#### A New Approach to Network Functions





Aurojit Panda







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NFs built to target hardware & software features.





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Core placement for performance.

#### Problems with the Current Approach

High overheads for isolation and chaining.



### High Overheads for Isolation

#### No Isolation



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# High Overheads for Chaining



Chain Length

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- **Programmers** responsible for meeting NF performance requirements.
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- Result: Largely written by the same companies that built middleboxes.
  - Hard for carriers or new entrants to develop NFs, limiting innovation.

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- Feature availability dictated by deployment environment, use dictated by vendor.
- Result: Delays before new features are used, increased cost for upgrades.



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NetBricks Addresses these Problems

A new execution environment and programming framework for NFs.

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### NetBricks: Execution Environment

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  - Consolidation: Maximize number of NFs that can be consolidated.
  - Isolation: Ensure NFs do not affect each other

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- **Performance Isolation**: One NF does not affect another's performance.
  - Ongoing work: Partition last level cache.

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- Enables consolidation for NFs like SSL proxies that have secrets.
  - Isolation is a building block for protecting secrets in applications.
- Simplifies programming: don't need to worry about other programs and NFs.
  - Lack of isolation between drivers is a major cause of crashes in OSes.





M/Container	VM/Container

Memory Isolation

Packet Isolation







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M/Container	

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### Packet Isolation







#### VM/Container



### Packet Isolation







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### Packet Isolation







### NetBricks: Low Overhead Isolation





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### Runtime mechanisms too expensive for NFV workloads.

- Process a packet approximately every 100ns (10 MPPS) or faster.
- Must rely on static compile-time checks for isolation.

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  - Check isolation almost entirely at compile time, limited runtime overhead.
- Built on Rust type checks, bound checks, no garbage collection.
- Framework designed to meet the rest of the memory isolation requirements.



# Approaches to Packet Isolation



- Existing approaches convert this **temporal problem** to a **spatial one**.
  - Copy packets from one packet space to another.



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- Solution: Use **linear types** (1990s) for isolation.
- Syntax marks argument that are moved.
  - Argument moved during calls.
    - Ownership is transferred to callee.
  - Moved variables can not be reused.

fn consume(a: Packet) { // Work with packet. }

// pkt is a packet consume(pkt);

pkt.set\_length(200)





NF B



NF A



NF B

NF A



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- API is designed so that safe code can never learn packet buffer address.
- Assuming compiler is sound packet isolation is guaranteed.





**DPDK:** Fast packet I/O.



#### NF Chains: Units of scheduling





-Do not preempt NF chain.



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  - -Reduces working set size.
- -Preemption points added using queues



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- Provides low overhead memory and packet isolation.
- Improved consolidation: multiple NFs can share a core.
  - Context switch (~1 $\mu$ s) vs function call to NF (~ few cycles = few ns).
- Reduce **memory** and **cache pressure**.
  - Zero copy I/O => do not need to copy packets around.

#### **VM/Container**



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#### NetBricks



#### **VM/Container**



**NetBricks** 



#### **NetBricks Multicore**





#### **VM/Container** NF2 vSwitch NF •• •• •• •• •• •• •• $\bullet$ $\bullet$ • •

**NetBricks** 



#### **NetBricks Multicore**







Processing Cycles Per Packet

## NetBricks: More Efficient


Processing Cycles Per Packet

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Processing Cycles Per Packet

### NetBricks: More Efficient

## Scaling with Chain Length



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NetBricks: Programming Environment

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- User defined functions (UDFs) provide flexibility.
  - Insight: customization is largely orthogonal to performance
- Framework can implement **global optimization**.



### Abstractions

Packet Processing Parse/Deparse Transform Filter





### State

Lookup Tables

LPM Tables

### Abstractions









### Abstractions and UDFs

.parse::<MacHeader>()

.parse::<IpHeader>()

.parse::<TcpHeader>()

f(pkt.header().src\_port())

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  - Reported: **2.6 MPPS**

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- More than **5x** better than commercial **EPCs** based on DPDK.

### Upgrading NFs through abstractions

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Warning: Future work ahead.


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- Choose which version to use at deployment time.
  - Choice depends on what is supported, and resource scheduling.



Lookup Table (Current)

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Adopting Sateless Abstraction

Lookup Table (Current)



Adopting Sateless Abstraction



Abstraction backed by remote KV-store

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Abstraction backed by remote KV-store

UDFs remain unchanged.

Adopting Sateless Abstraction



Abstraction backed by remote **KV-store** 

**UDFs** remain unchanged.

Adopting Stateless and Caching



Use consistency requirements to implement local caching.

UDFs remain unchanged.



Adopting Sateless Abstraction



#### Adopting new features requires no changes to NF code.

#### Becomes a **policy decision** made by **deployment**.

UDFs remain unchanged.

Adopting Stateless and Caching

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#### Upgrading Abstractions: Shuffles as RSS

Shuffle Abstraction

Partition traffic across cores



- For many UDFs can implement on NIC.
- Using receive side scaling (RSS).
- RSS can be used when shuffling by
  - TCP 5-tuple
  - Masked parts of the IP header.
- Currently implemented.
- Core 3 Significant performance benefits.

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Core 0 **UDF** dictates how traffic is Core 1 split. Core 2

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  - Relying on resource allocation policy to help with these questions.



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