



Post-Quantum

# Hybrid Quantum-Safe Key Exchange for IKEv2

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- **The viability of quantum computers is an engineering problem**
- **Threats to Diffie-Hellman key-exchange in IKEv2**
- **There are a number of algorithms that are believed to be quantum-safe.**
  - NIST call for proposals on quantum-safe algorithms
- **FIPS compliant**



- **Optional key exchange payload carrying quantum-safe public data**
- **Goals:**
  - To allow quantum-safe key exchange to be used alongside DH key-exchange while we are transitioning to a post-quantum era.
  - To keep the modifications to IKEv2 to a minimum
  - To maintain compatibility with IKEv2
  - To provide a path to phase out vulnerable DH key exchange in the future (if required)



Initiator

Responder

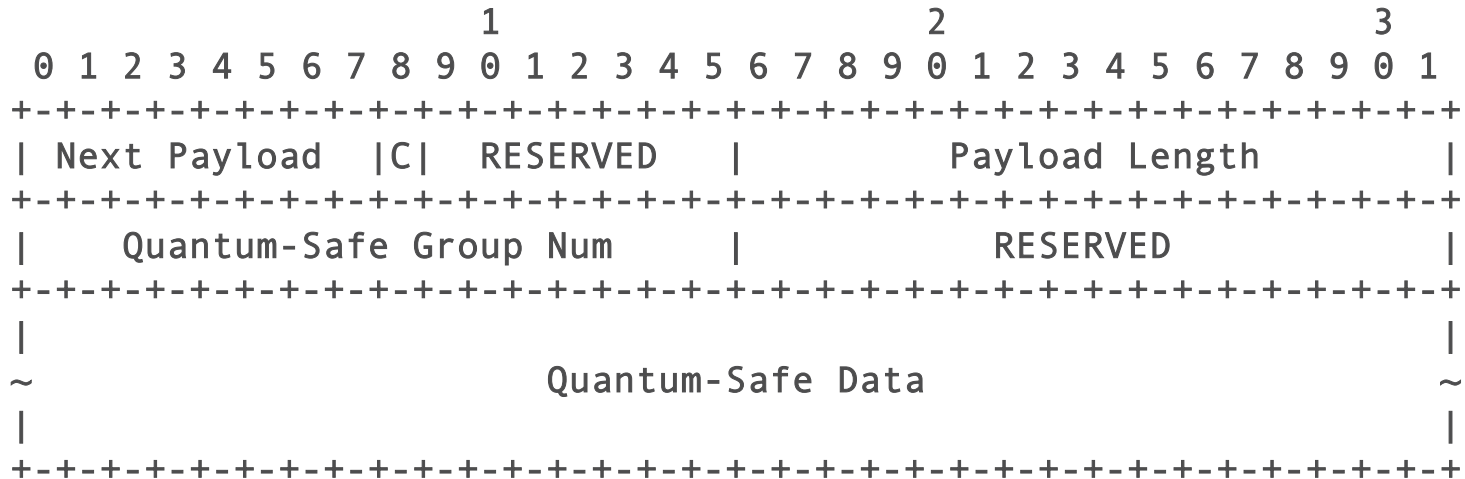
HDR, SAi1, KEi, [QSKEi,]

Ni

-->

<-- HDR, SAR1, KEr, [QSKEr,]  
Nr, [CERTREQ]

$$\text{SKEYSEED} = \text{prf}(\text{Ni} \mid \text{Nr}, \text{g}^{\text{ir}} \mid \text{QSSS})$$



- Initiator understands QSKE, but Responder does not.
- **Policy-based fallbacks: allows non quantum-safe SAs to be established**
  - SAi payload contains a combination of proposals, some offer KE and the rest offer both KE and QSKE.
  - QSKE payload is marked NON-CRITICAL
- **Policy-based fallback: only quantum-safe SAs are allowed**
  - SAi payload contains no proposal that offers KE only;
  - QSKE payload is marked CRITICAL (if required)
- **What current IKEv2 implementations will do if it receives a proposal containing a transform type that it doesn't understand?**



## Creating a new Child SA

```
HDR, SK{SA, Ni, [KEi,]
      [QSKEi,] TSi, TSr}    -->

                                <-- HDR, SK{SA, Nr, [KEr,]
                                [QSKEr,] TSi, TSr}

KEYMAT = prf+(SK_d, g^ir(new) | QSSS(new) | Ni | Nr)
```

## Rekeying a Child SA

```
HDR, SK{N(REKEY_SA), SA
      Ni, [KEi,] [QSKEi,]
      TSi, TSr}    -->

                                <-- HDR, SK{SA, Nr, [KEr,]
                                [QSKEr,] TSi, TSr}
```



## Rekeying an IKE SA

HDR, SK{SA, Ni, KEi,  
[QSKEi,]}

-->

<-- HDR, SK{SA, Nr,  
KEr, [QSKEr,]}

SKEYSEED = prf(SK\_d(old), g<sup>ir(new)</sup> | QSSS(new) | Ni | Nr)



- QSKE payload is large
- Fragmentation mitigation methods considered:
  - Send QSKE payload as part of IKE\_AUTH
  - Immediate CHILD\_SA rekeying right after IKE\_AUTH
  - Use QSKE payload of small size in IKE\_SA\_INIT, e.g. SIDH; more options in CREATE\_CHILD\_SA
  - Introduce a new state between IKE\_SA\_INIT and IKE\_AUTH that is for exchanging QSKE payload
  - TCP encapsulation



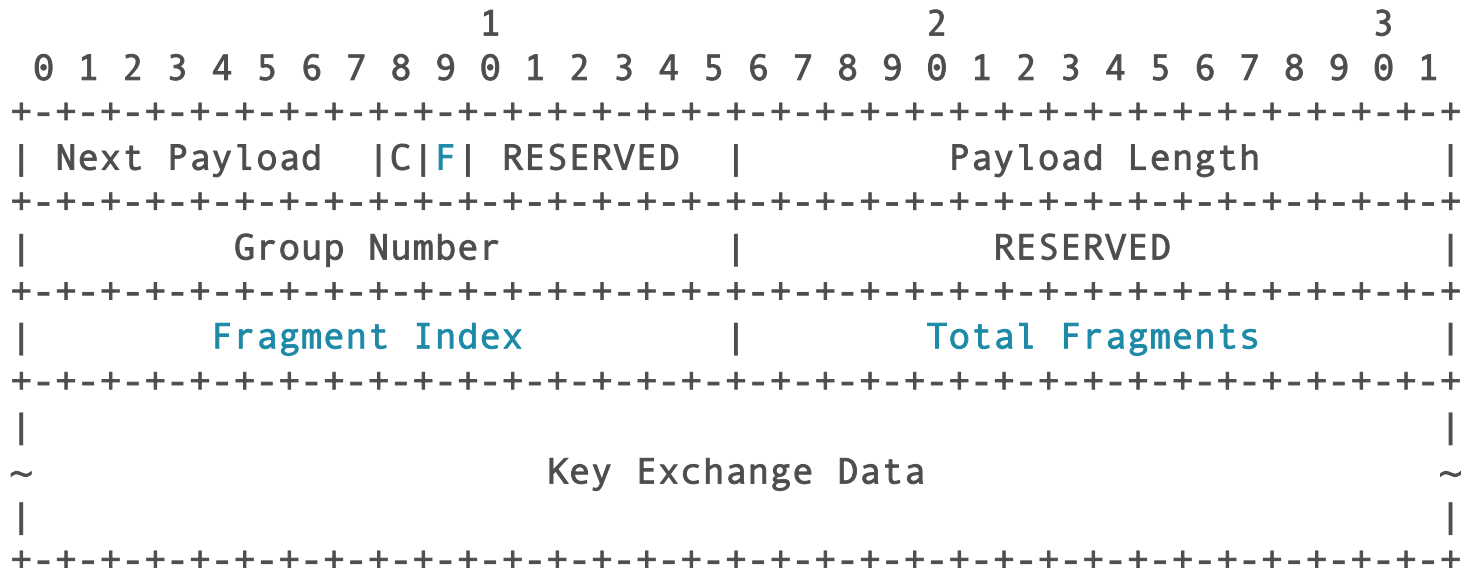


- Address fragmentations, crypto-agility and backward compatibility.

```

HDR, SAi, KEi, Ni, N(QSKE) -->
                                <-- HDR, SAR, N(COOKIE), [N(QSKE)]
HDR, N(COOKIE), SAi, KEi,
    Ni, QSKEi(1/2)             -->
HDR, QSKEi(2/2)               -->
                                <-- HDR, SAR, KEr, Nr, QSKEr(1/2)
                                <-- HDR, QSKEr(2/2)
                                IKE_AUTH exchange

```





- As a demonstration, check-out our forked of strongSwan: <https://github.com/post-quantum/strongswan/blob/qske/README.QSKE.md>
- Note that it is in **qske** branch