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CCNX INTEREST AGGREGATION

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PRINCIPLES

- Minimize coupling between consumer and forwarders
- No timers on forwarders
- Provable properties

DIFFERENCE FROM ICNRG -03 DRAFT

- Spelled out algorithms in more detail
- Adopted Interest HopLimit decrement idea [1]

[1] J.J. Garcia-Luna-Aceves and M. Mirzazad-Barijough. Enabling Correct Interest Forwarding and Retransmissions in a Content Centric Network. In *Proc. ACM ANCS'15*, pages 135–146, Oakland, California, 2015.

PROPERTIES (1-3)

1. **Aggregation:** If there are no losses and no retransmissions, then for each link in the forwarding path exactly one Interest and at most one ContentObject traverse the link.
2. **Retransmissions:** If there are retransmissions by one or more consumers and those retransmissions do not pass a ContentObject response in flight, then at most one ContentObject traverses the reverse path.
3. **Interest In Flight:** If a retransmission passes a ContentObject in flight on a given link, then that Interest will propagate only as far as the first cached copy of the response.

PROPERTIES (4-6)

4. **ContentObject In Flight:** If a retransmission passes a ContentObject in flight on a given link, then there will be a duplicate ContentObject sent on that link, but it will not propagate further than that link.
5. **Cycle Termination:** An Interest that travels in a cycle will not repeat the same cycle.

PROPERTY 6: CYCLE TERMINATION

- The **Cycle Termination** property is significantly different than prior CCNx 0.x and NDN cycle termination. In those protocols, because of the nonce, an Interest that visits the same node twice – assuming the nonce is remembered long enough – is dropped.
- This means that such a node cannot continue using other paths, even if those will reach the destination.
- In CCNx 1.0, the cycle termination property means that an Interest will not repeat the same cycle, but it could be forwarded along another shorter path.
- Described in more detail in a few slides.

HOPLIMIT

- An Interest predecessor may be remote (i.e. another node) or local (i.e. an application connecting directly to the forwarder).
- The HopLimit indicates the number of allowed remote node hops.
- Forwarding to and from an application does not count as a hop.
- This semantic allows an implementation, if desired, to indicate “for me” in the FIB table via routes with a 0 HopCount, which is the notional model we use here. As this is internal behavior of a forwarder, it is not standardized.

HopLimit must
decrease each hop.

Algorithm 1 Receive Interest

```
1: procedure RECEIVEINTEREST(Predecessor P, Interest I)
2:   if P is remote then
3:     if  $I.HopLimit = 0$  then
4:       Send InterestReturn (HopLimit)
5:       Drop Interest
6:     else
7:       Decrement  $I.HopLimit$ 
8:       end if
9:     end if
10:    if Satisfy I from ContentStore then
11:      Send ContentObject to P
12:    else
13:      Verdict  $\leftarrow$  AGGREGATE( $P, I$ )
14:      if Verdict is Forward then
15:        if FORWARDINTEREST( $P, I$ ) is false then
16:          Drop Interest
17:        end if
18:      end if
19:    end if
20:  end procedure
```

Only first Interest from
a new predecessor for
existing PIT entry is
aggregated.

FIB hop count must not
exceed the
decremented HopLimit
(Alg 1, Line 7)

Algorithm 2 Interest Aggregation

```
1: function AGGREGATE(Predecessor P, Interest I)
2:   if does not exist a PIT entry for I then
3:     Create PIT entry with Predecessor P
4:     Return Forward
5:   else if PIT entry exists with Predecessor P then
6:     Retransmission from P
7:     Return Forward
8:   else
9:      Add Predecessor P to PIT Entry
10:    Return Aggregate
11:   end if
12: end function
```

Algorithm 3 FIB Lookup

```
1: function FIBLOOKUP(Predecessor P, Interest I)
2:   Exclude FIB entries that point to P
3:   NextHops  $\leftarrow$  LongestPrefixMatch(I)
4:   for N  $\in$  NextHops do
5:      if I.HopLimit  $<$  N.HopCount then
6:       Remove N from NextHops
7:     end if
8:   end for
9:   Return NextHops
10: end function
```

Interest output has
HopLimit set to the FIB's
HopCount (which must
be less than the input
HopLimit, Alg 3 Line 5)

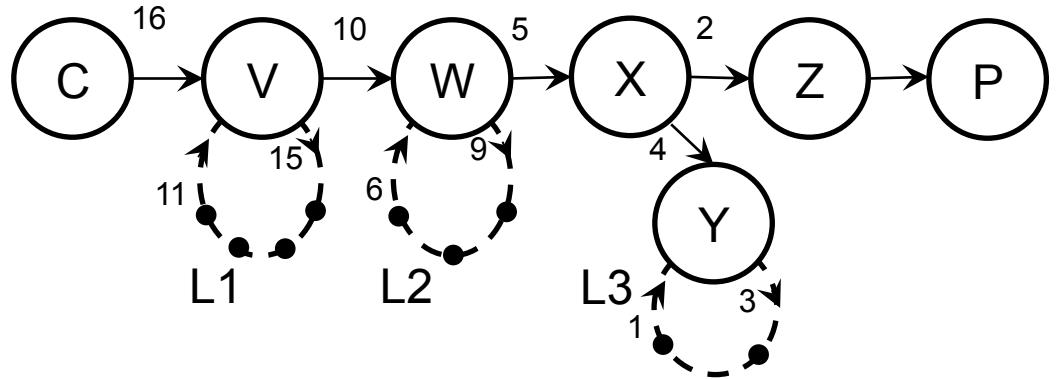
Algorithm 4 Forward Interest

```
1: function FORWARDINTEREST(Predecessor P, Interest I)
2:   NextHops  $\leftarrow$  FIBLOOKUP(P, I)
3:   if NextHops is not empty then
4:     for N  $\in$  NextHops do
5:       I.HopLimit  $\leftarrow$  N.HopCount
6:       Send I to N
7:     end for
8:     Return true
9:   else
10:    Send InterestReturn (NoRoute)
11:    Return false
12:  end if
13: end function
```

PROPERTY PROOFS

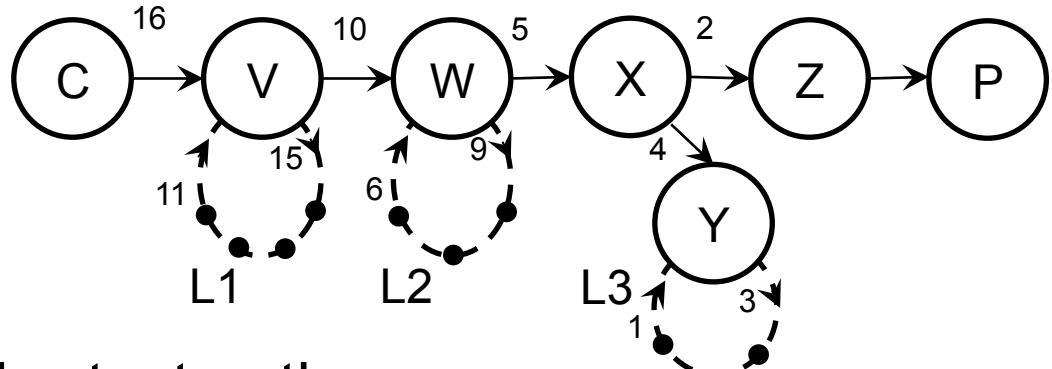
- Not presented here, written up in paper submission.
- Can sketch out the proofs in Q&A if you have specific questions.

CYCLES



- Use this topology for example.
- Assumes each forwarder uses unequal cost multipath.
- Numbers indicate the FIB HopCount of each successor.
- Interest sent from consumer C towards producer P.

CYCLES



- **C-V-W-X-Z-P:** The shortest path.
- **C-V-L1-V-W-X-Z-P:** The cycle L1 is permissible because the HopLimit will be 15 on entering the loop and 11 on exiting the loop, so the successor W is feasible.
- **C-V-W-L2-W-X-Z-P:** The cycle L2 is permissible because the HopLimit will be 9 on entering the loop and 6 on exiting the loop, so the successor X is feasible.
- **C-V-L1-V-W-L2-W-X-Z-P:** As in prior two cases.
- **C-V-W-X-Y-L3-Y-{}:** The path via Y will terminate after the first cycle because the HopLimit on exiting L3 will be 1 and Y has no successor with a feasible HopCount.

CONCLUSION

- CCNx 1.0 Interest aggregation algorithm
 - Does not use timers.
 - Does not depend on the ARQ mechanism (nor does the ARQ mechanism depend on the Interest aggregation scheme).
- Has desirable properties
 - At most 1 interest forwarder if no retransmissions.
 - At most 1 ContentObject downstream per link, even with retransmissions (ignoring in-flight misses).
 - Interest will not travel a cycle more than once.