

ACT Update

Draft progress and a path to general state machines

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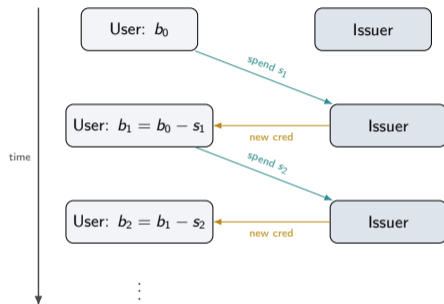
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What ACT is

A two-party protocol between a **user** and an **issuer**:

- ▶ User holds an *anonymous credential* carrying a hidden balance b .
- ▶ User presents an unlinkable spend proof that $b \geq s$; the issuer returns a refreshed credential.

The chain of refreshed credentials *is* the state machine.



What's changed in ACT since last time

Three updates, all landing in draft-act by IETF 126:

- 1 **Sigma proofs migration:** moving off bespoke sigma proofs onto the **CFRG sigma proofs draft**.
- 2 **SHAKE128** for hash-to-scalar and Fiat–Shamir, aligning with that draft.
- 3 **Flexible refunds:** the server credits back any $r \in [0, S_{\max}]$.

From scalar spend to general state machines

ACT today: $b' = b - s$, $s \in [0, S_{\max}]$, $b' \geq 0$.

Generalize the transition to a single *affine* step in a state vector:

$$\mathbf{s}' = T \mathbf{s} + \mathbf{u}, \quad \mathbf{s}, \mathbf{s}' \in \mathbb{Z}_q^k, \quad T \in \mathbb{Z}_q^{k \times k}, \quad \mathbf{u} \in \mathbb{Z}_q^k.$$

A valid transition proves, in zero knowledge:

- ▶ **Inputs:** each component of \mathbf{u} satisfies its declared role (user-chosen, server-chosen, or jointly-determined) with its own range or equality bound.
- ▶ **Consequence:** \mathbf{s} satisfies a conjunction of range / equality predicates over linear combinations of its entries.

Why? Epochs; access levels that control rate limits or rate consumption; and more generally, richer authorization policies expressed inside one credential.

Simple example: tiered per-spend cap

State vector $\mathbf{s} = (b, \ell)$: balance b and an immutable access tier ℓ baked in at issuance.

Policy constants α, β .

Single admitted transition, server-chosen spend s :

$$\underbrace{\begin{pmatrix} b' \\ \ell' \end{pmatrix}}_{\mathbf{s}'} = \underbrace{\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}}_T \underbrace{\begin{pmatrix} b \\ \ell \end{pmatrix}}_{\mathbf{s}} + \underbrace{\begin{pmatrix} -s \\ 0 \end{pmatrix}}_{\mathbf{u}(s)}.$$

Predicates proved in ZK (each a range proof on a linear combination):

- ▶ $0 \leq s$ (server input is well-formed)
- ▶ $s \leq \alpha \ell + \beta$ (per-spend cap scales with the hidden tier)
- ▶ $b' \geq 0$, i.e. $b - s \geq 0$ (solvency)

Questions

draft-act: github.com/samuelSchlesinger/draft-act