

# Distributed Name Rewriting

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# Overview

What is the common abstraction behind “distributed Internet infrastructures”?

Hypothesis: **Distributed Name Rewriting**

This talk revisits (in 15 min): DNS, BTC, ARP, DHT, NFN, Tangle, GitHub ...

Other possible names (instead of DNR) from this morning:

- *distributed secure mappings (bind identities to keys), potentially trustless*
- *NFaaS (securely map inputs to computation results)*

# DNS vs BTC

DNS:

lookup(*in*) --> out

e.g. '*in*' is hierarchical domain name, '*out*' is IP number

BTC:

lookup(*in*) --> out

*'in'* is random (account or tx) number, *'out'* is an account balance, a transaction, a smart contract, an insurance policy etc

# DNS vs BTC (contd)

At the end of the day, DNS and BTC both are:

- small to midsize databases (BTC is a ledger, after all)
- global (distributed or replicated)
- simply query interface, maps one name to another

There are differences, of course ...

# DNS vs BTC - somehow different

	DNS	BTC
<b>versioning (history)</b>	no	yes (IOTA has snapshots and forgets)
<b>consistency</b>	eventually cons. (dependent on caching params)	strong cons. IFF you are on the winning branch
<b>input names (how to prevent conflicts)</b>	unique because pre-coordinated	unique because random
<b>decentralized storage</b>	yes (iterative/recursive remote query)	no (full replica)

one very big difference, though:

# DNS vs BTC - really different

## **b/c of the UPDATE method**

DNS: "pre-established agreement on delegation"

updates only possible in delegated subtree, are independent and can be done in parallel

BTC: trustless process

Byzantine Agreement Protocol for global, synchronous consensus

Is "distributed name rewriting" still a good common abstraction? I think yes.

# More Name-Rewriting Infrastructures

Seen so far: DNS, BlockChain (BTC), Tangle (IOTA)

**ARP** - dynamic mapping

**Forwarding** - routing table with next-hop lookup

**DHT** - an index, beside DNS the other “exemplary lookup” infrastructure

**PKI** - `secured_lookup(some_public_key) -> signing_key`

**cloud computation** - `lookup( fct(in) ) -> result`

Web pages are computation results: lookup results are cacheable, see memcached

**NFN** (Named-Function-Networking) — `resolve( symbolic_expr ) -> result`

**scalable!** immutable inputs, *confluence of resolution strategies avoids need for consensus finding*

Again: “update” is probably the strongest differentiator

# Communication is Computation is Distributed Name Rewriting is Communication is ...

Notation used: `A(something)` means: "something is on host A"

`config[...]` represents global state

Story: We want to replicate an item, send a unicast datagram from A to B via X

`config[ A(srcA,nameB,item), B(srcB) ]`

-> name rewriting due to DNS: map nameB to B's IP address

`config[ A(srcA,nameB,dstB,item), B(srcB) ]`

-> name rewriting due to route table lookup: map dstB to gwX

-> name rewriting due to ARP: map gwX IP name to eth name

`config[ A(srcA,nameB,item), lan1(ethX,srcA,dstB,item), B(srcB) ]`

-> delivery at gateway X

`config[ A(srcA,nameB,item), X(pkt(srcA,dstB,item)), B(srcB) ]`

-> name rewriting due to route table lookup: map dstB to dstB

-> name rewriting due to ARP: map dstB IP name to eth name

`config[ A(srcA,nameB,item), lan2(ethB,srcA,dstB,item), B(srcB) ]`

-> delivery at B

`config[ A(srcA,nameB,item), B(srcB, pkt(srcA,dstB,item)) ]`

-> delivery at application level:

`config[ A(srcA,nameB,item), B(srcB,item) ]`

# voila: the item was replicated through DNR

# Name-ReWriting Service, the API

Name-Rewriting as an Abstract Data Type (ADT), basically a key-value store

```
class NaRW: # a name rewriting service, its interface
    def get(): # also called "lookup", "resolve", "compute"
    def put(): # also called "update", "define", "undefine"
    def items(): # also called "walk", "listdir"
```

DNS, BTC, etc are then subclasses, type refinements, interface implementors.

Goal of this “ADT talk” is to abstract away from the implementation details, define the ADT by its properties, not the implementation

# Name-ReWriting Service, the API

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```
class NaRW: # a name rewriting service
```

**NaRWaaS**



DNS, ...

Goal of ...  
define t...

... details,

... implementation

# Name-ReWriting Service, the API

Name-Rewriting as an Abstract Data Type (ADT), basically a key-value store

```
class NaRW: # a name rewriting service
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**NaRWaaS**



DNS, ...

Goal of ...  
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**A potential DIN result: Name-ReWriting-as-a-Service spec**

# Implementing NaRW with two sub-services

Hypothesis: DIN will revisit these two services over and over

- a) **Persistent storage** to store a new item (`lookup(id) -> data`)  
take some CRUD database (create,read,update,delete), potentially append-only
- b) **head- (or “tip”) service** -- points to the most recent versions of an item

The rest is chaining items to other items via hash pointers (= items' intrinsic names)

Intuition:

- GitHub, Blockchain (fuses a and b), IOTA's tangle has multiple tips
- DNS has/is only head-service, ICN offers only storage ...

# The sweet spot for scalability and trustlessness ?

..... NaRW API (put, get, items)

**“Conflict-Free Replicated Data Types” (CRDT):**  
deterministic eventual  
**consistency** without  
consensus, hence  
**scalable**

—— *medium guarantees* ——

—— *strong guarantees* ——

**DNS** scales but:  
no history, trust-based,  
no auto-conflict resolution

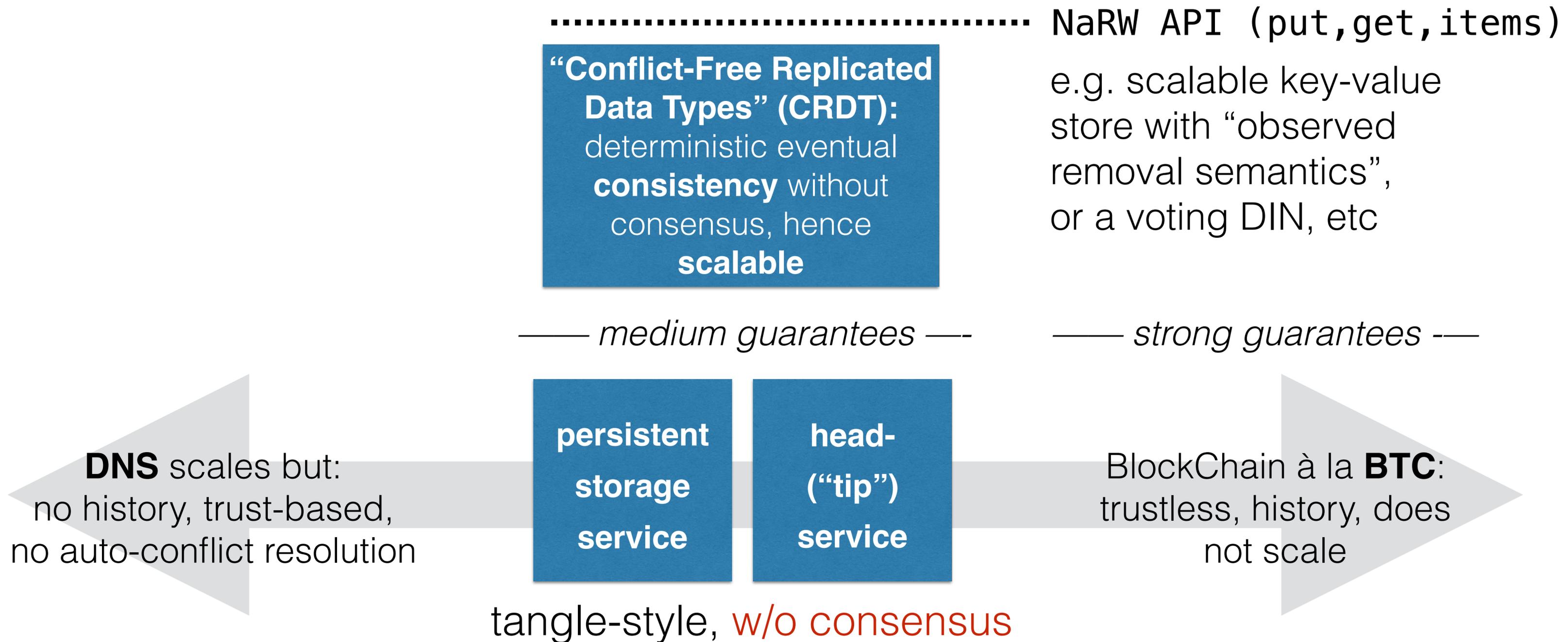
**persistent  
storage  
service**

**head-  
 (“tip”)  
service**

BlockChain à la **BTC**:  
trustless, history, does  
not scale

tangle-style, **w/o consensus**

# The sweet spot for scalability and trustlessness ?



# Questions

