Section 5.5.2 of [RFC7252].

7. IANA Considerations

7.1. Media Types Registry

IANA has added the media type ’application/edhoc’ to the Media Types registry:

Type name: application

Subtype name: edhoc

Required parameters: N/A

Optional parameters: N/A

Encoding considerations: binary

Security considerations: See Section 7 of this document.

Interoperability considerations: N/A

Published specification: [[this document]] (this document)

Applications that use this media type: To be identified

Fragment identifier considerations: N/A

Additional information:

* Magic number(s): N/A

* File extension(s): N/A

* Macintosh file type code(s): N/A

Person & email address to contact for further information:
Goeran Selander <goran.selander@ericsson.com>

Intended usage: COMMON

Restrictions on usage: N/A

Author: Goeran Selander <goran.selander@ericsson.com>

Change Controller: IESG

8. Security Considerations

EDHOC builds on the SIGMA-I family of theoretical protocols that provides perfect forward secrecy and identity protection with a minimal number of messages. The encryption algorithm of the SIGMA-I protocol provides identity protection, but the security of the protocol requires the MAC to cover the identity of the signer. Hence the message authenticating functionality of the authenticated encryption in EDHOC is critical: authenticated encryption MUST NOT be

replaced by plain encryption only, even if authentication is provided at another level or through a different mechanism.

EDHOC adds an explicit message type and expands the message authentication coverage to additional elements such as algorithms, application data, and previous messages. EDHOC uses the same Sign-then-MAC approach as TLS 1.3.

EDHOC does not include negotiation of parameters related to the ephemeral key, but it enables Party V to verify that the ECDH curve used in the protocol is the most preferred curve by U which is supported by both U and V.

Party U and V must make sure that unprotected data and metadata do not reveal any sensitive information. This also applies for encrypted data sent to an unauthenticated party. In particular, it applies to APP_1 and APP_2 in the asymmetric case, and APP_1 and KID in the symmetric case. The communicating parties may therefore anonymize KID.

Using the same KID or unprotected application data in several EDHOC sessions allows passive eavesdroppers to correlate the different sessions. Another consideration is that the list of supported algorithms may be used to identify the application.

Party U and V are allowed to select the session identifiers S_U and S_V, respectively, for the other party to use in the ongoing EDHOC protocol as well as in a subsequent traffic protection protocol (e.g. OSCOAP). The choice of session identifier is not security critical but intended to simplify the retrieval of the right security context in combination with using short identifiers. If the wrong session identifier of the other party is used in a protocol message it will result in the receiving party not being able to retrieve a security context (which will terminate the protocol) or retrieving the wrong security context (which also terminates the protocol as the message cannot be verified).

Party U and V must make sure that unprotected data does not trigger any harmful actions. In particular, this applies to APP_1 in the asymmetric case, and APP_1 and KID in the symmetric case. Party V should be aware that replays of EDHOC message_1 cannot be detected unless previous nonces are stored.

The availability of a secure pseudorandom number generator and truly random seeds are essential for the security of EDHOC. If no true random number generator is available, a truly random seed must be provided from an external source. If ECDSA is supported, "deterministic ECDSA" as specified in RFC6979 is RECOMMENDED.

Nonces MUST NOT be reused, both parties MUST generate fresh random nonces.

Ephemeral keys SHOULD NOT be reused, both parties SHOULD generate fresh random ephemeral key pairs. Party V MAY reuse the ephemeral key to limit the effect of certain DoS attacks. For example, to reduce processing costs in the case of repeated uncompleted protocol runs, party V MAY pre-compute its ephemeral key E_V and reuse it for a small number of concurrent EDHOC executions, for example until a number of EDHOC protocol instances has been successfully completed, which triggers party V to pre-compute a new ephemeral key E_V to use with subsequent protocol runs.

The referenced processing instructions in [SP-800-56a] must be complied with, including deleting the intermediate computed values along with any ephemeral ECDH secrets after the key derivation is completed.

Party U and V are responsible for verifying the integrity of certificates. The selection of trusted CAs should be done very carefully and certificate revocation should be supported.

The choice of key length used in the different algorithms needs to be harmonized, so that a sufficient security level is maintained for certificates, EDHOC, and the protection of application data. Party U and V should enforce a minimum security level.

Note that, depending on the application, the keys established through the EDHOC protocol will need to be renewed, in which case the communicating parties need to run the protocol again.

Implementations should provide countermeasures to side-channel attacks such as timing attacks.

9. Acknowledgments

The authors want to thank Dan Harkins, Ilari Liusvaara, Jim Schaad and Ludwig Seitz for reviewing intermediate versions of the draft and contributing concrete proposals incorporated in this version. We are especially indebted to Jim Schaad for his continuous reviewing and implementation of different versions of the draft.

TODO: This section should be after Appendices and before Authors' addresses according to RFC7322.

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Appendix A. Test Vectors

TODO: This section needs to be updated.

Appendix B. PSK Chaining

An application using EDHOC with symmetric keys may have a security policy to change the PSK as a result of successfully completing the EDHOC protocol. In this case, the old PSK SHALL be replaced with a new PSK derived using other = exchange_hash, AlgorithmID = "EDHOC PSK Chaining" and keyDataLength equal to the key length of AEAD_V, see Section 3.2.

Appendix C. EDHOC with CoAP and OSCOAP

C.1. Transferring EDHOC in CoAP

EDHOC can be transferred as an exchange of CoAP [RFC7252] messages, with the CoAP client as party U and the CoAP server as party V. By default EDHOC is sent to the Uri-Path: "/.well-known/edhoc", but an application may define its own path that can be discovered e.g. using resource directory [I-D.ietf-core-resource-directory].

In practice, EDHOC message_1 is sent in the payload of a POST request from the client to the server’s resource for EDHOC. EDHOC message_2 or the EDHOC error message is sent from the server to the client in the payload of a 2.04 Changed response. EDHOC message_3 or the EDHOC error message is sent from the client to the server’s resource in the payload of a POST request. If needed, an EDHOC error message is sent from the server to the client in the payload of a 2.04 Changed response

An example of successful EDHOC exchange using CoAP is shown in Figure 5.

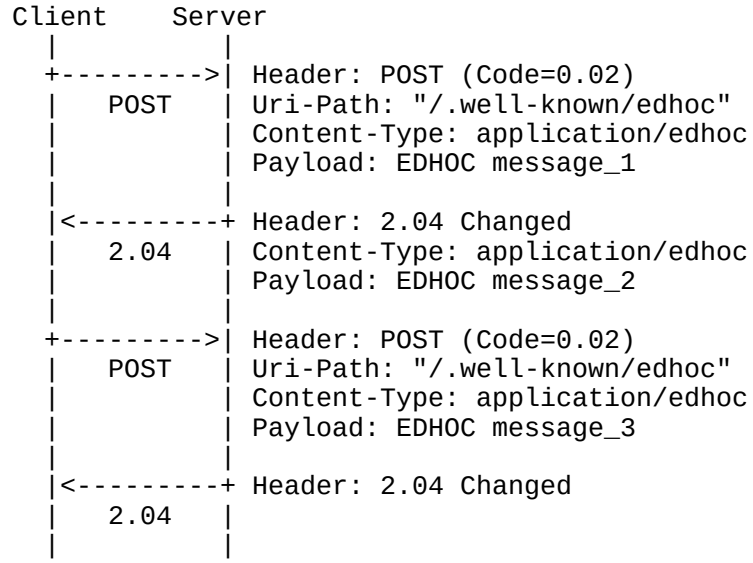

Figure 5: Transferring EDHOC in CoAP

C.2. Deriving an OSCOAP context from EDHOC

When EDHOC is use to derive parameters for OSCOAP [I-D.ietf-core-object-security], the parties must make sure that the EDHOC session identifiers are unique Recipient IDs in OSCOAP. In case that the CoAP client is party U and the CoAP server is party V:

- The AEAD Algorithm is AEAD_V, as defined in this document
- The KDF algorithm is HKDF_V, as defined in this document
- The Client’s Sender ID is S_V, as defined in this document
- The Server’s Sender ID is S_U, as defined in this document
- The Master Secret is derived as specified in Section 3.2 of this document, with other = exchange_hash, AlgorithmID = "EDHOC OSCOAP Master Secret" and keyDataLength equal to the key length of AEAD_V.
- The Master Salt is derived as specified in Section 3.2 of this document, with other = exchange_hash, AlgorithmID = "EDHOC OSCOAP Master Salt" and keyDataLength equal to 64 bits.

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January 25, 2017

MQTT-TLS profile of ACE
draft-sengul-ace-mqtt-tls-profile-00

Abstract

This document specifies a profile for the ACE (Authentication and Authorization for Constrained Environments) to enable authorization in an MQTT-based publish-subscribe messaging system. Proof-of-possession keys, bound to OAuth2.0 access tokens, are used to authenticate and authorize publishing and subscribing clients. The protocol relies on TLS for confidentiality and server authentication.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

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This Internet-Draft will expire on July 29, 2017.

Copyright Notice

Table of Contents

1. Introduction

This document specifies a profile for the ACE framework [I-D.ietf-ace-oauth-authz]. In this profile, clients and a resource server use MQTT to communicate. The protocol relies on TLS for communication security between entities. Protocol interactions follow MQTT v3.1 - OASIS Standard [MQTT-OASIS-Standard]. Future releases may enable improvements to the protocol operation (e.g., by allowing MQTT message return codes for authorization errors).

MQTT is a publish-subscribe protocol, and supports two types of client operation: publish and subscribe. Once connected, a client can publish to multiple topics, and subscribe to multiple topics; however for the purpose of this document these actions are described separately. The MQTT broker is responsible for distributing messages published by the publishers to the appropriate subscribers. Each publish message contains a topic, which is used by the broker to filter the subscribers for the message. Subscribers must subscribe to the topics to receive the corresponding messages.

In this document, message topics are treated as resources. Both publisher and subscriber clients use an access token, each bound to a key (the proof-of-possession key) to authorize with the MQTT broker their connection and publish/subscribe permissions to topics. In the context of this ACE profile, the MQTT broker acts as the resource server. In order to provide communication confidentiality and resource server authentication, TLS is used.

The publisher and subscriber clients use client authorization servers [I-D.ietf-ace-actors] to obtain tokens from the authorization server. The communication protocol between the client authorization server and the authorization server is assumed to be HTTPS. Also, if the broker supports token introspection, it is assumed to use HTTPS to communicate with the authorization server. These interfaces MAY be implemented using other protocols e.g., CoAP or MQTT. This document makes the same assumptions as the Section 4 of the ACE framework [I-D.ietf-ace-oauth-authz] in terms of client and RS registration with the AS and establishing of keying material.

This document describes authorization of the following exchanges between publisher and subscriber clients, and the broker.

- Connection establishment between the clients and the broker
- Publish messages from the publishers to the broker, and from the broker to the subscribers
- Subscribe messages from the subscribers to the broker

1.1. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC 2119 [RFC2119].

1.2. ACE-Related Terminology

The terminology for entities in the architecture is defined in OAuth 2.0 RFC 6749 [RFC6749] and ACE actors [I-D.ietf-ace-actors], such as "Client" (C), "Resource Server" (RS) and "Authorization Server" (AS).

The term "endpoint" is used following its OAuth definition, to denote resources such as /token and /introspect at the AS.

The term "Resource" is used to refer to an MQTT "topic", which is defined in Section 1.2. Hence, the "Resource Owner" is any entity that can authoritatively speak for the "topic".

Certain security-related terms such as "authentication", "authorization", "confidentiality", "(data) integrity", "message authentication code", and "verify" are taken from RFC 4949 [RFC4949].

1.3. MQTT-Related Terminology

The document describes message exchanges as MQTT protocol interactions. For additional information, please refer to the MQTT v3.1 - OASIS Standard [MQTT-OASIS-Standard].

Topic name
: The label attached to an application message, which is matched to a subscription.

Topic filter
: An expression that indicates interest in one or more topic names. Topic filters may include wildcards.

Subscription
: A subscription comprises of a Topic filter and a maximum quality of service (QoS).

Application Message
: The data carried by the MQTT protocol. The data has an associated QoS level and a Topic name.

MQTT sends various control messages across a network connection. The following is not an exhaustive list, and the control packets that are not relevant for authorization are not explained. These include, for instance, the PUBREL and PUBCOMP packets used in the 4-step handshake required for the QoS level 2.

CONNECT

Client request to connect to the broker. After a network connection is established, this is the first packet sent by a client.

CONNACK
: The broker connection acknowledgment. The first packet sent from broker to a client is a CONNACK packet. CONNACK packets contain return codes indicating either a success or an error state to the client.

PUBLISH
: Publish packet that can be sent from a client to the broker, or from the broker to the client.

PUBACK
: Response to PUBLISH packet with QoS level 1. PUBACK can be sent from the broker to the client or the client to the broker.

PUBREC
: Response to PUBLISH packet with QoS level 2. PUBREC can be sent from the broker to the client or the client to the broker.

SUBSCRIBE
: The client subscribe request.

SUBACK
: Subscribe acknowledgment.

2. Protocol Interactions

This document describes the following exchanges between publisher and subscriber clients, the broker, and the authorization server.

- Authorizing connection establishment between the clients and the broker
- Authorizing publish messages from the publishers to the broker, and from the broker to the subscribers
- Authorizing subscribe messages from the subscribers to the broker

Message topics are treated as resources. The publisher and subscriber clients are assumed to have identified the topics of interest out-of-band (topic discovery is not a feature of the MQTT protocol).

A connection request carries a token specifying the permissions that the client has (e.g., publish permission to a given topic). A resource owner can pre-configure policies at the AS that give clients publish or subscribe permissions to different topics.

2.1. Authorizing Connection Establishment

This section specifies how publishers and subscribers establish an authorized connection to an MQTT broker. The token request and response use the /token endpoint of the authorization server, as specified in Section 6 of the the ACE framework [I-D.ietf-ace-oauth-authz].

Figure 1 shows the basic protocol flow during connection establishment.

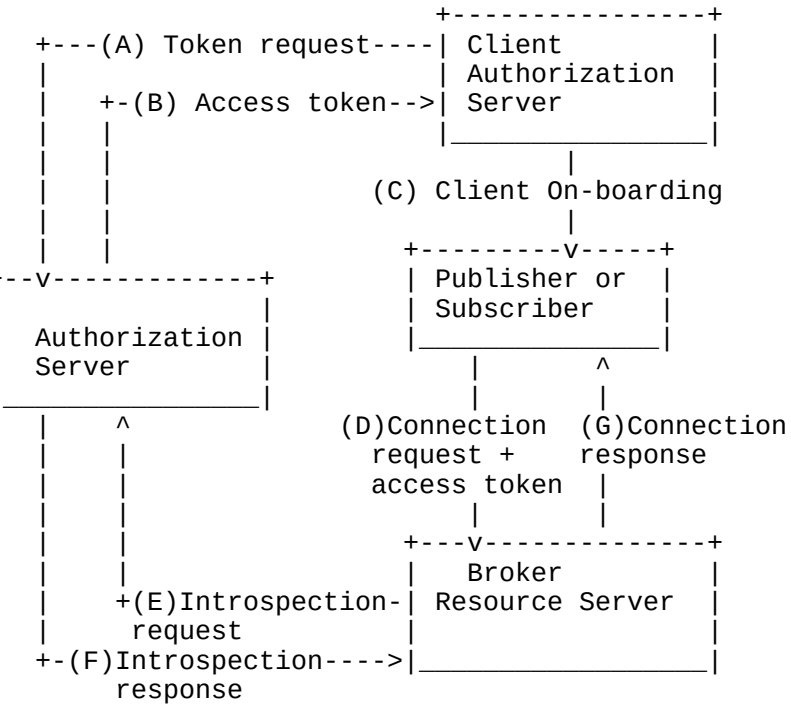

Figure 1: Connection establishment

2.1.1. Client Authorization Server (CAS) to Authorization Server (AS)

The first step in the protocol flow is token acquisition by the client authorization server (CAS) from the AS. If a client has

enough resources and can support HTTPS, or optionally the AS supports MQTT, these steps can instead be carried out by a client directly.

When requesting an access token from the AS, the CAS MAY include parameters in its request as defined in Section 6.1 of the ACE framework [I-D.ietf-ace-oauth-authz]. The content type is set to "application/json".

The response contains a token and a 'cnf' parameter with a symmetric or asymmetric proof-of-possession (PoP) key.

The token request is similar to the examples presented in Section 6.1. of the ACE framework [I-D.ietf-ace-oauth-authz] with a modified profile name 'mqtt_tls'.

2.1.2. Authorization Server (AS) to Client Authorization Server (CAS)

If the access token request has been successfully verified by the AS and the client is authorized to obtain a PoP token for the indicated audience (i.e., broker) and scopes (i.e., publish/subscribe to the requested topics), the AS issues an access token. The response includes the parameters described in Section 6.2 of the ACE framework [I-D.ietf-ace-oauth-authz].

In the case of an error, the AS returns error responses for HTTP-based interactions as ASCII codes in JSON content, as defined in Section 5.2 of RFC 6749 [RFC6749].

2.1.3. Client connection request to the broker

Having received the token, the client can use it to request an MQTT connection to the broker over a TLS session with server authentication. This document describes the client transporting the token to the broker (RS) via the CONNECT control message after the TLS handshake. This is similar to an earlier proposal by Freemantle et al. [freemantle14]. Alternatively, the token may be used for the TLS session establishment as described in the DTLS profile for ACE [I-D.gerdes-ace-dtls-authorize]. In this case, both the TLS PSK and RPK handshakes MAY be supported. This may additionally require that the client transports the token to the broker before the connection establishment. To this end, the broker MAY allow clients to publish to "authz-info" topic unauthorized, and in this case, "authz-info" topic SHOULD be publish only (i.e., the clients are not allowed to subscribe to it). Implementation of the public "authz-info" topic is discussed in Appendix B.

When the client wishes to connect to the broker, it uses the CONNECT message of MQTT. Figure 2 shows the structure of the MQTT CONNECT control message.

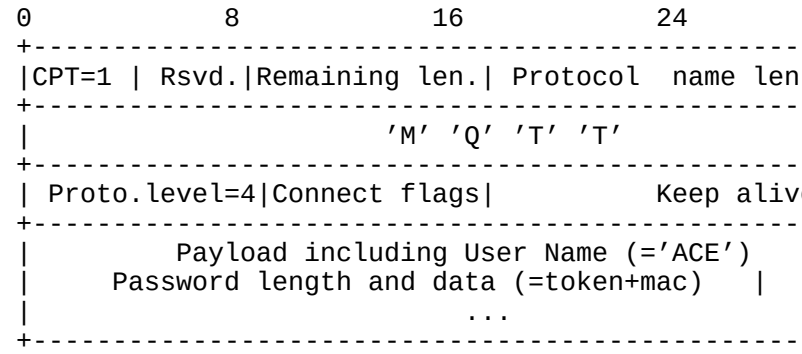

Figure 2: MQTT CONNECT control message. (CPT=Control Packet Type, Rsvd=Reserved, len.=length, Proto.=Protocol)

To communicate the necessary connection parameters, the Client uses the appropriate flags of the CONNECT message. Figure 3 shows how the MQTT connect flags MUST be set to initiate a connection with the broker.

User name flag	Pass. flag	Will retain	Will QoS	Will Flag	Clean	Rsvd.
1	1	X	X X	X	1	0

Figure 3: MQTT CONNECT flags. (Rsvd=Reserved)

In order to ensure that the client and the broker discard any previous session and start a new session, the Clean Session Flag MUST be set to 1.

The Will flag indicates that a Will message needs to be sent when a client disconnection occurs. The situations in which the Will message is published include disconnections due to I/O or network failures, and the server closing the networking connection due to a protocol error. The client may set the Will flag as desired (marked as 'X' in Figure 3). If the Will flag is set to 1 and the broker accepts the connection request, the broker must store the Will message, and publish it when the network connection is closed. The Will QoS specifies the QoS level of the Will message, and may be set to 0x00, 0x01 or 0x02. The Will retain flag may be set as desired.

If it is set to 1, the Will message will be delivered to all future subscribers whose subscriptions match the Will topic. Section 2.6 explains how the broker deals with the RETAINed messages in further detail.

Finally, Username and Password flags MUST be set to 1, which ensures that the Payload of the CONNECT message includes both Username and Password fields.

The Username and Password field are used to indicate to the resource server that the CONNECT message is carrying an ACE token and the MAC of the request. To this end, the Username is set to 'ACE', and the Password field is populated with a JSON object containing the token and the MAC. The Password length field MUST be set to the size of the JSON object. (The maximum size of the password field is defined as 65535 bytes by MQTT v3.1 - OASIS Standard [MQTT-OASIS-Standard].)

Figure 4 shows an example for setting the password field in an MQTT CONNECT message.

```
{
"token": b64'SLAV32hkKG....,
"mac": b64'kDZvddkndxvhGRXZhvuDjEWhGeE=,
}
```

Figure 4: Example token and MAC as password data in CONNECT message.

2.1.4. Token validation

RS MUST verify the validity of the token. This validation MAY be done locally or the RS MAY send an introspection request to the AS. If introspection is used, this section follows similar steps to those described in Sections 7.2 and 7.3 of the ACE framework [I-D.ietf-ace-oauth-authz]. The communication between AS and RS MAY be HTTPS, but it, in every case, MUST be confidential, mutually authenticated and integrity protected.

The broker MUST check if the token is active either using 'expires_in' parameter of the token or 'active' parameter of the introspection response.

The access token is constructed by the AS such that RS can associate the access token with the client key. This document assumes that the Access Token is a PoP token as described in [I-D.ietf-ace-oauth-authz]. Therefore, the necessary information is contained in the 'cnf' claim of the access token, and may use either public or shared key approaches. The client uses the 'mac' parameter in the password field to prove the possession of the key. The

resource server validates the 'mac' over the contents of the packet, authenticating the client.

The broker uses the scope field in the token (or in the introspection result) to determine the publish and subscribe permissions for the client. Scope strings MAY follow an application specific convention e.g., 'publish_topic1' or 'subscribe_topic2'. If the Will flag is set, then the broker MUST check that the token allows the publication of the Will message too.

The broker MAY cache the introspection result, because it will need to decide whether to accept subsequent PUBLISH and SUBSCRIBE messages, and these messages, which are sent after a connection is set-up, may not contain tokens. If the introspection result is not cached, then the RS needs to introspect the saved token for each request.

2.1.5. The broker's response to client connection request

Based on the validation result (obtained either via local inspection or using the /introspection interface of the authorization server), the broker MUST send a CONNACK message to the client.

The following responses may be returned to the client.

- If the broker accepts the connection, the broker MUST send a CONNACK message with Return Code 0x00 indicating 'Connection Accepted'.
- If the connection is denied, the broker MUST send a CONNACK message with 0x05 indicating 'Connection Refused, not authorized'.
- If the data in the user name or password is malformed, the broker MUST send a CONNACK message with 0x04 indicating 'Connection Refused, bad user name or password'.

It is not possible to support AS discovery via sending a tokenless CONNECT message to the broker. This is because a CONNACK packet does not have a payload. Therefore, AS discovery is assumed to have taken place out-of-band.

If the RS accepts the connection, it MUST store the token.

2.2. Authorizing PUBLISH messages

Figure 5 shows the PUBLISH message used in MQTT, which includes fixed and variable headers.

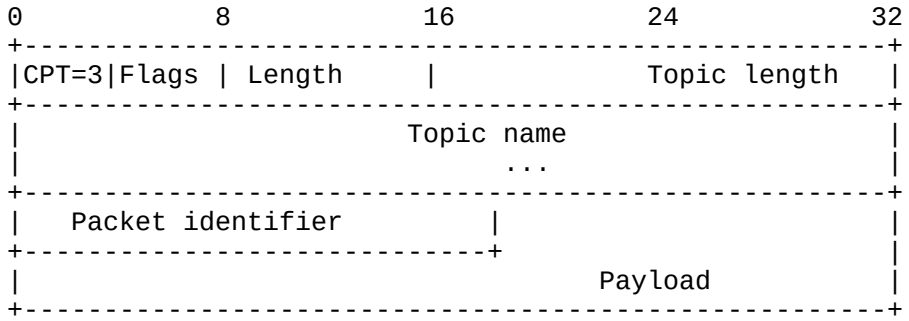

Figure 5: MQTT PUBLISH control message. (CPT=Control Packet Type)

The variable header includes flags for QoS and RETAIN. If RETAIN is set to 1, then the broker must store the most recent application message per topic, and its QoS, to forward to future subscribers. Other fields in the PUBLISH header are topic name and packet identifier. The topic name MUST NOT include wildcard characters according to MQTT v3.1 - OASIS Standard [MQTT-OASIS-Standard].

2.2.1. PUBLISH messages from the publisher client to the broker

The payload of PUBLISH messages contains an application message. The content and the format of the data is application specific. Therefore, the client MAY include its token inside the PUBLISH messages. The token could for example be included as:

```
{
"message":"topic message",
"token": b64'SLAV32hkKG....,
}
```

Figure 6: Example token in PUBLISH payload.

If the application message contains a token, the broker MAY locally inspect the token or MAY use the /introspect interface of the authorization server. The token received in the PUBLISH message MAY be different than the one stored after connection handshake. On receiving a new token, the RS discards any previously stored token for the client and MUST store the new token as not all PUBLISH messages may carry tokens.

If the application message does not contain a token, the broker MUST use the type of message (i.e., PUBLISH) and the topic name in the message header to compare against the 'scope' field of the cached introspection result for the client.

If the client is allowed to publish to the topic, the broker may return an acknowledgment message. This is determined by the QoS flags in the Flags field of the PUBLISH header. If QoS level is 0, the RS does not send back a response. If the QoS level is equal or greater than 1, the RS must respond with an acknowledgement message (i.e., PUBACK for QoS=1 or PUBREC for QoS=2). These messages can currently only indicate success, and there is no equivalent 'NACK' to indicate failure. Next, the RS must publish the message to all valid subscribers to the topic.

In the case of an authorization error, the broker SHOULD disconnect the client. Otherwise, it MUST silently ignore the message. In the MQTT v3.1 - OASIS Standard [MQTT-OASIS-Standard], the MQTT DISCONNECT messages are only sent from a client to the broker. So, server disconnection needs to take place below the application layer. Appendix C describes an alternative method for handling authorization errors, possibly avoiding disconnections.

2.2.2. PUBLISH messages from the broker to the subscriber clients

To forward PUBLISH messages to the subscribing clients, the broker identifies all the subscribers that have matching valid topic subscriptions (i.e., the tokens are valid and token scopes allow a subscription to the particular topic name). The broker sends a PUBLISH message with the topic name and the topic message to all the valid subscribers.

In MQTT, after connection establishment there is no way to inform a client that an authorization error has occurred for previously subscribed topics, e.g., token expiry. In the case of an authorization error, the broker has two options: (1) stop forwarding PUBLISH messages to the unauthorized client or (2) disconnect the client. In the MQTT v3.1 - OASIS Standard [MQTT-OASIS-Standard], the MQTT DISCONNECT messages are only sent from a client to the broker. Therefore, the server disconnection needs to take place below the application layer. Appendix C describes an alternative method, where disconnections may be avoided.

2.3. Authorizing SUBSCRIBE messages

In MQTT, a SUBSCRIBE message is sent from a client to the broker, to create one or more subscriptions to one or more topics.

Figure 7 shows the MQTT SUBSCRIBE message format. The SUBSCRIBE message may contain multiple topic filters. The topic filters may include wildcard characters.

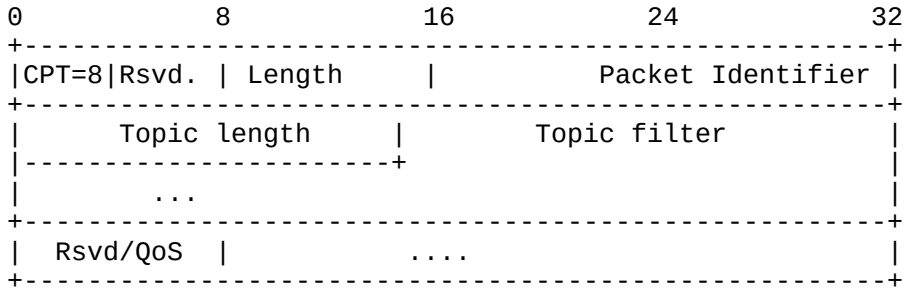

Figure 7: MQTT SUBSCRIBE control message. (CPT=Control Packet Type, Rsvd.=Reserved, QoS=Quality of Service)

The SUBSCRIBE message does not have any field suitable for including a token. Therefore, on receiving the SUBSCRIBE message, the broker MUST use the type of message (i.e., SUBSCRIBE) and the topic name in the message header to compare against the 'scope' field of the stored token or introspection result.

As a response to the SUBSCRIBE message, the broker issues a SUBACK message. For each topic filter, the SUBACK packet includes a return code. In the case of success, the return code must be either 0x00, 0x01 or 0x02, matching the QoS level for the corresponding topic filter.

In the case of failure, the return code must be 0x08 indicating 'Failure'. There is no other way to signal the reason for an authorization failure to the Subscriber, or to communicate further detail. Appendix C describes an alternative method where more detailed error messages can be provided to the client.

2.4. Token expiration

The broker checks for token expiration whenever a CONNECT, PUBLISH or SUBSCRIBE message is received or sent. The validation is done either by checking the 'exp' claim of a CWT/JWT or via performing an introspection request with the Authorization server as described in the Section 8.2 of the ACE framework [I-D.ietf-ace-oauth-authz]. Token expirations leads to disconnecting the associated client. Appendix C describes an alternative method, where clients are allowed to update tokens, avoiding disconnections.

2.5. Handling disconnections

According to MQTT v3.1 - OASIS Standard [MQTT-OASIS-Standard], only Client DISCONNECT messages are allowed. (This is expected to change in the future, enabling Server DISCONNECT messages.) In the case of

a Client DISCONNECT, due to the Clean Session flag, the broker deletes all session state but MUST keep the retained messages and send them according to methodology described in Section 2.6. The broker MUST continue publishing the retained messages as long as the associated tokens are valid.

In case of disconnections due to network errors, or Server disconnection due to a protocol error (which includes the authorization errors), the Will message must be sent if the client supplied a Will in the CONNECT request message (see Figure 3). According to the [MQTT-OASIS-Standard], if the CONNECT request is accepted, then any WILL message must be stored. The Will message must be published to the Will topic when the network connection is closed.

2.6. Handling retained messages

The broker treats retained messages according to the [MQTT-OASIS-Standard]. By setting a RETAIN flag in a PUBLISH message (or in a CONNECT message for the Will message), the publisher indicates to the broker that it should store the most recent message for the associated topic, so the broker can send the message to any future subscribers. Hence, the new subscribers can receive the last sent message from the publisher for that particular topic, without waiting for the next PUBLISH message.

According to the [MQTT-OASIS-Standard], if a publisher client disconnects, the retained messages do not form part of the session state and must not be deleted by the broker (since the Clean session flag set to 1 during the connection request, othersession state is deleted). Therefore, the broker MUST continue publishing retained messages for as long as the stored token for the client is valid: This applies to both the PUBLISH and WILL messages. However, if the disconnection is triggered by the broker due to an authorisation error, the broker MUST stop publishing all retained messages from that client.

3. IANA Considerations

This memo includes no request to IANA.

4. Security Considerations

TBD.

5. Privacy Considerations

TBD.

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Appendix A. Checklist for profile requirements

- AS discovery: The clients/client authorization servers need to be configured out-of-band. RS does not provide any hints to help AS discovery.
- Communication protocol between the client and RS: MQTT
- Security protocol between the client and RS: TLS
- Client and RS mutual authentication: RS provides a server certificate during TLS handshake. Client uses token and MAC fields in the MQTT connect message.
- Content format: For the HTTPS interactions with AS, "application/json". The MQTT payloads may be formatted JSON or CBOR.
- PoP protocols: Either symmetric or asymmetric keys can be supported.
- Unique profile identifier: mqtt_tls
- Token introspection: RS uses HTTPS /introspect interface of AS.
- Token request: CAS uses HTTPS /token interface of AS.
- /authz-info endpoint: It MAY be supported using the method described in Appendix B, not protected.
- Token transport: In MQTT CONNECT message or using the method described in Appendix B.

Appendix B. The ‘authorization information’ endpoint

The main document described a method where the access token is transported inside the MQTT CONNECT message. In this section, we describe an alternative method to transport the access token.

The method consists of the MQTT broker providing a public "authz-info" topic. A client using this method MUST first connect to the

broker, and publish the access token using the "authz-info" topic. The broker must verify the validity of the token (i.e., through local validation or introspection). After publishing the token, the client disconnects from the broker and is expected to try reconnecting over TLS.

After the client published to the "authz-info" topic, it is not possible for the broker to communicate the result of the verification. The response to a PUBLISH message may be a PUBACK or PUBREC, and these messages indicate successful reception of the PUBLISH message and cannot communicate authorization errors. However, the token failure will affect the TLS handshake, which may be used to prompt the client to obtain a valid token. In Appendix C, an alternative method for error handling is discussed.

Appendix C. Error handling and token updates

Section 2.1.3 uses the CONNECT message to transfer the PoP token to the broker. This is simple, with only two states: 'Disconnected' and 'Authorized' (see Figure 8). However, the result of an authorization error is a server side disconnection without any feedback or error message.

	State	
	0	1
Events	Disconnected	Authorized
CONNECT success	1	-
CONNECT failure	0	-
Authorization expires	-	0
other disconnection		

Figure 8: Client state machine - simple token transport

To enable token updates during the lifetime of a connection, and also to allow the broker to send error messages to a client, this section proposes using a client-specific "authz-info-${ClientId}" topic. This case requires three states: 'Disconnected', 'Connected' and 'Authorized', shown in Figure 9.

In the Disconnected state, as before, the client needs to transport the token and attempt to establish a connection. Token transport MAY be done using any of the methods mentioned in this document. The CONNECT payload SHOULD include a unique client identifier: Although in MQTT a broker may accept 0-byte client identifiers, in this use case the client would not be aware of its client identifier, so would be unable to update its token, or subscribe to the "authz-info-${ClientId}" topic to receive error messages.

If the CONNECT succeeds, then the client moves to the Authorized state. It SHOULD also subscribe to the topic "authz-info-${ClientId}" to be able to receive authorization errors. The subscription MUST fail if the topic name does not contain the ClientId established during the CONNECT handshake.

If the token validation fails,the broker MUST publish an authorization error to "authz-info-${ClientId}" and the client moves to the Connected state. In this state, the client MAY publish a new token or it MAY disconnect. If a new token is published, then the broker MUST verify the token and send an authorization response to "authz-info-${ClientId}" indicating success or failure. In case of success, the client moves to the "Authorized" state. In the case of failure, the client remains in "Connected" state but will not be able to publish or receive message from its subscribed topics due to authorisation problems. It MAY be able to publish/subscribe to public topics.

	State		
	0	1	2
Events	Disconnected	Connected	Authorized
CONNECT success	2	-	-
CONNECT failure	0	-	-
PUBLISH new token success	-	2	2
PUBLISH new token failure	-	1	1
Authorization expires	-	-	1
Other disconnection	-	0	0

Figure 9: Client state machine - with error handling and token update

While this solution allows better authorization error feedback, it is a more complex solution. The broker needs to maintain separate authz-info topics for separate clients.

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Joining of OSCOAP multicast groups in ACE
draft-tiloca-ace-oscoap-joining-00

Abstract

This document describes a method to join a multicast group where communications are based on CoAP and secured with Object Security of CoAP (OSCOAP). The proposed method delegates the authentication and authorization of client nodes that join a multicast group through a Group Manager server. This approach builds on the ACE framework for Authentication and Authorization, and leverages protocol-specific profiles of ACE to achieve communication security, proof-of-possession and server authentication.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

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This Internet-Draft will expire on January 3, 2018.

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Table of Contents

1. Introduction

The Constrained Application Protocol (CoAP) [RFC7252] supports also group communication scenarios, where request messages can be delivered to multiple recipients using CoAP on top of IP multicast [RFC7390].

Object Security of CoAP (OSCOAP) [I-D.ietf-core-object-security] is a method for application layer protection of CoAP messages, using the CBOR Object Signing and Encryption (COSE) [I-D.ietf-cose-msg], and enabling end-to-end security of CoAP payload and options.

OSCOAP may also be used to protect group communication for CoAP over IP multicast, as described in [I-D.tiloca-core-multicast-oscoap]. This relies on a Group Manager entity, which is responsible for managing a multicast group where members exchange CoAP messages secured with OSCOAP. In particular, the Group Manager coordinates the join process of new group members and can be responsible for multiple groups.

This document builds on the ACE framework for Authentication and Authorization [I-D.ietf-ace-oauth-authz] and specifies how a client joins an OSCOAP multicast group through a resource server acting as Group Manager. The client acting as joining node relies on an Access Token, which is bound to a proof-of-possession key and authorizes the access to a specific join resource at the Group Manager.

The client and the Group Manager leverage protocol-specific profiles of ACE such as [I-D.seitz-ace-oscoap-profile] and [I-D.ietf-ace-dtls-authorize], in order to achieve communication security, proof-of-possession and server authentication.

1.1. Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119]. These words may also appear in this document in lowercase, absent their normative meanings.

Readers are expected to be familiar with the terms and concepts described in the ACE framework for authentication and authorization [I-D.ietf-ace-oauth-authz]. Message exchanges are presented as RESTful protocol interactions, for which HTTP [RFC7231] provides useful terminology.

The terminology for entities in the considered architecture is defined in OAuth 2.0 [RFC6749] and [I-D.ietf-ace-actors]. In particular, this includes client (C), resource server (RS), and authorization server (AS). Terminology for constrained environments, such as "constrained device" and "constrained-node network", is defined in [RFC7228].

Readers are expected to be familiar with the terms and concepts related to the CoAP protocol described in [RFC7252][RFC7390]. Note that the term "endpoint" is used here following its OAuth definition, aimed at denoting resources such as /token and /introspect at the AS and /authz-info at the RS. This document does not use the CoAP definition of "endpoint", which is "An entity participating in the CoAP protocol".

Readers are expected to be familiar with the terms and concepts for protection and processing of CoAP messages through OSCOAP [I-D.ietf-core-object-security] also in group communication contexts [I-D.tiloca-core-multicast-oscoap]; and with the OSCOAP profile of ACE described in [I-D.seitz-ace-oscoap-profile].

Readers are expected to be familiar with the terms and concepts related to the DTLS protocol [RFC6347]; the support for DTLS handshake based on Raw Public Keys (RPK) [RFC7250] and on Pre-Shared Keys (PSK) [RFC4279]; and the CoAP-DTLS profile of ACE [I-D.ietf-ace-dtls-authorize].

This document refers also to the following terminology.

- Joining node: a network node intending to join an OSCOAP multicast group, where communication is based on CoAP [RFC7390] and secured with OSCOAP as described in [I-D.tiloca-core-multicast-oscoap].

- Join process: the process through which a joining node becomes a member of a multicast group. The join process is enforced and assisted by the Group Manager responsible for that group.

- Join resource: a protected resource hosted by the Group Manager, associated to a multicast group under that Group Manager. A joining node accesses the join resource in order to start the join process and become a member of that group.

- Join endpoint: an endpoint hosted by the Group Manager associated to a join resource.

2. Protocol Overview

Group communication for CoAP over IP multicast has been enabled in [RFC7390] and can be secured with Object Security of CoAP (OSCOAP) [I-D.ietf-core-object-security] as described in [I-D.tiloca-core-multicast-oscoap]. A network node explicitly joins an OSCOAP multicast group, by interacting with the responsible Group Manager. Once registered in the group, the new node can securely exchange (multicast) messages with other group members.

This specification describes how a network node joins an OSCOAP multicast group leveraging the ACE framework for authentication and authorization [I-D.ietf-ace-oauth-authz]. With reference to the ACE framework and the terminology defined in OAuth 2.0 [RFC6749]:

- The Group Manager acts as Resource Server (RS), and owns one join resource for each OSCOAP multicast group it manages. Each join resource is exported by a distinct join endpoint.

- The joining node acts as Client (C), and requests to join an OSCOAP multicast group by accessing the related join endpoint at the Group Manager.

- The Authorization Server (AS) enables and enforces the authorized access of joining nodes to join endpoints at the Group Manager. Multiple Group Managers can be associated to the same AS.

If authorized to join the multicast group, the joining node receives from the AS an Access Token bound with a proof-of-possession key. After that, the joining node provides the Group Manager with the Access Token. This step involves the opening of a secure

communication channel between the joining node and the Group Manager, in case they have not already established one.

Finally, the joining node accesses the join endpoint at the Group Manager, so starting the join process to become a member of the multicast group. A same Access Token can authorize the joining node to access multiple groups under the same Group Manager. In such a case, the joining node sequentially performs multiple join processes with the Group Manager, separately for each multicast group to join and by accessing the respective join endpoint.

The AS is not necessarily expected to release Access Tokens for any other purpose than accessing join resources on registered Group Managers. In particular, the AS is not necessarily expected to release Access Tokens for accessing protected resources at members of multicast groups.

The following steps are performed for joining an OSCOAP multicast group, by leveraging the CoAP-DTLS profile of ACE [I-D.ietf-ace-dtls-authorize] or the OSCOAP profile of ACE [I-D.seitz-ace-oscoap-profile].

1. The joining node retrieves an Access Token from the AS to access a join resource on the Group Manager (Section 3). The response from the AS enables the joining node to start a secure channel with the Group Manager, if not already established. The joining node can also contact the AS for updating a previously released Access Token, in order to access further groups under the same Group Manager (Section 6).

2. Authentication and authorization information is transferred between the joining node and the Group Manager, which establish a secure channel in case one is not already set up (Section 4). That is, a joining node MUST establish a secure communication channel with a Group Manager, before joining a multicast group under that Group Manager for the first time.

3. The joining node starts the join process to become a member of the multicast group, by accessing the related join resource hosted by the Group Manager (Section 4).

All communications between the involved entities rely on the CoAP protocol and MUST be secured. In particular, communications between the joining node and the AS (/token endpoint) and between the Group Manager and the AS (/introspection endpoint) can be secured by different means, e.g. with DTLS [RFC6347] or with OSCOAP (see Sections 3 and 4 of [I-D.seitz-ace-oscoap-profile]).

3. Joining Node to Authorization Server

This section considers a joining node that intends to contact the Group Manager for the first time. That is, the joining node has never attempted before to join a multicast group under that Group Manager. Also, the joining node and the Group Manager do not have a secure communication channel established.

In case the specific AS associated to the Group Manager is unknown to the joining node, the latter can rely on mechanisms like the one described in Section 2.2 of [I-D.ietf-ace-dtls-authorize] to discover the correct AS in charge of the Group Manager.

The joining node contacts the AS, in order to request an Access Token for accessing the join resource(s) hosted by the Group Manager. In particular, the Access Token request sent to the /token endpoint specifies the join endpoint(s) of interest at the Group Manager.

The AS is responsible for authorizing the joining node, accordingly to group join policies enforced on behalf of the Group Manager. In case of successful authorization, the AS releases an Access Token bound to a proof-of-possession key associated to the joining node. The same Access Token can authorize the joining node to access multiple groups under the same Group Manager.

Then, the AS provides the joining node with the Access Token, together with an Access Token response. In particular, the Access Token response indicates how to secure communications with the Group Manager, when accessing the join resource(s) for which the Access Token is valid. Specifically, the Access Token response MUST specify one of the following alternatives:

- CoAP over DTLS, i.e. coaps://, indicating to consider the CoAP-DTLS profile of ACE, with asymmetric or symmetric proof-of-possession key (see Section 3 and Section 4 of [I-D.ietf-ace-dtls-authorize], respectively).

- OSCOAP, indicating to consider the OSCOAP profile of ACE with asymmetric or symmetric proof-of-possession key, as described in Section 2.2 of [I-D.seitz-ace-oscoap-profile].

Consistently with the profiles of ACE specified in [I-D.ietf-ace-dtls-authorize] and [I-D.seitz-ace-oscoap-profile], a symmetric proof-of-possession key is generated by the AS, which uses it as proof-of-possession key bound to the Access Token, and provides it to the joining node in the Access Token response. Instead, in case of asymmetric proof-of-possession key, the joining node provides its own public key to the AS in the Access Token request. Then, the

AS uses it as proof-of-possession key bound to the Access Token, and provides the joining node with the Group Manager's public key in the Access Token response.

4. Joining Node to Group Manager

First, the joining node establishes a secure channel with the Group Manager, according to what is specified in the Access Token response. In particular:

- If the CoAP-DTLS profile of ACE is specified, the joining node MUST upload the Access Token to the /authz-info resource before starting the DTLS handshake. Then, the Group Manager processes the Access Token according to [I-D.ietf-ace-oauth-authz]. If this yields to a positive response, the joining node and the Group Manager establish a DTLS session, as described in Section 3 and Section 4 of [I-D.ietf-ace-dtls-authorize], in case of either asymmetric or symmetric proof-of-possession key, respectively.

- If the OSCOAP profile of ACE is specified, the joining node and the Group Manager establish an OSCOAP channel, as described in Section 2.2 of [I-D.seitz-ace-oscoap-profile]. In particular, if the EDHOC protocol [I-D.selander-ace-cose-ecdhe] is used to this end, the joining node MUST include the Access Token in the EDHOC message_1 sent to the /authz-info resource. The Group Manager processes the Access Token as specified in [I-D.ietf-ace-oauth-authz] and proceeds as defined in Section 2.2 of [I-D.seitz-ace-oscoap-profile].

Once the secure channel with the Group Manager has been established, the joining node requests to join the OSCOAP multicast groups of interest, by accessing the related join resources at the Group Manager. That is, the joining node performs multiple join processes with the Group Manager, separately for each multicast group to join and by accessing the respective join endpoint.

In particular, for each multicast group to join, the joining node sends to the Group Manager a confirmable CoAP request, using the method POST and targeting the join endpoint associated to that multicast group. The request payload specifies the intended role(s) of the joining node in the multicast group, i.e. multicaster and/or (pure) listener [I-D.tiloca-core-multicast-oscoap].

The Group Manager processes the request according to [I-D.ietf-ace-oauth-authz]. If this yields to a positive response, the Group Manager updates the group membership by registering the joining node as a new member of the group. Then, the Group Manager

replies to the joining node including the following pieces of information in the CoAP response payload:

- An OSCOAP endpoint ID, if the joining node is not configured exclusively as pure listener (see Section 3 of [I-D.tiloca-core-multicast-oscoap]). The Group Manager ensures that each OSCOAP endpoint ID in use is unique within a same multicast group.

- The OSCOAP Security Common Context associated to the joined multicast group (see Section 4 of [I-D.tiloca-core-multicast-oscoap]).

From then on, the joining node is registered as a member of the multicast group, and can exchange group messages secured with OSCOAP as described in Section 5 of [I-D.tiloca-core-multicast-oscoap].

5. Public Keys of Joining Nodes

Source authentication of OSCOAP messages exchanged within the multicast group is ensured by means of digital counter signatures [I-D.tiloca-core-multicast-oscoap]. Therefore, group members must be able to retrieve each other's public key from a trusted key repository, in order to verify the authenticity of incoming group messages. As also discussed in Section 7.4 of [I-D.tiloca-core-multicast-oscoap], the Group Manager can be configured to store public keys of group members and provide them upon request.

Upon joining a multicast group, a joining node is expected to make its own public key available to the other group members, either through the Group Manager or through another trusted, publicly available, key repository. However, this is not required, if at least one of the following conditions hold.

- The joining node joins a group exclusively as pure listener.

- The joining node joins a group where only group authentication of messages is provided (see Appendix C of [I-D.tiloca-core-multicast-oscoap]).

In case the Group Manager is not configured to store public keys of group members, a joining node SHOULD specify to the Group Manager the address of a trusted key repository where its own public key is available. In particular, upon performing a join process with a given Group Manager for the first time, the joining node additionally includes this information in the payload of the POST request

targeting the join endpoint. The Group Manager can then redirect group members to the correct key repository in case of need.

Instead, in case the Group Manager is configured to store public keys of group members, two main cases can occur.

- The joining node and the Group Manager have used an asymmetric proof-of-possession key to establish a secure communication channel. In this case, the Group Manager stores the proof-of-possession key conveyed in the Access Token as the public key of the joining node.

- The joining node and the Group Manager have used a symmetric proof-of-possession key to establish a secure communication channel. In this case, upon performing a join process with that Group Manager for the first time, the joining node includes its own public key in the payload of the POST request targeting the join endpoint. Then, the Group Manager MUST verify that the joining node actually owns the associated private key, for instance by performing a proof-of-possession challenge-response.

Furthermore, if the Group Manager is configured as key repository, it SHOULD provide a joining node with the public keys of the current members in the joined group. In particular, when providing the OSCOAP Endpoint ID and the OSCOAP Security Common Context as described in Section 4, the Group Manager additionally includes the following material in the response to the joining node:

- The public keys of the non-pure listeners currently in the joined multicast group, if the joining node is configured (also) as multicaster.

- The public keys of the multicasters currently in the joined multicast group, if the joining node is configured (also) as non-pure listener.

6. Updating Authorization Information

At any point in time, a node might want to join further OSCOAP multicast groups under the same Group Manager. In such a case, the joining node requests from the AS an updated Access Token for accessing the new multicast groups of interest.

The joining node uploads the new Access Token to the /authz-info resource at the Group Manager, using the already established secure channel. After that, the joining node performs the joining process described in Section 4, separately for each multicast group to join.

Since the joining node and the Group Manager already share a secure communication channel, they are not required to establish a new one. However, according to the specific profile of ACE in use, the joining node and the Group Manager may leverage the new Access Token to establish a new secure communication channel or update the currently existing one. For instance, Section 4.2 of [I-D.ietf-ace-dtls-authorize] describes how the new Access Token can be used to renegotiate an existing DTLS session or to establish a new one by performing a new DTLS handshake.

7. Security Considerations

This document relies on the security considerations included in Section 7 of [I-D.tiloca-core-multicast-oscoap], as to different management aspects related to OSCOAP multicast groups:

- Management of group keying material (Section 7.2). This includes the need to revoke and renew the keying material currently used in the multicast group, upon changes in the group membership. In particular, renewing the keying material is required upon a new node joining the multicast group, in order to preserve backward security. The Group Manager is responsible to enforce rekeying policies and accordingly update the keying material within the multicast groups of its competence.

- Synchronization of sequence numbers (Section 7.3). This concerns how a listener node that has just joined a multicast group can synchronize with the sender sequence number of multicasters in the same group. To this end, the new listener node performs a challenge-response with a multicaster node, leveraging the Repeat Option for CoAP [I-D.amsuess-core-repeat-request-tag].

- Provisioning of public keys (Section 7.4). This provides guidelines about how to ensure the availability of group members' public keys, possibly relying on the Group Manager as trusted key repository. Section 5 of this specification leverages and builds on such considerations.

Further security considerations are (going to be) inherited from the ACE framework for Authentication and Authorization [I-D.ietf-ace-oauth-authz], as well as from the CoAP-DTLS profile [I-D.ietf-ace-dtls-authorize] and the OSCOAP profile [I-D.seitz-ace-oscoap-profile] of ACE.

8. IANA Considerations

This document has no actions for IANA.

9. Acknowledgments

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EST over secure CoAP (EST-coaps)
draft-vanderstok-ace-coap-est-02

Abstract

Low-resource devices in a Low-power and Lossy Network (LLN) can operate in a mesh network using the IPv6 over Low-power Wireless Personal Area Networks (6LoWPAN) and IEEE 802.15.4 link-layer standards. Provisioning these devices in a secure manner with keys (often called secure bootstrapping) used to encrypt and authenticate messages, is the subject of Bootstrapping of Remote Secure Key Infrastructures (BRSKI) [I-D.ietf-anima-bootstrapping-keyinfra] and 6tisch Secure Join [I-D.ietf-6tisch-dtsecurity-secure-join]. Enrollment over Secure Transport (EST) [RFC7030], based on TLS and HTTP, is used in BRSKI. Low-resource devices often use the lightweight Constrained Application Protocol (CoAP) [RFC7252] for message exchanges. This document defines how low-resource devices are expected to use EST over secure CoAP (EST-coaps) for secure bootstrapping and certificate enrollment. 6LoWPAN fragmentation management and extensions to CoAP registries are needed to enable EST-coaps.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

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This Internet-Draft will expire on December 14, 2017.

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1. Introduction

IPv6 over Low-power Wireless Personal Area Networks (6LoWPANs) [RFC4944] on IEEE 802.15.4 [ieee802.15.4] wireless networks is becoming common in many industry application domains such as lighting controls. However, commissioning of such networks suffers from a lack of standardized secure bootstrapping mechanisms for these networks.

Although IEEE 802.15.4 defines how security can be enabled between nodes within a single mesh network, it does not specify the provisioning and management of the keys. Therefore, securing a 6LoWPAN network with devices from multiple manufacturers with different provisioning techniques is often tedious and time consuming.

Bootstrapping of Remote Secure Infrastructures (BRSKI) [I-D.ietf-anima-bootstrapping-keyinfra] addresses the issue of bootstrapping networked devices in the context of Autonomic Networking Integrated Model and Approach (ANIMA). [I-D.ietf-6tisch-minimal-security] and [I-D.ietf-6tisch-dtsecurity-secure-join] also address secure bootstrapping in the 6tisch context targeted to low-resource devices. BRSKI has not been developed specifically for low-resource devices in constrained networks. Constrained networks use DTLS [RFC6347], CoAP [RFC7252], and UDP instead of TLS [RFC5246], HTTP [RFC7230] and TCP.

BRSKI relies on Enrollment over Secure Transport (EST) [RFC7030] for the provisioning of the operational domain certificates. EST-coaps provides a subset of EST functionality and extends EST with BRSKI functions. EST-coaps replaces the invocations of TLS and HTTP by DTLS and CoAP invocations thus enabling EST and BRSKI for CoAP-based low-resource devices.

Although EST-coaps paves the way for the utilization of EST for constrained devices on constrained networks, some devices will not have enough resources to handle the large payloads that come with

EST-coaps. The specification of EST-coaps is intended to ensure that bootstrapping works for less constrained devices that choose to limit their communications stack to UDP/CoAP. It is up to the network designer to decide which devices execute the EST protocol and which not.

EST-coaps is designed for use in professional control networks such as Building Control. The autonomic bootstrapping is interesting because it reduces the manual intervention during the commissioning of the network. Typing in passwords is contrary to this wish. Therefore, the HTTP Basic authentication of EST is not supported in EST-coaps.

In the constrained devices context, it is very unlikely that full PKI request messages will be used. Therefore, full PKI request messages are not supported by EST-coaps.

Because the relatively large EST messages cannot be readily transported over constrained (6LoWPAN, LLN) wireless networks, this document specifies the use of CoAP Block-Wise Transfer ("Block") [RFC7959] to fragment EST messages at the application layer.

Support for Observe CoAP options [RFC7641] with BRSKI is not supported in the current BRSKI/EST message flows and is thus out-of-scope for this discussion. Observe options could be used by the server to notify clients about a change in the cacerts or csr attributes (resources) and might be an area of future work.

1.1. Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

Many of the concepts in this document are taken over from [RFC7030]. Consequently, much text is directly traceable to [RFC7030]. The same document structure is followed to point out the differences and commonalities between EST and EST-coaps.

The following terms are defined in the BRSKI protocol [I-D.ietf-anima-bootstrapping-keyinfra]: pledge, Join proxy (or Circuit Proxy?), Join Registrar, and Manufacturer Authorized Signing Authorities (MASA).

2. EST operational differences

Only the differences to EST with respect to operational scenarios are described in this section. EST-coaps server differs from EST server as follows:

- Replacement of TLS by DTLS and HTTP by CoAP, resulting in:
  - DTLS-secured CoAP sessions between EST-coaps client and EST-coaps server.
- Only certificate-based client authentication is supported, which results in:
  - The EST-coaps client does not support HTTP Basic authentication (as described in Section 3.2.3 of [RFC7030]).
  - The EST-coaps client does not support authentication at the application layer (as described in Section 3.2.3 of [RFC7030]).
- EST-coaps does not support full PKI request messages[RFC5272].
  - Consequently, the fullcmc request of section 4.3 of [RFC7030] and response MUST NOT be supported by EST-coaps].
- EST-coaps specifies the BRSKI extensions over CoAP as specified in sections 3.2, 3.4, 3.5, and 3.8.4 of [I-D.ietf-anima-bootstrapping-keyinfra].

3. Conformance to RFC7925 profiles

This section shows how EST-coaps fits into the profiles of low-resource devices as described in [RFC7925]. Within the bootstrap context a Public Key Infrastructure (PKI) is used, where the client is called "pledge", the Registration Authority (RA) is called Join Registrar, which acts at the front-end for the Certificate Authority (CA) and receives voucher feedback from as many Manufacturer Authorized Signing Authorities (MASA) as there are manufacturers. A Join Proxy (Circuit Proxy?)is placed between client and RA to receive join requests over a 1-hop unsecured channel and transmitted over the secure network to the EST-server. The EST-server of EST-coaps is placed between Join-Proxy (Circuit Proxy) and RA or is part of RA.

EST-coaps can transport certificates and private keys. Private keys can be transported as response to a request to a server-side key generation as described in section 4.4 of [RFC7030]. In the bootstrapping context, EST-coaps transport is limited to the EST certificate transport conformant to section 4.4 of [RFC7925]. For

BRSKI, outside the profiles of [RFC7925], EST-coaps transports vouchers, which are YANG files specified in [I-D.ietf-anima-voucher].

The mandatory cipher suite for DTLS is TLS_ECDHE_ECDSA_WITH_AES_128_CCM_8 defined in [RFC7251] which is the mandatory-to-implement cipher suite in CoAP. Additionally, the curve secp256r1 MUST be supported [RFC4492]; this curve is equivalent to the NIST P-256 curve. The hash algorithm is SHA-256. DTLS implementations MUST use the Supported Elliptic Curves and Supported Point Formats Extensions [RFC4492]; the uncompressed point format MUST be supported; [RFC6090] can be used as an implementation method.

The EST-coaps client MUST be configured with an explicit TA database or at least an implicit TA database from its manufacturer. The authentication of the EST-coaps server by the EST-coaps client is based on Certificate authentication in the DTLS handshake.

The authentication of the EST-coaps client is based on client certificate in the DTLS handshake. This can either be

- DTLS with a previously issued client certificate (e.g., an existing certificate issued by the EST CA); this could be a common case for simple re-enrollment of clients;

- DTLS with a previously installed certificate (e.g., manufacturer-installed certificate or a certificate issued by some other party);

4. Protocol Design and Layering

EST-coaps uses CoAP to transfer EST messages, aided by Block-Wise Transfer [RFC7959] to transport CoAP messages in blocks thus avoiding (excessive) 6LoWPAN fragmentation of UDP datagrams. The use of "Block" for the transfer of larger EST messages is specified in Section 4.5. The Figure 1 below shows the layered EST-coaps architecture.

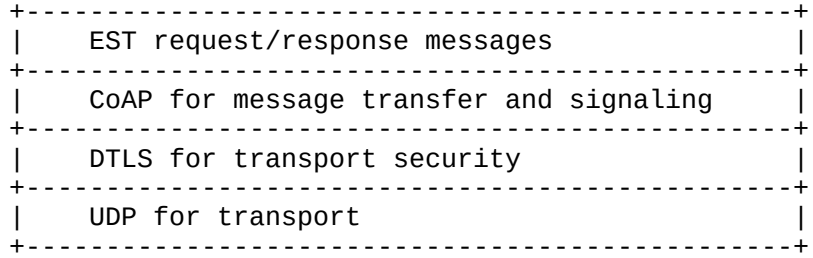


Figure 1: EST-coaps protocol layers

The EST-coaps protocol design follows closely the EST design, excluding some aspects that are not relevant for automatic bootstrapping of constrained devices within a professional context. The parts supported by EST-coaps are identified by their message types:

- Simple enroll and reenroll, for CA to sign public client-identity key.
- CA certificate retrieval, needed to receive the complete set of CA certificates.
- CSR Attributes request messages, informs the pledge of the fields to include in generated CSR.
- Server-side key generation messages, to provide a private client-identity key when the client is too restricted or because of lack of an entropy source. [Encrypting these keys is important. RFC7030 specifies how the private key can be encrypted with CMS using symmetric or asymmetric keys.]

4.1. Discovery and URI

EST-coaps is targeted to low-resource networks with small packets. Saving header space is important and the EST-coaps URI is shorter than the EST URI.

The presence and location of (path to) the management data are discovered by sending a GET request to "/.well-known/core" including a resource type (RT) parameter with the value "ace.est" [RFC6690]. Upon success, the return payload will contain the root resource of the EST resources. It is up to the implementation to choose its root resource; throughout this document the example root resource /est is used. The example below shows the discovery of the presence and location of management data.

```
  REQ: GET /.well-known/core?rt=ace.est

  RES: 2.05 Content
</est>; rt="ace.est"
```

The EST-coaps server URIs differ from the EST URI by replacing the scheme https by coaps and by specifying shorter resource path names:

coaps://www.example.com/est/short-name

Figure 5 in section 3.2.2 of [RFC7030] enumerates the operations and corresponding paths which are supported by EST. Table 1 provides the mapping from the EST and BRSKI URI path to the EST-coaps URI path.

BRSKI	EST	EST-coaps
	/cacerts	/crts
	/simpleenroll	/sen
	/simplereenroll	/sren
	/csrattrs	/att
	/serverkeygen	/skg
/requestvoucher		/rv
/voucher_status		/vs
/enrollstatus		/es

Table 1

/requestvoucher and /enrollstatus are needed between pledge and Registrar.

When discovering the root path for the EST resources, the server MAY return the full resource paths and the used content types. This is useful when multiple content types are specified for EST-coaps server. For example, the following more complete response is possible.

```
  REQ: GET /.well-known/core?rt=ace.est

  RES: 2.05 Content
</est>; rt="ace.est"
</est/crts>; rt="ace.est";ct=TBD1
</est/sen>; rt="ace.est";ct=TBD1 TBD4
</est/sren>; rt="ace.est";ct=TBD1 TBD4
</est/att>; rt="ace.est";ct=TBD4
</est/skg>; rt="ace.est";ct=TBD1 TBD4 TBD2
</est/rv>; rt="ace.est";ct=TBD5 TBD6
</est/vs>; rt="ace.est";ct=50
</est/es>; rt="ace.est";ct=50
```

ct=50 stands for the Content-Format "application/json"

The return of the content-types allows the client to choose the most appropriate one from multiple content types.

4.2. Payload format

The content-format (media type equivalent) of the CoAP message determines which EST message is transported in the CoAP payload. The media types specified in the HTTP Content-Type header (see section 3.2.2 of [RFC7030]) are in EST-coaps specified by the Content-Format Option (12) of CoAP. The combination of URI path-suffix and content-format used for coap MUST map to an allowed combination of path-suffix and media type as defined for EST. The required content-formats for these request and response messages are defined in Section 8. The CoAP response codes are defined in Section 4.4.

EST-coaps is designed for use between low-resource devices using CoAP and hence does not need to send base64-encoded data. Simple binary is more efficient (30% less payload compared to base64) and well supported by CoAP. Therefore, the content formats specification in Section 8 requires the use of binary for all EST-coaps Content-Formats.

4.3. Message Bindings

This section describes BRSKI to CoAP message mappings.

All /crts, /sen, /sren, /att, /skg, /rv, /vs, and /es EST-coaps messages expect a response, so they are all CoAP CON messages.

The Ver, TKL, Token, and Message ID values of the CoAP header are not influenced by EST.

CoAP options are used to convey Uri-Host, Uri-Path, Uri-Port, Content-Format and more in CoAP. The CoAP Options are used to communicate the HTTP fields specified in the BRSKI REST messages.

BRSKI URLs are HTTPS based (https://), in CoAP these will be assumed to be transformed to coaps (coaps://)

Appendix A includes some practical examples of EST messages translated to CoAP.

4.4. CoAP response codes

Section 5.9 of [RFC7252] specifies the mapping of HTTP response codes to CoAP response codes. Every time the HTTP response code 200 is specified in [RFC7030] in response to a GET request, in EST-coaps the equivalent CoAP response code 2.05 MUST be used. Response code HTTP 202 in EST is mapped to CoAP 2.06 as specified in [I-D.hartke-core-pending]. All other HTTP 2xx response codes are not used by EST. For the following HTTP 4xx error codes that may occur:

400, 401, 403, 404, 405, 406, 412, 413, 415; the equivalent CoAP response code for EST-coaps is 4.xx. For the HTTP 5xx error codes: 500, 501, 502, 503, 504 the equivalent CoAP response code is 5.xx.

4.5. Message fragmentation

DTLS defines fragmentation only for the handshake part and not for secure data exchange (DTLS records). [RFC6347] states "Each DTLS record MUST fit within a single datagram". To avoid using IP fragmentation, which is not supported by 6LoWPAN, invokers of the DTLS record layer MUST size DTLS records so that they fit within any Path MTU estimates obtained from the record layer. In addition, invokers residing on a 6LoWPAN over IEEE 802.15.4 network SHOULD attempt to size CoAP messages such that each DTLS record will fit within one or two IEEE 802.15.4 frames.

That is not always possible. Even though ECC certificates are small in size, they can vary greatly based on signature algorithms, key sizes, and OID fields used. For 256-bit curves, common ECDSA cert sizes are 500-1000 bytes which could fluctuate further based on the algorithms, OIDs, SANs and cert fields. For 384-bit curves, ECDSA certs increase in size and can sometimes reach 1.5KB. Additionally, there are times when the EST cacerts response from the server can include multiple certs that amount to large payloads. CoAP [RFC7252]’s section 4.6 describes the possible payload sizes: "if nothing is known about the size of the headers, good upper bounds are 1152 bytes for the message size and 1024 bytes for the payload size". Also "If IPv4 support on unusual networks is a consideration, implementations may want to limit themselves to more conservative IPv4 datagram sizes such as 576 bytes; per [RFC0791], the absolute minimum value of the IP MTU for IPv4 is as low as 68 bytes, which would leave only 40 bytes minus security overhead for a UDP payload". Thus, even with ECC certs, EST-coaps messages can still exceed sizes in MTU of 1280 for IPv6 or 60-80 bytes for 6LoWPAN [RFC4919] as explained in section 2 of [RFC7959]. EST-coaps needs to be able to fragment EST messages into multiple DTLS datagrams. Fine-grained fragmentation of EST messages is essential.

To perform fragmentation in CoAP, [RFC7959] specifies the "Block1" option for fragmentation of the request payload and the "Block2" option for fragmentation of the return payload of a CoAP flow.

The BLOCK draft defines SZX in the Block1 and Block2 option fields. These are used to convey the size of the blocks in the requests or responses.

The CoAP client MAY specify the Block1 size and MAY also specify the Block2 size. The CoAP server MAY specify the Block2 size, but not

the Block1 size. As explained in Section 1 of [RFC7959]), blockwise transfers SHOULD be used in Confirmable CoAP messages to avoid the exacerbation of lost blocks.

The Size1 response MAY be parsed by the client as a size indication of the Block2 resource in the server response or by the server as a request for a size estimate by the client. Similarly, Size2 option defined in BLOCK should be parsed by the server as an indication of the size of the resource carried in Block1 options and by the client as a maximum size expected in the 4.13 (Request Entity Too Large) response to a request.

Examples of fragmented messages are shown in Appendix B.

5. Transport Protocol

EST-coaps depends on a secure transport mechanism over UDP that can secure (confidentiality, authenticity) the CoAP messages exchanged.

5.1. DTLS

DTLS is one such secure protocol. Within BRSKI and EST when "TLS" is referred to, it is understood that in EST-coaps, security is provided using DTLS instead. No other changes are necessary (all provisional modes etc. are the same as for TLS).

CoAP was designed to avoid fragmentation. DTLS is used to secure CoAP messages. However, fragmentation is still possible at the DTLS layer during the DTLS handshake when using ECC ciphersuites. If fragmentation is necessary, "DTLS provides a mechanism for fragmenting a handshake message over a number of records, each of which can be transmitted separately, thus avoiding IP fragmentation" [RFC6347].

CoAP and DTLS can provide proof of identity for EST-coaps clients and server with simple PKI messages conformant to section 3.1 of [RFC5272]. EST-coaps supports the certificate types and Trust Anchors (TA) that are specified for EST in section 3 of [RFC7030].

Channel-binding information for linking proof-of-identity with connection-based proof-of-possession is optional for EST-coaps. When proof-of-possession is desired, a set of actions are required regarding the use of tls-unique, described in section 3.5 in [RFC7030]. The tls-unique information translates to the contents of the first "Finished" message in the TLS handshake between server and client [RFC5929]. The client is then supposed to add this "Finished" message as a ChallengePassword in the attributes section of the PKCS#10 Request Info to prove that the client is indeed in control of

the private key at the time of the TLS session when performing a /simpleenroll, for example. In the case of EST-coaps, the same operations can be performed during the DTLS handshake. In the event of handshake message fragmentation, the Hash of the handshake messages used in the MAC calculation of the Finished message

PRF(master_secret, finished_label, Hash(handshake_messages))
[0..verify_data_length-1];

MUST be computed as if each handshake message had been sent as a single fragment [RFC6347].

In a constrained CoAP environment, endpoints can’t afford to establish a DTLS connection for every EST transaction. Authenticating and negotiating DTLS keys requires resources on low-end endpoints and consumes valuable bandwidth. The DTLS connection SHOULD remain open for persistent EST connections. For example, an EST cacerts request that is followed by a simpleenroll request can use the same authenticated DTLS connection. Given that after a successful enrollment, it is more likely that a new EST transaction will take place after a significant amount of time, the DTLS connections SHOULD only be kept alive for EST messages that are relatively close to each other.

5.2. 6tisch approach

The 6tisch bootstrapping is targeted to the "imprinting" of the "pledge" with layer 2 keys. The content formats for the transport are being defined and may be expressed in a YANG module.

Instead of using transport security, the 6tisch approach relies on application security provided by OSCOAP [I-D.ietf-core-object-security] and EDHOC [I-D.selander-ace-cose-ecdhe]. [I-D.selander-ace-eals] uses OSCOAP to securely enroll certificates by using Certificate Management over CMS (CMC) (EST is profile of CMC).

It is suggested that the EST-coaps communication between pledge and registrar, specified in this document, can be freely exchanged with the same communication specified in [I-D.ietf-6tisch-dtsecurity-secure-join] and [I-D.ietf-6tisch-minimal-security].

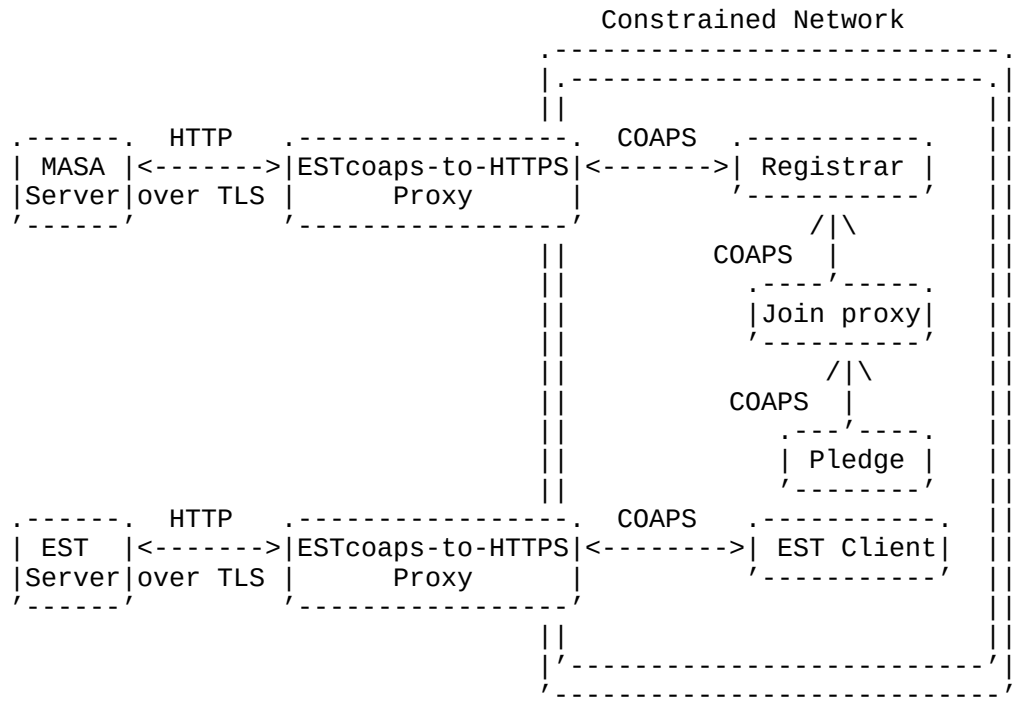
[EDNOTE: The evolution of this section depends on the directions taken by 6tisch and anima and the possible commonality that will be provided.]

6. Proxying

In real-world deployments, entities like the EST server, CA or MASA will not always reside within the COAP boundary. The MASA or a CA can exist outside the constrained network in a non-constrained network that supports TLS/HTTP. In such environments EST-coaps is used by the pledge within the COAP boundary and TLS is used to transport the EST/BRSKI messages outside the CoAP boundary. A proxy entity at the edge is required to operate between the COAP environment and the external HTTP network. The ESTcoaps-to-HTTPS proxy SHOULD terminate EST-coaps downstream and initiate EST/BRSKI connections over TLS upstream.

Two separate use-cases, shown in one figure below, are expected to be deployed in practice:

- A proxy between any EST-client and EST-server indepedent of BRSKI
- A proxy between Registrar and MASA



ESTcoaps-to-HTTPS proxy at the COAP boundary.

Table 1 contains the mapping between the EST-coaps and EST/BRSKI URIs the proxy SHOULD adhere to. Section 7 of [RFC8075] and Section 4.4 define the mapping between EST-coaps and HTTP response codes, that determines how a proxy translates COAP response codes from/to HTTP status codes. The mapping from Content-Type to media type is defined in Section 8. The conversion from binary to BSD64 needs to be done in the proxy. Conversion is possible because a TLS link exists between EST-coaps-to-HTTP proxy and HTTP MASA or EST server and a corresponding DTLS linked exists between EST-coaps-to-HTTP proxy and EST client or Registrar.

Due to fragmentation of large messages into blocks, an EST-coaps-to-HTTP proxy SHOULD reassemble the BLOCKs before translating the binary content to BSD64, and consecutively relay the message upstream into the HTTP environment.

For the discovery of the EST server by the EST client in the coap environment, the EST-coaps-to-HTTP proxy MUST announce itself according to the rules of Section 4.1. The available functions of the proxies MUST be announced with as many resource paths. The discovery of MASA and EST server in the http environment follow the rules specified in [I-D.ietf-anima-bootstrapping-keyinfra].

[EDNOTE: PoP will be addressed here.]

A proxy SHOULD authenticate the client downstream and it should be authenticated by the EST or BRSKI server or CA upstream. A trust relationship needs to be pre-established between the proxy and the TCP entities (EST, BRSKI servers) to be able to proxy these connections on behalf of various clients.

[EDNOTE: To add more details about trust relations in this section.]

7. Parameters

[EDNOTE: This section to be populated. It will address transmission parameters for BRSKI described in sections 4.7 and 4.8 of the CoAP draft. BRSKI does not impose any unique parameters that affect the CoAP parameters in Table 2 and 3 in the CoAP draft but the ones in CoAP could be affecting BRSKI. For example, the processing delay of CAs could be less then 2s, but in this case they should send a CoAP ACK every 2s while processing.]

8. IANA Considerations

8.1. Content-Format registry

Additions to the sub-registry "CoAP Content-Formats", within the "CoRE Parameters" registry are needed for the below media types. These can be registered either in the Expert Review range (0-255) or IETF Review range (256-9999).

1.

* application/pkcs7-mime
* Type name: application
* Subtype name: pkcs7-mime
* smime-type: certs-only
* ID: TBD1
* Required parameters: None
* Optional parameters: None
* Encoding considerations: binary
* Security considerations: As defined in this specification
* Published specification: [RFC5751]
* Applications that use this media type: ANIMA Bootstrap (BRSKI) and EST

2.

* application/pkcs8
* Type name: application
* Subtype name: pkcs8
* ID: TBD2
* Required parameters: None
* Optional parameters: None
* Encoding considerations: binary

   * Security considerations: As defined in this specification

   * Published specification: [RFC5958]

   * Applications that use this media type: ANIMA Bootstrap (BRSKI) and EST

3.

   * application/csrattrs

   * Type name: application

   * Subtype name: csrattrs

   * ID: TBD3

   * Required parameters: None

   * Optional parameters: None

   * Encoding considerations: binary

   * Security considerations: As defined in this specification

   * Published specification: [RFC7030]

   * Applications that use this media type: ANIMA Bootstrap (BRSKI) and EST

4.

   * application/pkcs10

   * Type name: application

   * Subtype name: pkcs10

   * ID: TBD4

   * Required parameters: None

   * Optional parameters: None

   * Encoding considerations: binary

   * Security considerations: As defined in this specification

* Published specification: [RFC5967]

* Applications that use this media type: ANIMA bootstrap (BRSKI) and EST

*

  + application/voucherrequest

  + Type name: application

  + Subtype name: voucherrequest

  + ID: TBD5

  + Required parameters: None

  + Optional parameters: None

  + Encoding considerations: binary

  + Security considerations: As defined in this specification

  + Published specification: BRSKI??

  + Applications that use this media type: ANIMA bootstrap (BRSKI)

*

  + application/voucher+cms

  + Type name: application

  + Subtype name: voucher+cms

  + ID: TBD6

  + Required parameters: None

  + Optional parameters: None

  + Encoding considerations: binary

  + Security considerations: As defined in this specification

  + Published specification: BRSKI??

+ Applications that use this media type: ANIMA bootstrap (BRSKI)

8.2. Resource Type registry

Additions to the sub-registry "CoAP Resource Type", within the "CoRE Parameters" registry are needed for a new resource type.

o rt="ace.est" needs registration with IANA.

9. Security Considerations

9.1. proxy considerations

In the BRSKI bootstrap protocol, there is a direct TLS connection from pledge to EST-server. With the EST-coaps specification a direct DTLS connection from pledge to EST-server is possible thus avoiding the placement of a https/coaps proxy between pledge and http EST-server. Such a https/coaps proxy presents a security issue because the proxy needs to make a TLS connection with the EST-server and a DTLS connection with the pledge.

In [I-D.ietf-anima-bootstrapping-keyinfra] the EST-server and Registrar are co-located on the same host, thus avoiding security connections between Registrar and EST-server. It is RECOMMENDED that the links "Registrar/MASA" and "Registrar/CA" use a http TLS connection, identical to BRSKI protocol. The consequence is that the Registrar host provides both a coaps and a https stack.

The proxies proposed in Section 6 must be deployed with great care, and only when the recommended connections are impossible.

9.2. EST server considerations

The security considerations of section 6 of [RFC7030] are only partially valid for the purposes of this document. As HTTP Basic Authentication is not supported, the considerations expressed for using passwords do not apply.

Given that the client has only limited resources and may not be able to generate sufficiently random keys to encrypt its identity, it is possible that the client uses server generated private/public keys to encrypt its certificate. The transport of these keys is inherently risky. A full probability analysis MUST be done to establish whether server side key generation enhances or decreases the probability of identity stealing.

When a client uses the Implicit TA database for certificate validation, the client cannot verify that the implicit data base can act as an RA. It is RECOMMENDED that such clients include "Linking Identity and POP Information" Section 5.1 in requests (to prevent such requests from being forwarded to a real EST server by a man in the middle). It is RECOMMENDED that the Implicit Trust Anchor database used for EST server authentication be carefully managed to reduce the chance of a third-party CA with poor certification practices from being trusted. Disabling the Implicit Trust Anchor database after successfully receiving the Distribution of CA certificates response (Section 4.1.3 of [RFC7030]) limits any vulnerability to the first DTLS exchange.

In accordance with [RFC7030], TLS cipher suites that include "_EXPORT_" and "_DES_" in their names MUST NOT be used. More information about recommendations of TLS and DTLS are included in [RFC7525].

As described in CMC, Section 6.7 of [RFC5272], "For keys that can be used as signature keys, signing the certification request with the private key serves as a POP on that key pair". The inclusion of tls-unique in the certification request links the proof-of-possession to the TLS proof-of-identity. This implies but does not prove that the authenticated client currently has access to the private key.

Regarding the CSR attributes that the CA may list for inclusion in an enrollment request, an adversary could exclude attributes that a server may want, include attributes that a server may not want, and render meaningless other attributes that a server may want. The CA is expected to be able to enforce policies to recover from improper CSR requests.

Interpreters of ASN.1 structures should be aware of the use of invalid ASN.1 length fields and should take appropriate measures to guard against buffer overflows, stack overruns in particular, and malicious content in general.

10. Acknowledgements

The authors are very grateful to Klaus Hartke for his detailed explanations on the use of Block with DTLS. The authors would like to thank Esko Dijk and Michael Verschoor for the valuable discussions that helped in shaping the solution. They would also like to thank Peter Panburana from Cisco for his feedback on technical details of the solution. Constructive comments were received from Eliot Lear and Julien Vermillard.

11. Change Log

-02:

binary instead of CBOR binary in mime types.

supported content types are discoverable.

DTLS POP text improved.

First version of Security considerations section written.

First version of Proxying section written.

Various text improvements.

-01:

Merging of draft-vanderstok-ace-coap-est-00 and draft-pritikin-coap-bootstrap-01

URI and discovery are modified

More text about 6tisch bootstrap including EDHOC and OSCOAP

mapping to DICE IoT profiles

adapted to BRSKI progress

12. References

12.1. Normative References

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Appendix A. EST messages to EST-coaps

This section takes all examples from Appendix A of [RFC7030], changes the payload from Base64 to binary and replaces the http headers by their CoAP equivalents.

A.1. cacerts

In EST-coaps, a coaps cacerts message can be:

GET coaps://[192.0.2.1:8085]/est/crts

The corresponding CoAP header fields are shown below. The use of block and DTLS are worked out in Appendix B.

```
  Ver = 1
  T = 0 (CON)
  Code = 0x01 (0.01 is GET)
  Options
   Option1 (Uri-Host)
     Option Delta = 0x3  (option nr = 3)
     Option Length = 0x9
     Option Value = 192.0.2.1
   Option2 (Uri-Port)
     Option Delta = 0x4  (option nr = 4+3=7)
     Option Length = 0x4
     Option Value = 8085
   Option3 (Uri-Path)
     Option Delta = 0x4   (option nr = 7+4= 11)
     Option Length = 0x9
     Option Value = /est/crts
  Payload = [Empty]
```

A 2.05 Content response with a cert in EST-coaps will then be:

2.05 Content (Content-Format: application/pkcs7-mime)
{payload}

with CoAP fields

```
  Ver = 1
  T = 2 (ACK)
  Code = 0x45 (2.05 Content)
  Options
    Option1 (Content-Format)
      Option Delta = 0xC  (option nr = 12)
      Option Length = 0x2
      Option Value = TBD1 (defined in this document)
```

Payload =

```
30233906092a6206734107028c2a3023260201013100300b06092a6206734107018
c0c3020bb302063c20102020900a61e75193b7acc0d06092a620673410105050030
1b3119301706035504031310657374578616d706c654341204f774f301e170d313
330353039303353333315a170d313430353039303353333315a301b3119301706
035504031310657374578616d706c654341204f774f302062300d06092a6206734
10101050003204f0030204a022041003a923a2968bae4aae136ca4e2512c5200680
358482ac39d6f640e4574e654ea35f48b1e054c5da3372872f7a1e429f4edf39584
32efb2106591d3eb783c1034709f251fc86566bda2d541c792389eac4ec9e181f4b
9f596e5ef2679cc321542b11337f90a44df3c85f1516561fa968a1914f265bc0b82
76ebe3106a790d97d34c8c37c74fe1c30b396424664ac426284a9f6022e02693843
6880adfcd95c98ca1dfc2e6d75319b85d0458de28a9d13fb16d620fff7541f6a25d
7daf004355020301000130b040300f0603551d130101f10530030101fc1d0603551
d0e04160414084d321ca0135e77217a486b686b334b00e0603551d0f0101f104030
20106300d06092a620673410105050003204100237 03b965746a0c2c978666d787a
94f89b495a11f0d369b28936ec2475c0f0855c8e83f823f2b871a1d92282f323c45
904ba008579216cf5223b8b1bc425a0677262047f7700240631c17f3035d1c3780b
2385241cba1f4a6e98e6be6820306b3a786de5a557795d1893822347b5f825d34a7
ad2876f8feba4d525b31066f6505796f71530003431a3e6bbfe788b4565029a7e20
a51107677552586152d051e8eebf383e92288983421d5c5652a4870c3af74b9bdbe
d6b462e2263d30f6d3020c330206bc20102020101300d06092a6206734101050500
301b3119301706035504031310657374578616d706c654341204f774f301e170d3
133303530393033353333325a170d313430353039303353333325a301b31193017
060355040313106573744578616d706c654341204e774f302062300d06092a62067
341010105000320410023703b965746a0c2c978666d787a
1f49e0421120e6b8384160f2bf02630ef544d5fd0d5623b35713c79a7229283a790
8751a634aa420a3e2a4b1f10519d046f02f5a5dd6d760c2a842356e067b7bd94338
d1faa3b3ddd4813060a207b0a097067007e45b052b60fdbae4656e11562c4f5abb7
b0cf87a79d221f1127313c53371ce1245d63db45a1203a23340ba08042c768d03b8
076a028d3a51d87d2ef107bbd6f2305ce5e67668724002fb726df9c14476c37de0f
55033f192a5ad21f9a2a71c20301000134b050300e0603551d0f0101f104030204c
1d0603551d0e04160414112966e304761732fbfe6a2c823c301f0603551d2304183
0165084d321ca0135e77217a486b686b334b00d06092a620673410105050003204 1
00b382ba3355a50e287bae15758b3beff63d34d3e357b90031495d018868e49589b
9faf46a4ad49b1d35b06ef380106677440934663c2cc111c183655f4dc41c0b3401
123d35387389db91f1e1b4131b16c291d35730b3f9b33c7475124851555fe5fc647
e8fd029605367c7e01281bf6617110021b0d10847dce0e9f0ca6c764b6334784055
172c3983d1e3a3a82301a54fcc9b0670c543a1c747164619101ff23b240b2a26394
c1f7d38d0e2f4747928ece5c34627a075a8b3122011e9d9158055c28f020c330206
bc20102020102300d06092a620673410105050030 1b3119301706035504031310 65
73744578616d706c654341204e774e301e170d313330353039303353333325a170
d313430353039303353333325a301b3119301706035504031310657374578616d
706c654341204f774e302062300d06092a620673410101050003204f0030204a022
041003a923a2968bae4aae136ca4e2512c5200680358482ac39d6f640e4574e654e
a35f48b1e054c5da3372872f7a1e429f4edf3958432efb2106591d3eb783c103470
9f251fc86566bda2d541c792389eac4ec9e181f4b9f596e5ef2679cc321542b1133
7f90a44df3c85f1516561fa968a1914f265bc0b8276ebe3106a790d97d34c8c37c7
```

```
4fe1c30b396424664ac426284a9f6022e026938436880adfcd95c98ca1dfc2e6d75
319b85d0458de28a9d13fb16d620fff7541f6a25d7daf004355020301000134b050
300e0603551d0f0101f104030204c1d0603551d0e04160414084d321ca0135e7721
7a486b686b334b01f0603551d2304183016531129 66e304761732fbfe6a2c823c30
0d06092a620673410105050003204 1002e106933a443070acf5594a3a584d08af7e
06c295059370a06639eff9bd418d13bc25a298223164a6cf1856b11a81617282e4a
410d82ef086839c6e2356903227630654553 51e4c596acc7c016b225dec094706c2
a10608f403b10821984c7c152343b18a768c2ad30238dc45dd653ee6092b0d5cd4c
2f7d236043269357f76d13f95fb5f00d0e19263c6833948e1ba612ce8197af650e2
5d882c12f4b6b9b67252c608ef064aca3f9bc867d71172349d510bb7651cd438837
73d927deb41c4673020bb302063c201020209009b9dda3324700d06092a62067341
01050500301b311930170603550403131065737445786 16d706c654341204e774e3
01e170d313330353039303353333325a170d313430353039303353333325a301b
311930170603550403131065737445786 16d706c654341204e774e302062300d060
92a620673410101050003204f0030204a02204100ef6b677a3247c1fc03d2b9baf1
13e5e7e11f49e0421120e6b8384160f2bf02630ef544d5fd0d5623b35713c79a722
9283a7908751a634aa420a3e2a4b1f10519d046f02f5a5dd6d760c2a842356e067b
7bd94338d1faa3b3ddd4813060a207b0a097067007e45b052b60fdbae4656e11562
c4f5abb7b0cf87a79d221f1127313c53371ce1245d63db45a1203a23340ba08042c
768d03b8076a028d3a51d87d2ef107bbd6f2305ce5e67668724002fb726df9c1447
6c37de0f55033f192a5ad21f9a2a71c20301000130b040300f0603551d130101f10
530030101fc1d0603551d0e04160414112966e304761732fbfe6a2c823c300e0603
551d0f0101f10403020106300d06092a620673410105050003204100423f06d4b76
0f4b42744a279035571696f272a0060f1325a40898509601ad14004f652db6312a1
475c4d7cd50f4b269035585d7856c5337765a66b38462d5bdaa7778aab24bbe2815
e37722cd10e7166c50e75ab75a1271324460211991e7445a2960f47351a1a629253
34119794b90e320bc730d6c1bee496e7ac125ce9a1eca595a3a4c54a865e6b623c9
247bfd0a7c19b56077392555c955e233642bec643ae37c166c5e221d797aea3748f
0391c8d692a5cf9bb71f6d0e37984d6fa673a30d0c006343116f58403100
```

A.2. csrattrs

In the following valid /csrattrs exchange, the EST-coaps client authenticates itself with a certificate issued by the connected CA.

The initial DTLS handhake is identical to the enrollment example. The CoAP GET request looks like:

GET coaps://[192.0.2.1:8085]/est/att

with CoAP header fields

```
Ver = 1
T = 0 (CON)
Code = 0x01 (0.01 is GET)
Options
 Option1 (Uri-Host)
   Option Delta = 0x3  (option nr = 3)
   Option Length = 0x9
   Option Value = 192.0.2.1
 Option2 (Uri-Port)
   Option Delta = 0x4  (option nr = 4+3=7)
   Option Length = 0x4
   Option Value = 8085
 Option3 (Uri-Path)
   Option Delta = 0x4   (option nr = 7+4= 11)
   Option Length = 0x8
   Option Value = /est/att
Payload = [Empty]
```

A 2.05 Content response contains attributes which are relevant for the authenticated client. In this example, the EST-coaps server two attributes that the client can ignore when they are unknown to him.:

```
2.05 Content (Content-Format: application/crsattrs)
   {payload}
```

with CoAP fields

```
Ver = 1
T = 2 (ACK)
Code = 0x45 (2.05 Content)
Options
  Option1 (Content-Format)
    Option Delta = 0xC  (option nr = 12)
    Option Length = 0x2
    Option Value = TBD3 (defined in this document)

Payload =
307c06072b060101010116302206038dc1311b131950617273652053455420617
320322e3939392e31206461746106092a620673410907302c06038dc231250603
8dc306038dc413195061727365205345542061732032 2e3939392e32206461746
106092b240303020801010b06096062016503040202
```

A.3. enroll / reenroll

[EDNOTE: We might need a new Option for the Retry-After response message. We might need a new Option for the WWW-Authenticate response.]

During the Enroll/Reenroll exchange, the EST-coaps client uses a CSR (PKCS#10) request in the POST request payload.

POST coaps://[192.0.2.1:8085]/est/sen
(Content-Format: application/pkcs10)

with CoAP header fields

```
     Ver = 1
     T = 0 (CON)
     Code = 0x02 (0.02 is POST)
     Options
      Option1 (Uri-Host)
        Option Delta = 0x3  (option nr = 3)
        Option Length = 0x9
        Option Value = 192.0.2.1
      Option2 (Uri-Port)
        Option Delta = 0x4  (option nr = 4+3=7)
        Option Length = 0x4
        Option Value = 8085
      Option3 (Uri-Path)
        Option Delta = 0x4   (option nr = 7+4= 11)
        Option Length = 0x8
        Option Value = /est/sen
      Option4 (Content-Format)
         Option Delta = 0x1  (option nr = 11+1 = 12)
         Option Length = 0x2
         Option Value = TBD5 (defined in this document)

     Payload =
   [EDNOTE: If POP is used, make sure tls-unique in the CSR is a
    valid HMAC output. ]
   30208530206d020100301f311d301b060355040313146465 6d6f7374657034203
   133363831343133353230206230 0d06092a620673410101050003204f0030204a
   022041005d9f4dffd3c5949f646a9584367778560950b355c35b8e34726dd3764
   54231734795b4c09b9c6d75d408311307a81f7adef7f5d241f7d5be85620c5d44
   38bbb4242cf215c167f2ccf36c364ea2618a62f0536576369d6304e6a96877224
   7d86824f079faac7a6f694cfda5b84c42087dc062d462190c525813f210a036a7
   37b4f30d8891f4b75559fb72752453146332d51c937557716ccec624f5125c3a4
   447ad3115020048113fef54ad554ee88af09a2583aac9024075113db4990b1786
   b871691e0f02030100018701f06092a620673410907311213102b72724369722f
   372b4559753530543430 0d06092a620673410105050003204100441b40177a3a6
   5501487735a8ad5d3827a4eaa867013920e2afcda87aa81733c7c0353be47e1bf
   a7cda5176e7ccc6be22ae03498588d5f2de3b143f2b1a6175ec544e8e7625af6b
   836fd4416894c2e55ea99c6606f69075d6d53475d410729aa6d806afbb9986caf
   7b844b5b3e4545f19071865ada007060cad6db26a592d4a7bda7d586b68110962
   17071103407553155cddc75481e272b5ed553a8593fb7e25100a6f7605085dab4
   fc7e0731f0e7fe305703791362d5157e92e6b5c2e3edbcadb40
```

After verification of the certificate by the server, a 2.05 Content response with the issued certificate will be:

```
2.05 Content (Content-Format: application/pkcs7-mime)
   {payload}
```

with CoAP fields

```
  Ver = 1
  T = 2 (ACK)
  Code = 0x45 (2.05 Content)
  Options
    Option1 (Content-Format)
      Option Delta = 0xC  (option nr = 12)
      Option Length = 0x2
      Option Value = TBD1 (defined in this document)

  Payload =
3020f806092a62067341070283293020e50201013100300b06092a62067341070
1830b3020c730206fc20102020115300d06092a620673410105050030 1b311930
17060355040313106573744578616d706c654341204e774e301e170d313330353
0393233313535335a170d313430353039323331353533 5a301f311d301b060355
0403131464656d6f737465703420313336383134313335323020623 00d06092a6
2067341010105000320 4f0030204a022041005d9f4dffd3c5949f646a95843677
78560950b355c35b8e34726dd376454231734795b4c09b9c6d75d408311307a81
f7adef7f5d241f7d5be85620c5d4438bbb4242cf215c167f2ccf36c364ea2618a
62f0536576369d6304e6a968772247d86824f079faac7a6f694cfda5b84c42087
dc062d462190c525813f210a036a737b4f30d8891f4b75559fb72752453146332
d51c937557716ccec624f5125c3a4447ad3115020048113fef54ad554ee88af09
a2583aac9024075113db4990b1786b871691e0f020301000134b050300e060355
1d0f0101f104030204c1d0603551d0e04160414e81d0788aa2710304c5ecd4d1e
065701f0603551d230418301653112966e304761732fbfe6a2c823c300d06092a
62067341010505000320 41002910d86f2ffeeb914c046816871de601567d291b4
3fabee0f0e8ff81cea27302a7133e20e9d04029866a8963c7d14e26fbe8a0ab1b
77fbb1214bbcdc906fbc381137ec1de685f79406c3e416b8d82f97174bc691637
5a4e1c4bf744c7572b4b2c6bade9fb35da786392ee0d95e3970542565f3886ad6
7746d1b12484bb02616e63302dc371dc6006e431fb7c457598dd204b367b0b3d3
258760a303f1102db26327f929b7c5a60173e1799491b69150248756026b80553
171e4733ad3d13c0103100
```

A.4. serverkeygen

During this valid /serverkeygen exchange, the EST-coaps client authenticates itself using the certificate provided by the connected CA.

The initial DTLS handhake is identical to the enrollment example. The CoAP GET request looks like:

POST coaps://[192.0.2.1:8085]/est/skg

with CoAP header fields

```
     Ver = 1
     T = 0 (CON)
     Code = 0x02 (0.02 is POST)
     Options
      Option1 (Uri-Host)
        Option Delta = 0x3  (option nr = 3)
        Option Length = 0x9
        Option Value = 192.0.2.1
      Option2 (Uri-Port)
        Option Delta = 0x4  (option nr = 4+3=7)
        Option Length = 0x4
        Option Value = 8085
      Option3 (Uri-Path)
        Option Delta = 0x4   (option nr = 7+4= 11)
        Option Length = 0x8
        Option Value = /est/skg
     Option4 (Content-Format)
   [EDNOTE: the client incudes a CSR with a public key  that the
    server should ignore, so we need a content-format here. ]
         Option Delta = 0x1  (option nr = 12)
         Option Length = 0x2
         Option Value = TBD5 (defined in this document)
     Payload =
   [EDNote: If POP is used, make sure tls-unique in the CSR
    is a valid HMAC output. ]
   3020813020690201003 05b313e303c06035504031335736572766572 4b6579476
   56e2072657120627920636c69656e7420696e2064656d6f207374657020313220
   3133363831343139353531193017060355040513105049443a576964676574205
   34e3a3130302062300d06092a620673410101050003204f0030204a02204100f4
   dfa6c03f7f2766b23776c333d2c0f9d1a7a6ee36d01499bbe6f075d1e38a57e98
   ecc197f51b75228454b7f19652332de5e52e4a974c6ae34e1df80b33f15f47d3b
   cbf76116bb0e4d3e04a9651218a476a13fc186c2a255e4065ff7c271cff104e47
   31fad53c22b21a1e5138bf9ad0187314ac39445949a48805392390e78c7659621
   6d3e61327a534f5ea7721d2b1343c7362b37da502717cfc2475653c7a3860c5f4
   0612a5db6d33794d755264b6327e3a3263b149628585b85e57e42f6b3277591b0
   2030100018701f06092a620673410907311213106446734 1586d4a6e6a6f6b427
   4447672300d06092a6206734101050500032041004 72d11007e5a2b2c2023d47a
   6d71d046c307701d8ebc9e47272713378390b4ee321462a3dbe54579f5a514f6f
   4050af497f428189b63655d03a194ef729f101743e5d03fbc6ae1e84486d1300a
   f9288724381909188c851fa9a5059802eb64449f2a3c9e441353d136768da27ff
   4f277651d676a6a7e51931b08f56135a2230891fd184960e1313e7a1a9139ed19
   28196867079a456cd2266cb754a45151b7b1b939e381be333fea61580fe5d25bf
   4823dbd2d6a98445b46305c10637e202856611
```

Without the DecryptKeyIdentifier atttribute, the response has no additional encryption beyond the DTLS one. The EST-coaps server response is:

```
   2.05 Content (Content-Format: application/pkcs8)
      {payload}
```

The response contains first a preamble that can be ignored. The EST-coaps server can use the preamble to include additional explanations, like ownership or support information

```
     Ver = 1
     T = 2 (ACK)
     Code = 0x45 (2.05 Content)
     Options
       Option1 (Content-Format)
         Option Delta = 0xC  (option nr = 12)
         Option Length = 0x2
         Option Value = TBD2 (defined in this document)

     Payload =
   30213e020100300d06092a620673410101050004212830212402010002204100 3
   c0bc2748f2003e3e8ea15f746f2a71e83f585412b92cf6f8e64de02e056153274
   dd01c95dd9cff3112aa141774ab655c3d56359c3b3df055294692ed848e7e30a1
   1bf14e47e0693d93017022b4cdb3e6d40325356152b213c8b535851e681a7074c
   0c6d2b60e7c32fc0336b28e743eba4e5921074d47195d3c05e43c527526e692d5
   45e562578d2d4b5f2191bff89d3eef0222764a2674637a1f99257216647df6704
   efec5adbf54dab24231844eb595875795000e673dd6862310a146ad7e31083010
   001022041004e6b3f78b7791d6377f33117c17844531c81111fb8000282816264
   915565bc7c3f3f643b537a2c69140a31c22550fa97e5132c61b74166b68626704
   260620333050f510096b6570f5880e7e1c15dc0ca6ce2b5f187e2325da14ab705
   ad004717f3b2f779127b5c535e0cee6a343b502722f2397a26126e0af606b5aa7
   f96313511c0b7eb26354f91b82269de62757e3def807a6afdf83ddcbb0614bb7c
   542e6975d6456554e7bd9988fbd1930cd44d0e01ee9182ca54539418653150254
   1ad1a2a11e5021040bfce554b642c29131e7d65455e83c5406d76771912f758f5
   ee3ee36af386f38ffa313c0f661880c5a2b0970485d36f528e7f77a2e55b4ad76
   1242d1c2f75939c8061217d31491d305d3e07d6161c43e26f7de4477b1811de92
   33dc75b426302104015bf48ac376f52887813461fc54635517bcb67293837053e
   8ce1a33da7a35565a75a370dc14555b5316cb55742380350774d769d151ff0456
   0214389a232a2258326163167504cfce44cd316f63bb8a52da53a4cb74fd87194
   c0844881f791f23b0813ea0921325edd14459d41c8a1593f04316388e40b35fef
   7d2a195a5930fa54774427ac821eee2c62790d2c17bd192af794c611011506557
   83d4efe22185cbd83368786f2b1e68a5a27067e321066f0217b4b6d7971a3c21a
   241366b7907187583b511102103369047e5cce0b65012200df5ec697b5827575c
   db6821ff299d6a69574b31ddf0fbe9245ea2f74396c24b3a7565067e41366423b
   5bdd2b2a78194094dbe333f493d159b8e07722f2280d48388db7f1c9f0633bb0e
   173de2c3aa1f200af535411c7090210401421e2ea217e37312dcc606f453a6634
   f3df4dc31a9e910614406412e70eec9247f10672a500947a64356c015a845a7d1
   50e2e3911a2b3b61070a73247166da10bb45474cc97d1ec2bc392524307f35118
   f917438f607f18181684376e13a39e07
   --estServerExampleBoundary
   Ver = 1
```

```
     T = 2 (ACK)
     Code = 0x45 (2.05 Content)
     Options
       Option1 (Content-Format)
         Option Delta = 0xC  (option nr = 12)
         Option Length = 0x2
         Option Value = TBD1 (defined in this document)
     Payload=
   3020c506092a620673410702833630 20f20201013100300b06092a62067341070
   183183020d430207cc20102020116300d06092a620673410105050030 1b311930
   17060355040313106573744578616d706c654341204e774e301e170d313330353
   0393233323535365a170d3134303530393233323535365a302c312a3028060355
   040313217365727665727369646520 6b65792067656e6572617465642072657 37
   06f6e7365302062300d06092a620673410101050003204f0030204a022041003c
   0bc2748f2003e3e8ea15f746f2a71e83f585412b92cf6f8e64de02e056153274d
   d01c95dd9cff3112aa141774ab655c3d56359c3b3df055294692ed848e7e30a11
   bf14e47e0693d93017022b4cdb3e6d40325356152b213c8b535851e681a7074c0
   c6d2b60e7c32fc0336b28e743eba4e5921074d47195d3c05e43c527526e692d54
   5e562578d2d4b5f2191bff89d3eef0222764a2674637a1f99257216647df6704e
   fec5adbf54dab24231844eb595875795000e673dd6862310a146ad7e310830100
   0134b050300e0603551d0f0101f104030204c1d0603551d0e04160414764b1bd5
   e69935626e476b195a1a8c1f0603551d23041830165311296 6e304761732fbfe6
   a2c823c300d06092a620673410105050003204100474e5100a9cdaaa813b30f48
   40340fb17e7d6d6063064a5a7f2162301c464b5a8176623dfb1a4a484e618de1c
   3c3c5927cf590f4541233ff3c251e772a9a3f2c5fc6e5ef2fe155e5e385deb846
   b36eb4c3c7ef713f2d137ae8be4c022715fd033a818d55250f4e6077718180755
   a4fa677130da60818175ca4ab2af1d15563624c51e13dfdcf381881b72327e2f4
   9b7467e631a27b5b5c7d542bd2edaf78c0ac294f3972278996bdf673a334ff74c
   84aa7d65726310252f6a4f41281ec10ca2243864e3c5743103100
```

A.5. enrollstatus

[EDNOTE: Include CoAP message examples.]

A.6. voucher_status

[EDNOTE: Include CoAP message examples.]

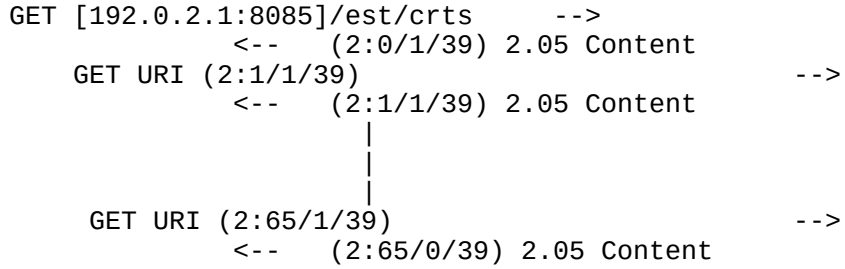
A.7. requestvoucher

[EDNOTE: Include CoAP message examples.]

Appendix B. EST-coaps Block message examples

This section provides a detailed example of the messages using DTLS and BLOCK option Block2. The minimum PMTU is 1280 bytes, which is the example value assumed for the DTLS datagram size. The example block length is taken as 64 which gives an SZX value of 2.

The following is an example of a valid /cacerts exchange over DTLS. The content length of the cacerts response in appendix A.1 of [RFC7030] is 4246 bytes using base64. This leads to a length of 2509 bytes in binary. The CoAP message adds around 10 bytes, the DTLS record 29 bytes. To avoid IP fragmentation, the CoAP block option is used and an MTU of 127 is assumed to stay within one IEEE 802.15.4 packet. To stay below the MTU of 127, the payload is split in 39 packets with a payload of 64 bytes each, followed by a packet of 13 bytes. The client sends an IPv6 packet containing the UDP datagram with the DTLS record that encapsulates the CoAP Request 40 times. The server returns an IPv6 packet containing the UDP datagram with the DTLS record that encapsulates the CoAP response. The CoAP request-response exchange with block option is shown below. Block option is shown in a decomposed way indicating the kind of Block option (2 in this case because used in the response) followed by a colon, and then the block number (NUM), the more bit (M = 0 means last block), and block size exponent (2**(SZX+4)) separated by slashes. The Length 64 is used with SZX= 2 to avoid IP fragmentation. The CoAP Request is sent with confirmable (CON) option and the content format of the Response is /application/ cacerts.



For further detailing the CoAP headers of the first two blocks are written out.

The header of the first GET looks like:

```
     Ver = 1
     T = 0 (CON)
     Code = 0x01 (0.1 GET)
     Options
      Option1 (Uri-Host)
        Option Delta = 0x3  (option nr = 3)
        Option Length = 0x9
        Option Value = 192.0.2.1
      Option2 (Uri-Port)
        Option Delta = 0x4   (option nr = 3+4=7)
        Option Length = 0x4
        Option Value = 8085
      Option3 (Uri-Path)
        Option Delta = 0x4    (option nr = 7+4=11)
        Option Length = 0x9
        Option Value = /est/crts
     Payload = [Empty]
```

The header of the first response looks like:

```
     Ver = 1
     T = 2 (ACK)
     Code = 0x45 (2.05 Content.)
     Options
       Option1 (Content-Format)
         Option Delta = 0xC  (option 12)
         Option Length = 0x2
         Option Value = TBD1
       Option2 (Block2)
         Option Delta = 0xB  (option 23 = 12 + 11)
         Option Length = 0x1
         Option Value = 0x0A (block number = 0, M=1, SZX=2)
     Payload =
30233906092a6206734107028c2a3023260201013100300b06092a6206734107018
c0c3020bb302063c20102020900a61e75193b7acc0d06092a6206734101
```

The second Block2:

```
     Ver = 1
     T = 2 (means ACK)
     Code = 0x45 (2.05 Content.)
     Options
       Option1 (Content-Format)
         Option Delta = 0xC   (option 12)
         Option Length = 0x2
         Option Value = TBD1
       Option2 (Block2)
         Option Delta = 0xB  (option 23 = 12 + 11)
         Option Length = 0x1
         Option Value = 0x1A (block number = 1, M=1, SZX=2)
     Payload =
   05050030
   1b311930170603550403131065737445786l6d706c654341204f774f301e170d313
   330353039303353333315a170d313430353039303353333315a
```

The 40th and final Block2:

```
     Ver = 1
     T = 2 (means ACK)
     Code = 0x21
     Options
       Option1 (Content-Format)
         Option Delta = 0xC   (option 12)
         Option Length = 0x2
         Option Value = TBD1
       Option2 (Block2)
         Option Delta = 0xB  (option 23 = 12 + 11)
         Option Length = 0x2
         Option Value = 0x272 (block number = 39, M=0, SZX=2)
     Payload = 73a30d0c006343116f58403100
```

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