

# Proposal for Fast Subflow Creation

---

Quentin De Coninck

quentin.deconinck@uclouvain.be

Universite Catholique de Louvain

IETF 99, Prague

From "Every Millisecond Counts: Tuning Multipath TCP for Interactive Applications on Smartphones", tech. report.

<http://hdl.handle.net/2078.1/185717>

## Why do we propose this?

Current implementations (e.g., Linux one) opens subflows on all available interfaces upon connection establishment

→ Great for bandwidth aggregation

But *Make-Before-Break* is sometimes useless

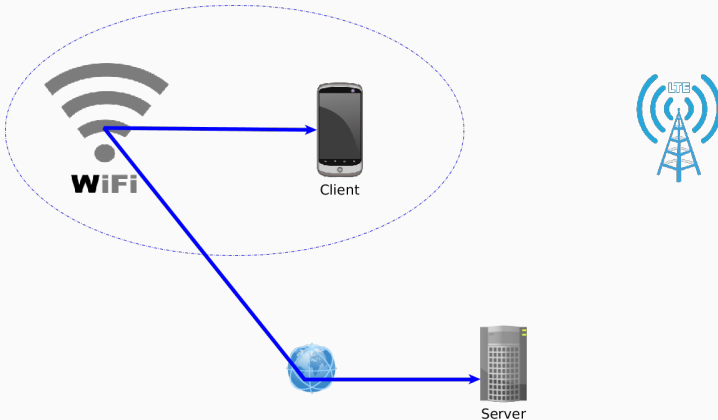
- Connection finished before establishment of additional subflows
- Application not requiring lot of capacity
  - But rather low latency...
  - ...and seamless network handover

Useless subflows waste network/energy resources

# The Smartphone Usecase

Typically wants to remain on the WiFi if good

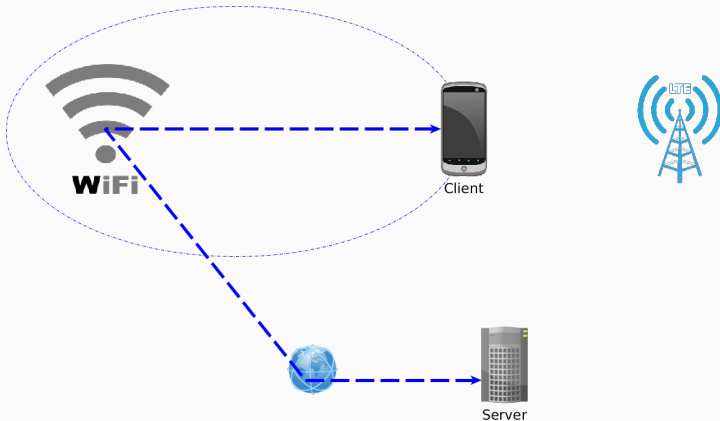
- Only create cellular subflows if bad/no WiFi
  - Or if traffic actually requires large bandwidth



# The Smartphone Usecase

Typically wants to remain on the WiFi if good

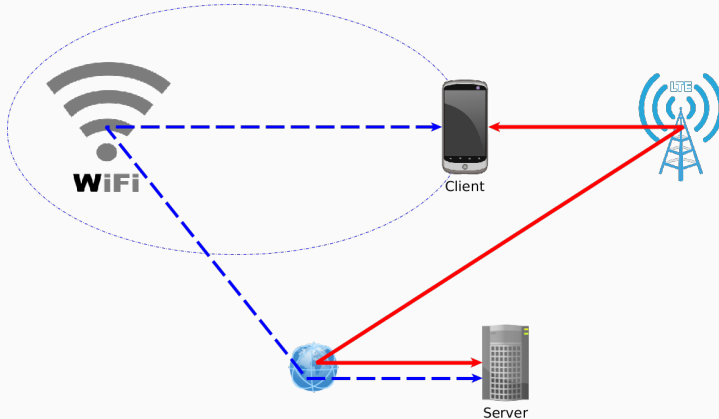
- Only create cellular subflows if bad/no WiFi
  - Or if traffic actually requires large bandwidth



# The Smartphone Usecase

Typically wants to remain on the WiFi if good

- Only create cellular subflows if bad/no WiFi
  - Or if traffic actually requires large bandwidth



## Proposing a *Break-Before-Make* Approach

We propose three mechanisms for the smartphone usecase

- "Global Scheduling": server following client's choices
- "Multipath TCP Oracle": detecting bad performing subflows
- **"Immediate Reinjections": fast subflow creation**

Only discussing the last one here, as it affects interoperability

# Limiting Handover Delay

The backup (cellular) path creation is delayed

- Nice from a energy consumption point of view...
- ...but incurs larger app perceived latency in mobility cases
  - Reactive approach: need to detect first bad network

Furthermore, additional Multipath TCP path creation takes time...

# Towards Fast Establishment of Additional Subflows

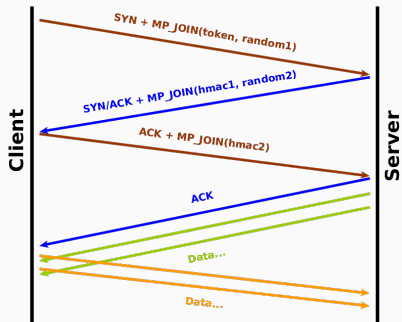


Figure 1: Normal JOIN.



# Towards Fast Establishment of Additional Subflows

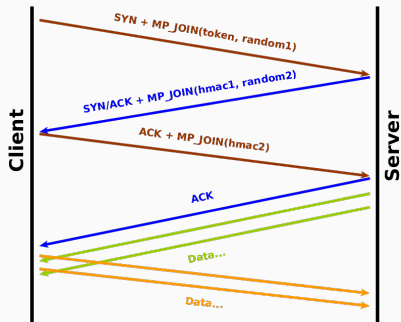
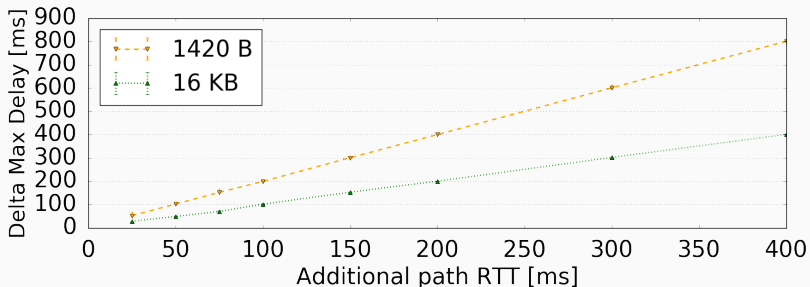


Figure 1: Normal JOIN.



Figure 2: Fast JOIN with data.

## Quantifying Latency Gains (Request/Response – Mininet Setup)



**Figure 3:** Latency gain between Fast and Normal Joins depending on the request size.

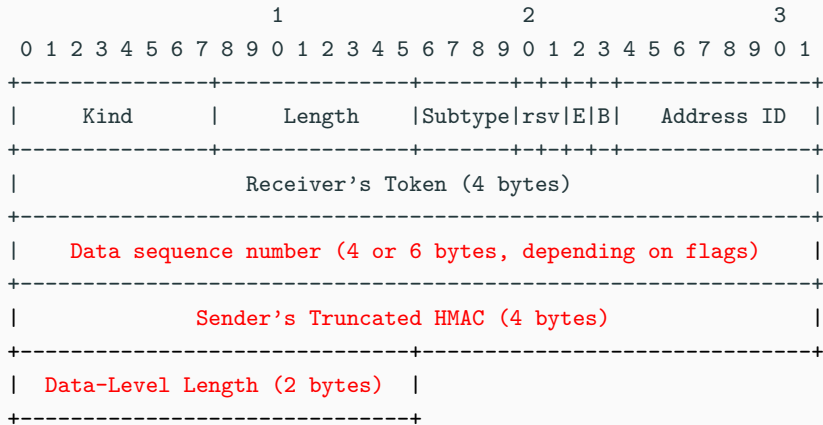
- Saving at least 1 RTT
- Saving 2 RTTs if request size < MSS

# Defining New Multipath TCP Options

Two new proposed options

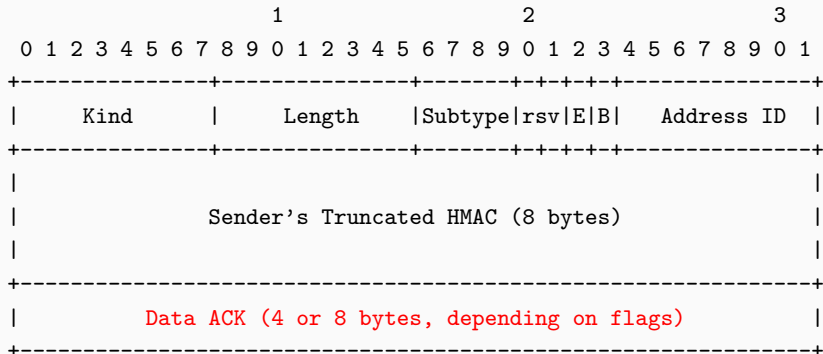
- FAST\_JOIN\_OUT
  - When the client has data to send
- FAST\_JOIN\_IN
  - When the client has no data to send
    - But knows that main path(s) failed
    - And the server might still have data to send (e.g., middle of bulk download)

# FAST\_JOIN\_OUT Option Format – Initial SYN



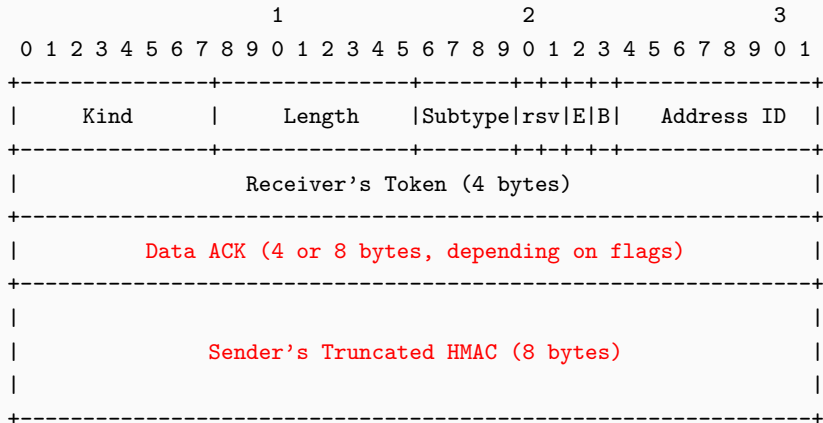
With  $HMAC_{CS}(DSN, token)$

## FAST\_JOIN\_OUT Option Format – SYN/ACK



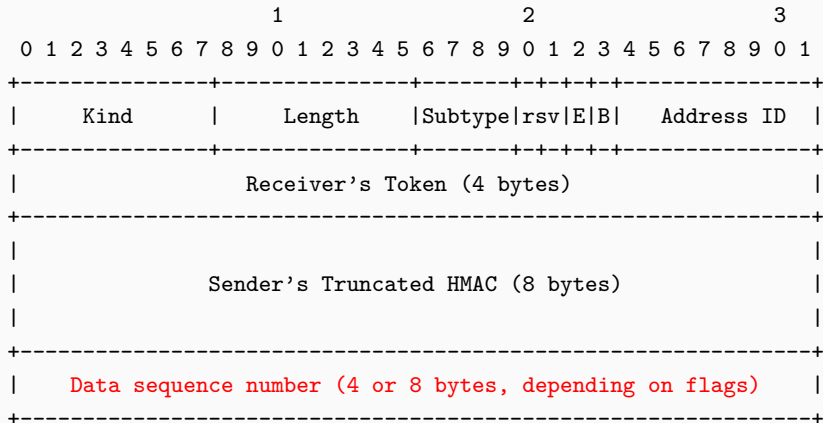
With  $HMAC_{SC}(DataACK, DSN)$

## FAST\_JOIN\_IN Option Format – Initial SYN



With  $HMAC_{CS}(DataACK, token)$

## FAST\_JOIN\_IN Option Format – SYN/ACK



With  $HMAC_{SC}(DSN, DataACK)$

# Possible Security Considerations

- Shortened HMAC size
  - Try to save as much TCP option space as possible
- SYN replay attacks
  - Prevent more than 1 subflow creation with a given DSN/Data ACK
  - Limit number of fast created subflows per connection



## To Summarize

### Tuning Multipath TCP for smartphone

- Cellular subflow consumes radio resources
- Only use cellular when needed

### Applying changes at MPTCP design to suit this usecase

- Implemented in Linux MPTCP v0.91
- And in Nexus 5 (with Android 6.0.1)

→ Interest in writing a draft about this?

**Thanks for your attention!**

**Feel free to ask questions or provide  
feedback 😊**