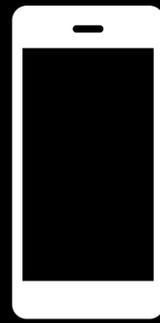


# Architecture

Chris Wood & Jana Iyengar

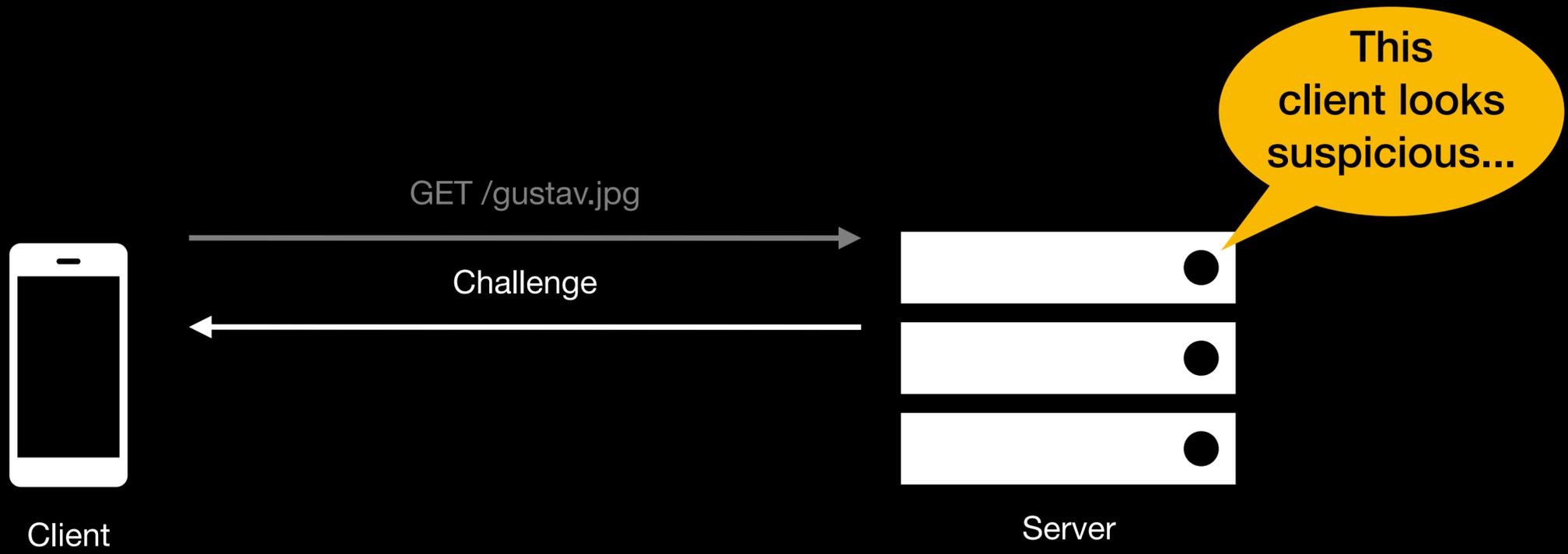


Client

GET /gustav.jpg



Server



Generate  
proof based on this  
challenge

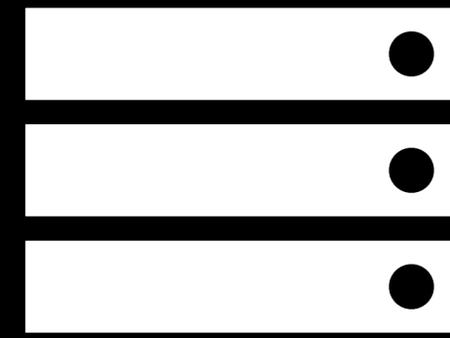


Client

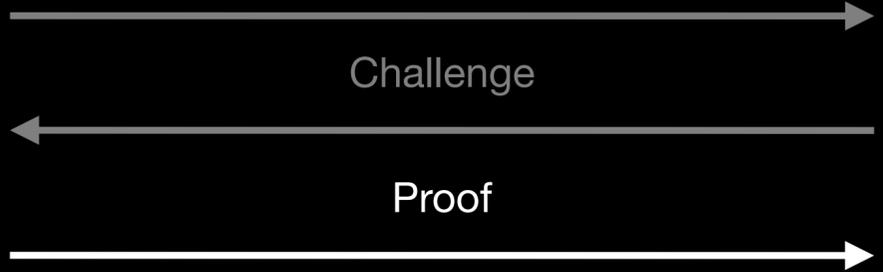
GET /gustav.jpg

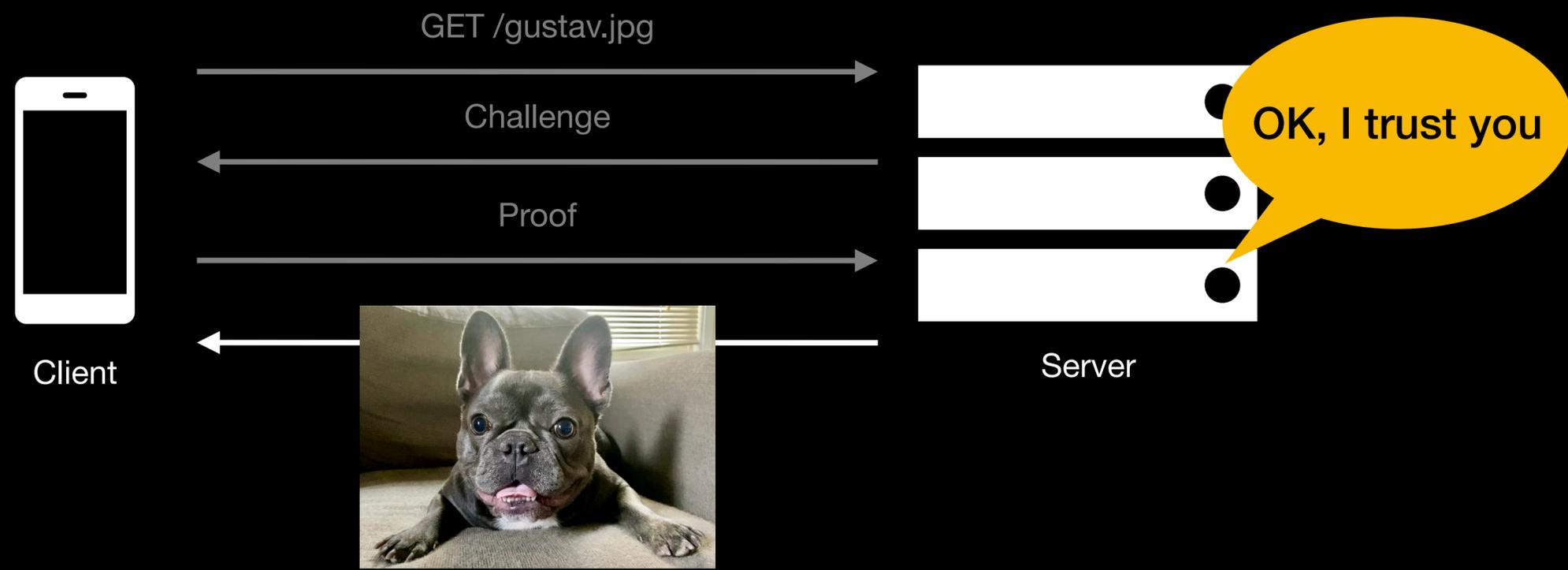
Challenge

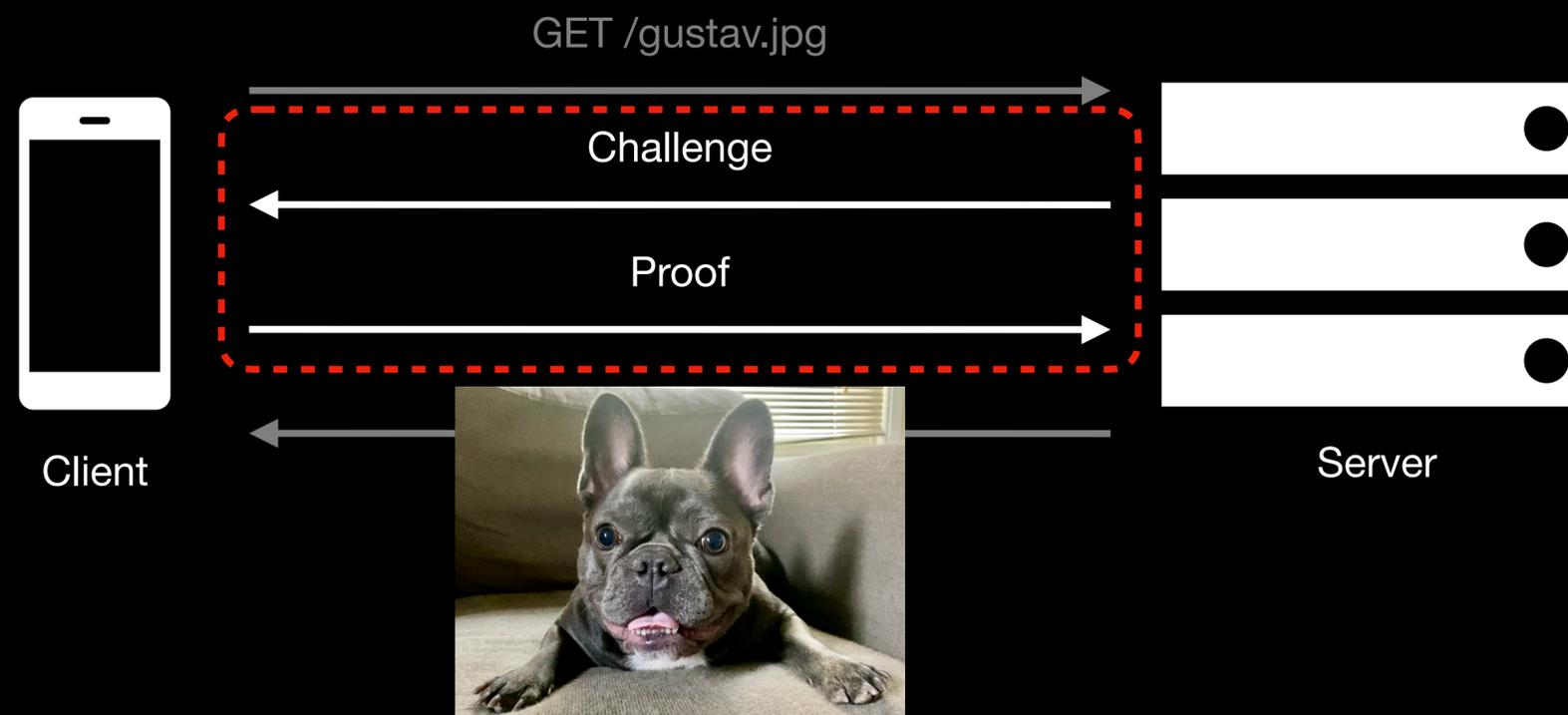
Proof



Server

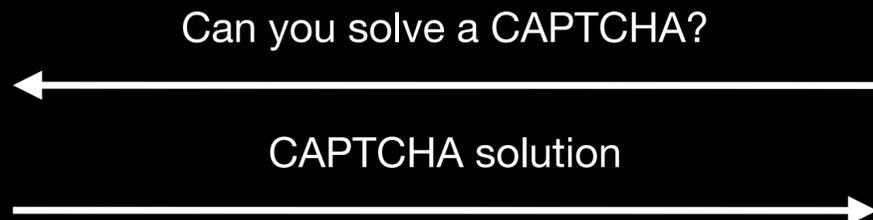






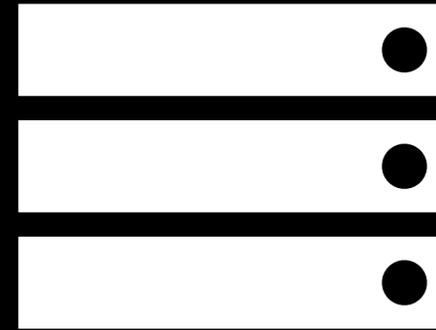


Client



Can you solve a CAPTCHA?

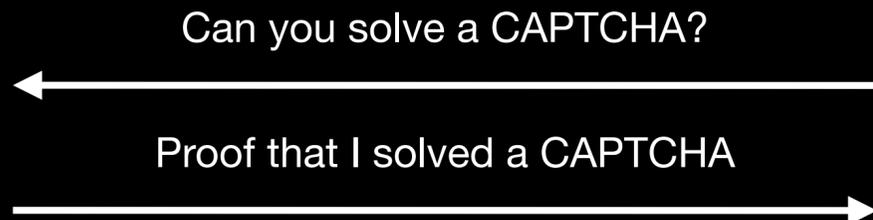
CAPTCHA solution



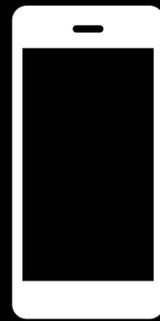
Server



Client



Server



Client

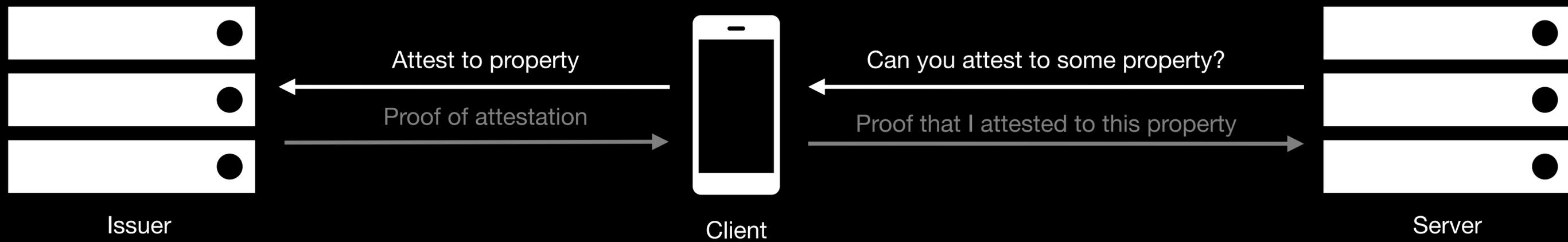
Can you attest to some property?

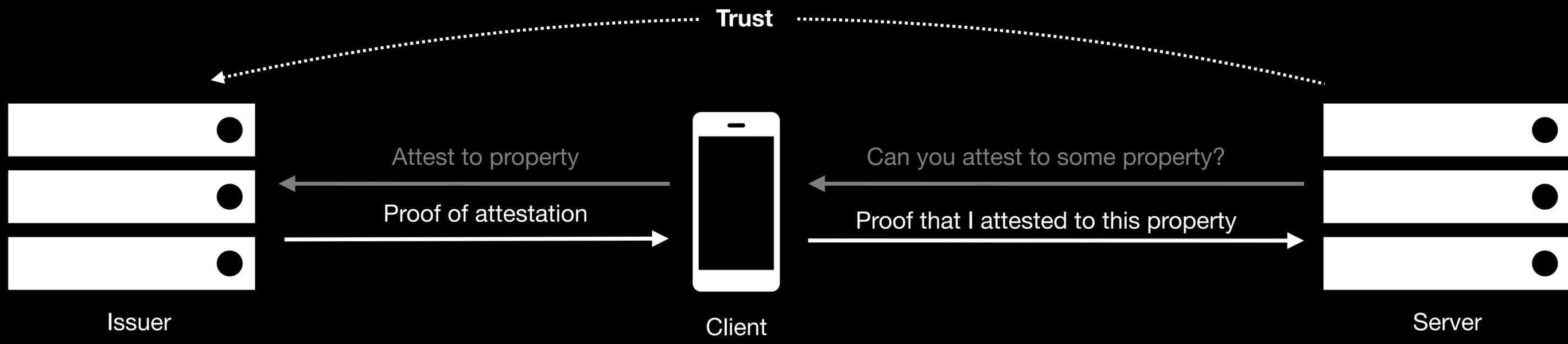


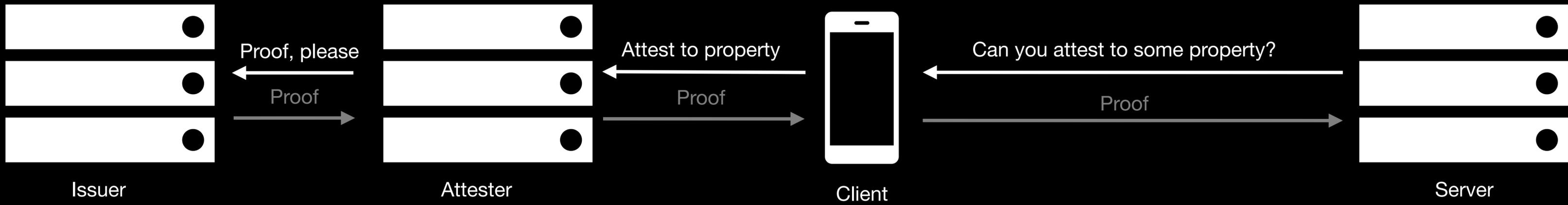
Proof that I attested to this property

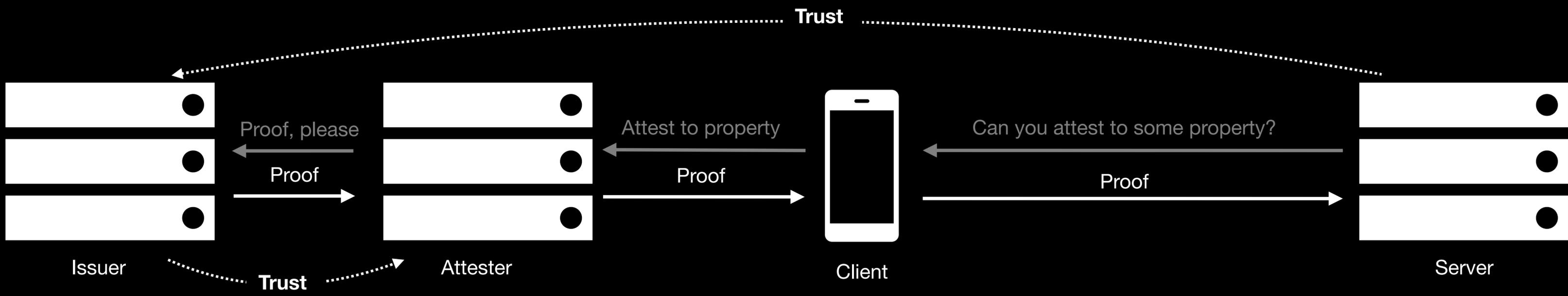


Server

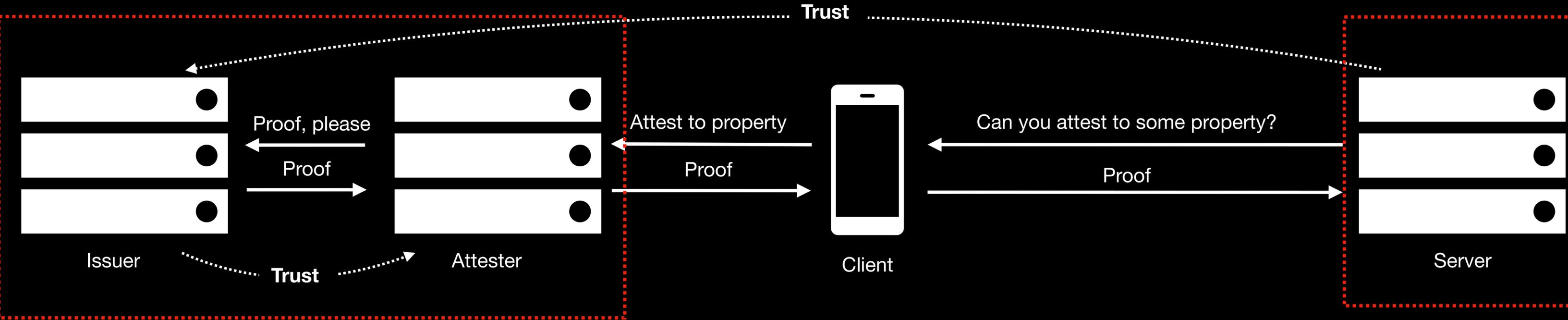








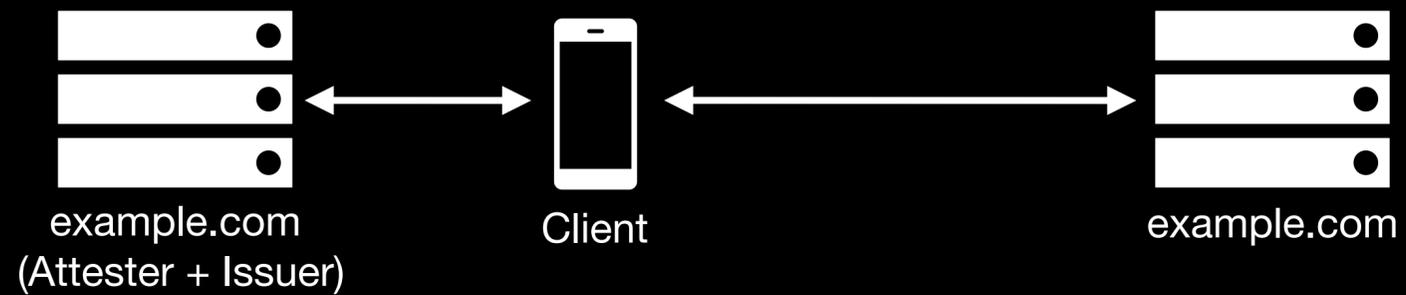
# Privacy Pass Architecture



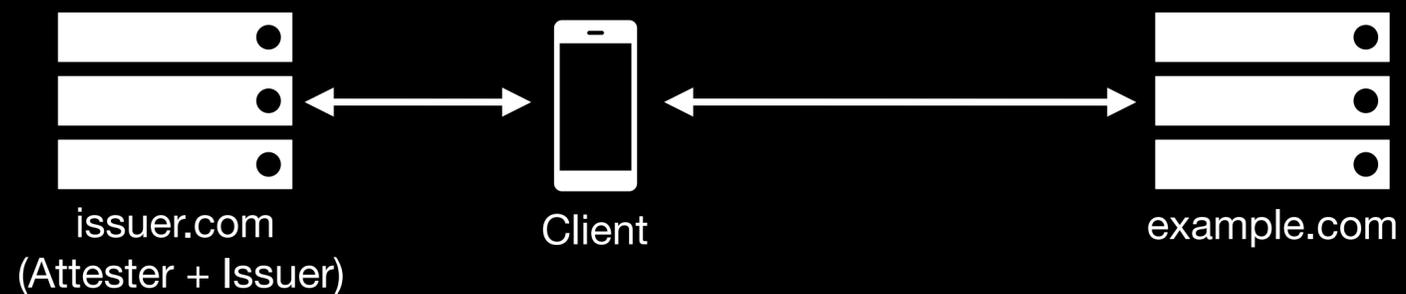
# Deployment Variations

The architecture can be instantiated in various ways

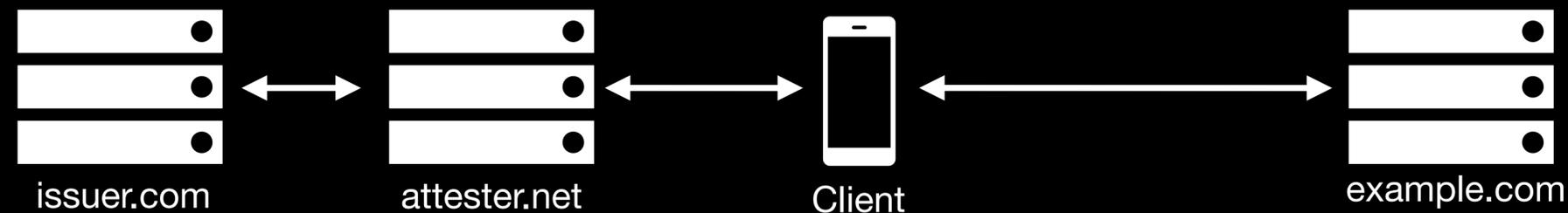
*Combined origin/attester/issuer (“single verifier”)*



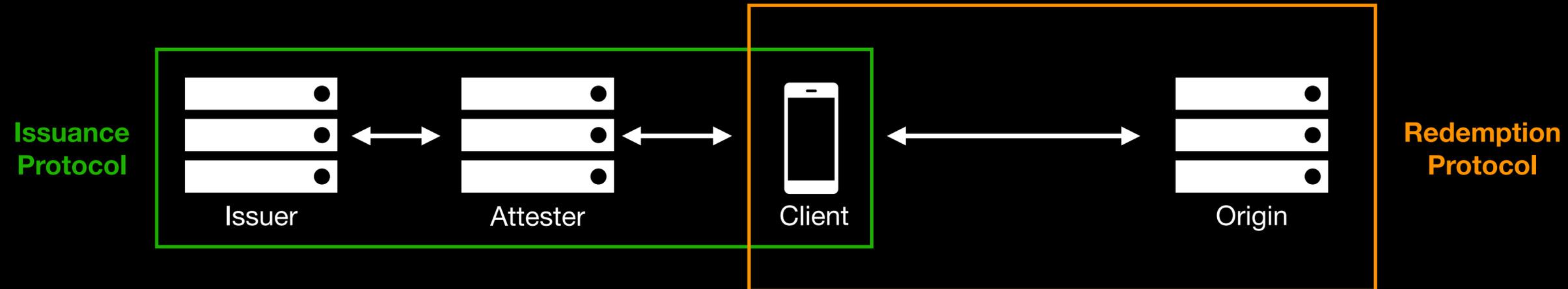
*Combined attester/issuer*



*All separate*



# Protocol Structure



Architecture describes two parts of the protocol, which are detailed in two separate documents:

**Redemption** is a consistent/unified API for redeeming tokens, along with the ability to challenge.

**Issuance** can support multiple types (VOPRF, publicly verifiable, etc). This is the exchange that can be extended or replaced for new deployment models.

# Big Picture Architecture

Some function *attests* to certain state or properties associated with a client

- Has this person solved a CAPTCHA?
- Does this person have a subscriber account?

Issuers that trust attesters produce proof -- tokens -- bound to these properties

Redeemers, or **origins**, consume tokens from trusted issuers

# Why rework the architecture?

Current architecture tightly couples *issuance* and *redemption*

Issuer and redeemer may be the same (as in Privacy Pass) but don't need to be

Separate roles allow for new deployment models and are more compatible with features like public verifiability

New architecture separates these functions and *shifts extensibility to issuance*

New extensions or features can be solved by new issuance protocols

Redemption is unchanged

Makes *attestation* explicit, but deployment specific

# Proposal

Define architecture in terms of functional roles (Client, Origin, Attester, Issuer)

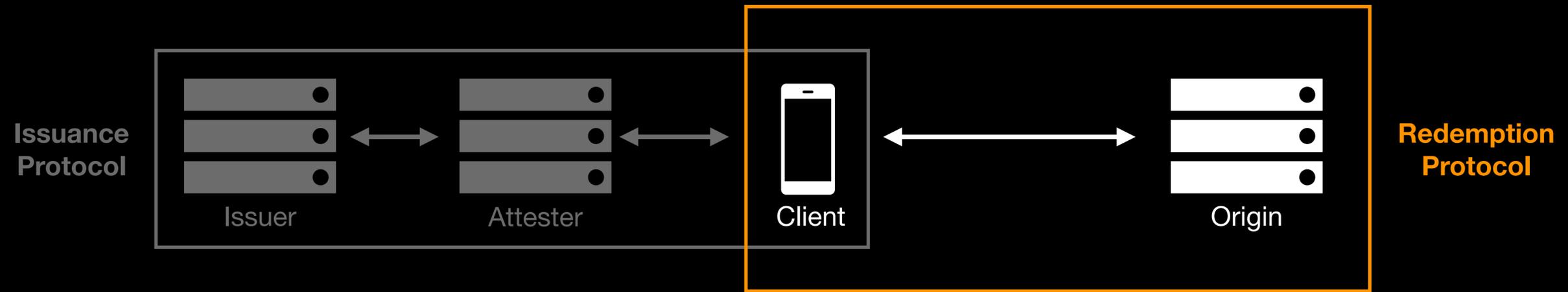
Define protocols in terms of Redemption and Issuance

Merge PR into architecture document

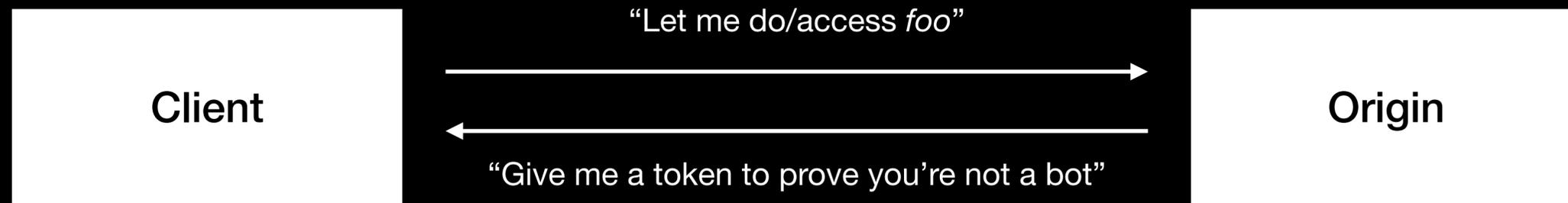
<https://github.com/ietf-wg-privacypass/base-drafts/pull/86>

# Challenge & Redemption

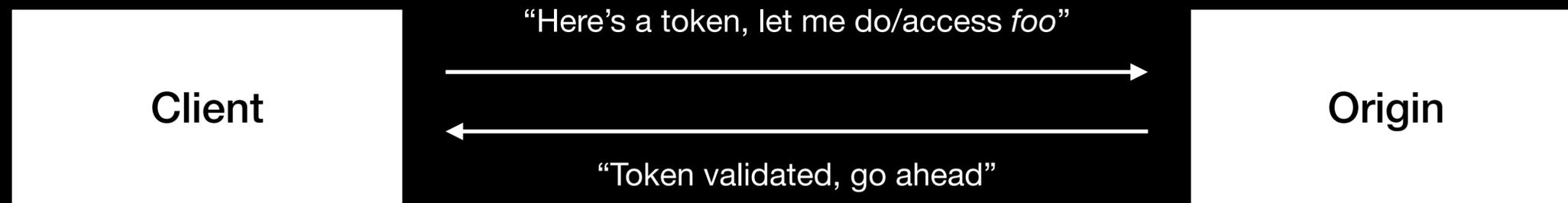
Tommy Pauly



# Challenge



# Redemption



# Challenge & Redemption

All token schemes involve token redemption

Token redemption is when a client presents a token to gain access, anonymously

Challenges are optional

Allows a server to indicate that it needs tokens

Indicates types of tokens and token issuers that are trusted

Allows for interactive tokens

# What was missing?

Previous design required Javascript APIs (W3C) work to functionally drive token interactions

No clear way to support new token types (POPRF vs publicly verifiable, etc)

HTTP authentication method allows a more standard definition

- Explicit support for different types of tokens, defined in their own contexts

- Works both in Javascript (W3C) and non-Javascript contexts

Authors proposing that this work replaces the HTTP API document

# Features

Define an IANA registry of token types, indicate in challenge & redemption

Indicate Issuer name(s) (who does the Origin trust to vend tokens?)

Allow for “interactive tokens” with a one-time nonce to prevent farming

Allow for binding tokens to an origin to prevent cross-origin spending

# Origin considerations

Make it easy for origins to adopt!

Origins don't need to do complex crypto, just need to verify

- Publicly verifiable types are simple (RSA signatures)

- Privately verifiable requires Issuer key (or a single HTTP request to the issuer)

Interactive tokens mitigate concerns about farming and double-spending

- Shifts server state from *redeemed tokens* (unbounded) to *number of outstanding challenges* (bounded by active sessions)

# Challenge

```
WWW-Authenticate: PrivateToken challenge=abc..., token-  
key=123...
```

```
struct {  
    uint16_t token_type; // Defines Issuance protocol  
    opaque issuer_name<1..2^16-1>;  
    opaque redemption_nonce<0..32>; // Optional  
    opaque origin_name<0..2^16-1>; // Optional  
} TokenChallenge;
```

Redemption nonce: If present, token presented must be *fresh* (interactively minted)

Origin name: If present, token is restricted to the origin, else it's *cross-origin*

# Redemption

Authorization: PrivateToken token=abc...

```
struct {
    uint16_t token_type; // Matches challenge
    uint8_t nonce[32]; // Client-generated nonce
    uint8_t context[32]; // Hash of TokenChallenge
    uint8_t token_key_id[Nid];
    uint8_t authenticator[Nk]; // From Issuance protocol
} Token;
```

Context: SHA256 hash of the corresponding challenge

Authenticator: Signature, POPRF output, etc

# Redemption Properties

## Security properties

**Redemption unlinkability:** Redeemer cannot link two tokens to the same client ✓

# Proposal

Replace HTTP API document with this HTTP auth scheme

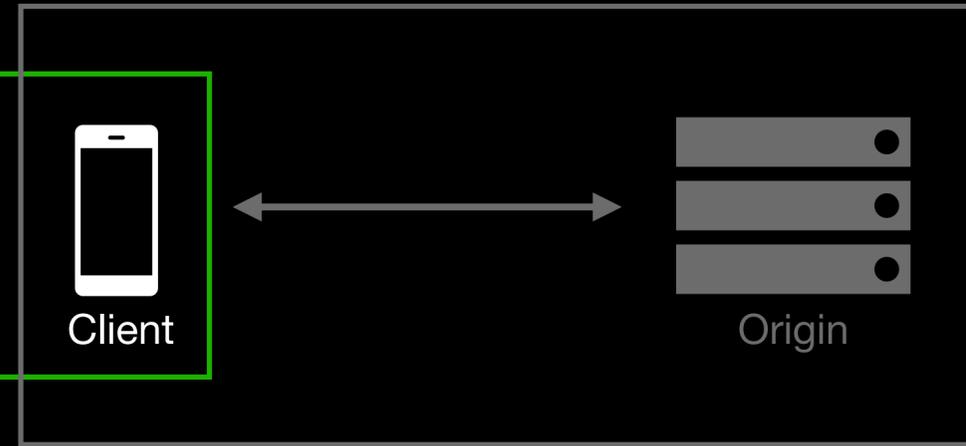
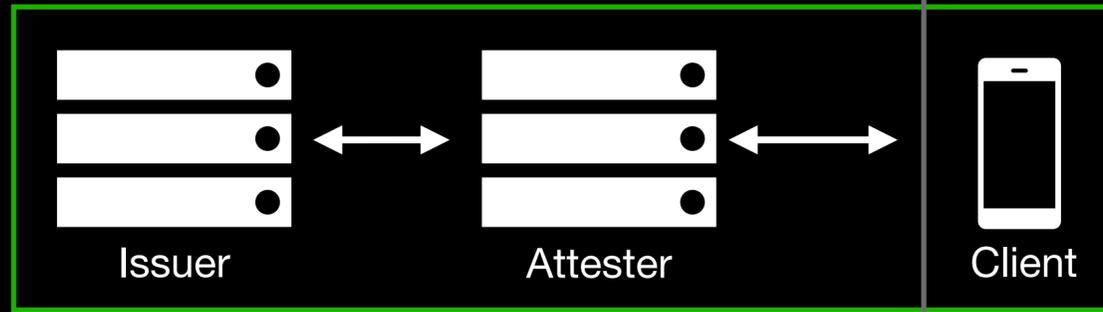
HTTP interactions with Issuers go to the Issuance Protocol document

Update W3C APIs to drive this HTTP API

# Issuance

Chris Wood

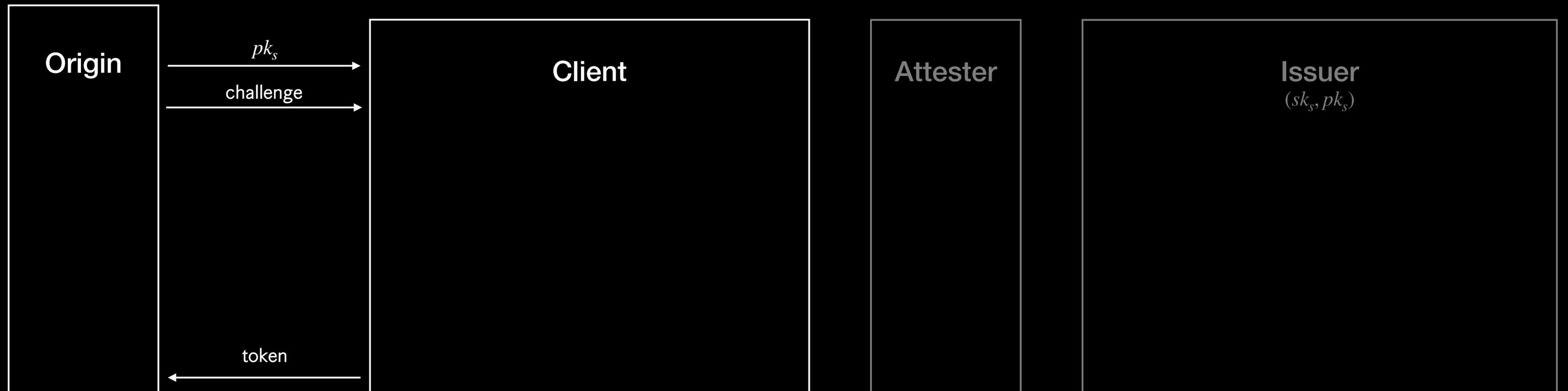
Issuance  
Protocol



Redemption  
Protocol

# Basic Tokens

## Issuance protocols



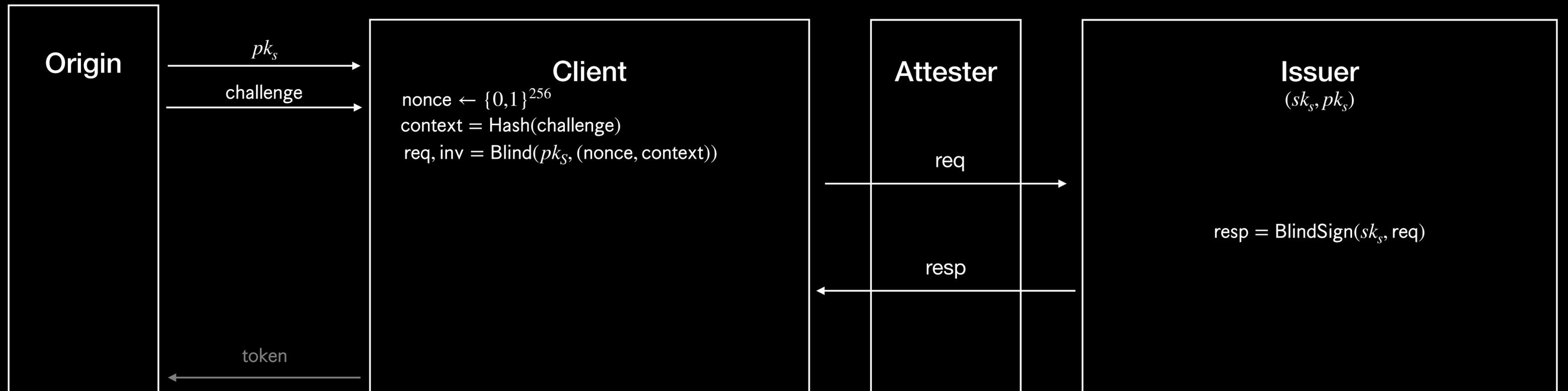
# Basic Tokens

## Issuance protocols



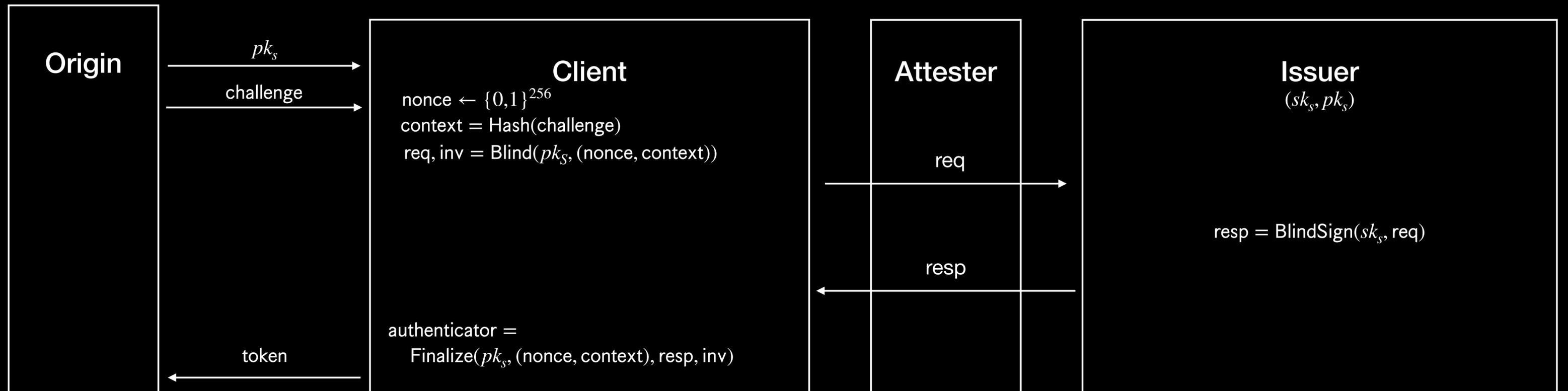
# Basic Tokens

## Issuance protocols



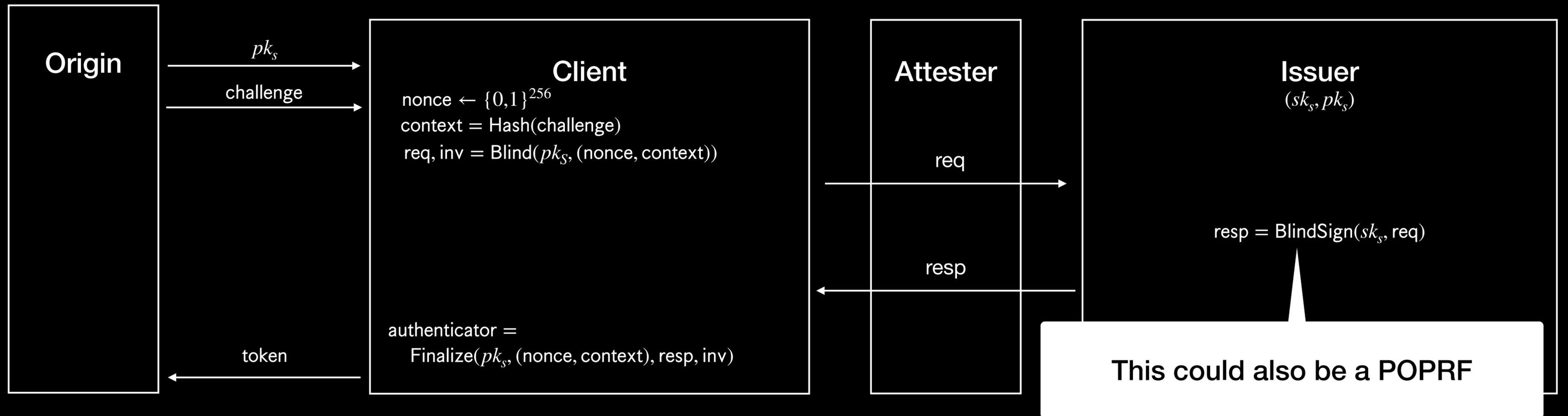
# Basic Tokens

## Issuance protocols



# Basic Tokens

## Issuance protocols



# Issuance Registry

## Issuance protocols

Value	Name	Publicly Verifiable	Public Metadata	Private Metadata	Authenticator Size
0x0001	POPRF(P-384, SHA-384)	N	Y	N	48
0x0002	Blind RSA (4096)	Y	N	N	512
...	...	...	...	...	...

Extensions for consideration:

Anonymous tokens with private metadata bit (<https://eprint.iacr.org/2020/072>)

Publicly verifiable anonymous tokens with private metadata bit (<https://eprint.iacr.org/2022/004>)

Blind BLS (<https://datatracker.ietf.org/doc/draft-irtf-cfrg-bls-signature/>)

# Issuance Considerations

## Issuance protocols

Issuance protocol is assumed to be stateless on the *Issuer*

Blind signature protocols that require multiple rounds and state are possible, but not specified

Compatible with deployment specific key consistency mechanisms

Issuer keys are discoverable such that applications can build consistency systems on top

# Issuance Properties

## Security properties

**One-more unforgeability:** Clients cannot forge tokens ✓

**Issuance secrecy:** Issuing parties cannot link per-client and per-origin state ✓

# Proposal

Replace existing protocol document with new issuance protocol details

Integrates with HTTP-based redemption protocol

Satisfies private and public verifiability (per the charter)

Makes issuance flow in the protocol document explicit and interoperable

# Questions for the WG

## Wrapping up

1. Are the document proposals clear?
2. Is there consensus in this new direction, which includes:
  1. Updates to *draft-ietf-privacypass-architecture* and *draft-ietf-privacypass-protocol*
  2. Adoption of *draft-pauly-privacypass-auth-scheme*

# Rate-Limited Issuance

Chris Wood

# Rate-Limited Tokens

## Issuance protocols

Rate-limited tokens *extend* the basic issuance protocol with new properties:

1. Issuers learn origin associated with a token challenge
2. Attesters learn *stable mapping* between per-client secret and per-origin secret, and no per-origin information
3. Token requests may fail if the per-origin rate limit is exceeded

Challenge: How to reveal only the origin to issuer, and only the mapping to attester?

# Rate-Limited Tokens

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# Detour: Stable Mappings

## Issuance protocols

A stable mapping is a deterministic function between per-client and per-origin information, e.g.,  $F(\text{client secret}, \text{origin secret})$

The mapping is used to enforce per-origin limits

Attester uses mapping as index into data structure tracking per-client state

Mapping	Count
....	...
<b>12311235123</b>	N
....	...

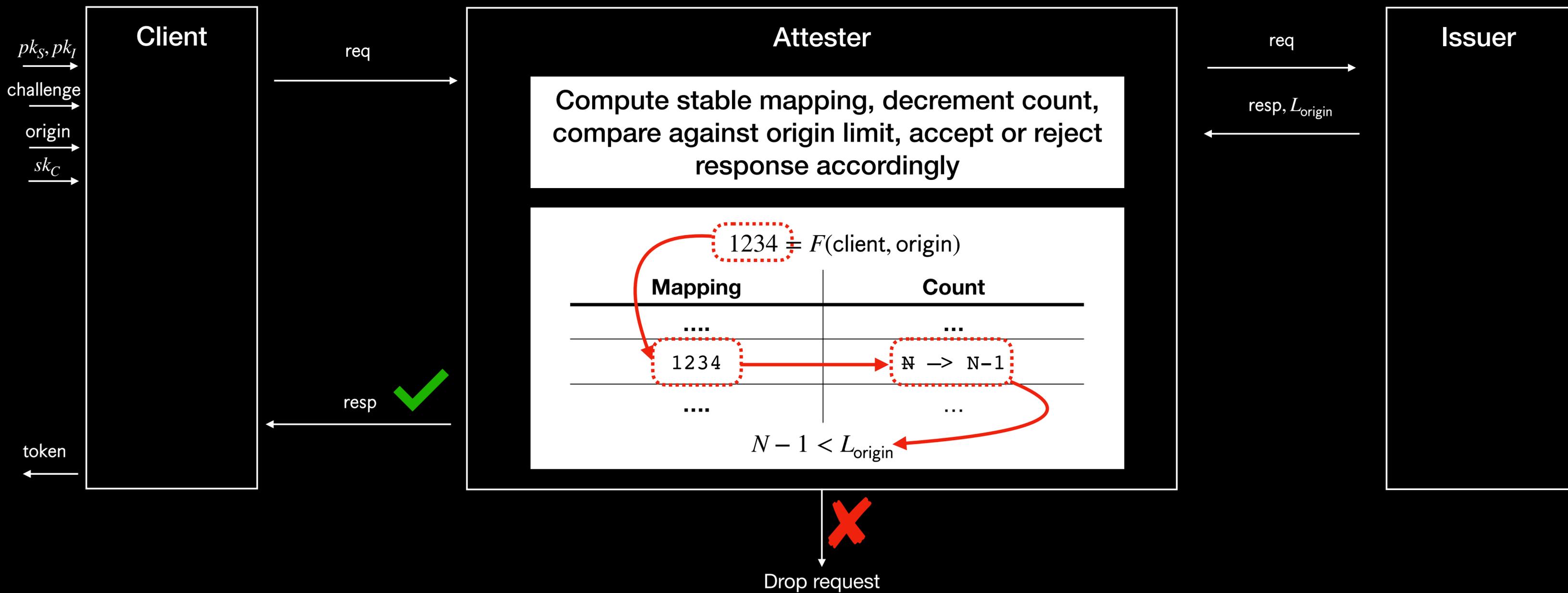
# Detour: Stable Mappings

## Issuance protocols



# Detour: Stable Mappings

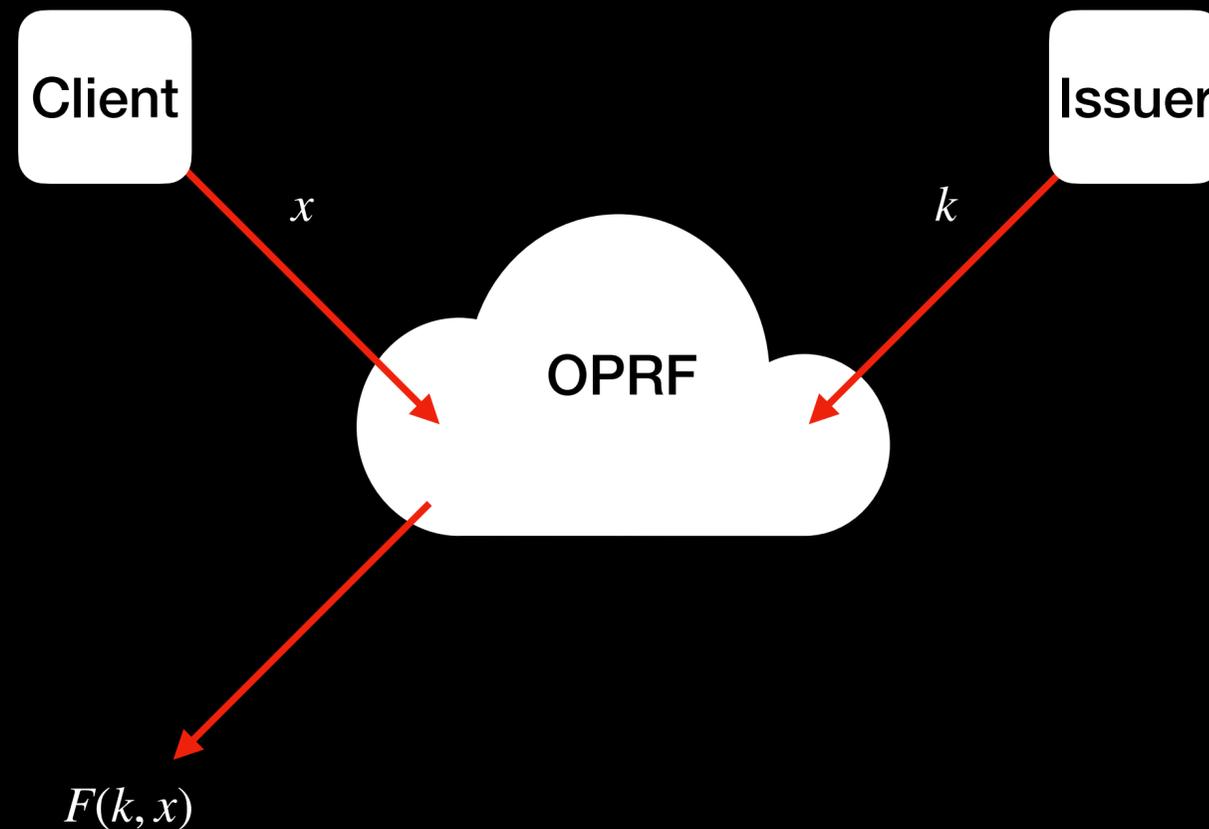
## Issuance protocols



# Detour: An OPRF Sketch

## Issuance protocols

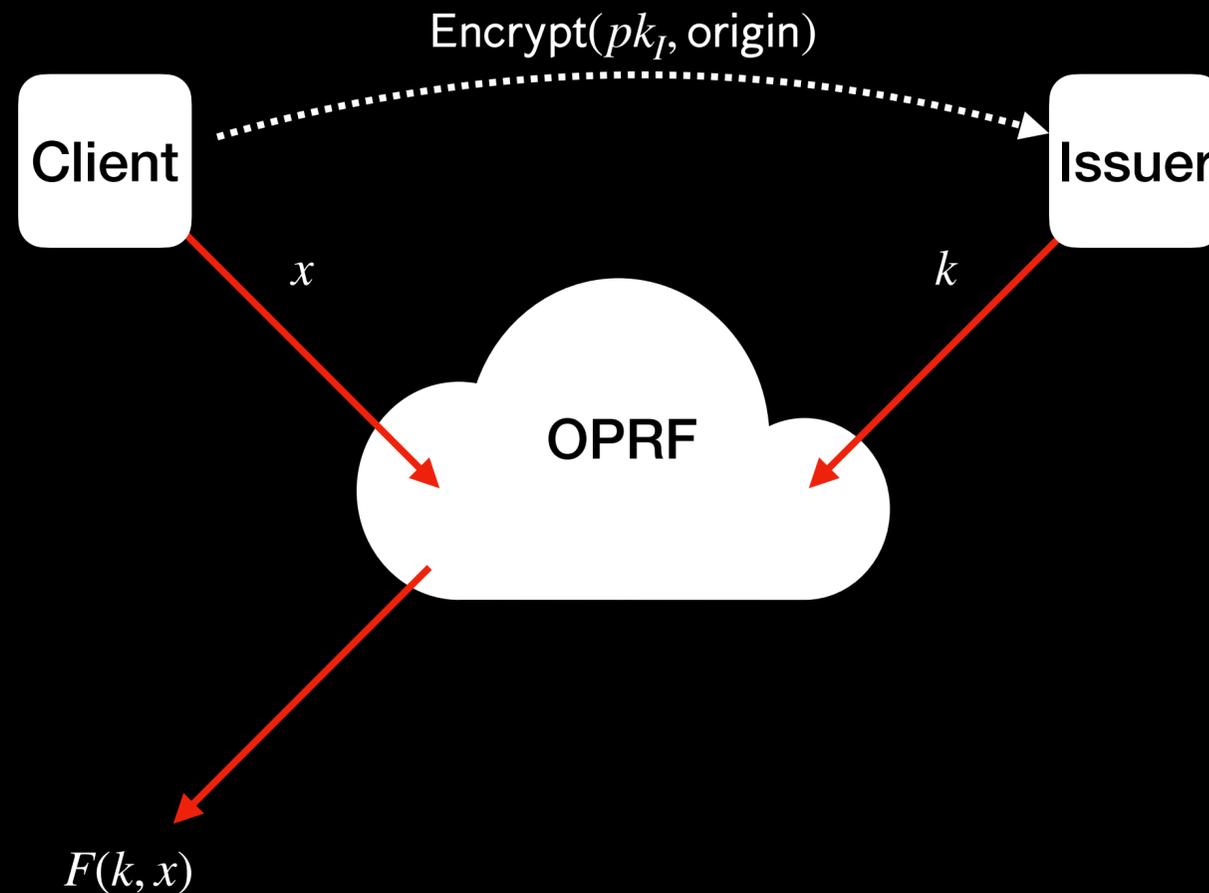
An OPRF protocol computes  $F(k, x)$  for per-origin  $k$  and per-client  $x$



# Detour: An OPRF Sketch

## Issuance protocols

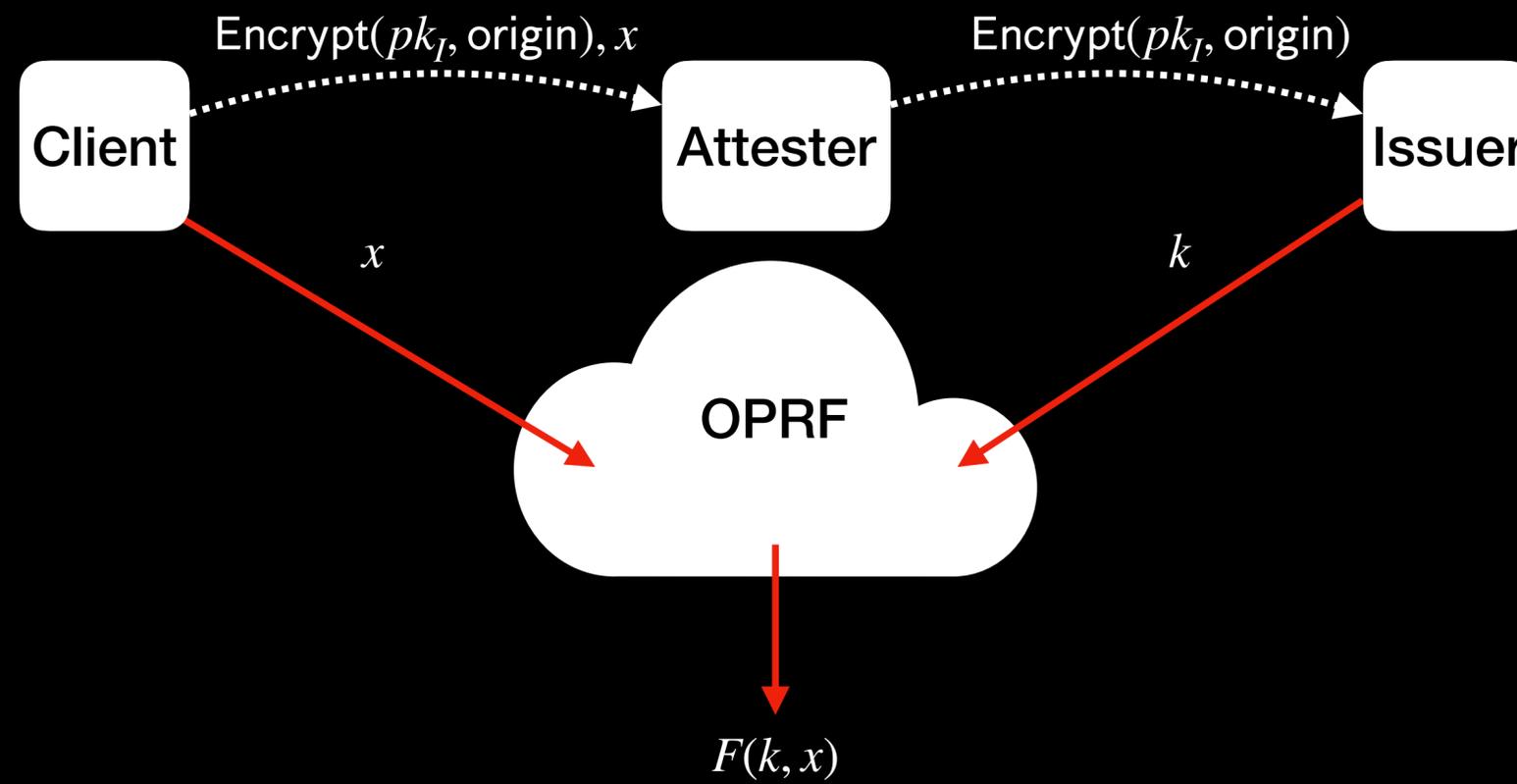
Clients can encrypt the origin identifier under the Issuer's public key



# Detour: An OPRF Sketch

## Issuance protocols

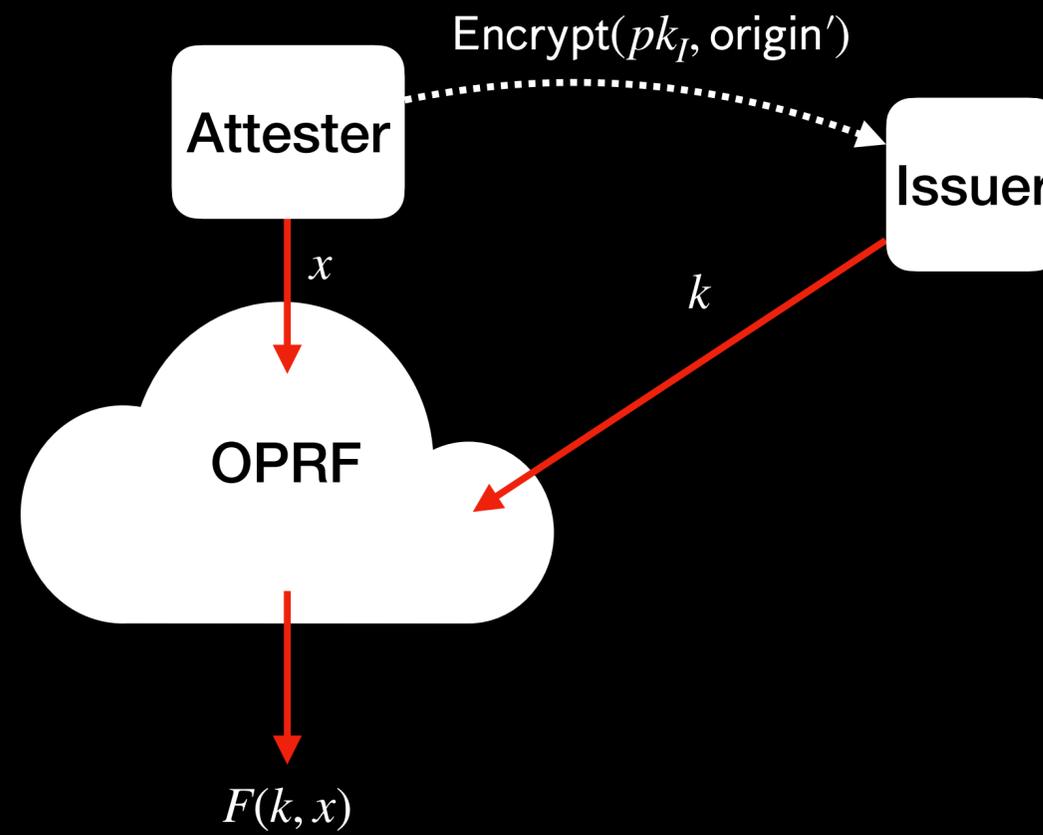
An Attester can relay the encrypted origin name and complete the OPRF



# Detour: An OPRF Sketch

## Issuance protocols

... Attester can perform a dictionary attack to learn  $F(k, x)$



# Rate-Limited Tokens

## Issuance protocols

Rate-limited tokens *extend* the basic issuance protocol with new properties:

1. Issuers learn origin associated with a token challenge
2. Attesters learn *stable mapping* between per-client secret and per-origin secret, and no per-origin information
3. Token requests may fail if the per-origin rate limit is exceeded

**Challenge 1:** How to reveal only the origin to issuer, and only the mapping to attester?

**Challenge 2:** How to ensure the attester cannot dictionary attack or replay client requests to learn per-origin information?

Proposed solution to both  
uses the same mechanism!

# Rate-Limited Tokens

## Issuance protocols

Cryptographic primitives:

- Blind RSA: Token request
- HPKE: Encrypting origin names from Client to Issuer
- EdDSA with key blinding: Signing Client requests *and* computing stable mappings

This is the interesting piece!

# Detour: EdDSA with Key Blinding

## Issuance protocols

Extend RFC8032 EdDSA with two functionalities

BlindPublicKey and UnblindPublicKey: Given public key and secret blind, produce *blinded public key*

$$\text{UnblindPublicKey}(\text{BlindPublicKey}(pkS, skB), skB) = pkS$$

BlindKeySign: Sign message with secret key and secret blind

$$\text{Verify}(\text{BlindPublicKey}(pkS), msg, \text{BlindKeySign}(skS, skB, msg)) = true$$

# Rate-Limited Tokens Issuance protocols

