



# Activity-Based Congestion Management for Fair Bandwidth Sharing in Trusted Packet Networks

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- ▶ The problem
- ▶ Definition of activity
- ▶ ABC principle
- ▶ Experiment setup
- ▶ Simulation results
- ▶ Related work
- ▶ Use cases
- ▶ Conclusion & outlook

Inspired by  
congestion  
policing using  
congestion  
exposure (ConEx)  
signals



## ► Scenario

- Bottleneck link shared by multiple users, congestion may occur
- Flows may be transmitted over different transport protocols  
⇒ UDP flows can achieve higher rates than TCP flows
- TCP users may have different number of flows  
⇒ Users with more flows can achieve higher rates (as TCP attempts to be per-flow fair)
- How to share bandwidth fairly among users?

## ► If pkts can be related to users

- Weighted fair queueing (WFQ) may enforce per-user fair bandwidth sharing

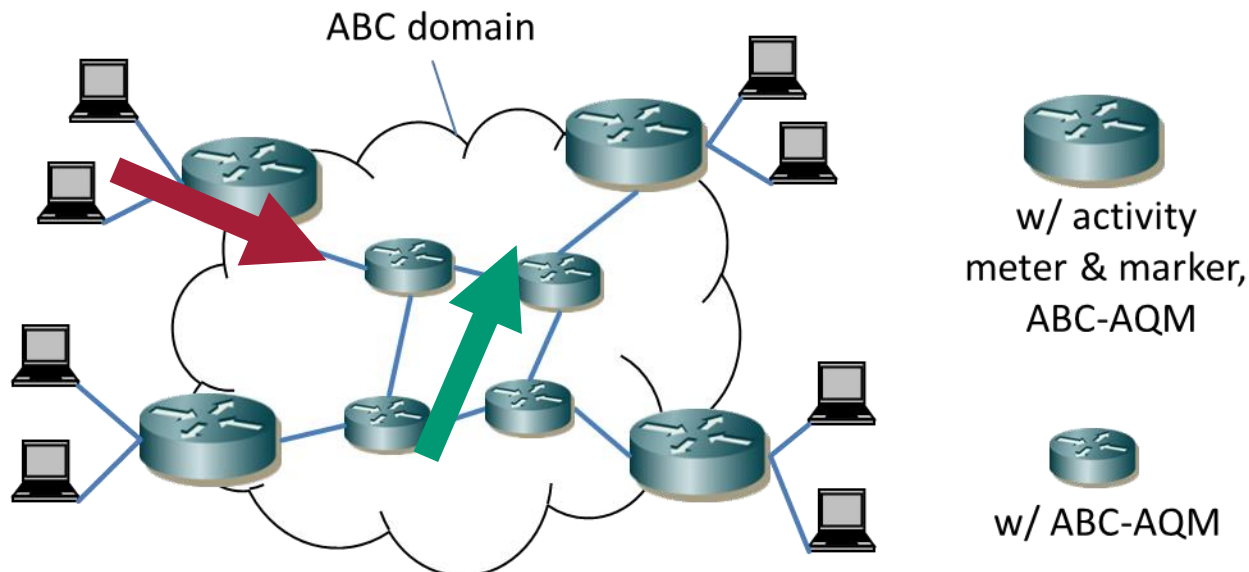
## ► If pkts cannot be related to users

- WFQ does not work
- **ABC may be used**



# ABC Principle – Architecture

- ▶ Per-domain concept
  - Configuration
  - Trust
- ▶ Access nodes
  - Connect user equipment
  - Have **activity meter**
  - Add activity  $A$  to packet headers
- ▶ Core and access nodes
  - Run **ABC-AQM** (active queue management)
  - Drop packets during congestion depending on their activity
- ▶ Core nodes do not hold per-user/flow states





## ► Definition of activity

- $R_r$ : a user's assigned reference rate
- $R_t$ : a user's recent transmission rate
- $A = \log_2 \left( \frac{R_t}{R_r} \right)$ : a user's activity
- Indicates how much a users transmission rate exceeds its reference rate

## ► Activity meter

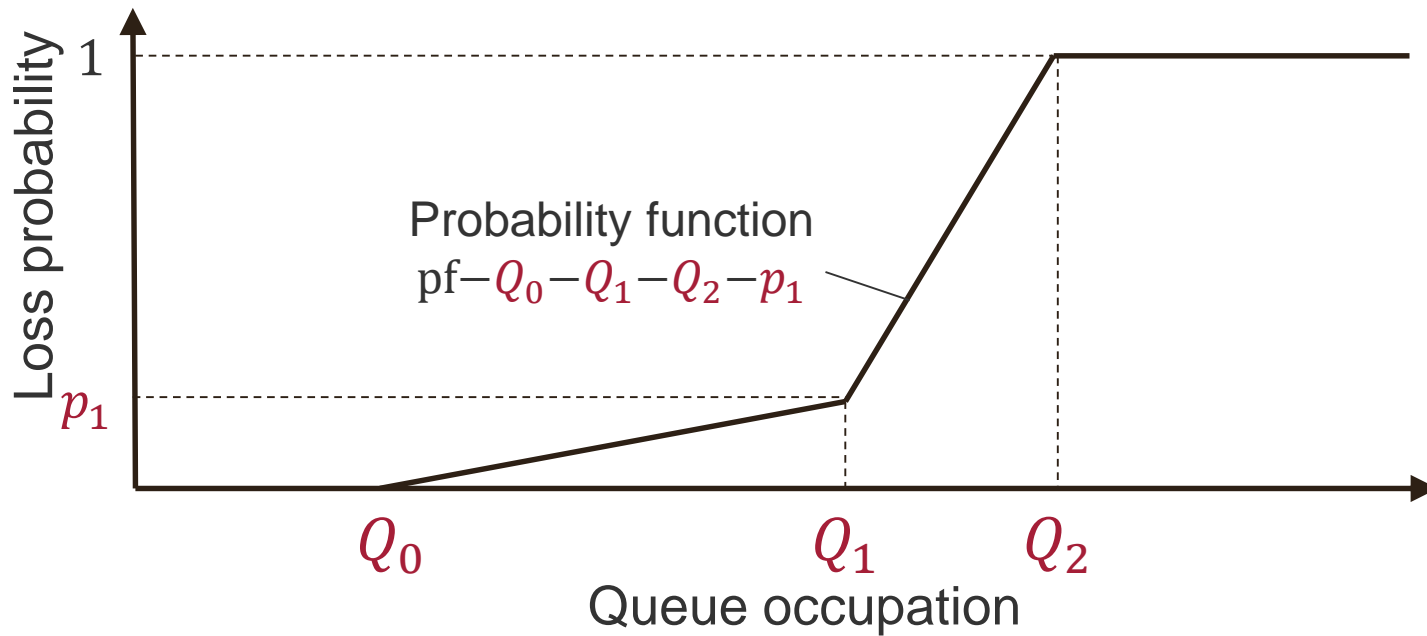
- Computes activity  $A$  per packet and adds it to the header
- Implemented through a token-based algorithm
- Configured with
  - Reference rate  $R_r$  (per user)
  - Inertia  $I$  (time scale parameter, per domain)
  - Burst allowance  $B$  (to ignore first  $B$  bytes after silence, per domain)



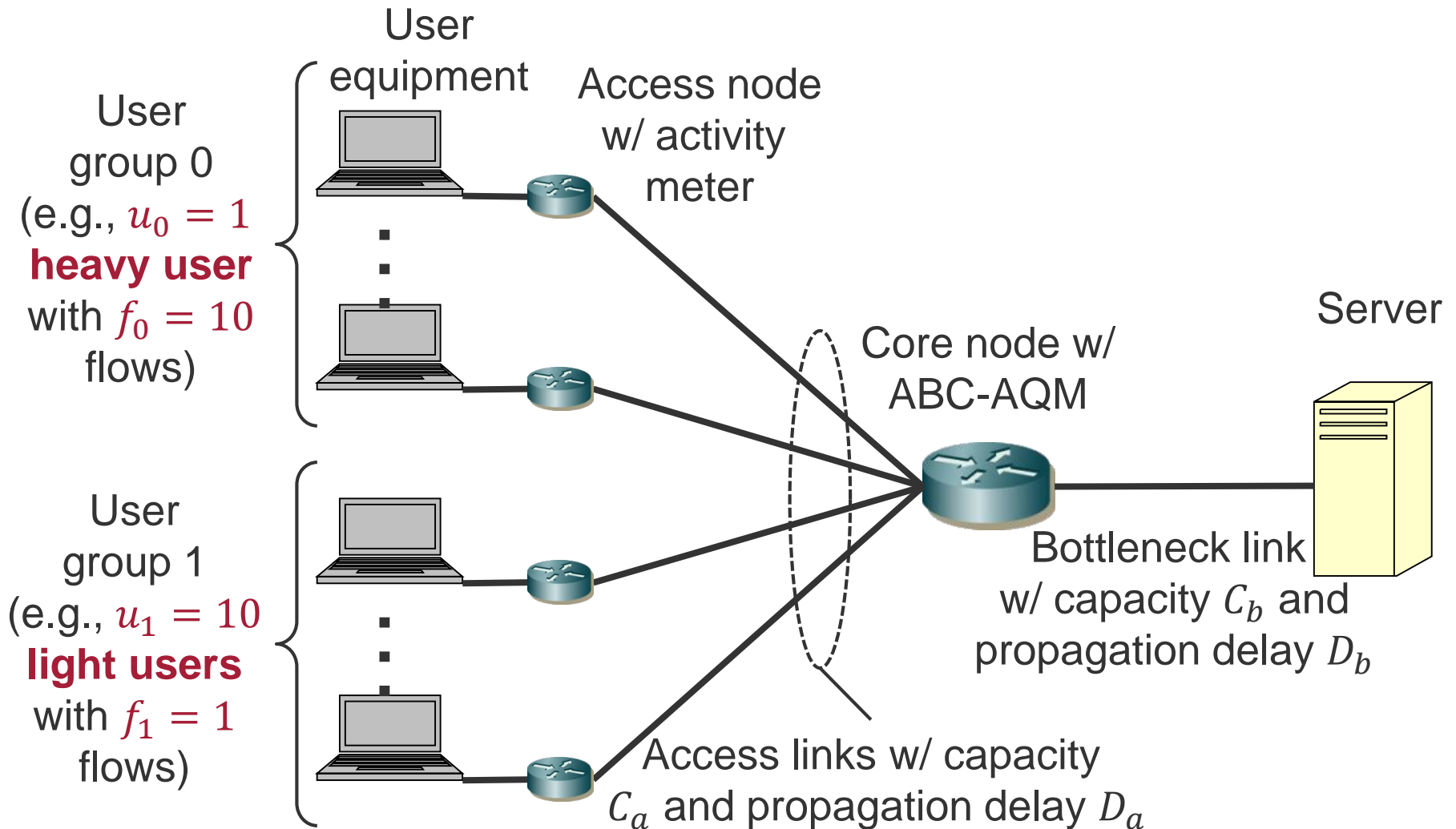
- ▶ Operation per packet arrival
  - Track average activity  $A_{avg}$
  - Drop packet depending on
    - Link congestion
    - Activity  $A$  in packet header
  
- ▶ Average activity:  $A_{avg} = w_A \cdot A_{avg} + (1 - w_A) \cdot A$
  
- ▶ Use some existing AQM based on drop probabilities
  
- ▶ Adapt drop probability:  $p_A = p^{2^{-\gamma \cdot (A - A_{avg})}}$ 
  - Differentiation factor  $\gamma \geq 0$
  - $p_A > p$  for large activities
  - $p_A < p$  for small activities



# AQM Used in the Study (Exchangeable)



$$p(Q) = \begin{cases} 0 & Q < Q_0 \\ \frac{Q - Q_0}{Q_1 - Q_0} \cdot p_1 & Q_0 < Q \leq Q_1 \\ p_1 + \frac{Q - Q_1}{Q_2 - Q_1} \cdot (1 - p_1) & Q_1 < Q \leq Q_2 \\ 1 & Q_2 < Q \end{cases}$$







# Impact of Differentiation Parameter $\gamma$

## ► Experiment

- User group 0: 1 heavy user with 10 TCP flows
- User group 1: 10 light users with only 1 TCP flow

## ► Performance metrics

- Average throughput per user
  - $\bar{T}_0$ : for user group 0
  - $\bar{T}_1$ : for user group 1
- Throughput ratio  $T_R = \frac{\bar{T}_0}{\bar{T}_1}$

Throughput ratio  $T_R$  depending on differentiation factor  $\gamma$ .

No differentiation, no fairness

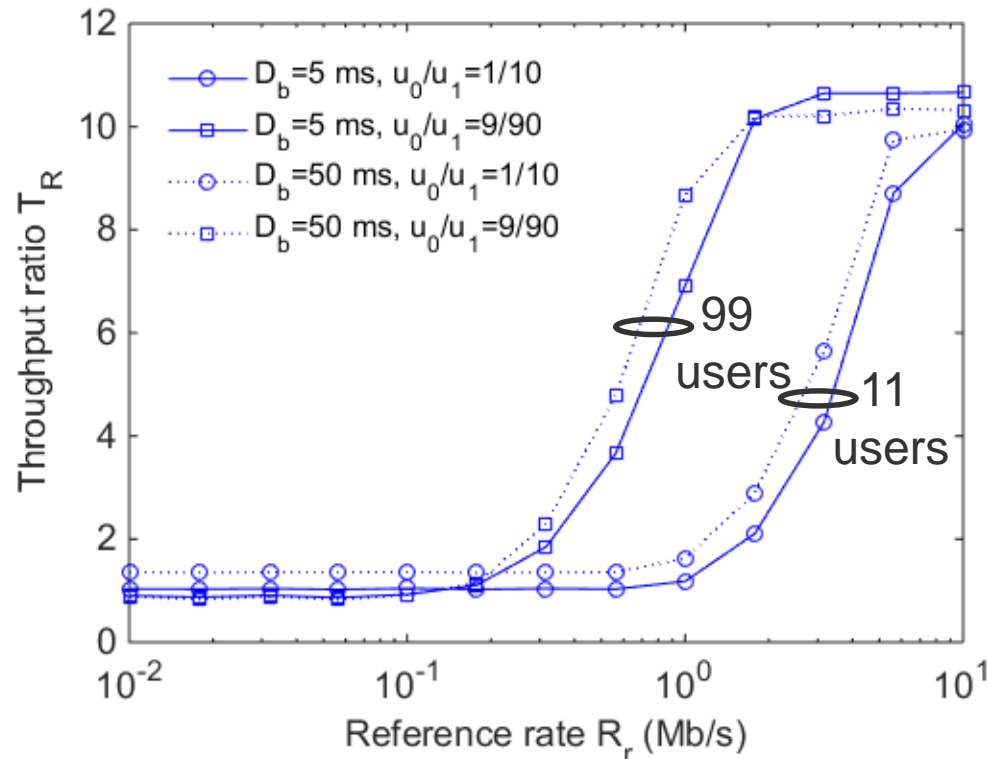
Almost fair

$D_b$ (ms)	users $u_0/u_1$	differentiation factor $\gamma$					
		0	1	2	3	4	5
50	1/10	10.02	2.68	1.71	1.36	1.20	1.10
	9/90	10.20	1.51	1.03	0.88	0.82	0.80
5	1/10	10.14	1.90	1.26	1.03	0.94	0.89
	9/90	10.76	1.37	1.01	0.91	0.87	0.84

- Use  $\gamma = 3$  to achieve good fairness, confirmed by more experiments

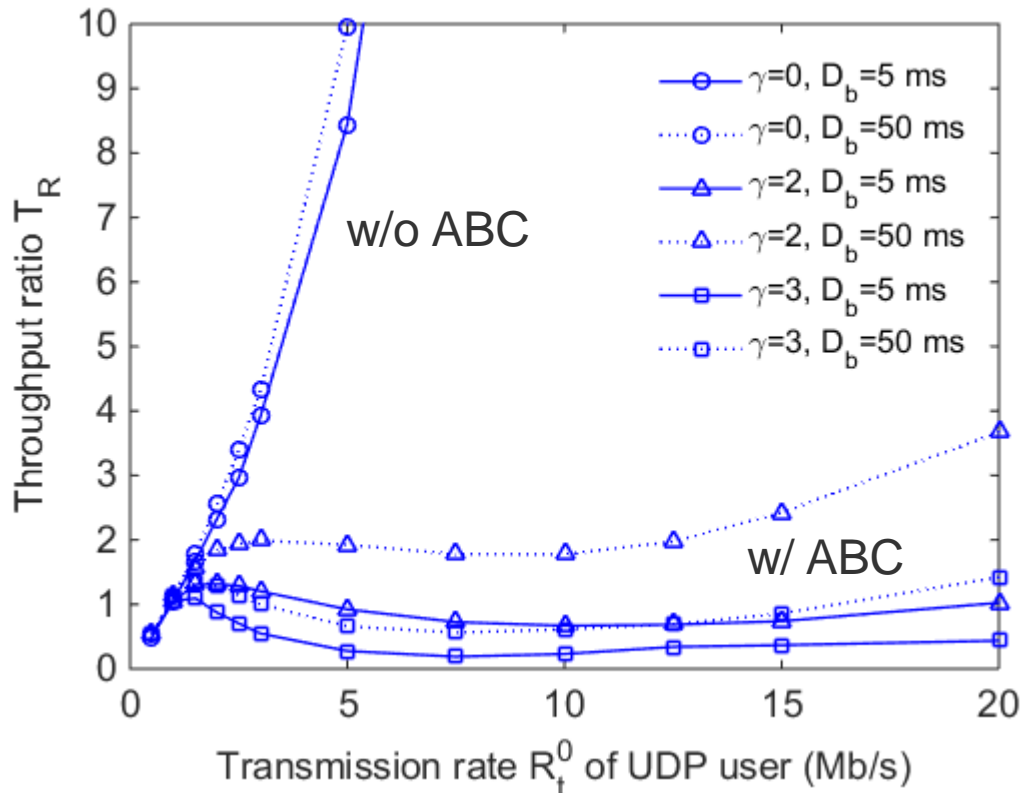


- ▶ Vary reference  $R_r$  rate per user
- ▶ ABC enforces fairness if  $R_r \leq \frac{C}{u}$ 
  - $C$ : bottleneck bandwidth
  - $u$ : number of users on bottleneck link





- ▶ 1 UDP user vs. 10 users with 1 TCP flow
- ▶ ABC protects TCP users against UDP user
- ▶  $\gamma = 3$  best suited
- ▶ Improvement through better AQM possible





## ► Experiment setup

- $u_0 = 1$  **probe user** sending a probe of 1 MB
- $u_1$  **background users** with  $f_1$  flows each
- Performance metric: upload time
- Link delay  $D_b$ 
  - 50 ms
  - 5 ms

## ► Upload time

- Increases with  $u_1$  and  $f_1$
- Shorter with ABC
- Decreases with increasing **burst tolerance  $B$**

$D_b$ (ms)	ABC	$B$ (MB)	$f_1 = 1$		$f_1 = 4$	
			$u_1 = 10$	$u_1 = 20$	$u_1 = 10$	$u_1 = 20$
50	no	n/a	20.3	29.2	47.7	85.2
	yes	0	9.6	17.9	12.3	21.7
		0.25	8.2	15.7	10.4	17.6
		0.5	6.1	10.7	8.3	14.0
		0.75	2.2	5.7	6.7	8.1
		1	2.4	2.7	3.3	5.4
		1.25	2.1	3.3	3.3	5.2
5	no	n/a	10.3	20.2	34.2	61.4
	yes	0	9.6	17.6	10.1	18.8
		0.25	6.9	13.7	8.3	14.0
		0.5	4.6	9.5	5.6	10.0
		0.75	1.3	5.4	4.0	6.4
		1	1.0	1.3	1.0	1.4
		1.25	1.0	1.0	1.0	1.0



## ▶ CSFQ

- [28] I. Stoica, S. Shenker, and H. Zhang: „**Core-Stateless Fair Queuing**: A Scalable Architecture to Approximate Fair Bandwidth Allocations in High-Speed Networks“, IEEE/ACM Transactions on Networking, 11(1), Feb. 2003
- Similar to ABC but measures traffic rate on link to determine drop probabilities  $\Rightarrow$  slower reaction than ABC

## ▶ ABC

- Uses buffer occupation and AQM to determine drop probabilities
- Extensible towards low-delay communication  
 $\Rightarrow$  Enforce both fairness and low-delay



- ▶ ABC fulfills objectives of congestion policing using congestion exposure (ConEx)
  - Fair rate sharing
  - Improved QoE for light users (reduced upload times)
  
- ▶ ABC supports the same uses cases as ConEx
  - Datacenter  
draft-briscoe-conex-data-centre
  - Mobile access networks  
RFC7778: Mobile Communication Congestion Exposure Scenario
  
- ▶ ABC protects light users against heavy UDP users
  - Another use case: defense against denial of service attacks



- ▶ Activity-based congestion management (ABC)
  - Access node add activity into to packet headers
  - Access and core nodes use it for preferential dropping
  - Enforces per-user fairness (heavy and light users)
  - Protocol-independent, stateless core network
- ▶ Presented simulation results
  - Impact of differentiation factor, reference rate
  - Protection against non-responsive traffic
  - Reduced upload time
- ▶ ABC supports
  - Different users profiles (reference rates)
  - ConEx use cases (and more)
- ▶ Outlook
  - Combine ABC with PIE to achieve fairness and low latency
  - Adapt ABC towards service classes



- ▶ M. Menth and N. Zeitler: Activity-Based Congestion Management for Fair Bandwidth Sharing in Trusted Packet Networks, in Proceedings of the IEEE/IFIP Network Operations and Management Symposium (NOMS), April, 2016, Istanbul, Turkey
- ▶ Preprint:  
<https://atlas.informatik.uni-tuebingen.de/~menth/papers/Menth16b.pdf>