Key Update for OSCORE (KUDOS)

draft-ietf-core-oscore-key-update-02

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Recap

- > OSCORE (RFC8613) uses AEAD algorithms
 - Need to follow limits in number of encryptions and failed decryptions, before rekeying
 - Excessive use of the same key can enable breaking security properties of the AEAD algorithm*
- > (1) Key Update for OSCORE (KUDOS) ==> Today's main focus
 - Renew the Master Secret and Master Salt; derive new Sender/Recipient keys
 - No change to the ID Context; can achieve Perfect Forward Secrecy
 - Loosely inspired by Appendix B.2 of OSCORE
- > (2) AEAD Key Usage Limits in OSCORE
 - Defining appropriate limits for OSCORE, for a variety of algorithms
 - Defining counters for key usage; message processing details; steps when limits are reached

Rekeying procedure

- Key Update for OSCORE (KUDOS)
 - Client and server exchange nonces N1 and N2
 - UpdateCtx() function for deriving new OSCORE Security Context using the nonces
 - Extended OSCORE Option





// The actual key update process ends here.

// The two peers can use the new Security Context CTX_NEW.



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Open points for today

- > Flag bits in the OSCORE Option
 - First byte
 - Bit 'd' in the new second byte
- > Single method to update the key material
- > No runtime "negotiation" of FS mode or no-FS mode
- > Content about key usage limits
- > Learning KUDOS support through EDHOC EAD items
- > Where to define the update of OSCORE Sender/Recipient IDs

OSCORE flag bits

- Bit 15, namely 'd', has been registered
 If set to 1, it is a KUDOS message
- > Current situation: bits 0 and 1 are Reserved
 - Current text: define bit 1 for signaling a second flag bytes (as intended by RFC 8613)
- > Alternative approach discussed on the mailing list [1]
 - Define bit 0 for signaling a second flag byte
 - Change the status of bit 1 to "Unassigned" Possible to add to the already
 - No real plan for bit 0 otherwise --- Only old though som an end of the set COSE Object
 - Nice to have a consistent "extension pattern" through bits 0/8/16/24/...

> Ok with the alternative approach? If yes:

- Do Early Allocation of bit 0?
- Register bits 8/16/24/... already?

[1] <u>https://mailarchive.ietf.org/arch/msg/core/x_lx5a4PV-XcrvmLECtsC_CmoYs/</u>

0 1 2 3 4 5 6 7 8 9 10) 11 12 13	14 15	<> n bytes>			
0 <mark>1 0 h k n 0 0 0</mark>	0 0 0	0 d	Partial IV (if any)			
····						
<- 1 byte -> < s bytes	;> <- 1	. byte ->	< m + 1 bytes>			
s (if any) kid context (if any) x (if any)	nonce (if any)			
+	/		++ \			
/ / 01234567						
++ kid (if any)	+-+-+ 0 0	-+-+-+-+ blpl m	-+-+			
++	+-+-+	·-+-+-+-+	-+-+			

Single method for context update

- > Current method: updateCtx() has two internal paths for key update
 - One based on EDHOC-KeyUpdate() (Method 1)
 - > When EDHOC was used at first
 - One based on a HKDF Extract and Expand (Method 2)
 - > When EDHOC was not used at first
 - Method 1 implies that the EDHOC session is still valid
 - > Otherwise, need to dynamically fallback to Method 2
- > From IETF 114: then why not only Method 2?
 - No additional benefits from EDHOC-KeyUpdate
 - Building X_N becomes simpler
- > Proposed change: updateCtx() uses only Method 2

> Objections?

	1+ <the context="" edhoc="" established="" original="" security="" through="" was=""> { </the>
	// METHOD 1
	<pre>// Update the EDHOC key PRK_out, and use the // new one to update the EDHOC key PRK_exporter (new PRK_out, new PRK_exporter) = EDHOC-KeyUpdate(X_N)</pre>
	<pre>MSECRET_NEW = EDHOC-Exporter(0, h'', oscore_key_length)</pre>
	oscore_salt_length = < Size of CTX_IN.MasterSalt in bytes >
2	<pre>MSALT_NEW = EDHOC-Exporter(1, h'', oscore_salt_length)</pre>
	}
	else {
	// METHOD 2
	// MEIHOD 2 Label = "key update"
	<pre>// MEIHOD 2 Label = "key update" MSECRET_NEW = HKDF-Expand-Label(CTX_IN.MasterSecret, Label,</pre>
	<pre>// MEIHOD 2 Label = "key update" MSECRET_NEW = HKDF-Expand-Label(CTX_IN.MasterSecret, Label,</pre>
	<pre>// MEIHOD 2 Label = "key update" MSECRET_NEW = HKDF-Expand-Label(CTX_IN.MasterSecret, Label,</pre>

"Negotiation" of FS/no-FS mode

- > Mode currently signaled through the 'p' bit in the 'x' byte of the OSCORE Option
 - 'p' set to 0 ==> sender's wish to run KUDOS in FS mode (original mode)
 - 'p' set to 1 ==> sender's wish to run KUDOS in no-FS mode
 - If p = 0 in both KUDOS messages ==> use the FS mode
 - If p = 1 in both KUDOS messages ==> use the no-FS mode
- > If the initiator uses p = 0 and the responder uses p = 1
 - Abort KUDOS; from now on, the initiator uses p = 1
 - The initiator might not know the responder's capabilities from the start
- > Is the above possible, and thus an agreed fallback necessary? (issue #54)
- > Does an OSCORE Security Context also have information:
 - On the other peer's support for KUDOS? (answer: "maybe")
 - If yes, also on the other peer's support for the FS mode? (answer: "maybe")
 - If no, should it? That pre-knowledge may not be possible

012345678	9 10 11	12 13 14	15 < n	bytes>
+-+-+-+-+-+-+-+-+-+-	++	+	++	+
0 1 0 h k n 0	0 0 0 0	0 0 0	d Partial	. IV (if any)
+-+-+-+-+-+-+-+-+-	++	++	++	+
<- 1 byte -> < s	bytes	> <- 1 by	te -> < m	+ 1 bytes>
s (if any) kid cor	text (if an	y) x (if	any) nonce	(if any)
+		/	\	
		/		
		/ 0123	4567	
++		+-+-+-	*-+-+-+ I	
kid (i† any)		0 0 b p	m	
++		+-+-+-	+-+-+-+	

Split out update of OSCORE IDs?

- > Defined method for updating the peers' OSCORE Sender/Recipient IDs
 - Based on earlier discussions on the mailing list [1][2] and on [3]
 - This procedure can be embedded in a KUDOS execution or run standalone
 - This procedure can be initiated by a client or by a server
- > Properties
 - The sender indicates its new wished Recipient ID in the new Recipient-ID Option (class E)
 - Both peers have to opt-in and agree in order for the IDs to be updated
 - Changing IDs practically triggers derivation of new OSCORE Security Context

> From IETF 114: split out as a separate draft?

- This is strictly related to OSCORE, but ...
- ... not strictly related to KUDOS functionality
- Thus the KUDOS draft can focus on KUDOS!

	No.	+ C	+ U	+ N	+ R	Name	Format	Length	Default
	TBD1				 	Recipient-ID	 opaque 	0-7	(none)
+	+++++++++++								

[1] https://mailarchive.ietf.org/arch/msg/core/GXsKO4wKdt3RTZnQZxOzRdIG9QI/

[2] https://mailarchive.ietf.org/arch/msg/core/ClwcSF0BUVxDas8BpgT0WY1yQrY/

[3] https://github.com/core-wg/oscore/issues/263#issue-946989659

Signal KUDOS support in EDHOC

- > We can register EDHOC an EAD item for signaling KUDOS support
 - A peer learns if the other peer supports KUDOS (and which modes) during EDHOC execution
- > Possible semantics:
 - Value 1 -> "Tell me about what you support"
 - Value 2 -> "I do not support KUDOS"
 - Value 3 -> "I support KUDOS in both modes; tell me about you if you haven't already"
 - Value 4 -> "I support KUDOS only in no-FS mode; tell me about you if you haven't already"

Should we do it? Comments?

Relocate content related to limits?

- > Current structure of Section 2
 - Section 2.1 Overview of key usage limits; specific values to follow --- This builds on [1]
 - Section 2.2 Extensions of the OSCORE Security Context
 - > 'exp' in the Common Context; limits and counters in Sender/Recipient Context
 - Section 2.3 Extensions of the OSCORE message processing
 - > On incrementing the counters and when stopping using the current keys

> How to proceed?

- 1. Keep as is
- 2. Move content to an Appendix
 - > 2.1 ==> Appendix A
 - > Appendix A ==> Appendix A.1

3. Move content to a new draft

> The whole Section 2? Only part of it?

It was agreed to elaborate on limits and to have all this content in this same document [2]

[1] <u>https://datatracker.ietf.org/doc/draft-irtf-cfrg-aead-limits/</u>

2] https://datatracker.ietf.org/doc/minutes-interim-2021-core-04-202104281600/

Thank you!

Comments/questions?

https://github.com/core-wg/oscore-key-update