

# HopAuth: Hop-by-Hop Authentication with Suspension Chain Model

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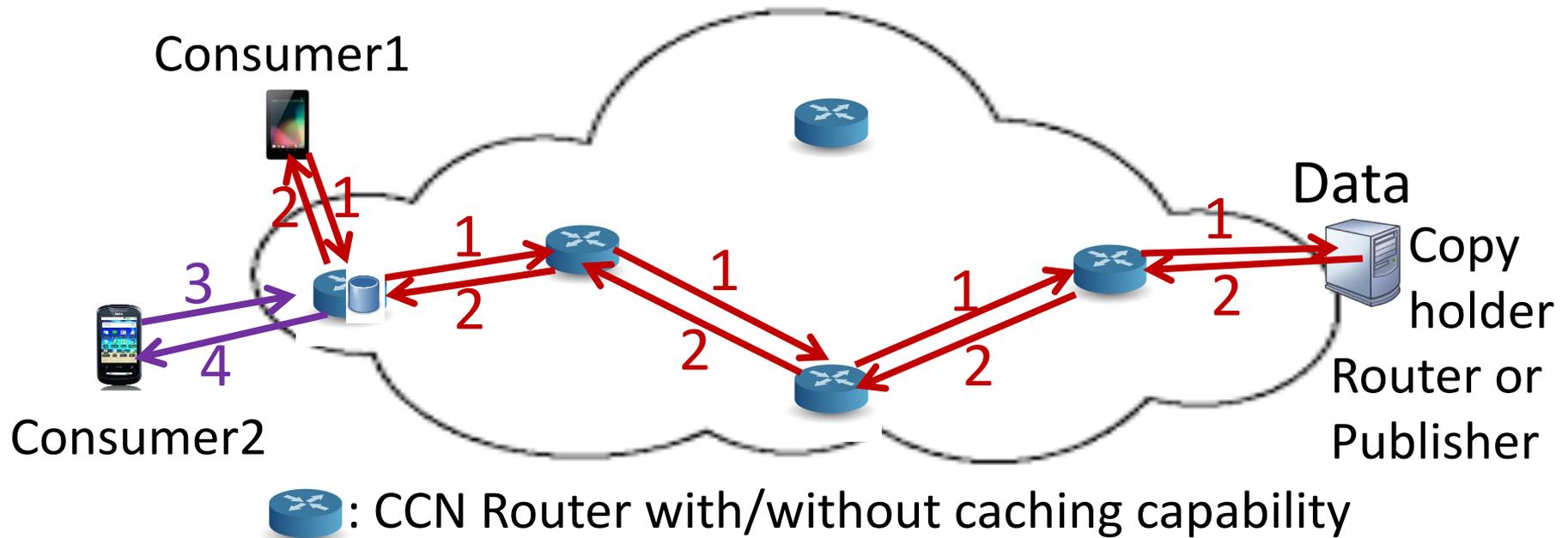
The scientific version of this proposal is published as;  
“DCAuth: Data-Centric Authentication for Secure In-Network Big-Data Retrieval”,  
*IEEE Transactions on Network Science and Engineering (TNSE)*, Sep. 2018.

# Outline

- Content-Centric Network
- Adversary Model
- Traditional Approach
- Proposed Hop-by-Hop Authentication Mechanism (HopAuth)
- Conclusions

# Content-Centric Network (CCN)

Packet Types: Interest/Data



**Publisher:** the entity that publishes data in network.

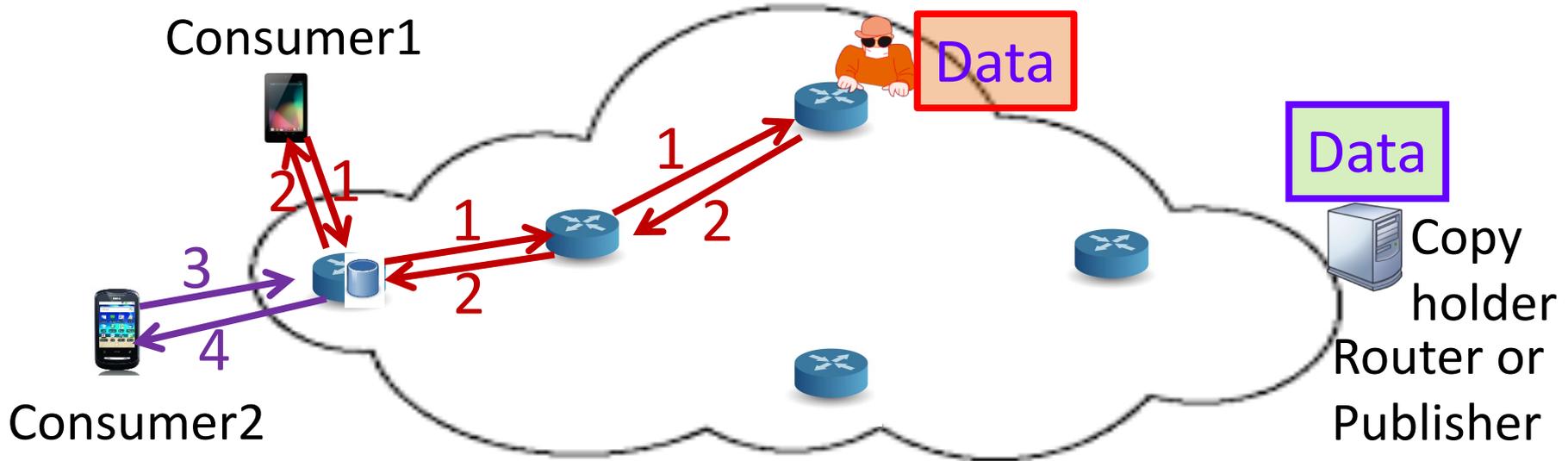
**Consumer:** the entity that retrieves data from network.

**Copyholder:** the entity that provides data to network. (Caching router or Publisher)

# Adversary Model

- **A1 (Content Poisoning Attack): Impersonate a Copy holder** to provide fake data
  - Currently, the content is only signed with the key of the entity who publishes it.
  - Impossible to verify the publishers' validity unless all routers use the authentication service of CA for all forwarded/cached data
- **A2 (Interest Flooding Attack): Impersonate a Consumer** to request data
  - Much existing work on restricting the Interest sending rate
  - Impossible to verify the consumer's validity unless all the Copyholders (Router or Publisher) use the authentication service of CA

# Content Poisoning Attack



Fake/corrupted data are cached along the path

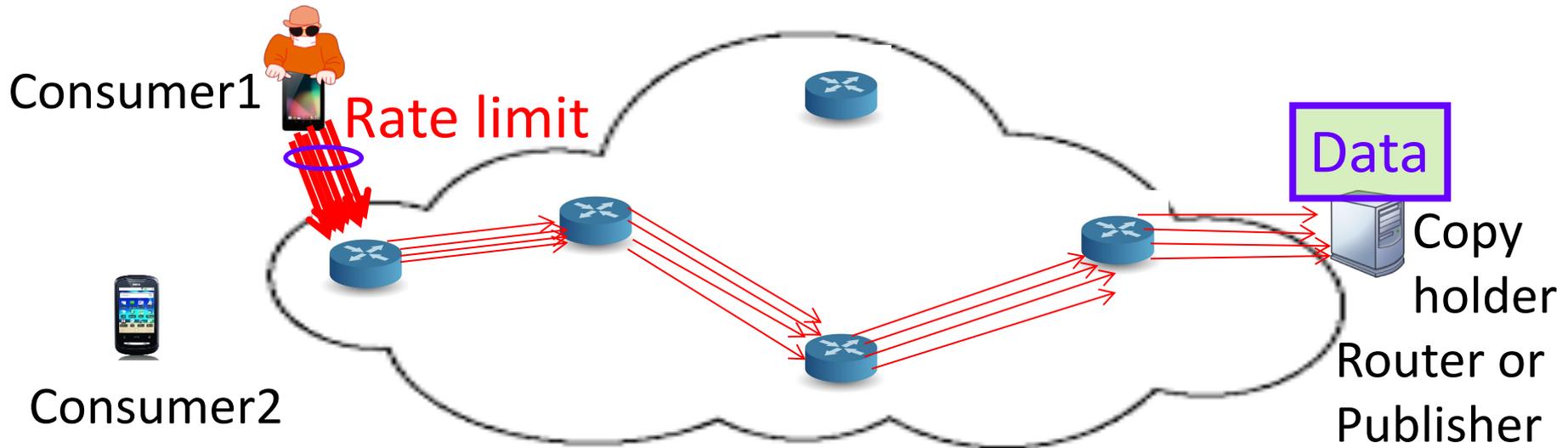
Consumers who use the path **always retrieve the wrong data**, because the router does not detect the cached data validity (as it's signed by attacker correctly)

Fake data are further cached, which **pollute the routers as virus spreads**.

Consumer 2 and other potential consumers will also retrieve the fake data from routers.

Routers need to **verify the data before caching**. Consumers verify copyholder and path to identify the polluted entities besides data verifications.

# Interest Flooding Attack



Consumer1 floods the Interests to the network to malfunction routers.



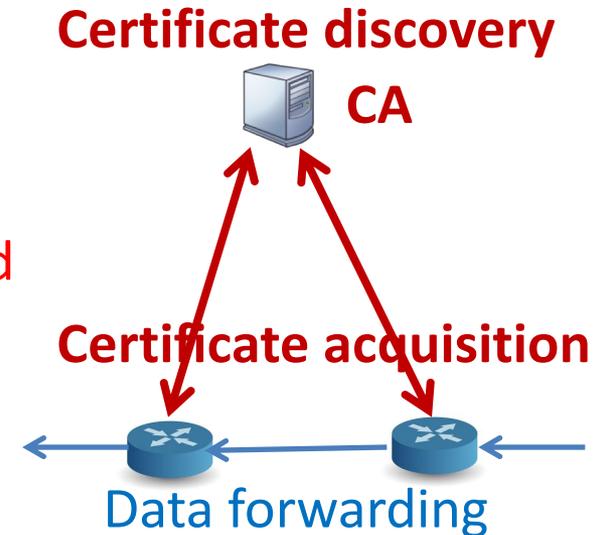
Even malicious Interests can be reduced by rate limiting, some malicious Interests still can reach the copyholder.



It will be effective to inhibit the Interest flooding attack, if routers **verify the Interests before replying** the data.

# Traditional Authentication Approach

- CA-based PKI (Certificate Authority-based Public Key Infrastructure)
  - Routers rely on centralized server to discover the required certificates.
  - Routers need interact with CA to discover and acquire the corresponding certificates using additional messaging for authentication, which leads it to be infeasible during data retrieval.



- Scalability Problem

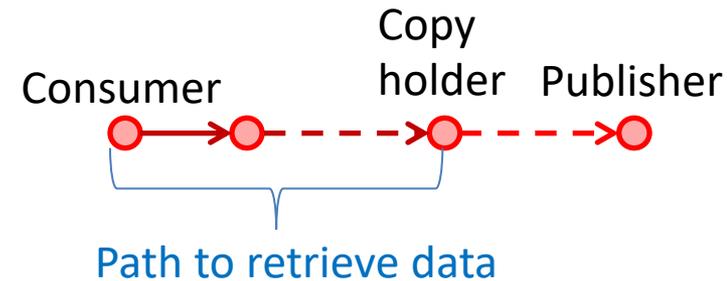
For router, it is highly cost if each router always needs to verify the signature of every data and every Interest.

For consumer, the public key of all potential copy holders should be provided to users beforehand.

# Hop-by-Hop Authentication Mechanism (HopAuth)

- **Key idea**

- HopAuth constructs a certificate chain **between copyholder and consumer along the path**.
- HopAuth can be CA-independent, yet also can partially use CA-based trust to **shorten the length of certificate chain** to form a suspension-chain based on the **web-of-trust** concept
- Packet forwarding with collection of certificate chain without extra messages
- **Extend the entity trust graph** with the application trust graph

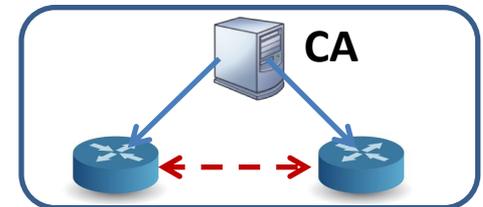


- **Effects**

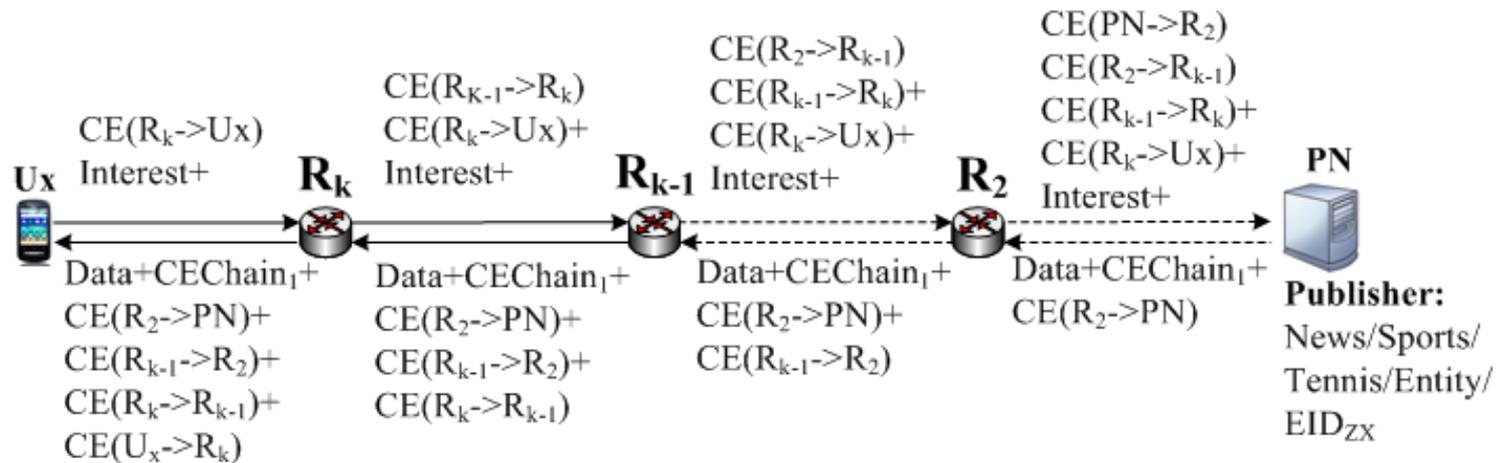
**Routers cache data only after authentication.** Authentication can be performed offline. (Currently caching data without authentication)  
Publisher authentication, copyholder authentication, consumer authentication, and path authentication can be achieved as required.

# Certificate Management

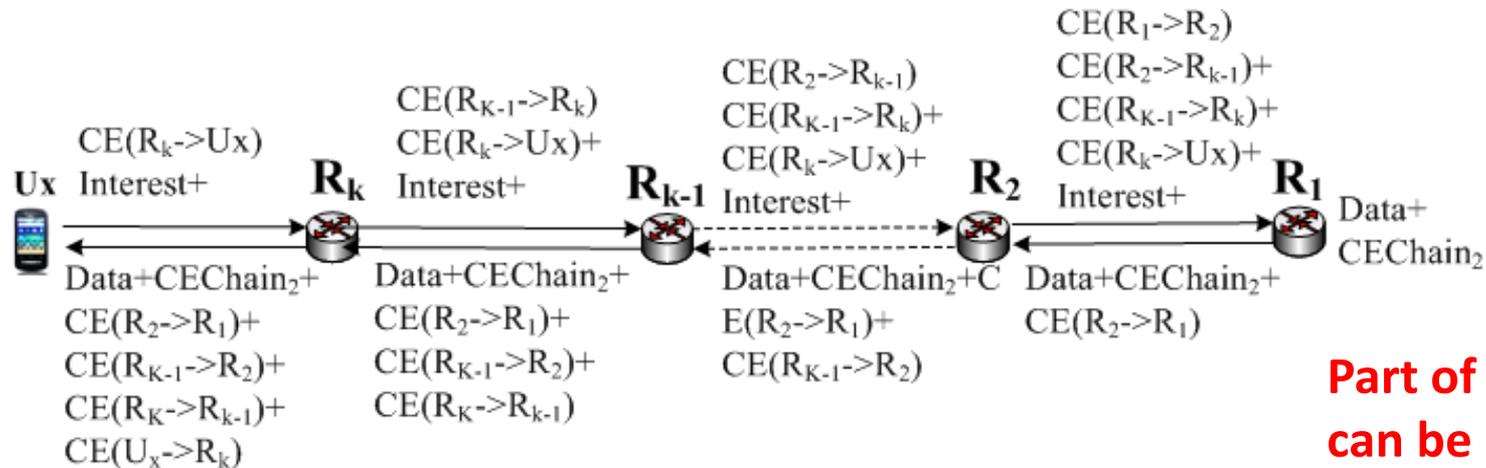
- **Physical entity:** the entity that communicates using a physical device. **Logical entity:** the entity that is involved in an application. (authorizer, sub-authorizer, publisher)
- In HopAuth, certificates are issued based on trust relationships, namely as **neighbor-based trust**, **CA-based trust** for physical entities, and **authorization relationships in applications** for logical entities.
- Physical Entity Trust
  - **Neighbor-based trust:** Routers issue certificates to each other in the neighborhood.
  - **CA-based trust:** CA issues certificates to highly trustable physical entities, which form highly trustable router group (HTRG).
- Logical Entity Trust
  - **Authorization relation-based trust:** If logical entity A authorizes the right for entity B to manage a sub-category or publish data, A should provide a certificate for the true public key of B.



# Forwarding-Integrated Authenticable Data Retrieval



(a) Data Retrieval from the Publisher



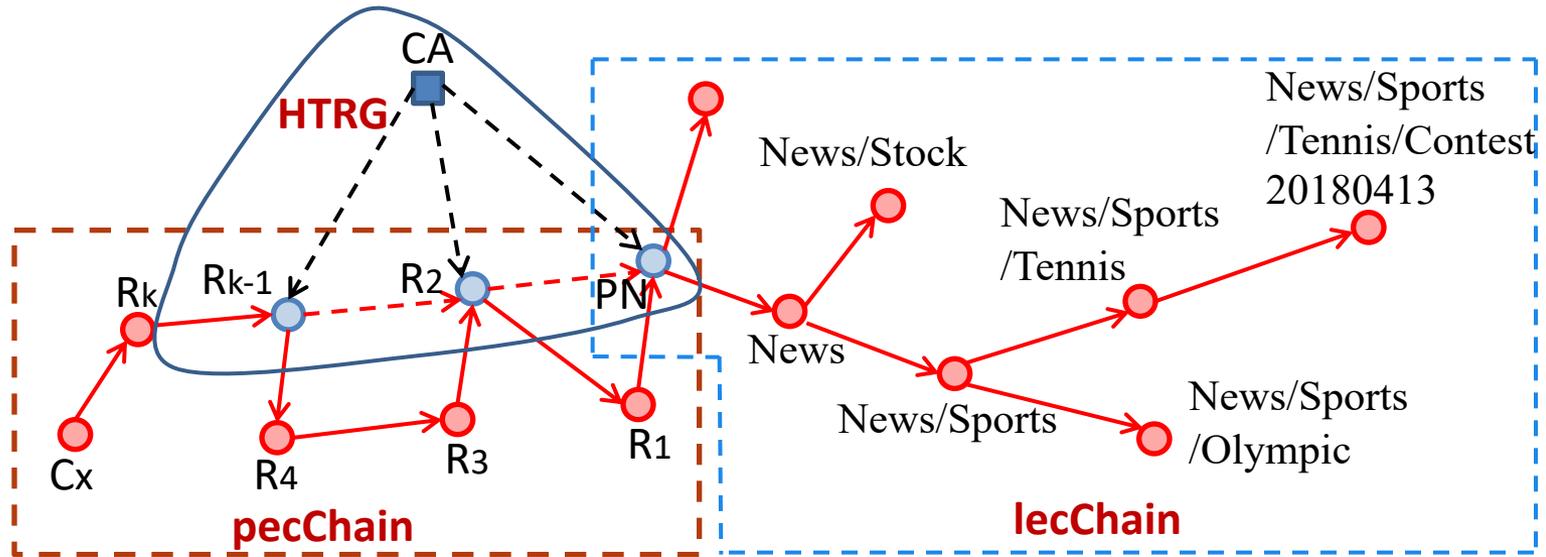
Part of CEChain can be cached.

(b) Data Retrieval from the intermediate physical entity

CEChain<sub>1</sub>: Certificate chain from PN to Publisher

CEChain<sub>2</sub>: Certificate chain from R<sub>1</sub> to PN and PN to Publisher

# Suspension Chain Model (SCM)



Cx: Consumer x    Ri: Router i    PN: Publisher Node    CA: Certificate Authority

→ : Neighbor-based trust relation    - - - - -> : CA-based trust relation in HTIG

→ : Suspension CA trust relation

**CA:** Certificate Authority                      **PN:** Publisher Node

**HTRG:** Highly Trustable Router Group

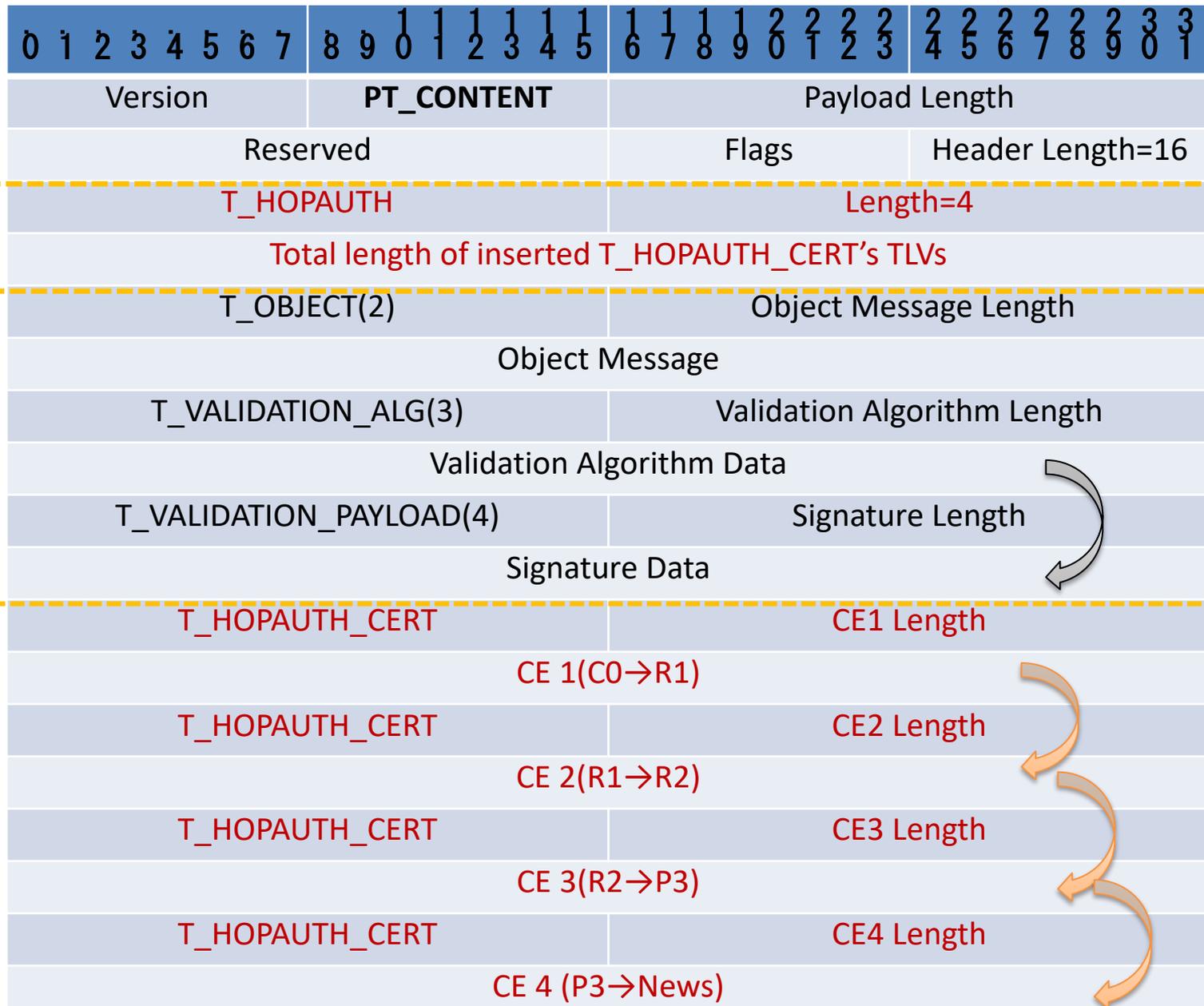
**pecChain:** Physical Entity Certificate Chain

**lecChain:** Logical Entity Certificate Chain

# Security Levels

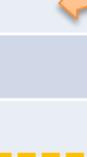
- We identify **three security levels** as follows.
  - **Security Level 1 (Low):** Publisher Authentication
    - Authenticate the entity who publishes data.
    - For one piece of data, the suspension chain will be constructed only once.
  - **Security Level 2 (Medium):** Copyholder Authentication
    - Publisher Authentication + Authenticate the router who provides data.
    - For each copyholder during data retrieval, the suspension chain will be constructed once.
  - **Security Level 3 (High):** Path Authentication
    - Copyholder Authentication + Authenticate the path from which data are retrieved.
    - For the first time data chunk retrieval from one path, the suspension chain will be constructed.
    - For the following data chunk retrieval from the same path, a chain of certificate name will be appended.

# (Potential) HopAuth Packet Format



Hop-by-hop hdr.

CEChain



# Certificate Size

- DTLS (Datagram Transport Layer Security in Constrained Environments) [REF.1]

- ECDSA P-256: **91 bytes**
- ECDSA P-384: **120 bytes**
- ECDSA P-521: **156 bytes**

For 2K bytes size data chunk, **546 bytes (26.7%)** are used to accommodate 6 certificates, if ECDSA P-256 is employed.

RSA Public Key Length (bit)	ECDSA Public Key Length (bit)
1024	160
2048	224
<b>3072</b>	<b>256</b>
7680	384
15360	512

[REF.1] Datagram Transport Layer Security in Constrained Environments  
<https://www.ietf.org/proceedings/83/slides/slides-83-lwig-2.pdf>

# Conclusions

- As an authentication protocol, we propose **HopAuth** to enable **hop-by-hop certificate collection and authentication**.
- As a trust model, we propose **SCM**, which merges **web-of-trust and partially CA-based trust** to construct a suspension chain to enable content-centric trust.
- The scientific version of this proposal is published as;
  - “DCAuth: Data-Centric Authentication for Secure In-Network Big-Data Retrieval”, *IEEE Transactions on Network Science and Engineering (TNSE)*, Sep. 2018.
- We will write an I-D for HopAuth if icnrg thinks it's useful / valuable.
- Any comment?

Thank you!